

# **USING ND-LAPLACE TO TRAIN PRIVACY-PRESERVING CLUSTER ALGORITHMS ON DISTRIBUTED N-DIMENSIONAL DATA**

by

**Tjibbe van der Ende**

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Student number: 852372917  
Course code: IMA0002  
Thesis committee: Dr. Ir. Mina Sheikhalishahi (chairman), Open University  
Dr. Ir. Clara Maathuis (supervisor), Open University

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# 1

## INTRODUCTION

### 1.1. RESEARCH QUESTIONS

#### Main question:

*How can the nD-Laplace algorithm be applied in training privacy-preserving clustering algorithms on distributed n-dimensional data?*

1. RQ1: How can 2D-Laplace be used to protect the data privacy of 2-dimensional data which is employed for training clustering algorithms?
2. RQ2: How can 3D-Laplace be extended to protect the data privacy of n-dimensional data which is employed for training clustering algorithms?
3. RQ3: What is the impact of different privacy budgets, dataset properties, and other clustering algorithms on the research conducted for research question 2?

# 2

## LITERATURE REVIEW

The current chapter provides an explanation of the theoretical aspects of the research and evaluates similar and previous studies. First, an explanation is given of the concept of Differential Privacy and its various types. Next, we examine cluster algorithms and the methods for evaluating their performance. Regarding the related literature, we first look at studies that have used differential privacy and then the studies that have combined it with cluster methods.

## 2.1. DIFFERENTIAL PRIVACY

In practice, data is often sent to a central storage point. This requires trust, and because all data is collected in one place, the risk of private data leakage becomes very high. By applying differential privacy, noise can be added to the data to protect it. This principle is illustrated in figure 2.1 with the following actors:

1. Trusted curator: The system that receives data from users. It is assumed in this setting that the system is trustworthy and that the data is securely stored.
2. Adversary: An adversary is someone who uses the data. This could be, for example, a data scientist who wants accurate results, or an attacker who wants to obtain as much data as possible.
3. The users are clients (for example, websites or mobile apps) who entrust their data to a central server.

With the introduction of differential privacy, the privacy of a user would be ensured (to a certain extent). This will be further explained in the next section.

Although differential privacy solves many problems (as mentioned earlier in the introduction), it remains difficult to calibrate the mechanism. There is an important trade-off between utility and privacy for the adversary, where a data scientist wants accurate data while the noise must be sufficient to prevent an attacker from obtaining too much information. For this reason, the following chapters will be devoted to outlining the mathematical background of differential privacy. We will examine which factors influence this calibration and whether other methods contribute to it. Afterward, we will also further explain other types of differential privacy (local and geo-indistinguishable) in the same way.

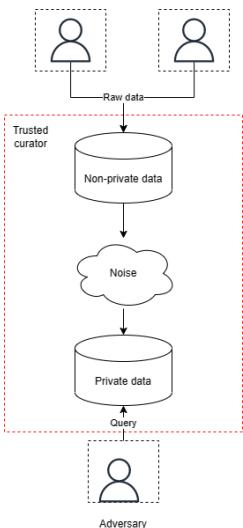


Figure 2.1: General approach for setting up (central) differential privacy.

### 2.1.1. DEFINITIONS

We examine the different notations and types of differential privacy we consider in this research.

#### $\epsilon$ -DIFFERENTIAL PRIVACY

Dwork et al formulated the notion of privacy as: Participating in a database should not significantly increase the risk of an individual's privacy being compromised [Dwork \[2006\]](#). This is mathematically formulated in the same research with the name **Differential Privacy (DP)**. Using the definition of privacy, it is formulated as the maximal possible change when adding or removing a single record [\[Dwork, 2006; Friedman and Schuster, 2010\]](#). This is reflected using the formal mathematical formulation as formulated by dwork et al:

$$\Pr[K(D_1) \in S] \leq \exp(\epsilon) \times \Pr[K(D_2) \in S] \quad (2.1)$$

So, given a randomization function  $K$ , it gives  $\epsilon$ -differential privacy if dataset  $D_1$  and  $D_2$  are differing at most one element [\[Dwork, 2006\]](#). The  $\epsilon$  determines the amount of noise (privacy budget) [\[Friedman and Schuster, 2010\]](#). The lower the value of  $\epsilon$  means, the higher the privacy guarantee. In this regard, it is important for a method that ensures differential privacy to take this into account. For this reason, a common way to determine the  $\epsilon$  is to calculate the sensitivity. This value is calculated based on the impact of a function or query on the data. For example, if there is a method called *sum* for the summation of data points, the sensitivity of the method is 1. This is because removing one data point would greatly affect the outcome and  $\epsilon$ -differential privacy could no longer be guaranteed. It is also mathematically defined by dwork et al:

$$\Delta f = \max_{D_1, D_2} \|f(D_1) - f(D_2)\|_1 \quad (2.2)$$

#### $(\epsilon, \delta)$ -DIFFERENTIAL PRIVACY

The formal notion of differential privacy has only  $\epsilon$  input. This formulation is really strict, but most methods relax this a little which is defined as  $(\epsilon, \delta)$ .

$$\Pr[K(D_1) \in S] \leq \exp(\epsilon) \times \Pr[K(D_2) \in S] + \delta \quad (2.3)$$

This means the sensitivity (now denoted as delta  $\delta$ ) is used to loosen up the definition. The purpose of  $\epsilon$  now is to calibrate the desired amount of privacy. To some extent, the delta represents the probability of the algorithm leaking information [\[Aitsam, 2021\]](#). With  $\epsilon$ -differential privacy, there would be no difference in the case of information leakage (delta = 0). However, with  $(\epsilon, \delta)$ -differential privacy, the information can leak up to the probability of delta.

### $\epsilon$ -LOCAL DIFFERENTIAL PRIVACY

As the name suggests, **Local Differential Privacy (LDP)** is executed on the client-side instead of on the server, as was the case in figure 2.1. This is illustrated in figure 2.2. Local differential privacy was introduced to remove the "trusted" curator, preventing sensitive data leakage even if an attacker gains access to the dataset [Xiong et al., 2020a]. The definition for LDP is the same as for equation 2.3 with  $\delta = 0$  being equal to equation 2.1.

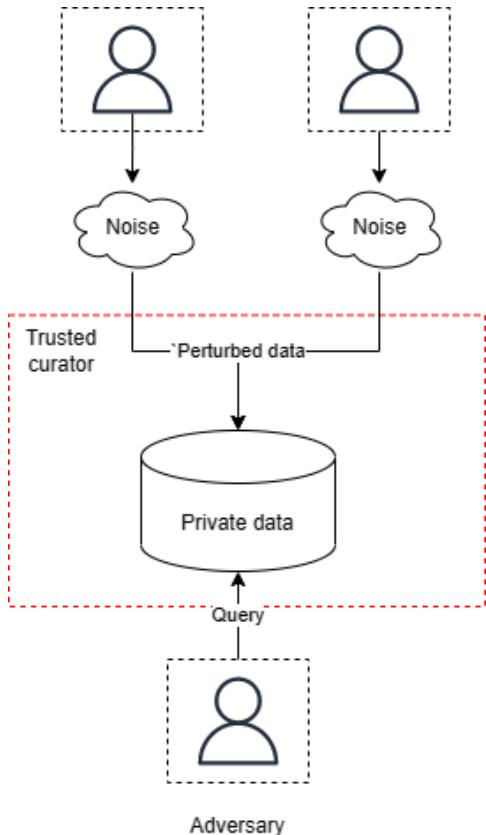


Figure 2.2: Local differential privacy, which moves the noise-adding step to the client-side.

### Describe issues with local differential privacy

Local differential privacy has two different frameworks. It can be set up as interactive or non-interactive [Xiong et al., 2020a]. Where interactive can be distinguished as either sequentially interactive or fully interactive <sup>1</sup>:

<sup>1</sup>Image source: [https://www.majos.net/focs\\_19\\_talk.pdf](https://www.majos.net/focs_19_talk.pdf)

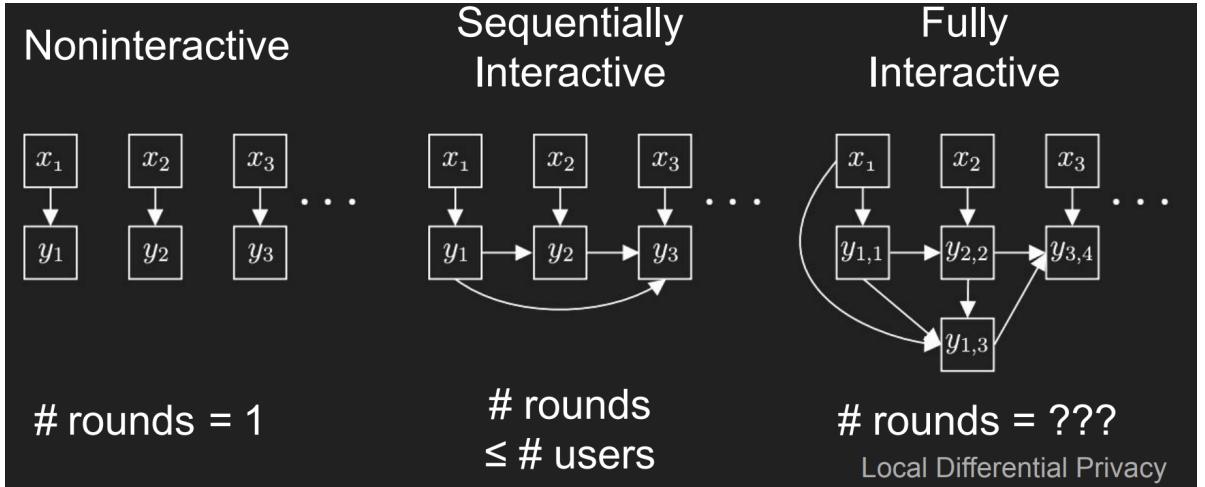


Figure 2.3: Non-interactive versus sequentially interactive versus fully interactive local differential privacy [Joseph et al., 2019]

In the non-interactive framework, the privacy mechanism operates on a single data instance (e.g. users' smartphone) without requiring any interaction with other clients or between data-points. This means that each data point is perturbed independently, and the privacy guarantees are achieved without any communication or coordination between the data points. In the interactive framework, the privacy mechanism involves a process of iterative interactions between users' clients. Each data point takes into account the responses of other data points in order to determine its own perturbation or privacy level. The interactions can occur through a communication channel, where data points exchange information with each other or a trusted mediator. The iterative nature of the interactive framework allows for refining and improving the privacy guarantees based on the global knowledge of the data set.

In figure 2.3 the different interactive types are visualised, where the most notable difference is the around of rounds [Xiong et al., 2020b]. A sequentially interactive mechanism passes at most one message between data-points. A fully interactive mechanism could provide an unknown and unlimited amount of messages between data-points. As the interactive mechanism has unlimited access to other data-points, it is most suitable for practical appliances [Xiong et al., 2020b]. This is because the mutual correlation between the data points also contributes [Wang et al., 2020].

#### $\epsilon$ -GEO-INDISTINGUISHABILITY

The last and for this study's most important type of differential privacy is **Geo-indistinguishability (GI)**. GI can be applied to preserve the privacy using a differential privacy method specific to spatial data [Andrés et al., 2012]. Consider a local based system (LBS) that provides a map service to its users (e.g. Google Maps). The users can query the LBS for a route from their current location to a destination. This information is strictly personal, and the LBS should not be able to track the user's location. Instead, the LBS should only receive an approximate location enough to provide the service at an adequate level. This happens completely locally, so there is not a central server inbetween like classical DP.

A client (e.g. smartphone) is eligible to define a privacy radius  $r$  around their location  $x$ . Within this radius a fake location is generated, based on the Euclidean distance  $d$ . This is

based on the privacy level  $l$ , which is combined with  $r$  to obtain a privacy budget  $\epsilon$  using the formula  $\epsilon = \frac{l}{r}$ :

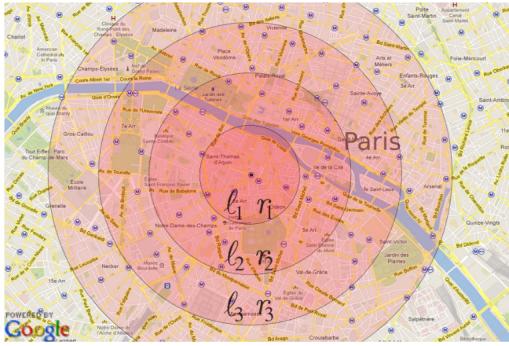


Figure 2.4: Visualisation of radius  $r$  with  $l$  to provide Geo-indistinguishability mechanism [Andrés et al., 2012]

The formula for GI to measure if an algorithm preserves  $\epsilon$ -geo-indistinguishability can be expressed as:

$$K(x)(y) \leq e^{\epsilon * d(x, x')} K(x')(y) \quad (2.4)$$

Where  $K$  is a probability method reporting  $x, x' \in X$  as  $z \in Z$ . The idea of this algorithm looks a lot like that of differential privacy using the La Place method; but includes distance. The intuition for this is that it displays the distinguishability level between two secret locations/points  $x$  and  $x'$  [Chatzikokolakis et al., 2015]. An extension of this is called  $d_x$ -privacy and is a more general notation of distance-aware differential privacy. Their definition for GI is, therefore,  $d_2$ -privacy, but is essentially the same as the proof provided for GI.

### 2.1.2. MECHANISMS

In this section, we will explain the various mechanisms that are used to achieve differential privacy. The different mechanisms are not limited to the type of DP and some can be applied to multiple types.

#### RANDOMIZED RESPONSE MECHANISM

The random response method is a relatively simple method and was first applied in 1965 by Warner et al. It was originally used to mask the answers of individuals by randomly switching the answers with predictable randomness [Warner, 1965]. Therefore, the method is mainly used for categorical data. This method satisfies its own set of requirements for LDP [Xiong et al., 2020a], which differs from the formal definition that was mentioned earlier

Since then, it is still one of the better-known methods, and larger organizations such as Google use it [Erlingsson et al., 2014]. They have named their extension RAPPOR and expanded it with bloom filters to be able to collect numerical data as well. It has been possible to ensure  $\epsilon$ -differential privacy in this way, and it is also possible to preserve LDP [Xiong et al., 2020a].

#### LAPLACE MECHANISM

The method that was originally proposed in the differential privacy paper by Dwork et al. is the Laplace algorithm [Dwork, 2006]. Therefore, the method works by configuring the  $\epsilon$  and  $\delta$ . A shorthand definition is provided by Rey et al [Xiong et al., 2020a]:

$$M(f(x), \epsilon) = f(x) + (Z_1, \dots, Z_d) \quad (2.5)$$

The mechanism is based on the Laplace distribution with scale  $\lambda f / \epsilon$ , where  $\lambda f$  is the same as in equation 2.2. Therefore, the Laplace mechanism is tightly linked to the definition of pure-dp and so it is  $\epsilon$ -differential privacy [Dwork, 2006]. It is also suitable for preserving LDP [Xiong et al., 2020a]. One disadvantage is that sensitivity is always required, and this parameter can sometimes be difficult to configure. Especially when there is no clear function and the entire dataset is perturbed. This can make it challenging to find the right balance between ensuring privacy and utility.

To this end, the sensitivity can be calculated in two forms: global and local. Global sensitivity is calculated over two different datasets and is part of the original definition of DP [Dwork, 2006]. It is called global because it is independent of the queried dataset. Usually, this is not the desired situation, since local sensitivity always has more context of the dataset in question [Nissim et al., 2007]. As a result, the trade-off for noise is much more precise, and the balance between utility and privacy is much better. This local sensitivity definition [Nissim et al., 2007].

$$LS(f, x) = \max_{x' : d(x, x') \leq 1} |f(x) - f(x')| \quad (2.6)$$

A methodology to calculate the local sensitivity is by using the smooth sensitivity method, which was proposed by the same authors [Nissim et al., 2007]. This method aims at smoothing out the local sensitivity by focusing on reducing the noise. The goal is to smooth out the amount of noise, to reduce the risk of something being revealed in the data. Due to this, a disadvantage of this method is that it can be computationally expensive. Also, the introduction of this method makes Laplace preserve  $(\epsilon, \delta)$ -DP instead of pure  $\epsilon$ -DP.

## 2.2. CLUSTERING

Explain why these three algorithms

### 2.2.1. METHODS

In this chapter, a brief description is given of each clustering algorithm on how it works. The different parameters for the algorithm are also highlighted.

#### K-MEANS

The K-Means algorithm is based on the algorithm of Loyd et al. The starting point is determined randomly by choosing  $k$  number of centroids [Lloyd, 1982]. Each point is then assigned to a centroid based on the Euclidean distance. A new centroid is then determined based on the cluster average. This is repeated until the given number of iterations is reached or until the results are stable.

**Choosing K:** The most important parameter of the K-Means algorithm is the value of  $k$ . This value determines the number of clusters to consider and has a big influence on the results [Ahmed et al., 2020].

The first method is called an "elbow" plot [Kodinariya and Makwana, 2013]. This method can be used to determine the best  $k$  by applying the algorithm multiple times and estimating the best  $k$ .

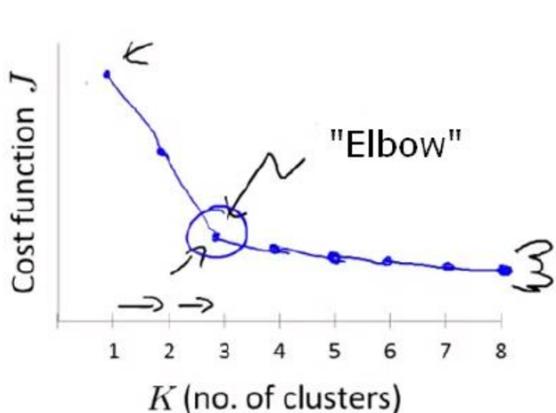


Figure 2.5: Illustration of determining  $k$  using the "elbow" method [Kodinariya and Makwana, 2013]

In this situation, a Silhouette plot can be used to determine the  $k$ . This method uses the Silhouette coefficient for each cluster [SAPUTRA et al., 2020]. It takes into consideration the separation and cohesion of the cluster. The plot can then be made by plotting the silhouette coefficient on the y-axis and  $k$  on the x-axis [SAPUTRA et al., 2020]. Then the  $K$  with the highest coefficient can be selected based on that.

Another popular method is the Gap statistic method [Yuan and Yang, 2019]. It compares the total within-cluster variation for different values of  $k$  with their expected values under a null reference distribution of the data [Tibshirani et al., 2001]. A practical appliance of this method uses a line plot for comparing the  $k$ -value and gap value [Yuan and Yang, 2019]. Based on the line's visual change, someone can select the best  $k$ .

All in all, there is no fixed method to choose a good  $k$  for K-Means. The elbow method is common in the existing literature and is very popular due to its simplicity. However,

one disadvantage is that it can be challenging to determine the "elbow" point, as it is not always present [Kodinariya and Makwana, 2013]. In that case, the silhouette or gap statistic method can be chosen, with the algorithm for silhouette being the most obvious choice due to its simplicity.

#### AFFINITY PROPAGATION

**Affinity Propagation (AP)** is an algorithm that clusters data points by iteratively passing messages between them. Each point sends and receives messages about the attractiveness of other points as cluster centers (exemplars) and the suitability of itself as a center [Keller et al., 2021]. The method was introduced by Frey et al. and does not require any hyperparameters [Frey and Dueck, 2007]. Still, there are important properties that could potentially impact the clustering [Wang et al., 2007].

**Choosing preference( $p$ ):** Indicates the preference for selecting a data point as cluster center [Wang et al., 2007]. It highly influences the number of clusters; a high one would lead to more clusters and a small one to less [Mojane and Machado, 2018]. Depending on the data, a good choice is to set the  $p$  to the median of all data similarities [Wang et al., 2007]. But, the effectiveness of this could be highly influenced based on the dataset. Analyzing the silhouette coefficient [Mojane and Machado, 2018] to validate if the preference is correctly set is possible.

**Choosing damping factor( $lam \in [0, 1]$ ):** The damping factor is used to improve the stability (convergence) of the algorithm [Wang et al., 2007]. By default, this value is 0.5 and can be increased to 1 to reduce the impact of numerical oscillations. This can be applied manually by re-running the algorithm, finding the optimal, or increasing. However, both approaches take a lot of time, especially for bigger datasets [Wang et al., 2007].

To conclude on this, damping is important if big datasets are considered. However, this research does not use large datasets or consider time complexity as a metric. The preference on the other hand could influence the results a lot. For this, the silhouette coefficient can be evaluated to choose the best option. In general, it should be sufficient to take the median.

#### DBSCAN

**Density-based spatial clustering of applications with noise (DBSCAN)** was introduced by [Ester et al.] and works by drawing a radius (neighborhood) around data points. It then groups all points within this radius as clusters. The main advantage is its ability to find arbitrarily shaped clusters and detect outliers [Liu et al., 2012]. To do this, the DBSCAN algorithm uses the inputs  $minPts$ ,  $radius(\epsilon)$  and a distance function [Schubert et al., 2017]. The  $\epsilon$  is used to draw a neighborhood and the  $minPts$  is used as a weight to evaluate which points should be inside the neighborhood. For the distance function, the Euclidean distance is used, to be consistent with the other algorithms.

**Choosing minimum points ( $minPts$ ):** This hyperparameter is considered by a paper written by Sander et al. The work describes a way of calculating this parameter by doing two times the feature amount [Sander et al., 1998]. So, using this approach a dataset with two features will have an  $minPts$  of four. This is confirmed by Schubert et al. to use the default  $minPts = 4$  for a 2-dimensional dataset [Schubert et al., 2017].

**Choosing radius( $\epsilon'$ ):** The desired  $\epsilon'$  (not to be confused with the privacy budget  $\epsilon$ ) can be calculated using the K-NearestNeighbours algorithm [Ester et al.; Schubert et al., 2017]. The general approach for this is to choose a  $K = 2 * N - 1$  (where  $N$  is the number of features) and plot the distance for each point. This can then be plotted using a k-dist plot and the best "elbow" can be chosen for deciding the  $\epsilon$  (similar to choosing the  $k$  for K-means) [Elbatta and Ashour, 2013].

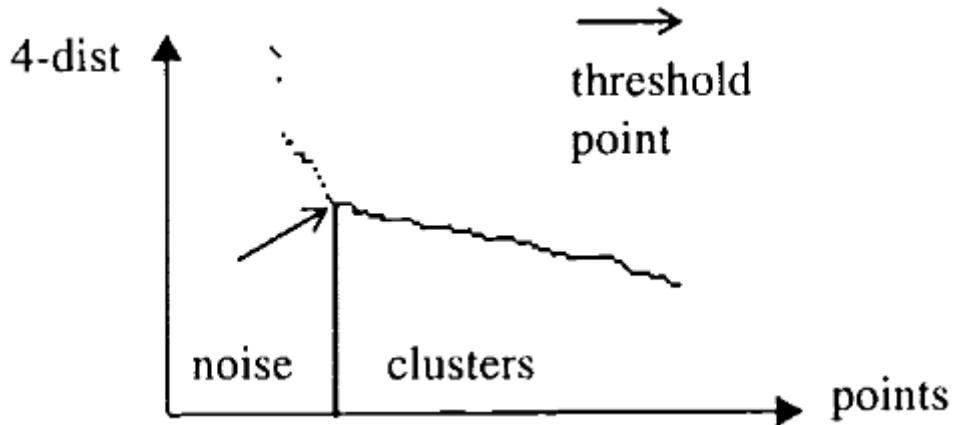


Figure 2.6: K-dist plot example for  $\text{min}Pts = 4$  based on a 2-dimensional dataset [Ester et al.]

Since we employ various types of datasets and algorithms, it is desirable to have an automated method for determining the epsilon (radius) value. The epsilon could possibly influence the results as the clusters are determined by the radius. For this reason, research has also been conducted on an extension of DBSCAN called OPTICS.

This algorithm attempts different epsilon values to achieve the best result [Ankerst et al.]. Instead of directly assigning data points to clusters, certain distance calculations are first stored in a list. For each data point, two values are tracked: core-distance and reachability-distance. These represent the shortest distance to make a data point,  $p$ , a core point and the shortest distance from  $p$  to another core point,  $p'$ , respectively. A specific ordering is maintained to ensure that clusters with higher density values are processed first. By using this ordering, along with the  $\text{min}Pts$  parameter and the two data point attributes, a DBSCAN algorithm can be constructed in an hierarchical way [Schubert et al., 2017].

In conclusion, both K-Means and Affinity Propagation have clear methods for determining the hyperparameters. For K-Means, the elbow method is the most common and for Affinity Propagation, the median is used for the preference. DBSCAN is a little harder due to the variety of datasets and noise altering mechanisms we experiment with. This is why we use OPTICS to determine the best  $\epsilon$ , and choose  $\text{min}Pts$  based on the number of features times two.

### 2.2.2. EVALUATION METHODS

Clustering comparison measures are important in cluster analysis for external validation by comparing clustering solutions to a "ground truth" clustering [Vinh et al.]. These external validity indices are a common way to assess the quality of unsupervised machine learning methods like clustering [Warrens and van der Hoef, 2022]. A method that could be

used for this is the Rand Index [Rand, 1971]. It is a commonly applied method for comparing two different cluster algorithms [Wagner and Wagner]. An improvement of this method is adjusted for chance by considering the similarity of pairwise cluster comparisons [Vinh et al.]. Both the Rand Index (RI) and Adjusted Rand Index (ARI) [Hubert and Arabie, 1985] report a value between 0 and 1. Where 0 is for no-similarity and 1 for identical clusters. Alternatives for RI are the Fowlkes-Mallows Index and Mirkin Metric. However, these two methods have their disadvantages. Respectively, being sensitive to a few clusters and cluster sizes [Wagner and Wagner]. The ARI metric suffers from cluster size imbalance as well, so it only provides not a lot of information on smaller clusters [Warren and van der Hoef, 2022]. Instead, they recommend using the cluster index metric that was proposed by Fränti et al. [Fränti et al., 2014].

Another popular group of methods is the information theoretic-based measures [Vinh et al.]. This metric measures the information between centroids; the higher the value, the better [Vinh et al.]. Mutual Information (MI) is such metric, which calculates the probability of an element belonging to cluster  $C$  or  $C'$ . But, is not easy to interpret as it does not have a maximum value [Wagner and Wagner]. To this end, Normalized Mutual Information (NMI) can be used to report a value between 0 and 1 using the geometric mean [Strehl and Ghosh, 2002]. The metric exists also in an adjusted version as Adjusted Mutual Information (AMI). This works in the same way as for the Adjusted Rank Index (ARI) and is mostly needed if the number of data items is small in comparison to the number of clusters [Vinh et al.].

Besides the external validity measurements for clustering, it is also possible to use internal validation methods. These metrics focus entirely on the intrinsic dataset properties, instead of relying on an external baseline cluster algorithm [Craenendonck and Blockeel]. Assessing two important concepts of clustering: compactness and separation [Hassani and Seidl, 2017]. Both studies, consider three different metrics and measure both concepts at the same time [Hassani and Seidl, 2017]:

1. Calinski-Harabasz Index (CHI) [Caliński and Harabasz, 1974] is used to measure the cluster variance (well-separated clusters) and low variance within the clusters (tightly coupled data). A high score indicates better clustering.
2. Silhouette Index [Rousseeuw, 1987] this metric is similar, by also measuring cohesion within clusters and separation of clusters. However, this metric uses the pairwise distance Hassani and Seidl [2017]. A score of -1 indicates incorrect clustering and +1 for dense clusters Rousseeuw [1987].
3. Davies-Bouldin [Davies and Bouldin, 1979] uses the average distance between centroids. A lower score indicates good clustering.

K-Means scores relatively high for CHI [Craenendonck and Blockeel; Hassani and Seidl, 2017] and SI [Craenendonck and Blockeel]. The same applies to DBSCAN, which scores relatively high on SI and DB due to the sensitivity of noise [Craenendonck and Blockeel].

## EXISTING LITERATURE

Comparable studies with differential privacy use external validation [Sun et al., 2022; Xia et al., 2020]. Their experiment setup uses a so-called non-private cluster algorithm as ex-

ternal validation. This cluster algorithm is trained without the perturbed data and compared with the same clustering algorithm that is trained with perturbed data. Thus, the non-private variant functions as an external validation by providing the ground truth.

They compare the mutual information between a baseline cluster algorithm using **AMI** [Huang et al., 2021] or **NMI** [Sun et al., 2022; Xia et al., 2020]. Another study for evaluating **DP** with **AP** uses both **ARI** and **AMI**. In addition to mutual information and rand index scores, it is also not uncommon to calculate the error between the two cluster algorithm's centroids [Huang et al., 2021; Xia et al., 2020]. These two studies used Relative Error (RE) for this.

