

# Signed Binary Representations and Arithmetic

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# Motivation

- Computers use binary to represent data and perform arithmetic.
- To handle **positive and negative numbers**, we use signed representations.
- Three main systems:
  - ① Sign-Magnitude
  - ② 1's Complement
  - ③ 2's Complement
- We will discuss how they represent numbers and perform arithmetic operations.

# Sign-Magnitude Representation

- MSB is the **sign bit**:
  - 0 = Positive
  - 1 = Negative
- Remaining bits represent magnitude (as unsigned binary).
- **Example (4 bits)**:
  - $+5 = 0101, -5 = 1101$

## Issues:

- Two zeros: 0000 (+0), 1000 (-0)
- Arithmetic requires extra logic for comparing and handling sign.

# 1's Complement Representation

- Positive numbers: normal binary
- Negative numbers: invert (flip) all bits of the positive number
- **Example (4 bits):**
  - $+5 = 0101$
  - $-5 = 1010$  (1's complement of 0101)

## Features:

- Two zeros: 0000 (+0), 1111 (-0)
- Subtraction becomes addition by complementing
- **End-around carry** must be added back

# 2's Complement Representation

- Positive numbers: unchanged
- Negative numbers: flip all bits and add 1
- **Example (4 bits):**
  - $+5 = 0101$
  - $-5: \text{flip} \rightarrow 1010, \text{add } 1 \rightarrow 1011$

## Advantages:

- Only one zero (0000)
- Simple arithmetic – same circuit for add/sub
- Most widely used method in computers today

# Sign-Magnitude Arithmetic

## Rules:

- Same signs → add magnitudes, keep the sign
- Different signs → subtract smaller from larger, use larger's sign

## Example: $-5 + 3$

- $-5 = 1101$ ,  $+3 = 0011$
- Magnitude:  $5 - 3 = 2$
- Result:  $-2 = 1010$

**Note:** Requires magnitude comparison and conditional logic

# 1's Complement Arithmetic

## Rules:

- Add both numbers (including sign bits)
- If there is a carry from MSB, add it back (end-around carry)

**Example:**  $-3 + 2$  (4-bit)

- $+2 = 0010$
- $-3 = 1100$  (1's complement of 0011)
- Add:  $0010 + 1100 = 1110$  (no carry)
- Result:  $1110 = -1$

**Note:** Still has two zeros. End-around carry must be handled.

# 1's Complement Arithmetic (Example 2 with End-Around Carry)

**Example 2:**  $-3 + 4$  (4-bit)

- $+4 = 0100$
- $-3 = 1100$  (1's complement of 0011)
- Add:  $0100 + 1100 = 10000$  (5 bits: carry out is 1)
- End-around carry:  $0000 + 1 = 0001$
- Result:  $0001 = +1$

**Note:** End-around carry must be added back for correct result.

# 1's Complement Arithmetic (Example 3 with End-Around Carry)

**Goal:** Add  $-2 + (-1)$  using 4-bit 1's complement

- $-2$ :
  - $+2 = 0010 \rightarrow 1\text{'s complement} = \mathbf{1101}$
- $-1$ :
  - $+1 = 0001 \rightarrow 1\text{'s complement} = \mathbf{1110}$
- Add:  $1101 + 1110 = \mathbf{1 1011}$  (5 bits; carry = 1)
- End-around carry:  $1011 + 1 = \mathbf{1100}$
- Result:  $\mathbf{1100} = -3$  (1's complement of 0011)

**Note:** End-around carry must be added back. MSB = 1 indicates a negative result.

# 1's Complement: Circular Number Line (Part 1)

**Why Circular?** In 1's complement, the number line wraps around — overflow loops back to the start like a circle.

## 4-bit 1's Complement Representations:

Binary	Decimal	Binary	Decimal
0000	+0	1000	-7
0001	+1	1001	-6
0010	+2	1010	-5
0011	+3	1011	-4
0100	+4	1100	-3
0101	+5	1101	-2
0110	+6	1110	-1
0111	+7	1111	-0

