



# ECE291

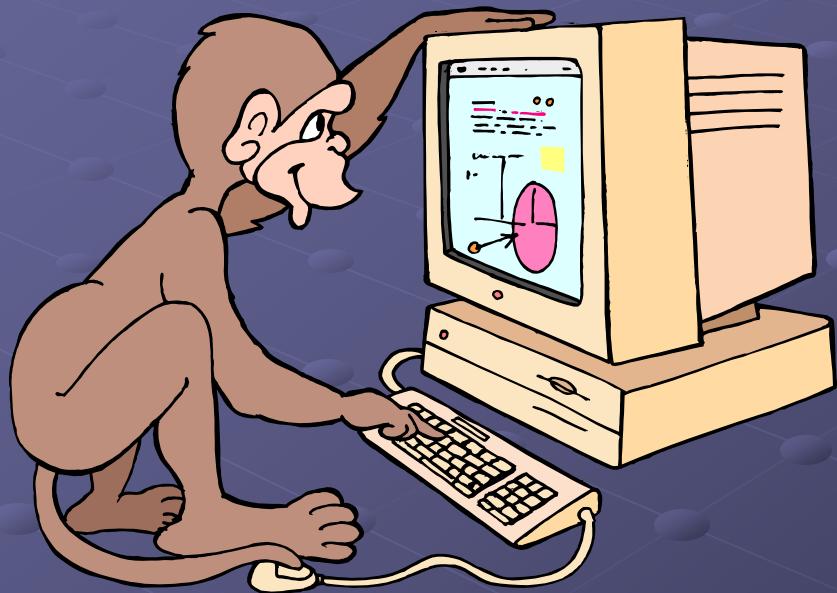
## Lecture 1

### *All about computers*

# Lecture outline

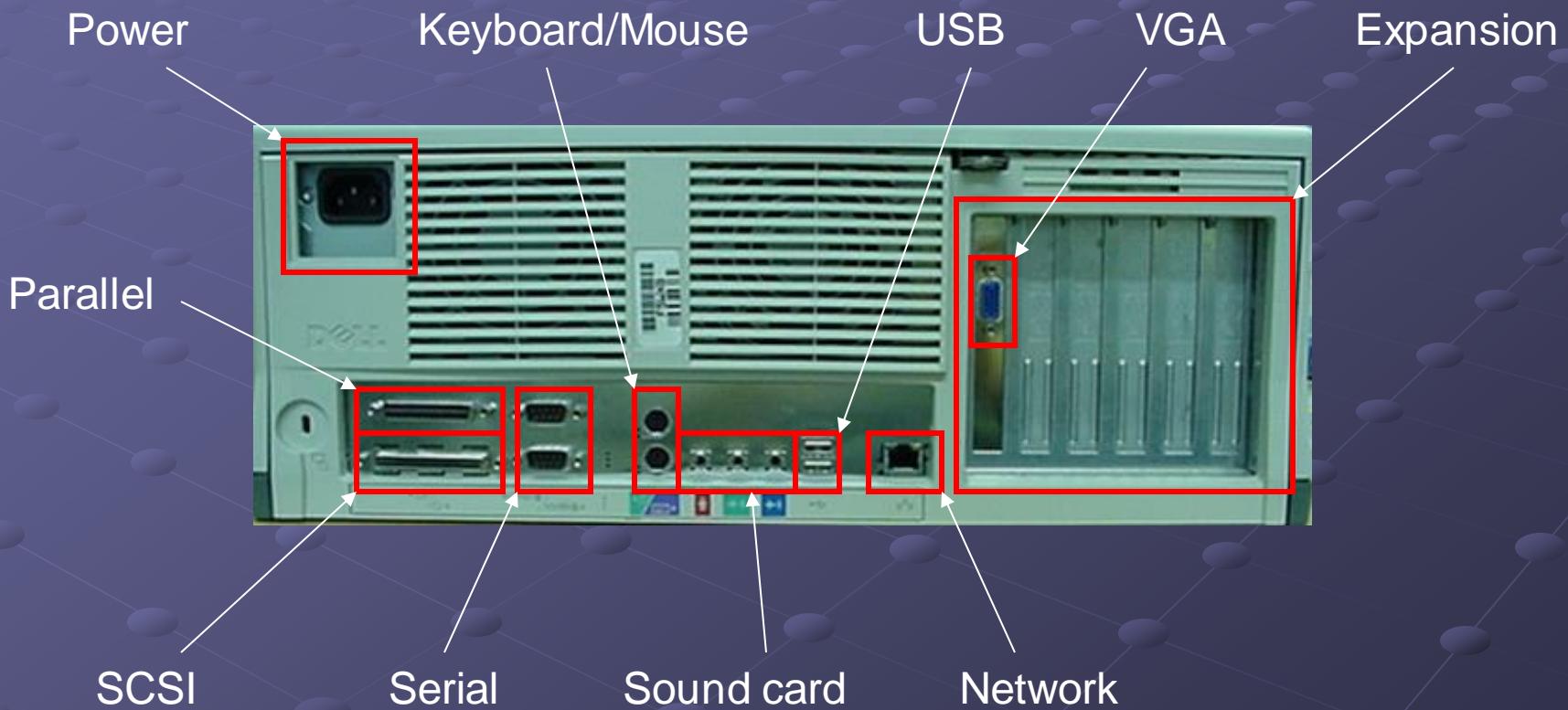
- Top-down view of computers
- Processor architecture
- Registers
- Memory
- Peripherals
- Assembly primer
- Assignments