## 2.3. LITERATURE REVIEW

In the search for related literature, we focused mainly on (L)DP mechanisms that can be used for general purposes (e.g. not only mean estimation), as this is the most comparable to our mechanism. The related literature is divided into two parts:

1. Differential privacy methods
2. Cluster methods with (L)DP

Afterward, we provide a summary for both in the form of a table. This table includes components such as the type (LDP or DP) and whether there is a public code available.

### 2.3.1. DIFFERENTIAL PRIVACY METHODS

As was discussed in earlier sections, the Laplace method was the first one to establish DP [Dwork, 2006].

The first paper we discuss is provided by Soria-Comas et al. and considers the distribution of the dataset for generating [Soria-Comas and Domingo-Ferrer, 2013]. Their work claims the Laplace mechanism is not optimal for a univariate function and aims at improving it by introducing their mechanism based on Laplace. The proposed method performs slightly better than Laplace on multivariate/multiple queries.

Quan et al. also proposed a new method to extend Dwork et al.'s Laplace algorithm [Geng et al., 2015]. They introduced the staircase mechanism for 1-dimensional noise, which was later extended to support multidimensional data [Geng et al., 2015]. The mechanism aims to improve utility by adding the same level of privacy while adding less noise. It is represented as a staircase-shaped probability density function called Staircase Mechanism (SM). This mechanism accepts three configurable parameters comparable to those of the Laplace mechanism. The authors' work can handle multidimensional numerical data and preserve the  $(\epsilon)$ -differential privacy, where  $\epsilon$  is the privacy budget (see 2.1). In the previous two paragraphs, we have mainly focused on the interesting literature regarding differential privacy. Next, the succeeding paragraphs will mainly center on related literature concerning local differential privacy.

Another paper introduces a new local differential privacy (LDP) mechanism for working with numerical data [Nguyn et al., 2016]. Their primary focus is on estimating means and frequencies, with a particular emphasis on machine learning techniques such as Support Vector Machines (SVM) and linear regression using Empirical Risk Minimization. Initially, the authors analyze Duchi et al.'s method [Duchi et al., 2013] and highlight several shortcomings. To address these issues, they introduce Harmony as a mechanism for LDP perturbation. This mechanism is able to perturb both categorical and numerical data and provides high accuracy for classification and regression tasks. To compare Harmony to other methods they introduce the Hybrid Mechanism (HM), which is a combination of two existing methods for categorical and numerical data. For this purpose they extend other work [Bassily and Smith, 2015] for perturbing multidimensional categorical attributes and use Duchi et al.'s method for numerical data. This allows them to compare their Harmony mechanism to the hybrid mechanism and measure the utility/accuracy differences.

Duchi et al. improved their method by proposing a formalization of the trade-off between statistical utility and (local) privacy, analyzing multiple types of estimation problems [Duchi et al., 2017]. Examples include mean, median, and density estimation. To achieve

this, they use minimax, a technique for finding the worst-case probability distribution. Additionally, they focus on existing work and propose several optimization strategies for it.

Duchi et al.'s method was then extended by adding support for bounded and multi-dimensional data [Wang et al., 2019]. The authors introduce the Piecewise Mechanism (PM) to handle both numeric and categorical data and the Hybrid Mechanism (HM) which combines PM and Duchi et al.'s method for 1-dimensional data. The effectiveness of their method is demonstrated using Support Vector Machines (SVM), linear regression, and mean estimation. PM and HM are compared to Laplace and Duchi et al.'s solutions. Optimized Unary Encoding (OUE) [Wang et al.] is also used for comparison, but for categorical data only, as this is not supported by the other methods.

### 2.3.2. CLUSTER METHODS WITH (L)DP

This chapter examines the various studies that have been conducted on clustering in combination with differential privacy. Initially, we looked at the most fundamental papers in this field. Subsequently, the focus shifted toward researching well-known papers that have been published since 2020.

The first work we highlight was proposed by Nissim et al. and aims at improving differential privacy methods, such as Laplace, which uses sensitivity to compensate the noise for a function [Nissim et al., 2007]. In addition to compensating the function, they also consider the dataset itself. The algorithm for this is called "smooth sensitivity" and is used for instance-specific noise. To apply it, the authors introduce a method/framework to effectively calculate it. To demonstrate the effectiveness of the method, they use K-means, among other cluster algorithms. Their method requires the calculation of cluster distances, using Wasserstein distance instead of Euclidean distance.

Another study focuses on both interactive and non-interactive approaches for differential privacy in K-Means [Su et al., 2015]. The study builds upon the work that was done for DPLlloyd, an interactive privacy extension of K-Means described by Blum et al. [Blum et al., 2005]. The DPLlloyd mechanism partitions an  $n$ -dimensional dataset into a grid and releases the count for each grid by adding Laplacian noise to each count. Another part of their research focuses on determining the width of the data cells. The grid estimation method used in their research is called the extended uniform grid approach (EUG), and the complete K-Means method is called EUGkM. The experiment consists of evaluating it against the DPLlloyd mechanism, which performs better in an interactive setting. Therefore, they combine their algorithm for combining both aspects into a hybrid approach (EUGkM + DPLlloyd approach) and show a better final performance.

The study of Nissim et al. researches the idea of finding the smallest possible radius in the Euclidean space  $R^d$  for a set of  $n$  points [Nissim and Stemmer, 2018]. They propose a new solution that uses locality-sensitive hashing (LSH) for differential privacy and use it to find 1-cluster in the  $d$ -dimensional Euclidean space. This method works for differential privacy (LSH-GoodCenter), but they also extend this to the local model (LDP-GoodCenter). The algorithm to find this radius is used to count the points enclosed by the radius and Laplace noise is added to the count to preserve differential privacy. The mechanism is combined and applied to work with the K-Means algorithm (LDP-K-Mean). This mechanism was extended later in a paper proposed by Kaplan et al. and introduces a similar LDP method [Kaplan and Stemmer, 2018]. They aim to reduce the number of interactions needed between the server and users to one, instead of the  $O(k \log n)$  required for Nissam

et al.'s solution [Nissim and Stemmer, 2018]. To increase the success probability, they use the same idea but extend it to have multiple centers instead of a single large one. They call it the LSH-Procedure and the algorithm Private-Centers is applied to generate centers to use with K-Means. Then, they apply the same method to the LDP method that was also originally proposed by Nissam et al. (LDP-GoodCenter). The most recent work by Stemmer et al. focuses on improving the work that was done by Kaplan et al. [Kaplan and Stemmer, 2018; Stemmer, 2021]. Because the original mechanism has a higher additive error, which means the noise that is added introduces a lot of error. To solve this, the authors aim at reducing this error by improving the original GoodCenter algorithm [Nissim and Stemmer, 2018]. Their extended method is called WeightedCenters and also adds weights to candidate centers. In the final iteration, the weights are used to create the K-Means or K-Median clusters.

Sun et al. proposed a mechanism for distributed clustering using local differential privacy (LDP) to preserve distance-based information. They claim to have the first non-interactive LDP algorithm for clustering [Su et al., 2015]. This means, they are being able to perturb the data locally at once and sent it to the server to cluster with both K-Means and DBSCAN. They encode the client-side data into an anonymous hamming space using Bit Vector (BV) and modify the encoding to preserve Euclidean distance. As their mechanism only shares distance information they were not able to use K-Means directly. To overcome this, they modified the algorithm and called it K-Cluster. Finally, the method is evaluated using Normalized Mutual Information (NMI) and Average Estimated Error (AEE).

Xia et al. noticed the shortcoming of Sun et al.'s work which is the need to share privacy-sensitive distance information [Xia et al., 2020]. Therefore, they propose a new interactive method for distributed K-means clustering using LDP. The method converts features to binary strings and uses the Random Response mechanism (RR) to perturb each feature into a feature vector. The privacy cost depends on the length of the bits of each feature transformation, meaning that a longer length yields more information at the cost of the privacy budget. In each iteration, the serverside calculates and sends K-means centroids to each user, who recalculates distances until the centroids become stable. The approach has the disadvantage of a high correlation between user data and the clusters. To solve this problem, the algorithm is improved by having the client-side send not only the user data but also a set of random zero strings. The server side then performs similar calculations to determine the true cluster. Huang et al. propose a private distributed K-means clustering algorithm for interval data that addresses a shortcoming in Xia et al.'s work by using Condensed Local Differential Privacy (CLDP) for small-scale values and LDP for large-scale values [Huang et al., 2021]. They preserve distance using a Square Wave (SW) mechanism and apply a classical K-Means algorithm on the server side to the perturbed data.

A very recent mechanism that also builds around K-Means to preserve LDP, is called the LDPK mechanism [Yuan et al., 2021]. As K-Means works only with numerical data, they use K-prototypes for supporting mixed data types. The LDPK mechanism perturbs the user data first locally and interactively exchanges information with the server to complete the clustering process. The mechanism they use for perturbation is the Harmony algorithm, which was proposed earlier by [Nguyễn et al., 2016]. To also support categorical data the S-Hist method is used, which was also introduced by Nguyen et al. But the author replaces this algorithm with OUE [Wang et al.] to improve accuracy. Due to the correlation between the cluster centroids and the real data, the server could still infer the correct information.

Therefore, the authors also disturb the user’s cluster information with an extra extension to the LDPK method, called ELDPK. To this end, they perturb the clusters with the GRR (Generalized Random Response) algorithm. Their evaluation focuses on the privacy budget and the amount of data points. They show that if the amount of data points increases the clustering quality does as well.

Most existing work focuses on (L)DP in combination with K-Means. Finally, two interesting studies focus on differential privacy for AP or DBSCAN. A study conducted by Cai et al. focuses on AP [Cai et al., 2020]. Their method involves adding Laplace noise to the responsibility matrix. For each sample data, a neighborhood is specified using a radius around the data point. This area is called the neighborhood density, and each sample point’s preference value is adjusted according to its density value. Higher density yields a higher chance of belonging to a cluster center and then being ranked based on size. The perturbed responsibility matrix and densities are combined and used to run AP. They evaluated their method using the ARI, Fowlkes-Mallows Index (FMI), and AMI.

Another recent study focuses on differential privacy for DBSCAN [Bozdemir et al.]. The proposed solution involves clustering data between two or more parties using two servers. Secure two-party computation (S2PC) is used to achieve this. Using S2PC, both servers receive a random-looking secret share. To recover the original data, both servers would need to combine their shares using S2PC, which combines the data without the servers having access to the full value. The proposed protocol is named privacy-preserving DBSCAN (ppDBSCAN). The calculations in this study are based on squared Euclidean distance (SED) and are evaluated using different methods. To evaluate the performance of ppDBSCAN, the study compares its Adjusted Rand Index (ARI) to that of K-means.

year	Name	Data type	Dataset	Code implementations	Preserving	Type	Interactive	Methods
2022 [Sun et al., 2022]	PrivBV: Distance-aware encoding for distributed...	-	Synthetic dataset	-	$(\epsilon, \delta)$ -LDP	K-Means	Non interactive	-
2021 [Huang et al., 2021]	Private distributed K-means clustering on inter...	-	-	LDP	K-Means	Interactive	-	
2021 [Stemmer, 2021]	Locally Private k-Means Clustering	n- numerical n-dimensional numeric & categorical	-	-	LDP	K-Means	Interactive	-
2021 [Yuan et al., 2021]	Privacy-preserving mechanism for mixed data clu...	-	Adult dataset, US Census dataset	-	LDP	K-Prototypes	Interactive	LDPK and ELDPK
2021 [Bozdemir et al.]	Privacy-preserving Density-based Clustering	-	Deer dataset, Lsun dataset, S1	-	DP	DBSCAN	-	ppDBSCAN
2020 [Cai et al., 2020]	DP-AP: Differential Privacy-Preserving Affinity...	-	Iris dataset, Seeds dataset	-	DP	AffinityPropagation	-	DP-AP
2020 [Xia et al., 2020]	Distributed K-Means clustering guaranteeing loc...	n-dimensional numerical data	3D Road Network, CarGPS	-	LDP	K-Means	Interactive	LDPKmeans
2019 [Sun et al., 2019]	Distributed Clustering in the Anonymized Space...	n-dimensional numerical data	Aggregation dataset, Digit dataset, Pathbased d...	-	LDP	DBSCAN, K-Means	Non interactive	Distance Aware Bit Vector (DPBV)
2018 [Nissim and Stemmer, 2018]	Clustering algorithms for the centralized and l...	n-dimensional numerical data	-	$(\epsilon, \delta)$ -LDP	K-Means	Interactive	-	LDP-GOODCenter
2018 [Nissim and Stemmer, 2018]	Differentially private K-means with constant mu...	-	-	-	K-Means	Interactive	-	LSH-Procedure & Private-Centers
2015 [Su et al., 2015]	Differentially Private k-Means Clustering	2 - 10-dimensional numerical data	Adult dataset, Gowalla dataset, Image dataset, ... <sup>2</sup>	DP	K-Means	Both	-	EUGkM and hybrid EUGkM + DPLloyd
2007 [Nissim et al., 2007]	Smooth sensitivity and sampling in private data...	n-dimensional numeric	-	-	$(\epsilon, \delta)$ -LDP	K-Means	Non interactive	Smooth sensitivity

Table 2.1: Summary table of the literature review for (L)DP clustering algorithms.

year	Name	Data type	Dataset	Code implementations	Preserving	Interactive	Methods
2019 [Wang et al., 2019]	Collecting and Analyzing Multidimensional Data ...	n-dimensional (PM), but HM is 1-dimensional and...	BR, MR	<a href="https://github.com/forestneo/sunPytools/blob/m...">https://github.com/forestneo/sunPytools/blob/m...</a>	LDP	-	- Piecewise Mechanism (PM)
2017 [Duchi et al., 2017]	Minimax optimal procedures for locally private ...	1-dimensional numerical data	-	<a href="https://github.com/forestneo/sunPytools/blob/m...">https://github.com/forestneo/sunPytools/blob/m...</a>	LDP	-	- Hybrid Mechanism (HM)
2016 [Nguyễn et al., 2016]	Collecting and analyzing data from smart device...	numerical, binary and categorical data. Domain ...	BR	-	$\epsilon$ -LDP	-	Harmony
2015 [Geng et al., 2015]	The staircase mechanism in differential privacy	n-dimensional numerical data	-	<a href="https://github.com/IBM/differential-privacy-lib...">https://github.com/IBM/differential-privacy-lib...</a>	DP	Differential privacy method	Staircase mechanism (SM)
2013 [Geng and Viswanath, 2013]	Optimal data-independent noise for differential...	n-dimensional	-	-	$\epsilon$ -DP	Non interactive	-

Table 2.2: Summary table of the literature review for (L)DP algorithms.

number	name	samples	features	target	Source	Realworld data?
1	Adult	48,842	14 numerical/categorical/boolean	income (>50k, <= 50k)	UCI	Yes
2	Seeds	210	7 numerical	type	UCI	Yes
3	Iris	150	4 numerical	class (type of iris)	UCI	Yes
4	CarGPS	17,785,500	3 geographical data	-	-	Yes
5	3D Road Network	434,874	3 geographical data	-	-	Yes
6	Pathbased	300	2 numerical	ground truth clusters	PapersWithCode	No
7	Aggregation	788	2 numerical	-	-	Unknown
8	Digit	1797	8x8 numerical	number	Scikit-learn	Yes
9	Lsun	400	2 numerical	ground truth clusters	-	No
10	S1	1500	2 numerical	ground truth clusters	-	No
11	Deer	20,033	2 numerical	-	-	Yes
13	Gowalla	6,442,890	5 (geographical data, ids and time)	-	<a href="https://snap.stanford.edu/data/loc-gowalla.html">https://snap.stanford.edu/data/loc-gowalla.html</a>	Yes

Table 2.3: The different datasets used in the related literature.

### 2.3.3. EVALUATION

For this research, it is important to compare our nd-Laplace method with that of other studies. When looking at cluster methods that use an LDP (Local Differential Privacy) mechanism, it is common to see the application of K-Means. In terms of studies that are similar to ours, they mostly involve interactive LDP methods [Huang et al., 2021; Xia et al., 2020; Yuan et al., 2021]. Among these methods, there is one that is non-interactive, but it requires adjustments to the cluster algorithm [Sun et al., 2019].

For this reason, we mainly focus on general-purpose LDP methods. This category includes mechanisms that are not dependent on a sensitivity property, as it needs to be adjusted based on the function being calibrated (e.g., average). The LDP method's from this category can be used to add noise to the input data of a cluster algorithm, similar to what we are doing in this research with the nd-Laplace mechanism.

Based on the literature study, two LDP mechanisms fall into this category: Harmony and Piecewise [Nguyễn et al., 2016; Wang et al., 2019]. The source code has been published for Piecewise, but no source code is available for Harmony. Therefore, we will only focus on comparing with the Piecewise mechanism. The following sections will be focused on investigating the Piecewise mechanism.

#### DUCHI ET AL.'S MECHANISM

As mentioned earlier, the Piecewise mechanism is based on Duchi et al.'s mechanism for one-dimensional data. Therefore, we start by explaining this mechanism first. The latter is a relatively simple method based on the Bernoulli distribution [Duchi et al., 2017]. This distribution yields either a 0 (negative) or a 1 (false) based on specified probabilities<sup>3</sup>. The mechanism works on a tuple of the domain [-1, 1] and therefore returns either -1 or 1. Hence, the probability density function (PDF) for this function is as follows<sup>4</sup>:

$$P(n) = \begin{cases} 1-p & \text{if } n = 0 \\ p & \text{if } n = 1 \end{cases} \quad (2.7)$$

The probabilities are calculated based on the input value, and can then be used to determine the estimate a mean value. In comparison to Laplace, Duchi et al.'s performs better in terms of variance for epsilons smaller than 2 [Wang et al., 2019]. However, it performs worse when the epsilon value is higher because the algorithm does not take into account the privacy budget for values that are 0. The multidimensional variant looks much like the one-dimensional variant, but samples the noise each data-point independently.

#### PIECEWISE MECHANISM

?? The authors aim to create a method that combines the advantages of the Laplace mechanism and Duchi et al.'s methods [Wang et al., 2019]. The goal is to reduce the variance for a wider range of privacy budgets. Similar to the aforementioned mechanism, the Piecewise mechanism also only accepts data within the range of -1 to 1. We first explain the one-dimensional variant, and then the multidimensional variant.

The mechanism for one-dimensional has an output of  $c \in [-C, C]$ . Which is defined based on the privacy budget:

$$C = \frac{\exp(\epsilon/2) + 1}{\exp(\epsilon/2) - 1} \quad (2.8)$$

---

<sup>3</sup><https://mathworld.wolfram.com/BernoulliDistribution.html>

<sup>4</sup><https://mathworld.wolfram.com/BernoulliDistribution.html>

The noise is sampled from this distribution, by conditionally executing one of two algorithms (see for reference the Wang et al.'s paper) [Wang et al., 2019]. Due to the symmetric nature of the probability density function (PDF) of their mechanism, the authors can handle the value of 0. Imagine it as a histogram with the distribution centered around 0. The histogram consists of three "parts" on the left, right, and center (0). This allows for the determination of the probability of 0 based on a certain likelihood, unlike in Duchi et al.'s solution.

For the multidimensional approach, the authors generate a randomly sampled set of values  $k$ . Next, for each value from  $k$ , the value is sampled using the one-dimensional Piecewise mechanism. One important factor in the overall calculation is the scale factor  $C_d$ , which, in addition to  $C$  (referring to Equation 2.8), also takes into account the number of dimensions  $d$ . Finally, the authors also provide an extension to support categorical data. This is beyond the scope of this thesis for us.

# 3

## ND-LAPLACE

In this chapter, we delve deeper into geo-indistinguishability and the various mechanisms that work with it. This is done in the order of the number of dimensions supported by the mechanism:

1. 2D-Laplace
2. 3D-Laplace
3. nD-Laplace

For each mechanism, we explain the equation for GI, the mechanism, and the truncation of data.

### 3.1. 2D-LAPLACE

The idea of GI was introduced to solve the issue of privacy and location data [Andrés et al., 2012]. To recall what we discussed earlier (Equation 2.4) It offers an alternative approach for achieving differential privacy for geographical data (latitude/longitude). The mechanism achieves this by adding noise to the location locally before sending it to a location-based system (LBS). This section starts with an introduction to mathematics, and for each of the different subsections, we visualize and explain open challenges and theoretic for applying them for clustering.

#### 3.1.1. PLANAR AND POLAR LAPLACE

The idea of planar Laplace is to generate an area around  $x_0 \in X$  according to the multivariate Laplace distribution. The mechanism of planar Laplace is a modification of the Laplace algorithm to support distance [Andrés et al., 2012]. This distance method  $dist(x, x')$  is defined as the Euclidean distance between two points or sets. Recalling the definition of Laplace, this method  $|x - x'|$  is replaced by the distance metric. Hence, the definition of the Probability Density Function (pdf) by Andrés et al. is:

$$\frac{\epsilon^2}{2 * \pi} e(-\epsilon d(x_0, x)) \quad (3.1)$$

Which is the likelihood a generated point  $z \in Z$  is close to  $x_0$ . The method works for Cartesian coordinates but was modified to support polar coordinates by including  $\theta$ . So each

point is reflected as  $(r, \theta)$  and can be modified by using a slight modification to work for polar Laplace. A point  $z \in Z$  where  $z = (r, \theta)$  is randomly generated using two separate methods for calculating  $r$  and  $\theta$ .

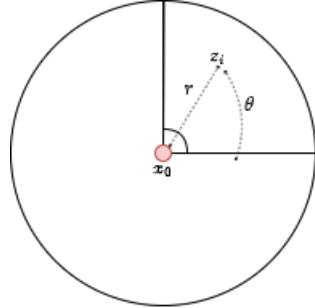


Figure 3.1: Representation of the generated  $z = r\theta$  and original point  $x_0$ .

**Calculating  $r$ :** This variable is described as  $dist(x_0, z)$  and can be randomly drawn by inverting the CDF for the Laplace distribution:

$$C_e^{-1}(p) = -\frac{1}{e} (W_{-1}(\frac{p-1}{e}) + 1) \quad (3.2)$$

For this equation,  $W_{-1}$  is a Lambert W function with a -1 branch. The Lambert w function also called the product logarithm, is defined as  $W(x)e^{W(x)} = x$  [Lehtonen, 2016]. The purpose of the Lambert w function is to invert the CDF of the Laplace distribution to generate random noise for one of the coordinates ( $r$ ) using the random value of  $p$ .

**Calculating  $\theta$ :** The other coordinate ( $\theta$ ) is defined as a random number  $[0, 2\pi]$ . To visualize these methods, it is necessary to convert the polar coordinates for  $z = (r, \theta)$  back to a plane  $(x, y)$ . This is described as step 4 of the planar Laplace algorithm [Andrés et al., 2012] and visualized using figure 3.2.

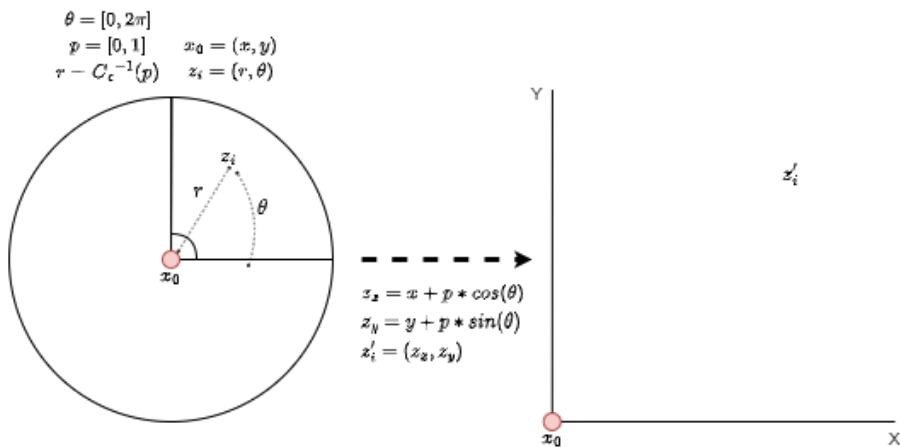


Figure 3.2: Representation of converting the perturbed point  $z = (r, \theta)$  to a point  $z_x, z_y$

### 3.1.2. TRUNCATION

The truncation is an essential part of the mechanism to ensure the data is contained within the domain of the original data  $X$ . If this is not the case, the data is easily distinguished by an unwanted adversary [Andrés et al., 2012; Min et al., 2022]. We assume a user has a set of data points with a range of  $[-1, 1]$ . If noise is added, it cannot be ensured that this data falls outside this range (figure 3.3).

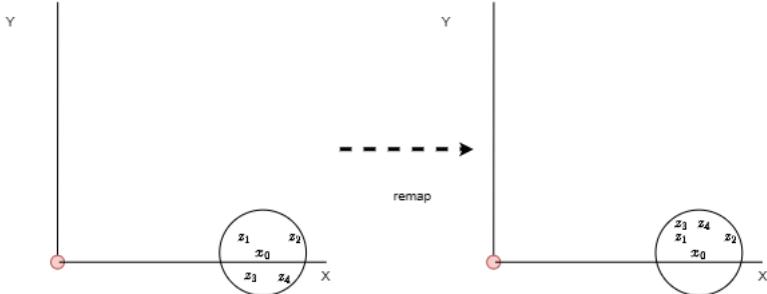


Figure 3.3: Representation of truncation of data points for 2-dimensional Laplace mechanism.

A solution was described by Andres et al. in step 5 of the Laplacian mechanism for 2D space [Andrés et al., 2012]. The idea is to create a grid around the diameter of the set of points  $X = R^2$  that belong to the user. This grid is defined as  $G = \{g_1, g_2, \dots, g_n\}$  where  $g_i$  is a grid point. So, a point generated outside the given domain is remapped to the closest in  $X \subset G$ .

This was later improved by Chatzikokolakis et al., introducing an optimized way of remapping [Chatzikokolakis et al., 2017]. The algorithm uses the Bayesian rule to minimize the loss of utility while remapping the data. Instead of remapping to the closest point, it remaps to a location where the loss is minimal. To decrease the performance impact of this algorithm, it is possible to only consider a certain region around the perturbed point  $z$ . The disadvantage of this method is the need for a prior set of data points to calculate the optimal remapping. It does not work for new users and extends the training period.

### 3.1.3. OPTIMIZING FOR CLUSTERING

The decision of the parameters for the algorithm is straightforward as it depends on the  $\epsilon$ . This constant is calculated by defining the radius  $r$ , and the desired level of privacy  $l$  and  $\epsilon$  is calculated using  $l/r$ . The  $l$  is a predefined constant  $l \in R^+$  but usually will be below 10. For geographical data, the  $r$  can be configured using meters as a unit of measure. So, for example,  $r = 200$  corresponds to a radius of 200m around point  $x_0$ . Without a unit, it is a challenge to define a reasonable radius.

In that regard, the radius can also be a flexible value that is defined based on the crowdedness of a region [Chatzikokolakis et al., 2015]. If a user is located in a crowded area, the radius can be smaller than if the user is located in a rural area (because the user's location is indistinguishable due to the overlap of other users' locations). Instead of providing GI, the authors introduce a more flexible privacy definition  $d_x$ -privacy. The  $r$  is calculated based on the mass of other locations in the region  $r$ , which they call *privacy mass*. So, the total mass of a set  $A$  is defined as  $M(A) = \sum_{x \in A} m(x) = a + q(x)b$ . Where  $m(x)$  is the mass of a location  $x$  and is of value  $[0, 1]$ . For this formula,  $a$  is the number of points assigned to each location. The authors define  $q(x)$  as the "quality" of a point, which is essentially the

number of other users that are also interested in the same point (e.g. a mall).

The  $a$  is defined as a Euclidean ball  $B_r = x' | d_{euc}(x, x') \leq r$ , which returns all locations within the radius  $r$ . To retrieve a value within  $[0, 1]$ , the authors use the following formula:

$$a = \frac{1}{|B_r|} \quad (3.3)$$

If the locations are only considered for space and not quality,  $q(x)$  is defined as 0 [Chatzikokolakis et al., 2015].

Although the method is an interesting approach to increasing utility while not reducing privacy, it is hard to adopt for the purpose of clustering. For applying it for LDP, it is required to supply each user with prior knowledge of the dataset. This would require an interactive setup of the mechanism instead of a non-interactive one.

