CSCE 625 Homework 1

UIN: 322007460

Name: Nagaraj Thenkarai Janakiraman

1. a) Following are the literals that we use to solve the tennis ball problem.

• Observations by drawing tennis ball:

Box-1	Box-2	Box-3
O1W	O2W	O3W
O1Y	O2Y	O3Y

• Correct labels for the boxes:

Box-1 Box-2	Box-3
C1W $C2W$ $C1Y$ $C2Y$ $C2P$	C3W $C3Y$
C1B $C2B$	C3B

• Wrong labels of each box

Box-1	Box-2	Box-3
L1W	L2W	L3W
L1Y	L2Y	L3Y
L1B	L2B	L3B

• Conditions for observations:

Box-1	Box-2	Box-3
$\mathbf{c1}: O1Y \implies C1Y \vee C1B$	$\mathbf{c3}: O2Y \implies C2Y \vee C2B$	$\mathbf{c5}: O3Y \implies C3Y \vee C3B$
$c2: O1W \implies C1W \lor C1B$	$\mathbf{c4}: O2W \implies C2W \vee C2B$	$\mathbf{c6}: O3W \implies C3W \lor C3B$

• Rules that imply wrong labels:

Box-1	Box-2	Box-3
$\mathbf{c7}: L1Y \implies \neg C1Y$	$\mathbf{c10}: L2Y \implies \neg C2Y$	$\mathbf{c13}: L3Y \implies \neg C3Y$
$c8: L1W \implies \neg C1W$	$\mathbf{c11}: L2W \implies \neg C2W$	$\mathbf{c14}: L3W \implies \neg C3W$
$\mathbf{c9}: L1B \implies \neg C1B$	$\mathbf{c12}: L2B \implies \neg C2B$	$\mathbf{c15}: L3B \implies \neg C3B$

• Constraint that there is at least one box of each kind (W, Y, B)

 $\begin{array}{l} \mathbf{c16} : C1Y \lor C1W \lor C1B \\ \mathbf{c17} : C2Y \lor C2W \lor C2B \\ \mathbf{c18} : C3Y \lor C3W \lor C3B \end{array}$

• Conditions to capture No repetitions in the box type (exactly one of each kind)

Box-1	Box-2	Box-3
, , , , , , , , , , , , , , , , , , , ,	$\mathbf{c22}: C2Y \implies \neg C1Y \land \neg C3Y$ $\mathbf{c23}: C2W \implies \neg C1W \land \neg C3W$, , , , , , , , , , , , , , , , , , , ,
$\mathbf{c21}: C1B \implies \neg C2B \land \neg C3B$	$\mathbf{c24}: C2B \implies \neg C1B \land \neg C3B$	$\mathbf{c27}: C3B \implies \neg C2B \land \neg C1B$

1. b)

Aim: To prove by using the Natural deduction $KB \models C2W$

Consider the initial KB observed from the given scenario:

* Initial observations:

R1 : O1Y

 $\mathbf{R2}:O2W$

R3 : O3Y

 $\mathbf{R4}:L1W$

 $\mathbf{R5}: L2Y$

R6 : L3B

- * $\mathbf{R7}: C3Y \vee C3B$ (Applying Modus Ponens to R3 and c5)
- * $\mathbf{R8} : \neg C3B$ (Applying Modus Ponens to R6 and c)
- * $\mathbf{R9}:C3Y$ (Resolving R7 and R8)
- * $\mathbf{R}\mathbf{10} : \neg C1Y \wedge C2Y$ (Applying Modus Ponens to R9 and c25)
- * **R11** : $C1Y \lor C1B$ (Applying Modus Ponens to R1 and c1)
- * $\mathbf{R}\mathbf{12} : \neg C1Y$ (Applying And-Elimination on R10)
- * **R13** :*C1B* (Resolving R11 and R12)
- * $\mathbf{R14}: C2W \wedge C2B$ (Applying Modus Ponens to R2 and c4)
- * $\mathbf{R15} : \neg C2B \land \neg C3B$ (From R13 and c21)
- * **R16** : $\neg C2B$ (From R15)
- * **R17** : C2W (Resolving R14 and R16)

Hence, we have $KB \models C2W$.

