ZPU Reference Handbook



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Introduction

The Zylin ZPU is the worlds smallest 32 bit CPU with GCC tool chain. The ZPU is a small CPU in two ways: it takes up very little resources and the architecture itself is small. The latter can be important when learning about CPU architectures and implementing variations of the ZPU where aspects of CPU design is examined. In academia students can learn VHDL, CPU architecture in general and complete exercises in the course of a year. The current ZPU instruction set and architecture has not changed for the last couple of years and can be considered quite stable. This shall be presented in detail the following chapters.

Part of this work is based on previous work done by Álvaro Lopes - <u>alvieboy@alvie.com</u> (see legal notice) on the ZPUino – a derivative work of the original ZPU core by Øyvind Harboe - <u>oyvind.harboe@zylin.com</u>. The original ZPUino can be found on the internet on the following website: http://www.alvie.com/zpuino. Furthermore, the original ZPU and the "ZPU Project" can also be found on the internet on the following website: https://github.com/zylin/zpu.

Instruction Set Summary

Stack Operation Definitions

```
TOS = Top Of Stack = SP
mem[SP] = valid data = stackA

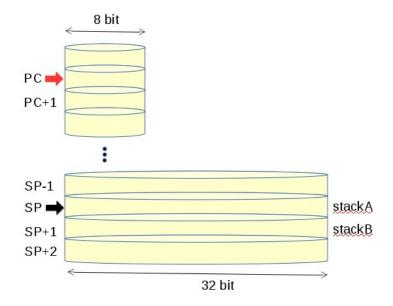
PUSH:
    SP = SP - 1;
    mem[SP] = data;

POP:
    data = mem[SP];
    SP = SP + 1;
```

Memory Operations

```
PC: Accesses memory in 8 bit cells SP: Accesses memory in 32 bit cells
```

NOTE: although PC points to 8 bit cells, the ZPU state machine always fetches 32 bit words and internally breaks down the words into bytes. The Stack Pointer, however, is a pointer to 32 bit cells which are aligned on 4-byte boundary, i.e. SP results in a memory fetch to address 4*SP and.



NOTE: the CPU implementation in VHDL is such that the TOS

(stackA) and mem[SP+1] (stackB), i.e. both instruction operands, are normally not immediately written back to memory in order to save CPU cycles. Care should be taken while reading the code to exactly understand when (due to state transitions) the stackA and/or stackB need to be written back to memory. See also instruction description below for a better understanding.

This means that stackA and stackB internal variables are actually cached versions of the corresponding memory positions. When SP changes, the stackA and stackB have to be updated accordingly and so does the memory positions corresponding to stackA (SP) and/or stackB (SP+1) <u>before</u> SP is updated.



Core instructions summary

Mnemonic	Opcode	Description
BREAKPOINT	0000 0000	Sets 'break' line to logic '1'
IM x	1xxx xxxx	
STORESP x	010x xxxx	
POP	0101 0000	Implemented using STORESP 0
POPDOWN	0101 0001	Implemented using STORESP 1
LOADSP x	011x xxxx	
DUP	0111 0000	Implemented using LOADSP 0
DUPSTACKB	0111 0001	Implemented using LOADSP 1
ADDSP x	0001 xxxx	
SHIFT	0001 0000	Implemented using ADDSP 0
ADDTOP	0001 0001	Implemented using ADDSP 1
EMULATE x	001x xxxx	
POPPC	0000 0100	
LOAD	0000 1000	
STORE	0000 1100	
PUSHSP	0000 0010	
POPSP	0000 1101	
ADD	0000 0101	
AND	0000 0110	
OR	0000 0111	
NOT	0000 1001	
FLIP	0000 1010	
NOP	0000 1011	

Optional instructions (emulated)

