Data Structure and Algorithm HW1

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1. What if you became a DSA TA?

Proof of conjecture 1

Consider $f(n) \geq 0$ and $g(n) \geq 0$ that $\lim_{n\to\infty} \frac{f(n)}{g(n)} = c$, where $c \in R$
To prove $\lim_{n\to\infty} \frac{f(n)}{g(n)} \geq 0$ for all $n \in R$, we can consider h(n) = 0
Since $f(n) \geq 0$ and $g(n) \geq 0$ for all $n \in R$, $\frac{f(n)}{g(n)} \geq h(n)$ for all n
therefore,

$$\lim_{n\to\infty} h(n) \geq \lim_{n\to\infty} \frac{f(n)}{g(n)}$$

$$\lim_{n\to\infty} h(n) = 0 \Rightarrow \lim_{n\to\infty} \frac{f(n)}{g(n)} = c \geq 0$$

By difinition of limit, there exists $\epsilon > 0$, $n_0 > 0$ such that for all $n > n_0$,

$$|\frac{f(n)}{g(n)} - c| < \epsilon$$

$$-\epsilon < \frac{f(n)}{g(n)} - c < \epsilon$$

$$-\epsilon + c < \frac{f(n)}{g(n)} < \epsilon + c$$

In case that $0 < c < \epsilon$

$$-\epsilon + c < 0 < \frac{f(n)}{g(n)} < \epsilon + c$$

In case that $c \geq \epsilon$

$$0 \leq -\epsilon + c < \frac{f(n)}{g(n)} < \epsilon + c$$

$$0 \leq f(n) < (\epsilon + c)g(n)$$

Let $(\epsilon + c) = c_1$

$$0 \leq f(n) \leq c_1 g(n)$$

which shows $f(n) = O(g(n))$, QED

1.1 - all by myself

Consider $c_1, c_2 > 0$ for all $n > 0$
if

$$\frac{c_1}{c_2} n \geq 1$$

then

$$c_1 n^3 \geq c_2 n^2$$

let

$$n_0 = \frac{c_2}{c_1}$$

$$\frac{c_1}{c_2} n_0 = 1$$

consider $k > 0$

$$\frac{c_1}{c_2}(n_0 + k) = 1 + \frac{c_1}{c_2} k$$

where

$$\frac{c_1}{c_2} k > 0$$

that is, for all $n \geq n_0$,

$$0 \leq c_2 n^2 \leq c_1 n^3$$

therefore, if

$$0 \leq f(x) \leq c_3 n^2$$

for all $n > n_1$ where $n_1, c_3 > 0$
then

$$0 \leq f(n) \leq c_3 n^2 \leq \frac{c_3 c_1}{c_2} n^3$$

for all $n > max(n_0, n_1)$, QED

1-2 - all by myself

Let $t_m = the\ number\ of\ while\ checks$, $k = the\ index\ of\ key$, $d_n = timecostofeachlinetoexecuteonce$
Since the value of m is assigned to l in the beginning of each while loop, m will start from 1 and increase by 1 when while loop is executed once, moreover, $A[m] \leq key$ will always be satisfied.
therefore,

$$t_m = k$$

in worst case, key is not in the array, in this case

$$T(n) = d_1 n + d_2(n-1) + d_3(n-1) + d_4 * 0 + d_5(n-1) + d_6 * 0 + d_7(n-1) + d_8(n-1) + d_9$$

$$T(n) = c_1 n + c_2$$

where c_1, $c_2 > 0$

$$\lim_{n \to \infty} \frac{c_1 n + c_2}{n} = c_1$$

By conjecture 1, time complexity of the algorithm is $O(n)$

1-3

ref: `https://math.stackexchange.com/questions/925053/using-limits-to-determine-big-o-big-omega-and-big-theta#comment6149810_925053`

$$f(n) = \Theta(n^2) \iff there\ exists\ positive\ (n_0, c_1, c_2)\ such\ that\ c_1 n^2 \leq f(n) \leq c_2 n^2\ for\ all\ n \geq n_0$$

Let $c_1 = 1$, $c_2 = 10$ and $f(n) = n^2(sin(n) + 2) + n$
for all $n > 0$, $n^2 \leq f(n) \leq 10n^2$
Suppose

$$\lim_{n\to\infty} f(n) = L$$

where $L \in R$, then there exist $\epsilon > 0$, $n > 0$ such that for all $n > n_0$,

$$|f(n) - L| < \epsilon$$

Let $f(n_1) = \epsilon_1$ where $n_1 > n_0$ and $\epsilon_1 < \epsilon$

$$f(n_1 + 2\pi k) - f(n_1) = ((n_1 + 2\pi k)^2 - n_1^2)(sin(n_1) + 2) + 2\pi k$$

where $k \in n$

$$\lim_{k\to\infty} f(n_1 + 2\pi k) - f(n_1) \neq 0$$
$$\Rightarrow |f(n + 2\pi k) - L| > \epsilon$$

for some $k > k_0$ where $k_0 \in n$
That is, $\lim_{n\to\infty} \frac{f(n)}{n^2}$does not exists, therefore, the proposition is WRONG.

