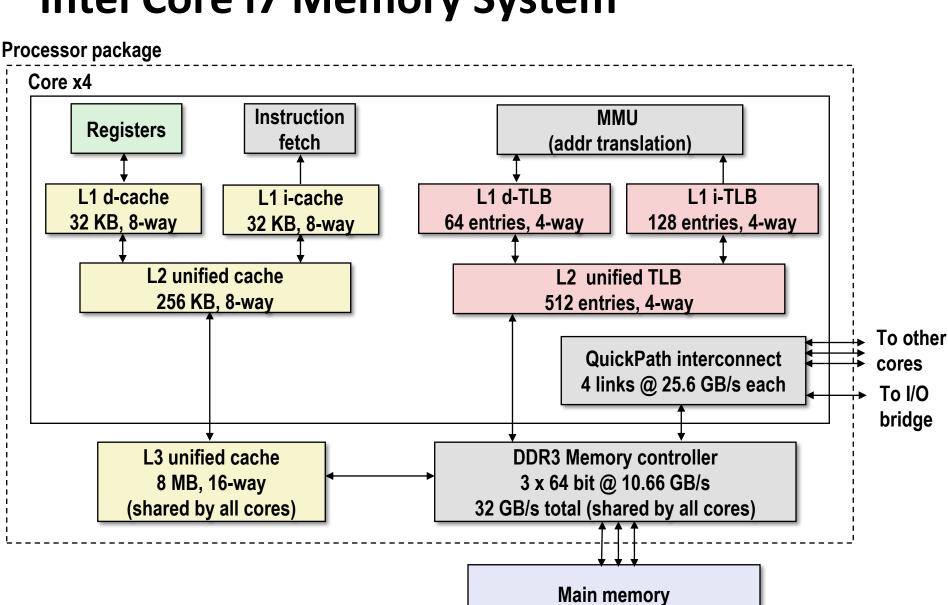
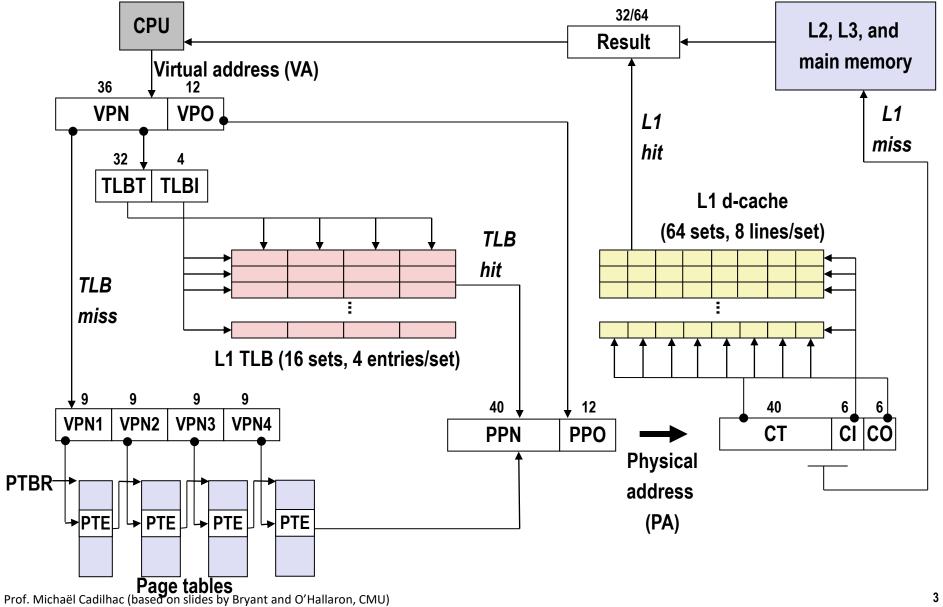
Virtual Memory: Case study: Core i7/Linux memory system

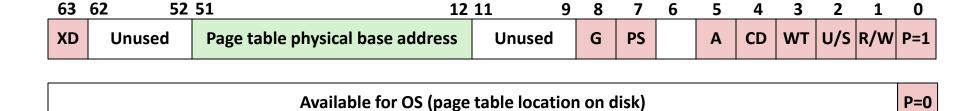
Intel Core i7 Memory System



End-to-end Core i7 Address Translation



Core i7 Level 1-3 Page Table Entries



Each entry references a 4K child page table. Significant fields:

P: Child page table present in physical memory (1) or not (0).

