



MAKE  
SCHOOL

# HEAPS

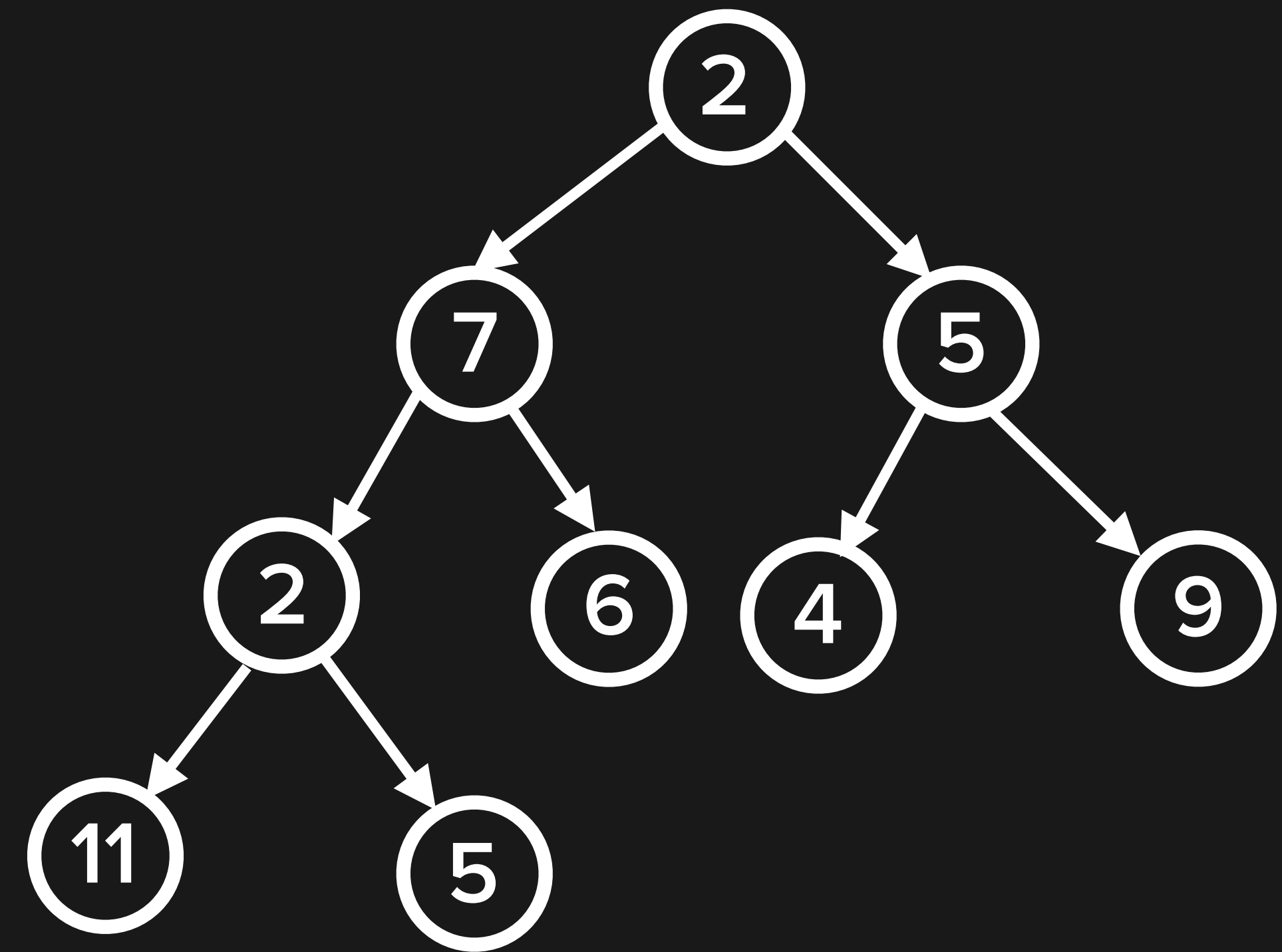
*Heaps more fun than a barrel of monkeys*

# COMPLETE BINARY TREE

*(REVIEW)*

Every level except  
possibly last is completely  
filled and nodes are as far  
left as possible

Height:  $O(\log_2 n)$



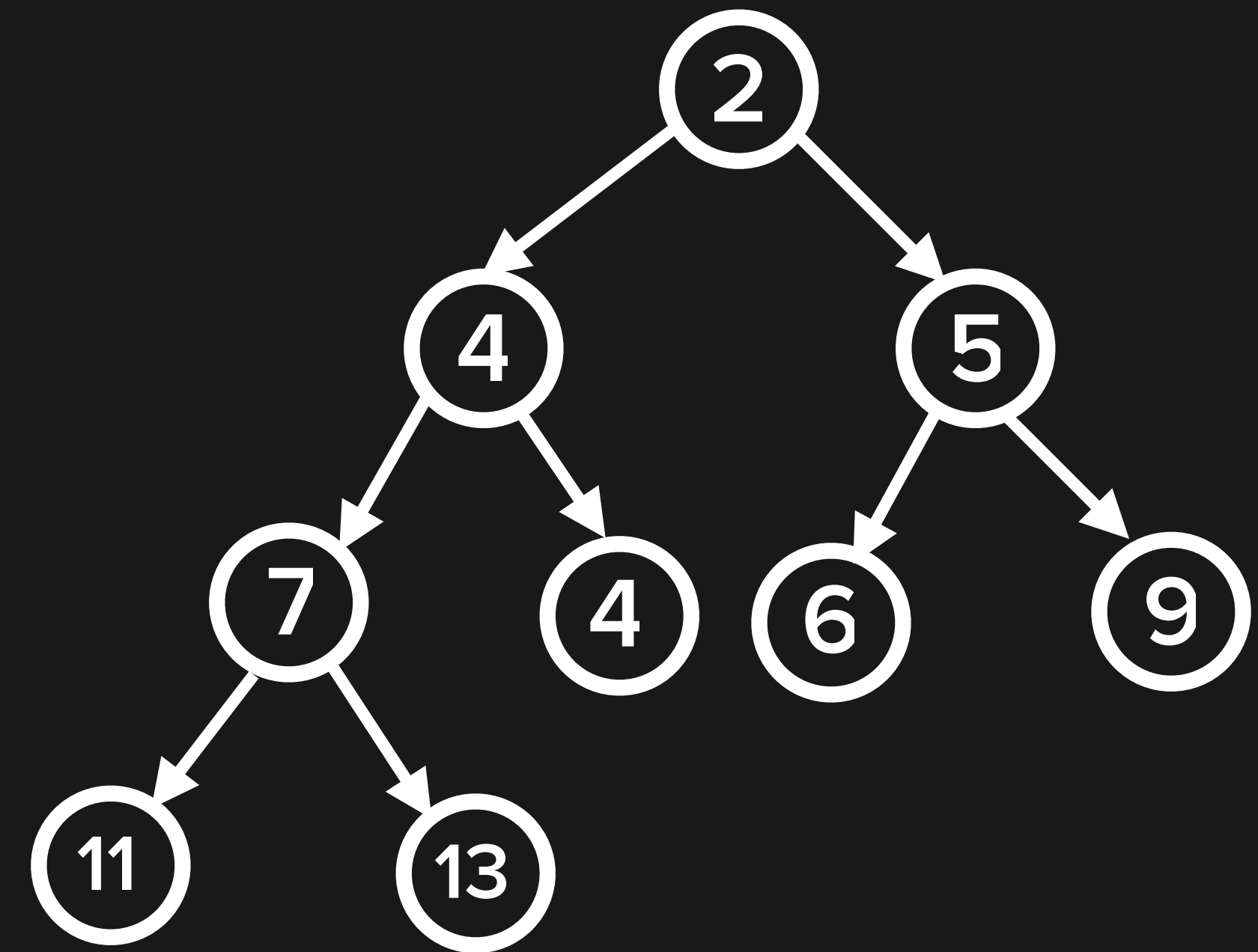
# BINARY HEAP DEFINITION

A complete binary tree

Satisfies heap ordering property

*min-heap* - each node is greater than or equal to its parent (min value is root)

