

## 1 (5') Stack, Queue and Complexity Analysis

Each question has one or more correct answer(s). Select all the correct answer(s). For each question, you get 0 point if you select one or more wrong answers, but you get 0.5 point if you select a non-empty subset of the correct answers.

Note that you should write your answers of section 1 in the table below.

Question 1	Question 2	Question 3	Question 4	Question 5
D	AD	ABCD	A	ABC

$$c = a/b$$

$$r = a - c \cdot b$$

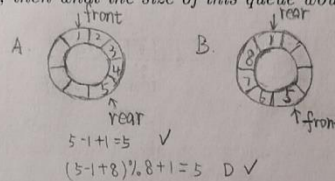
**Question 1.** In the lectures of Week 2, suppose we implement a circular queue by using an array with the index range from 1 to  $n$ , then what the size of this queue would be? We assume that the queue is non-empty.

(A)  $rear - front + 1$  ✗

(B)  $(rear - front + 1) \% n$  ✗

(C)  $(rear - front + n) \% n$  ✗

(D)  $(rear - front + n) \% n + 1$  ✓



$$1 - 5 + 1 = -3$$

$$-3 \% 8 = 5$$

$$-3 / 8 = 1 \dots 5$$

**Question 2.** Which of the following is known to be correct?

(A) Stack is a linear data structure and the operations on stacks are more restricted, the same is true for queue. ✓

(B) Lists store elements in sequential locations in memory. ✗

(C) Both stacks and queues allow us insert or delete an element at the front. ✗

(D) We can use two queues to implement stack. ✓

**Question 3.** Which of the following is/are applications of queue and stack respectively?

(A) Queue: A resource shared by multiple users/processes; Stack: Handling function calls ✓

(B) Queue: Loading Balancing; Stack: Reverse-Polish Notation ✓

(C) Queue: Handling of interrupts in real-time systems; Stack: Compilers/Word Processors ✓

(D) Queue: IO Buffers; Stack: Arithmetic expression evaluation ✓

**Question 4.** Read the following code, what function does it realize?

```
void Q4(Queue &Q)
{
    Stack S;
    int d;
    InitStack(S);
    while (!QueueEmpty(Q))
    {
```

```

DeQueue(Q, d)
Push(S, d);
}
while (!StackEmpty(S))
{
    Pop(S, d);
    EnQueue(Q, d);
}
}

```

Q6:

$$8 \ 2 \ 3 \wedge / 2 \ 3 \ * + 5 \ 1 \ * -$$

$$\begin{array}{l} \underbrace{2^3=8} \\ \underbrace{8/8=1} \quad \underbrace{2*3=6} \\ \underbrace{1+6=7} \quad \underbrace{5*1=5} \\ \underbrace{7-5=2} \end{array}$$

$$8 / 2^3 + 2 * 3 - 5 * 1 = 2$$

- (A) Use stack to reverse the queue. ✓  
 (B) Use queue to reverse the stack.  
 (C) Use stack to implement the queue.  
 (D) Use queue to implement the stack.

Question 5. Which of the following comparison is correct?

- (A)  $n^2 + n^3 = O(n^4)$  ✓  $\lim_{n \rightarrow \infty} \frac{n^2 + n^3}{n^4} = 0$   
 (B)  $\log_2 n = \Theta(\log n)$  ✓  $\lim_{n \rightarrow \infty} \frac{\log n}{\log n} = 1$   
 (C)  $\log^2 n = \Omega(\log \log n)$  ✓  $\lim_{n \rightarrow \infty} \frac{\log^2 n}{\log \log n} = \infty$   
 (D)  $n! = \omega(n^n)$  ✗  $\lim_{n \rightarrow \infty} \frac{n!}{n^n} = 0$

$$\begin{array}{l} \Omega(g(n)) \\ O(g(n)) \\ \Theta(g(n)) \end{array}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \quad 0 < c < \infty \quad \infty$$

$$\frac{1}{n!} \sim \frac{e^n}{n! 2\pi n} \quad (n \rightarrow +\infty)$$

## 2 (10') Stack and Queue

Question 6. (2') The following post-fix expression (Reverse-Polish Notation) with single digit operands is evaluated using a stack:

$$8 \ 2 \ 3 \wedge / 2 \ 3 \ * + 5 \ 1 \ * -$$

Note that  $\wedge$  is the exponentiation operator. Please write down the corresponding in-fix notation A and the final result:

$$8 / 2^3 + 2 * 3 - 5 * 1 = 2 \quad (\text{The process is on the above})$$

Question 7. (4') Describe how to implement a queue using a singly-linked list. You can use pseudocode or natural language to describe all the operations, especially the key operations.

First: define a class Single-list, including some functions: `bool empty() const` ..... determine whether it is empty  
`Type front() const` ..... return the data stored in the first node; `void push-back(Type const &)` ..... insert a new element at the back of the singly-linked list; `void pop-front()` ..... pop the data in the first node.