R/W: Read-only or read-write access access permission for all reachable pages.

U/S: user or supervisor (kernel) mode access permission for all reachable pages.

WT: Write-through or write-back cache policy for the child page table.

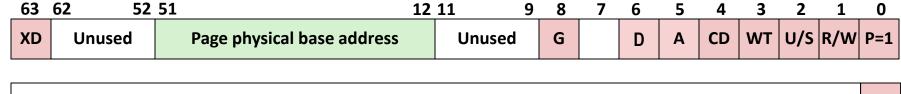
A: Reference bit (set by MMU on reads and writes, cleared by software).

PS: Page size either 4 KB or 4 MB (defined for Level 1 PTEs only).

Page table physical base address: 40 most significant bits of physical page table address (forces page tables to be 4KB aligned)

XD: Disable or enable instruction fetches from all pages reachable from this PTE.

Core i7 Level 4 Page Table Entries



Available for OS (page location on disk)

P=0

Each entry references a 4K child page. Significant fields:

P: Child page is present in memory (1) or not (0)

R/W: Read-only or read-write access permission for child page

U/S: User or supervisor mode access

WT: Write-through or write-back cache policy for this page

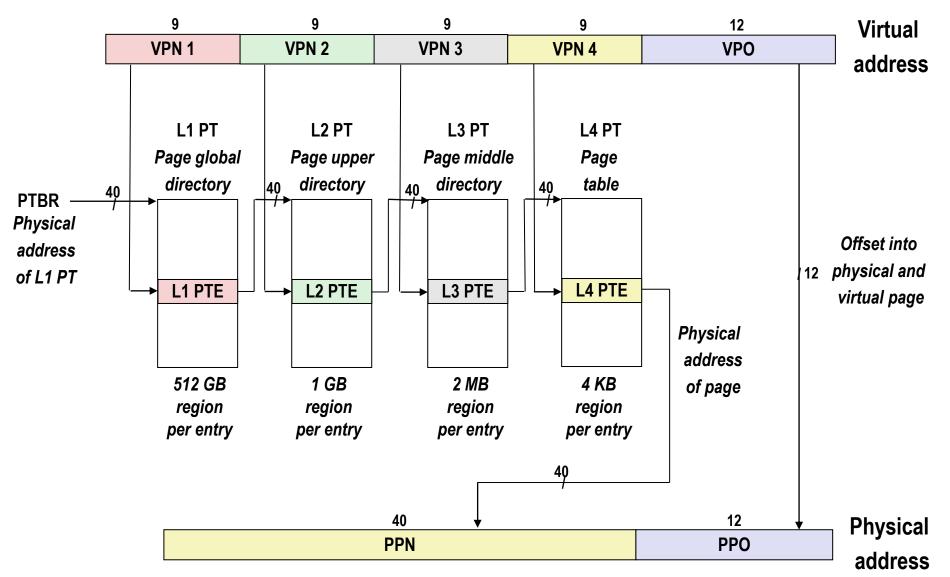
A: Reference bit (set by MMU on reads and writes, cleared by software)

D: Dirty bit (set by MMU on writes, cleared by software)

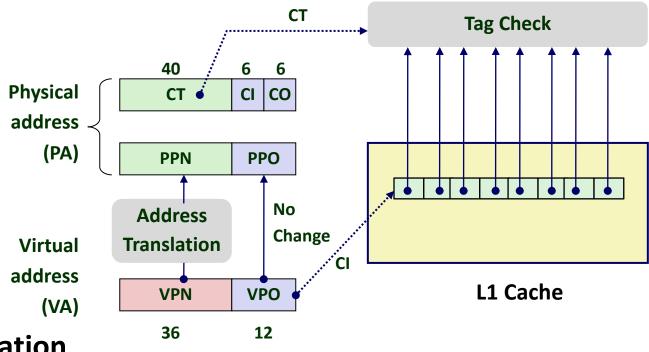
Page physical base address: 40 most significant bits of physical page address (forces pages to be 4KB aligned)

XD: Disable or enable instruction fetches from this page.