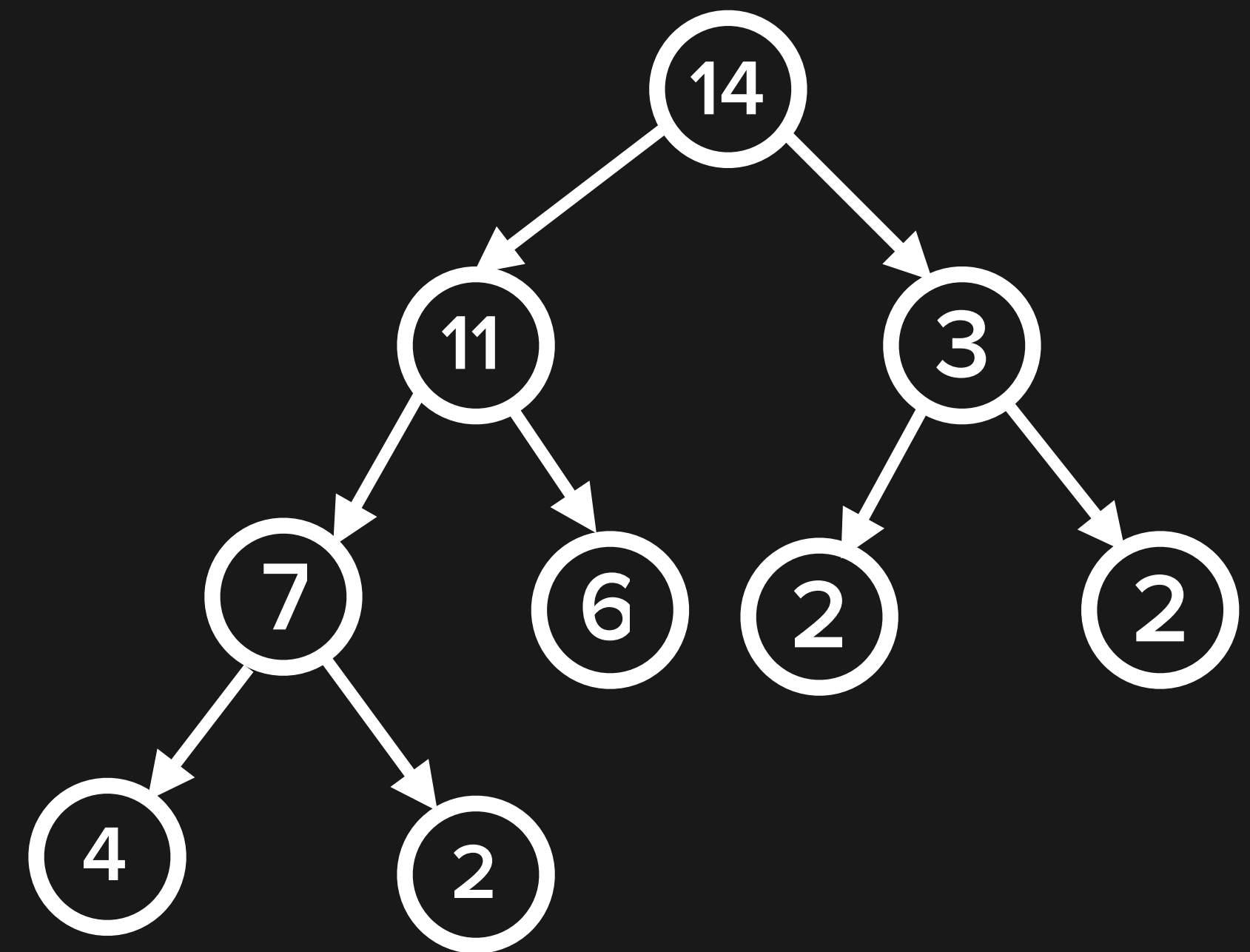
*max-heap* - each node is less than or equal to its parent (max value is root)



*min-heap*

# CAREFUL

Heaps are *not* sorted, instead they are considered “partially ordered”



*max-heap*

# PRIORITY QUEUES

Almost always implemented with a heap

Elements with smaller numbers are higher priority

Elements are inserted in  $O(\log n)$  time instead of  $O(n)$  time for a sorted array or linked list

Ordering happens with each insertion, so the cost of ordering is distributed across insertion instead of in one big chunk

# PRIORITY QUEUE APPLICATIONS

Prioritizing data packets in routers

Tracking unexplored routes in path-finding

Bayesian spam filtering

Data compression

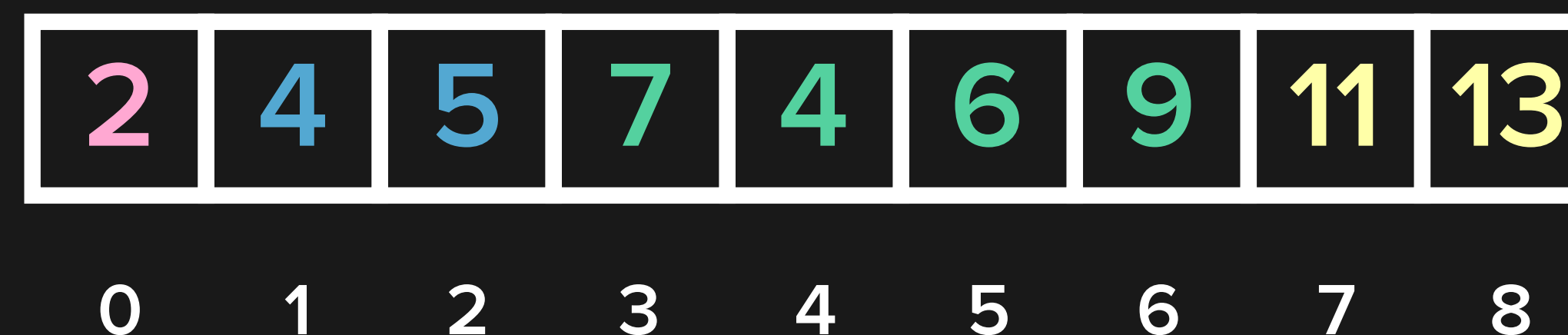
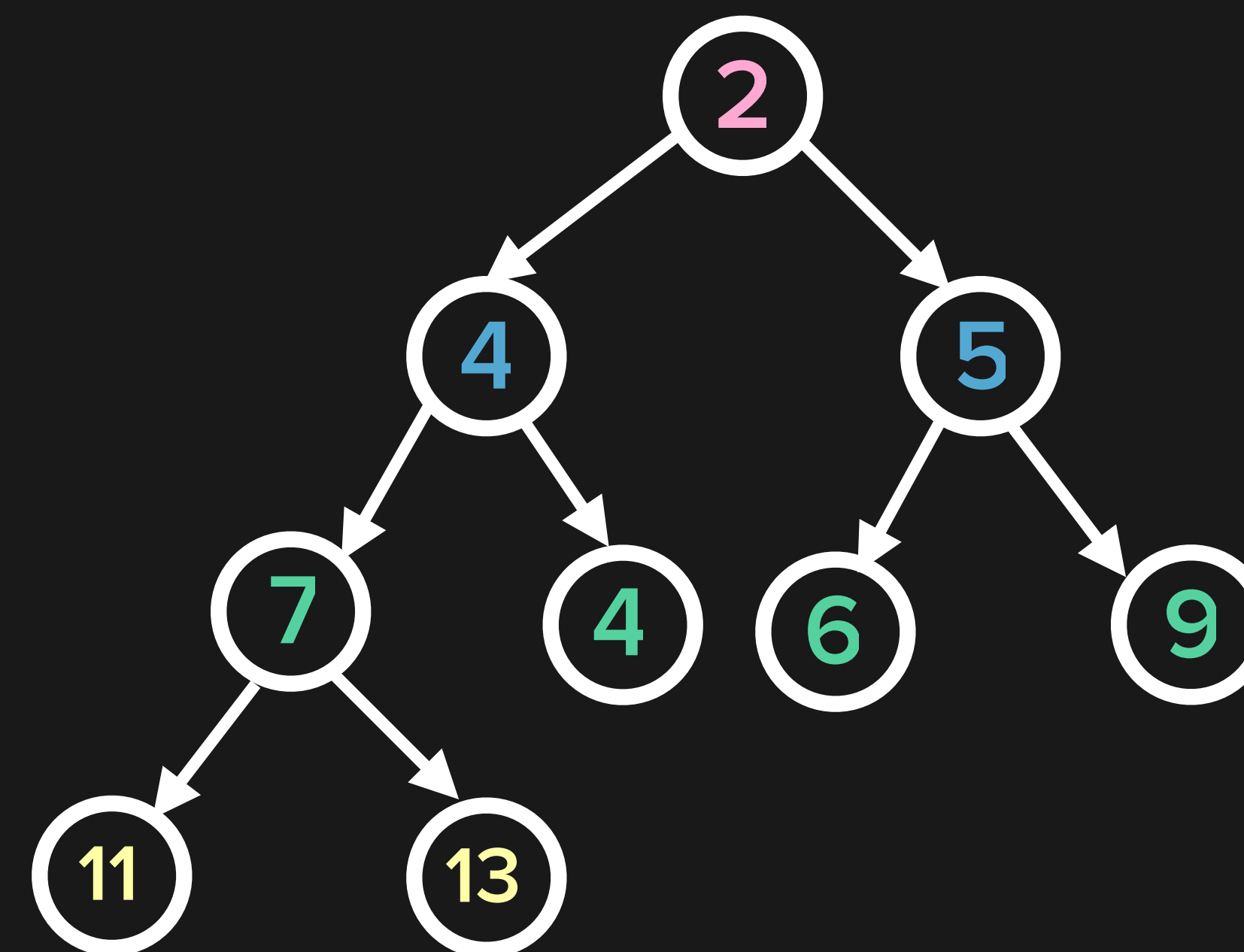
OS: load balancing, interrupt handling