1. c)

• Knowledge Base for the given scenario:

R1 : O1Y

 $\mathbf{R2}:O2W$

R3 : O3Y

 $\mathbf{R4}:L1W$

R5:L2Y

R6 : L3B

• Converting all constraints to CNF (Conjunctive Normal Form):

- Solving for the query q as C2W, Getting the negation of the query R13:¬C2W
- R14: $C3Y \lor C3B$ (Resolution on R3 and R7)
- $\mathbf{R15} : \neg C3B$ (Resolution on R6 and R10)
- **R16** :*C3Y* (Resolution on R14 and R15)
- $\mathbf{R17} : \neg C1Y$ (Resolution on R16 and R12b)
- $\mathbf{R18}:C1Y \wedge C1B$ (Resolution on R1 and R8)
- $\mathbf{R19}:C1B$ (Resolution on R17 and R18)
- $\mathbf{R20} : \neg C2B$ (Resolution on R19 and R11a)
- $\mathbf{R21}: C2W \wedge C2B$ (Resolution on R2 and R9)
- $\mathbf{R22}:C2W$ (Resolution on R2 and R9)
- $\mathbf{R23}:\Phi$ (Resolution on R22 and R13)

Since we end up in an empty clause, the solution to the entailment is true. In other words $KB \models C2W$

2.

We consider the following 16 literals to construct the propositional logic for 4-Queens problem.

```
QA1, QA2, QA3, QA4
QB1, QB2, QB3, QB4
QC1, QC2, QC3, QC4
QD1, QD2, QD3, QD4
```

Note: For example, QB3 = true, refers to a queen present in position B3.

Knowledge Base (KB):

• Constraints to capture at least one Queen in each row.

```
egin{array}{llll} {f z1}: QA1 & \lor & QB1 & \lor & QC1 & \lor & QD1 \\ {f z2}: QA2 & \lor & QB2 & \lor & QC2 & \lor & QD2 \\ {f z3}: QA3 & \lor & QB3 & \lor & QC3 & \lor & QD3 \\ {f z4}: QA4 & \lor & QB4 & \lor & QC4 & \lor & QD4 \\ \end{array}
```

• There are no repeating queens in any row.

• There are no repeating queens in any column.

• There are no repeating queens in any diagonal.

$\begin{array}{lll} \mathbf{D1}: & \neg QA1 \ \lor \ \neg QB2 \\ \mathbf{D2}: & \neg QA1 \ \lor \ \neg QC3 \\ \mathbf{D3}: & \neg QA1 \ \lor \ \neg QD4 \end{array}$	$\begin{array}{llll} {\bf D15}: & \neg QB2 \ \lor \ \neg QA3 \\ {\bf D16}: & \neg QB2 \ \lor \ \neg QC3 \\ {\bf D17}: & \neg QB2 \ \lor \ \neg QD4 \\ \end{array}$
$\begin{array}{lll} \mathbf{D4}: & \neg QB1 \ \lor \ \neg QA2 \\ \mathbf{D5}: & \neg QB1 \ \lor \ \neg QC2 \\ \mathbf{D6}: & \neg QB1 \ \lor \ \neg QD3 \end{array}$	$\begin{array}{lll} {\bf D18:} & \neg QC2 \ \lor \ \neg QB3 \\ {\bf D19:} & \neg QC2 \ \lor \ \neg QA4 \\ {\bf D20:} & \neg QC2 \ \lor \ \neg QD3 \\ \end{array}$
$\begin{array}{lll} \mathbf{D7}: & \neg QC1 \ \lor \ \neg QB2 \\ \mathbf{D8}: & \neg QC1 \ \lor \ \neg QA3 \\ \mathbf{D9}: & \neg QC1 \ \lor \ \neg QD2 \end{array}$	$\begin{array}{lll} \textbf{D21}: & \neg QD2 \ \lor \ \neg QC3 \\ \textbf{D22}: & \neg QD2 \ \lor \ \neg QB4 \end{array}$
$\begin{array}{lll} \mathbf{D10}: & \neg QD1 \ \lor \ \neg QC2 \\ \mathbf{D11}: & \neg QD1 \ \lor \ \neg QB3 \\ \mathbf{D12}: & \neg QD1 \ \lor \ \neg QA4 \end{array}$	$\begin{array}{llllllllllllllllllllllllllllllllllll$
$\begin{array}{lll} \mathbf{D13}: & \neg QA2 \ \lor \ \neg QB3 \\ \mathbf{D14}: & \neg QA2 \ \lor \ \neg QC4 \end{array}$	$\begin{array}{lll} \mathbf{D26}: & \neg QC3 \lor \neg QB4 \\ \mathbf{D27}: & \neg QC3 \lor \neg QD4 \\ \mathbf{D28}: & \neg QCD \lor \neg QC4 \end{array}$