Core i7 Page Table Translation



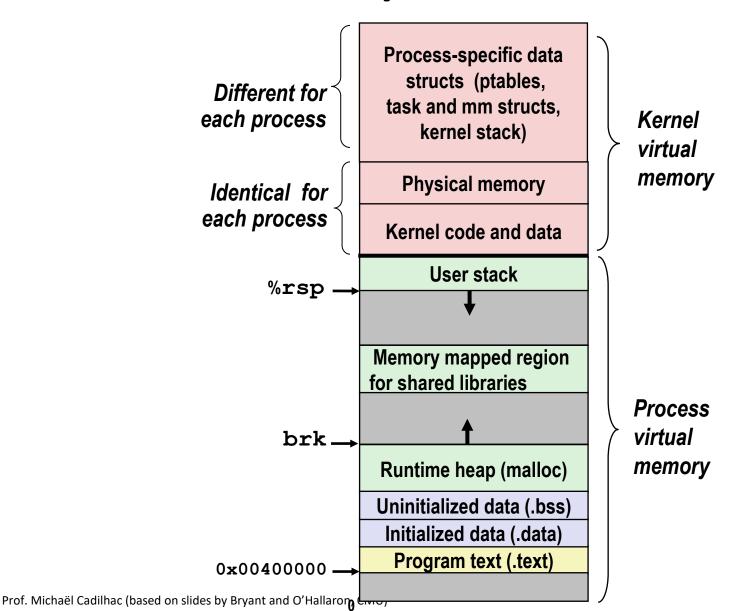
Cute Trick for Speeding Up L1 Access



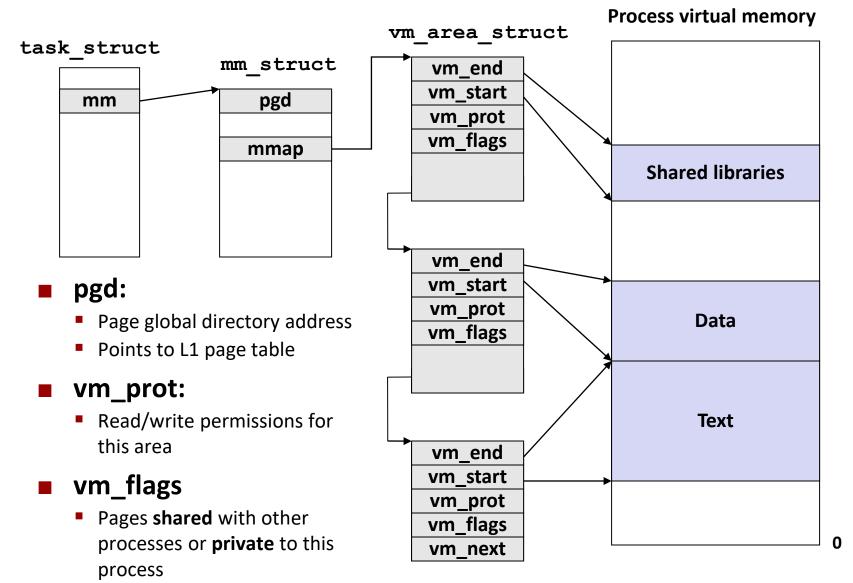
Observation

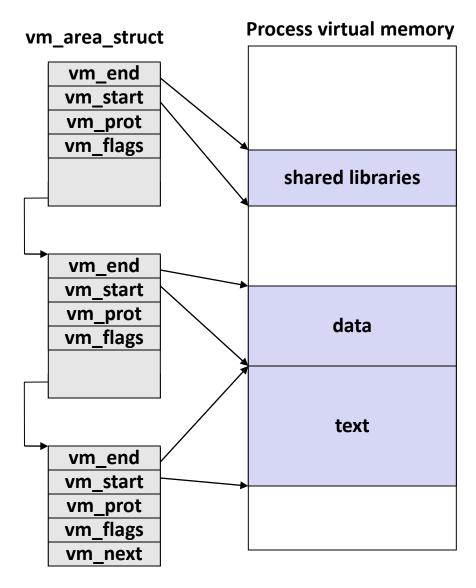
- Bits that determine CI identical in virtual and physical address
- Can index into cache while address translation taking place
