

MTHE 474 Notes

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Fall 2022

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1 Chapter 2 Topics

1.1 Information Measures for Discrete Systems

1.1.1 Definitions

- **Definition 2.2:** Entropy of discrete random variable X with pmf $P_X(\cdot)$ is defined as

$$H(X) := - \sum_{x \in \mathcal{X}} P_X(x) \log_2 P_X(x)$$

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- **Definition 2.8 (Joint entropy):**

$$H(X, Y) := - \sum_{(x, y) \in \mathcal{X} \times \mathcal{Y}} P_{X, Y}(x, y) \log_2 P_{X, Y}(x, y)$$

- **Definition 2.9 (Conditional entropy):**

$$H(Y|X) := \sum_{x \in \mathcal{X}} P_X(x) \left(- \sum_{y \in \mathcal{Y}} P_{Y|X}(y|x) \log_2 P_{Y|X}(y|x) \right)$$

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1.1.2 Lemmas/Theorems

- **Lemma 2.4 (Fundamental Inequality):** $\forall x > 0$ and $D > 1$ we have

$$\log_D(x) \leq \log_D e * (x - 1)$$

- **Lemma 2.5 (Non-negativity):** $H(X) \geq 0$

- **Lemma 2.6 (Entropy Upper-Bound):** $H(X) \leq \log_2 |\mathcal{X}|$ where random variable X takes values from finite set \mathcal{X}

- **Lemma 2.7 (Log-Sum inequality):** For nonnegative numbers, a_1, a_2, \dots, a_n and b_1, b_2, \dots, b_n

$$\sum_{i=1}^n (a_i \log_D \frac{a_i}{b_i}) \leq (\sum_{i=1}^n a_i) \log_D \frac{\sum_{i=1}^n a_i}{\sum_{i=1}^n b_i}$$

with equality iff for all $i = 1, \dots, n$

$$\frac{a_i}{b_i} = \frac{\sum_{j=1}^n a_j}{\sum_{j=1}^n b_j}$$

is constant and does not depend on i

- **Theorem 2.10 (Chain rule for entropy):** $H(X, Y) = H(X) + H(Y|X)$

1.2 Mutual Information

1.2.1 Definitions

- **Definition 2.2.1 (Mutual Information):**

$$I(X; Y) := H(X) - H(X|Y)$$

- **Definition 2.2.2 (Conditional Mutual Information):**

$$I(X; Y|Z) := H(X|Z) - H(X|Y, Z)$$

1.2.2 Lemmas

- **Lemma 2.15 (Properties of Mutual Information):**

$$1. I(X; Y) = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P_{X,Y}(x, y) \log_2 \frac{P_{X,Y}(x, y)}{P_X(x)P_Y(y)} \quad (1)$$

$$2. I(X; Y) = I(Y; X) = H(Y) - H(Y|X) \quad (2)$$

$$3. I(X; Y) = H(X) + H(Y) - H(X, Y) \quad (3)$$

$$4. I(X; Y) \leq H(X) \text{ equality iff } X \text{ is a function of } Y \quad (4)$$

$$5. I(X; Y) \leq 0 \text{ with equality iff } X \text{ and } Y \text{ are independent} \quad (5)$$

$$6. I(X; Y) \leq \min\{\log_2 |\mathcal{X}|, \log_2 |\mathcal{Y}|\} \quad (6)$$

- **Lemma 2.16 (Chain Rule for Mutual Information):**

$$I(X; Y, Z) = I(X; Y) + I(X; Z|Y) = I(X; Z) + I(X; Y|Z)$$

- **Theorem 2.17 (Chain Rule for entropy):** $X^n := (X_1, \dots, X_n)$ and $x^n := (x_1, \dots, x_n)$

