

# SIDs: Source and Sink Independent Drivers

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## Note

This memo documents a part of TinyOS for the TinyOS Community, and requests discussion and suggestions for improvements. Distribution of this memo is unlimited. This memo is in full compliance with TEP 1.

## Abstract

This memo documents a set of hardware- and sensor-independent interfaces for data sources and sinks in TinyOS 2.x.

## 1. Introduction

Sensing is an integral part of any sensor network application. The diversity of sensors can lead to a wide variety of different software interfaces to these sensors. However, the burden of connecting a general sensor data management application to every one of these different interfaces suggests that sensors also provide a simple, general-purpose interface for data acquisition. Therefore, TinyOS 2.0 has telescoping sensor abstractions, providing both sensor-independent and sensor-specific interfaces. This memo documents a set of hardware- and sensor-independent interfaces for data sources and sinks in TinyOS 2.x.

## 2. Sensors in TinyOS 1.x

Early TinyOS sensors were generally analog. To sample one of these sensors, an application makes an analog-to-digital conversion using the MCU ADC. Because all early sensors required ADC conversions, the ADC interface has become the de-facto 1.x sensor interface. However, the ADC interface was originally designed for inexpensive, interrupt-driven sampling. All of its commands and events are async and sensor values are always 16 bits, although only some subset of the bits may be significant (e.g., a 12-bit value).

Because sensing is an integral part of high-level application logic, having asynchronous events means that high-level components must work with atomic sections, even if the sampling rate is very low (e.g., every five minutes) and so could be easily placed in a task. Race conditions are problematic and possible in any real time multi-tasking design. Race conditions are a failure in design, and especially difficult to detect at low sampling rates.

Additionally, not all sensors require ADC conversions from the MCU. Many sensors today are digital. To sample these sensors, the MCU sends a sample command and receives the corresponding data over a bus (e.g., SPI, I2C). The latency involved, combined with possible Resource arbitration<sup>1</sup>, means that these bus operations are often synchronous code. In the command direction, this can force a task allocation to convert async to sync; in the event direction, the application has to deal with async code even though the event is, in practice, in a task.

Finally, the simplicity of the ADC interface has led many sensor modules to introduce several new interfaces for calibration and control, such as `Mic` and `MagSetting`. Because ADCs generally do not have error conditions, the ADC interface has no way to signal that a sample failed. This turns out to be important for sensors where the sampling request is split-phase, such as sensors over a bus. In these cases, it is possible that the driver accepts the request to sample, but once acquiring the bus discovers something is wrong with the sensor. This property has led bus-based sensors to also have a separate `ADCErrors` interface; this interface breaks the basic TinyOS pattern of a tight coupling between split-phase commands and their completion events, as the command is in ADC but the completion event is in `ADCErrors`.

All of these complications provide the context of the challenge to write high-level code that is sensor independent. Sensors, when possible, should follow an approach similar to the `HAA[_haa]`, where they have sensor- or sensor-class-specific interfaces for high performance or special case use, but also simple and common interfaces for basic and portable use. Providing a telescoping sensor abstraction allows both classes of use.

### 3. Sensors in TinyOS 2.x

TinyOS 2.x contains several nesC interfaces that can be used to provide sensor-independent interfaces which cover a range of common use cases. This document describes these interfaces, and explains how to use these interfaces to write a Source- or Sink-Independent Driver (SID). A SID is source/sink independent because its interfaces do not themselves contain information on the sort of sensor or device they sit on top of. A SID SHOULD provide one or more of the interfaces described in this section.

This table summarizes the SID interfaces:

Name	Phase	Data type	Section
Read	Split	Scalar	3.1
Get	Single	Scalar	3.2
Notify	Trigger	Scalar	3.3
ReadStream	Split	Stream	3.4

### 3.1 Split-Phase Small Scalar I/O

The first set of interfaces can be used for low-rate scalar I/O:

```
interface Read<val_t> {  
    command error_t read();  
    event void readDone( error_t result, val_t val );  
}
```

If the `result` parameter of the `Read.readDone` and `ReadWrite.readDone` events is not `SUCCESS`, then the memory of the `val` parameter **MUST** be filled with zeroes.

If the call to `read` has returned `SUCCESS`, but the `readDone` event has not yet been signaled, then a subsequent call to `read` **MUST** return `EBUSY` or `FAIL`. This simple locking technique, as opposed to a more complex system in which multiple `read/readDone` pairs may be outstanding, is intended to reduce the complexity of `SID` client code.

Examples of sensors that would be suited to this class of interface include many basic sensors, such as photo, temp, voltage, and ADC readings.