**Note:** Two representations for zero: 0000 (+0) and 1111 (-0)

# 1's Complement: Circular Number Line (Part 2)

## Circular Behavior in Action:

- $0111 (+7) + 0001 = 1000 (-7)$  — wraps around after +7
- $1111 (-0) + 0001 = 0000 (+0)$  — end-around behavior

## Why Add Carry Back?

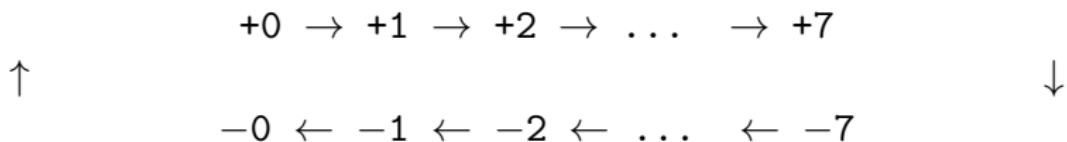
- Carry out from MSB is not overflow — it's a valid part of the result.
- 1's complement uses **modulo** ( $2^n - 1$ ) arithmetic.
- To stay on the circular number line, the carry must be **added back** (end-around carry).

## Conclusion:

1's complement arithmetic only gives correct results when the end-around carry is included.

# 1's Complement: Circular Number Line Visualization

1's Complement wraps like a circle — values move clockwise through positives and wrap to negatives.



## Explanation:

- After +7 (0111), next value is 1000 (-7)
- Carry from MSB wraps to LSB (end-around carry)
- Ensures arithmetic continuity on a circular number line

# Problem with Two Representations of Zero in 1's Complement

**Key Issue:** In 1's complement, zero has **two representations**:

- Positive zero: 0000 0000
- Negative zero: 1111 1111

**Example: Adding +4 and -4 (4-bit representation)**

$$+4 = 0100$$

$$-4 = 1011 \quad (\text{1's complement of } 0100)$$

$$\text{Sum} = 0100 + 1011 = 1111 \quad (\text{Negativezero})$$

**Problems Caused:**

- Ambiguity in checking for zero: 0000 vs 1111
- Comparison logic must account for both
- Requires **end-around carry** in addition
- One representation wasted

# Disadvantages of 1's Complement Representation

- **Two representations of zero:**
  - Positive zero: 0000 0000
  - Negative zero: 1111 1111
  - Leads to ambiguity and extra handling in comparisons.
- **Extra logic for equality checks:**
  - Both 0000 and 1111 must be treated as zero.
- **Addition requires end-around carry:**
  - If there is a carry from the MSB, it must be added to the LSB.
  - Slows down arithmetic operations.
- **Wasted encoding:**
  - One bit pattern wasted due to dual zero representation.
  - Only  $2^n - 1$  distinct integers can be represented.
- **Sign bit interpretation issues:**
  - MSB is 1 for both negative zero and negative numbers.
  - Negative zero looks like a negative number.

# 2's Complement Arithmetic

## Rules:

- Add both numbers directly
- Ignore any final carry out of MSB

**Example:**  $-3 + 2$  (4-bit)

- $+2 = 0010$
- $-3 \rightarrow 3 = 0011 \rightarrow \text{flip} = 1100 \rightarrow \text{add } 1 = 1101$
- Add:  $0010 + 1101 = 1111$
- Result:  $1111 = -1$

**Benefit:** No special handling; works like unsigned addition.

**Note:** In 2's complement, **ignore the carry-out**. No end-around carry is needed, unlike 1's complement.

## 2's Complement Arithmetic (Carry Ignored)

**Example:**  $-5 + 7$  (4-bit)

- $+7 = 0111$
- $-5:$ 
  - $5 = 0101$
  - Flip:  $1010$
  - Add 1: **1011**
- Add:  $0111 + 1011 = \mathbf{1\ 0010}$  (carry out = 1)
- Ignore carry-out  $\rightarrow$  result = **0010**
- Final result: +2

**Note:** Carry-out from MSB is discarded. Result is still correct.

# Comparison Summary

Property	Sign-Magnitude	1's Comp	2's Comp
Two zeros	Yes	Yes	No
Easy arithmetic	No	No	Yes
Used today	No	Rarely	Yes
End-around carry	No	Yes	No
Same HW for add/sub	No	No	Yes