

1 Introduction

In this report, we want to introduce a method to find the most-probable history of a Markov chain, that is, to find the path which has highest probability.

First, we discuss the problem in a non-hidden Markov chain. And later, the same problem in a hidden Markov chain will be covered. The solution is well-known as the Viterbi algorithm.

To make the problem simple, we discuss in a Markov chain which has only finite number of states.

2 Non-Hidden Cases

Let $(\Omega, \mathcal{F}, \mathcal{P})$ be a probability space.

Suppose that we have a Markov chain $(X_n)_{n \geq 0}$ with finite state-space and the transition matrix P .

Since the state space is finite, there must be an one-to-one correspondence between the state space and $I = \{1, 2, \dots, N\}$.

Thus it is reasonable to consider $X_n : \Omega \rightarrow I$ for all $n \in \mathbb{N}_0$, and $P = p[i, j]_{i, j \in I}$.

Here are some notations we use:

- \mathbb{N} denotes the set $\{1, 2, 3, \dots\}$
- \mathbb{N}_0 denotes the set $\{0\} \cup \mathbb{N}$
- $x^{(n)} = (x_0, x_1, \dots, x_n) \in I^{n+1}$ is an apparent variable denoting a state sequence of length n for some $n \in \mathbb{N}_0$
- $\mathbf{X}^{(n)} = (X_0, X_1, \dots, X_n)$ denotes a Markov history from time 0 to some $n \in \mathbb{N}_0$

Our problem is described as follow:

Given an initial distribution $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_N)$ and a termination time $T \in \mathbb{N}$, what is the deterministic state sequence $x^* = (x_0^*, x_1^*, x_2^*, \dots, x_T^*)$ satisfying

$$\mathcal{P}(\mathbf{X}^{(T)} = x^*) = \max \left\{ \mathcal{P}(\mathbf{X}^{(T)} = x^{(T)}) \mid x^{(T)} \in I^{T+1} \right\} \quad (1)$$

Assume that such x^* is unique, the representation of equation (1) is equivalent to

$$x^* = \arg \max_{x^{(T)}} \left\{ \mathcal{P}(\mathbf{X}^{(T)} = x^{(T)}) \right\} \quad (2)$$

Now we want to show that the sequence x^* can be found recursively.

First, we know the (a priori) probability of a deterministic path can be computed by

$$\begin{aligned} \mathcal{P}(\mathbf{X}^{(n)} = x^{(n)}) &= \mathcal{P}(X_0 = x_0, X_1 = x_1, \dots, X_n = x_n) \\ &= \mathcal{P}(X_0 = x_0) \mathcal{P}(X_1 = x_1 \mid X_0 = x_0) \cdots \mathcal{P}(X_n = x_n \mid X_0 = x_0, \dots, X_{n-1} = x_{n-1}) \\ &= \mathcal{P}(X_0 = x_0) \mathcal{P}(X_1 = x_1 \mid X_0 = x_0) \cdots \mathcal{P}(X_n = x_n \mid X_{n-1} = x_{n-1}) \quad (\text{by Markov assumption}) \\ &= \mathcal{P}(X_0 = x_0) p[x_0, x_1] \cdots p[x_{n-1}, x_n] \end{aligned}$$

Before we move forward, we need a convenient notation.

For $n \in \mathbb{N}_0$ and $i \in I$, define $y(n, i)$ and $a(n, i)$.

- $y(n, i) = (y_0, y_1, \dots, y_{n-1}, i) \in I^{n+1}$ denotes the state sequence satisfies

$$\mathcal{P}(\mathbf{X}^{(n)} = y(n, i)) = \max \left\{ \mathcal{P}(\mathbf{X}^{(n-1)} = x^{(n-1)}, X_n = i) \mid x^{(n-1)} \in I^n \right\}$$

- $a(n, i)$ denotes the probability $\max \left\{ \mathcal{P}(\mathbf{X}^{(n-1)} = x^{(n-1)}, X_n = i) \mid x^{(n-1)} \in I^n \right\}$

Given $n \in \mathbb{N}_0$ and $i \in I$, since the state-space is finite, $y(n, i)$ is guarantee to exist.