# Computers from 50,000 feet

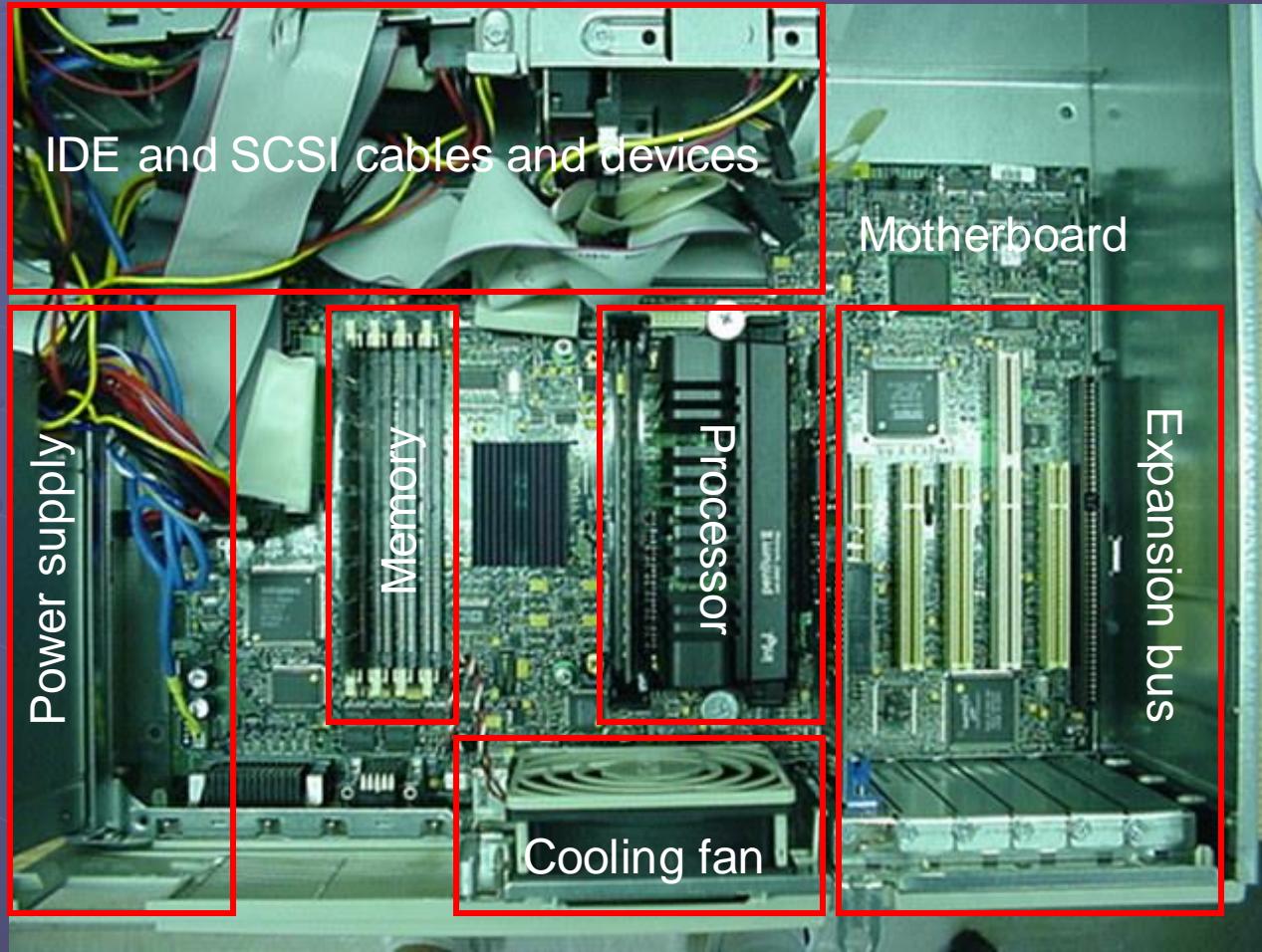


- Basic components
  - Mouse
  - Keyboard
  - Monitor
  - Printer
  - Scanner
  - Joystick
  - “The box”