# ARRAY REPRESENTATION

Items stored in (dynamic) array following level-order traversal

Calculate parent-child index relationships with arithmetic

- Left child index:  $2n + 1$
- Right child index:  $2n + 2$
- Parent index:  $(n - 1) / 2$

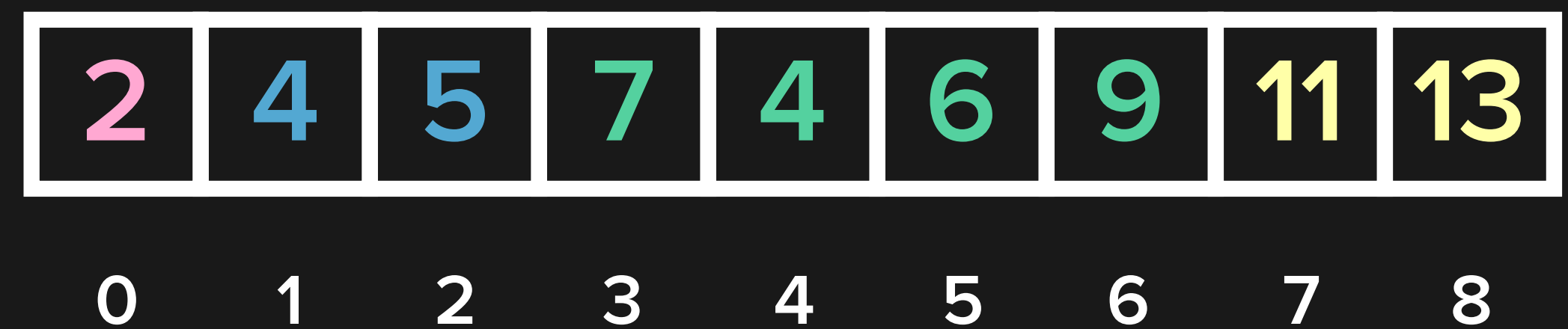
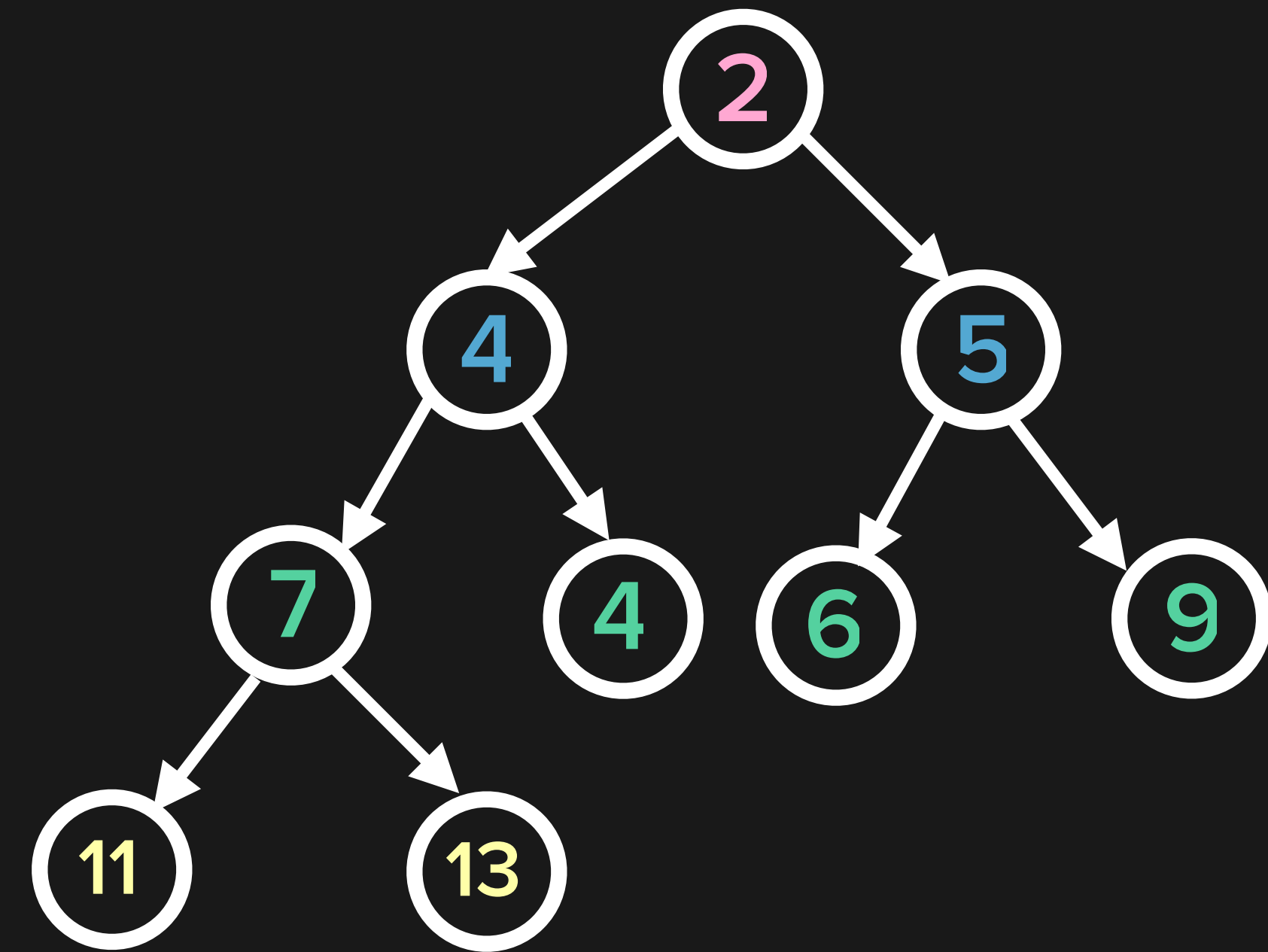




