

# Data Structure and Algorithm HW1

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## 1. What if you became a DSA TA?

### Proof of conjugate 1

Consider  $f(n) \geq 0$  and  $g(n) \geq 0$  that  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = c$ , where  $c \in R$   
By definition of limit, there exists  $\epsilon > 0$ ,  $n_0 > 0$  such that for all  $n > n_0$ ,

$$\left| \frac{f(n)}{g(n)} - c \right| < \epsilon$$

$$-\epsilon < \frac{f(n)}{g(n)} - c < \epsilon$$

$$-\epsilon + c < \frac{f(n)}{g(n)} < \epsilon + c$$

In case that  $c < \epsilon$

$$-\epsilon + c < 0 < \frac{f(n)}{g(n)} < \epsilon + c$$

In case that  $c \geq \epsilon$

$$0 \leq -\epsilon + c < \frac{f(n)}{g(n)} < \epsilon + c$$

$$0 \leq f(n) < (\epsilon + c)g(n)$$

Let  $(\epsilon + c) = c_1$

$$0 \leq f(n) \leq c_1 g(n)$$

which shows  $f(n) = O(g(n))$ , QED

### 1.1

Consider  $c_1, c_2 > 0$

for all  $n > 0$

if

$$\frac{c_1}{c_2} n \geq 1$$

then

$$c_1 n^3 \geq c_2 n^2$$

let

$$n_0 = \frac{c_2}{c_1}$$

$$\frac{c_1}{c_2}n_0 = 1$$

consider  $k > 0$

$$\frac{c_1}{c_2}(n_0 + k) = 1 + \frac{c_1}{c_2}k$$

where

$$\frac{c_1}{c_2}k > 0$$

that is, for all  $n \geq n_0$ ,

$$0 \leq c_2n^2 \leq c_1n^3$$

therefore, if

$$0 \leq f(x) \leq c_3n^2$$

for all  $n > n_1$  where  $n_1, c_3 > 0$   
then

$$0 \leq f(n) \leq c_3n^2 \leq \frac{c_3c_1}{c_2}n^3$$

for all  $n > \max(n_0, n_1)$ , QED

## 1-2

Let  $t_m = \text{the number of while checks}$ ,  $k = \text{the index of key}$ ,  $d_n = \text{time cost of each line to execute once}$   
Since the value of  $m$  is assigned to  $l$  in the beginning of each while loop,  $m$  will start from 1 and increase by 1 when while loop is executed once, moreover,  $A[m] \leq \text{key}$  will always be satisfied.  
therefore,

$$t_m = k$$

in worst case, key is not in the array, in this case

$$T(n) = d_1n + d_2(n-1) + d_3(n-1) + d_4 * 0 + d_5(n-1) + d_6 * 0 + d_7(n-1) + d_8(n-1) + d_9$$

$$T(n) = c_1n + c_2$$

where  $c_1, c_2 > 0$

$$\lim_{n \rightarrow \infty} \frac{c_1n + c_2}{n} = c_1$$

By conjecture 1, time complexity of the algorithm is  $O(n)$

## 1-3

$$f(n) = \Theta(n^2) \iff \text{there exists positive } (n_0, c_1, c_2) \text{ such that } c_1n^2 \leq f(n) \leq c_2n^2 \text{ for all } n \geq n_0$$