

Q: What are the three popular strategies for allocating free memory blocks to processes in dynamic memory partitioning? Explain briefly how each strategy works.

A: First-fit: chooses the first free block in the list that is large enough for the request. Best-fit: chooses the free block that is closest in size to the request. Next-fit: chooses the first free block that is large enough for the request and comes after the "Last Allocated Block" in the list.

Q: T/F: Buddy Strategy always allocates memory in chunks of size power of two and uses a data structure based on a binary tree.

A: True.

Q: Give an example drawing of a partially allocated memory and appropriate pointers in which each strategy ends up allocating a different free memory block to satisfy a request for 16 MB. Indicate clearly which strategy allocates which memory block.

A: Refer to the example in Chapter 7 slides.

Q: What interrupt is created when a desired page frame is not currently resident in RAM?

A: Page fault trap.

Q: How does the hardware 'know' that a desired page frame is not currently resident in RAM?

A: Valid bit.

Q: What precisely does it mean if the 'dirty bit' is set for a page frame?

A: The page frame has been modified.

Q: What is 'good' vs. 'bad' program locality?

A: 'Good' locality means that the process executes in clustered pages. 'Bad' locality means that the process executes in scattered pages.

Q: Explain when/how internal fragmentation may occur.

A: When fixed-sized pages are used, the last page of a program may be partially filled. This is called internal fragmentation.

Q: Explain when/how external fragmentation may occur.

A: Segmentation system breaks up the memory space into variable-sized pieces. After a sequence of allocation and deallocations, free memory may get fragmented into small pieces. Even if the total size of free memory is large enough to satisfy large memory requests, a large request may not be met due to the lack of continuity between small fragments. This is called external fragmentation. Compaction is needed to put free blocks into one large memory block.

Q: What is a global allocation scheme?

A: Global replacement allows a process to select a replacement frame from the set of all frames, even if that frame is currently allocated to some other process; one process can take a frame from another.

Q: What is a working set model?

A: The working set model assumes that processes execute in localities. The working set is the set of pages in the current locality. Accordingly, each process should be allocated enough frames for its current working set.

Q: Comparing global allocation vs. working set allocation, which would be more adversely affected by a program with 'bad' locality? and why would that be true?

A: Working set allocation would be more adversely affected by a program with 'bad' locality. This is because the program with 'bad' locality has poorly defined working sets and therefore, many page faults are likely to occur.

Q: Which one of the following is not among the set of events that may take place between the time a page fault occurs and the time the faulting process resumes execution?

(a) OS blocks the process and puts it into a wait queue.

(b) One of the processes in the ready queue is selected to run.

(c) A DMA is initiated to load the page from disk into main memory.

(d) A page replacement strategy is used to find a page frame to load the new page.

(e) Page table is updated to reflect the change.

(f) none of the above

A: f

Q: Which one of the following is not among the set of events that may take place between the time a page fault occurs and the time the faulting process resumes execution?

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(b) One of the processes in the ready queue is selected to run.

(c) A DMA is initiated to load the page from disk into main memory.

(d) The last page that the faulting process was executing is replaced with the newly loaded page.

(e) Page table is updated to reflect the change.

(f) none of the above

A: d

Q: T/F: While DMA (Direct Memory Access) is taking place, processor is free to do other things. The processor is only involved at the beginning and end of the DMA transfer.

A: True

Q: T/F: DMA uses "cycle stealing" to transfer data on the system bus. Each time cycle stealing is used, CPU is interrupted.

A: False. The only interrupts occur during DMA transfer is at the beginning and at the end of DMA (no interrupts in between).

Q: What is the "largest" program that could execute on a machine with a 24-bit virtual address?

A: 2^24 bytes

Q: What is the "largest" program that could execute on a machine with a 24-bit physical address?

A: Can't tell. Need to know the size of the virtual (logical) address.

Q: The address contained in a TLB entry is (physical / logical).

A: physical

Q: List at least 3 flags that are contained in a PTE.

A: Valid bit, Reference bit, Dirty bit.

Q: Define hit-ratio in a memory management context.

A: In a two-level memory (cache-RAM or RAM-Harddisk), the fraction of all memory accesses that are found in the master memory (i.e. the cache).