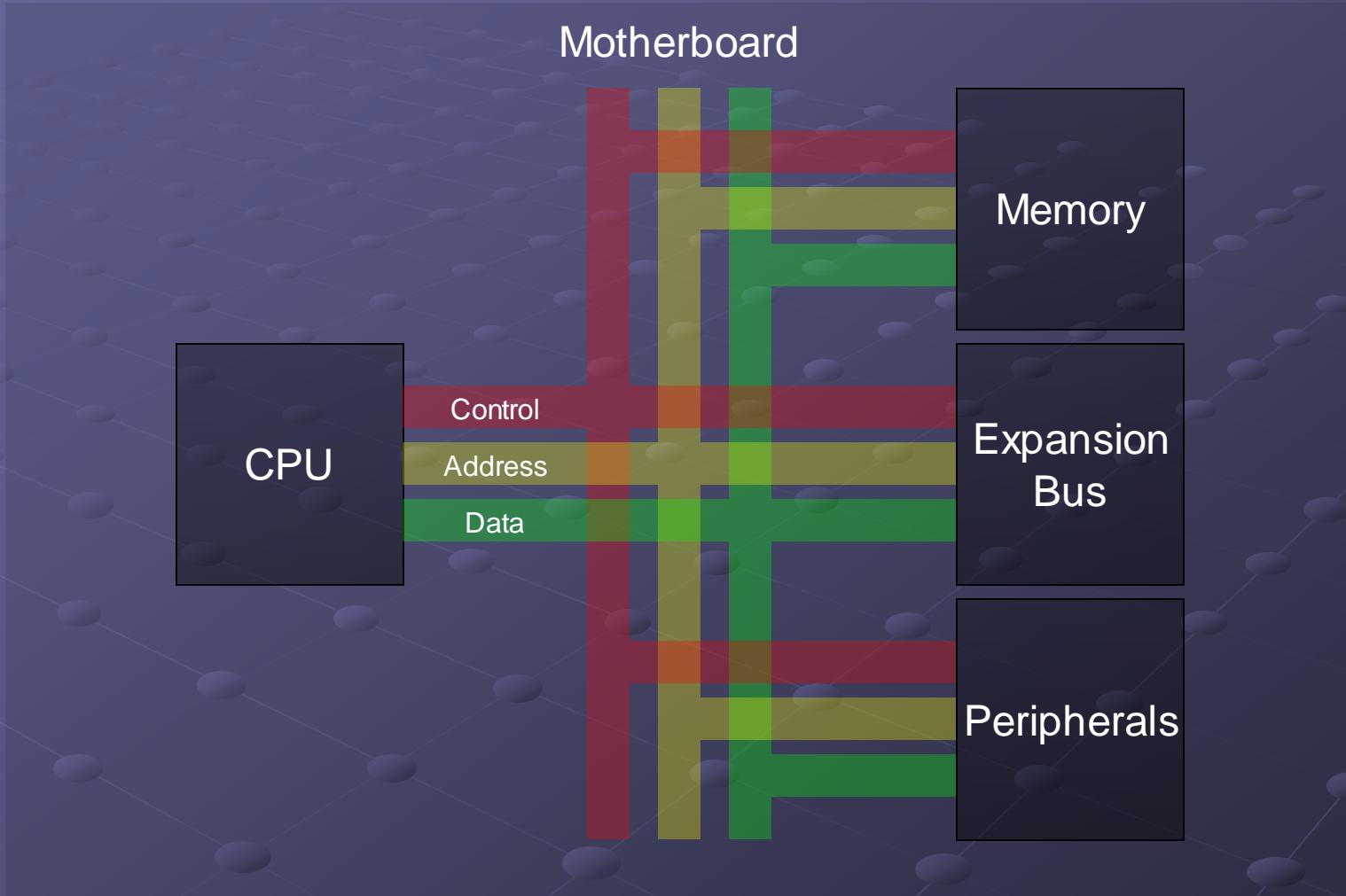
# Computers from 10,000 feet



# Computers from 100 feet



# Computers from 10 feet

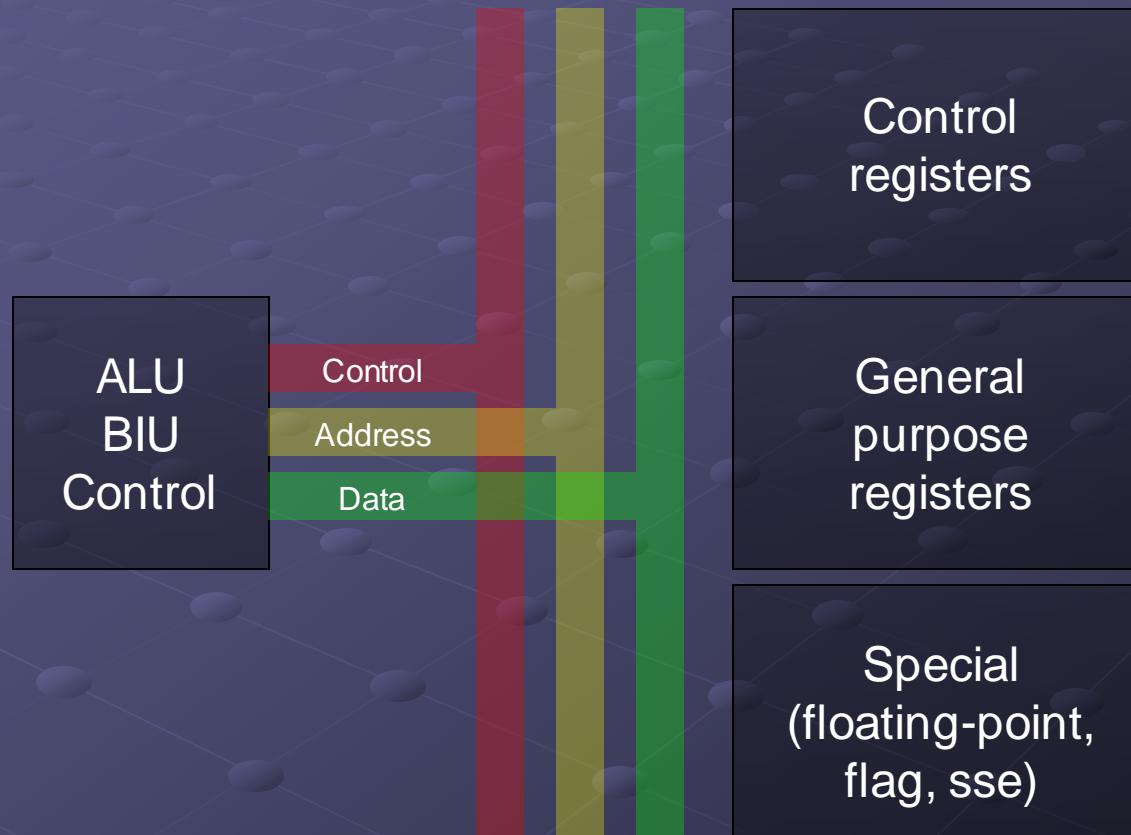


# Computers from 10 feet

Processor	Motherboard			Max Addressable Memory
	Data Bus	Address Bus		
8088	8	20	1,048,576	(1Mb)
8086	16	20	1,048,576	(1Mb)
80286	16	24	16,777,216	(16Mb)
80386dx	32	32	4,294,976,296	(4Gb)
80486	32	32	4,294,976,296	(4Gb)
80586/Pentium (Pro)	64	32	4,294,976,296	(4Gb)

# Computers from 1 micron

## Central Processing Unit – The Processor



# The processor

- It's a state machine
- Has a set of hard-coded operations
  - Each operation has an associated code; an “opcode”
  - Some examples are MOV, ADD, SUB, AND
- All it can do is move data or perform some calculation on data
- Very simple, right?