$$H(X^n) = \sum_{i=1}^n H(X_i | X^{i-1})$$

- **Theorem 2.18 (Chain Rule for conditional entropy):**

$$H(X_1, X_2, \dots, X_n | Y) = \sum_{i=1}^n H(X_i | X_{i-1}, \dots, X_1, Y)$$

- **Theorem 2.19 (Chain Rule for Mutual information):**

$$I(X_1, X_2, \dots, X_n; Y) = \sum_{i=1}^n I(X_i; Y | X_{i-1}, \dots, X_1)$$

Where $I(X_i; Y | X_{i-1}, \dots, X_1) := I(X_i, Y)$ for $i = 1$

1.3 Conditional Divergence

1.3.1 Definitions

1.3.2 Theorems

1.4 Data Processing Inequality

1.4.1 Definitions

- **Lecture 7 Definition (Markov Chain):** Three jointly distributed random variables X, Y, Z are said to form a Markov Chain (in that order), denoted by $X \rightarrow Y \rightarrow Z$ if:

$$P_{XZ|Y}(x, y, z) = P_{X|Y}(x|y)P_{Z|Y}(z, y) \iff P_{Z|XY}(z|x, y) = P_{Z|Y}(z|y)$$

$$\forall x \in X, y \in Y, z \in Z$$

1.4.2 Theorems

- **Lecture 7 Theorem (Data Processing Inequality):** If $X \rightarrow Y \rightarrow Z$, then

$$I(X; Y) \leq I(X; Z)$$

1.5 Convex/Concavity of Information Measures

1.5.1 Definitions

- **Lecture 6 Definition (Convex Set):**

A subset K of \mathbb{R} is called convex if the line segment joining any two points in K also lies in K

- **Lecture 6 Definition (Convex Function):** The function $f : k \rightarrow \mathbb{R}$ where k is a convex subset of \mathbb{R}^n , is called convex on k if $\forall x_1, x_2 \in k$ and $\lambda \in [0, 1]$,

$$f(\lambda x_1 + (1 - \lambda)x_2) \leq \lambda f(x_1) + (1 - \lambda)f(x_2)$$

Strict equality holds whenever $x_1 \neq x_2$ and $0 < \lambda < 1$ then f is called strictly convex

- **Lecture 6 Definition (Concave Function):** $f : k \rightarrow \mathbb{R}$ is concave on k (where $k \subseteq \mathbb{R}^n$ is a concave subset) if $-f$ is convex. In other words: if $\forall x_1, x_2 \in k$ and $\lambda \in [0, 1]$,

$$f(\lambda x_1 + (1 - \lambda)x_2) \geq \lambda f(x_1) + (1 - \lambda)f(x_2)$$

1.5.2 Theorems

- **Lecture 6 Theorem (Jensen's Inequality):** Let $K \subseteq \mathbb{R}$ (where K is a convex set?) and let $f : k \rightarrow \mathbb{R}$ be a convex function. Also let x be a RV with alphabet $\mathcal{X} \subseteq k$ and finite mean, then

$$E[f(x)] \leq f(E[x])$$

Also if f is strictly convex, then the inequality is strict unless x is deterministic

- **Lecture 7 Theorem (Convexity/Concavity of Information Measures):**

i. $D(p||q)$ is convex in the pair (p, q) (ie: if p_1, q_1 and p_2, q_2 are two pairs of PMFs defined on \mathcal{X}) then:

$$D(\lambda p_1 + (1 - \lambda)p_2 || \lambda q_1 + (1 - \lambda)q_2) \leq \lambda D(p_1 || q_1) + (1 - \lambda)D(p_2 || q_2)$$

$$\forall \lambda \in [0, 1]$$

ii. if $x \sim P_x$, then

$$H(x) = H(p_x) \text{ is concave in } P_x$$

iii. If $(x, y) \sim P_X P_{Y|X}$, then $I(X; Y) = I(P_X, P_{Y|X})$ is concave in P_X for fixed $P_{Y|X}$ and convex in $P_{Y|X}$ for fixed P_X

2 Tutorial Proofs

2.1 Week 2 Tutorial

- Given 2 discrete RVs, X, Y we have that

$$H(Y|X) = 0 \iff Y \text{ is a function of } X$$

- Given RV X with alphabet \mathcal{X} and function $f : x \rightarrow \mathbb{R}$

$$H(X) \leq H(f(X))$$