Q: T/F: If a virtual page number X generates a miss in the TLB (Translation Lookaside Buffer), then the corresponding physical page number for X is guaranteed to be found in the Page Table Entry.

A: False. Not necessarily. X may not be resident in RAM (still in secondary memory).

Q: True or False? It is possible that page tables are stored in virtual (secondary) memory?

A: True. when multi-level paging schemes are used

Q: T/F: In a virtual memory system with paging, page size must be large enough to offset the high cost of page faults.

A: True

Q: T/F: The Least Recently Used (LRU) page replacement strategy is based on the principle of spatial locality (locality in space) as opposed to temporal locality (locality in time).

A: False. LRU strategy is based on the temporal locality.

Q: T/F: Consider "clock policy" for page replacement; a newly arrived page (i.e., just swapped in) will not get replaced before the clock pointer makes two full rotations in the circular buffer of candidate page frames.

A: True. use=1 at arrival, use=0 after the first rotation of the clock pointer. If use=0 remains true after the second rotation, it may get replaced; otherwise, it stays on the clock.

Q: T/F: In a virtual memory system with paging, you can run a program whose size is larger than the size of Main Memory.

A: True.

Q: Draw a picture which shows the relationship between "Virtual Address (virtual page#, offset)", "Translation Lookaside Buffer", "Page Table", "Real Address (frame#, offset)", "Main Memory", and "Secondary Memory".

A: Slide #18 in CHAPS.ppt

Q: Draw the flowchart for paging using TLB, Page Table, and DMA.

A: Slide #19 in CHAPS.ppt

Q: How does the kernel 'know' where on disk the desired information is for a non-resident frame?

A: If valid bit=0, Page Table Entry should contain the Disk address.

Q: Describe what demand paging means.

A: The technique of only loading virtual pages into memory as they are accessed is known as demand paging. If the demand pages are not in memory, a page fault trap happens, and the operation system swaps them in.

Q: Describe what prepaging means.

A: Prepaging brings in more consecutive pages than needed. If a virtual page X causes a pagefault, then virtual page (X+1) is also brought in along with X. It is less overhead to bring in pages that reside contiguously on the disk

-> Must know slides 14 through 20 (in CHAPS.ppt) in detail

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Please fill in the blanks for the following questions:

-----Q: What is a working set model?

A: The working set model assumes that processes execute in localities. The working set is the set of pages in the current locality. Accordingly, each process should be allocated enough frames for its current working set.

A: A Page Fault

Q: Since paging system uses _____ sized pages, _____ fragmentation may occur.

A: fixed size; internal

Q: If a memory management system uses dynamic partitioning, _____ fragmentation may occur.

A: External

Q: _____ is a form of I/O in which a special module controls the exchange of data between main memory and an I/O device. During this I/O transfer, CPU is free to do other computation.

A: DMA (Direct Memory Access)

Q: The least recently used (LRU) page replacement strategy is based on the principle of _____ as opposed to _____.

A: temporal locality; spatial locality

Q: The top four levels in the memory hierarchy, starting with the fastest, are: _____, _____, _____, _____.

A: Registers; cache memory; RAM; Disk.

Q: Swapping out a piece of a process (i.e. process) to disk just before that piece is needed is called _____.

A: Thrashing

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Q: T/F: A unix socket is used for communication between processes running on the same machine. On the other hand, an Internet socket can not be used for communication between processes running on the same machine.

A: False. Internet Socket can be used for processes running on the same machine as well as on different machines

Q: T/F: If clients are connected to a server through "connect()" and "accept()" calls and the server calls "listen(soc,2)" before "accept()", then at most 2 clients can get connected to the server at any time.

A: False. listen() determines the size of the wait queue before the connections take place, not the max. number of clients that can be connected.

Q: T/F: In RAID (Redundant Array of Independent Disks) Level 1, every disk in the array has a mirror disk that contains the same data.

A: True.

Q: T/F: In Client/Server architectures, OS and the platforms in the client and server machines must be the same.

A: FALSE

Q: T/F: In Client/Server applications, there is heavy emphasis on providing a user-friendly Graphical User Interface (GUI) on the client side.