# The processor - components

- CPU Registers

- Special memory locations constructed from flip-flops and implemented on-chip
- E.g., accumulator, count register, flag register

- Arithmetic and logic Unit (ALU)

- Where most of the action takes place inside the CPU

- Bus Interface Unit (BIU)

- Responsible for controlling the address and data busses when accessing main memory and data in the cache

- Control Unit and Instruction Set

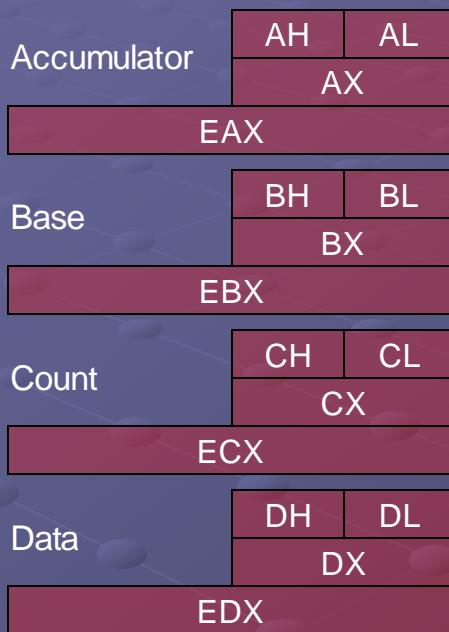
- CPU has a fixed set of instructions
- E.g. MOV, CMP, ADD, JMP

# The registers

- Small memory locations that are quickly accessible by the processor
- Hold data and results for operations
- Point to memory locations
- Contain status information about the results of operations

# 80x86 registers

## General Purpose



## Special Registers



## Index Registers



## Segment Registers



# Register specifics

- Accumulator (AL, AH, AX, EAX)

- Usually stores results from the ALU
- It “accumulates” totals from math operations
- General purpose

- Base (BL, BH, BX, EBX)

- Usually holds addresses – points to memory locations

# Register specifics

- Count (CL, CH, CX, ECX)
  - Counting and looping
  - Certain instructions automatically decrement the count register
- Data (DL, DH, DX, EDX)
  - Usually contains data for instructions
  - Contain the most significant bytes of 16-bit Mul/div's

# Register specifics

- ➊ Stack pointer (SP, ESP)
  - Used by stack operations
  - Usually you don't want to mess with this one
- ➋ Base pointer (BP, EBP)
  - Addresses stack memory
  - Can be used to read subroutine arguments

# Register specifics

- Source index (SI, ESI)

- Often used to address a source array or string data type

- Destination index (DI, EDI)

- Often used to address a destination array or string data type

# Register specifics

- Code segment (CS)

- Points to the area in memory where instructions are stored

- Instruction pointer (IP)

- An index into the code segment that points to the next instruction to be executed

- Data segment (DS)

- Points to the area in memory where data needed by a program is stored

# Register specifics

- Stack segment (SS)

- Points to the area in memory containing the system's FIFO stack. We'll have a whole lecture on this.

- ES, FS, GS

- Extra segment registers available to point to additional data segments should that be necessary.

