

Liveness:

```
spin -a file
gcc -o pan pan.c
pan -a -f or ./pan -a -f
spin -t -p -l -g -r -s file
```

Spin arguments

-a	generate verifier and syntax check
-i	interactive simulation
-I	display Promela program after preprocessing
-nN	seed for random simulation
-t	guided simulation with trail
-tN	guided simulation with Nth trail
-uN	maximum number of steps is N
-f	translate an LTL formula into a never claim
-F	translate an LTL formula in a file into a never claim
-N	include never claim from a file
-l	display local variables
-g	display global variables
-p	display statements
-r	display receive events
-s	display send events

Compile arguments

-DBFS	breadth-first search
-DNP	enable detection of non-progress cycles
-DSAFETY	optimize for safety
-DBITSTATE	bitstate hashing
-DCOLLAPSE	collapse compression
-DHC	hash-compact compression
-DMA=n	minimized DFA with maximum n bytes
-DMEMLIM=N	use up to N megabytes of memory

Pan arguments

-a	find acceptance cycles
-f	weak fairness
-l	find non-progress cycles

-cN	stop after Nth error
-c0	report all errors
-e	create trails for all errors
-i	search for shortest path to error
-I	approximate search for shortest path to error
-mN	maximum search depth is N
-wN	2^N hash table entries

-A	suppress reporting of assertion violations
-E	suppress reporting of invalid end states

Caveats

- Expressions must be side-effect free.
- Local variable declarations always take effect at the beginning of a process.
- A `true` guard can always be selected; an `else` guard is selected only if all others are false.
- Macros and `inline` do *not* create a new scope.
- Place labels before an `if` or `do`, *not* before a guard.
- In an `if` or `do` statement, interleaving can occur between a guard and the following statement.
- Processes are activated and die in LIFO order.
- Atomic propositions in LTL formulas must be identifiers starting with lowercase letters and must be boolean variables or symbols for boolean-valued expressions.
- Arrays of `bit` or `bool` are stored in bytes.
- The type of a message field of a channel cannot be an array; it can be a `typedef` that contains an array.
- The functions `empty` and `full` cannot be negated.

References

- G. J. Holzmann. *The Spin Model Checker: Primer and Reference Manual*, Addison-Wesley, 2004.
<http://spinroot.com>.
- M. Ben-Ari. *Principles of the Spin Model Checker*, Springer, 2008.
<http://www.springer.com/978-1-84628-769-5>.

Spin Reference Card

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September 29, 2008

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Datatypes

`bit` (1 bit)
`bool` (1 bit)
`byte` (8 bits unsigned)
`short` (16* bits signed)
`int` (32* bits signed)
`unsigned` ($\leq 32^*$ bits unsigned)
 * - for a 32-bit machine.
`pid`
`chan`
`mtype` = { `name`, `name`, ... } (8 bits)
`typedef` `typename` { sequence of declarations }

Declaration - `type var [= initial value]`

Default initial values are zero.

Array declaration - `type var[N] [= initial value]`

Array initial value assigned to all elements.

Operators (descending precedence)

()	[]	.
!	~	++ --
*	/	%
+	-	
<<	>>	
<	<=	> >=
==	!=	
&		
^		

```
|
&&
||
( ... -> ... : ... ) conditional expression
=
```

Predefined

Constants - true, false

Variables (read-only except _):

_ - write-only hidden scratch variable
 _nr_pr - number of processes
 _pid - instantiation number of executing process
 timeout - no executable statements in the system?

Preprocessor

```
#define name (arguments) string
#undef, #if, #ifdef, #ifndef, #else, #endif
#include "file name"
inline name (arguments) { ... }
```

Statements

Assignment - var = expression, var++, var--
 assert(expression)

printf, printm - print to standard output
 %c (character), %d (decimal), %e (mtype),
 %o (octal), %u (unsigned), %x (hex)
 scanf - read from standard input in simulation mode

skip - no operation
 break - exit from innermost do loop
 goto - jump to label
 Label prefixes with a special meaning:
 accept - accept cycle
 end - valid end state
 progress - non-progress cycle

atomic { ... } - execute without interleaving
 d_step { ... } - execute deterministically (no jumping in or out; deterministic choice among true guards; only the first statement can block).

{ ... } unless { ... } - exception handling.

Guarded commands

```
if :: guard -> statements :: ... fi
do :: guard -> statements :: ... od
else guard - executed if all others are false.
```

Processes

Declaration - proctype procname (parameters) { ... }
 Activate with prefixes - active or active[N]
 Explicit process activation - run procname (arguments)
 Initial process - init { ... }
 Declaration suffixes:
 priority - set simulation priority
 provided (e) - executable only if expression e is true

Channels

chan ch = [capacity] of { type, type, ... }

```
ch ! args      send
ch !! args     sorted send
```

```
ch ? args      receive and remove if first message matches
ch ?? args     receive and remove if any message matches
ch ? <args>    receive if first message matches
ch ?? <args>   receive if any message matches
ch ? [args]    poll first message (side-effect free)
ch ?? [args]   poll any message (side-effect free)
```

Matching in a receive statement: constants and mtype symbols must match; variables are assigned the values in the message; eval(expression) forces a match with the current value of the expression.

len(ch) - number of messages in a channel
 empty(ch) / nempty(ch) - is channel empty / not empty?
 full(ch) / nfull(ch) - is channel full / not full?

Channel use assertions:

```
xr ch - channel ch is receive-only in this process
xs ch - channel ch is send-only in this process
```

Temporal logic

!	not
&&	and
	or
->	implies
<->	equivalent to
[]	always
<>	eventually
X	next
U	strong until
V	dual of U defined as pVq <-> !(pU!q)

Remote references

Test the control state or the value of a variable:

```
process-name @ label-name
proctype-name [ expression ] @ label-name
process-name : label-name
proctype-name [ expression ] : label-name
```

Never claim

never { ... }.

Predefined constructs that can only appear in a never claim:

```
_last - last process to execute
enabled(p) - is process enabled?
np_ - true if no process is at a progress label
pc_value(p) - current control state of process
remote references
```

See also trace and notrace.

Variable declaration prefixes

hidden - hide this variable from the system state
 local - a global variable is accessed only by one process
 show - track variable in Xspin message sequence charts

Verification

Safety:

```
spin -a file
gcc -DSAFETY -o pan pan.c
pan or ./pan
spin -t -p -l -g -r -s file
```