# ADVANTAGES

Uses less memory than binary tree represented with nodes  
(avoids node objects containing 3 pointers: data, left, right child)

Allows sorting an array in-place  
(*heapsort*)

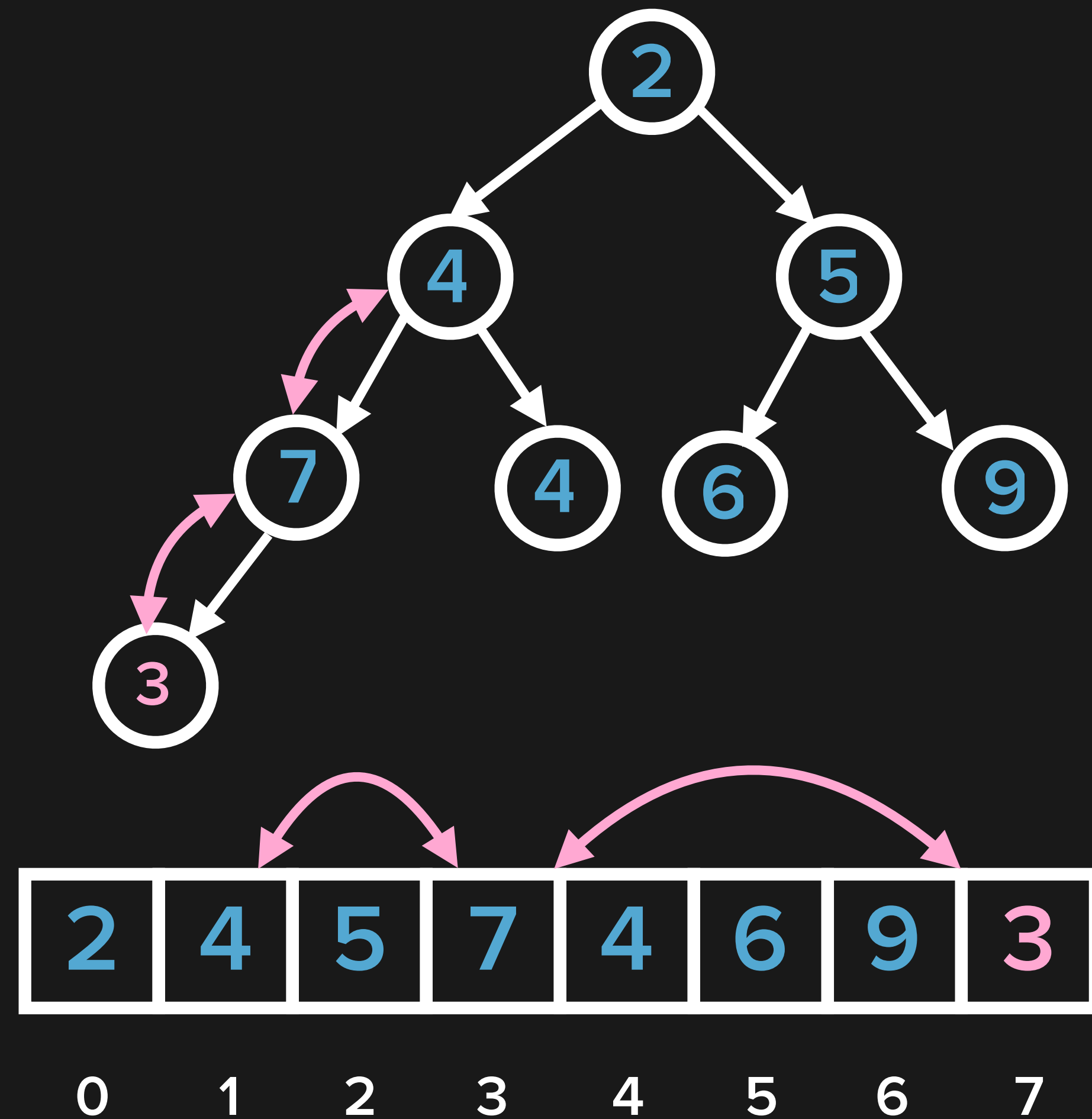


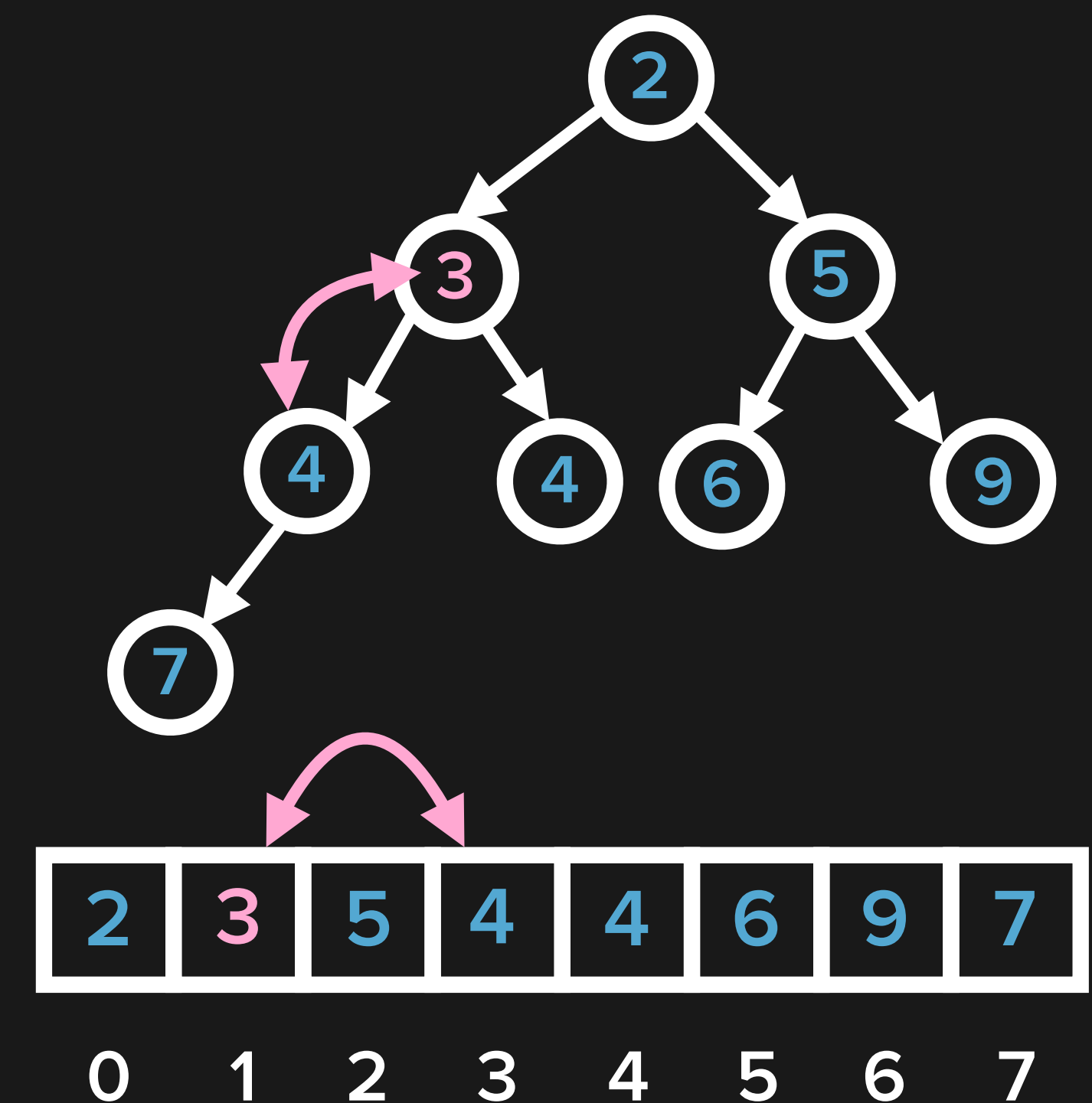
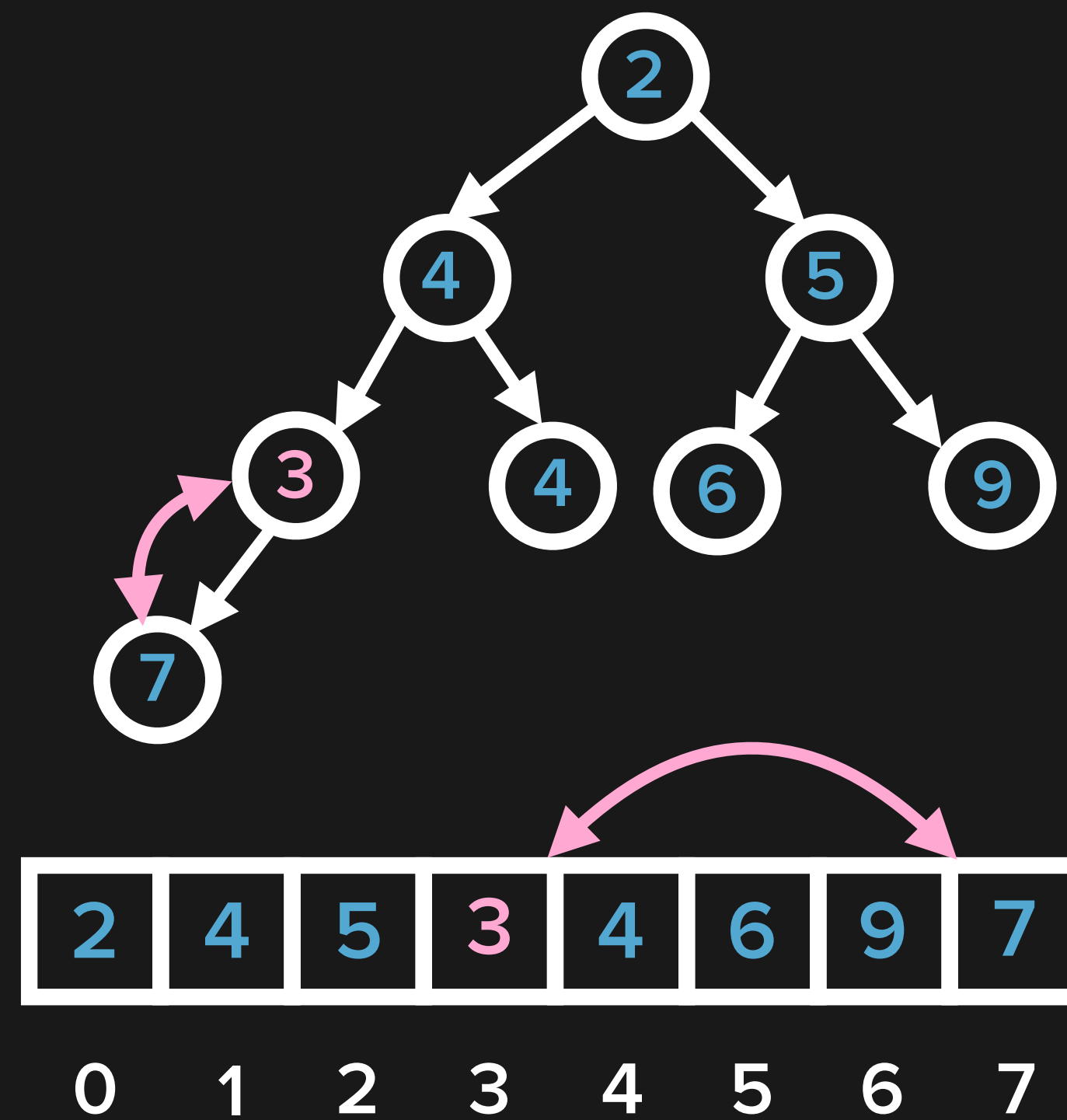
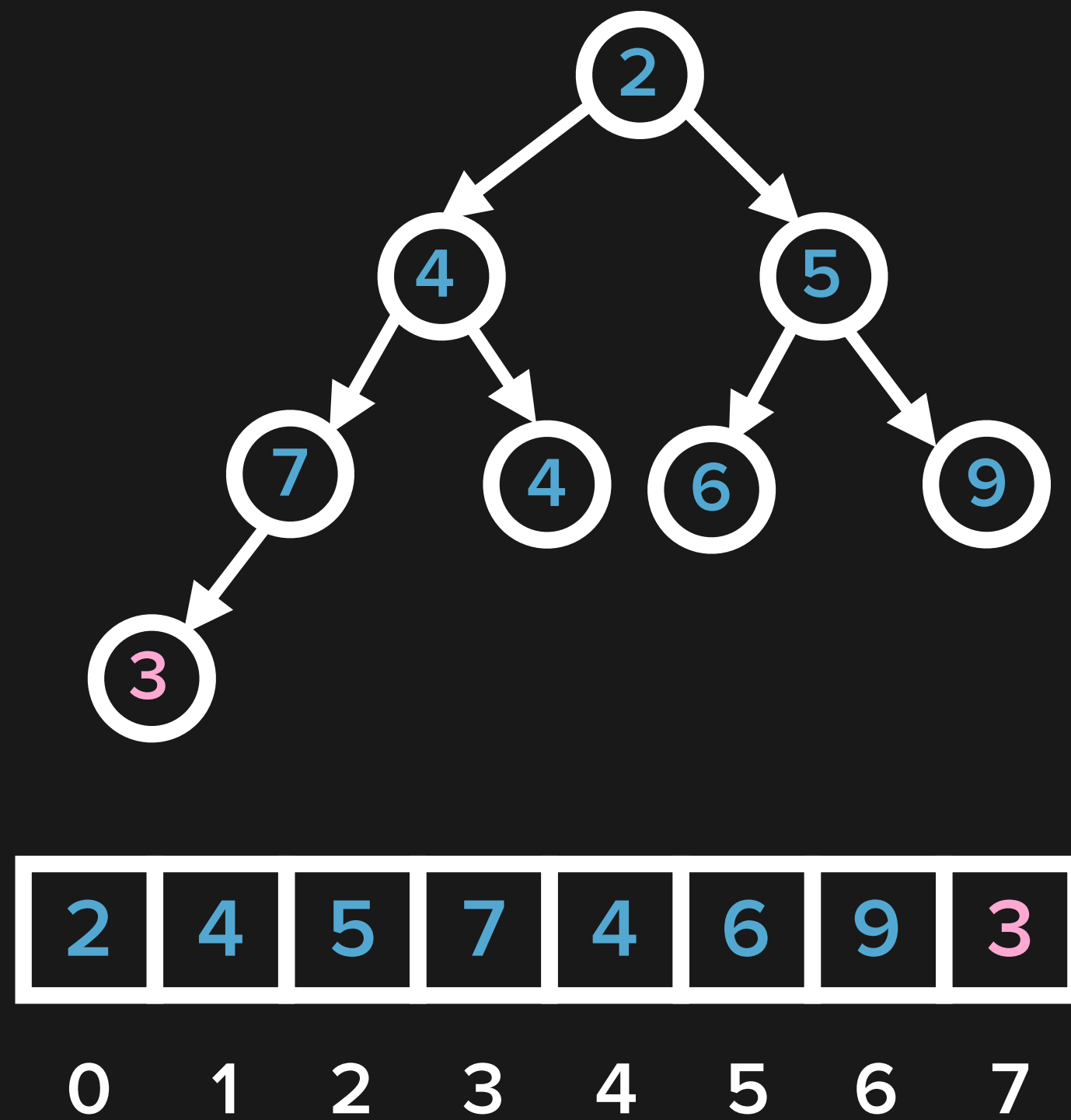
# INSERT