# Memory

- System memory or RAM holds data and instruction opcodes that are being used by the CPU
- Organized in banks, usually even and odd
- Always byte addressable
- More on all this tomorrow

# Memory access example

CPU wants to read memory at address 12345h



# Memory organization



# Memory organization

- In Real Mode:

- The address bus is 20-bits for memory operations.
- $2^{20} = 1\text{MB}$  = how much memory an x86 processor can address in real mode
- The data bus is 16-bits which means it can transfer two bytes in one operation
- Limited to 16-bit registers

# Memory organization

- In Protected Mode:

- The memory address bus is 32-bits
- $2^{32} = 4\text{GB}$  = the maximum addressable amount of memory in modern x86 processors
- Data bus is also 32-bits meaning four bytes can be transferred in a single operation
- Gain access to the 32-bit extended registers

# Peripherals

- In our context, things like printers, joysticks, and scanners aren't peripherals
- Some examples are parallel, serial and USB ports, the internal timers, interrupt controllers, network interface cards
- Many helper devices inside the processor, on the mother board, or in expansion slots

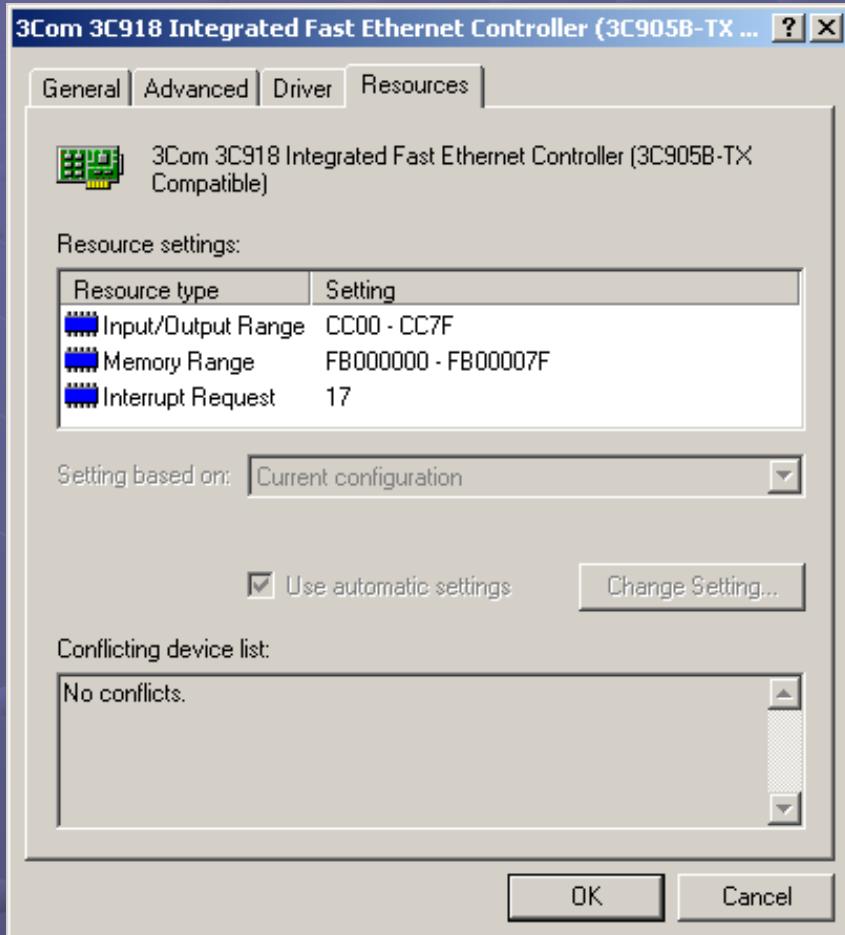
# Peripheral I/O bus

- There is a separate set of address/data busses commonly used for communication with peripheral devices
- In reality, they're the same busses, just different signals in the control bus cause different things to happen
- The I/O address bus is 16-bits wide, meaning you can address 65,535 different I/O ports.
- The I/O data bus is 16-bits

# Peripheral I/O

- Peripheral devices don't *have* to use the I/O bus. They can be memory mapped.
- Video cards are an example
- They have a large amount of memory (VRAM) and that is typically mapped to a range of addresses in the memory space

# Peripheral example



- Network card
- Input/Output Range  
CC00-CC7F
- Memory Range  
FB000000-FB00007F
- This one is both  
memory mapped *and*  
I/O mapped

