

EE 231A: Information Theory



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EE 231E: Information Theory Lecture 1

- A. Introduction
- B. Entropy
- C. Relative Entropy
- D. Mutual Information

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Mutual information

- $I(X;Y)$ describes how much information X contains about Y and vice versa.
- Formal Definition is (2.28) on page 20, but that is not the first expression to learn.

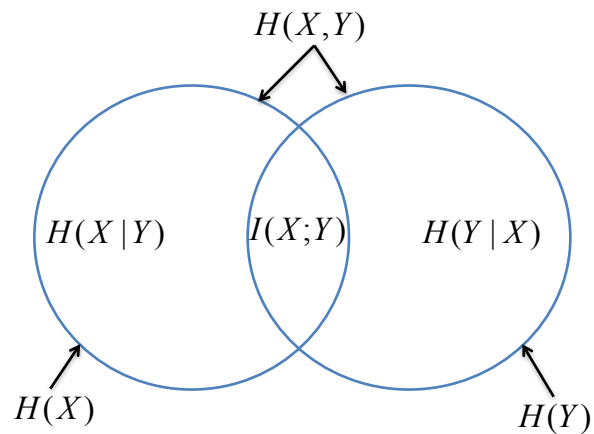
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Operational Definition of Mutual information

$$\begin{aligned}
 I(X;Y) &= I(Y;X) \\
 &= H(X) - H(X|Y) \\
 &= H(Y) - H(Y|X) \\
 &= H(X) + H(Y) - H(X,Y) \quad \text{Why?} \\
 &= H(X) + H(Y) - \underbrace{H(Y) - H(X|Y)}_{-H(X,Y)} \\
 &= H(X) + H(Y) - H(Y) - H(X|Y)
 \end{aligned}$$

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Venn diagram



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A special case of $D(p||q)$

- What if $q(x, y) = p(x)p(y)$

$$\begin{aligned}
 D(p(x, y) \parallel p(x)p(y)) &= E_{p(x,y)}[-\log p(x)p(y)] - E[-\log p(x, y)] \\
 &= H(X) + H(Y) - H(X, Y) \\
 &= I(X; Y)
 \end{aligned}$$

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Mutual information and relative entropy

$$\begin{aligned} I(X;Y) &\triangleq \sum_x \sum_y p(x,y) \log \frac{p(x,y)}{p(x)p(y)} \\ &= D(p(x,y) \parallel p(x)p(y)) \end{aligned}$$

- So $I(X;Y)$ measures the penalty associated with compressing two dependent R.V.'s as if they were independent.
- i.e. $I(X;Y)$ measures how much information X has about Y and vice versa.

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Information vs. Probability of Error

- The concept of mutual information provides a way for us to evaluate how informative some data is.
- i.e. Does this data tell me something I don't know?
- On your homework you will explore whether you should choose the weather forecaster having the lower probability of a wrong forecast or the higher mutual information.

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Mutual information and entropy

- Let's use the formal definition to show

$$I(X;Y) = H(X) - H(X|Y)$$

- Proof:

$$\begin{aligned}
 I(X;Y) &\triangleq \sum_x \sum_y p(x,y) \log \frac{p(x,y)}{p(x)p(y)} \\
 &= \sum_x \sum_y p(x,y) \log \frac{p(x|y)}{p(x)} \\
 &= \underbrace{\sum_x \sum_y p(x,y) \log p(x|y)}_{-H(X|Y)} - \underbrace{\sum_x \sum_y p(x,y) \log p(x)}_{p(x)} \\
 &= H(X) - H(X|Y)
 \end{aligned}$$

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Conditional mutual information

- $I(X;Y|Z) = H(X|Z) - H(X|Y,Z)$

- Chain rule for mutual information

$$I(X_1, X_2, \dots, X_n; Y) = \sum_{i=1}^n I(X_i; Y | X_{i-1}, X_{i-2}, \dots, X_1)$$

- Proof: $I(\underbrace{X_1, X_2, \dots, X_n}_{X^n}; Y)$

$$\begin{aligned}
 &= H(Y) - H(Y|X^n) \\
 &= H(X^n) - H(X^n|Y) \quad \begin{array}{l} \swarrow \text{Which decomposition} \\ \searrow \text{Is helpful?} \end{array} \\
 &= \sum_{i=1}^n H(X_i | X_{i-1}, X_{i-2}, \dots, X_1) - \sum_{i=1}^n H(X_i | X_{i-1}, X_{i-2}, \dots, X_1, Y) \\
 &= \sum_{i=1}^n I(X_i; Y | X_{i-1}, X_{i-2}, \dots, X_1)
 \end{aligned}$$

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