- Generally we hit in TLB, so PPN bits (CT bits) available next
- "Virtually indexed, physically tagged"
- Cache carefully sized to make this possible

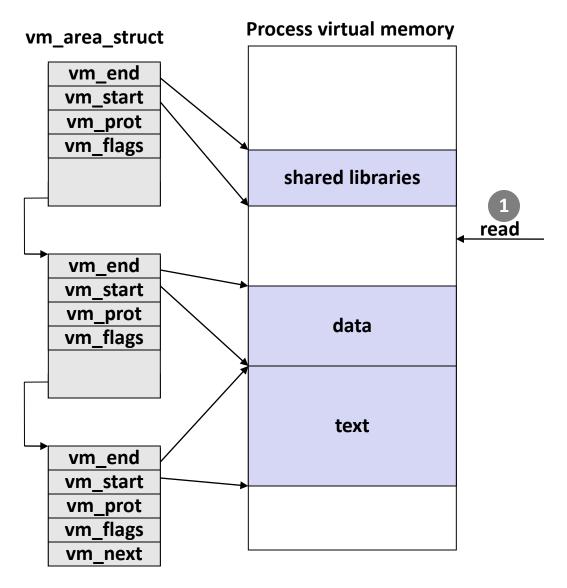
Virtual Address Space of a Linux Process

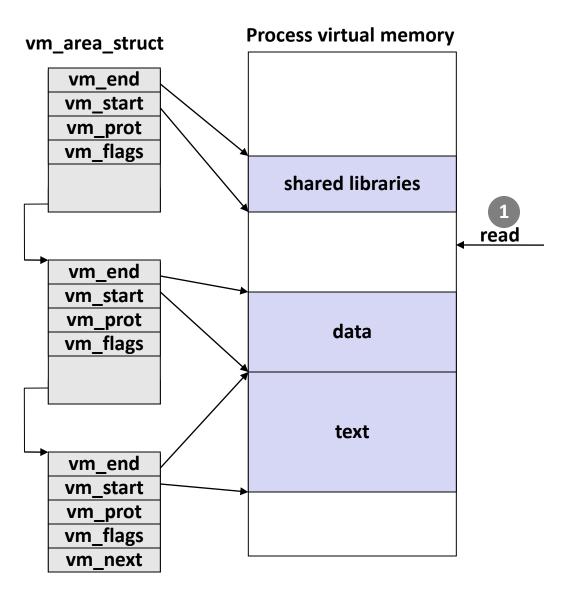


Linux Organizes VM as Collection of "Areas"

