## CS24 Elementary Computer Organization Appendix A Exercises: A.1, A.3, A.5, A.7, A.17

[A.1] [10] <A.2> In addition to the basic laws we discussed in this section, there are two important theorems, called DeMorgan's theorems:

$$\overline{A+B} = \overline{A} \times \overline{B}$$
 and  $\overline{A \times B} = \overline{A} + \overline{B}$ 

Prove DeMorgan's theorems with a truth table.

| A | В | $\overline{A}$ | $\overline{B}$ | A + B | $\overline{A} \times \overline{B}$ | $\overline{A \times B}$ | $\overline{A} + \overline{B}$ |
|---|---|----------------|----------------|-------|------------------------------------|-------------------------|-------------------------------|
| 0 | 0 | 1              | 1              | 1     | 1                                  | 1                       | 1                             |
| 0 | 1 | 1              | 0              | 0     | 0                                  | 1                       | 1                             |
| 1 | 0 | 0              | 1              | 0     | 0                                  | 1                       | 1                             |
| 1 | 1 | 0              | 0              | 0     | 0                                  | 0                       | 0                             |

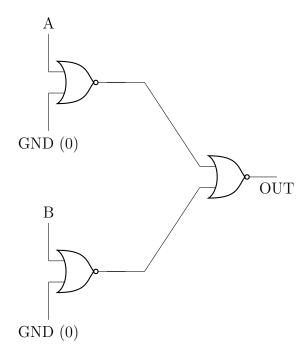
We can conclude that DeMorgan's theorems are correct given that the bit sequences for  $\overline{A} + \overline{B}$  are in fact equal to  $\overline{A} \times \overline{B}$  and the bit sequences for  $\overline{A} \times \overline{B}$  are in fact equal to  $\overline{A} + \overline{B}$ 

 $\boxed{\textbf{A.3}}$  [10] <A.2> Show that there are  $2^n$  entries in a truth table for a function with n inputs.

If you have a decoder, you **must** have  $2^n$  outputs for every n inputs. The reason behind this is due to the fact that with n inputs, you can represent every base 10 number up until  $2^n$  with a binary representation that is n bits long. In the context of a decoder, you can decode  $2^n$  numbers with n single bit inputs.

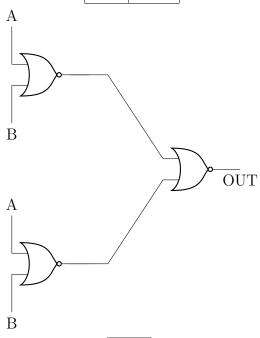
[A.5] [15] <A.2>Prove that the NOR gate is universal by showing how to build the AND, OR, and NOT functions using a two input NOR gate.

| AND |     |  |  |
|-----|-----|--|--|
| АВ  | OUT |  |  |
| 0.0 | 0   |  |  |
| 0.1 | 0   |  |  |
| 1 0 | 0   |  |  |
| 1 1 | 1   |  |  |



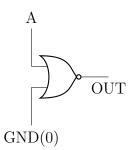
## OR

| АВ  | OUT |
|-----|-----|
| 0.0 | 0   |
| 0.1 | 1   |
| 1 0 | 1   |
| 1 1 | 1   |



NOT

| A | OUT |
|---|-----|
| 0 | 1   |
| 1 | 0   |



[A.7] [10] <A.2, A.3>Construct the truth table for a four input odd parity function (see page A-65 for a description of parity)

A parity is defined as a count of 1's in a word. The word has odd parity if the number of 1's is odd. We will base our truth table off of this information.

| ABCD    | Odd Number of 1's? |
|---------|--------------------|
| 0 0 0 0 | 0                  |
| 0 0 0 1 | 1                  |
| 0 0 1 0 | 1                  |
| 0 0 1 1 | 0                  |
| 0 1 0 0 | 1                  |
| 0 1 0 1 | 0                  |
| 0 1 1 0 | 0                  |
| 0 1 1 1 | 1                  |
| 1000    | 1                  |
| 1 0 0 1 | 0                  |
| 1010    | 0                  |
| 1 0 1 1 | 1                  |
| 1 1 0 0 | 0                  |
| 1101    | 1                  |
| 1 1 1 0 | 1                  |
| 1111    | 0                  |

[A.17] [5] <A.2, A.3>Show a truth table for a multiplexor (inputs A, B, and S; output C), using don't cares to simplify the table where possible.

Assuming the 'on' channel for A is set to 0 and the 'on' channel for B is set to 1, we will have the following truth table for a two input (A and B), one selector (S), and one output (C), multiplexor:

| S A B | С |
|-------|---|
| 0 0 0 | 0 |
| 0 0 1 | 0 |
| 0 1 0 | 1 |
| 0 1 1 | 1 |
| 1 0 0 | 0 |
| 1 0 1 | 1 |
| 1 1 0 | 0 |
| 1 1 1 | 1 |
|       |   |