Though $y(n, i)$ may not be unique, all of them would have the same probability, that is, $a(n, i)$ is unique. Therefore, we can assign any sequence $x^{(n)} \in I^{n+1}$ with $x_n = i$ satisfying $\mathcal{P}(\mathbf{X}^{(n)} = x^{(n)}) = a(n, i)$ as $y(n, i)$.

Based on the knowing of $y(n, j) \forall j \in I$, we want to obtain $y(n+1, i)$.

We have

$$\mathcal{P}(X_0 = x_0, X_1 = x_1, \dots, X_{n+1} = x_{n+1}) = \mathcal{P}(X_0 = x_0, X_1 = x_1, \dots, X_n = x_n) p[x_n, x_{n+1}] \quad (3)$$

For $i \in I$, by definition, $y(n+1, i)$ satisfies

$$\mathcal{P}(y(n+1, i)) = \max \left\{ \mathcal{P}(\mathbf{X}^{(n+1)} = x^{(n+1)} \mid x^{(n+1)} = (x_0, x_1, \dots, x_n, i) \right\} \quad (4)$$

Since every path 123

$$\begin{aligned} \mathcal{P}(y(n+1, i)) &= \max \left\{ \mathcal{P}(\mathbf{X}^{(n)} = x^{(n)}, X_{n+1} = i \mid x^{(n+1)} \in I^{n+1}) \right\} \\ &= \max \left\{ \mathcal{P}(\mathbf{X}^{(n-1)} = x^{(n-1)}, X_n = j, X_{n+1} = i \mid x^{(n-1)} \in I^n, j \in I) \right\} \quad \text{by (3)} \\ &= \max \left\{ \mathcal{P}(\mathbf{X}^{(n-1)} = x^{(n-1)}, X_n = j) \mathcal{P}(X_{n+1} = i \mid X_n = j) \mid x^{(n-1)} \in I^n, j \in I \right\} \quad (\text{by Markov assumption}) \\ &= \max \left\{ \mathcal{P}(\mathbf{X}^{(n-1)} = x^{(n-1)}, X_n = j) p[j, i] \mid x^{(n-1)} \in I^n, j \in I \right\} \quad (5) \end{aligned}$$

Since we already know the sequence to maximize $\mathcal{P}(\mathbf{X}^{(n-1)} = x^{(n-1)}, X_n = j)$ for each $j \in I$, we do not need to go through all I^{n+1} possible cases. We can search among the cases conditioning on the last state n . Hence,

$$\begin{aligned} \mathcal{P}(y(n+1, i)) &= (5) \\ &= \max \left\{ \mathcal{P}(\mathbf{X}^{(n)} = y(n, j)) p[j, i] \mid j \in I \right\} \quad (6) \end{aligned}$$

Now, we need a starting point to complete the recursive algorithm.

Since we have the initial distribution of the chain, the algorithm is complete, as follow:

Initialization:

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for  $i \in I$  do
   $y(0, i) \leftarrow \lambda_i$ 
end for
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Recursion:

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for  $n = 1$  to  $n = T$  do
  for  $i \in I$  do
     $k \leftarrow \arg \max_j \left\{ \mathcal{P}(\mathbf{X}^{(n)} = y(n, j)) p[j, i] \mid j \in I \right\}$ 
     $y(n, i) \leftarrow (y(n-1, k), i)$ 
  end for
end for
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Termination:

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 $x^* \leftarrow \arg \max_{y(T, i)} \left\{ \mathcal{P}(\mathbf{X}^{(T)} = y(T, i)) \mid i \in I \right\}$ 
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3 Hidden Cases

In a hidden Markov chain, we can never know the state sequence even when the chain has terminated.

Hence there is a small difference between the problem we are discussing in these two situations.

- In the non-hidden cases, we try to predict the most probable result before we would see it.
- In hidden cases, we try to figure out what state sequence has gone through after we have seen some outcome.