

- Introduction to Prime Numbers
- Get all primes 1 to n
- Smallest Prime Factor for 2 to n
- Prime Factorization
- Get the no. of factors / divisors

Prime Numbers

Numbers that have 2 factors \rightarrow 1 and itself.

eg:- 2, 3, 5, 7, 11...

Q1) Given a number n , check if it is prime or not.

$n = 3 \rightarrow \text{True}$

$n = 6 \rightarrow \text{False}$

```
boolean checkPrime(n) {  
    count = 0  
    for (i = 1; i * i <= n; i++) {  
        if (n % i == 0) {  
            count++  
            if (i != n / i) {  
                count++  
            }  
        }  
    }  
    return (count == 2)  
}
```

Alt. approach

```
if (n < 2)  
    return false  
for (i = 2; i * i <= n; i++) {  
    if (n % i == 0)  
        return false  
}  
return true
```

$O(\sqrt{n})$ TC

$O(1)$ S.C.

Q2) Given n , print all the prime nos. from 1 to n .

$n=10 \rightarrow 2, 3, 5, 7$

$n=20 \rightarrow 2, 3, 5, 7, 11, 13, 17, 19.$

```
void printAllPrime(n){  
    for(i → 2 to n){  
        boolean isPrime = true  
        for(j = 2; j * j ≤ i; j++){  
            if(i % j == 0){  
                isPrime = false  
                break  
            }  
        }  
        if(isPrime)  
            print(i)  
    }  
}
```

$$\begin{aligned} & \sqrt{2} + \sqrt{3} + \sqrt{4} + \dots + \sqrt{n} \\ &= O(n\sqrt{n}) \text{ T.C.} \\ & O(1) \text{ S.C.} \end{aligned}$$

Sieve of Eratosthenes

- Assume that all nos. 2 to n are prime
- Start with the first prime number and mark all its multiples as non-prime.
- Move to the next unmarked no. (not marked in red), this is prime, mark all its bigger multiples as non-prime.
- Repeat the process till the end of array
- All unmarked nos. are prime.

1	2	3	4	5	6	7	8	9	10
11	12	13	14	15	16	17	18	19	20
21	22	23	24	25	26	27	28	29	30
31	32	33	34	35	36	37	38	39	40
41	42	43	44	45	46	47	48	49	50

i	# iterations
2	$n/2$
3	$n/3$
4	\times
5	$n/5$
6	\times
7	$n/7$
8	\times

$$\begin{aligned}
 & \frac{n}{2} + \frac{n}{3} + \frac{n}{5} + \frac{n}{7} + \frac{n}{11} + \frac{n}{13} + \dots \\
 &= n \left(\frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \frac{1}{7} + \frac{1}{11} + \frac{1}{13} + \dots \right) \\
 & \quad \text{sum of reciprocals of prime nos. upto } n. \\
 &= O(n (\log(\log n))) \text{ T.C.}
 \end{aligned}$$

printAllPrimes(n) {

boolean isPrime[n+1] = {true}

isPrime[0/1] → false

for (i = 2; i * i ≤ n; i++) {

if (isPrime[i]) {

for (j = i * i; j ≤ n; j += i) {

isPrime[j] = false

}

}

}

for (i = 2; i ≤ n; i++) {

if (isPrime[i]) {

print(i)

}

}

}

$O(n \log(\log n))$ T.C

$O(n)$ S.C

7
 7 * 2 → 2
 7 * 3 → 3
 7 * 4 → 2
 7 * 5 → 5
 7 * 6 → 2, 3
 7 * 7 →
 7 * 8
 7 * 9

for (i = 2; i ≤ a; i++) {
 if (isPrime[i]) {
 for (j = i * 2; j ≤ b; j++) {
 ...
 }

}

i = 2 → j → b/2

i = 3 → j → b/3

i = 5 → j → b/5

$$\frac{b}{2} + \frac{b}{3} + \frac{b}{5} + \dots + \frac{b}{\sqrt{a}}$$

$$= b \left(\frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \dots + \frac{1}{\sqrt{a}} \right)$$

$$= b \log(\log a)$$

$$a = \sqrt{n}, b = n$$

$$n \log(\log \sqrt{n}) = O(n \log(\log n))$$

$$a = n, b = n$$

$$n \log(\log n)$$

[Break till 10:30 PM]

Q3) Given n , return the smallest prime factor for each no. from 2 to n .