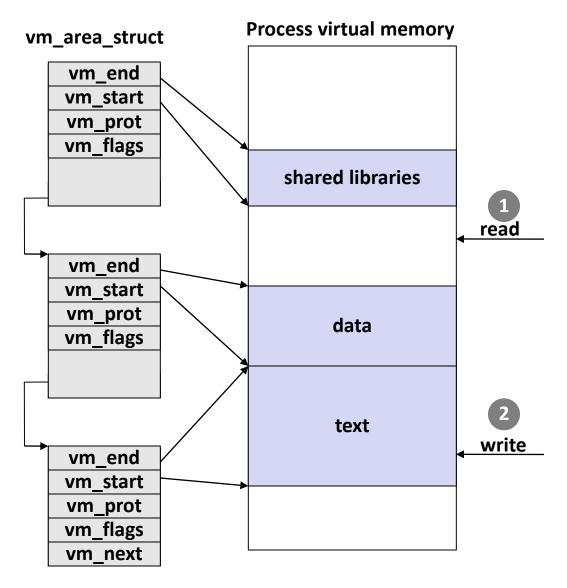




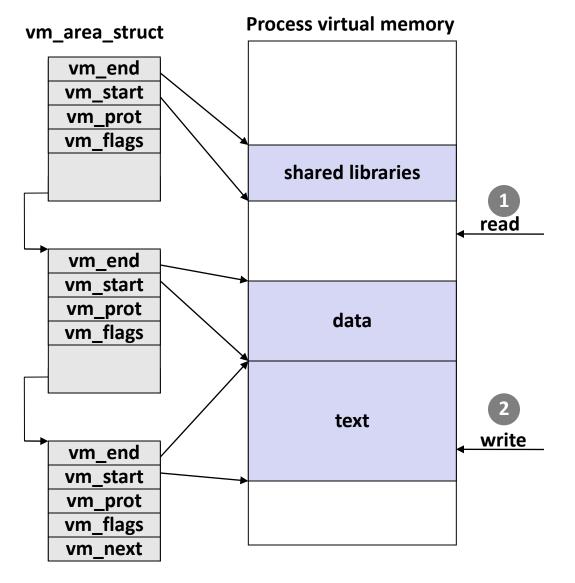




Segmentation fault: accessing a non-existing page



Segmentation fault: accessing a non-existing page

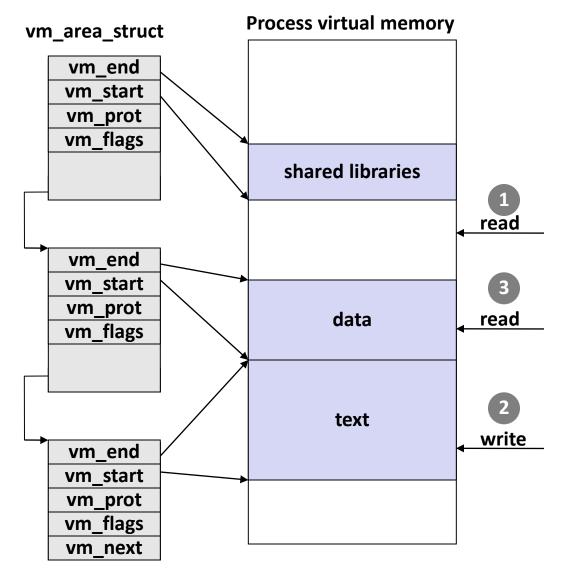


Segmentation fault:

accessing a non-existing page

Protection exception:

e.g., violating permission by writing to a read-only page (Linux reports as Segmentation fault)

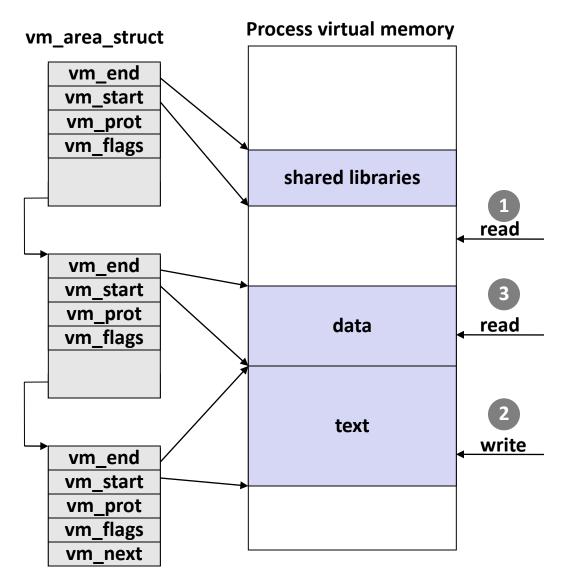


Segmentation fault:

accessing a non-existing page

Protection exception:

e.g., violating permission by writing to a read-only page (Linux reports as Segmentation fault)



Segmentation fault: accessing a non-existing page

Normal page fault

Protection exception:

e.g., violating permission by writing to a read-only page (Linux reports as Segmentation fault)