Mnemonic	Opcode	Decimal	Description
?	0010 0000	32	
N/A	0010 0001	33	
LOADH	0010 0010	34	
STOREH	0010 0011	35	
LESSTHAN	0010 0100	36	
LESSTHANOREQUAL	0010 0101	37	
ULESSTHAN	0010 0110	38	
ULESSTHANOREQUAL	0010 0111	39	
SWAP	0010 1000	40	
MULT	0010 1001	41	
LSHIFTRIGHT	0010 1010	42	
ASHIFTLEFT	0010 1011	43	
ASHIFTRIGHT	0010 1100	44	
CALL	0010 1101	45	
EQ	0010 1110	46	
NEQ	0010 1111	47	
NEG	0011 0000	48	
SUB	0011 0001	49	
XOR	0011 0010	50	
LOADB	0011 0011	51	
STOREB	0011 0100	52	
DIV	0011 0101	53	
MOD	0011 0110	54	
EQBRANCH	0011 0111	55	
NEQBRANCH	0011 1000	56	
POPPCREL	0011 1001	57	
CONFIG	0011 1010	58	
PUSHPC (a)	0011 1011	59	
SYSCALL (a)	0011 1100	60	
PUSHSPADD	0011 1101	61	
HALFMULT	0011 1110	62	
CALLPCREL	0011 1111	63	

Instruction Mapping

71	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001	1010	1011	1100	1101	1110	1111
0000	BRK	NA4	PUSHSP	NA3	POPPC	ADD	AND	OR	LOAD	NOT	FLIP	NOP	STORE	POPSP	NA2	NA
0001	ADDTOP	SHIFT	ADDSP													
0010	?	N/A	EMU													
0011	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU	EMU
0100	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP	STORESP
0101	POP	POPDOWN	STORESP													
0110	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP	LOADSP
0111	DUP	DUPSTACKB	LOADSP													
1000	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1001	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1010	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1011	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1100	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1101	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1110	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM
1111	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM	IM

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Implemented Instructions

7	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001	1010	1011	1100	1101	1110	1111
0000	BRK	NA4	PUSHSP	NA3	POPPC	ADD	AND	OR	LOAD	NOT	FLIP	NOP	STORE	POPSP	NA2	NA
0001	ADDSP															
0010	?	N/A	EMU													
0011	EMU															
0100	STORESP															
0101	STORESP															
0110	LOADSP															
0111	LOADSP															
1000	IM															
1001	IM															
1010	IM															
1011	IM															
1100	IM															
1101	IM															
1110	IM															
1111	IM															

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Instructions Specification

IM

OPCODE	IM x					
MACHINE CODE	1xxxxxxx					
IMPLEMENTED	YES					
EMULATED	NO					
SP ACTION	single PUSH					
DESCRIPTION	Pushes immediate value into TOS					
PSEUDOCODE	<pre>if (idim='0') { // no previous IM idim = 1; mem[sp+1] = stackB; stackB = stackA; stackA = {mem[sp][24:0], x[6:0]}; // sign extend sp = sp - 1; } else {</pre>					
EQUIVALENT CODE	<pre>if (!idim) { push(x); } else { idim = 1; a = pop(); push(a << 7 + x); }</pre>					
INTERNAL	PREVIOUS "IM" (idim='1'): Before					
	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$					

EMULATE

OPCODE	EMULATE x					
MACHINE CODE	001xxxxx					
IMPLEMENTED	YES					
EMULATED	NO					
SP ACTION	PUSH					
DESCRIPTION	If the instruction is not implemented in hardware, this instruction will fired-up the microcode implementation of the function. $0 <= x <= 31$					
PSEUDOCODE	<pre>mem[sp+1] = stackB; // sp = sp - 1; // stackB = stackA; stackA = pc + 1; // pc = 32*x; fetch();</pre>	/ make room for push				
EQUIVALENT CODE	call(32*x);					
INTERNAL	Before	After				
LAYOUT	SP → [] ← A	SP → [] ← A = return_address SP+1 → [] ← B SP+2 → [B] ← B				

STORESP

OPCODE	STORESP x					
MACHINE CODE	010xxxx					
IMPLEMENTED	YES					
EMULATED	NO					
SP ACTION	POP					
DESCRIPTION	Pop TOS and store it at mem[SP+x]					
PSEUDOCODE	<pre>(storeSP) mem[sp+x] = stackA; // NOTE: x is always unsigned stackA = stackB; sp = sp + 1; (storeSP2) stackB = mem[sp+1];</pre>					
EQUIVALENT	mem[SP+x] = TOS;					
CODE	pop();					
INTERNAL	Before After					
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$					