$n=10$

	2	3	4	5	6	7	8	9	10
ans →	2	3	2	5	2	7	2	3	2

1	2	3	4	5	6	7	8	9	10
11	12	13	14	15	16	17	18	19	20
21	22	23	24	25	26	27	28	29	30
31	32	33	34	35	36	37	38	39	40
41	42	43	44	45	46	47	48	49	50

$\begin{matrix} i \\ i \end{matrix} \} \text{ Prime No.}$

```

fn smallestPrimeFactor(n) {
    int spf[n+1]
    for (i → 2 to n)
        spf[i] = i
    for (i = 2; i * i ≤ n; i++) {
        if (spf[i] == i) {
            for (j = i * i; j ≤ n; j += i) {
                if (spf[j] == j)
                    spf[j] = i
            }
        }
    }
    return spf
}

```

$O(n \log(\log n))$ T.C
 $O(n)$ S.C

Prime Factorization

$$n=48$$

$$\begin{array}{r} 2 \overline{)48} \\ 2 \overline{)24} \\ 2 \overline{)12} \\ 2 \overline{)6} \\ 3 \overline{)3} \\ 1 \end{array}$$

$$2^4 * 3^1$$

$2^0, 2^1, 2^2, 2^3, 2^4$ (5 terms) and $3^0, 3^1$ (2 terms) = 10.

$$48 = 2 * 2 * 2 * 2 * 3$$

$$= 2^4 * 3^1$$

$$\text{Nr. of divisors} = (4+1) * (1+1)$$

$$= 5 * 2 = 10.$$

$$1, 2, 3, 4, 6, 8, 12, 16, 24, 48.$$

10 divisors.

$$2^0 * 3^0 = 1$$

$$2^0 * 3^1 = 3$$

$$2^1 * 3^0 = 2$$

$$2^1 * 3^1 = 6$$

$$2^2 * 3^0 = 4$$

$$\vdots$$

$$n=300$$

$$\begin{array}{r} 2 \overline{)300} \\ 2 \overline{)150} \\ 3 \overline{)75} \\ 5 \overline{)25} \\ 5 \overline{)5} \\ 1 \end{array}$$

$$300 = 2^2 * 3^1 * 5^2$$

$$2^{0-2} * 3^{0-1} * 5^{0-2}$$

$$\# \text{ divisors} = (2+1) * (1+1) * (2+1)$$

$$= 18.$$

$$n=560$$

$$\begin{array}{r} 2 \overline{)560} \\ 2 \overline{)280} \\ 2 \overline{)140} \\ 2 \overline{)70} \\ 5 \overline{)35} \\ 7 \overline{)7} \\ 1 \end{array}$$

$$560 = 2^4 * 5^1 * 7^1$$

$$d \rightarrow 2^{\leq 4} * 5^{\leq 1} * 7^{\leq 1}$$

$$(4+1) * (1+1) * (1+1) = 5 * 2 * 2 = 20.$$

$$n = p_1^{a_1} * p_2^{a_2} * p_3^{a_3} * \dots * p_k^{a_k}$$

Primes

$$\Rightarrow \# \text{ factors of } n = (a_1+1) * (a_2+1) * (a_3+1) * \dots * (a_k+1).$$

$$20 = 2^2 * 5^1$$

$$(2+1) * (1+1) = 6 \text{ factors.}$$

$$n = 360$$

$$\text{spf}[360] = 2$$

$$\downarrow /2$$

$$\text{spf}[180] = 2$$

$$\downarrow /2$$

$$\text{spf}[90] = 2$$

$$\downarrow /2$$

$$\text{spf}[45] = 3$$

$$\downarrow /3$$

$$\text{spf}[15] = 3$$

$$\downarrow /3$$

$$\text{spf}[5] = 5$$

$$\downarrow /5$$

$$1 \quad \checkmark$$

$$2^3 * 3^2 * 5^1$$

$$(3+1) * (2+1) * (1+1)$$

$$= 4 * 3 * 2 = 24.$$

```

fn factors(n) {
    ans = 1
    s = spf[n]
    while (n > 1) {
        cnt = 0
        while (n % s == 0) {
            n = n / s
            cnt += 1
        }
        ans = ans * (cnt + 1)
        s = spf[n]
    }
    return ans
}

```

$O(\log n)$ T.C.
after computing
shf[].

$$O(n \log(\log n)) \quad + \quad O(\log n)$$

for calculating
spf

for finding
factors of n .

Find count of factors for m diff. values, each upto 10^6 , $m \rightarrow 10^5$. using

$$[O(\text{max} * \log(\log \text{max})) + n * O(\log \text{max})]$$

for calculating shf for finding # factors of n.

using sqrt
approach
↓
 $O(m * \sqrt{max})$
T.C

$$= O(\max \cdot \log(\log(\max)) + \underbrace{m \cdot \log(\max)}_{T.C})$$