

Few Terms that you shall see/hear throughout the course:

PSP (Problem Solving Percentage) - Solved Assignment Problems / Total

Open Assignment Problems

- There are two types of section - Assignment and Additional. Assignment section consists of implementation of the problems done in class. PSP is calculated based on only Assignment Problems.
- Additional Problems are slight modifications of assignment problem, they are not part of PSP but once you're done with assignment, we highly recommend to complete additional problems as well.
- Try to keep PSP least 85% no matter what. It shall really help you to stay focused and we have seen in the past that people with $\geq 85\%$, do well in Interviews.

Attendance

- Try to maintain at-least 75% attendance either through live classes or by watching recording.
- Though I will recommend you to come to classes regularly because otherwise it may create backlogs.
- So, I expect all of you to attend live classes and if for any reason you are unable to, then please send me a message stating the reason.

Intermediate Module Description

- Introduction to Problem Solving
- Time Complexity
- Introduction to Arrays
- Prefix Sum
- Carry Forward
- Subarrays
- 2D Matrices
- Sorting Basics
- Hashing Basics
- Strings Basics
- Bit Manipulation Basics
- Interview Problems
- Contest [covers Full Intermediate DSA]

Note:

In Intermediate, we shall be learning the concepts around different topics and how to work with certain data structures.

- This module is dedicated to make you comfortable with Programming.

Contest will be organised after Intermediate Module.

- It'll will be for 1.5 hours and will be conducted within class duration followed by Contest Discussion (Instructor shall be discussing contest problems).
- It'll consist of 3 questions and we expect you to solve ≥ 2 problems. If for any reason you are unable to solve, then we shall also be having re-attempts as well.(We'll provide more info on re-attempts moving forward)
- Contests are critical to retaining what you have learnt and measuring where you need improvement. Please take contests seriously.

Be consistent in solving problems. If stuck, please post the issue in your WA/Slack group and let's make it a habit of helping each other as it will eventually help you to be better.

FAQs :

- Notes will be uploaded after the class.
- Assignments will be unlocked after the class ends.
- There is no deadline for assignments.
- If asking a question, ask in public chat.
- If answering a question, answer in private chat.

Content

- Count the factors
- Optimisation for count the factors
- Check if a number is prime
- Sum of N natural numbers
- Definition of AP and GP
- How to find the no. of times a code runs
- How to compare two algorithms.

What is a factor ?

N is number

$\frac{N}{x}$ is an integer

x

$\rightarrow x$ is a factor of N .

24

is 2 a factor ?

Remainder when 24 is divided by 2 = 0

How to check for a factor programmatically

$N \% x == 0$ { x is a factor of N }

Count factors of a given number N $N > 0$

$N = 24$ 1 2 3 4 6 8 12 24 = 8

$N = 10$ 1 2 5 10 = 4

$N = 4$ 1 2 4 = 3

What is the minimum value of a factor ?

1

What is the maximum value of a factor ?

N

Pseudocode

```
count = 0
for i → 1 to N {
    // check if i is a factor of N
    if (N % i == 0) {
        count += 1
    }
}
print(count)
```

iterations = N

what is iteration ?

No. of times a for loop runs

Any machine takes 1 sec to process 10^8 iterations

| N | iteration | time |
|--------|-----------|--------|
| 10^8 | 10^8 | 1 sec |
| 10^9 | 10^9 | 10 sec |

10^8 iterations → 1 sec
1 iteration → $1/10^8$ sec.
 10^9 iteration → ?
$$\frac{10^9 \times 1}{10^8} = 10$$

10^9 iteration = $10 * 10^8$ iterations
 $10 * 1$ sec
10 sec

| N | iteration | time |
|-----------|-----------|----------------------------------|
| 10^8 | 10^8 | 1 sec |
| 10^9 | 10^9 | 10 sec |
| 10^{18} | 10^{18} | 10^{10} sec. \approx 317 yrs |

$$\begin{aligned}
 10^8 \text{ iterations} &\longrightarrow 1 \text{ sec} \\
 1 \text{ iteration} &\longrightarrow 1/10^8 \text{ sec.} \\
 10^{18} \text{ iteration} &\longrightarrow \frac{10^{18} \times 1}{10^8} = 10^{10}
 \end{aligned}$$

$$a^b * a^c = a^{b+c}$$

$$\frac{a^b}{a^c} = a^{b-c}$$

Optimization for counting factors.

$$a * b = N$$

a is a factor of N

b is also a factor of N

$$N = 24$$

$$a \quad b \quad a < b$$

| | | | |
|---|---|----|------|
| 1 | * | 24 | True |
| 2 | * | 12 | True |
| 3 | * | 8 | True |
| 4 | * | 6 | True |

| | | | |
|----|---|---|-------|
| 6 | * | 4 | False |
| 8 | * | 3 | |
| 12 | * | 2 | |
| 24 | * | 1 | |

Repetition of factors

$$a * b = N \longrightarrow N/a$$

→ if a is a factor of N

$b = \frac{N}{a}$ is also a factor of N

Pseudocode

```

count = 0
for (a = 1 ; a < b ; a++) {
    // if a is a factor
    if (N % a == 0) {
        // ⇒ b is also a factor ie N/a
        count += 2
    }
}
print(count)

```

$a < \frac{N}{a} \longrightarrow a * a < N$

| $N = 4$ | a | b | $a \leq b$ |
|---------|-----|-----|------------|
| | 1 | 4 | True |
| | 2 | 2 | True |
| | 4 | 1 | False |