A drawn area, as shown in 3.1, can be expressed as a perturbation area  $P_{area}$  [Yan et al., 2022]. This metric was formulated as follows:

$$P_{area} = \left\{ center = x_0, radius = \frac{1}{N} \times \sum_{i=1}^N r_i \right\} \quad (3.4)$$

The method loops through each perturbed point  $r$  on center  $x_0$  (recall 3.1) and calculates the Euclidean distance for an  $n$  amount of perturbation points. Although the method does not contribute to the Laplace algorithm, it is useful for visualization purposes. This method can also be applied for efficiently calculating the grid points for the truncation method (recall 3.1.2).

### 3.1.4. FINAL MECHANISM

Finally, we provide as means of a summary the final algorithm for the Laplace mechanism for 2D space

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**Algorithm 1** Full algorithm for perturbing training data for 2D-clustering using planar/2D-Laplace [Andrés et al., 2012]

---

```

Input:  $x \in X$                                      ▷ 2D array of points
Input:  $l \in R^+$ 
Input:  $r \in R^+$                                 ▷ sensitivity
Output:  $z \in Z$                                  ▷ 2D array of perturbed points
 $\epsilon = \frac{l}{r}$                                ▷ Calculating privacy budget [Andrés et al., 2012]
 $x_{min} \leftarrow min(X)$ 
 $x_{max} \leftarrow max(X)$ 
 $Z \leftarrow []$ 
for  $point_i \in X$  do
     $\theta \leftarrow [0, \pi/2]$                          ▷ Random noise for  $\theta$ 
     $p \leftarrow [0, 1]$ 
     $z_i \leftarrow C_\epsilon^{-1}(p)$                   ▷ formula 3.2
     $x_{perturbed} \leftarrow point_{i_x} + (z_{i_x} * \cos(\theta))$  ▷ add noise to x-coordinate
     $y_{perturbed} \leftarrow point_{i_y} + (z_{i_y} * \sin(\theta))$  ▷ add noise to y-coordinate
    append  $x_{perturbed}, y_{perturbed}$  to Z
end for
return Z

```

---

## 3.2. 3D-LAPLACE

The previous sub-section focused on describing the use of 2-dimensional noise on geographical data. This approach has recently been extended to support 3-dimensional data, which benefits indoor navigation [Min et al., 2022]. The method is similar to the 2D approach but includes the azimuth angle  $\psi$ , in addition to the polar angle  $\theta$  and radius  $r$ .

### 3.2.1. GEO-INDISTINGUISHABILITY

To establish the same privacy guarantees for 3-dimensional data as for 2-dimensional data, the original equation 2.4 is extended [Min et al., 2022].

$$K(x_1)(z) \leq e^{\epsilon * d_3(x_1, x_2)} K(x_2)(z) \quad (3.5)$$

Where  $x_1$  and  $x_2$  are two real data points in the same dataset  $X$ .

### 3.2.2. SPHERICAL LAPLACE

The implementation of Min et al. projects the dimensions onto a sphere instead of a circle [Min et al., 2022]. This sphere is a unit sphere calculated with a radius of 1. Based on this sphere the polar angle  $\theta$  and azimuth angle  $\psi$  are randomly calculated.

**Calculating  $\theta$  and  $\psi$ :** Both are drawn from the unit sphere using the following equations:

$$\theta = \frac{1}{\pi} \quad (3.6)$$

$$\psi = \frac{1}{2\pi} \quad (3.7)$$

The tuple  $U = (\theta, \psi)$  is randomly selected based on the uniform distribution of the unit sphere [Min et al., 2022].

**Calculating  $r$ :** The radius  $r$  (distance from the center) is calculated using the following equation:

$$r = \frac{1}{2} \epsilon^3 * r^2 * e^{-\epsilon * r} \quad (3.8)$$

Where the gamma scale is the same as for 2D-Laplace, but with a shape of 3 instead of 2. The noise is added to the original location  $x$  to obtain the perturbed location  $z = x + U * r$ . A clear example of the noise that is generated by this method is shown in figure 3.4. Finally, we convert this to the Cartesian coordinate system to obtain the final location  $z$ :

$$\begin{aligned} z_x &= r * \sin(\theta) * \sin(\psi) \\ z_y &= r * \sin(\theta) * \cos(\psi) \\ z_z &= r * \cos(\theta) \end{aligned}$$

This is also visualized in figure 3.5.

Kernel density plot of a point  $x_0$  (centre) with 50 random generated points according to the formula for 3D-Laplace in a projected sphere with epsilon 3

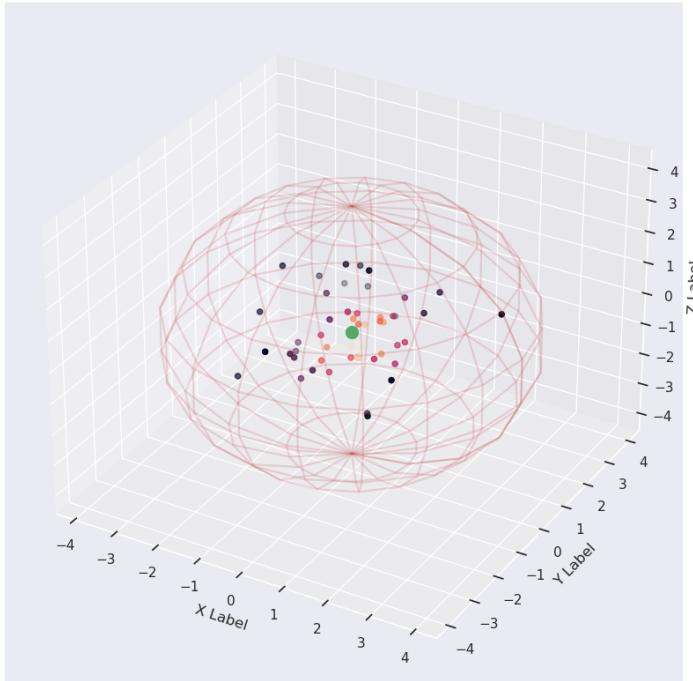


Figure 3.4: 50 random noise samples generated around point  $x_0$  (green dot) using the 3D-Laplace noise method [Min et al., 2022] plotted on a sphere.

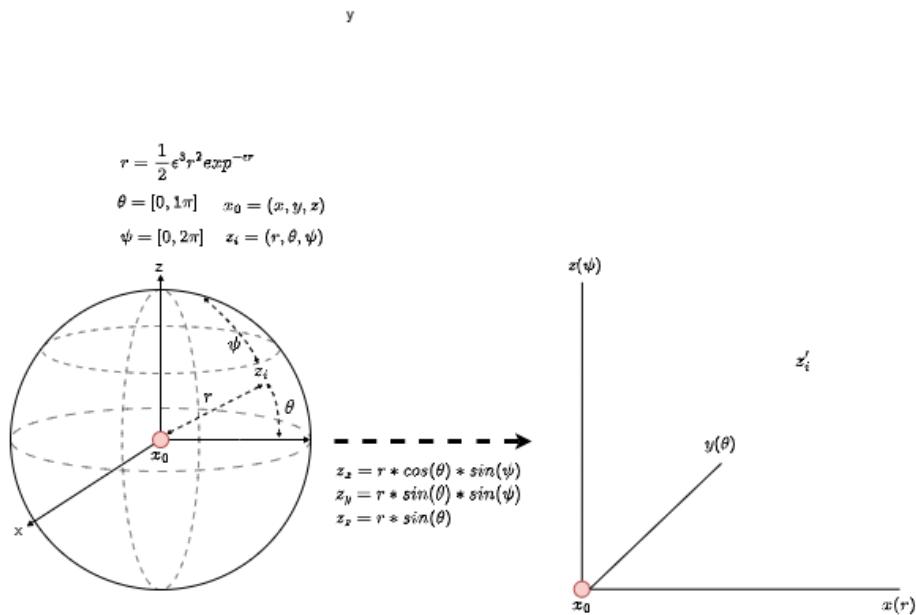


Figure 3.5: 3D-Laplace noise distribution according to the method proposed by Min et al. [Min et al., 2022]

### 3.2.3. TRUNCATION

As with the 2D-Laplace method, the 3D-Laplace method also has a truncation method. This truncation method is also based on the same method as the 2D-Laplace method. Instead of a plane grid, a cuboid grid is used for 3-dimensional space. This also remaps the noise to the closest grid point or existing point. To demonstrate this, we plotted example data points on a 3-dimensional grid in figure 3.6.

Example of generating noise for a dataset  
and remapping it to  $X \subset G$  when outside the domain

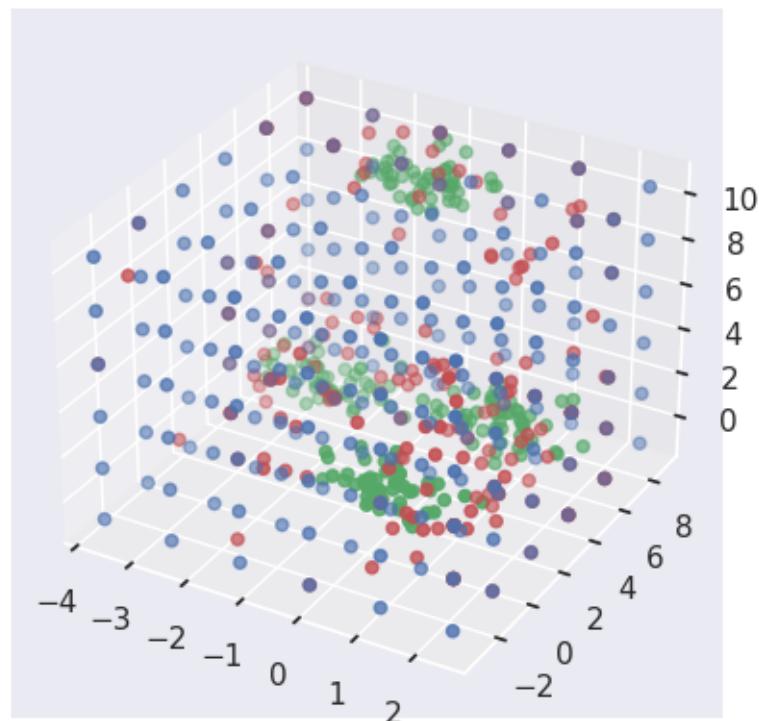


Figure 3.6: Applying 3-dimensional noise with  $\epsilon = 1$  (red dots) to a dataset  $X$  (green dots). Demonstrating remapping to the closest grid point (blue) or  $X$ .

### 3.2.4. FINAL MECHANISM

Finally, we provide as means of a summary the final algorithm for the Laplace mechanism for 3D space

---

**Algorithm 2** Full algorithm for perturbing training data for 3D-clustering using planar/2D-Laplace [Andrés et al., 2012]

---

```

Input:  $x \in X$                                      ▷ 3D array of points
Input:  $l \in R^+$ 
Input:  $r \in R^+$ 
Output:  $z \in Z$                                      ▷ 3D array of perturbed points
 $\epsilon \leftarrow \frac{l}{r}$                            ▷ Calculating privacy budget [Andrés et al., 2012]
 $Z \leftarrow []$ 
for  $point_i \in X$  do
     $\theta \leftarrow 1, \pi 2$                          ▷ Random noise according to equation 3.7
     $p \leftarrow \frac{1}{2\pi}$                         ▷ Random noise according to equation 3.8
     $r \leftarrow \frac{1}{2}\epsilon^3 * r^2 * e^{-\epsilon * r}$  ▷ Draw  $r$  based on equation 3.8
     $z_x \leftarrow r * \sin(\theta) * \sin(p)$ 
     $z_y \leftarrow r * \sin(\theta) * \cos(p)$ 
     $z_z \leftarrow r * \cos(\theta)$ 
     $Z ::= Add(z)$                                 ▷ Adds z to the list Z.
end for
return Z

```

---

### 3.3. ND-LAPLACE

As mentioned in the previous chapter, the paper that was introduced by Min et al. is be-able to handle 3-dimensional data. A small recap: a point  $(r, \theta, \psi)$  gives us the spherical coordinates of a given 3-dimensional sphere. An important property for this is the fact that each of these coordinates can be generated separately [Andrés et al., 2012; Min et al., 2022]. The  $r$  gives us the radius or distance from  $(\theta, \psi)$  to the center of the sphere<sup>1</sup>. So, instead of having just these two coordinates, we are be-able to extend this to n-dimensions by considering an n-hypersphere [Fernandes et al., 2019; Min et al., 2022]. To this end, besides points  $\theta$  and  $\psi$  we also consider  $\theta \in S^n$ , where  $S$  is a unit hypersphere.

The first step to generate the noise is first to select the  $r$ . This method is almost identical to the one for 3-dimensional (3.8). But, instead of applying a scale of 3, the scale will be  $n$  for the number of dimensions in the data [Fernandes et al., 2019]:

$$\gamma(n, 1/\epsilon) \quad (3.9)$$

For the other dimensions, we consider a vector  $U = (\theta_1, \theta_2, \theta_n)$  which is uniformly selected based on a unit  $n$ -hypersphere  $S^n$  [Fernandes et al., 2019]. We consider the work that was proposed by Marsaglia et al. for 4-sphere that can be used for selecting points from an n-hypersphere [Marsaglia, 1972]. This method resolves around selecting points from a hypersphere by using a uniform distribution for the domain  $[0, 1]$ . We adopted the approach that uses the Gaussian distribution<sup>2</sup>.

#### 3.3.1. CARTESIAN COORDINATES

As with the 2/3D-Laplace, the spherical coordinates need to be converted to Cartesian to be able to cluster. It is comparable to the way it was done in the previous chapters, however, as there are an  $n$ -amount of angles the equation is repeated and slightly different:

$$x_1 = r * \cos(\theta_1) \quad (3.10)$$

$$x_2 = r * \sin(\theta_1) * \cos(\theta_2) \quad (3.11)$$

$$x_n = r * \sin(\theta_1) \dots \sin(\theta_{n-2}) * \cos(\theta_{n-1}) \quad (3.12)$$

$$x_n = r * \sin(\theta_{n-1}) * \sin(\theta_{n-2}) * \sin(\theta_{n-1}) \quad (3.13)$$

If we combine sections 1 and 2 of this chapter, we are being able to give a good overview of the solution using a similar image as for the 2D and 3D variants (figure 3.7).

---

<sup>1</sup><https://mathworld.wolfram.com/SphericalCoordinates.html>

<sup>2</sup><https://mathworld.wolfram.com/SpherePointPicking.html>

$$r = \gamma(4, 1/\epsilon)$$

$$U = \frac{1}{1\sqrt{2\pi}} \exp^{-\frac{1}{2}(\frac{x-0}{1})^2}$$

$$\theta_n \in U$$

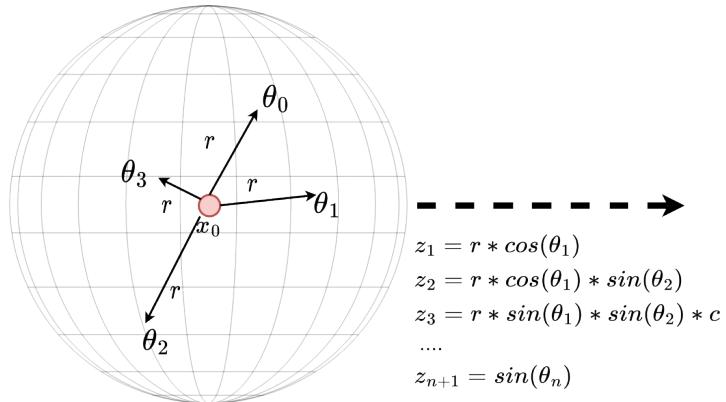


Figure 3.7: Overview of the nD-Laplace mechanism

### 3.3.2. PRIVACY VERSUS UTILITY

If we continue adding dimensions, we notice the noise is shrinking proportionally. To understand this behavior, we first have to examine the formula for a hypersphere's volume.

$$S_n = \frac{2\pi^{n/2}}{\gamma(\frac{1}{2}n)} \quad (3.14)$$

Where  $\gamma$  is the gamma distribution that is determined based on the number of dimensions <sup>3</sup><sup>3</sup>. As the amount of dimensions increases, the most volume is located on the hypersphere surface. When we convert the points to Cartesian coordinates, some will be located at the center (e.g., 0.5), while others will be close to the surface (e.g., 0.0). However, as the number of dimensions increases, the majority will be close to the surface (e.g., 0.99). The decreasing amount of volume is illustrated using this figure:

---

<sup>3</sup><https://mathworld.wolfram.com/Hypersphere.html>

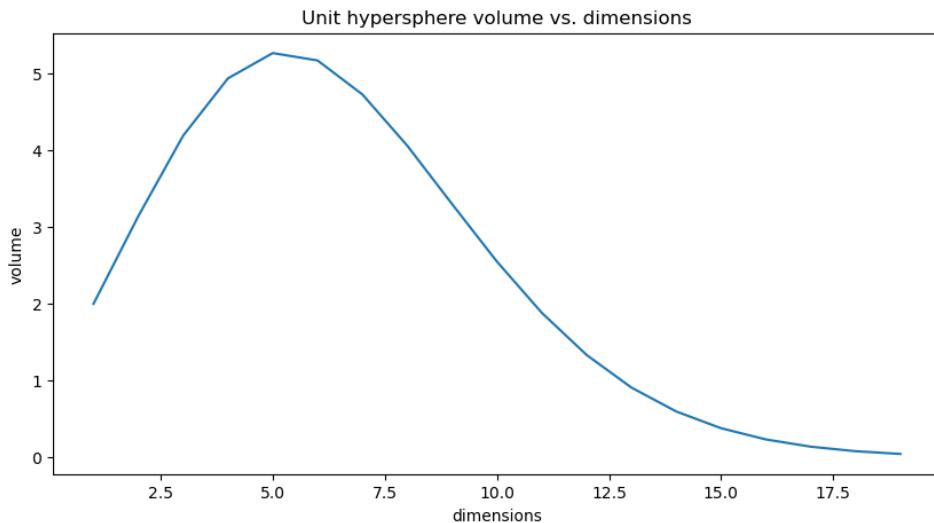


Figure 3.8: Illustration of the decreasing volume while increasing the number of dimensions

The noise decreases as the dimensions increase, increasing utility. Hence, it is intriguing to observe the behavior of privacy relative to utility. This behavior will be further emphasized in a later stage of this research.

### 3.3.3. TRUNCATION

For 2D/3D Laplace, a grid and cuboid were respectively introduced to truncate noise mechanisms [Andrés et al., 2012; Min et al., 2022]. This section extends the work done for 2D and 3D Laplace (3.1.2, 3.1.3)

This section introduces an extension for handling any number of dimensions, which can also substitute for 2D/3D. We want to note that this section focuses on improving utility with remapping. Any remap function preserves geo-indistinguishability Chatzikokolakis et al. [2017], so the required privacy is still preserved.

Recalling both mechanisms, the 2D version operates on a plane and approximates on a grid  $G$ , while the 3D version works in a 3D space using a cuboid grid. Given a set of input points  $X \subset R^2$ , we can truncate points that are outside the domain by remapping them to points within  $G$  ( $Z = X \cap G$ ) [Andrés et al., 2012]. Here,  $X$  represents other data points reported locally by the same user. To extend this approach to n-dimensional data, we need an efficient way to search points in an n-dimensional hypersphere. To do this, we adopt the idea proposed by Chatzikokolakis et al. of using a kd-tree for efficient searching of the grid [Chatzikokolakis et al., 2017]. In their research, they describe the utilization of a kd-tree for searching nearby points for a given point. For this reason, we also use a kd-tree for the following tasks:

1. Finding nearby points for  $z \in G$  (section: [Grid with kd-tree remapping](#)).
2. Finding nearby points for  $x \in X$  and  $z \in Z$  (section [Optimal remapping](#)).

For visualization purposes, this section will primarily focus on 2D data. However, it is important to emphasize that the same algorithm will also be applied to 3D and nD data. The underlying principles and steps of the algorithm remain the same.

We first give an introduction to kd-trees on the next page and then explain how we apply them for the two tasks.

## KD-TREES

A kd-tree is an algorithm that can be used to search a grid for nearby points [Bentley, 1975]. It is capable of doing so, by recursively splitting the grid into a binary tree to search for grid coordinates [Washington, 2]. In addition to this, it preserves spatial information of the data so it can be utilised to find nearby points using Euclidean distance (nearest neighbor search). The following example provides an idea on how this works (Figure 3.9):

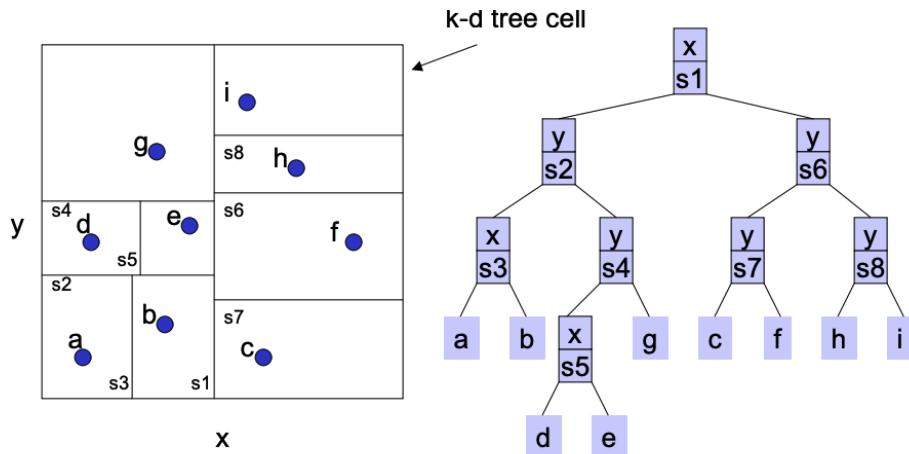


Figure 3.9: Representation of constructing a kd-tree with 2 dimensions [Washington, 2].

Take, for example, the 2D Laplace algorithm that utilizes a plane (left side). The data points can be divided based on their x and y coordinates. Each coordinate becomes a node in the binary tree, and the grid is divided based on these splits. The binary tree allows us to efficiently search the grid. An example of this is provided in the following image (Figure 3.10):

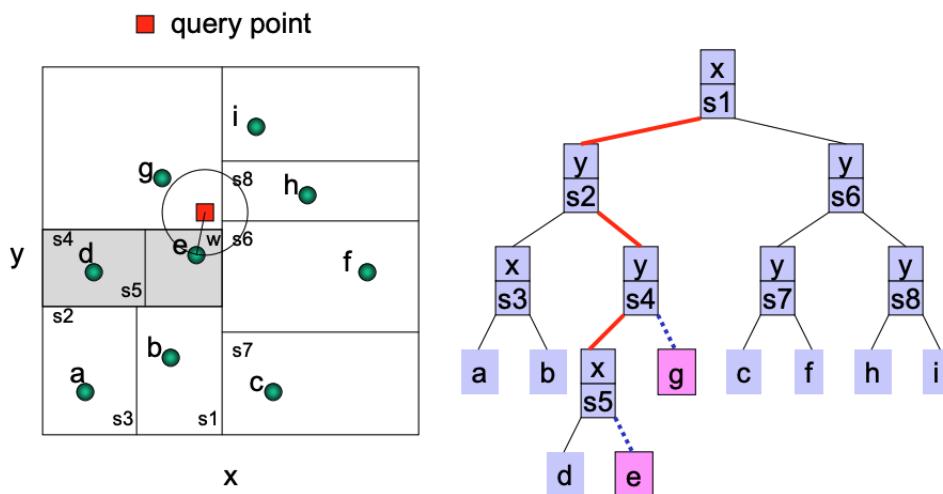


Figure 3.10: Representation of searching a kd-tree with 2 dimensions [Washington, 2].

In the example, we are searching for all points that fall within the radius of a random query point. Thanks to the grid being divided into a binary tree, a portion of the grid can

be efficiently searched, evaluated and referenced. The biggest advantage is that this greatly reduces the complexity of searching. Constructing the kd-tree only costs  $O(kn)$ , where  $k$  is the number of dimensions and  $n$  is the dataset size. Searching for a nearest neighbor is a little less efficient, with a time complexity of  $O(\log n)$  [Washington, 2]. A reference to the Big O notation can be found in the attachment.

### GRID WITH KD-TREE REMAPPING

As explained in the previous paragraph, a kd-tree can be used to perform a nearest neighbor search. This is highly relevant to our research as it proves to be beneficial for the optimizations we are striving for. To this end, we adopt this approach for remapping the perturbed datapoints  $z \in Z$  to a grid  $G$ .

We have illustrated the three steps required for this below:

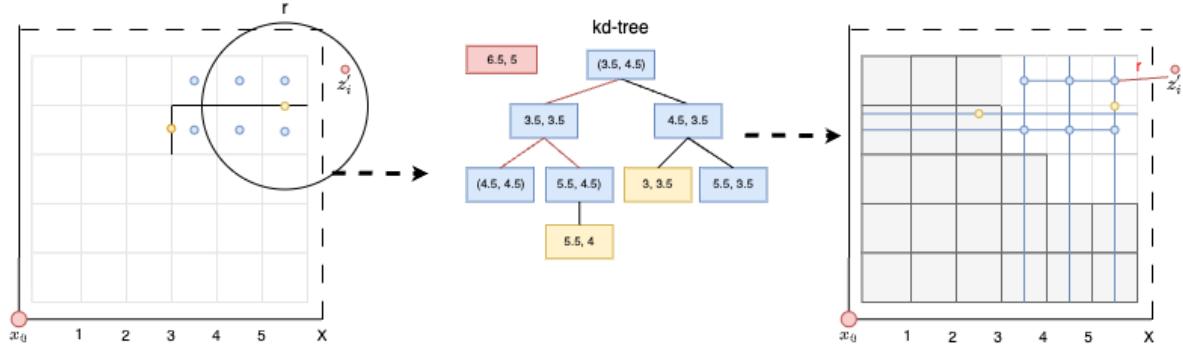


Figure 3.11: Representation of a kd-tree with 2 dimensions to remap based on a grid.

The above illustration presents the grid-remapping algorithm. Firstly, a grid is generated, where each (blue) point represents the center of a grid cell. Together, these centroids form the grid dataset, denoted as  $G$ . The yellow points and  $x_i$  are part of the original collection, denoted as  $X$ . Here,  $r$  represents the radius used to generate a private version of the data point  $x_i$ , named  $z_i$ , based on 2D-Laplace (in this case). In the illustration, you can observe that  $z_i$  falls outside the original domain of  $X$ , and for this reason, it needs to be remapped. We accomplish this by utilizing the nearest-neighbor search from the kd-tree algorithm, allowing us to search in  $X \cup G$ . Using this algorithm, we can effectively remap point  $z \in Z$  to either  $X$  or  $G$  based on the closest Euclidean distance (Algorithm 4 and Algorithm 3).

The utility of this method depends on the number of grid cells in  $G$ , since a smaller distance will result in more frequent mapping to the surface of the grid. When  $\epsilon$  is very low (and thus farther away), the data points are more likely to map to the grid surface (3.4, 3.6). Increasing the number of grid cells can improve the utility, but this comes at the cost of significantly increased space complexity for  $k$  dimensions. This is because a grid of  $n * m$  dimensions has a complexity of  $O(n^2)$ . Therefore, we also explore the optimal remapping algorithm proposed by Chatzikokolakis et al [Chatzikokolakis et al., 2017].

Clarify algorithms / rewrite algorithms

---

**Algorithm 3** Algorithm for finding points outside the domain of  $X$ .

---

**Input:**  $x \in X$  ▷ original dataset  
**Input:**  $z \in Z$  ▷ perturbed dataset

$tree \leftarrow KDTree(X)$  ▷ construct a KDTree from the original data.  
 $X_{domain} \leftarrow \text{KDTree::QUERY}(Z)$  ▷ find the closest points.  
 $X_{features} \leftarrow X.\text{features}$  ▷ retrieve dataset dimensions  
 $X_{outside-domain} \leftarrow []$

**for**  $feature \in X_{features}$  **do**

**for**  $index \in X[feature]$  **do**

**if**  $Z[index][feature] \leq X.\text{MIN}(Z)$  **then**

append  $row$  to  $X_{outside-domain}$

**end if**

**if**  $Z[index][feature] \geq X.\text{MAX}(Z)$  **then**

append  $row$  to  $X_{outside-domain}$

**end if**

**end for**

**end for**

**return**  $X_{outside-domain}$  ▷ The index of points outside the domain of  $X$ .

---

---

**Algorithm 4** Algorithm for generating and remapping to a grid.

---

**Input:**  $x \in X$  ▷ original dataset  
**Input:**  $z \in Z$  ▷ perturbed dataset  
**Input:**  $grid$  ▷ grid structure ( $n * m$ )

$d_X = dist(Z, X)$  ▷ euclidean distances  
 $d_{grid} = dist(Z, grid)$

$Z_{out-domain} \leftarrow FindPointsOutsideDomainX(X, Z)$  ▷ Algorithm 3

$grid\_tree \leftarrow KDTree(grid)$

$grid\_mask \leftarrow KDTree::query(Z)$  ▷ find indices of  $z \in Z$  that are closeby grid cells.