Note: I have tried to combine several clauses using AND into a row/column clause for better readability and tractability. If we write them separately, we get a total of 80 different clauses.

- 2 a. Table. 1 gives detailed calculations of DPLL algorithm with NO heuristics.
- 2 b. Table. 2 gives detailed calculations of DPLL algorithm with NO heuristics.

Note: Total number of iterations in naive DPLL is 2*41 = 82. Whereas, the DPLL with heuristics takes 37 iterations, which is almost a factor of half reduction in number of iterations.

3.

We consider the following literals to define the propositional logic for the tic-tac-toe problem.

• Positions of player X and O and void spaces (?):

X11, X12, X13	O11, O12, O13	?11, ?12, ?13
X21, X22, X23	O21, O22, O23	?21, ?22, ?23
X31, X32, X33	O31, O32, O33	?31, ?32, ?33

• Move positions:

```
moveX11, moveX12, moveX13

moveX21, moveX23, moveX31, moveX31, moveX32, moveX33
```

• canWin positions:

• forcedMove positions:

```
forcedMoveX11, forcedMoveX12, forcedMoveX13 forcedMoveX21, forcedMoveX22, forcedMoveX31, forcedMoveX32, forcedMoveX31
```

Knowledge Base (KB):

• We need at least one move by X:

c1 :move X11 \lor move X12 \lor move X21 \lor move X22 \lor move X23 \lor move X31 \lor move X32 \lor move X33

• can Win moves for X's and O's:

Completing a row pattern:

Table 1: DPLL with NO Heuristics

Comments	QA1 = T	QA2 = F by C1	QA3 = F by C1	QA4 = F by C1	QB2 = F by R1	QB2 = F by D1	QB3 = T	QB4 = F by C2	QC1 = F by R1	QC2 = F by D18	QC3 = F by R3	BackTrack to Iter. 7	QB3 = F	QB4 = T	QC1 = F by R1	QC2 = T	QC3 = F by C3	QC4 = F by C3	QD1 = F by R1	QD2 = F by R2	QD3 = F by D20	QD4 BackTrack To Iter. 16	QC2 = F	QC3 = F by D26	BackTrack to Iter. 1	QA1 = F	QA2 = T	QA3 = F by C1	QA4 = F by C1	QB1 = F by D4	QB2 = F by R2	QB3 = F by D13	QB4 = T	QC1 = T	QC2 = F by C3	QC3 = F by C3	QC4 = F by C3	QD1 = F by R1	QD2 = F by D9	\square	QD4 = F by C4
QD4	-	ı	,	,	,	,	ı	1	,	,	1	ı	ı	ı	-	1	ı	,		,	,	দ	1	1	,	ı	ı	ı	-	ı	ı	ı	ı	ı	ı	1	ı	ı	ı	ĮŢ.	[파
QD3	-	1	1	1		1	ı	ı	1	1	1	ı	1	ı	-	-	1	1		1	ഥ	Ē	1	-	1	1	1	1	1	ı	ı	1	1	-	1	1	1	1	ı	T	П
QD2	-	ı	ı	1	ı	ı	ı	1	1	1	ı	ı	ı	ı	-	-	ı	1	1	দ	দ	দ	ı	1	ı	ı	ı	ı	1	ı	1	ı	ı	-	ı	ı	1	ı	ഥ	ъ	ĹΉ
QD1	ı	ı	ı	1	ı	1	ı	ı	ı	1	ı	ı	ı	ı	-	ı	1	1	H	ГH	ГH	ГH	ı	1	ı	ı	ı	ı	-	ı	ı	ı	ı	ı	1	ı	-	F	ഥ	F	দ
QC4	-	1	1	1			1		1		1	ı	ı	1	-	-	1	দ	Ē	দ	伍	伍		1	ഥ	ı	ı	1	-	ı	ı	ı	ı	-	1	1	F	F	ഥ	H	[파
QC3	-	ı	ı	1	ı	1	ı	ı	1		দ	ĽΉ	ı	ı	-	-	দ	দ	ഥ	দ	দ	댄	ı	H	ĮΉ	ı	ı	ı	-	ı	ı	ı	ı	ı	ı	দ	F	H	ഥ	ГŦ	ഥ
QC2	-	ı	ı	1			ı	ı	1	Ή	দ	দ	ı	ı	-	${ m L}$	Т	L	L	Т	L	Т	দ	T	T	ı	ı	ı	-	ı	ı	ı	ı	ı	ഥ	দ	F	F	ഥ	H	ĹΉ
QC1	-	ı	ı	1	ı	ı	ı	1	ഥ	ഥ	ഥ	ĹΉ	ı	ı	F	F	ഥ	দ	ഥ	দ	দ	댄	ĮΉ	H	ĮΉ	ı	ı	ı	-	ı	ı	ı	ı	${ m L}$	T	T	${ m L}$	\mathbf{T}	Н	L	H
QB4	-	1	1				ı	ഥ	ഥ	ഥ	ഥ	ĹΉ	1	Τ	${ m L}$	${ m L}$	T	T	L	L	L	L	T	T	T	ı	1	1	-	ı	ı	1	T	${ m L}$	T	T	${ m L}$	\mathbf{T}	Н	L	H
QB3	-	1	1	-	1	1	H	Т	Н	Н	Т	Τ	Ŀı	Ē	Ł	H	Ē	Ē	ഥ	Ē	দ	Ŀı	Ē	H	Ŀı	1	1	1	-	ı	1	Ŀ	Ŀ	H	Ē	Ē	F	H	伍	ഥ	ഥ
QB2	-	1		-		H	伍	H	H	দ	ĿΊ	দ	দ	দ	F	F	দ	H	H	H	H	Г	দ	F	দ	1	ı	1	-	ı	伍	দ	H	F	দ	H	F	F	伍	H	ĹΤι
QB1	-	ı	ı	1	F	F	F	F	FI	F	H	H	H	H	А	А	H	F	F	FI	FI	H	낸	F	H	ı	-	-	-	伍	H	Н	Э	А	H	H	F	F	F	H	ĮΞ
QA4	ı	ı	ı	দ	伍	দ	伍	ഥ	দ	伍	뇬	দ	দ	দ	Н	Н	ഥ	দ	দ	দ	দ	দ	দ	ഥ	伍	ı	-	-	F	ഥ	ഥ	伍	伍	H	ഥ	দ	Ł	F	ഥ	F	ĮΞι
QA3	_	1	伍	伍	দ	냰	伍	H	놴	伍	Ħ	Ŀ	Ŀ	H	F	F	H	H	Ή	H	H	H	H	F	H	1	1	H	F	ഥ	ഥ	H	F	F	H	H	F	F	伍	H	ഥ
QA2	-	ഥ	ഥ	ഥ	ഥ	伍	ഥ	ഥ	ഥ	ഥ	ഥ	ĹΉ	ĹΉ	দ	F	Ā	ഥ	ഥ	ഥ	ഥ	দ	댄	ĹΉ	H	ĹΉ	1	T	L	m L	L	П	T	T	m L	T	L	${ m L}$	T	H	L	H
QA1	Т	Τ	L	Т	П	L	L	П	Т	L	Т	Т	Т	Т	L	Τ	Т	Т	L	Т	L	L	Т	T	Т	ഥ	ഥ	ഥ	F	ഥ	ഥ	伍	Ħ	F	ഥ	ഥ	F	F	ഥ	H	[I
Iteration No	1	2	33	4	ಬ	9	7	∞	6	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41