POP

OPCODE	POP / STORESP 0
MACHINE CODE	01010000
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP
DESCRIPTION	Pops value from stack
PSEUDOCODE	<pre>stackA = stackB; sp = sp + 1; stackB = mem[sp+1];</pre>
EQUIVALENT CODE	pop();
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

POPDOWN

OPCODE	POPDOWN / STORESP 1
MACHINE CODE	01010000
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP
DESCRIPTION	Pops two values from stack and pushes first value back to stack.
PSEUDOCODE	<pre>sp = sp + 1; stackB = mem[sp+1];</pre>
EQUIVALENT CODE	<pre>a = pop(); pop(); push(a);</pre>
INTERNAL LAYOUT	Before

LOADSP

OPCODE	LOADSP x				
MACHINE CODE	011xxxxx				
IMPLEMENTED	YES				
EMULATED	NO				
SP ACTION	PUSH				
DESCRIPTION	Push value at mem[SP+x] into stack				
PSEUDOCODE	<pre>(LoadSP) mem[sp+1] = stackB; // writeback cached stackB sp = sp - 1; // required for push (LoadSP2) read = mem[SP+x+1]; // fetch mem[SP+x] (LoadSP3) stackB = stackA; // stackB is now = old stackA stackA = read; // stackA = mem[SP+x]</pre>				
EQUIVALENT CODE	<pre>a = mem[sp+x]; push(a);</pre>				
INTERNAL LAYOUT	Before After SP-1 → [] SP → [] → A (= v) SP → [] ← A SP+1 → [] ← B (= prev A) SP+1 → [] ← B SP+2 → [] SP+x → [v] SP+x+1 → [v]				



DUP

OPCODE	DUP / LOADSP 0
MACHINE CODE	01110000
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	PUSH
DESCRIPTION	Push TOS again into stack
PSEUDOCODE	<pre>mem[sp+1] = stackB;</pre>
EQUIVALENT CODE	<pre>a = pop(); push(a); push(a);</pre>
INTERNAL LAYOUT	Before After

DUPSTACKB

OPCODE	DUPSTACKB / LOADSP 1
MACHINE CODE	01110001
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	PUSH
DESCRIPTION	Push stackB again into stack
PSEUDOCODE	A = stackA; // save old stackA B = stackB; // save old stackB stackA = stackB; // new stackA = old stackB stackB = A; // new stackB = old stackA mem[sp+1] = B; // writeback cached (old) stackB sp = sp - 1; // required for push
EQUIVALENT CODE	<pre>a = pop(); // get stackA b = pop(); // get stackB push(b); // save back stackB push(a); // save back stackA push(b); // duplicate stackB</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

ADDSP

OPCODE	ADDSP x	
MACHINE CODE	0001xxxx	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + PUSH	
DESCRIPTION	TOS = TOS + mem[SP+x]	
PSEUDO CODE	<pre>a = mem[SP+x]; stackA = stackA + a;</pre>	
EQUIVALENT CODE	<pre>a = mem[SP+x]; b = pop(); push (a+b);</pre>	
INTERNAL LAYOUT		→ [] → [] ← A (+= mem[SP+x]) → [] ← B

ADDTOP

OPCODE	ADDTOP / ADDSP 1
MACHINE CODE	00010001
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	TOS = TOS + mem[SP+1]
PSEUDO CODE	stackA = stackA + stackB;
EQUIVALENT CODE	<pre>A = pop(); // get stackA B = pop(); // get stackB push(B); // push back stackB push(A+B); // push A+B</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

SHIFT

OPCODE	SHIFT / ADDSP 0
MACHINE CODE	00010000
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	TOS = 2*TOS
PSEUDO CODE	stackA = stackA << 1;
EQUIVALENT CODE	<pre>A = pop(); // get stackA push(A<<1); // push 2*A</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

BREAKPOINT

OPCODE	BREAK
MACHINE CODE	0000000
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	None
DESCRIPTION	Sets the break output line to '1' for one clock cycle.
PSEUDO CODE	
EQUIVALENT CODE	
INTERNAL	Before After