### 3.2 Single-Phase Scalar I/O

Some devices may have their state cached or readily available. In these cases, the device can provide a single-phase instead of split-phase operation. Examples include a node's MAC address (which the radio stack caches in memory), profiling information (e.g., packets received), or a GPIO pin. These devices **MAY** use the `Get` interface:

```
interface Get<val_t> {  
    command val_t get();  
}
```

### 3.3 Notification-Based Scalar I/O

Some sensor devices represent triggers, rather than request-driven data acquisition. Examples of such sensors include switches, passive-IR (PIR) motion sensors, tone detectors, and smoke detectors. This class of event-driven sensors can be presented with the `Notify` interface:

```
interface Notify<val_t> {  
    command error_t enable();  
    command error_t disable();  
    event void notify( val_t val );  
}
```

The `Notify` interface is intended for relatively low-rate events (e.g., that can easily tolerate task latencies). High-rate events may require more platform- or hardware-specific async interfaces.

The `enable()` and `disable()` command enable and disable notification events for the interface instance used by a single particular client. They are distinct from the sensor's power state. For example, if an enabled sensor is powered down, then when powered up it will remain enabled.

If `enable` returns `SUCCESS`, the interface **MUST** subsequently signal notifications when appropriate. If `disable` returns `SUCCESS`, the interface **MUST NOT** signal any notifications.

The `val` parameter is used as defined in the Read interface.

### 3.4 Split-Phase Streaming I/O

Some sensors can provide a continuous stream of readings, and some actuators can accept a continuous stream of new data. Depending on the rate needed and jitter bounds that higher level components can tolerate, it can be useful to be able to read or write readings in blocks instead of singly. For example, a microphone or accelerometer may provide data at a high rate that cannot be processed quickly enough when each new reading must be transferred from asynchronous to synchronous context through the task queue.

The `ReadStream` interface **MAY** be provided by a device that can provide a continuous stream of readings:

```
interface ReadStream<val_t> {  
  
    command error_t postBuffer( val_t* buf, uint16_t count );  
  
    command error_t read( uint32_t usPeriod );  
  
    event void bufferDone( error_t result,  
                           val_t* buf, uint16_t count );  
  
    event void readDone( error_t result );  
}
```

The `postBuffer` command takes an array parameterized by the sample type, and the number of entries in that buffer. A driver can then enqueue the buffer for filling. The client can call `postBuffer()` more than once, to “pre-fill” the queue with any number of buffers. The size of the memory region pointed to by the `buf` parameter **MUST** be at least as large as the size of a pointer on the node architecture plus the size of the `uint16_t` `count` argument. This requirement supports drivers that may store the queue of buffers and count sizes by building a linked list.

After posting at least one buffer, the client can call `read()` with a specified sample period in terms of microseconds. The driver then begins to fill the buffers in the queue, signaling the `bufferDone()` event when a buffer has been filled. The client **MAY** call `postBuffer()` after `read()` in order to provide the device with new storage for future reads.

If the device ever takes a sample that it cannot store (e.g., due to buffer underrun), it **MUST** signal `readDone()` with an appropriate failure code. If an error occurs during a read, then the device **MUST** signal `readDone()` with an appropriate failure code. Before a device signals `readDone()`, it **MUST** signal `bufferDone()` for all outstanding buffers. If a `readDone()` is pending, calls to `postBuffer` **MUST** return `FAIL`.

In the `ReadStream` interface, `postBuffer` returns `SUCCESS` if the buffer was successfully added to the queue, `FAIL` otherwise. A return value of `SUCCESS` from `read` indicates reading has begun and the interface will signal `bufferDone` and/or `readDone` in the future. A return value of `FAIL` means the read did not begin and

the interface MUST NOT signal `readDone` or `bufferDone`. Calls to `read` MAY return `EBUSY` if the component cannot service the request.

## 4. Implementation

An implementation of the Read interface can be found in `tos/system/SineSensorC.nc` and `tos/system/ArbitratedReadC.nc`.

An implementation of the Get interface can be found in `tos/platforms/telosb/UserButtonC.nc`.

An implementation of the ReadStream interface can be found in `tos/sensorboards/mts300/MageXStreamC.nc`.

Implementations of the Notify interface can be found in `tos/platforms/telosb/SwitchToggleC.nc` and `tos/platforms/telosb/UserButtonP.nc`.

## 5. Summary

According to the design principles described in the HAA[\_haa], authors should write device drivers that provide rich, device-specific interfaces that expose the full capabilities of each device. In addition, authors can use the interfaces described in this memo to provide a higher-level device-independent abstractions: SIDs. By providing such an abstraction, driver authors can support developers who only need simple interfaces, and can reduce the effort needed to connect a sensor into a more general system.

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## 7. Citations

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<sup>1</sup>TEP 108: Resource Arbitration.