$Z_{grid-mask} \leftarrow Z_{out-domain} \cup d_{grid} < d_X$  ▷ All points  $z \in Z$  that are closeby grid cells and are outside domain.

$Z' \leftarrow Z[grid[grid\_mask][Z_{grid-mask}]]$  ▷ combinate masks to set appropriate indexes to  $g \in grid$ .

**return**  $Z'$

---

### OPTIMAL REMAPPING

As we discussed, the remapping will be performance intensive to be able to provide good utility and that is why we adopt the optimal remapping [Chatzikokolakis et al., 2017]. Consider the grid that was proposed in 3.11. After remapping point  $z_i$ , it is mapped to the center for the grid cell. Based on the cell width, the distance to the original point  $x_i$  and  $z_i$  could be really large. Our goal is to remap the center of the grid cell (now  $z_i$ ) to a point that is closer to  $x_i$  (if applicable).

This is visualized by zooming into the last step of the grid-remapping (Figure 3.11).

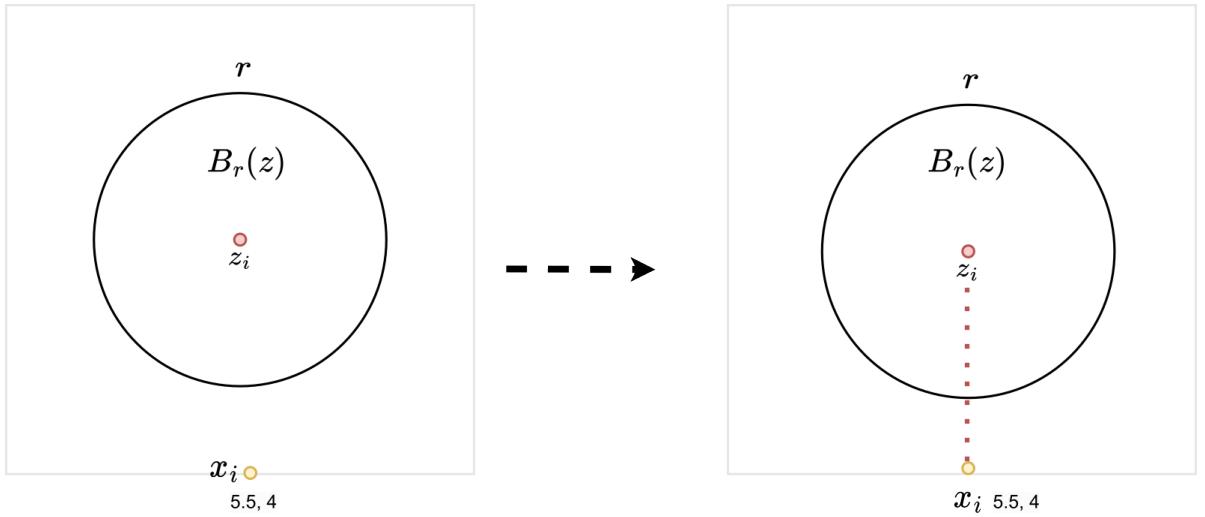


Figure 3.12: Representation of optimal remapping [Chatzikokolakis et al., 2017], where  $z_i$  is remapped to  $x_i$  using  $\sigma(x)$  instead of the center of the grid cell.

In Figure 3.12 we can observe that  $z_i$  is remapped to the center of the grid cell as a consequence of the grid-remapping. For the reasons mentioned earlier, this is not optimal, and we can further optimize it by utilizing the other data points. The remapping algorithm works on the idea of crowded places 3.1.3, with the intuition that a crowded place leverages indistinguishability by crowdedness [Chatzikokolakis et al., 2017]. For the remainder of this thesis, we will use the term "density" instead of "crowdedness" because it better aligns with the clustering of data.

The first step is to calculate  $B_r(z)$ , which refers to all the data points that fall within the original radius  $r$  around the data point  $z_i$ . The next step in the algorithm is to collect the data points around  $x_i$ , in order to calculate how closely it can be remapped to  $z_i$  while preserving distinguishability. This is the collection (convex hull) of all original data points  $x \in X$  that are close to  $x_i$ , determined based on the radius  $r$  around  $x_i$ . Finally, we combine both sets to obtain  $Q_r = B_r \cap X$ . Now that we have the sets of points around  $x_i$  and  $z_i$ , we can calculate the density for each point  $q \in Q_r$  [Chatzikokolakis et al., 2017]:

$$\forall x \in Q_r \quad \sigma(x) = \frac{w(x)e^{-\epsilon d(x,z)}}{\sum_{q \in Q_r} w(q)e^{-\epsilon d(q,z)}} \quad (3.15)$$

Where  $w(q)$  is the weight of a point  $q$  in  $Q_r$ , and can be seen as points that are visited earlier by the user or other users (e.g. point of interests). We will revisit this topic in the next paragraph when discussing the practical implementation of the nd-Laplace algorithm. The

same applies to  $w(x)$ , but for an individual point  $q \in Q_r$  instead of the summation. The outcome of the formula is a collection of values that indicate the degree of density. Which we call  $\sigma \in S$ . With this data, we are be-able to calculate a new  $z'$  that is closer to  $x_i$  to minimize the expected loss of utility [Chatzikokolakis et al., 2017].

The collection  $S$  can be seen as the coefficient for each point  $x \in X$ , that can be used as a scale to apply for each point  $X$  [Chatzikokolakis et al., 2017]:

$$\bar{\sigma} = \sum_{\sigma \in S} \sigma(x) * x \quad (3.16)$$

Then next, the probabilities are calculated:

$$W = \forall \sigma \in S = \frac{\sigma(x)}{\bar{\sigma}} \quad (3.17)$$

Finally, the  $z'$  is calculated by calculating the average with the probabilities as weight:

$$z' = \frac{\sum_{weight \in W} x * weight}{\sum_{weight \in W} weight} \quad (3.18)$$

## PRACTICAL IMPLEMENTATION

It is difficult to interpret  $w(q) \in Q_r$  beforehand based on other users, as we do not have this information (there is a way, but we explain this at the end of this section). To this end, we will interpret  $w(x)$  as the number of points within the radius  $r$  around a point  $x \in Q_r$ . Afterwards, it is possible to divide the outcome of this value by the sum of these points (as done in Algorithm 3.15). We therefore remain to the same algorithm as proposed by Chatzikokolakis et al. but interpret the weight differently.

It is however still possible to interpret  $w(q)$  as the weight based on other user's data-points. This requires us to implement the mechanism interactively. In this approach, all clients perturb their data and sent it to the server. The server clusters the private data and calculates weight based on cluster information (e.g., crowdedness/density) and shares it with the clients. The clients then use the optimal remap and share their private information with the server again. Although this system requires only a single round-trip between server and clients, it reveals cluster information, so we prefer the non-interactive setup.

---

**Algorithm 5** Algorithm to implement the optimal remapping of  $z \in Z$  to be in the domain of  $x \in X$

---

**Input:**  $x \in X$  ▷ n-dimensional array of original points  
**Input:**  $z \in Z$  ▷ n-dimensional array of grid-remapped perturbed points  
**Input:**  $\epsilon_{\text{epsilon}}$  ▷ privacy budget  
**Output:**  $z' \in Z$  ▷ n-dimensional array of optimal-remapped perturbed points

$$Z' = \text{FindRemappedPoints}(Z)$$

$$\text{tree} \leftarrow \text{KDTree}(X) \quad \text{▷ construct a KDTree from the original data.}$$

**for**  $z' \in Z'$  **do**

$$r = \text{FINDRADIUS}((z')) \quad \text{▷ Get original radius } r.$$

$$X_r \leftarrow \text{KDTREE::QUERY}(x) \quad \text{▷ find } q \in X \text{ around } x \text{ with radius } r.$$

$$B_r \leftarrow \text{KDTREE::QUERY}(z')$$

$$\sigma(x) = []$$

$$Q_r = X_r \cap B_r$$

**for**  $q \in Q_r$  **do**

$$q \leftarrow \text{KDTREE::QUERY}(q)$$

$$w_x = \text{LENGTH}(X_r, B_r) \quad \text{▷ weight is simply adding density of } X_r \text{ and } B_r.$$

$$w_q = \text{LENGTH}(Q_r) \quad \text{▷ } w_q \text{ is the density of each point } q \text{ within } Q_r.$$

$$\sigma(w_x) \leftarrow \text{REMAP}(w_x, \epsilon_{\text{epsilon}}) \quad \text{▷ Use equation 3.15 to remap } w_x.$$

$$\sigma(x) \leftarrow \text{APPEND}(\sigma(x), \frac{\sigma(w_x)}{w_q}) \quad \text{▷ add to the list } \sigma(x).$$

**end for**

$$z' \leftarrow \text{AVERAGE}(\sigma(x), P) \quad \text{▷ Calculate } z' \text{ using the equations: 3.16, 3.17 and 3.18.}$$

**end for**

---

### 3.3.4. PUTTING IT TOGETHER

---

**Algorithm 6** Full algorithm for perturbing training data for nD-clustering using planar/2D-Laplace [Andrés et al., 2012]

---

**Input:**  $x \in X$  ▷ n-dimensional array of original points  
**Input:**  $\epsilon$  ▷ privacy budget  
**Output:**  $z \in Z$  ▷ n-dimensional array of optimal-remapped perturbed points

$$\text{sphere} = \text{GENERATEUNITSPHERE}(x) \quad \text{▷ construct a sphere around } x.$$

**for**  $row \in X$  **do**

$$d \leftarrow \text{LENGTH}(row) \quad \text{▷ amount of dimensions}$$

$$r \leftarrow \text{GENERATERADIUS}(d) \quad \text{▷ generate radius } r \text{ using Equation 3.9.}$$

$$\text{sphere} \leftarrow \text{GENERATETHETA}(d) \quad \text{▷ perturb the sphere Figure 3.7.}$$

$$\text{noise} \leftarrow \text{CARTESIAN}(row, \epsilon_{\text{epsilon}}) \quad \text{▷ use Equation 3.13 to perturb } row \text{ and convert to cartesian.}$$

$$z = x + \text{noise}$$

$$\text{APPEND}(Z, z)$$

**end for**

**return**  $Z$

---

### 3.3.5. MECHANISM FLOWCHART

All formulas and theories are established for 2D, 3D and nD-Laplace, so the mechanism design is applicable to all three variants:

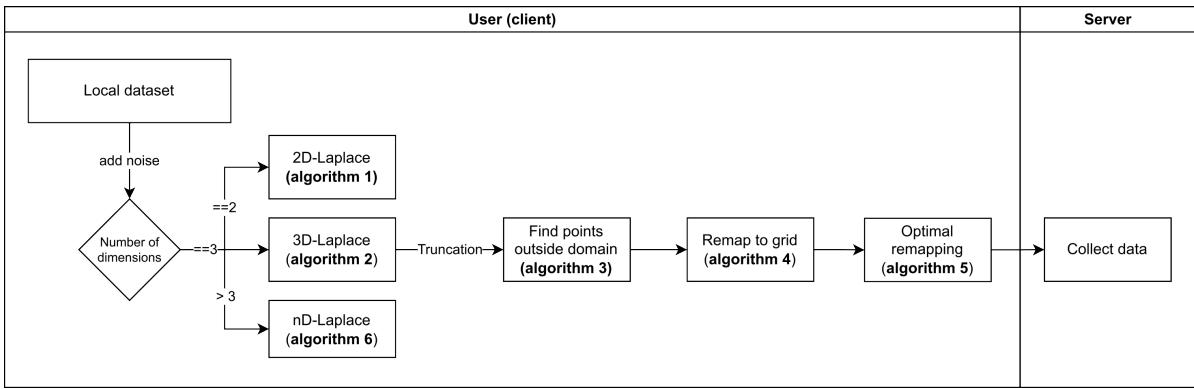


Figure 3.13: Non-interactive mechanism design for nD-Laplace.

For easy navigation, we provide a list of all algorithms:

1. 2D-Laplace: [1](#)
2. 3D-Laplace: [2](#)
3. nD-Laplace: [6](#)
4. Find points outside domain: [3](#)
5. Grid remapping: [4](#)
6. Optimal remapping: [5](#)

#### PRACTICAL EXAMPLE

The shape of the dataset is important for the usefulness of clustering. With our algorithm, there are four different shapes/variants of the dataset. To provide an example, this has been visualized using a 3D dataset based on the Cardiotocography dataset ([5.1](#)). The goal for our mechanism is to provide privacy, but also to preserve the shape of the dataset to benefit the utility of clustering. Grid remapping and ultimately optimal remapping are used to achieve this goal.

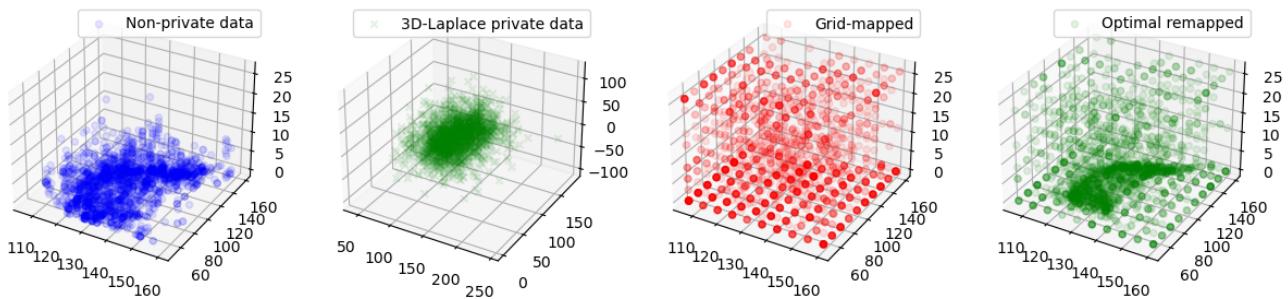


Figure 3.14: Example of optimal remapping for the 3D-dataset: Cardiotocography. The example shows the different steps of the mechanism in sequence for a dataset perturbed with a privacy budget of 0.1.

1. Dataset: the blue dots represent the original dataset without any modifications.
2. Adding noise: the green crosses represent the dataset after adding noise, for this particular example this is 3D-Laplace (Algorithm 2): As can be observed, the data is generated from the center, causing many data points to fall outside the original domain of the dataset.
3. Grid-remapping: the red dots represent the dataset after grid-remapping (Algorithm 4) After performing the grid remapping algorithm, all points within the domain are plotted. However, the original shape of the data is mostly lost. This makes it challenging to cluster the data as was possible with the original data.
4. Optimal-remapping: the green dots represent the dataset after optimal-remapping (Algorithm 5). After completing the previous step, the data points that were remapped are now again remapped based on the (original) density. This results in the restoration of the original shape of the data and consequently the clusters.

# 4

## ATTACKS ON PRIVACY

This chapter is devoted to investigating and evaluating attacks on machine learning models. Differential privacy protects the centrally stored dataset from leaking sensitive information. Assessing the performance of the algorithm is, therefore, best measurable using common attacks [Jayaraman and Evans]. In this context, we evaluate attacks specifically aimed at uncovering training data from a privately trained model. We consider two types of attacks:

1. **Membership inference attack:** An adversary attempts to infer whether a data point was used for training.
2. **Reconstruction attack:** An adversary attempts to reconstruct the training data using the model.

The knowledge of the attacker (adversarial knowledge) is an important factor to consider. This can be divided into white-box and black-box approaches [Hu et al., 2022].

1. **White-box:** The attacker has all the data that is needed. Including target model parameters, the training dataset and even the architecture [Hu et al., 2022].
2. **Black-box:** The attacker has a limited amount of information, like training data distribution and the trained model [Hu et al., 2022].

We will discuss both type of attacks and types in the next two sections.

### 4.1. MEMBERSHIP INFERENCE ATTACKS

An attack model that plays a big role in machine learning is a membership inference attack (MIA). With this attack, an adversary attempts to infer the training data  $x \in X$  (member) from a given data point  $z \in Z$  (non-member). The attack happens exclusively on supervised learning models, which either predict labels or probabilities. Most attacks on models trained on a centralized dataset occur during the inference phase, where the trained model is used to make predictions. [Rigaki and Garcia, 2021]. This is also why we are primarily interested in this phase, as we are not using a distributed learning model.

The most well-known member inference attack is training shadow models [Rigaki and Garcia, 2021]. In this attack, an attacker trains multiple models. These models do not necessarily have to be the same as the original model, and the focus is mainly on the data

input/output. It is a black-box attack, but often the attacker also needs knowledge of the data distribution to create a good shadow dataset [Rigaki and Garcia, 2021].

One of the earlier works that used this attack was Shokri et al. [Shokri et al., 2017]. An attacker trains multiple models (shadow model), with as goal to overfit the original modal. This idea is based on the fact that the model gives higher scores to the data on which it was trained (overfitting). Using this approach, attackers can retrieve the training data (member data) from the model by injecting a large amount of fake data (non-membership data).

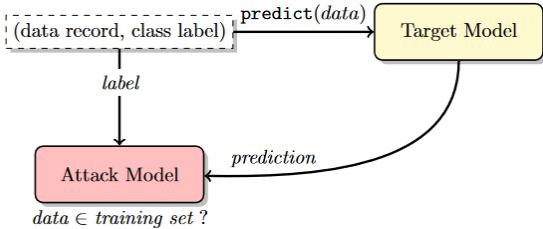


Figure 4.1: Black-box MIA attack on a machine learning model [Shokri et al., 2017]

Another approach to a black-box attack was introduced by Peng et al. and only considers that the attacker has access to the already trained model. They rescale the probabilities first using temperature scaling, to compensate for models that are overconfident [Peng et al.]. So instead of having a probability between two classes with for example 99% against 1% it will be more evenly distributed based on the training data. They then proceed in clustering the probabilities into two clusters using K-Means and label the higher confidence scores as members.

The above attacks do rely on the model to also provide the confidence or probabilities of the predictions. This is often not the case for the practical appliance of a model, and therefore Choquette-Choo et al. introduced a label-only attack. While the existing models exploit the probability output for MIA, they solely rely on labels [Choquette-Choo et al., 2021]. For this, they make use of the "HopSkipJump" attack; a so-called decision-based attack [Chen et al., 2020]. Choquette-Choo et al. consider a more semi-black-box approach, for which the attacker still requires access to a subset of the original training data and the trained model. Another paper that also uses "HopSkipJump" requires only the trained model and achieves higher accuracy by using an approach with random data [Li and Zhang, 2021].

Access to only the output of the model is a typical characteristic of black-box attacks. If the attacker also knows architecture, for example, it is referred to as a white-box attack. Another take on this is prediction and confidence-based MIA which are both proposed by [Yeom et al., 2018]. They assume that an attacker knows the standard error and has access to the perturbation dataset. The algorithm is then able to extract the truth label by minimizing the loss.

## 4.2. RECONSTRUCTION ATTACK

The concept of reconstruction attacks predates differential privacy, as this principle also gave rise to the idea of necessary database privatization [Dinur and Nissim, 2003]. Using a reconstruction attack, an adversary could reconstruct training data from a given (classifier) model. The research they did evaluates the perturbation that is needed to be added to a database to protect it versus a reconstruction attack.

A general reconstruction attack for our use-case is the attribute inference attack [Dwork et al., 2017] or model inference [Rigaki and Garcia, 2021]. Both terms are essentially the same, and we have chosen to name attribute inference attack, as it is the most common one in most literature [Jegorova et al., 2022]. Attribute inference focuses on a reconstruction attack with the adversary's goal of retrieving the secret from each user [Dwork et al., 2017]. For example, an attacker may attempt to reconstruct information about someone's heart disease using the individual's properties.

A practical implementation of the attack was provided by Fredrikson et al. as a way to infer sensitive features [Fredrikson et al., 2015]. To accomplish this, they used a decision tree attack, which was a white-box approach as they also accessed the count of instances for each decision tree branch. They also considered a black-box approach with access only to the target model (ML-as-a-service in this example). The attack targets gradient descent, which is used to optimize the input data of an attack to mimic the original data. It has only been shown to work on a neural network for face identification images (named MIFace) [Fredrikson et al., 2015], but it could be extended to another machine learning classifier if it uses gradient descent (e.g. Support Vector Machines) [Nicolae et al., 2019]. A more general approach was undertaken by Yeom et al., building upon the same research discussed in the previous section (Member inference). In their work, Yeom et al. proposed a membership inference attack, which can also be utilized for attribute inference in a white-box setting. The attribute inference attack follows a similar methodology, wherein the target model is queried to assess the loss of synthetic data, based on adversarial knowledge. By repeating the process, the attack identifies and selects the value with the highest prior probability, incorporating membership information [Jayaraman and Evans, 2022; Yeom et al., 2018]. Evaluating the effectiveness of such attacks faces a significant challenge, as it relies on the correlation between attributes, irrespective of whether the data belongs to the training or test dataset [Zhao et al., 2021]. Therefore, Yeom et al. focus on the similarity between both membership/attribute- inference and combine them to evaluate the attribute inference attack [Yeom et al., 2018].

The objective of differential privacy is to introduce sufficient noise to mitigate the risk of various attacks [Dwork et al., 2017; Jayaraman and Evans]. Assessing the privacy leakage of our model can be effectively accomplished by employing attribute/membership inference techniques.

Within this thesis, our focus is mainly on Membership Inference attacks. We aim to establish a quantifiable measure for our mechanism, and combining both attribute and membership attacks is not beneficial since they essentially capture the same information. To reach this goal, it is sufficient to concentrate on membership inference attacks.

## 4.3. ATTACK EVALUATION

In this section, we evaluate the membership inference attack and evaluate as it is the most appropriate for this study (See previous chapter). We assess whether differential privacy provides protection for the attack and discuss how this can be measured.

### 4.3.1. MEMBER INFERENCE ATTACKS

Most current research for MIA is evaluated for neural networks [Rigaki and Garcia, 2021]. Just a small percentage evaluates this attack for supervised learning, with the majority using classification with decision trees. For these attacks, most studies have used a black-box approach [Rigaki and Garcia, 2021]. This is not surprising, as these attacks have a high success rate and pose a greater risk of being exploited.

The introduction of differential privacy reduces the impact of a member inference attack [Hu et al., 2022; Rigaki and Garcia, 2021]. This is because the input to the model is perturbed. While it is still possible to retrieve the training data, the leaked privacy is significantly reduced. A simple, but effective way to measure the privacy leakage is by calculating the accuracy of correctly predicting membership by the adversary [Choquette-Choo et al., 2021]. Yeom et al. created a metric specifically for membership inference attacks which can be measured using a metric called "adversarial advantage." This metric describes the percentage of privacy that is compromised in the event of a member inference attack [Yeom et al., 2018]. This is calculated by subtracting the False Positive Rate from the True Positive Rate. The TPR represents the number of correctly predicted member data (training data) and the FPR represents the number of correctly predicted non-member data. Although these metrics are commonly applied in the literature for MIA, they do not provide enough information [Carlini et al., 2022]. Both metrics do not take into account the imbalance between the TPR and FPR. The metric should emphasize the TPR, as this is the percentage of correctly predicted member data. Therefore, they propose to use a ROC curve to show the effectiveness of a membership inference attack [Carlini et al., 2022].

In conclusion, the attacks that use member inference attacks are all based on supervised machine learning. However, in this study, we use cluster algorithms. Therefore, a semi-supervised approach can be used, as illustrated in this figure:

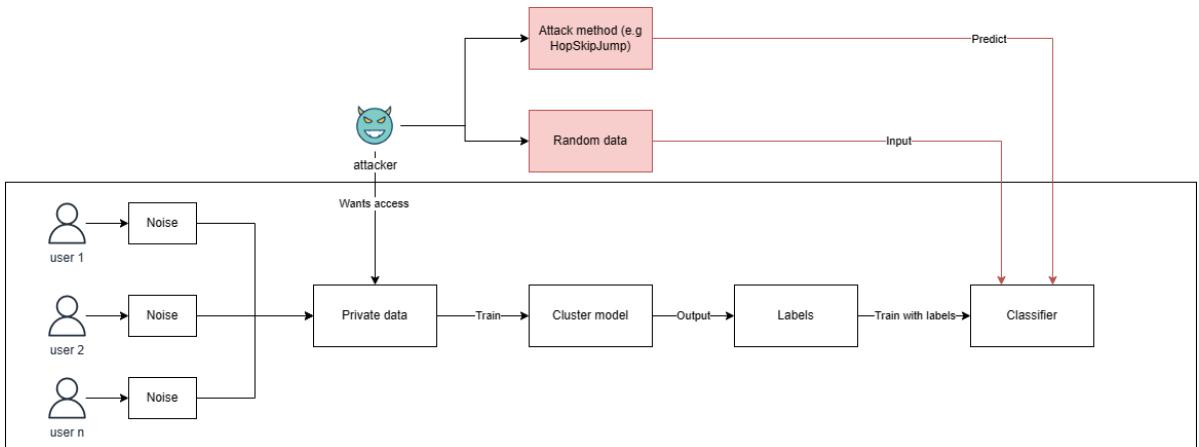


Figure 4.2: Semi-supervised black-box approach to execute a member inference attack.

# 5

## METHODOLOGY

We conducted experiments to gain insights into the proposed methods for researching the appliance of (ND)-Laplace for cluster algorithms. The experiment results are used to evaluate our method against other literature. In this chapter, we explain:

1. Datasets
2. Environmental setup.
3. For each research question: Description of the different experiments.
4. For each research question: Results.

### 5.1. DATASETS

For this research, we selected datasets based on the related papers (2.3). The datasets are sourced from the UCI Machine Learning Repository [noa].

1. **Seeds dataset**<sup>1</sup>: This dataset was used in several related works and contains 210 samples with 7 (numerical) attributes. The dataset contains information about seeds, like kernel width and density. We conducted experiments with 2, 3, and 7 dimensions and decided to use the following features (based on the correlation between the features):
  - (a) 2-dimensional data: area and perimeter.
  - (b) 3-dimensional data: area, perimeter, and length of the kernel.
  - (c) 7-dimensional data: All numerical features.
2. **Heart dataset**<sup>2</sup>: This dataset is selected because of the mixed data and amount of instances. It has 23 attributes, ten numerical attributes, and 2126 samples. The dataset contains information about measurements of fetal heart rate (FHR) and uterine contraction (UC). We conducted experiments with 2, 3, and 10 dimensions and decided to use the following features (based on the correlation between the features):
  - (a) 2-dimensional data: FHR baseline and histogram-min.

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<sup>1</sup><http://archive.ics.uci.edu/ml/datasets/seeds>

<sup>2</sup><https://archive.ics.uci.edu/ml/datasets/cardiotocography>

- (b) 3-dimensional data: FHR baseline, histogram-min, and accelerations.
- (c) 10-dimensional data: All numerical features.

The following datasets are for a component investigated in research question 3. Further detail is provided in Section: [5.3.4](#).

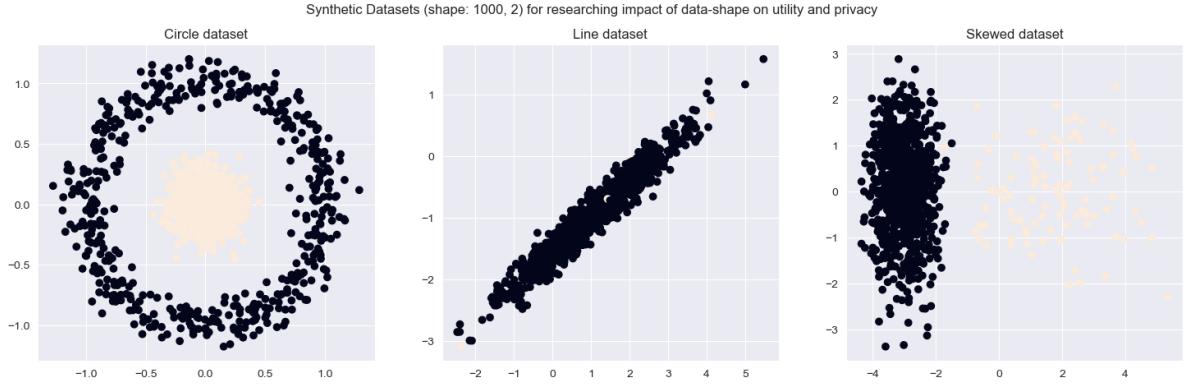


Figure 5.1: Synthetic datasets with 1000 samples and 2-dimensions

1. **circle dataset:** Dataset with a circle shape.
2. **line dataset:** Dataset is a line (shape is much like the seeds dataset).
3. **skewed dataset:** Dataset samples are skewed/concentrated to one side (shape is much like the heart dataset).

