

Booth's Multiplication Algorithm

Worked Example: Multiply -3×2
(4-bit)

Intuition Behind Booth's Algorithm

- Handles signed binary multiplication efficiently using 2's complement.
- Reduces operations for runs of 1s in multiplier.
- Uses a combination of add/subtract and arithmetic right shift.
- Based on detecting transitions in Q_0 and Q_{-1} .

Booth's Algorithm Steps

- 1. Initialize $A = 0$, $Q = \text{multiplier}$, $Q-1 = 0$, $M = \text{multiplicand}$.
- 2. For n bits:
 - a. If $Q_0Q_{-1} = 10$: $A = A - M$
 - If $Q_0Q_{-1} = 01$: $A = A + M$
 - b. Perform arithmetic right shift on $[A, Q, Q-1]$
- 3. Final result is in $[A, Q]$ (concatenated).

Initial Setup (4-bit Example)

- $M = -3 \rightarrow 1101$
- $Q = 2 \rightarrow 0010$
- $A = 0000, Q-1 = 0$
- $-M$ (2's complement of 1101) = 0011

Step 1

- $Q_0 Q_{-1} = 00 \rightarrow$ No Operation
- Right shift $[A|Q|Q-1] = 0000\ 0010\ 0$
- Result: $A = 0000$, $Q = 0001$, $Q-1 = 0$

Step 2

- $Q_0Q_{-1} = 10 \rightarrow A = A - M \rightarrow 0000 - 1101 = 0011$
- Right shift $[A|Q|Q-1] = 0011\ 0001\ 0$
- Result: $A = 0001, Q = 1000, Q-1 = 1$

Step 3

- $Q_0Q_{-1} = 01 \rightarrow A = A + M \rightarrow 0001 + 1101 = 1110$
- Right shift $[A|Q|Q-1] = 1110\ 1000\ 1$
- Result: $A = 1111, Q = 0100, Q-1 = 0$

Step 4

- $Q_0 Q_{-1} = 00 \rightarrow$ No Operation
- Right shift $[A|Q|Q-1] = 1111\ 0100\ 0$
- Result: $A = 1111$, $Q = 1010$, $Q-1 = 0$

Final Result

- Final $[A|Q] = 11111010$ (8-bit 2's complement)
- \rightarrow 2's complement of $11111010 = 00000110$
 $\rightarrow 6$
- \rightarrow Final result = -6
-  Verified: $-3 \times 2 = -6$