Table 2: DPLL with Pure Symbol and Unit Clause heuristics

Comments	Q1A = T	Unit Clause: $QA2 = F \text{ by } C1$	Unit Clause: $QA3 = F \text{ by } C1$	Unit Clause: $Q4A = F \text{ by C1}$	Unit Clause : $QB1 = F \text{ by } R1$	Unit Clause: $QB2 = F \text{ by } D1$	Unit Clause : $QC1 = F \text{ by } R1$	Unit Clause: $QC3 = F \text{ by } D2$	Unit Clause : $QD1 = F \text{ by } R1$	Unit Clause : $QD4 = F \text{ by } D3$	Q3B = T	Unit Clause: $QB4 = F \text{ by } C2$	Unit Clause: $QC2 = F \text{ by D18}$	Unit Clause: QD3 = F by R3	Q4C = F Backtrack to Iter. 11	Q3B = F	Q4B = T	Unit Clause: $QC4 = F \text{ by } R4$	Unit Clause: $QD2 = F \text{ by } D22$	Unit Clause: $QC2 = F \text{ by D18}$	Q3D = T/F Backtrack to Iter. 1	QA1 = F	QA2 = T	Unit Clause: $QA3 = F \text{ by } C1$	Unit Clause: $QA4 = F \text{ by C1}$	Unit Clause: $QA4 = F by D4$	Unit Clause: $QB2 = F \text{ by } R2$	Unit Clause: $QB3 = F \text{ by D13}$	Unit Clause: $QC2 = F \text{ by } R2$	Unit Clause: $QC4 = F$ by D14	Unit Clause: $QD2 = F \text{ by } R2$	$\mathrm{QB4} = \mathrm{T}$	Unit Clause: $QC3 = F$ by D26	Unit Clause: $QD4 = F \text{ by } D27$	Τ	QD1 =	Pure Symbol: QD3 = T
QD4	1	1	1	1	1	1	1	1	1	Ħ	ഥ	Ħ	Ħ	Ħ	Ħ	দ	দ	Ħ	Ħ	H	1	1	1	1	1	1	1	1	-	1	1	1	1	ഥ	F	F	伍
QD3	1	1	1		1	1	1	-	1	1	1	1	1	Ħ	1			1	1	1	1	1	1	1	1	1	1	1	-	1	1	-	1	1	1	1	П
QD2	1	-	-	1	1	1	1	-	1	1	ı	1	1	1	1	1	1	1	FI	FI	-	1	1	1	1	ı	1	ı	-	ı	H	F	F	H	Э	F	H
QD1	ı	ı	ı	ı	ı	ı	ı	ı	দ	দ	দ	দ	দ	দ	দ	ഥ	ഥ	দ	দ	দ	ı	ı	ı	ı	ı	ı	1	ı	-	ı	ı	1	ı	ı	ı	H	ഥ
QC4	ı	ı	ı	ı	ı	ı	ı	1	1	1	ı	1	1	ı	1	ı	1	দ	দ	দ	ı	1	1	1	ı	ı	1	ı	-	দ	ഥ	দ	দ	ഥ	H	H	Έų
QC3	1	1	1	'	'	,	,	伍	ഥ	দ	Ŀı	দ	됴	Ē	됴	ഥ	ഥ	됴	Ē	Ē	1	1	'	1	1	1	1	1	-	1	1	,	Ē	Ē	Ŀ	Ŀ	ഥ
QC2	ı	ı	1	1	1	1	1	1	1	1	ı	1	ГH	Έų	1			1	1	L	1	1	1	1	1	ı	1	ı	H	伍	伍	ГH	ĹΉ	伍	Ъ	伍	Ħ
QC1	1	1	'	,	,	,	伍	দ	Ŀ	Ŀ	দ	ъ	H	ГŦ	Ħ	伍	伍	ГŦ	ГŦ	ГH	'	,	,	,	,	,	1	1	-	1	1	,	1	1	Τ	T	H
QB4	ı	1	1	'	'	1	,	'	1	1	ı	দ	দ	দ	1	1	L	L	L	L	1	1	1	1	1	ı	1	ı	-	ı	ı	L	L	T	T	T	H
QB3	1	1	'	'	'	'	'	'	'	1	H	L	L	L	L	伍	伍	伍	Ē	Ē	'	'	'	'	'	'	1	냰	F	놴	伍	Ħ	Ē	伍	H	F	ഥ
QB2	ı	ı	ı	1	1	伍	伍	伍	দ	দ	ഥ	দ	伍	伍	伍	伍	伍	伍	伍	伍	ı	1	1	1	1	ı	দ	ഥ	H	ഥ	ഥ	দ	伍	ഥ	H	伍	Ħ
QB1	ı	1	1	,	দ	伍	伍	দ	দ	দ	伍	ГŦ	Ħ	দ	Ħ	伍	伍	Ē	Ē	伍	ı	,	,	1	1	伍	দ	দ	F	দ	伍	Ħ	Ē	伍	H	伍	伍
QA4	1	1	'	ഥ	ഥ	ഥ	伍	伍	Ē	伍	Ŀı	伍	伍	Ē	伍	ഥ	伍	伍	Ŀ	Ŀ	'	'	'	'	Ē	Ŀı	伍	Ŀı	H	Ē	ഥ	দ	Ē	Ŀı	ഥ	H	ഥ
QA3	1	1	ഥ	ഥ	দ	ഥ	ഥ	দ	伍	দ	দ	দ	দ	দ	দ	ഥ	ഥ	দ	ഥ	ഥ	'	-	'	ഥ	দ	দ	দ	দ	ഥ	দ	দ	দ	ഥ	দ	দ	伍	[Ŧı
QA2	1	ഥ	ഥ	ഥ	ഥ	ഥ	ഥ	ഥ	ഥ	됴	ഥ	됴	ഥ	됴	ഥ	ഥ	ഥ	ഥ	ഥ	ഥ	1	-	L	L	L	T	L	T	T	L	Τ	L	L	L	T	T	H
QA1	L	L	L	L	L	L	L	L	L	L	Τ	Τ	Τ	Τ	Τ	L	L	L	L	L	L	됴	দ	দ	দ	দ	দ	দ	伍	দ	দ	দ	দ	ഥ	伍	伍	Œ