## 5.2. ENVIRONMENTAL SETUP

For running the experiments, we use 16GB RAM and an i7-10750H 2.6Ghz processor. The experiments are run using a Docker container which runs a pre-configured distribution of Linux Alpine. It includes a pre-installed Anaconda environment for Python <sup>3,4</sup>. We run the container using the dev-container feature for visual-studio code <sup>5</sup>. This allows us to create a reproducible experiment environment.

### 5.2.1. LIBRARIES & CODE VERSIONS

We use Python version 3.9.13 with Jupyter Notebook for creating a reproducible experimental environment. The packages for Python are:

1. Scikit-learn: 1.0.\*
2. Yellow-brick: 1.5
3. Numpy: 1.24.\*
4. Pandas: 2.0.\*

<sup>3</sup><https://github.com/devcontainers/images/tree/main/src/anaconda>

<sup>4</sup>tag: mcr.microsoft.com/devcontainers/anaconda:0-3

<sup>5</sup><https://code.visualstudio.com/docs/devcontainers/containers>

5. Seaborn: 0.11.\*
6. Mathplotlib: 3.5.\*

## 5.3. METHODS

This section explains what methods/ algorithms we used and how we evaluate them.

### 5.3.1. PRIVACY MECHANISMS

We will evaluate two different privacy frameworks called kd-Laplace (our method) and piecewise (2.8). For the first one, we test three configurations:

1. kd-Laplace: Plain mechanism without any additions for truncation.
2. kd-Laplace/grid: Mechanism with truncation using grid-remapping (See Equation: 4).
3. kd-Laplace/grid/optimal: Mechanism with truncation using grid-remapping in combination with optimal-remapping (See Equation: 5).

The variants of kd-Laplace are different on the implementation level for 2, 3, and n-dimensional data (See figure 3.13). Piecewise, however, will be the same for any data dimension.

We evaluate all privacy mechanisms by comparing them in utility and privacy. To reduce the measurement bias of results, we executed them ten times for multiple privacy budgets and reported the average for each [Huang et al., 2021].

1. The experiments run ten times, and we report the mean.
2. All experiments run for multiple epsilons: 0.1, 0.5, 1, 2, 3, 5, 7, 9.

Firstly, we primarily focus on evaluating the utility in terms of clustering. Then, we shift our focus to comparing the privacy mechanisms themselves. Lastly, we compare the four mechanisms in terms of privacy.

### 5.3.2. CLUSTERING METHODS

For the three different algorithms: K-Means, AP, and OPTICS, we analyzed the most important decisions regarding parameter selection (See section 3.1.3). This section gives a short list and explanation of the parameters we used throughout the experiments. In addition to this, we also provide the corresponding algorithms as implemented by the Scikit-learn package.

#### K-MEANS

$$\sum_{i=0}^n \min_{\mu_j \in C} (\|x_i - \mu_j\|^2) \quad (5.1)$$

Parameter	Description	Value	Dataset
K-value	Calculated based on an "elbow" plot.	4 (see figure 1)	Seeds dataset
K-value	""	<b>TODO</b>	Heart dataset

Table 5.1: K-Means hyperparameters for dataset 1 - 3

#### AFFINITY PROPAGATION

As specified in section 2.2.1, the clustering algorithm has two types of similarity. The following formula calculates the responsibility:

$$r(i, k) \leftarrow s(i, k) - \max[a(i, k') + s(i, k') \forall k' \neq k] \quad (5.2)$$

Then the availability is given using this formula:

$$a(i, k) \leftarrow \min[0, r(k, k) + \sum_{i' \text{ s.t. } i' \notin \{i, k\}} r(i', k)] \quad (5.3)$$

And iteratively calculated while considering the damping factor.

Parameter	Description	Value	Dataset
Preference	We decided to use the median similarity as described in section 2.2.1	Median	Seeds dataset
Preference	""	Median	Heart dataset
Damping factor	Default value as specified in section 2.2.1	0.5	Seeds dataset
Damping factor	""	0.5	Heart dataset

Table 5.2: Affinity Propagation hyperparameters for the datasets

#### OPTICS

As previously explained in section 2.2.1, we use OPTICS to automatically implement DBSCAN to determine the epsilon/radius. Instead of referring to DBSCAN, we will now refer to OPTICS in the remaining parts of this thesis.

Provide algorithm implementation

Parameter	Description	Value	Dataset
Minimum points	Decided using the formula $minPts = n * 2$ , where n is the number of features (2.2.1)	4	Seeds dataset (2-dimensions)
Minimum points	""	6	Seeds dataset (3-dimensions)
Minimum points	""	14	Seeds dataset (7-dimensions)
Minimum points	Decided using the formula $minPts = n * 2$ , where n is the number of features (2.2.1)	4	Heart dataset (2-dimensions)
Minimum points	""	6	Heart dataset (3-dimensions)
Minimum points	""	20	Heart dataset (10-dimensions)

Table 5.3: DBSCAN hyperparameters for datasets 1 - 3

### 5.3.3. UTILITY AND PRIVACY EVALUATION

#### UTILITY

To measure cluster utility, both internal and external validation methods are used:

1. **External validation:** The external validation is measured by comparing the labels of non-private trained cluster algorithms with those trained using a privacy mechanism. The outcome is between 0 and 1, where 1 indicates the highest similarity (thus the best result). AMI (Adjusted Mutual Information) and ARI (Adjusted Rand Index) are used to assess the external validity of the cluster algorithms (See section 2.2.2).
2. **Internal validation:** The internal validation measures the intrinsic properties of the clustering algorithms. The outcome is a value between -1 and 1 for the SC evaluation. Where -1 indicates incorrect clustering and 1 dense clustering. Another metric we use is the CH metric, where a higher value suggests good clustering.

Because the cluster algorithms rely on Euclidean distance, we need to apply some data standardization. For this purpose, we use standard scaling provided by the Scikit-learn package<sup>6</sup>.

#### PRIVACY

Privacy is hard to quantify, but we can measure the privacy loss/gain by calculating the Euclidean distance between the non-perturbed and perturbed data. In addition, we evaluate privacy by simulating a membership inference attack and calculating the adversary advantage.

1. **Privacy distance:** The first measurement we evaluate is the distance between the dataset's non-private and private variants. This metric gives us a sense of how much extra distance (and privacy) the privacy mechanisms offer compared to the non-private variant. To this end, the average Euclidean distance is measured and reported per epsilon.
2. **Adversarial advantage:** We used the Membership Inference Attack (MIA) that Shokri et al. proposed with the implementation that was provided by Adversarial Robustness Toolkit (ART) [Nicolae et al., 2019]. An earlier study also explored this attack to

<sup>6</sup><https://scikit-learn.org/stable/modules/preprocessing.html>

evaluate differential privacy. Similar to this attack, we train a classifier with perturbed data and evaluate it using non-perturbed data (test-data / shadow-data) [Zhao et al., 2020]. We consider a semi-supervised setup (figure 4.2), where we train a classifier with the perturbed data and evaluate it using non-perturbed data (test-data / shadow-data). For the classifier, we use a RandomForestClassifier [Rigaki and Garcia, 2021]. Finally, we evaluate the adversary advantage (percentage) using  $TPR - FPR$  [Yeom et al., 2018]. Figure: 5.2 provides a visual setup of the experiment.

Although the adversary advantage is a proven method for measuring MI attacks, we also include the TPR as an essential measurement. This metric is most interesting to us because this displays the percentage of the actual leaked labels (see Section 4.3).

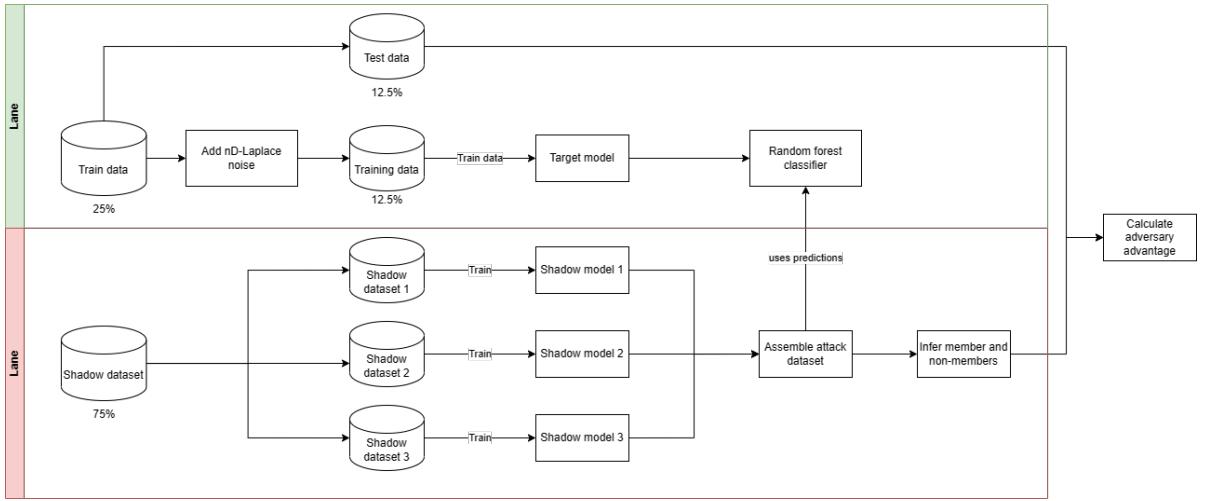


Figure 5.2: Member inference attack using shadow models. The green swim lane illustrates the normal setup, and the red swim lane projects the adversary steps.

Unlike the other utility metrics, the lower the adversary score and TPR, the better.

#### **5.3.4. RESEARCH QUESTION 3**

In this research question, we compare the behavior of our mechanism with what we have observed in the literature for similar mechanisms. To do this, we have formulated several hypotheses.

1. *Adding optimal remapping improves utility without sacrificing privacy:* Since the noise is updated based on the density of the data points, the clusters stay preserved (5). As a result, the utility increases while the data points remain indistinguishable. Therefore, we have extended our mechanism with optimal remapping and named it kd-Laplace/grid/optimal.
2. *The privacy leakage (adversary advantage) increases for a higher number of dimensions:* The hypothesis arose from our observation while investigating/implementing kd-laplace/grid/optimal (3.3.2). Adding more dimensions ( $5 >$ ) decreases the added noise gradually. The study is conducted for the kd-Laplace/grid/optimal mechanism and compared with the Piecewise mechanism.
3. *The shape of the data negatively impacts the kd-Laplace mechanism in terms of privacy:* Because the mechanism heavily relies on Euclidean distance, the shape (distance between the data points) can ultimately harm utility. We have already observed this difference when comparing the heart and seed datasets in research questions 1 and 2. But, to rule out the potential impact of the number of data points (the heart dataset has 2126 samples and the seed-dataset 210), we generate three synthetic datasets, each with 1000 samples (See section 5.1).

# 6

## RESULTS

This chapter aims to present the results to the reader. We tested four mechanisms for their utility and privacy, as indicated in the methodology. For all four mechanisms, we applied a prefix to indicate the number of dimensions on which the mechanism was used. For example, we named the piecewise mechanism used on 2-dimensional data as "2d-piecewise," even though the underlying method is not different from the original piecewise method. The results are presented in the following order:

1. Cluster utility: We use external validation (AMI/ARI) to compare the results between the three different cluster algorithms when trained privately with nD-laplace (3.13) and piecewise.
2. Mechanism utility: We compare the three types of Laplace (normal, grid-based, and optimal truncation) and piecewise, using external validation.
3. Privacy: We compare the three types of nD-laplace and piecewise using the ROC-curve for evaluating the membership inference attack (MIA).
4. Dimensionality: We evaluate the number of dimensions' influence on utility and privacy.
5. Shape: We are investigating the shape of the data to measure its impact on the utility and privacy of the Laplace mechanism.

Points 1 to 3 are structured according to the number of dimensions. So we start with 2-dimensional data (2D-Laplace), then 3-dimensional data (3D-Laplace), and finally n-dimensional data (nD-Laplace). See the mechanism flowchart for a reference structure: 3.13. All results are reported for each dataset (See methodology: 5.1) separately. And after each section, a summary is provided for the results.

## 6.1. CLUSTER UTILITY

The results below display the difference between the three cluster algorithms for Piecewise and Laplace with optimal truncation for 2/3/n-dimensional data.

The first figure shows the seeds dataset, and the second shows the heart dataset. Displaying the different cluster algorithms and their corresponding hyperparameters is done using the legend. The x-axis shows the privacy budget, and the y-axis shows the adjusted mutual information (AMI) and the adjusted rand index (ARI). We conducted this research for the following cluster algorithms.

1. **K-Means:** Displayed as a blue line.
2. **Affinity Propagation:** Displayed as a red line. For the larger datasets ( $>1000$  data points with 3 > dimensions), the experiments for the affinity propagation cluster algorithm were omitted. This algorithm required too much computational power for our experimental setup.
3. **OPTICS:** Displayed as a green line.

Please refer to the plots in the appendix for internal validation and the other variants (Laplace / Laplace-truncated): [3](#).

### 6.1.1. 2-DIMENSIONAL DATA

Figure 6.1: External validation piecewise & kd-Laplace/grid/optimal for the 2-dimensional data seeds-dataset

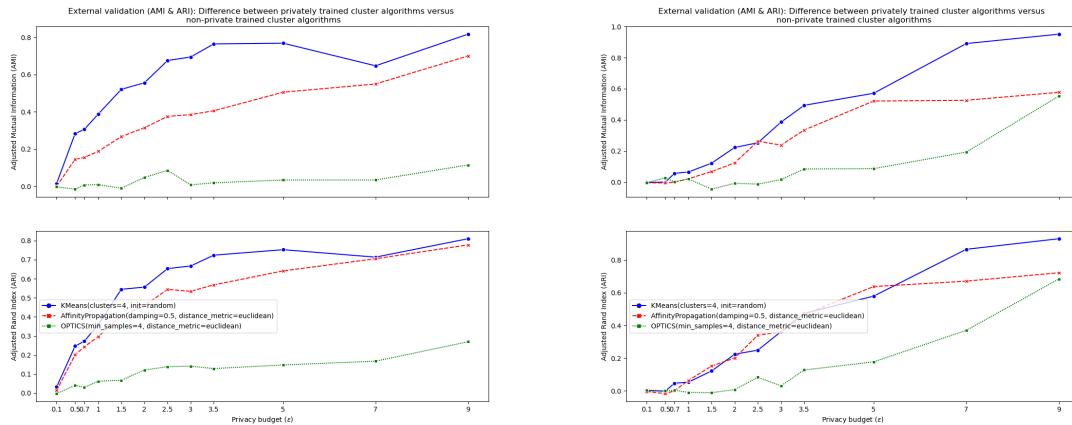


Figure 6.2: External validation (ARI/ AMI) for the 2-dimensional data seeds-dataset for kd-Laplace with optimal truncation

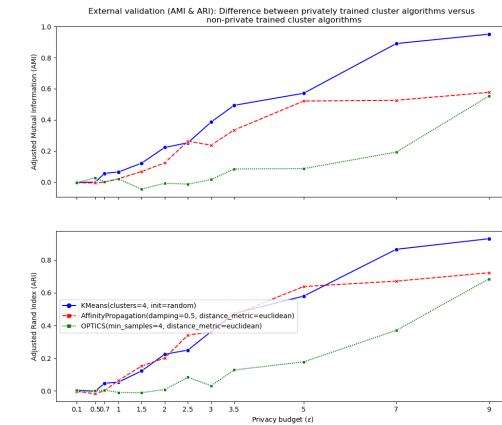


Figure 6.3: External validation (ARI/ AMI) for the 2-dimensional data seeds-dataset for piecewise mechanism

The above plots show the AMI and ARI scores for the seeds-dataset with kd-laplace/grid/optimal (left-side) and piecewise (right-side). Piecewise shows a higher AMI / ARI for epsilons 7 and 9. In contrast, the kd-Laplace mechanism scores higher for all the other epsilons (0.1 onward 7). K-Means achieves the best score for kd-Laplace with a slight margin compared to AP. However, for the Piecewise mechanism, K-Means is the clear winner. For both the mechanisms, OPTICS underperforms heavily but still scores the same as AP for Piecewise.

Figure 6.4: External validation piecewise & kd-Laplace/grid/optimal mechanisms for the 2-dimensional data heart-dataset

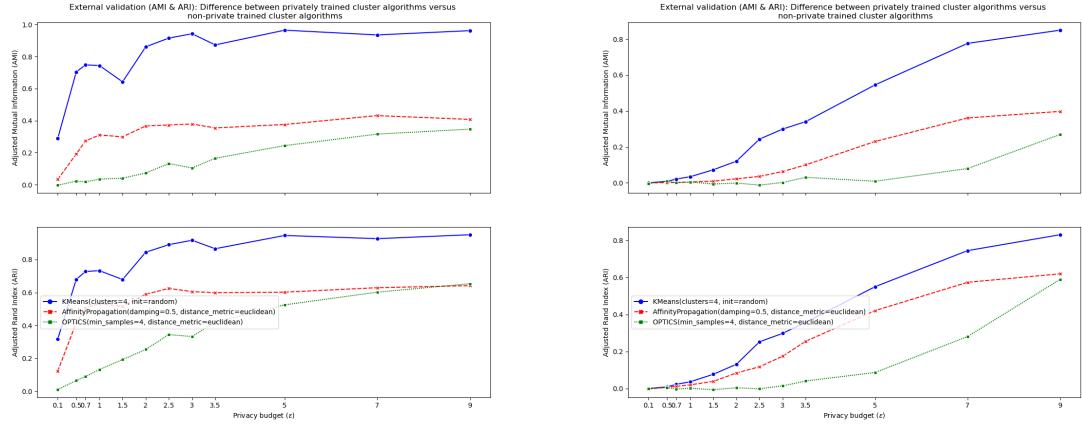


Figure 6.5: External validation (AMI/ ARI) for the 2- dimensional data heart-dataset for kd-Laplace with optimal truncation

Figure 6.6: External validation (AMI/ ARI) for the 2-dimensional data heart-dataset for piecewise mechanism

The above plots show the AMI and ARI scores for the heart dataset with kd-laplace/grid/optimal (left-side) and piecewise (right-side). The kd-Laplace mechanism performs better (0.8 - 1.0 from epsilon 1.0 to 9) for all epsilons for K-Means. AP and OPTICS show the same trend, reaching their peaks around epsilon 7 and 9. Both underperform by delivering maximum results for these epsilons below 0.3-0.4 AMI/ARI.

### 6.1.2. 3-DIMENSIONAL DATA

Figure 6.7: External validation piecewise & kd-Laplace/grid/optimal for the 3-dimensional data seeds-dataset

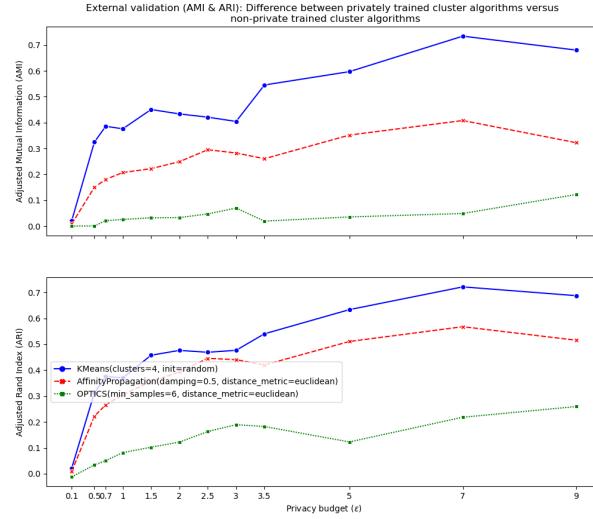


Figure 6.8: External validation (ARI/ AMI) for the 3-dimensional data seeds-dataset for kd-Laplace

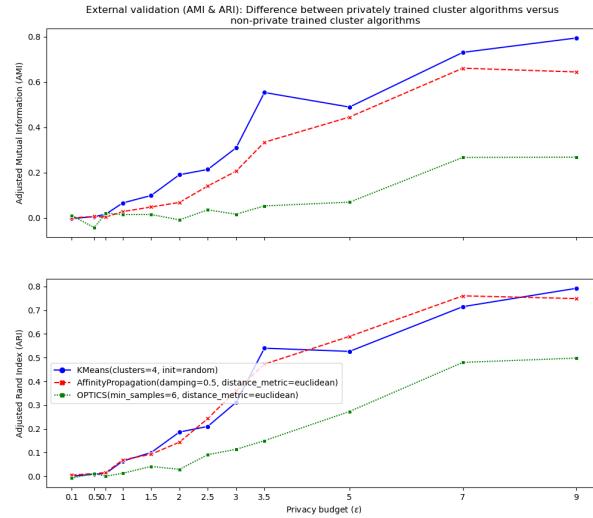


Figure 6.9: External validation (ARI/ AMI) for the 3-dimensional data seeds-dataset for piecewise mechanism

The above plots show the AMI and ARI scores for the seeds dataset with kd-laplace/grid/optimal (left-side) and piecewise (right-side). Piecewise performs better for epsilon 9 but worse for the lower epsilons (0.1 - 5). K-Means scores 0.7 - 0.8 AMI/ARI, followed by AP, which is 0.6. The kd-Laplace mechanism, on the other hand, shows more difference (K-Means 0.65 - 0.7 and AP 0.3 - 0.2). OPTICS scores again low but shows a slight upwards trend when the epsilon increases.

Figure 6.10: External validation piecewise & kd-Laplace/grid/optimal mechanisms for the 3-dimensional data heart-dataset

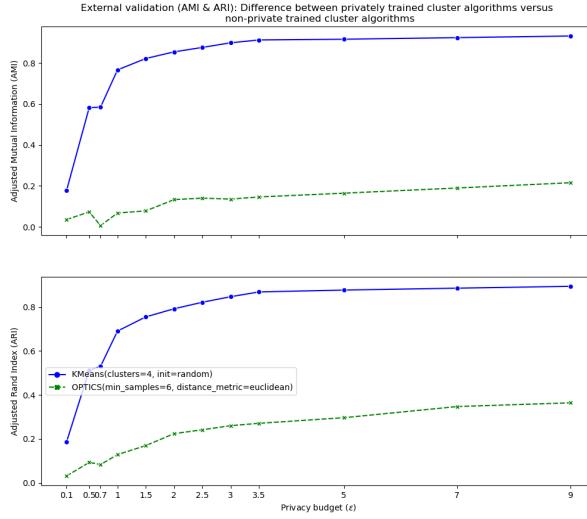


Figure 6.11: External validation (ARI/ AMI) for the 3-dimensional data heart-dataset for kd-Laplace with optimal truncation

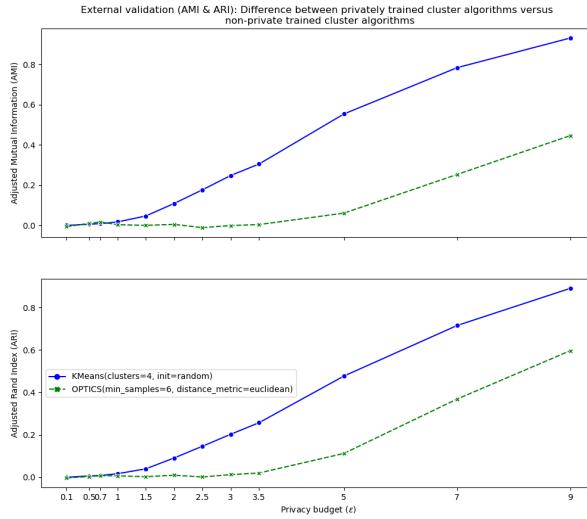


Figure 6.12: External validation (ARI/ AMI) for the 3-dimensional data heart-dataset for piecewise mechanism

The above plots show the AMI and ARI scores for the heart dataset with kd-laplace/grid/optimal (left-side) and piecewise (right-side). Piecewise for K-Means performs well (0.78 and 0.93 ARI) for epsilon 7 and 9, respectively. However, from epsilon 1 onwards kd-Laplace scores 0.77 to 0.93 ARI for K-Means. On the other hand, OPTICS scores low for kd-Laplace. It shows an upward trend for Piecewise as the algorithm scores 0.45 ARI and 0.6 AMI for epsilon 9. While kd-Laplace scores 0.22 ARI and 0.36 AMI for epsilon 9. OPTICS scores below 0.4 ARI/AMI for both mechanisms if the epsilon is  $< 9$ .

### 6.1.3. N-DIMENSIONAL DATA

Figure 6.13: External validation piecewise & kd-Laplace/grid/optimal mechanisms for the n-dimensional data seeds-dataset

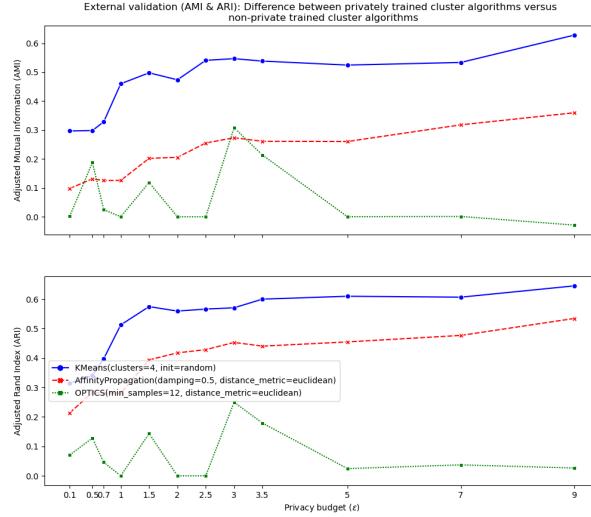


Figure 6.14: External validation (ARI/ AMI) for the n-dimensional data seeds-dataset for kd-Laplace with optimal truncation

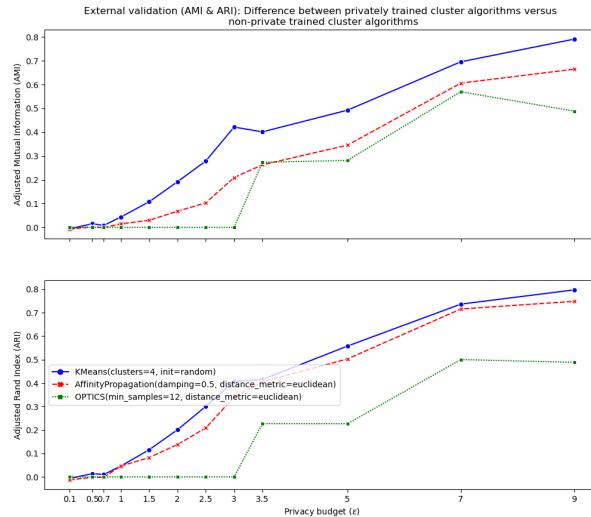


Figure 6.15: External validation (ARI/ AMI) for the n-dimensional data seeds-dataset for piecewise mechanism

The above plots show the ARI and AMI scores for the heart dataset with kd-laplace/grid/optimal (left-side) and piecewise (right-side). Piecewise performs well (0.89 and 0.87 ARI) for epsilon 9 and 7 for K-Means, respectively. In comparison, kd-Laplace scores worse for the same epsilon values (0.56 and 0.51 ARI); on the other hand, kd-Laplace scores better for all the other epsilons. After K-Means, the best scoring cluster algorithm is OPTICS for Piece-

wise, with a score of 0.62 ARI for 9 epsilon. However, the other epsilons AP scores better. For kd-Laplace, AP is second after K-Means, and OPTICS scores much worse for AMI. But, for ARI, the scores for OPTICS and AP lay close together for epsilon 9.