Add element to end

*Sift up (aka bubble up,  
percolate up, trickle up)*

Swap with parent up to  
the root until path fulfills  
the ordering property





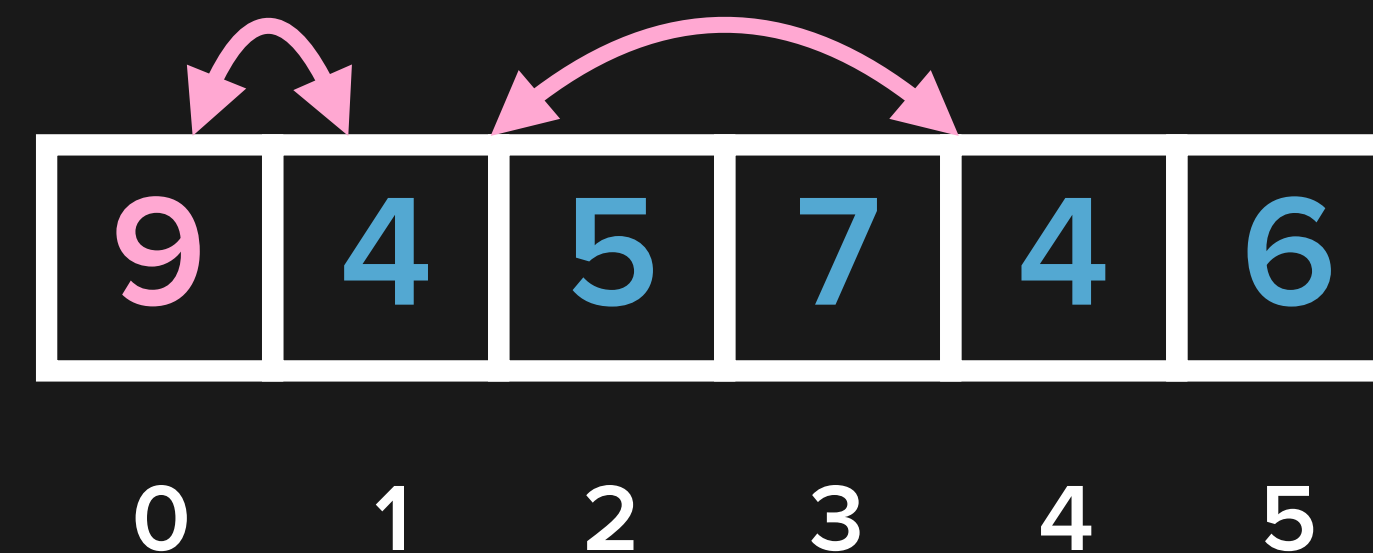
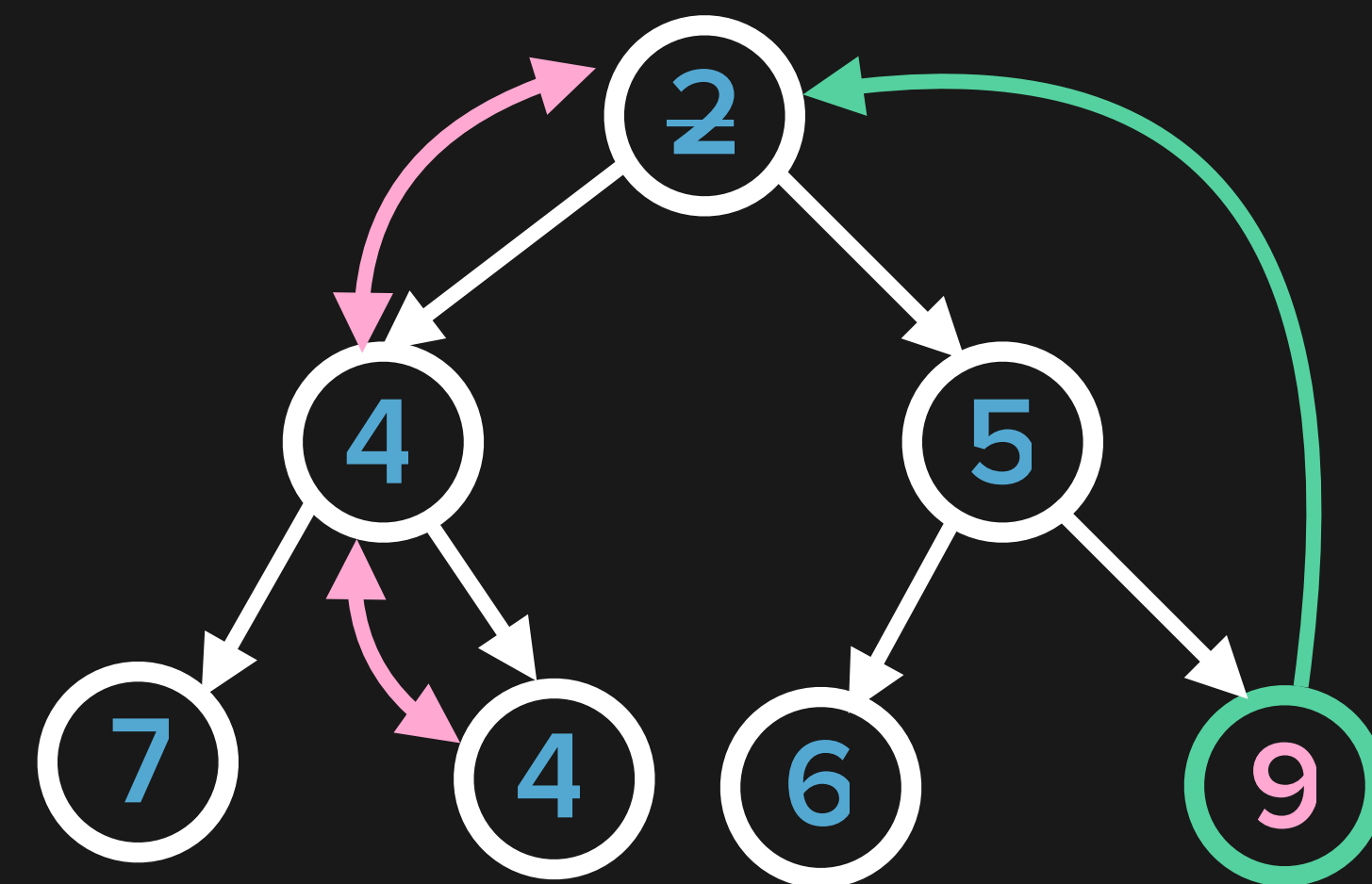
*Insert*