SHIFTLEFT

OPCODE	SHIFTLEFT
MACHINE CODE	0000001
DESCRIPTION	

PUSHSP

OPCODE	PUSHSP
MACHINE CODE	0000010
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	PUSH
DESCRIPTION	This instruction pushes the SP value into the stack
PSEUDO CODE	<pre>push(SP);</pre>
EQUIVALENT CODE	<pre>mem[SP+1] = stackB; stackB = stackA; stackA = SP; SP = SP - 1;</pre>
INTERNAL LAYOUT	Before

POPINT

OPCODE	POPINT
MACHINE CODE	0000011
DESCRIPTION	<pre>pc = mem[sp]; sp = sp + 1; fetch(); decode(); clear_interrupt_flag();</pre>

POPPC

OPCODE	POPPC
MACHINE CODE	00000100
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP
DESCRIPTION	Sets PC to value popped from stack.
PSEUDO CODE	<pre>pc = mem[sp]; sp = sp + 1; (resynch) stackA = mem[SP]; stackB = mem[SP+1];</pre>
EQUIVALENT CODE	pc = pop();
INTERNAL LAYOUT	Before After

ADD

OPCODE	ADD	
MACHINE CODE	00000101	
IMPLEMENTED	YES	
EMULATED	NO	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: X and Y. tIt then pushes back into the stack the value given by A+B.	
PSEUDO CODE	<pre>a = pop(); b = pop(); b = b + a; push(b);</pre>	
EQUIVALENT CODE	<pre>mem[sp+1] = mem[sp+1] + mem[sp]; sp = sp + 1;</pre>	
INTERNAL LAYOUT	Before After	

AND

OPCODE	AND
MACHINE CODE	00000110
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. tIt then pushes back into the stack the value given by A AND B.
PSEUDO CODE	<pre>a = pop(); b = pop(); b = b AND a; push(b);</pre>
EQUIVALENT CODE	mem[sp+1] = mem[sp+1] AND mem[sp]; sp = sp + 1;
INTERNAL LAYOUT	Before After

OR

OPCODE	OR
MACHINE CODE	00000111
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. tIt then pushes back into the stack the value given by A OR B.
PSEUDO CODE	<pre>a = pop(); b = pop(); b = b OR a; push(b);</pre>
EQUIVALENT CODE	mem[sp+1] = mem[sp+1] OR mem[sp]; sp = sp + 1;
INTERNAL LAYOUT	Before After

LOAD

OPCODE	LOAD
MACHINE CODE	00001000
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	PUSH
DESCRIPTION	Push value at mem[stackA] into stack
PSEUDOCODE	<pre>mem[SP+1] = stackB; SP = SP-1; stackB = stackA; stackA = mem[stackA];</pre>
EQUIVALENT CODE	<pre>a = pop(); b = mem[a]; push(b);</pre>
INTERNAL	Before After

NOT

OPCODE	NOT
MACHINE CODE	00001001
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	This instruction pops one value from stack (X). It then pushes back into the stack the value NOT X;
PSEUDO CODE	<pre>a = pop(); push(~a);</pre>
EQUIVALENT CODE	stackA = NOT stackA;
INTERNAL	Before After

FLIP

OPCODE	FLIP
MACHINE CODE	00001010
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	This instruction pops one value from stack (X). It then pushes back into the stack the value flip(X). The flip function rearranges the bits of X such that bit $\{n\} = bit\{L-1-n\}$, where L is the word length in bits. $\begin{array}{c ccccccccccccccccccccccccccccccccccc$
PSEUDO CODE	<pre>a = pop(); push(flip(a));</pre>
EQUIVALENT CODE	<pre>mem[sp] = flip(mem[sp]);</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

NOP

OPCODE	NOP
MACHINE CODE	00001011
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	none
DESCRIPTION	No operation
PSEUDOCODE	
EQUIVALENT CODE	
INTERNAL LAYOUT	Before After SP-1 \rightarrow [] SP-1 \rightarrow [] SP \rightarrow [] \leftarrow A SP \rightarrow [] \leftarrow A SP+1 \rightarrow [] \leftarrow B SP+1 \rightarrow [] \leftarrow B