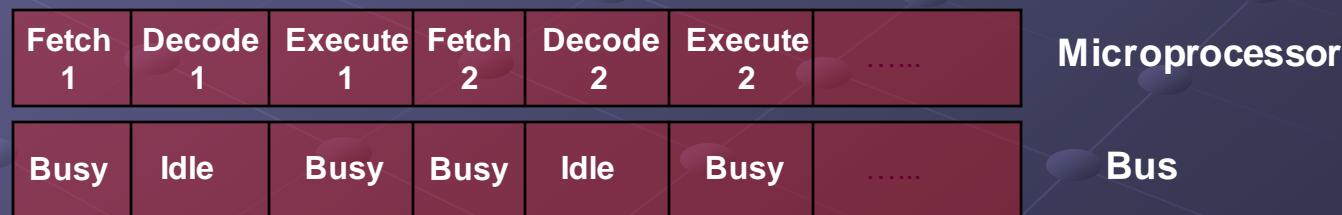
# Peripherals

- All this really means to you is that sometimes you'll use one set of instructions to access a device, and sometimes you'll use a different set.
- We'll talk about this again when we cover serial/parallel ports, timers, interrupts, and the video graphics adapter.

# Instruction processing

- Processing of an instruction by the microprocessor consists of three basic steps
  - fetch instruction from the memory
  - decode the instruction
  - execute (usually involves accessing the memory for getting operands and storing results)

- Operation of an early processor like the Intel 8085



# Instruction processing

- Modern microprocessors can process several instructions simultaneously at various stages of execution
  - this ability is called pipelining
- Operation of a pipelined microprocessor like the Intel 80486



# Assembly primer

- ➊ Everything in your assembly file is one of the following:
  - A comment
  - A label
  - An instruction or directive
  - Data

# Assembly primer - comments

- ; Comments are denoted by semi-colons.
- ; These are comments!!! Yeah!!!!.
- ; Please comment your MP's thoroughly.
- ; It helps us figure out what you were doing
- ; It also helps you figure out what you were doing when you look back at code you
- ; wrote more than two minutes ago.

# Assembly primer

- You might be asking, “What about variables? Why aren’t they on the bulleted list?”
- Variables are nothing more than labels that mark specific memory locations
- Labels can also mark the beginning of procedures or “jump to” locations in your code

# Assembly primer - labels

; Labels can be global

MyGlobalLabel:

MyOtherGlobalLabel

; They can be local too

.MyLocalLabel

; Local labels are good only back to the previous un-dotted label

# Assembly primer - labels

```
MyBigGlobalLabel
```

```
    .local_label1  
    .local_label2
```

```
MyNextBigGlobalLabel
```

```
    .local_label1      ; these are distinct  
    .local_label2      ; to the ones above
```

# Assembly primer - variables

```
; Let's do something with labels now
```

```
VAR1    DB    255
```

```
VAR2    DB    0FFh
```

```
VAR3    DW    1234h
```

```
ARR1    DB    12, 34, 56, 78, 90
```

```
ARR2    DW    12, 34, 56, 78, 90
```

```
STR1    DB    '291 is awesome!!!!', 0
```

```
BUF1    RESB 80
```

labels

directives

data

# Assembly primer - directives

- **EXTERN** – allows you to declare external procedures so you can use them in your program.
- **SEGMENT** – lets you declare the beginning of a memory segment. We'll talk in detail about memory segments tomorrow.
- **EQU** – lets you define pre-processor constants.

# Assembly primer - directives

```
; Let's declare some external functions  
EXTERN kbdine, dspout, dspmsg, dosxit
```

```
; Let's begin our code segment  
SEGMENT code
```

```
; I'd normally put my variables right here
```

```
..start ; This is a special label that  
; denotes the execution starting  
; point. The next instruction  
; after ..start will be the first  
; to run when you execute your program
```

# Assembly primer - instructions

```
; Here are some simple instructions  
mov      ax, [VAR1] ; notice the brackets  
mov      dx, STR1   ; notice the not brackets  
call     dspmsg  
jmp     done  
mov      bx, [VAR2] ; this will never happen  
  
done
```

# Example

```
SEGMENT code

var1 db 55h ; 55
var2 db 0AAh ; AA
var3 db 12h ; 12
str1 db "Hello", 0 ; 48 65 6C 6C 6F 00

..start
    mov al, [var1] ; A0 00 00
    mov bx, var3 ; BB 02 00
    inc bx ; 43
    mov al, [bx] ; 8A 07
```

# Assignments

- Keep working on MP0 and HW0