Then: we can use this singly-linked list to implement a queue.

template <typename Type>

class Queue {

private:  
 Single-list list;

public:  
 bool empty() const {  
 return the function empty() of the list; }  
 Type front() const {  
 return the function front() of the list; }  
 };

void push (Type const &S)

{ return the function push-back() of the list;

};

Type pop()

{ return the function pop-front() of the list;

};

The template of the Single\_list and the implementations are as follows:

```

template <class T>
class Single_node {
private:
    T element;
    Single_node *next_node;

public:
    Single_node( int e, Single_node *n ) : element(e), next_node(n);
    T retrieve() const;
    Single_node *next() const;
};

```

```

template <class T>
class Single_list {
private:
    Single_node *list_head;

public:
    Single_list();
    ~Single_list();

    int size() const; /* return the length of the list */
    bool empty() const; /* return true when list is empty */
    T front() const; /* return the data in the first node */
    T back() const; /* return the data in the last node */
    Single_node * head() const; /* return the first node */
    Single_node * tail() const; /* return the last node */
    int count ( T const & ) const; /* count the number of instances of data */

    void push_front( T const & ); /* insert a node as the first node */
    void push_back( T const & ); /* insert a node as the last node */
    T pop_front(); /* return the data in the first node and delete the first node */
    int erase( T const & ); /* remove the nodes containing the integer */
};

```

```

template <class T>
class Stack {
private:
    Single_list list;
public:
    bool empty() const;
    T top() const;
    void push( T const & );
    T pop();
};

T Single_node :: retrieve() const
{
    return element;
}

Single_node *Single_node :: next() const
{
    return next_node;
}

Single_list :: Single_list : list_head(nullptr)
{
    return;
}

bool Single_list :: empty() const
{
    return ( list_head == nullptr );
}

```

```

bool Single_list :: empty() const
{
    return ( list_head == nullptr );
}

Single_node *Single_list :: head() const
{
    return list_head;
}

template <class T>
T Single_list :: front() const
{
    if ( empty() )
    {
        throw underflow();
    }
    return head() -> retrieve();
}

```

```

template <class T>
void Single_list :: push_front( T const &n )
{
    list_head = new Single_node( n, list_head );
}

template <class T>
int Single_list :: pop_front()
{
    if ( empty() )
    {
        throw underflow();
    }
    int e = front();
    Single_node *ptr = list_head;
    list_head = list_head -> next();
    delete ptr;
    return e;
}

```

```
template <class T>
T Stack :: top() const
{
    if ( empty() )
    {
        throw underflow();
    }
    return list.front();
}

T Stack :: pop()
{
    if ( empty() )
    {
        throw underflow();
    }
    return list.pop_front();
}
```



Question 8. (1') If we use an array with size  $N$  to implement a normal queue, it gets full when the index  $Back$  pointing to the index =  $N-1$

Question 9. (1') By implementing the following operations on stack, the value of  $x$  is 0  
 $InitStack(st)$ ;  $Push(st,a)$ ;  $Push(st,b)$ ;  $Pop(st,x)$ ;  $Top(st,x)$ ;

Question 10. (2') What dose "stack overflow" and "stack underflow" mean? (give a short explanation)

Stack overflow: <sup>push()</sup> push something into the stack when it is full  
 Stack underflow: <sup>pop()</sup> pop something out of the stack when it is empty  
 or  $top()$

### 3 (8') Complexity Analysis

Question 11. (3') Given a fraction of a code as the following, write down the time complexity for each for loop.

```
for ( i=1; i<n; i*=2) {
    for ( j=n; j>0; j/=2) {
        for ( k=j; k<n; k+=2) {
            sum += (i + j*k)
        }
    }
}
```

$O(n \log n)$   
 $O(n \log n)$   
 $O(n)$

Question 12. (5') Calculate the average processing time  $T(n)$  of the following recursive algorithm. Suppose that it takes one unit time for  $random(int n)$  to return a random integer which is uniformly distributed in the range  $[0,n]$ . Also note that  $T(0) = 0$ .

Hints: The equation  $\frac{1}{1*2} + \frac{1}{2*3} + \dots + \frac{1}{n*(n+1)} = \frac{n}{n+1}$  might be needed.

```
int hw( int n) {
    if ( n <= 0 ) return 0;
    else {
        int i = random( n-1 );
        return hw( i ) + hw( n-1-i );
    }
}
```

$T(n) = 0 \cdot \frac{1}{n+1} + \frac{n}{n+1} (1 + T(i) + T(n-1-i))$   
 $= \frac{n}{n+1} + 2 \sum_{k=0}^{n-1} T(k)$   
 $\Rightarrow hw(n) = n$

$T(n) = 0 + T(i) + T(n-1-i) \quad i \in [0, n-1]$   
 Guess  $T(n) = O(n) \exists c, T(n) \leq cn$   
 $T(n) \leq 1 + ci + c(n-1-i)$   
 $= cn + (1-c)$   
 $\exists c \geq 1 \quad T(n) \leq cn$   
 $\therefore T(n) = O(n)$

$nT(n) = n \sum_{m=0}^{n-1} T(m) + n$   
 $\frac{1}{1*2} + \frac{1}{2*3} + \dots + \frac{1}{n*(n+1)} = \frac{n}{n+1}$