STORE

OPCODE	STORE
MACHINE CODE	00001100
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP + POP
DESCRIPTION	Pop memory address A and value B from stack. Write to memory at address A the value B.
PSEUDOCODE	<pre>mem[stackA] = stackB; SP = SP+2; resynch(); // reloads stackA and stackB</pre>
EQUIVALENT CODE	<pre>A = pop(); // memory address B = pop(); // value to write mem[A] = B; // write value to memory</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

POPSP

OPCODE	POPSP
MACHINE CODE	00001101
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	new stack frame
DESCRIPTION	Load SP with TOS
PSEUDOCODE	<pre>sp = stackA; resynch();</pre>
EQUIVALENT CODE	SP = TOS;
INTERNAL LAYOUT	Before After

COMPARE

OPCODE	COMPARE / IPSUM
MACHINE CODE	00001110
DESCRIPTION	<pre>c = mem[sp]; s = mem[sp+1]; sum = 0; while (c>0){ sum += halfword(mem[s],s); s += 2; }; sp = sp+1; mem[sp] = sum; (overwrites mem[0] & mem[4] words)</pre>

SNCPY

```
OPCODE

MACHINE CODE

00001111

DESCRIPTION

c = mem[sp];
d = mem[sp+1];
s = mem[sp+2];
while ( *(char*)s != 0 && c>0){
    *((char*)d++) =* ((char*)s++));
    c--
};
sp = sp+3;
(overwrites mem[0] & mem[4] words)
```

SNCPY2

WCPY

```
OPCODE

MCPY

MACHINE CODE

00100001

v = mem[sp];
c = mem[sp+1];
d = mem[sp+2];
while (c-->0) {
    mem[d++] = v;
}
sp = sp+3;
(overwrites mem[0] & mem[4] words)
```



LOADH

OPCODE	LOADH
MACHINE CODE	00100010
DESCRIPTION	<pre>mem[sp] = halfword[mem[sp]];</pre>

STOREH

OPCODE	STOREH
MACHINE CODE	00100011
DESCRIPTION	<pre>halfword[mem[sp]] = (mem[sp+1] & 0xFFFF); sp = sp + 2;</pre>

LESSTHAN

OPCODE	LESSTHAN
MACHINE CODE	00100100
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value 1 if A <b 0;<="" back="" comparison;="" it="" otherwise="" pushes="" signed="" td="" the="" using="" value="">
PSEUDO CODE	<pre>a = pop(); b = pop(); if ((signed)a < (signed)b) { push(1); } else { push(0); }</pre>
EQUIVALENT CODE	<pre>if (mem[sp] < mem[sp+1]) { // signed comparison a = 1; } else { a = 0; } mem[sp+1] = a; sp = sp + 1</pre>
INTERNAL LAYOUT	Before After

LESSTHANOREQUAL

OPCODE	LESSTHANOREQUAL
MACHINE CODE	00100101
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value 1 if A<=B using signed comparison; otherwise it pushes back the value 0;
PSEUDO CODE	<pre>a = pop(); b = pop(); if ((signed)a <= (signed)b) { push(1); } else { push(0); }</pre>
EQUIVALENT CODE	<pre>if (mem[sp] <= mem[sp+1]) { // signed comparison a = 1; } else { a = 0; } mem[sp+1] = a; sp = sp + 1</pre>
INTERNAL LAYOUT	Before After

ULESSTHAN

OPCODE	ULESSTHAN
MACHINE CODE	00100101
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value 1 if A <b 0;<="" back="" comparison;="" it="" otherwise="" pushes="" td="" the="" unsigned="" using="" value="">
PSEUDO CODE	<pre>a = pop(); b = pop(); if ((unsigned)a < (unsigned)b) { push(1); } else { push(0); }</pre>
EQUIVALENT CODE	<pre>if (mem[sp] < mem[sp+1]) { // unsigned comparison a = 1; } else { a = 0; } mem[sp+1] = a; sp</pre>
INTERNAL LAYOUT	Before After