1-4 - all by myself

$$lg(n) = 2lg(\sqrt{n}) = 2\frac{ln(\sqrt{n})}{\ln(2)}$$

Consider

$$\lim_{n\to\infty} \frac{lg(n)}{\sqrt{n}} = \lim_{n\to\infty} \frac{2}{ln(2)} \frac{ln(\sqrt{n})}{\sqrt{n}}$$

Let $\sqrt{n} = x$

$$\lim_{n\to\infty} \frac{2}{ln(2)} \frac{ln(\sqrt{n})}{\sqrt{n}} = \lim_{x\to\infty} \frac{2}{ln(2)} \frac{ln(x)}{x}$$

By L'Hopital's rule

$$\lim_{x\to\infty} \frac{2}{ln(2)} \frac{ln(x)}{x} = \lim_{x\to\infty} \frac{2}{ln(2)} \frac{\frac{d}{dx} ln(x)}{\frac{d}{dx} x} = \lim_{x\to\infty} \frac{2}{ln(2)} \frac{\frac{1}{x}}{1} = 0 \in R$$

By conjecture 1, $lg(n) = O(\sqrt{n})$, QED

1-5 - all by myself

Consider

$$\frac{\lim_{n\to\infty}\Sigma_{i=1}^{n} i^n}{\lim_{n\to\infty}\Sigma_{i=1}^{n} n^n} = \lim_{n\to\infty}\Sigma_{i=1}^{n}(\frac{i}{n})^n$$

and

$$\lim_{n\to\infty}\int_1^n (\frac{i}{n})^n di$$

and

$$\lim_{n\to\infty}\int_0^{n-1} (\frac{i}{n})^n di$$

By definition of Σ,

$$\lim_{n\to\infty}\int_0^{n-1} (\frac{i}{n})^n di < \lim_{n\to\infty}\Sigma_{i=1}^{n}(\frac{i}{n})^n < \lim_{n\to\infty}\int_1^n (\frac{i}{n})^n di$$

$$\lim_{n\to\infty}\int_1^n (\frac{i}{n})^n di = \lim_{n\to\infty}\frac{1}{n^n}\int_1^n i^n di = \lim_{n\to\infty}\frac{1}{n^n}\ [i^n]_1^n = \lim_{n\to\infty}\frac{n^n}{n^n} - \lim_{n\to\infty}\frac{1^n}{n^n} = 1-0=1$$

$$\lim_{n\to\infty}\int_0^{n-1} (\frac{i}{n})^n di = \lim_{n\to\infty}\frac{1}{n^n}\int_0^{n-1} i^n di = \lim_{n\to\infty}\frac{1}{n^n}\ [i^n]_0^{n-1} = \lim_{n\to\infty}\frac{(n-1)^n}{n^n} - \lim_{n\to\infty}\frac{1^n}{n^n} = 1-0=1$$

By squeeze theorem,

$$\lim_{n\to\infty}\Sigma_{i=1}^{n}(\frac{i}{n})^n = 1 \in R$$

By conjecture 1, $\Sigma_{i=1}^{n} i^n = O(n^n)$

1-6 - all by myself

In second last line, since c can be positive or negative, it is possible that $\frac{1}{2^c} > 2^c$ and therefore $f(n) > \frac{1}{2^c}$
The correct proof:

$$|lg(f(n)) - lg(g(n))| = c \Rightarrow lg(\frac{f(n)}{g(n)}) = c \Rightarrow \frac{f(n)}{g(n)} = 2^c$$

$$\Rightarrow |lg(f(n)) - lg(g(n))| = c \Rightarrow lg(\frac{f(n)}{g(n)}) = c_1 \Rightarrow \frac{f(n)}{g(n)} = 2^{c_1}$$

where $c_1 = \pm c$
Take $c' = 2^c - 1$, we have $f(n) \le c'g(n)$ for all $n > n_0$, QED

2. DSA Judge

2-1

ref: chatGPT psuedo code: ID-search(A, key, l, r)

$$while(r - l > 1)$$

$$m = (r - l)/2$$

procedure BinarySort(A, key, l, r)
    $m = (l + r)/2$
    while $(r - l) > 1$ **do**
        if $m > key$ **then**
            $l = m$
        else if $m < key$ **then**
            $sdafsd$
        else
            return m
        end if
        $i \leftarrow i - 1$
    end while
    $A[i + 1] \leftarrow key$
end procedure