- 1/2 dimension, all elements have some type (homogeneous), are ordered (start at a[o] to a[n-1])
-1 dim = vector - 2 dim = matrix (array of vectors) -> 2 dim = tensor (arrays of arrays...) -a(-z) -> makes sense w/ pointers but don't use this - Declaring array. (type) name [count] = { Initalize list} - don't need count of you initiatize list & vice versa - Matrix: (type) m[3][3] = D elements m[0][0] to [2][2] - C stores memory row by now - Memory = 1 dim, rous are sequential - arrays are large, don't copy on to stack -> relate to pointers - name of array = pointer to element 0 of array (base address)

- size of operator tells variable's # of bytes
-gizeof (array) = size of array
$-\frac{9i}{2}eof(array) = 5ize of array$ $-\frac{9i}{2}eof(pointer) = 5ize of pointer$ $-\alpha[i] == *(a+i)$
-a[c] == *(a+i)
-a[i] == *(a+i) -if matrix allocated at compile time, laid out contiguously -if allocated dynamically, it's array of pointers to vectors => matrix = array of pointage
=> matrix = array of pointers
Geording for beggest/smallest/an Hem Lecture?
- if array ordered, get min & max in constant time -binomy search fastest of ordered
-Strings = arrays of characters, ends in 10
-char $s[] = "word"$ char $*g = word $ char $s[] = \{ word $ char $s[] = \{ w , o , r , d \}$
-stromp(): Go means gct >0 means s>t
= 0 means s==t -strlen 1): # of non-null characters
- strappyll: copy str it not null bremember to tack 'V'atend
the state of the s

-stracting (): checks to make sure we didn't go too few (depends on knowing how large array is) - strocpyl): must account for null at the end -Sorting -> place into defined order -ex) lexicographical, natural order, total a partial ordering -sorting adds into to our data

(can make assortions about before, after, lesser, greater, etd) -Menumerate all possible orderings of nobjects'=> inefficient of tooling

- selection sort

- find smallest item by looking at all n of item them

- place in first stot

- if there's anything in that slot, swarp it {- array of 1 element = gorted {- array of 2 elements = D check before or after 19t element - array of 3 elements = D check possible public possible public if -> not efficient, goes up to we large # of elements -Bubble Sort - if 1st element > 2nd element then swap them - if 2nd > 3rd element, swap them too, etc. - last element es largest -> any book at n-1 pairs -no swaps = already sorted -to make code efficient, check if sorted already

	-Merge Sort -split elements into disjoint pairs (act), aci] d a[z], a[3])
All Carrier Service Constitution of the Consti	-split elements into disjoint pairs (als), als a de situation
	- n # of pairs but still n # of elements
	-place elements of each pair in order of call it a run
and the Montant of the same popular	-then take pair of runs of merge them ento runs of 4
and Special contractions of the special specia	- can double 1 before we exceed n by boyn
and him a make	-sort funishes in time proportional to inlogin
(SORT)	$O(n^2)$ sorts $O(n^{5/3})$ $O(n\log n)$
Line	Selection shell (good) merge heap guick (average)
	insertion heap
	selection quick (average)
	queck (worst one)
	for any choice of perotes, there are arrays
	-Heaps
	- types: min/max, leftist, binomial, fibonacci, brodal, radix, unial
	- max hear (put largest child & swap of parent)
	-single node = heap
	- hear of parent is bear of trees nooted at booth
	children one heaps
	- pavent's value (key) > either child's
	- no order among children
July July and A	
	radix sort runs propertional to # of degits * # of records
	netter algorithm > faster computer

POINTERS (ptr) - Pointer = variable, holds memory address - points to location of object in memory -not all pointers contain an address (can be set to NULL) - NULL for macros can be: ((wid*)0), 0, 01 -memory stored in registers, accessed by address -each byte has unique address
-byte = word of size 4 7 pointers point to address they're assigned - assign pointer the address of a variable using & Lo multiple pointers can point to same address process = dereference or indirection (use * operator) -useful for manipulating values of several variables w/ call by ret -CX) int fob = 13; int & bor = & foo; I foo's address Stoned in pointer bar -pointers have addresses of EX) int ** bor 2 = & bor "ban's address in ptr bar 2 Tuse ptr for dynamic data & large #s of It returns - Passing by value = May displicates passed values onto stack -passing by reference = duplicates ptr onto stack Li returns multiple values, passes large data queckly - pointer arithmetic = same L> 5x new) array [i] == * (a+i) of ptr

(1M2) 1000 1001 1002 1003 1004 1 4005 4006 4007 1008 arr[2] Sylvan bytes per mt > street() M O LILLA -addray/subtracting ent w/ ptr / -subtracting z ptr / fort mush -adding 2 ptr × 3 don't do Mese and o have where of fining organization - awrays can be written using ptrs

- declaring array allocases function on stack

- global array in data area

- dynamic array dec. to get ptr allocates it on heap -str = ptr to array of an chars Lis can be indexed, passed by ref -multi dun array -> EX mt a3 = {50,1,23, -can be of any dum \$3,4,53. £3,4,5%, £ 8 5 3, 86, 7333; - Function Ptrz

- pourts to executable code in memory (not data value)

- deref func ptr yeelds referenced func as use ()