ULESSTHANOREQUAL

OPCODE	ULESSTHANOREQUAL
MACHINE CODE	00100110
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value 1 if A<=B using unsigned comparison; otherwise it pushes back the value 0;
PSEUDO CODE	<pre>a = pop(); b = pop(); if ((unsigned)a <= (unsigned)b) { push(1); } else { push(0); }</pre>
EQUIVALENT CODE	<pre>if (mem[sp] <= mem[sp+1]) { // unsigned comparison a = 1; } else { a = 0; } mem[sp+1] = a; sp</pre>
INTERNAL LAYOUT	Before After

SWAP

OPCODE	SWAP
MACHINE CODE	00101000
DESCRIPTION	•

MULT

OPCODE	MULT
MACHINE CODE	00101001
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value X*Y. Note, only the lower order bits are pushed into the stack.
PSEUDO CODE	<pre>a = pop(); b = pop(); push(a*b);</pre>
EQUIVALENT CODE	<pre>stackA = stackA * stackB; mem[SP+1] = stackB; sp = sp + 1;</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$

LSHIFTRIGHT

OPCODE	LSHIFTRIGHT
MACHINE CODE	00101010
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP + PUSH
DESCRIPTION	This instruction pops two values from stack: the shift parameter N and the shift variable V. It then pushes back into the stack the right logic shift value given by V>>N. The figure below shows an example for only 8 bit (this CPU works on 32bit).
PSEUDO CODE	<pre>a = pop();</pre>
EQUIVALENT CODE	<pre>if (TOS!=0) { while (stackA) { stackA = stackA - 1; stackB = ('0', stackB[31:1]); } } stackA = stackB; stackA = stackB; stackB = mem[SP+2]; SP = SP + 1;</pre>
INTERNAL	Before After

ASHIFTLEFT

OPCODE	ASHIFTLEFT	
MACHINE CODE	00101011	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: the shift parameter N and the shift variable V. It then pushes back into the stack the left arithmetic shift value given by V< <n. (this="" 0="" 1="" 1<="" 2="" 3="" 32bit).="" 4="" 5="" 6="" 7="" 8="" an="" below="" bit="" cpu="" example="" figure="" for="" on="" only="" shows="" td="" the="" works=""></n.>	
PSEUDO CODE	<pre>a = pop(); // nr of shifts b = pop(); // value to shift b = b << a; push(b);</pre>	
EQUIVALENT CODE	<pre>if (TOS!=0) { while (stackA) { stackA = stackA - 1; stackB <<= 1; } } stackA = stackB; stackA = stackB; stackB = mem[SP+2]; SP = SP + 1;</pre>	
INTERNAL LAYOUT	Before After SP → [] ← A SP-1 → [] SP+1 → [] ← B SP → [] ← A SP+2 → [] SP+1 → [] → B	

ASHIFTRIGHT

OPCODE	ASHIFTRIGHT	
MACHINE CODE	00101100	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: the shift parameter N and the shift variable V. It then pushes back into the stack the left arithmetic shift value given by V>>N. The figure below shows an example for only 8 bit (this CPU works on 32bit). $ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	
PSEUDO CODE	<pre>a = pop(); // nr of shifts b = pop(); // value to shift b = b >> a; push(b);</pre>	
EQUIVALENT CODE	<pre>if (TOS!=0) { while (stackA) { stackA = stackA - 1; stackB >>= 1; } } stackA = stackB; stackA = stackB; stackB = mem[SP+2]; SP = SP + 1;</pre>	
INTERNAL	Before After	

CALL

OPCODE	CALL
MACHINE CODE	00101101
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	jumps to memory address TOP, while adjusting the stack with the return address pc+1
PSEUDO CODE	<pre>x = stackA[L-1:0]; // (POP) L: mem address lines stackA = pc + 1; // (PUSH) return address pc = x; // jump to x</pre>
EQUIVALENT CODE	<pre>x = pop(); call(x); // pushes PC+1 into stack</pre>
INTERNAL LAYOUT	Before After