Iteration	П	2	က	4	ಬ	9	7	∞	6	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37

```
\mathbf{c8}: X11 \land ?12 \land X13 \implies \operatorname{canWin}X12
                                                                          \mathbf{c11}: O11 \land ?12 \land O13 \implies \mathrm{canWin}O12
\mathbf{c9}: X21 \land ?22 \land X23 \implies \mathrm{canWin}X22
                                                                          \mathbf{c12}: O21 \land ?22 \land O23 \implies \mathrm{canWin}O22
\mathbf{c10}: X31 \land ?32 \land X33 \implies \mathrm{canWin}X32
                                                                          \mathbf{c13}: O31 \land ?32 \land O33 \implies \mathrm{canWin}O32
\mathbf{c14}: ?11 \land X12 \land X13 \Longrightarrow \mathrm{canWin}X11
                                                                          c17:?11 \land O12 \land O13 \implies canWinO11
\mathbf{c15}: ?21 \land X22 \land X23 \implies \mathrm{canWin}X21
                                                                          \mathbf{c18}: ?21 \land O22 \land O23 \implies \mathrm{canWin}O21
\mathbf{c16}: ?31 \land X32 \land X33 \implies \mathrm{canWin}X31
                                                                          \mathbf{c19}: ?31 \land O32 \land O33 \implies \mathrm{canWin}O31
Completing a column pattern:
                                                                          \mathbf{c23}: O11 \land O21 \land ?31 \implies canWinO31
\mathbf{c20}: X11 \land X21 \land ?31 \implies canWinX31
\mathbf{c21}: X12 \land X22 \land ?23 \implies canWinX23
                                                                          \mathbf{c24}: O12 \land O22 \land ?23 \implies canWinO23
\mathbf{c22}: X13 \land X23 \land ?33 \implies canWinX33
                                                                          \mathbf{c25}: O13 \land O23 \land ?33 \implies canWinO33
\mathbf{c26}: X11 \land ?21 \land X31 \implies canWinX21
                                                                          \mathbf{c29}: O11 \land ?21 \land O31 \implies canWinO21
\mathbf{c27}: X12 \land ?22 \land X23 \implies canWinX22
                                                                          \mathbf{c30}: O12 \land ?22 \land O23 \implies canWinO22
\mathbf{c28}: X13 \land ?23 \land X33 \implies canWinX23
