



Summary

Alice has her pair of keys ($sk_A + pk_A$).

Message d is encrypted with a symmetric key DEK.

$$D = \text{encrypt_data}(d, \text{DEK})$$

Message D is an encrypted version of message d with the DEK key.

Now, it's time for Alice to encrypt the DEK key, and put it into a capsule only she can open while decrypting it with her secret key sk_A .

$$(c_A, k_A) = \text{encrypt}(\text{DEK}, pk_A)$$

There is a ciphertext c_A and a capsule k_A . The capsule can only be opened with a secret key sk_A which is a counterpart to pk_A .

$$\text{DEK} = \text{decrypt}(c_A, k_A, sk_A)$$

In order to enable Bob to open the capsule, Alice has to provide a re-encryption key $\Gamma_{A \rightarrow B}$, such that allow to transform (c_A, k_A) into (c_B, k_B) .

$$\Gamma_{A \rightarrow B} = \text{generate_reencryption_key}(sk_A, pk_B)$$

$$(c_B, k_B) = \text{reencrypt}(c_A, k_A, \Gamma_{A \rightarrow B})$$

Bob would now be able to retrieve DEK.

$$\text{DEK} = \text{decrypt}(c_B, k_B, sk_B)$$

To improve the level of security, Alice needs to split the reencryption key $\Gamma_{A \rightarrow B}$ and send parts to multiple locations. This process makes it possible to reduce trust required from a re-encryption proxy. If it doesn't know the whole $\Gamma_{A \rightarrow B}$, then it has no incentive to misbehave.

Instead of generating a single $r_{A \rightarrow B}$, we will have it split into multiple fragments, called $k\text{frags}$.

$$k\text{frags}_{A \rightarrow B} = \text{split_reencryption_key}(sk_A, pk_B, M, N)$$

$$\text{where } k\text{frags}_{A \rightarrow B} = \sum_{i=1}^N k\text{frag}_{A \rightarrow B, i}$$

Once Alice has $k\text{frags}_{A \rightarrow B}$ generated, she needs to publish it to the network of proxies. These proxies will later use partial reencryption keys $k\text{frags}_{A \rightarrow B}$ to transform c_A into $c\text{frags}_{A \rightarrow B}$.

$$c\text{frags}_{A \rightarrow B} = \sum_{i=1}^N \text{reencrypt_frag}(c_A, k_A, k\text{frag}_{A \rightarrow B, i})$$

Bob will need exactly M of all distinct $c\text{frags}_{A \rightarrow B}$ available in order to reveal DEK. The order of $c\text{frags}$ doesn't matter.

$$(c_B, k_B) = \text{merge}(c_A, k_A, c\text{frags}_{A \rightarrow B}^*)$$

And now, finally, Bob can access DEK

$$\text{DEK} = \text{decrypt}(c_B, k_B, sk_B)$$

Following DEK access, it's time to see what is hidden under the ciphertext D .

$$d = \text{decrypt_data}(D, \text{DEK})$$

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