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- Insertions are at the rear of the queue and removals are at the front of the queue. You do not add, remove or access from any other point in the queue.
- Like stacks, queues are an Abstract Data Type (ADT) with more than one implementation.

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  - integer size(): returns the number of elements stored.
  - boolean isEmpty(): indicates whether no elements are stored.

# Example

Operation	Return value	Queue
enqueue(5)	-	(5)
enqueue(3)	-	(5, 3)
dequeue()	5	(3)
enqueue(7)	-	(3, 7)
dequeue()	3	(7)
first()	7	(7)
dequeue()	7	()
dequeue()	null	()
<pre>isEmpty()</pre>	true	()
enqueue(9)	-	(9)
enqueue(7)	-	(9, 7)
size()	2	(9, 7)
enqueue(3)	-	(9, 7, 3)
enqueue(5)	-	(9, 7, 3, 5)
dequeue()	9	(7, 3, 5)

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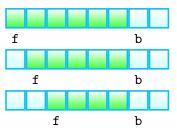
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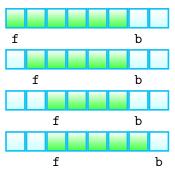
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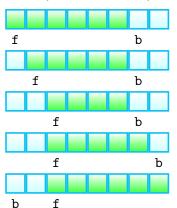
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- Two variables keep track of the front and size:
  - front: index of the front element
  - size: number of stored elements
- When the queue has fewer than N elements, array location back = (front + size) mod N is the first empty slot past the rear of the queue.

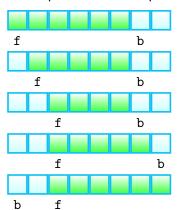




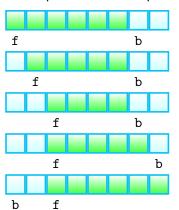




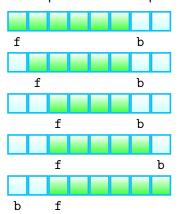


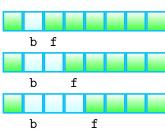


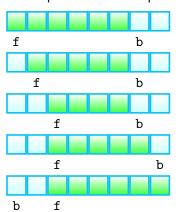


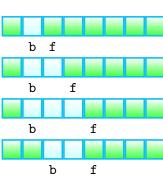












```
procedure ENQUEUE(o)

back \leftarrow (front + size) \mod N

arr[back] \leftarrow o

size \leftarrow size + 1

end procedure
```

Keep in mind that you need to consider the case when the queue is full and how to handle it.

```
procedure DEQUEUE item \leftarrow arr[f]front \leftarrow (front + 1) \bmod Nsize \leftarrow size - 1return itemend procedure
```

Notice that the cell in the array, arr[f] is not set to null after it is dequeued. Consider the case, when the queue is empty and you attempt to dequeue.

## Linked List-backed Queue

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- A linked list could also be used as the backing data structure of a queue.
- Elements would be added and removed to opposite ends of the list. The most efficient way to do this would depend on the list.

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  - Each operations runs in time O(1).

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  - During the process of resizing, the queue would be "unwound", so that the front of the queue would be in index 0 of the resized array.
- Linked list-backed queue do not have either of the above two limitations.

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- We can implement a round robin scheduler using a queue processes by repeatedly performing the following steps:
  - 1. process = processes.dequeue()
  - 2. Do work on process
  - 3. processes.enqueue(process)