Figure 6.16: External validation piecewise & kd-Laplace/grid/optimal mechanisms for the n-dimensional data heart-dataset

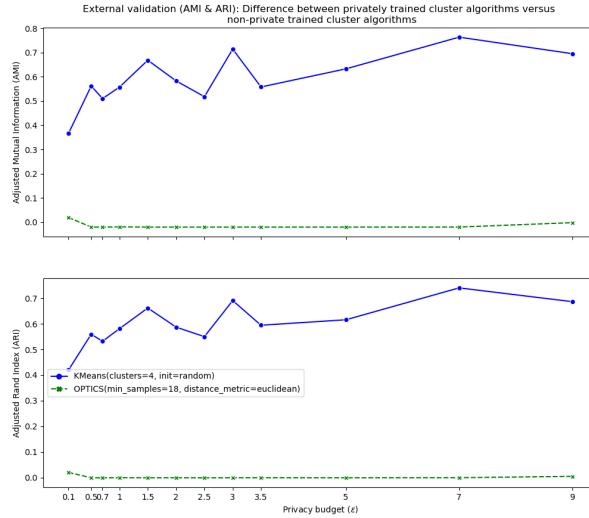


Figure 6.17: External validation (ARI/ AMI) for the n-dimensional data heart-dataset for kd-Laplace with optimal truncation

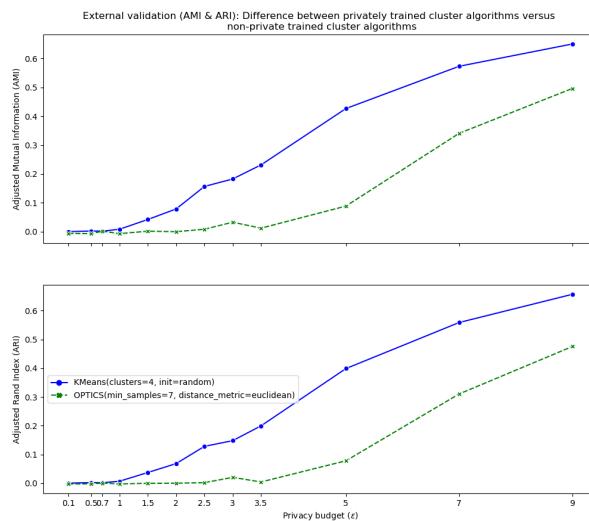


Figure 6.18: External validation (ARI/ AMI) for the n-dimensional data heart-dataset for piecewise mechanism

The above plots show the ARI and AMI scores for the heart dataset with kd-laplace/grid/optimal (top) and Piecewise (bottom). Piecewise shows similar trends for K-Means and OPTICS. But, K-Means scores better (0.65 ARI) for epsilon 9, while OPTICS peaks at 0.5 ARI. However,

Piecewise scores (<0.2 ARI) for epsilon 0.1 til 3.5. In contrast, kd-Laplace scores more evenly (0.6 - 0.7 ARI) between epsilon 3.5 and 9. Conversely, OPTICS scores lower (< 0.2 ARI) for all epsilon values.

#### 6.1.4. SUMMARY

The seeds-dataset and heart-dataset show similarities. K-Means performs the best, followed by AP and OPTICS. Noteworthy is the fact that kd-Laplace overall performs better on the heart dataset than the seeds dataset.

Still has to be written

## 6.2. MECHANISM UTILITY

In the images below, the different types of Laplace mechanisms and piecewise are compared using a barplot. Since ARI and AMI provide the same information, we have shown only the Adjusted Mutual Information (AMI) to avoid redundant information. For the same reason, only the K-Means algorithm was used to calculate the AMI. The x-axis displays the privacy budget, and the y-axis shows the corresponding AMI. The privacy mechanisms are displayed as follows:

1. **Piecewise:** Displayed as a yellow bar or line.
2. **kd-Laplace:** Displayed as a red bar or line.
3. **kd-Laplace/grid:** Displayed as a blue bar or line.
4. **kd-Laplace/grid/optimal:** Displayed as a green bar or line.

Please refer to the plots in the appendix for internal validation: [.3](#).

### 6.2.1. 2-DIMENSIONAL DATA

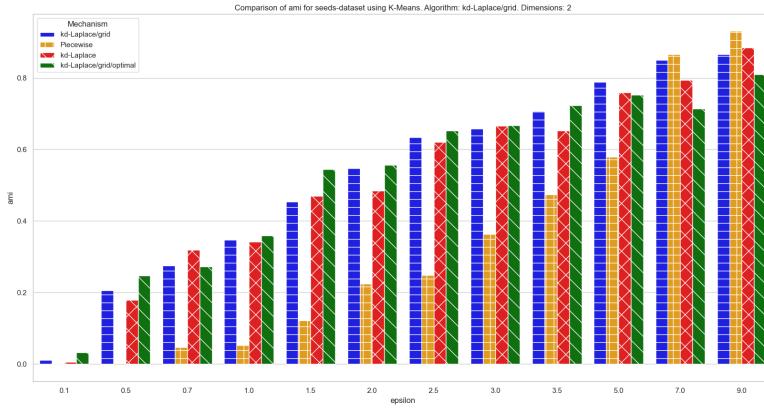


Figure 6.19: Adjusted Mutual Information comparison for the seeds-dataset

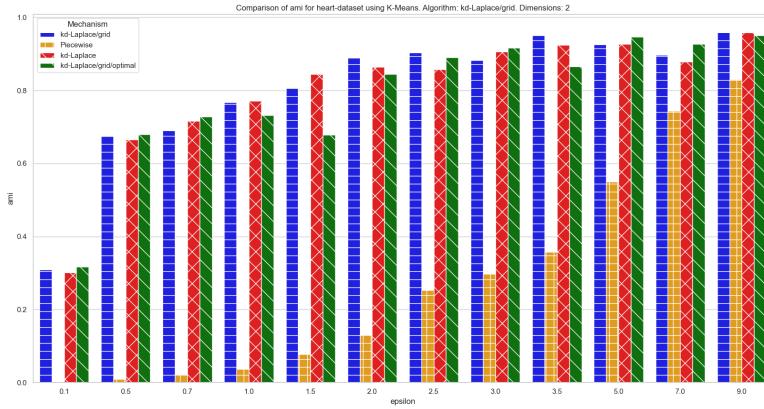


Figure 6.20: Adjusted Mutual Information comparison for the heart-dataset

The two plots compare external validation (AMI) between the two privacy frameworks, Piecewise, and kd-Laplace, with three variants. A clear trend in AMI can be observed in the top plot, with the lowest epsilon (0.1) yielding the lowest score for the mechanisms. Piecewise consistently scores lower for all epsilons, except for epsilons 7 and 9. Additionally, we notice minimal differences among the variants of kd-Laplace, but kd-Laplace/grid/optimal performs slightly better for epsilon 0.1 while scoring worse for 7 and 9.

### 6.2.2. 3-DIMENSIONAL DATA

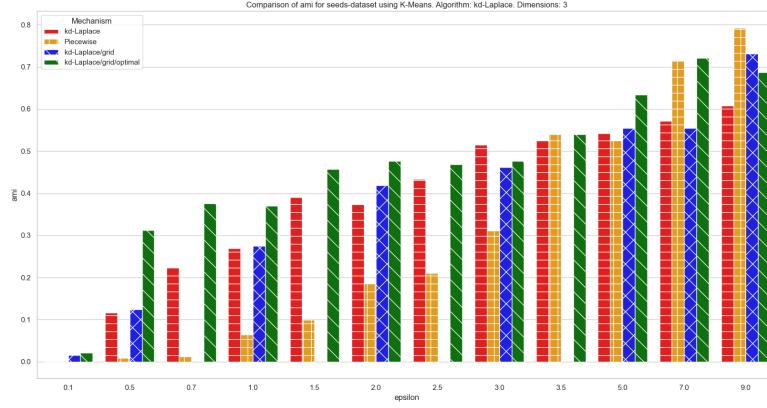


Figure 6.21: Adjusted Mutual Information comparison for the 3-dimensional seeds-dataset

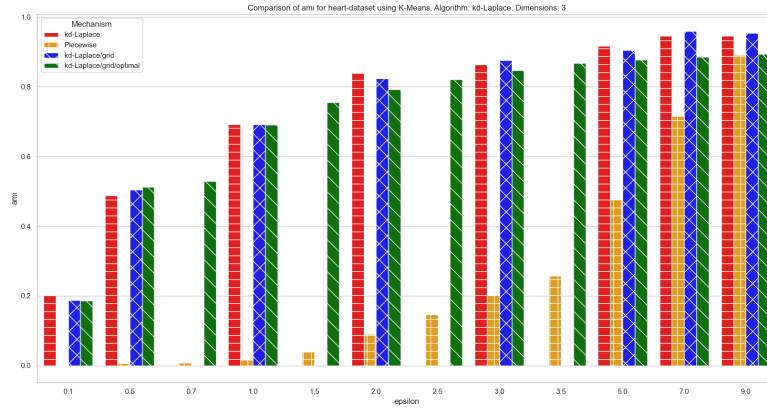


Figure 6.22: Adjusted Mutual Information comparison for the 3-dimensional heart-dataset

Between the seeds and heart datasets, the kd-Laplace mechanism with variants scored better overall than the piecewise mechanism. The Piecewise mechanism scores low (< 0.5 ARI) for the epsilons 0.1 to 3.5 in the seeds dataset. In contrast, kd-Laplace with variants reach this score for around epsilon 3.0. For epsilon 7 and 9, the Piecewise mechanism catches up with kd-Laplace and scores around equal (0.7 & 0.8 ARI). Kd-Laplace/grid/optimal scores are a little better for epsilon 5 and 7, but the difference between the variants is minimal overall.

The kd-Laplace mechanism with variants score > 0.6 AMI for epsilon 1 and higher. As with the seeds dataset, the Piecewise mechanism scores low (< 0.5 ARI) for the epsilons 0.1 to 7.0. For epsilon 9, the Piecewise mechanism scores around 0.8 AMI, which equals kd-Laplace. Between the different kd-Laplace variants, the difference is small. But, kd-Laplace/grid/optimal scores are slightly worse for all epsilons.

### 6.2.3. N-DIMENSIONAL DATA

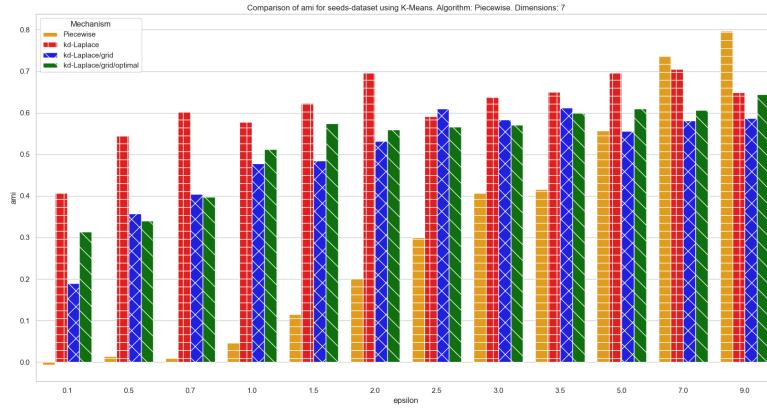


Figure 6.23: Adjusted Mutual Information comparison for all numerical features for the seeds-dataset

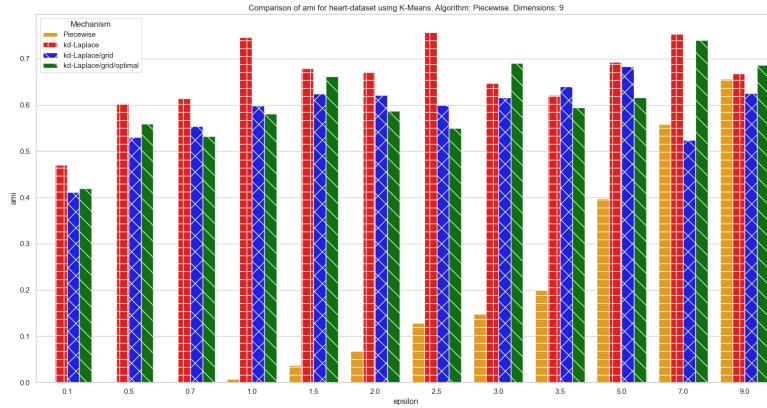


Figure 6.24: Adjusted Mutual Information comparison for all numerical features for the heart-dataset

The above plot represents the Adjusted Mutual Information (AMI) for n-dimensional data. The seeds dataset (top) shows a clear peak for kd-Laplace without any optimization. Piecewise consistently scores lower than the kd-Laplace variants for most epsilon values, except for epsilon 7 and 9 where Piecewise performs better (AMI > 0.7). The latter performs slightly better among kd-Laplace/grid and kd-Laplace/grid/optimal variants.

For the heart dataset (bottom graph), the scores are relatively similar across all kd-Laplace variants. Piecewise scores below 0.5 AMI until epsilon 7 and performs at an equivalent level as the kd-Laplace variants for epsilon 9. The kd-Laplace variants score significantly better, as they already surpass 0.5 AMI from epsilon 0.5. Among the kd-Laplace variants, kd-Laplace without optimization stands slightly higher (AMI > 0.7).

### 6.3. PRIVACY

The images below display a ROC curve plot of membership inference attacks as described in the methodology. There are three baseline values: The red and green dashed lines represent the baseline True Positive Rate (TPR) and False Positive Rate (FPR), respectively. These baselines represent the attack without using any privacy mechanism. The "random guess" is the baseline for the ROC curve and reflects the 50% chance that an adversary has in determining whether a data point is a member. Anything above this baseline can be considered an improvement in that chance. The y-axis represents the TPR, and the x-axis represents the FPR. For each mechanism, a ROC curve is shown using different lines. Finally, we provide a table with the mechanism and mean of the epsilon that is outside the specification of the TPR (TPR of the non-private dataset).

We use the same color scheme in the previous section to display the different mechanisms. With as addition, a red line is drawn to indicate the baseline TPR (the non-private dataset's TPR).

Please refer to the plots in the appendix for comparing the other variants (laplace / laplace-truncated): [.3](#).

Add interpretation of results

### 6.3.1. 2-DIMENSIONAL DATA

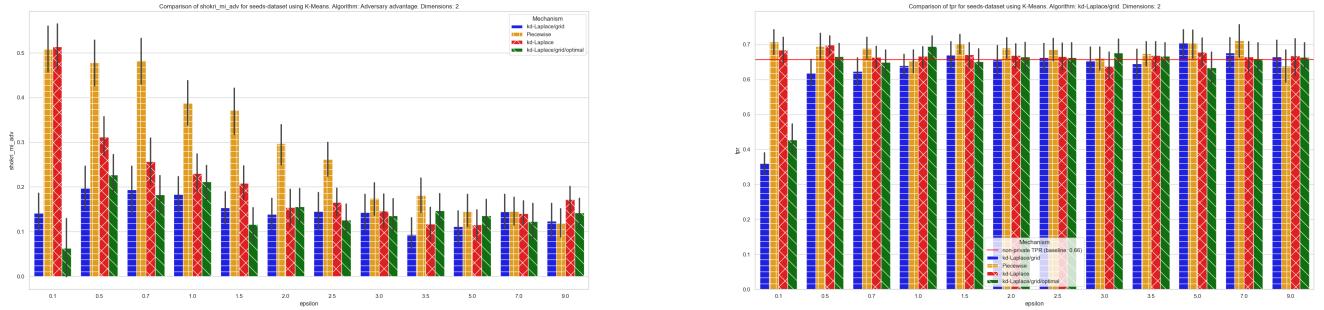


Figure 6.25: Barplot for adversary advantage (left) and TPR (right) per privacy mechanism for seeds-dataset.

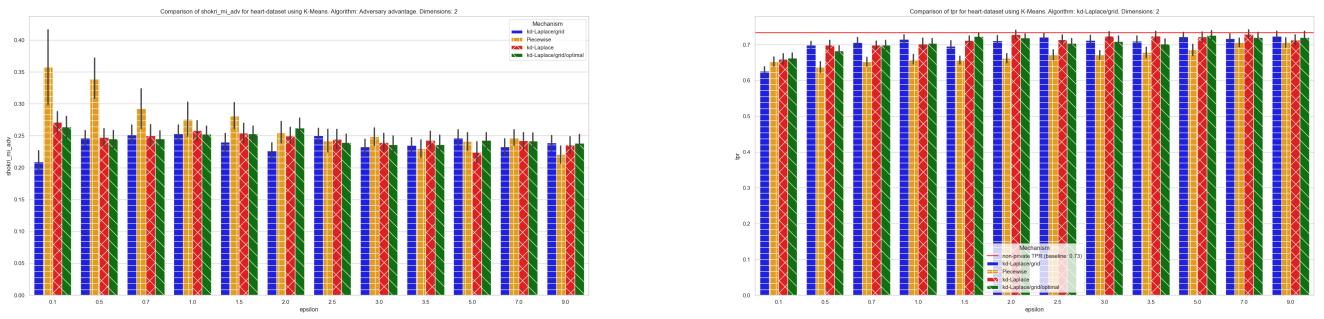


Figure 6.26: Barplot for adversary advantage (left) and TPR (right) per privacy mechanism for heart-dataset.

The above bar plot shows the adversary advantage based on the membership inference attack. The first graph displays the attack on the seeds dataset, where a clear peak of 0.5 adversary advantage is visible for both Piecewise and kd-Laplace at epsilon 0.1. The other variants of kd-Laplace score similarly, and all fall below 0.2. Piecewise reaches below 0.2 at epsilon 3, while kd-Laplace reaches it at epsilon 2.

For the heart dataset, the difference is slightly smaller, but Piecewise spikes to 0.36 and 0.34 adversary advantage at epsilon 0.1 and 0.5, respectively. Furthermore, kd-Laplace and Piecewise hover around 0.25 adversary advantage for the remaining epsilons.

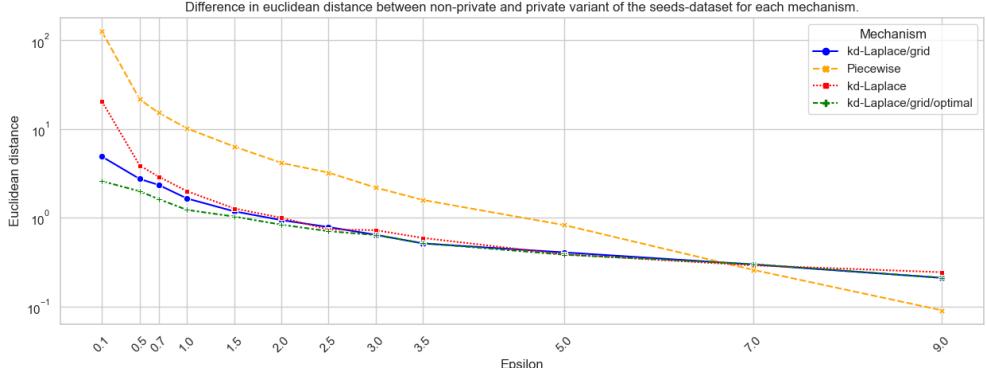


Figure 6.27: Privacy distance for each mechanism for seeds-dataset.

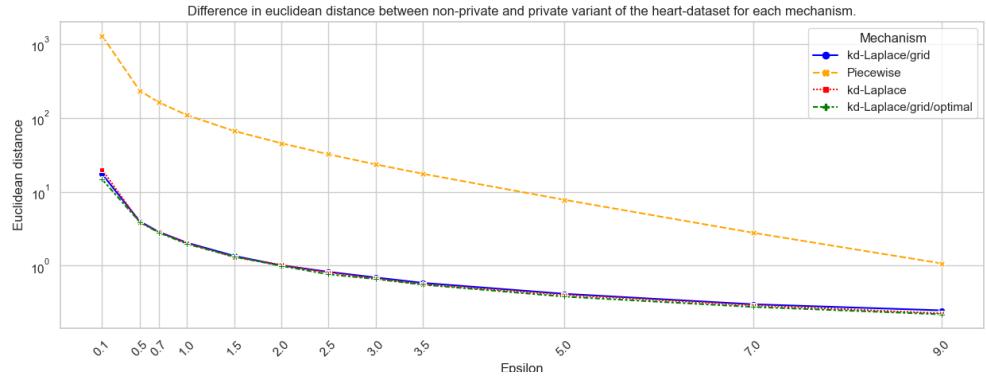


Figure 6.28: Privacy distance for each mechanism for heart-dataset.

The above graph represents the average Euclidean distance added by each privacy mechanism compared to the two datasets. We use a logarithmic scale for the Euclidean distance to better visualize the results. For both datasets, the Piecewise mechanism adds the most privacy. When considering the kd-Laplace variants on the seeds dataset, kd-Laplace without optimizations adds the most distance, followed by kd-Laplace/grid and kd-Laplace/grid/optimal. For the heart dataset, all variants of kd-Laplace are nearly equal.

### 6.3.2. 3-DIMENSIONAL DATA

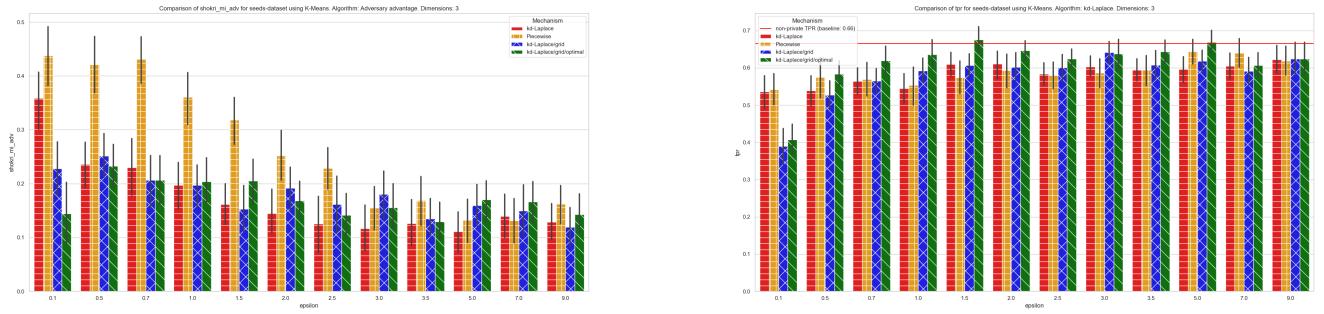


Figure 6.29: Barplot for adversary advantage (left) and TPR (right) per privacy mechanism for seeds-dataset.

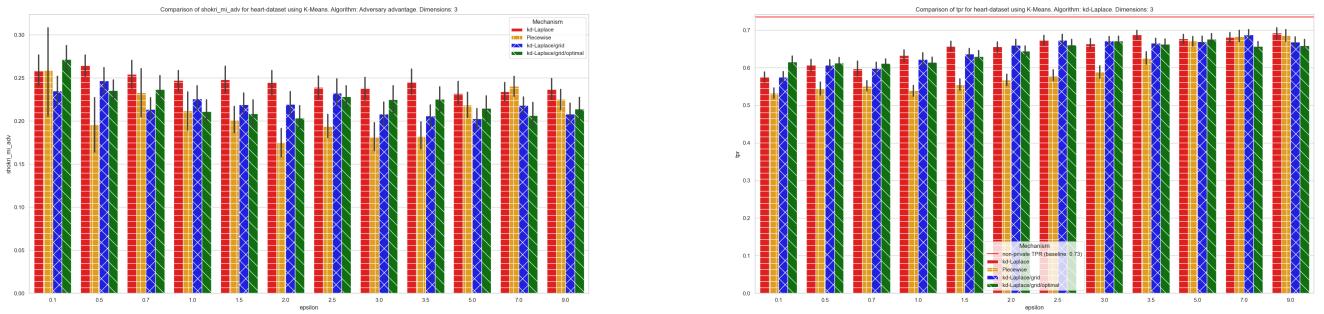


Figure 6.30: Barplot for adversary advantage (left) and TPR (right) per privacy mechanism for heart-dataset.

The graphs above show the adversary advantage (left) and TPR (right) for the seeds dataset (top) and heart dataset (bottom). The first dataset we analyze is the seeds dataset. We can observe a clear pattern for the Piecewise mechanism based on the adversary advantage plot for the seeds dataset. It scores between 0.4 and 0.5 for epsilon values ranging from 0.1 to 0.7 and drops below 0.2 for epsilon values higher than 3. The kd-Laplace mechanism, particularly the variant without optimizations (red), has a high adversary advantage (0.35) for epsilon 0.1. After that, all kd-Laplace variants perform similarly. When we compare them to the TPR, we still see that Piecewise and the regular kd-Laplace variant have higher scores for epsilon 0.1 (0.5+). After that, the mechanisms perform similarly, but we notice that kd-Laplace/grid/optimal consistently outperforms the others. Additionally, the latter is above the baseline value for epsilon 1.5.

Now, let's turn our attention to the heart dataset. Here, we can see that the adversary advantage shows less variation than the seeds dataset. For all epsilon values, kd-Laplace without optimizations stands out. The mechanisms perform similarly, but the Piecewise mechanism performs better for epsilon values between 1.5 and 5.0. For the TPR, all mechanisms are below the baseline value. They score the same, but the Piecewise mechanism is approximately 0.05 TPR lower than the kd-Laplace variants for epsilon values 0.1 to 3.5.

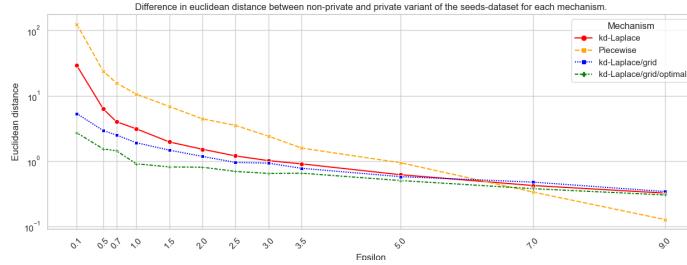


Figure 6.31: Privacy distance for each mechanism for 3D seeds-dataset.

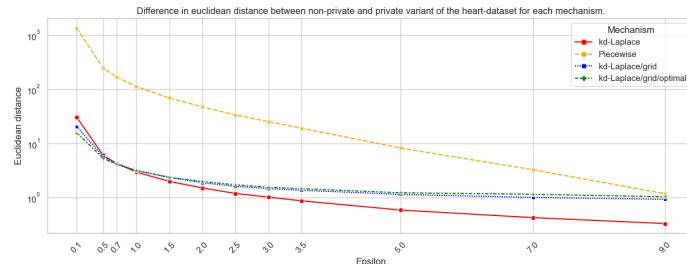


Figure 6.32: Privacy distance for each mechanism for 3D heart-dataset.

At epsilon values 7 and 9, the Piecewise mechanism adds the least distance. Among the various kd-Laplace variants, the variant without optimizations adds the most distance. This trend continues until epsilon 5, after which the variants become equal.

Similarly, a noticeable difference is observed between the Piecewise mechanism and kd-Laplace for the heart dataset. At epsilon 9, the Piecewise mechanism has the same score as kd-Laplace. Among the kd-Laplace variants, the variant without optimizations adds the least distance, but it is slightly higher than the other variants for epsilon 0.1.

### 6.3.3. N-DIMENSIONAL DATA

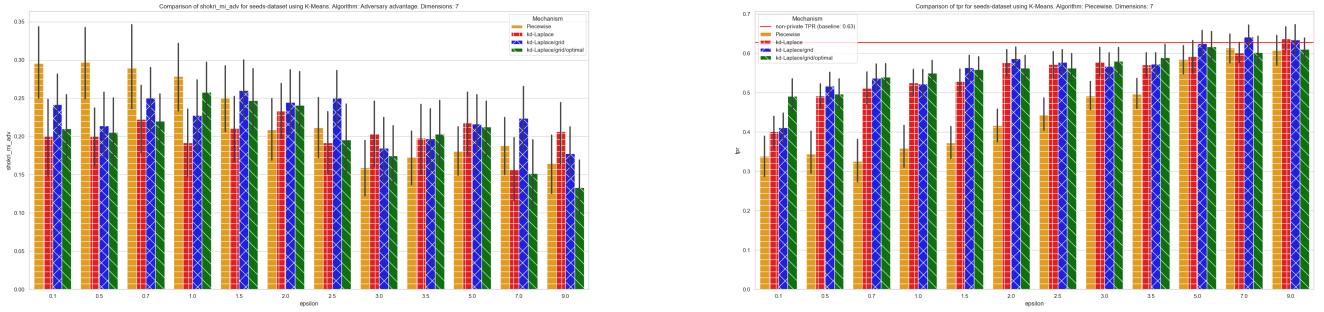


Figure 6.33: Barplot for adversary advantage (left) and TPR (right) per privacy mechanism for seeds-dataset.

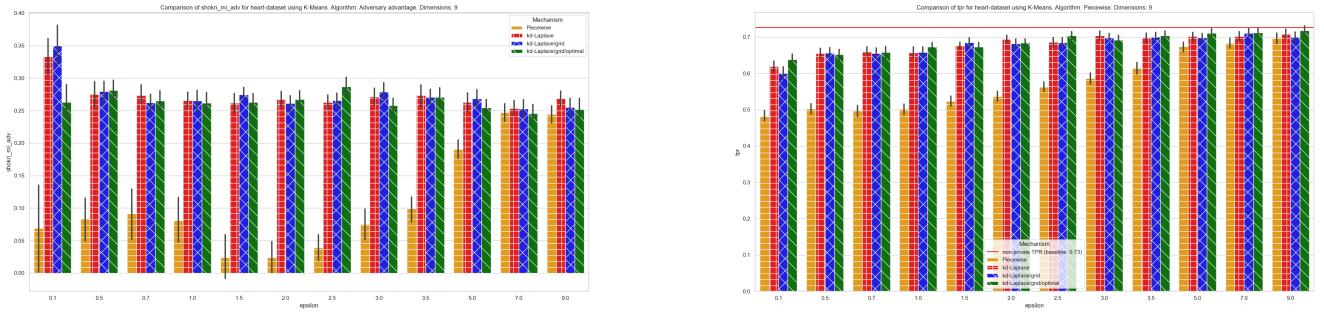


Figure 6.34: Barplot for adversary advantage (left) and TPR (right) per privacy mechanism for heart-dataset.

The seeds dataset shows that Piecewise has a higher adversary advantage for epsilons ranging from 0.1 to 1. However, the TPR for Piecewise is consistently lower than that of kd-Laplace. There is no clear distinction among the kd-Laplace variants. However, for epsilon values 7 and 9, kd-Laplace/grid/optimal does not exceed the baseline for TPR, while the other variants do.

For the heart dataset, the Piecewise mechanism scores below 0.10 for adversary advantage for epsilons 0.1 to 3. In contrast, the kd-Laplace variants yield values above 0.25 for the same epsilon values. The Piecewise mechanism also scores lower for epsilon values 3.5 and 5, but they have nearly equal scores for epsilon values 7 and 9. Among the variants of kd-Laplace, kd-Laplace, and kd-Laplace/grid perform slightly worse for epsilon 0.1, but for the other epsilon values, they function similarly. The TPR follows a similar trend to the adversary advantage, except for epsilon 0.1, where the kd-Laplace variants have equal scores.

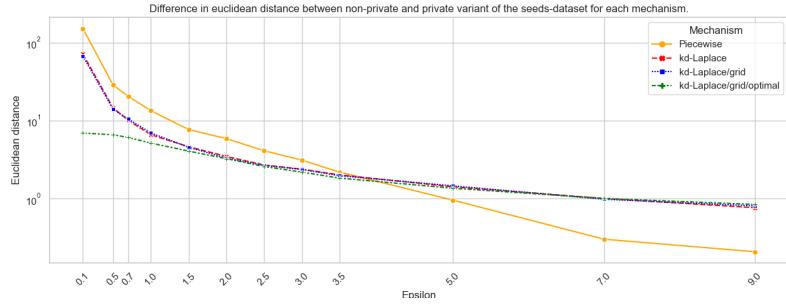


Figure 6.35: Privacy distance for each mechanism for nD seeds-dataset.

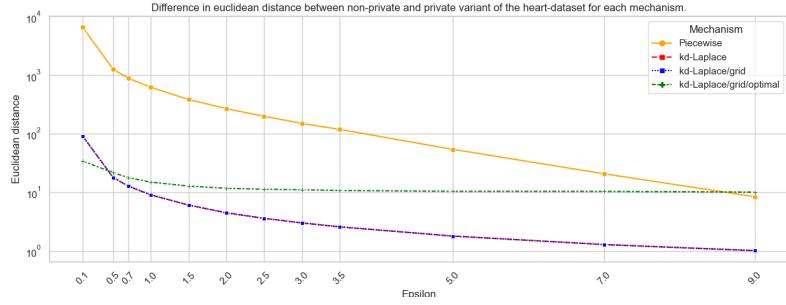


Figure 6.36: Privacy distance for each mechanism for nD heart-dataset.

The Piecewise mechanism for epsilon 0.1 to 3.5 for the seeds dataset adds the most Euclidean distance. After that, the Piecewise mechanism decreases significantly and scores lower than the kd-Laplace variants. For kd-Laplace/grid/optimal, the privacy distance starts lowest up to epsilon 1.5. After that, the variants of the kd-Laplace score are almost the same.

For the heart dataset, the Piecewise mechanism also adds the most Euclidean distance, only now for all epsilons. The kd-Laplace/grid/optimal mechanism starts as the lowest again but is then the highest of all variants, and scores for epsilon 9 are almost equal to the Piecewise mechanism. The other two variants of kd-Laplace (grid / no optimization) score the same.

## 6.4. DIMENSIONALITY

The chart below provides two heat maps for the seeds dataset (top) and the heart dataset (bottom). The y-axis column represents the privacy budget ( $\epsilon$ ), and the x-axis represents the dimensions. In each cell of the matrix, the TPR (True Positive Rate) is indicated, so the darker the cell, the higher the TPR. A higher value implies that, on average, more information is leaked.

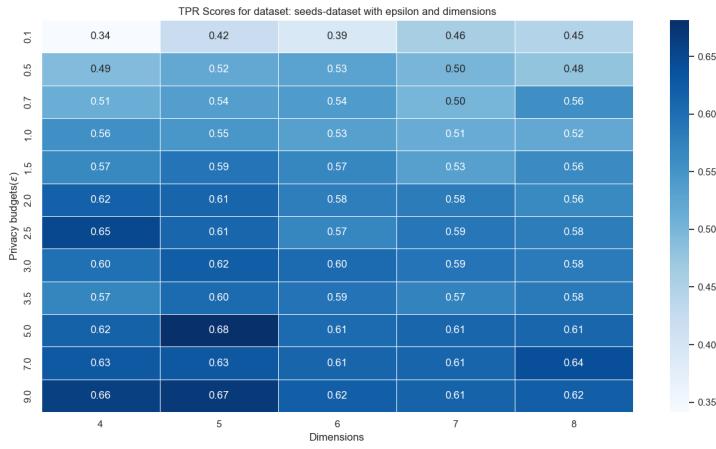


Figure 6.37: Heatmap for TPR and dimensionality for the seeds-dataset for kd-Laplace/grid/optimal

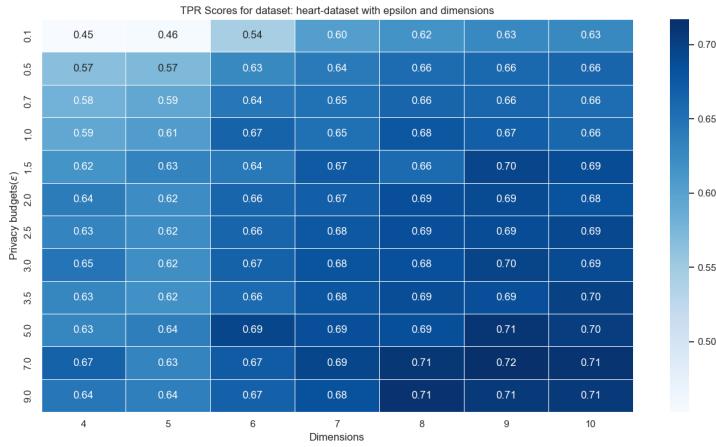


Figure 6.38: Heatmap for TPR and dimensionality for the heart-dataset for kd-Laplace/grid/optimal

Generally, a higher  $\epsilon$  value corresponds to a higher TPR for the seeds dataset. Dimensions 4 and 5 have the highest scores ( $0.60 >$ ) starting from  $\epsilon = 0.5$ . The bottom row achieves the highest scores, and for dimensions 4, 5, and 6, TPR values less than 0.40 are reported for  $\epsilon = 0.1$ . No clear trend is visible for the remaining  $\epsilon$  values based on increasing dimensions.

The heart dataset's lowest scores ( $< 0.50$ ) are also observed for  $\epsilon = 0.1$  and dimensions 4 and 5. From  $\epsilon = 0.2$  onwards, the values increase ( $> 0.50$  TPR) for dimensions 4 and 5. The heatmap becomes darker for dimensions higher than 5, indicating TPR values higher than 0.60. From  $\epsilon = 0.6$  and 8 dimensions, the scores exceed 0.70 TPR.

## 6.5. SHAPE

This chapter examines three datasets with a specific shape: circle, line, and left-skewed. The adversary advantages (privacy) and AMI (utility) are compared between the mechanisms for all three datasets. We compare kd-Laplace/grid/optimal (green) and Piecewise (yellow).

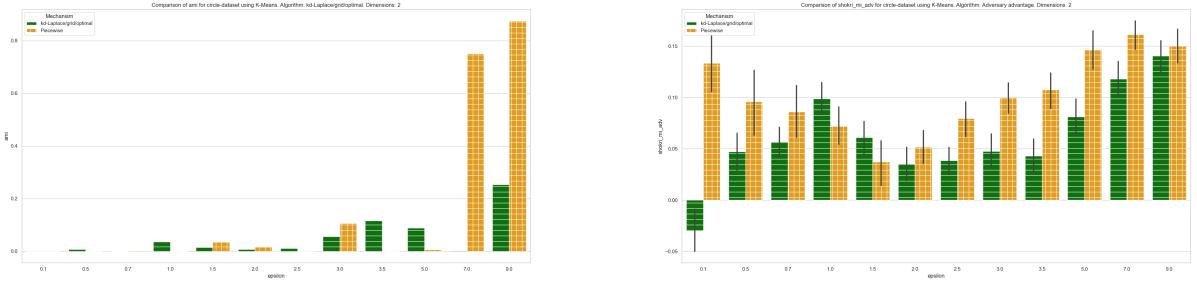


Figure 6.39: The AMI (left) and adversary advantage (right) for the circle-dataset

There's a noticeable difference between the Piecewise and kd-Laplace/grid/optimal mechanisms in the circle dataset. For the AMI, Piecewise scores are significantly higher at epsilon 7 to 9. In comparison, kd-Laplace/grid/optimal scores are lower than 0.2 for most other epsilons. Regarding the adversary advantage, kd-Laplace/grid-optimal scores are lower than Piecewise, except for epsilon 1 and 1.5.

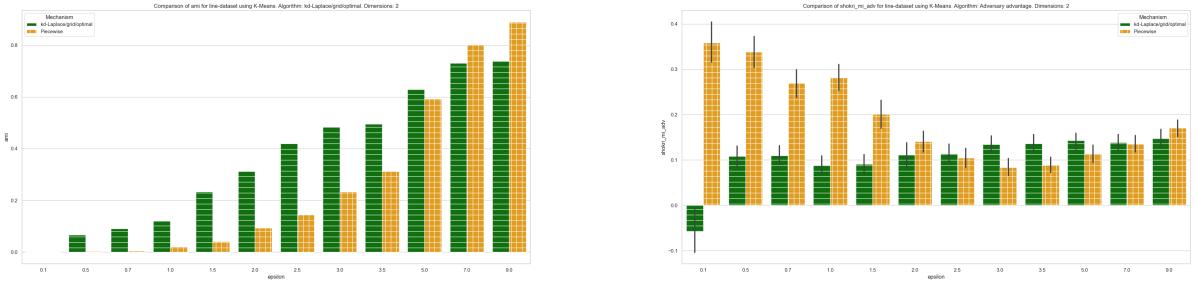


Figure 6.40: The AMI (left) and adversary advantage (right) for the line dataset

In the line dataset, Piecewise outperforms kd-Laplace/grid/optimal for epsilon values above 5 for the AMI metric. For epsilons between 0.1 and 5, kd-Laplace/grid/optimal scores are higher. Regarding adversary advantage, Piecewise performs worse for epsilons between 0.1 and 1.5, while kd-Laplace/grid/optimal scores worse for epsilons 2 to 7.

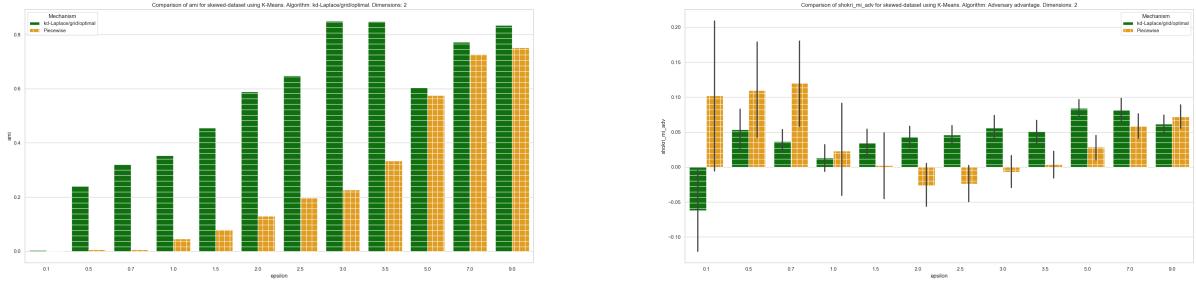


Figure 6.41: The AMI (left) and adversary advantage (right) for the skewed dataset

Kd-Laplace/grid/optimal is better than Piecewise for skewed datasets, across all epsilon values, with AMI scores ranging from 0.6 to 0.8. For adversary advantage, Kd-Laplace/grid/optimal outperforms Piecewise between 0.1 and 1.0 epsilon values. The adversary advantage stays low for both mechanisms (below 0.1).

# 7

## DISCUSSION

In progress

# **CONCLUSION**

In progress

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# HYPERPARAMETERS

## .1. K-MEANS

For selecting the appropriate amount of clusters, we used an "elbow" plot in combination with the silhouette score.

1. Seeds dataset: 4 clusters (see figure: 1)
2. Heart dataset: TODO

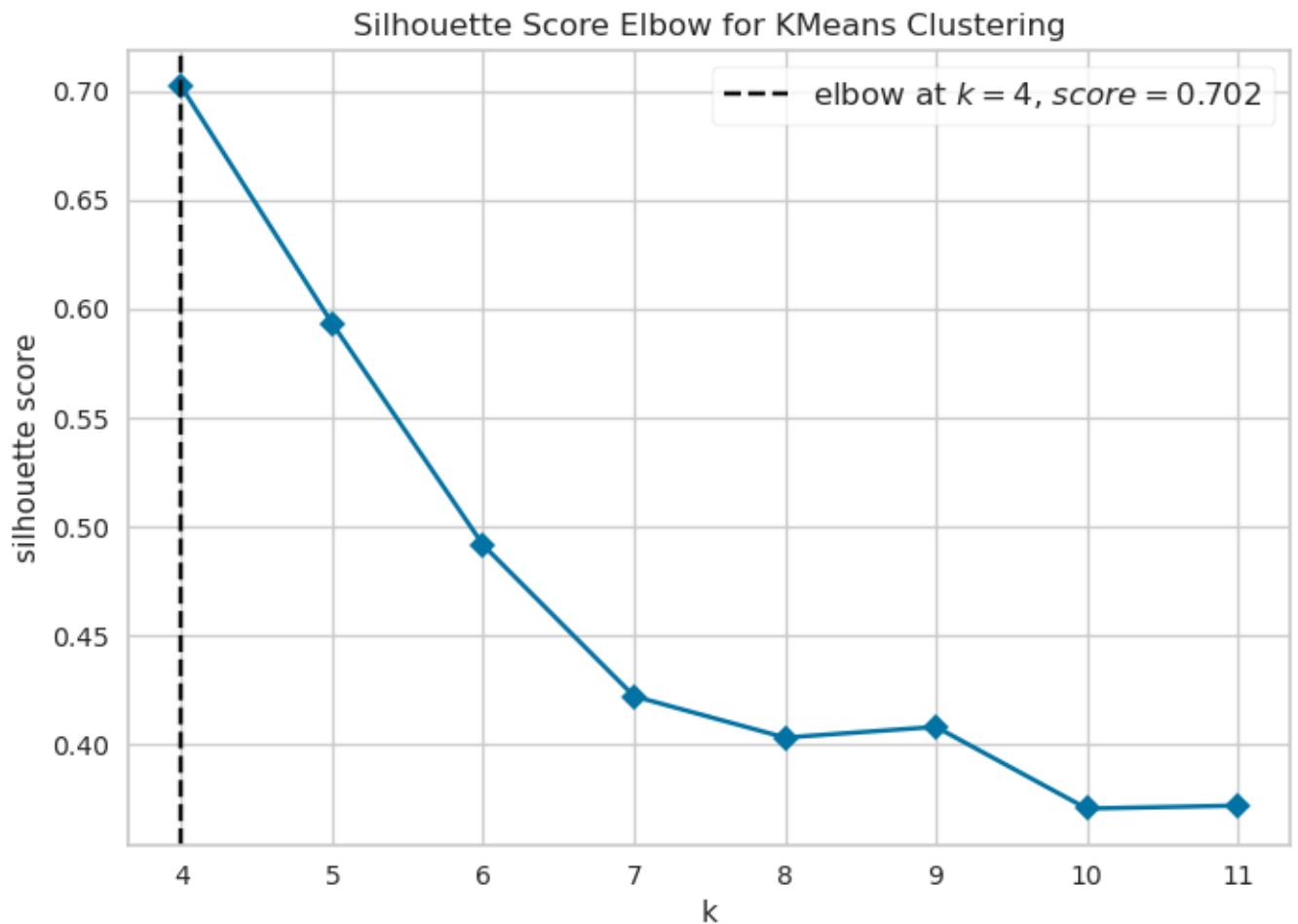


Figure 1: Selecting the  $k$  for K-Means for dataset 1 using the "elbow plot" using section 2.2.1

# THEORY

## .2. BIG O NOTATION

The big O notation is a common way to describe the complexity of an algorithm. It is used to describe the worst-case scenario of an algorithm. For example, if an algorithm has a complexity of  $O(n)$ , it means that the algorithm will take at most  $n$  steps to complete. Below, this is illustrated using a graph.

Figure 2: Graphical representation of the big O notation

# RESULTS

## .3. CLUSTER UTILITY

### .3.1. 2-DIMENSIONAL DATA

Figure 3: Internal validation for all mechanisms the 2-dimensional data heart-dataset

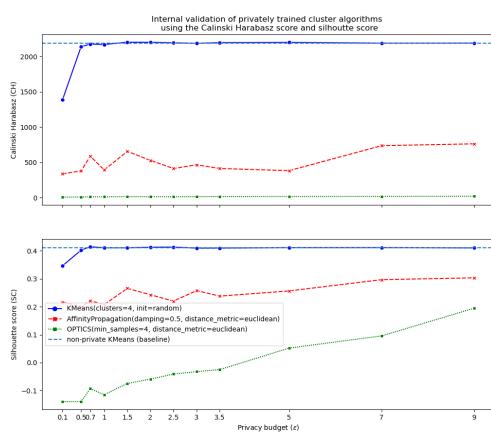


Figure 4: Internal validation (CH/ SC) for the 2-dimensional data heart-dataset for laplace.

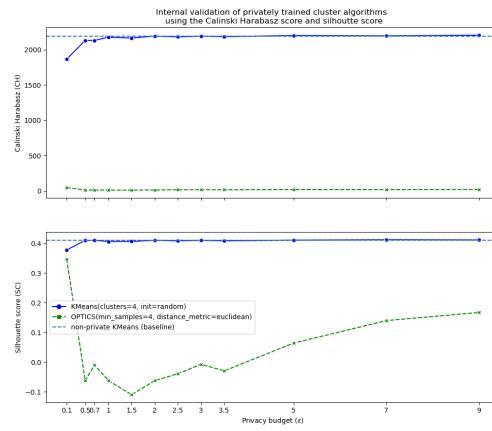


Figure 5: Internal validation (CH/ SC) for the 2-dimensional data heart-dataset for laplace with truncation.

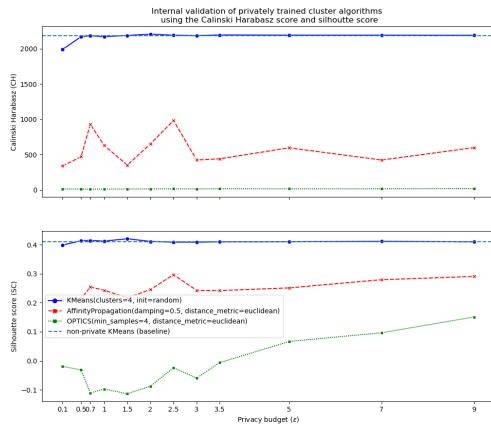


Figure 6: Internal validation (CH/ SC) for the 2-dimensional data heart-dataset for laplace with op-truncation

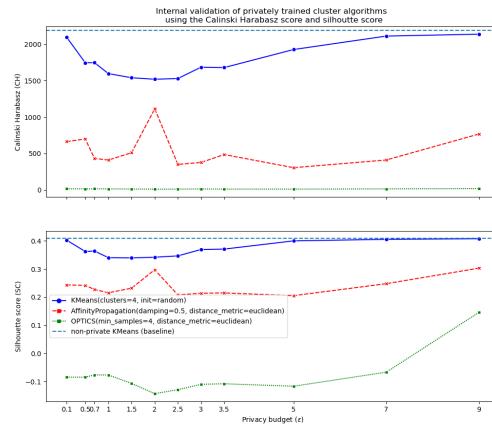


Figure 7: Internal validation (CH/ SC) for the 2-dimensional data heart-dataset for piecewise mechanism

Figure 8: Internal validation for all mechanisms the 2-dimensional data seeds-dataset

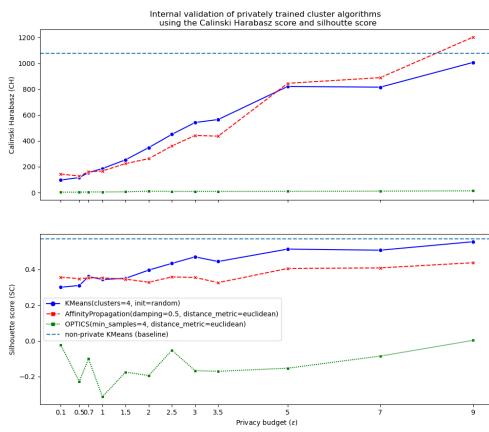


Figure 9: Internal validation (CH/ SC) for the 2-dimensional data seeds-dataset for laplace.

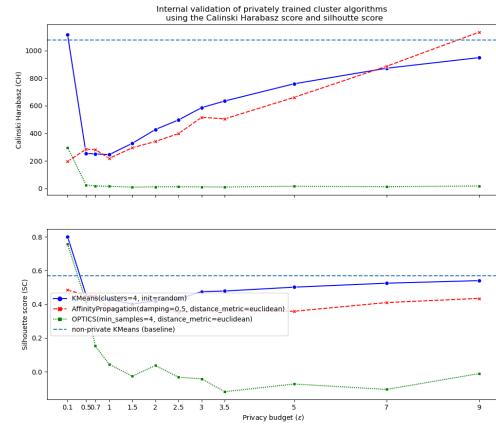


Figure 10: Internal validation (CH/ SC) for the 2-dimensional data seeds-dataset for laplace with truncation.

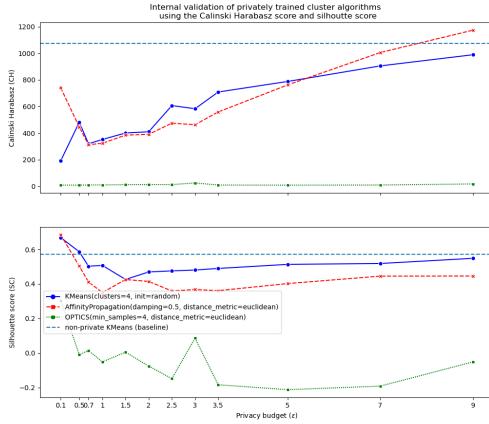


Figure 11: Internal validation (CH/ SC) for the 2-dimensional data seeds-dataset for op-timal truncation

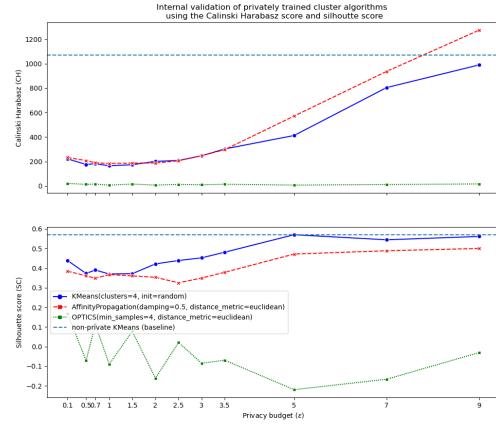


Figure 12: Internal validation (CH/ SC) for the 2-dimensional data seeds-dataset for piecewise mechanism

### 3.2. 3-DIMENSIONAL DATA

Figure 13: Internal validation for all mechanisms the 3-dimensional data seeds-dataset

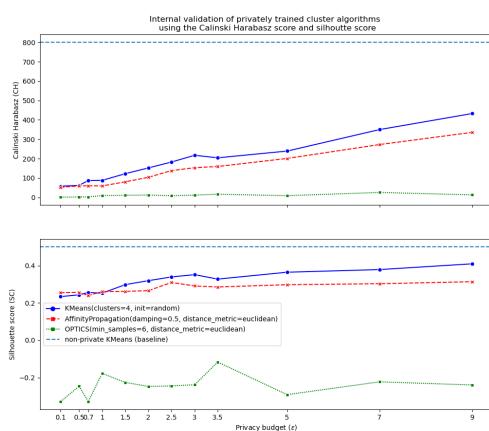


Figure 14: Internal validation (CH/ SC) for the 3-dimensional data seeds-dataset for laplace.

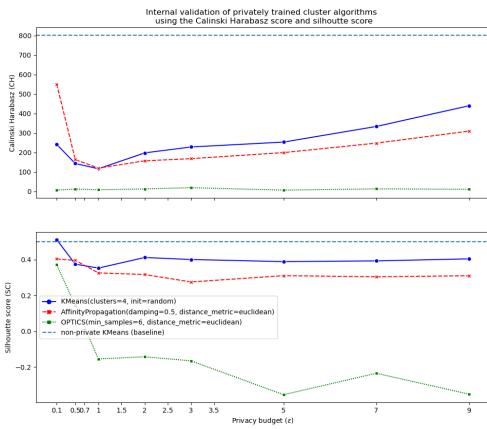


Figure 15: Internal validation (CH/ SC) for the 3-dimensional data seeds-dataset for laplace with truncation.

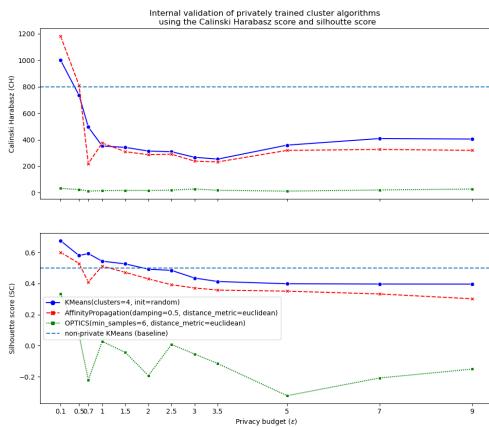


Figure 16: Internal validation (CH/ SC) for the 3-dimensional data seeds-dataset for laplace with op-truncation

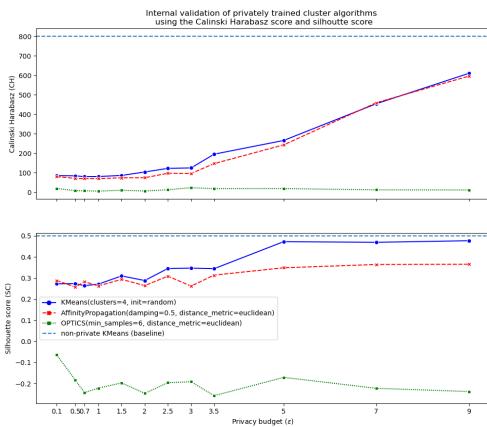


Figure 17: Internal validation (CH/ SC) for the 3-dimensional data seeds-dataset for piecewise mechanism

Figure 18: Internal validation for all mechanisms the 3-dimensional data heart-dataset

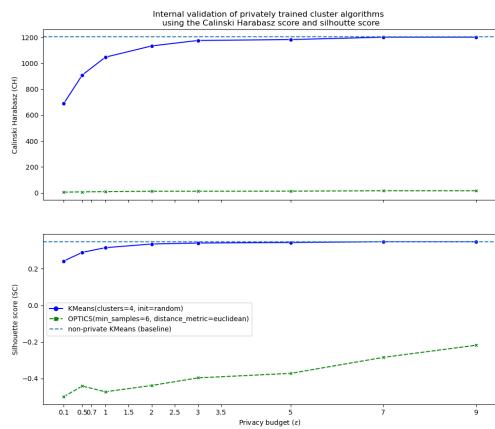


Figure 19: Internal validation (CH/ SC) for the 3-dimensional data heart-dataset for laplace.

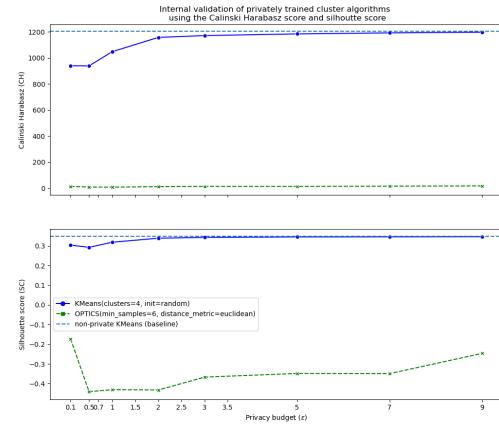


Figure 20: Internal validation (CH/ SC) for the 3-dimensional data heart-dataset for laplace with truncation.