If $a == b$ then increment the count by 1


```

count = 0
for (a = 1 ; a*a <= N ; a++) {
    // if a is a factor
    if (N % a == 0) {
        // => b is also a factor ie N/a
        b = N/a
        if (a == b) count += 1
        else count += 2
    }
}
print(count)

```

$a \leq \sqrt{N}$

iterations = \sqrt{N}

| | | | |
|----------|---|----|-------|
| $N = 12$ | a | b | count |
| | 1 | 12 | 2 |
| | 2 | 6 | 4 |
| | 3 | 4 | 6 |
| | 4 | | |

since $a*a > 12$

$N = 12$ 1 2 3 4 6 12

| | | | |
|----------|---|----|-------|
| $N = 36$ | a | b | count |
| | 1 | 36 | 2 |
| | 2 | 18 | 4 |
| | 3 | 12 | 6 |
| | 4 | 9 | 8 |

$$7 * 7 > 36$$

Assumption \longrightarrow It takes 1sec to process 10^8 iterations

| | | |
|-----------|-------------------------|---------|
| N | iteration | time |
| 10^{18} | $\sqrt{10^{18}} = 10^9$ | 10 secs |

317 yrs \longrightarrow 10 sec.

Break \rightarrow 22:38

Given N , check if N is prime or not?

A prime no. has exactly two factors

10 11 23 2 25 27 31

```
bool isPrime (int N) {
    if (countFactors(N) == 2) {
        return true
    }
    else {
        return false
    }
}
```

Sum of all no. from 1 to 100

$$\frac{n*(n+1)}{2}$$

$$S = 1 + 2 + \dots + 99 + 100$$

$$S = 100 + 99 + \dots + 2 + 1$$

$$2*S = \underbrace{101 + 101 + \dots + 101 + 101}_{100 \text{ times}}$$

$$2S = 101 * 100$$

$$S = \frac{101 * 100}{2} = 5050$$

$$S = 1 + 2 + 3 + \dots + n-1 + n$$

$$S = n + n-1 + \dots + 2 + 1$$

$$2S = \underbrace{(n+1) + (n+1) + \dots + (n+1) + (n+1)}_{n \text{ times}}$$

$$2S = (n+1) * (n)$$

$$S = \frac{(n+1) * (n)}{2}$$

Basic Math Concepts

$[1, 5]$ \longrightarrow All the values from 1 to 5 inclusive of 1 and 5

$[1, 5)$ \longrightarrow All the values from 1 to 4 include 1 and exclude 5

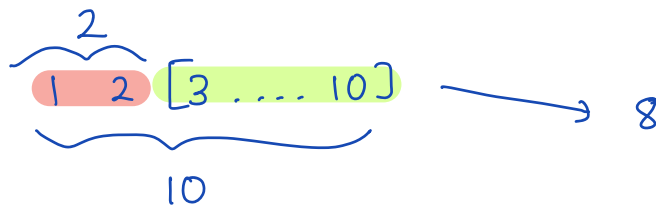
$[a, b]$ \longrightarrow All values from a to b, including a and b

(a, b) \longrightarrow All values from a to b, excluding a and b

$[3, 10]$ \longrightarrow 3 4 5 6 7 8 9 10

$[1, 10]$ \longrightarrow 1 10 { 10 values }

$[1, 2]$ \longrightarrow 1 2 { 2 values }



$$[a, b] \longrightarrow \overbrace{b} - (a-1) = b - a + 1$$
$$[1, b] - [1, a-1]$$

$$[2, 3] \longrightarrow \begin{array}{c|c|c} \cancel{b-a} & \cancel{b-a-1} & b-a+1 \\ \cancel{1} & \cancel{0} & 2 \end{array}$$

Calculate the iterations for below code

```
for (i = 1; i <= N; i++) {  
    |  
    if (i == N) break  
}
```

iterations = N

```
for (i = 0; i <= 100; i++) {  
    |  
    s = s + i + i^2  
}
```

→ [0, 100]

↓
 $100 - 0 + 1 = 101$

iterations 101

```
for (i = 1; i <= N; i++) {  
    |  
    if (i % 2 == 0) print(i)  
}
```

N

```
for (i = 1; i <= M; i++) {  
    |  
    if (i % 2 == 0) print(i)  
}
```

M

total
N + M

Geometric Progression

$$\begin{array}{ccccccc} & a & & & & & \\ & 5 & 10 & 20 & 40 & 80 & \\ & \underbrace{\quad} & \underbrace{\quad} & \underbrace{\quad} & \underbrace{\quad} & & \\ & \frac{10}{5} & \frac{20}{10} & \frac{40}{20} & \frac{80}{40} & & \\ r & 2 & 2 & 2 & 2 & & \end{array}$$

$$r = \frac{A_{i+1}}{A_i}$$

$a \longrightarrow$ first term of GP

$r \longrightarrow$ common ratio

Sum of a GP

$$a, a*r, a*r^2, \dots, a*r^{n-1}$$

$$\text{sum of } n \text{ terms of GP} = \frac{a * (r^n - 1)}{r - 1}$$

$$r \neq 1$$

if $r = 1$ then $\text{sum} = a*n$

Aswin

Algo 1

15 sec



windows XP



mac M3 pro

7 sec



C++



Hot environment
rajath



manali

5 sec

Suyash

Algo 2

10 sec



mac M3 pro

10 sec



python



C++

5 sec



cold environment
manali

5 sec

Conclusion → Running time of algo depends on lot of factors

⇒ It's always better to compare the no. of iterations

```
for (i = L; i <= N; i++) {  
    |   if (i == N) break  
    |  
    }  
    3
```

iterations = N