                                                                          \mathbf{c31}: O13 \land ?23 \land O33 \implies canWinO23
\mathbf{c32}: ?11 \land X21 \land X31 \implies canWinX11
                                                                          \mathbf{c35}: ?11 \land O21 \land O31 \implies canWinO11
\mathbf{c33}: ?12 \land X22 \land X23 \implies canWinX12
                                                                          \mathbf{c36}: ?12 \land O22 \land O23 \implies canWinO12
\mathbf{c34}: ?13 \land X23 \land X33 \implies canWinX13
                                                                          \mathbf{c37}: ?13 \land O23 \land O33 \implies canWinO13
Completing a diagonal pattern:
\mathbf{c38}: X11 \land X22 \land ?33 \implies canWinX33
                                                                          \mathbf{c40}: O11 \land O22 \land ?33 \implies canWinO33
\mathbf{c39}: X13 \land X22 \land ?31 \implies canWinX31
                                                                          \mathbf{c41}: O13 \land O22 \land ?31 \implies canWinO31
\mathbf{c42}: X11 \land ?22 \land X33 \implies canWinX22
                                                                          \mathbf{c44}: O11 \land ?22 \land O33 \implies canWinO22
\mathbf{c43}: X13 \land ?22 \land X31 \implies canWinX22
                                                                          \mathbf{c45}: O13 \land ?22 \land O31 \implies canWinO22
\mathbf{c46}: ?11 \land X22 \land X33 \implies canWinX11
                                                                          \mathbf{c48}: ?11 \land X22 \land X33 \implies canWinO11
\mathbf{c47}: ?13 \land X22 \land X31 \implies canWinX13
                                                                          \mathbf{c49}: ?13 \land X22 \land X31 \implies canWinO13
\mathbf{c50}: canWinX11 \lor canWinX12 \lor canWinX13
     \lor canWinX21 \lor canWinX22 \lor canWinX23
     \lor canWinX31 \lor canWinX32 \lor canWinX33 \implies canWinX
\mathbf{c51} : \operatorname{canWin}O11 \lor \operatorname{canWin}O12 \lor \operatorname{canWin}O13
    \vee \operatorname{canWin}O21 \vee \operatorname{canWin}O22 \vee \operatorname{canWin}O23
    \vee \operatorname{canWin}O31 \vee \operatorname{canWin}O32 \vee \operatorname{canWin}O33 \implies \operatorname{canWin}O
```

\bullet Conditions for winning moves by X:

\bullet Conditions for forcedMoves by X: