

1 (5') Stack, Queue and Complexity Analysis

Each question has one or more correct answer(s). Select all the correct answer(s). For each question, you get 0 point if you select one or more wrong answers, but you get 0.5 point if you select a non-empty subset of the correct answers.

Note that you should write your answers of section 1 in the table below.

Question 1	Question 2	Question 3	Question 4	Question 5
D	AD	ABCD	A	ABC

$$c = a/b$$

$$r = a - c \cdot b$$

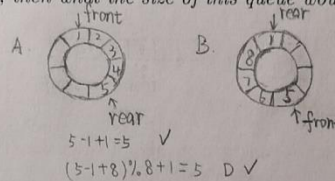
Question 1. In the lectures of Week 2, suppose we implement a circular queue by using an array with the index range from 1 to n , then what the size of this queue would be? We assume that the queue is non-empty.

(A) $rear - front + 1$ ✗

(B) $(rear - front + 1) \% n$ ✗

(C) $(rear - front + n) \% n$ ✗

(D) $(rear - front + n) \% n + 1$ ✓



Question 2. Which of the following is known to be correct?

(A) Stack is a linear data structure and the operations on stacks are more restricted, the same is true for queue. ✓

(B) Lists store elements in sequential locations in memory. ✗

(C) Both stacks and queues allow us insert or delete an element at the front. ✗

(D) We can use two queues to implement stack. ✓

Question 3. Which of the following is/are applications of queue and stack respectively?

(A) Queue: A resource shared by multiple users/processes; Stack: Handling function calls ✓

(B) Queue: Loading Balancing; Stack: Reverse-Polish Notation ✓

(C) Queue: Handling of interrupts in real-time systems; Stack: Compilers/Word Processors ✓

(D) Queue: IO Buffers; Stack: Arithmetic expression evaluation ✓

Question 4. Read the following code, what function does it realize?

```
void Q4(Queue &Q)
{
    Stack S;
    int d;
    InitStack(S);
    while (!QueueEmpty(Q))
    {
```

```

DeQueue(Q, d)
Push(S, d);
}
while (!StackEmpty(S))
{
    Pop(S, d);
    EnQueue(Q, d);
}
}
    
```

Q6:

$$8 \ 2 \ 3 \wedge / 2 \ 3 \ * + 5 \ 1 \ * -$$

$$\begin{array}{l} \underbrace{2^3=8} \\ \underbrace{8/8=1} \quad \underbrace{2*3=6} \\ \underbrace{1+6=7} \quad \underbrace{5*1=5} \\ \underbrace{7-5=2} \end{array}$$

$$8 / 2^3 + 2 * 3 - 5 * 1 = 2$$

- (A) Use stack to reverse the queue. ✓
 (B) Use queue to reverse the stack.
 (C) Use stack to implement the queue.
 (D) Use queue to implement the stack.

Question 5. Which of the following comparison is correct?

- (A) $n^2 + n^3 = O(n^4)$ ✓ $\lim_{n \rightarrow \infty} \frac{n^2 + n^3}{n^4} = 0$
 (B) $\log_2 n = \Theta(\log n)$ ✓ $\lim_{n \rightarrow \infty} \frac{\log n}{\log n} = 1$
 (C) $\log^2 n = \Omega(\log \log n)$ ✓ $\lim_{n \rightarrow \infty} \frac{\log^2 n}{\log \log n} = \infty$
 (D) $n! = \omega(n^n)$ ✗ $\lim_{n \rightarrow \infty} \frac{n!}{n^n} = 0$

$$\begin{array}{l} \Omega(g(n)) \\ O(g(n)) \\ \Theta(g(n)) \\ \omega(g(n)) \end{array}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \quad 0 < c < \infty \quad \infty$$

$$\frac{1}{n!} \sim \frac{e^n}{n! 2\pi n} \quad (n \rightarrow +\infty)$$

2 (10') Stack and Queue

Question 6. (2') The following post-fix expression (Reverse-Polish Notation) with single digit operands is evaluated using a stack:

$$8 \ 2 \ 3 \wedge / 2 \ 3 \ * + 5 \ 1 \ * -$$

Note that \wedge is the exponentiation operator. Please write down the corresponding in-fix notation A and the final result:

$$8 / 2^3 + 2 * 3 - 5 * 1 = 2 \quad (\text{The process is on the above})$$

Question 7. (4') Describe how to implement a queue using a singly-linked list. You can use pseudocode or natural language to describe all the operations, especially the key operations.

First: define a class Single-list, including some functions: bool empty() const determine whether it is empty
 Type front() const return the data stored in the first node; void push-back(Type const &) insert a new element at the back of the singly-linked list; void pop-front() pop the data in the first node.
 Then: we can use this singly-linked list to implement a queue.

template <typename Type>

class Queue {

private:

Single-list list;

public:

bool empty() const {

return the function empty() of the list; }

Type front() const {

return the function front() of the list;

};

void push (Type const &S)

{ return the function push-back() of the list;

};

Type pop() {

{ return the function pop-front() of the list;

};

Question 8. (1') If we use an array with size N to implement a normal queue, it gets full when the index **Back** pointing to the index = $N-1$

Question 9. (1') By implementing the following operations on stack, the value of x is a
`InitStack(st); Push(st,a); Push(st,b); Pop(st,x); Top(st,x);`

Question 10. (2') What dose "stack overflow" and "stack underflow" mean? (give a short explanation)

Stack overflow: push something into the stack when it is full

Stack underflow: pop something out of the stack when it is empty

3 (8') Complexity Analysis

Question 11. (3') Given a fraction of a code as the following, write down the time complexity for each **for** loop.

<code>for (i=1; i<n; i*=2) {</code>	<u>$O(\log n)$</u>
<code>for (j=n; j>0; j/=2) {</code>	<u>$O(\log n)$</u>
<code>for (k=j; k<n; k+=2) {</code>	<u>$O(n)$</u>
<code>sum += (i + j*k)</code>	
<code>}</code>	
<code>}</code>	
<code>}</code>	

Question 12. (5') Calculate the average processing time $T(n)$ of the following recursive algorithm. Suppose that it takes one unit time for `random(int n)` to return a random integer which is uniformly distributed in the range $[0, n]$. Also note that $T(0) = 0$.

Hints: The equation $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \dots + \frac{1}{n \cdot (n+1)} = \frac{n}{n+1}$ might be needed.

```
int hw( int n ) {
    if ( n <= 0 ) return 0;
    else {
        int i = random( n-1 );
        return hw( i ) + hw( n-1-i );
    }
}
```

$$T(n) = 0 + T(i) + T(n-1-i) \quad i \in [0, n-1]$$

$$\text{Guess } T(n) = O(n) \quad \exists c, T(n) \leq cn$$

$$T(n) \leq 1 + c_i + c(n-1-i)$$

$$= cn + (1-c)$$

$$\exists c \geq 1 \quad T(n) \leq cn$$

$$\therefore T(n) = O(n)$$

$$nT(n) = n \sum_{m=0}^{n-1} T(m) + n$$

$$\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \dots + \frac{1}{n \cdot (n+1)} = \frac{n}{n+1}$$