EQ

OPCODE	EQ	
MACHINE CODE	00101110	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value 1 if A==B; otherwise it pushes back the value 0;	
PSEUDO CODE	<pre>a = pop(); b = pop(); if (a==b) { push(1); } else { push(0); }</pre>	
EQUIVALENT CODE	<pre>if (mem[sp] == mem[sp+1]) { a = 1; } else { a = 0; } mem[sp+1] = a; sp = sp + 1</pre>	
INTERNAL	Before After	

NEQ

OPCODE	NEQ	
MACHINE CODE	00101111	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value 0 if A==B; otherwise it pushes back the value 1;	
PSEUDO CODE	<pre>a = pop(); b = pop(); if (a!=b) { push(1); } else { push(0); }</pre>	
EQUIVALENT CODE	<pre>if (mem[sp] != mem[sp+1]) { a = 1; } else { a = 0; } mem[sp+1] = a; sp = sp + 1</pre>	
INTERNAL	Before After	
LAYOUT	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	



NEG

OPCODE	NEG	
MACHINE CODE	00110000	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + PUSH	
DESCRIPTION	_	ops one value from stack s back into the stack the
PSEUDO CODE	<pre>a = pop(); push(-a);</pre>	
EQUIVALENT CODE	stackA = 1 + NOT stack	ςA;
INTERNAL	Before	After
LAYOUT	SP → [] ← A SP+1 → [] ← B SP+2 → []	SP → [] ← A SP+1 → [] ← B SP+2 → []

SUB

OPCODE	SUB	
MACHINE CODE	00110001	
IMPLEMENTED	YES	
EMULATED	NO	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: X and Y. It then pushes back into the stack the value given B-A.	
PSEUDO CODE	<pre>a = pop(); b = pop(); b = b - a; push(b);</pre>	
EQUIVALENT CODE	mem[sp+1] = mem[sp+1] - mem[sp]; sp = sp + 1;	
INTERNAL	Before After	
LAYOUT	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	

XOR

OPCODE	XOR	
MACHINE CODE	00110010	
IMPLEMENTED	YES	
EMULATED	YES	
SP ACTION	POP + POP + PUSH	
DESCRIPTION	This instruction pops two values from stack: X and Y. tIt then pushes back into the stack the value given by A XOR B.	
PSEUDO CODE	<pre>a = pop(); b = pop(); b = b XOR a; push(b);</pre>	
EQUIVALENT CODE	mem[sp+1] = mem[sp+1] XOR mem[sp]; sp = sp + 1;	
INTERNAL LAYOUT	Before After	



LOADB

OPCODE	LOADB
MACHINE CODE	00110011
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP + PUSH
DESCRIPTION	Pops the memory address A and pushes back the byte at mem[stackA] as a 32bit uint into stack
PSEUDOCODE	<pre>mem[sp] = byte[stackA];</pre>
EQUIVALENT CODE	<pre>a = pop(); b = BYTE{ mem[a] }; push((uint32)b);</pre>
INTERNAL LAYOUT	Before After

STOREB

OPCODE	STOREB
MACHINE CODE	00110100
IMPLEMENTED	YES
EMULATED	NO
SP ACTION	POP + POP
DESCRIPTION	Pops the memory address A and byte B (as 32bit) and writes back a byte with value B to memory at address A
PSEUDOCODE	<pre>mem[stackA] = byte[stackB];</pre>
EQUIVALENT CODE	<pre>a = pop(); b = pop(); BYTE{ mem[A] } = BYTE{ b & 0x000000FF };</pre>
INTERNAL LAYOUT	Before

DIV

OPCODE	DIV
MACHINE CODE	00110101
DESCRIPTION	•

MOD

OPCODE	мор
MACHINE CODE	00110110
DESCRIPTION	•

EQBRANCH

OPCODE	EQBRANCH
MACHINE CODE	00110111
DESCRIPTION	<pre>if (mem[sp+1] == 0) { pc = pc + mem[sp]; sp = sp + 2; }</pre>

NEQBRANCH

OPCODE	NEQBRANCH
MACHINE CODE	00111000
DESCRIPTION	<pre>if (mem[sp+1] != 0) { pc = pc + mem[sp]; sp = sp + 2; }</pre>

