

Info retrieval sheet 4

Exercise 1

To show: for any prefix free encoding:

$$\mathbf{E} L_x \geq H(X) = \sum_i p_i * \log(1/p_i)$$

Given formula:

$$\sum_i 2^{-L_i} \leq 1$$

We need to build the lagrange function:

$$\rightarrow L(p_1, \dots, p_k, L_1, \dots, L_k, \lambda) = \sum_i p_i * L_i + \lambda * (\sum_i 2^{-L_i} - 1)$$

$$dL/dL_i = p_i - \lambda * \ln(2) * 2^{-L_i} \stackrel{!}{=} 0$$

$$\leftrightarrow 2^{L_i} = p_i / (\lambda * \ln(2))$$

$$\leftrightarrow -L_i = \log_2(p_i) - \log_2(\lambda * \ln(2))$$

Differentiate after lambda :

$$dL/d\lambda = \sum_i 2^{-\log_2(p_i / (\lambda * \ln(2)))} - 1 \stackrel{!}{=} 0$$

$$\leftrightarrow \sum_i p_i / (\lambda * \ln(2)) = 1$$

$$\rightarrow (\lambda * \ln(2)) = 1, \text{ because } \sum_i p_i = 1$$

$$\rightarrow L_i = \log_2(1/p_i) + \log_2(1) = \log_2(p_i)$$

$$\rightarrow \mathbf{E} L_x = \sum_i p_i * \log_2(1/p_i) = H(X)$$

Now we know that this is an extreme point, we only need to show that it is a minimum:

$$d^2L/dL_i^2 = \lambda * \ln(2)^2 * 2^{-L_i} > 0$$

→ It is a minimum