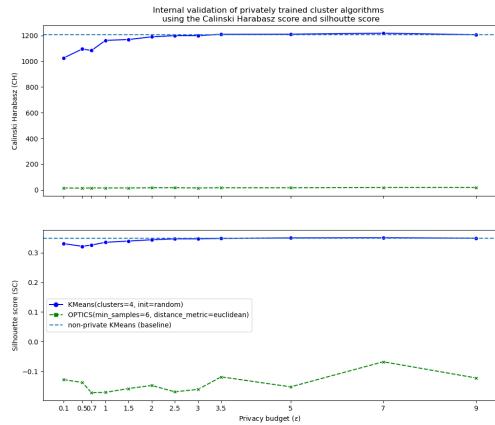


Figure 21: Internal validation (CH/ SC) for the 3-dimensional data heart-dataset for op-timal truncation

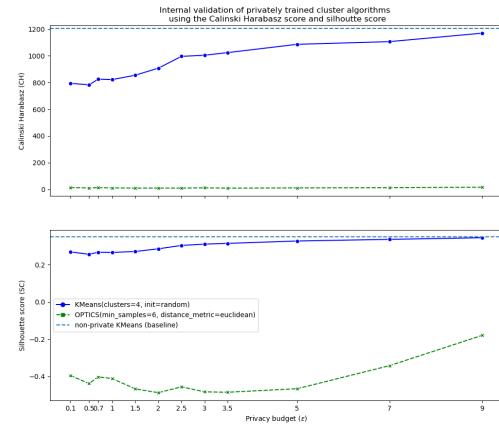


Figure 22: Internal validation (CH/ SC) for the 3-dimensional data heart-dataset for piecewise mechanism

### 3.3. N-DIMENSIONAL DATA

Figure 23: Internal validation for all mechanisms the n-dimensional data seeds-dataset

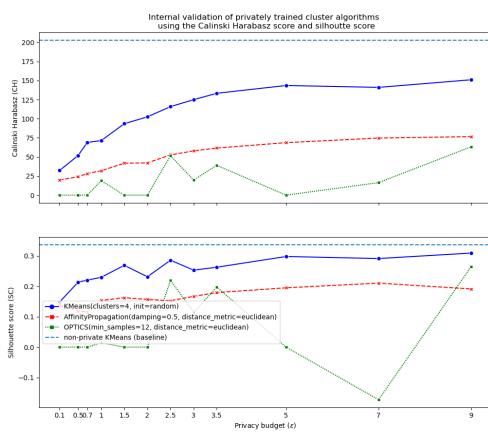


Figure 24: Internal validation (CH/ SC) for the n-dimensional data seeds-dataset for laplace.

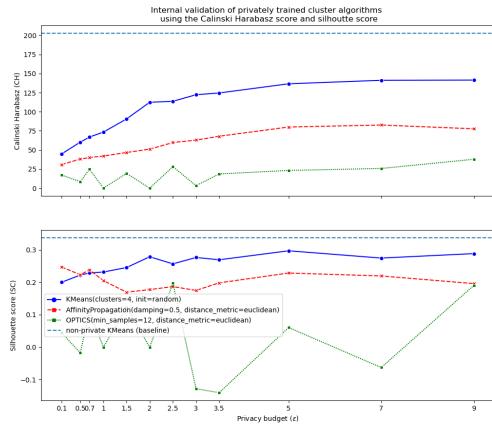


Figure 25: Internal validation (CH/ SC) for the n-dimensional data seeds-dataset for laplace with truncation.

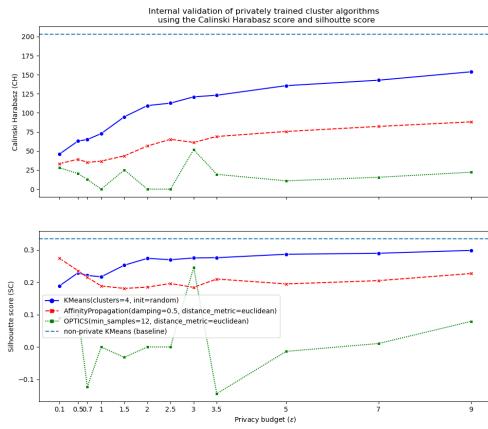


Figure 26: Internal validation (CH/ SC) for the n-dimensional data seeds-dataset for laplace with op-truncation

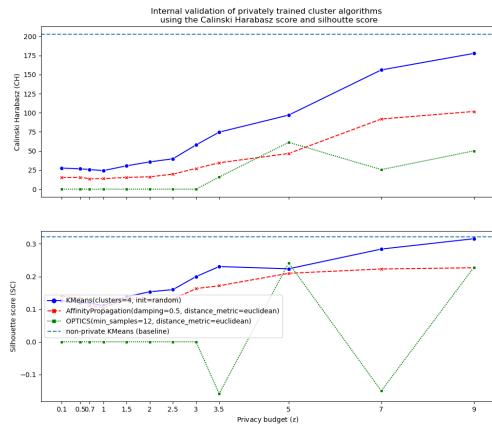


Figure 27: Internal validation (CH/ SC) for the n-dimensional data seeds-dataset for piecewise mechanism

Figure 28: Internal validation for all mechanisms the n-dimensional data heart-dataset

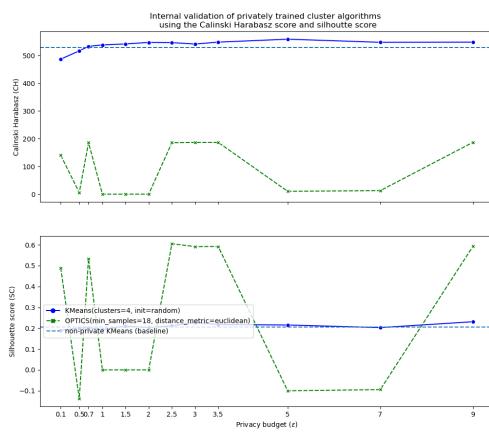


Figure 29: Internal validation (CH/ SC) for the n-dimensional data heart-dataset for laplace.

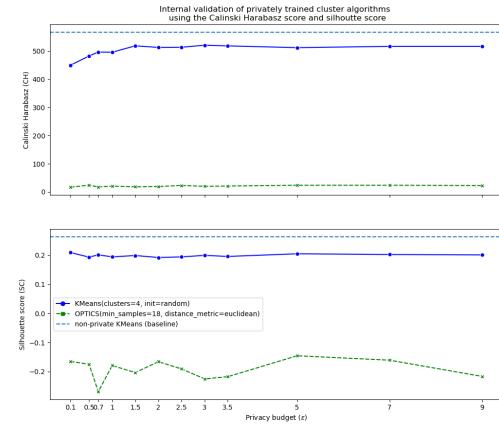


Figure 30: Internal validation (CH/ SC) for the n-dimensional data heart-dataset for laplace with truncation.

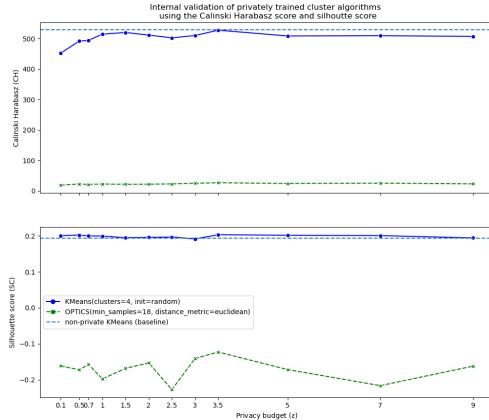


Figure 31: Internal validation (CH/ SC) for the n-dimensional data heart-dataset for laplace with optimal truncation

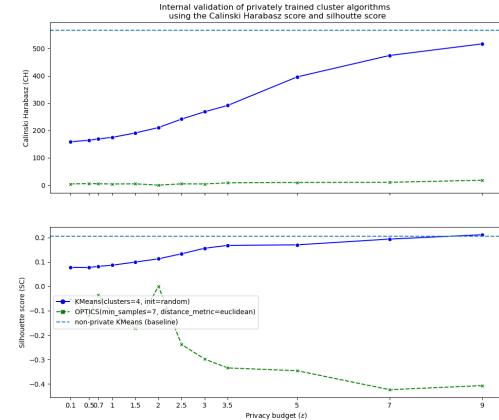


Figure 32: Internal validation (CH/ SC) for the n-dimensional data heart-dataset for piecewise mechanism

## 4. MECHANISM UTILITY

### 4.1. 2-DIMENSIONAL DATA

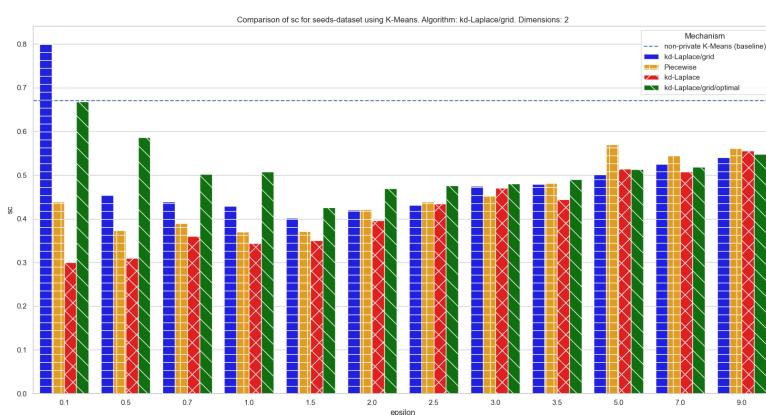


Figure 33: Silhouette score comparison for the 2D seeds-dataset

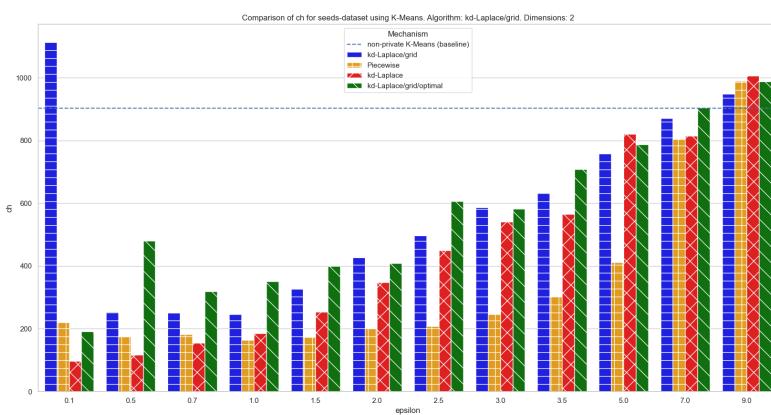


Figure 34: Calinski Harabasz score comparison for the 2D seeds-dataset

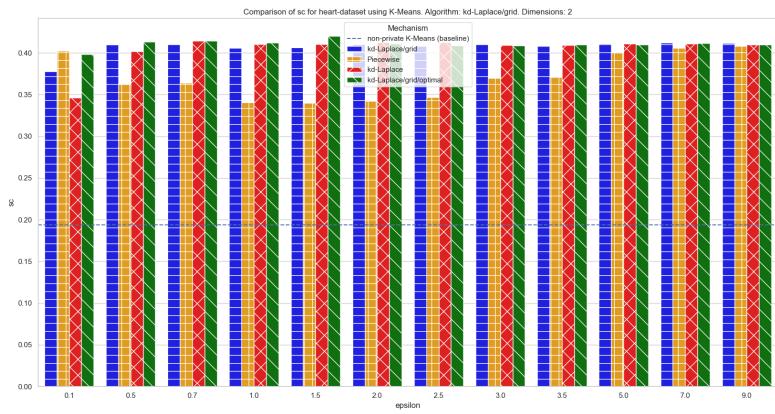


Figure 35: Silhouette score comparison for the 2D heart-dataset

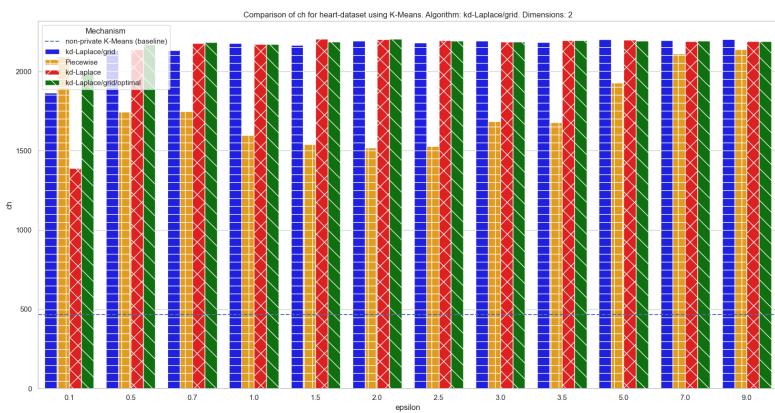


Figure 36: Calinski Harabasz score comparison for the 2D heart-dataset

## 4.2. 3-DIMENSIONAL DATA

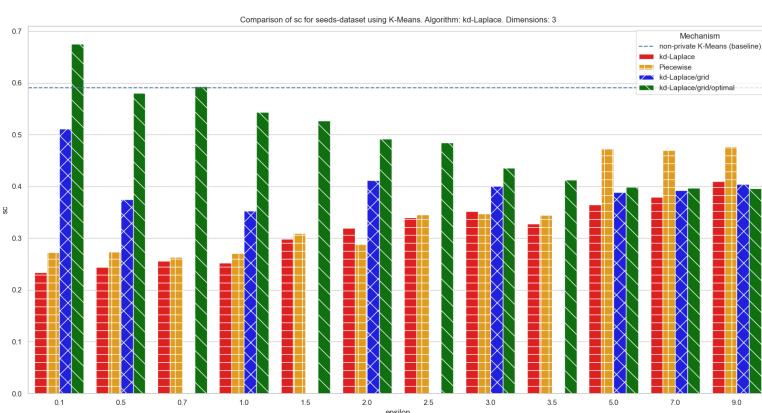


Figure 37: Silhouette score comparison for the 3D seeds-dataset

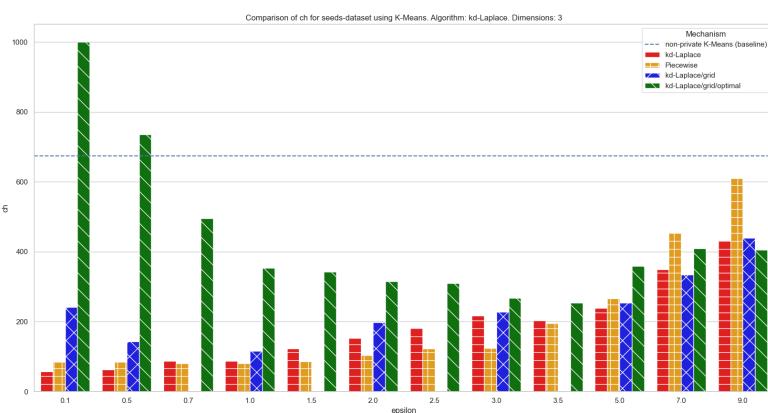


Figure 38: Calinski Harabasz score comparison for the 3D seeds-dataset

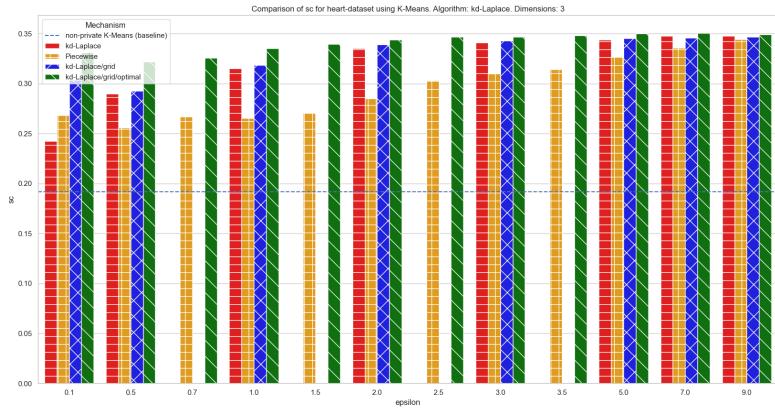


Figure 39: Silhouette score comparison for the 3D heart-dataset

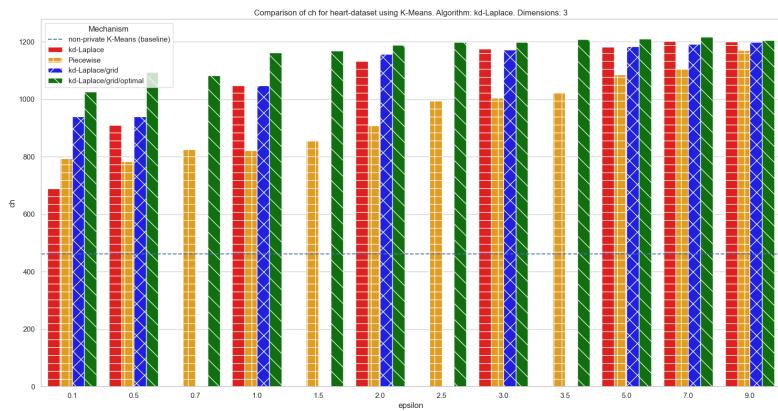


Figure 40: Calinski Harabasz score comparison for the 3D heart-dataset

### 4.3. N-DIMENSIONAL DATA

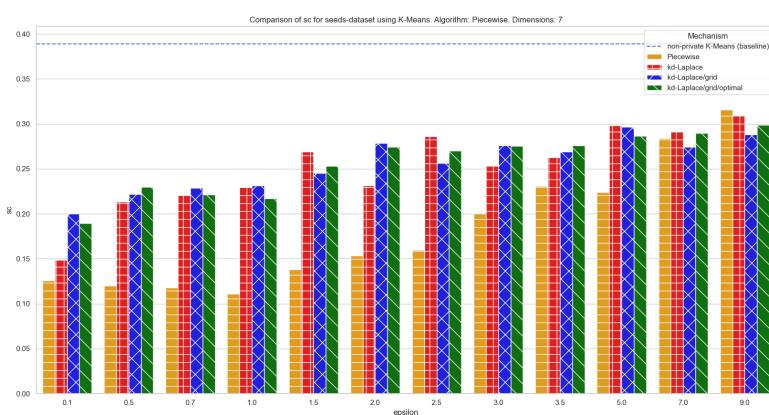


Figure 41: Silhouette score comparison for the nd seeds-dataset

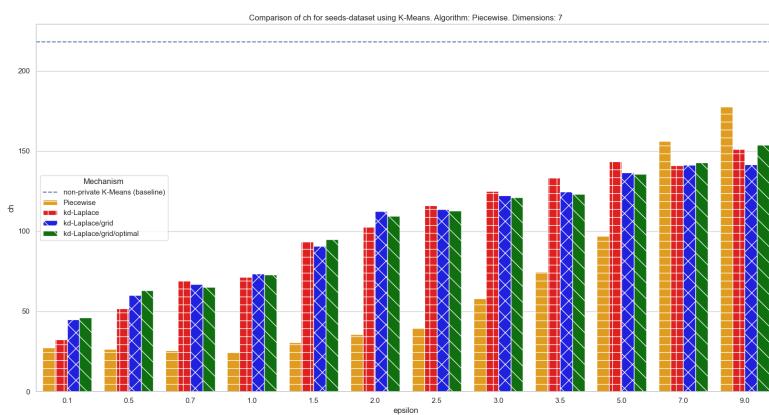


Figure 42: Calinski Harabasz score comparison for the nd seeds-dataset

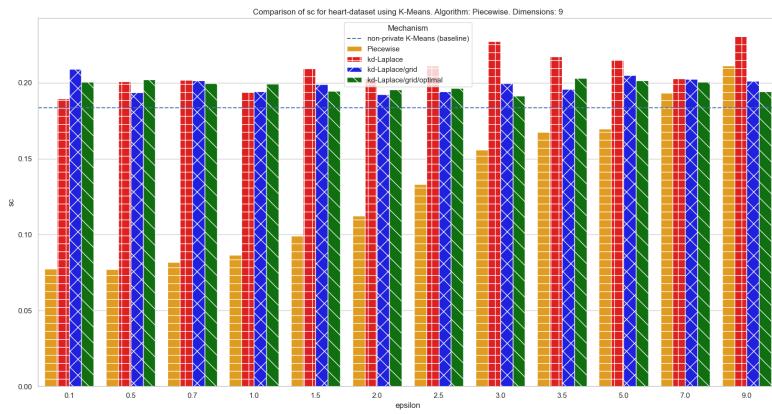


Figure 43: Silhouette score comparison for the nd heart-dataset

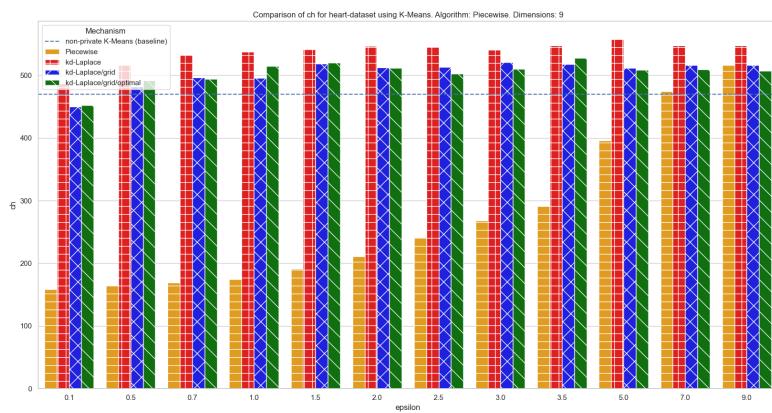


Figure 44: Calinski Harabasz score comparison for the nd heart-dataset

## .5. PRIVACY

### .5.1. 2-DIMENSIONAL DATA

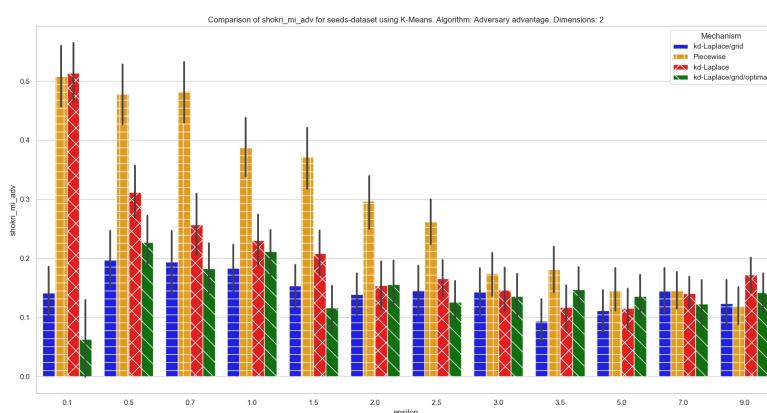


Figure 45: Adversary advantage comparison for the 2D seeds-dataset

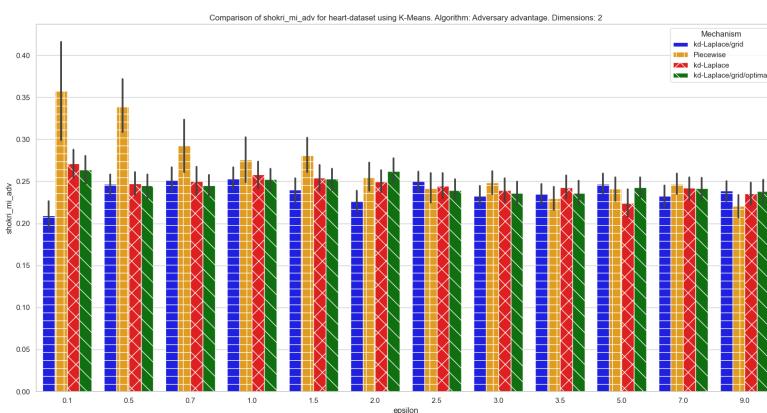


Figure 46: Adversary advantage comparison for the 2D heart-dataset

## .5.2. 3-DIMENSIONAL DATA

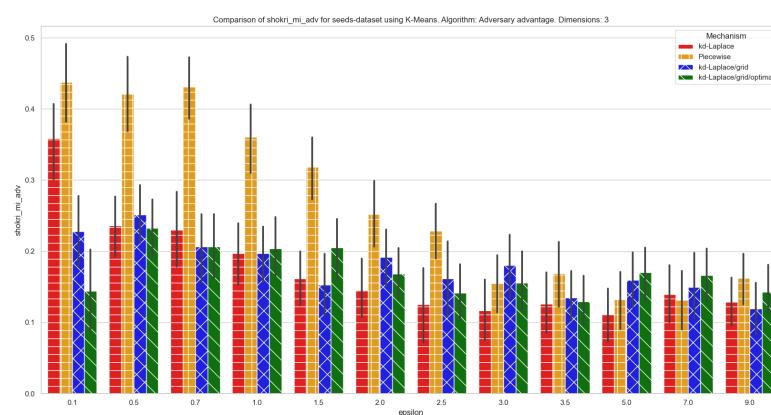


Figure 47: Adversary advantage comparison for the 3D seeds-dataset

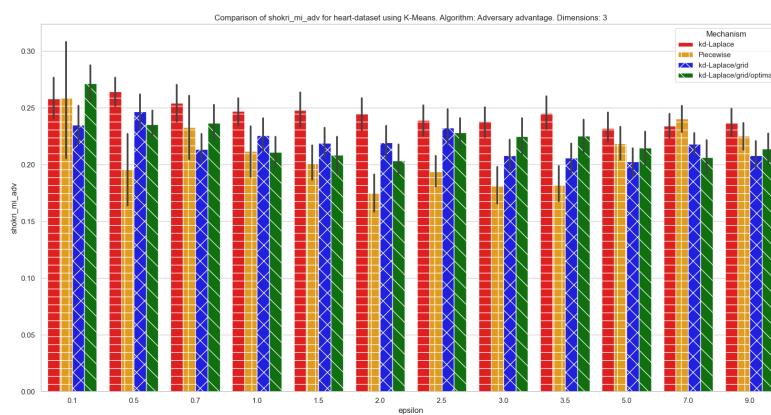


Figure 48: Adversary advantage comparison for the 3D heart-dataset

### 5.3. N-DIMENSIONAL DATA

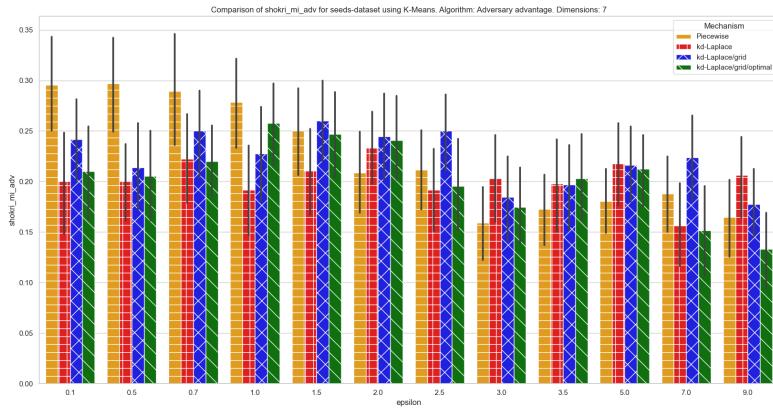


Figure 49: Adversary advantage comparison for the nD seeds-dataset

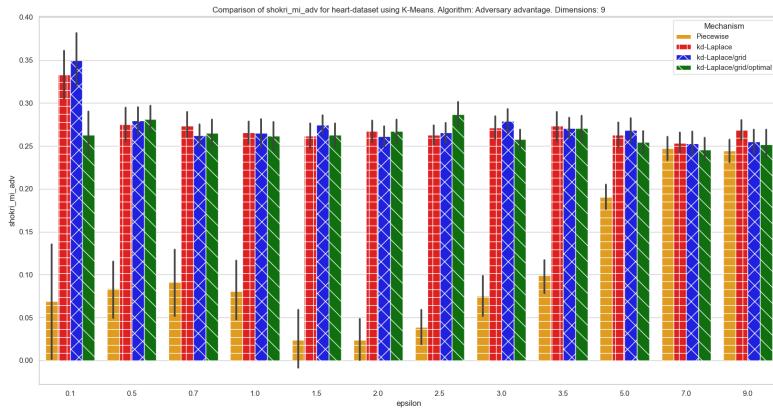


Figure 50: Adversary advantage comparison for the nD heart-dataset