OPCODE	NEQBRANCH
MACHINE CODE	00111000
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + POP
DESCRIPTION	This instruction pops two values from stack: X and Y; and adjusts the stack pointer accordingly. If Y!=0, then load PC with PC+X
PSEUDO CODE	<pre>A = pop(); B = pop(); SP = SP + 2; if (B!=0) { PC = PC + A; }</pre>
EQUIVALENT CODE	<pre>if (stackB != 0) { pc = pc + stackA; } sp = sp + 2; resynch();</pre>
INTERNAL LAYOUT	Before After

POPPCREL

OPCODE	POPPCREL
MACHINE CODE	00111001
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP
DESCRIPTION	Adds to PC the value popped from stack.
PSEUDO CODE	<pre>pc = pc + mem[sp]; sp = sp + 1; (resynch) stackA = mem[SP]; stackB = mem[SP+1];</pre>
EQUIVALENT CODE	<pre>pc = pop();</pre>
INTERNAL LAYOUT	Before After

CONFIG

OPCODE	CONFIG
MACHINE CODE	00111010
DESCRIPTION	

PUSHPC

OPCODE	PUSHPC
MACHINE CODE	00111011
DESCRIPTION	<pre>sp = sp - 1; mem[sp] = pc;</pre>

SYSCALL

OPCODE	SYSCALL
MACHINE CODE	00111100
DESCRIPTION	•

PUSHSPADD

OPCODE	PUSHSPADD
MACHINE CODE	00111101
DESCRIPTION	mem[sp] = sp + (stackA << 2)

OPCODE	PUSHSPADD
MACHINE CODE	00111101
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	This instruction replaces TOS with SP+4*A
PSEUDO CODE	A = pop(); push(SP+4*A);
EQUIVALENT CODE	mem[sp] = sp + (stackA << 2)
INTERNAL	Before After
LAYOUT	$\begin{array}{cccccccccccccccccccccccccccccccccccc$

OPCODE	PUSHSPADD
MACHINE CODE	00111101
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	TOS = 2*TOS
PSEUDO CODE	stackA = stackA << 1;
EQUIVALENT CODE	<pre>A = pop(); // get stackA push(A<<1); // push 2*A</pre>
INTERNAL	Before After
LAYOUT	$\begin{array}{cccccccccccccccccccccccccccccccccccc$

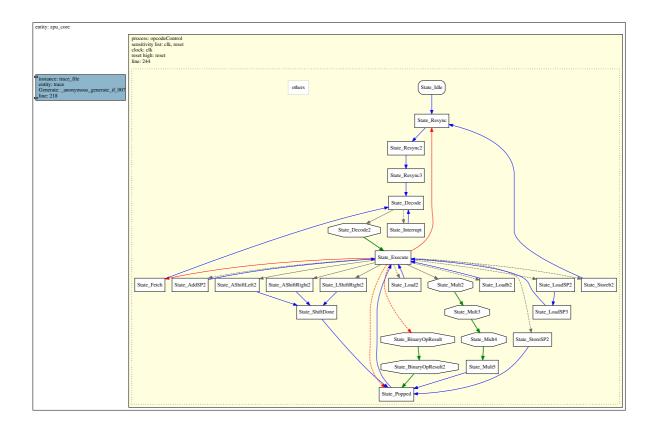
HALFMULT

OPCODE	HALFMULT
MACHINE CODE	00111110
DESCRIPTION	<pre>mem[sp+1] = 16bits(mem[sp]) * 16bits(mem[sp+1]); sp = sp + 1;</pre>

CALLPCREL

OPCODE	CALLPCREL
MACHINE CODE	00111110
IMPLEMENTED	YES
EMULATED	YES
SP ACTION	POP + PUSH
DESCRIPTION	jumps to memory address PC + TOP, while adjusting the stack with the return address pc+1
PSEUDOCODE	<pre>x = stackA[L-1:0]; // (POP) L: mem address lines stackA = pc + 1; // (PUSH) return address pc = pc + x; // jump to pc+x</pre>
EQUIVALENT CODE	<pre>x = pop(); call(pc+x); // pushes PC+1 into stack</pre>
INTERNAL LAYOUT	Before After

State Machine Transition Diagram



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