A: True.

Q: T/F: In client/server applications, fat client models can not take advantage of the desktop power and therefore can only serve a small number of clients.

A: True.

Q: T/F: First-Come-First-Served (FCFS) process scheduling favors I/O-bound processes.

A: FALSE

Q: True or False? Most antivirus software is based on program emulation and virus signature analysis.

A: True.

Q: T/F: RAID 2 (Redundant Array of Independent Disks with Level 2) is designed to provide error detection/correction.

A: True.

Q: T/F: User Datagram Protocol (UDP) provides unduplicated and reliable packet delivery.

A: FALSE

Q: T/F: Two periodic real-time processes A and B have periods T_a=0.2 ms and T_b=0.5 ms respectively. Furthermore, their execution times are a₁=10 micro sec. and C_b=40 micro sec. respectively. If Rate Monolithic

scheduling is used A has higher priority than B.

A: True. False. "short-term scheduler (dispatcher)" executes most frequently.

Q: T/F: The following I/O devices are sorted correctly in decreasing order according to the typical data rates that they can sustain: Gigabit Ethernet, firewall 800, laser printer, hard disk, keyboard, and modem.

A: FALSE

Q: T/F: DMA uses "cycle stealing" to transfer data on the system bus. Each time cycle stealing is used, CPU is interrupted.

A: FALSE

Q: Which of the following strategies is not used in a Disk Scheduling Algorithm?

(a) First in first out (FIFO)

(b) Last in first out (LIFO)

(c) Shortest service time first (SSTF)

(d) Longest service time first (LSTF)

(e) Back and forth over disk (SCAN)

A: d.

Q: Explain what the following C calls do both when the call is successful and when it is unsuccessful.

1. socket(AF_INET, SOCK_STREAM, 0)

2. bind(sd, (struct sockaddr*)&server_addr, sizeof(server_addr))

3. socket(AF_INET, SOCK_DGRAM, 0)

4. accept(sd, (struct sockaddr*)&client_addr, &client_len)

A: 1. creates an Internet stream (TCP) socket and returns the socket descriptor. If the call fails, it returns -1.

2. binds the definition of a socket (socket descriptor) to a port number. If the call fails, it returns -1.

3. creates an Internet datagram (UDP) socket and returns the socket descriptor. If the call fails, it returns -1.

4. blocks execution until a client connection is received. When that happens, it returns a descriptor for the connection. If the call fails, it returns -1.

5. binds the definition of a socket (socket descriptor) to a port number. If the call fails, it returns -1.

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short-term), long-term scheduler executes most frequently.

A: False. "short-term scheduler (dispatcher)" executes most frequently.

Q: T/F: Among the short-term scheduling policies, feedback policy penalizes jobs that have been running longer.

A: True

Q: Which decisions are made by Long-term, Medium-term, and Short-term scheduling? Be brief.

A: Long-term scheduling: Determines which programs are admitted to the system for processing; Controls the degree of multiprogramming

Medium-term scheduling: Determines which programs will be resident. Part of the swapping function; Swapping-in decision is based on the need to manage the degree of multiprogramming

Short-term scheduling: Determines which program will be executed on CPU next. Known as the dispatcher; Executes most frequently

Q: Name 3 things that are essential to launching a "botnet" attack:

A: (1) attack software

(2) a large number of vulnerable machines

(3) locating these machines (scanning or fingerprinting)

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Please fill in the blanks for the following questions:

-----Q: The two lowest layers in the 7-layer ISO Open Systems Interconnect (OSI) model are _____ and _____.

A: Physical; Data Link; signaling technology; frame management.

Q: Two transport protocols, _____ and _____, are defined and handled at the Transport Layer.

A: Transmission Control Protocol (TCP); User Datagram Protocol (UDP)

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Miscellaneous questions related to the VIDEO: TRIUMPH OF THE MERDS

-----Q: Fill in the blanks.

1. _____ and _____ are generally credited with the invention of the microcomputer.

A: Dennis Ritchie, Ken Thompson.

2. _____ and _____ started Microsoft in 19 _____.

A: Bill Gates, Paul Allen, 1975.

3. What corp./laboratory may