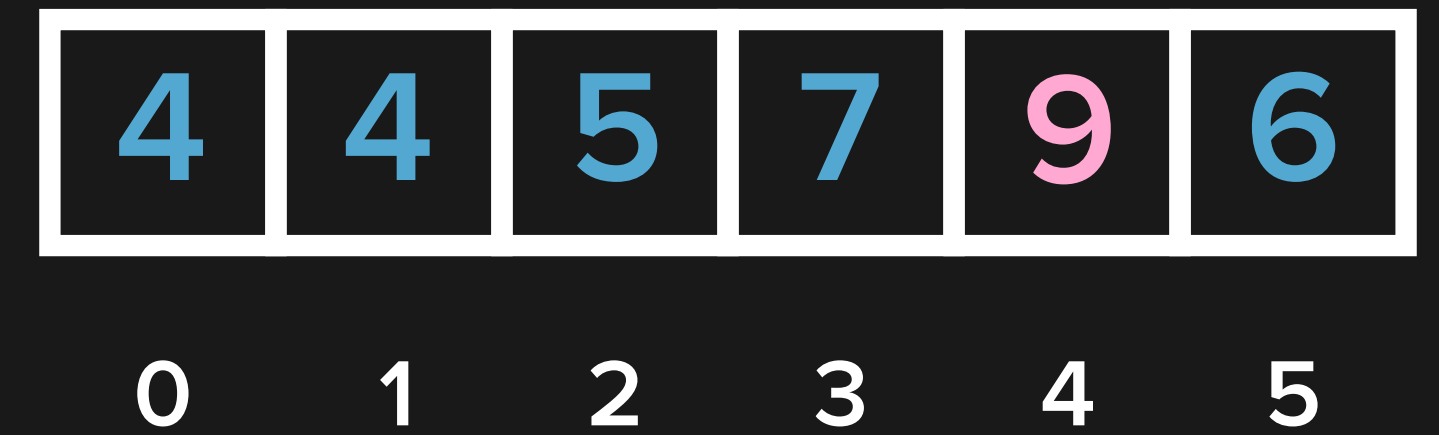
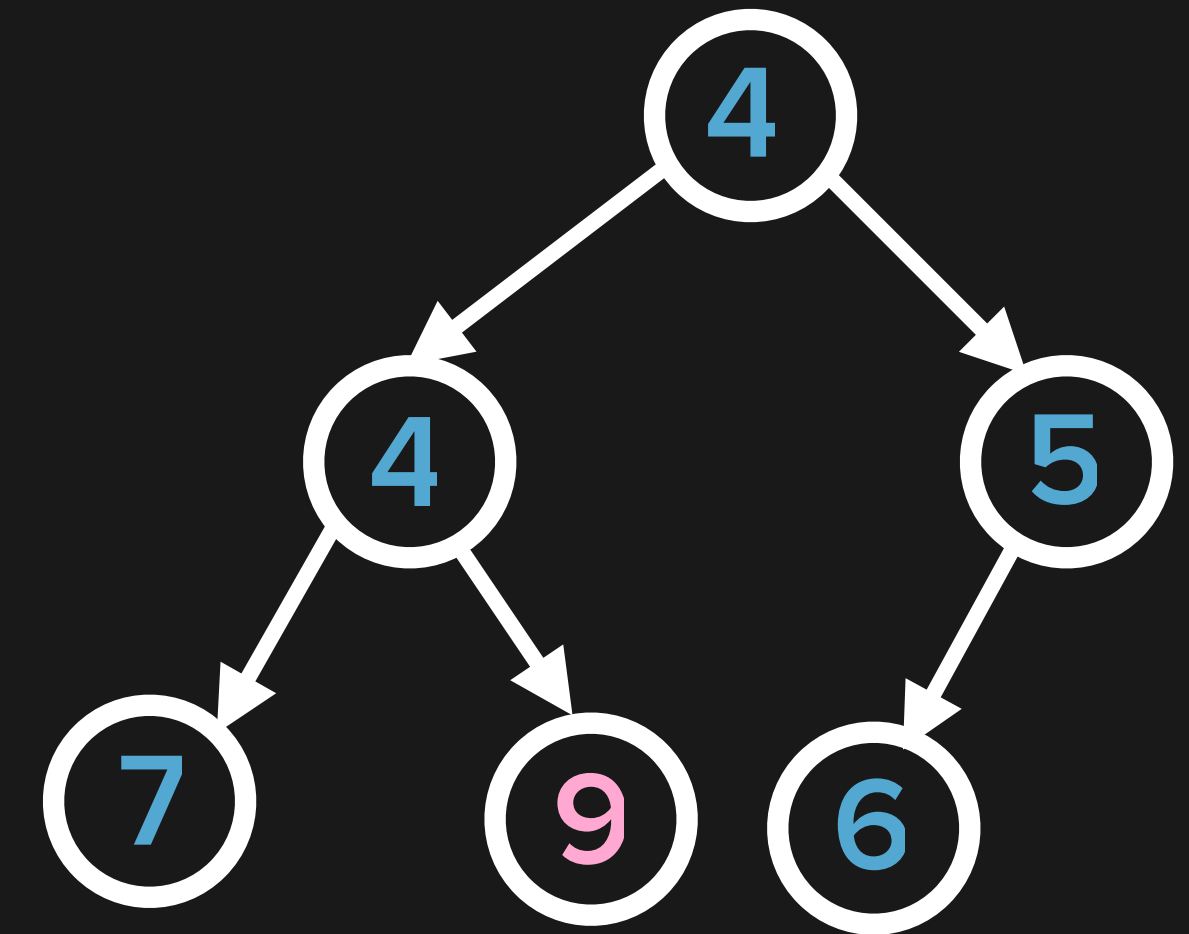
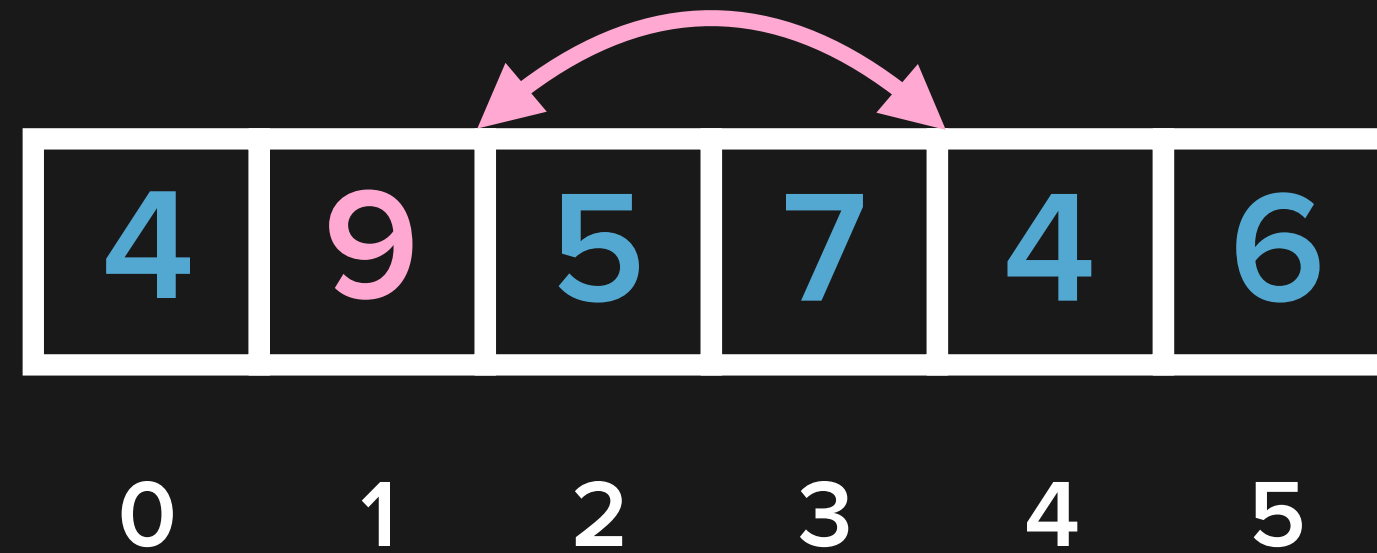
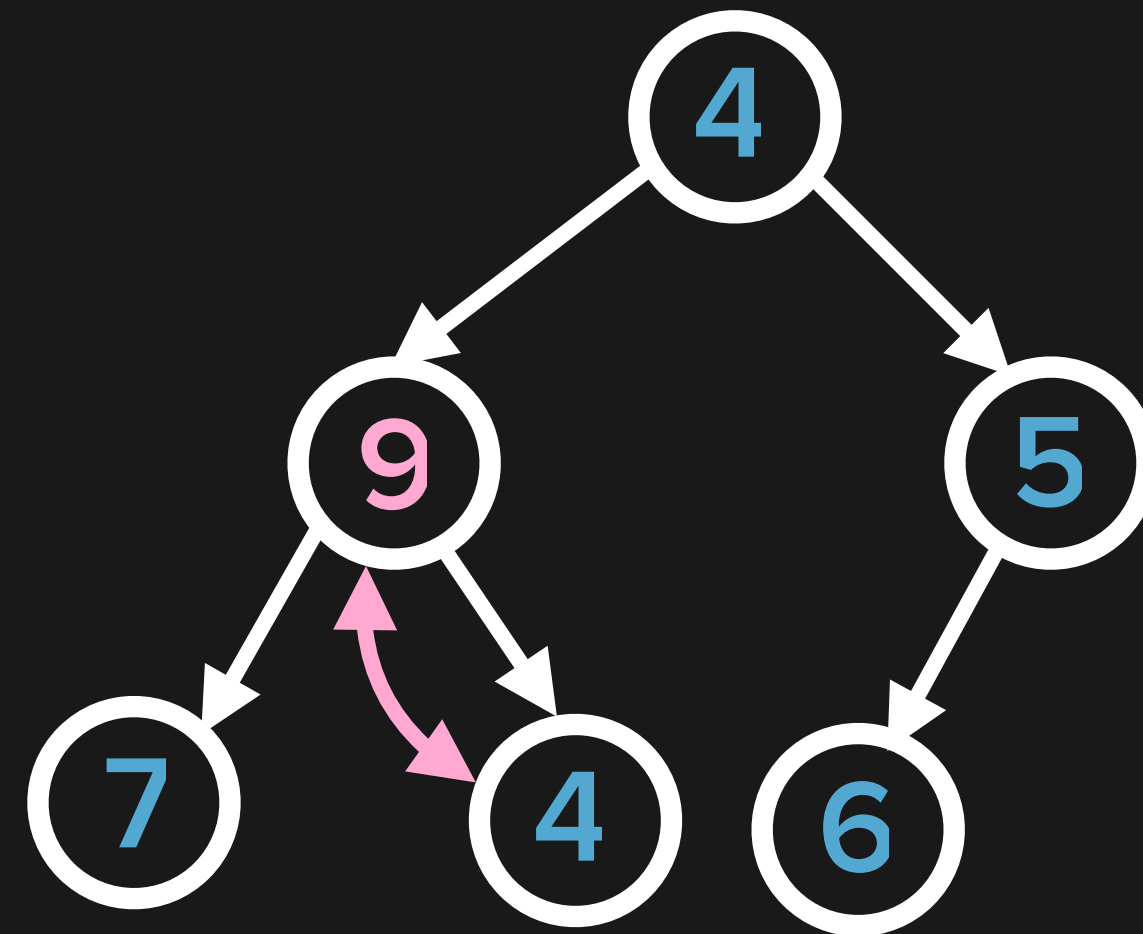
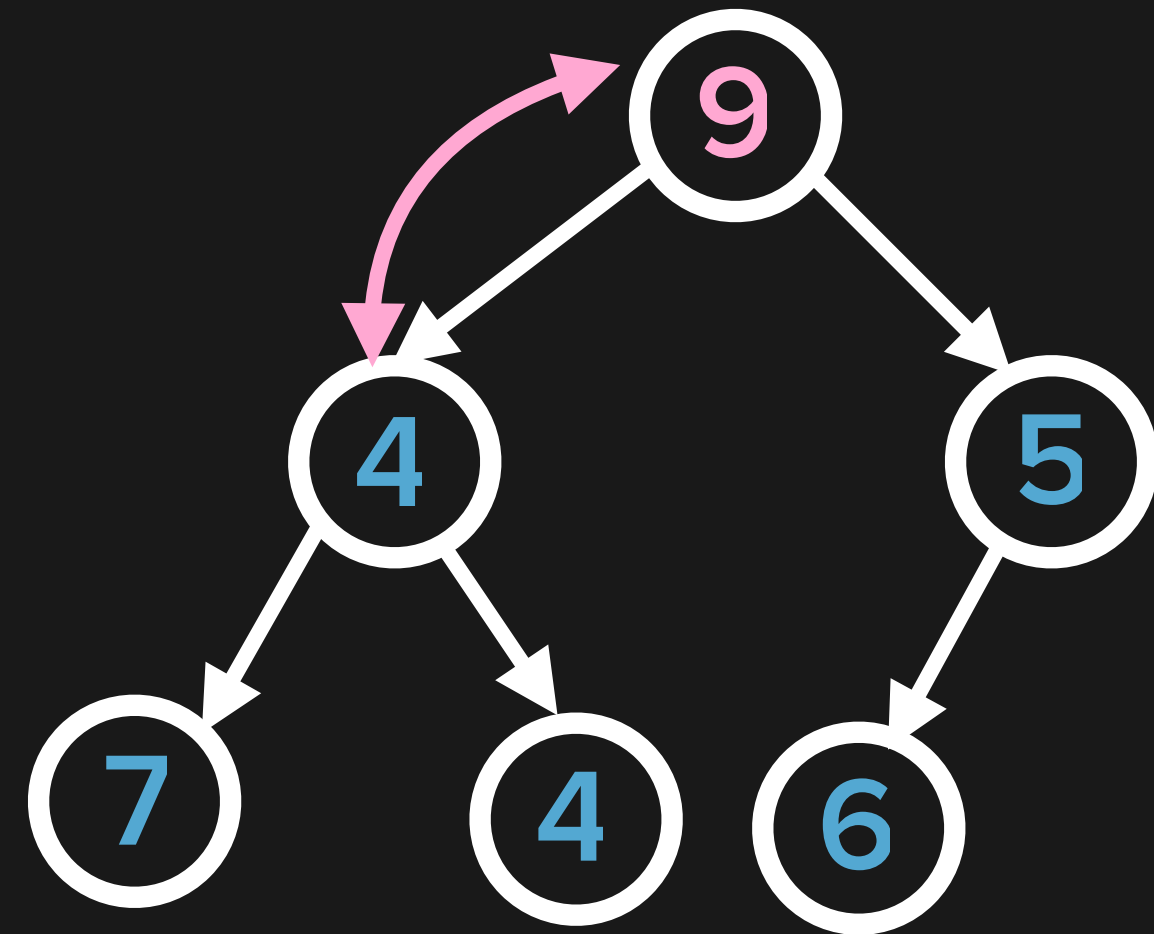
# DELETE MIN / MAX

Replace root with last element

*Sift down (aka bubble down, percolate down, trickle down)*

Swap with smaller child (min) or larger child (max) until trio fulfills the ordering property





*Delete min*

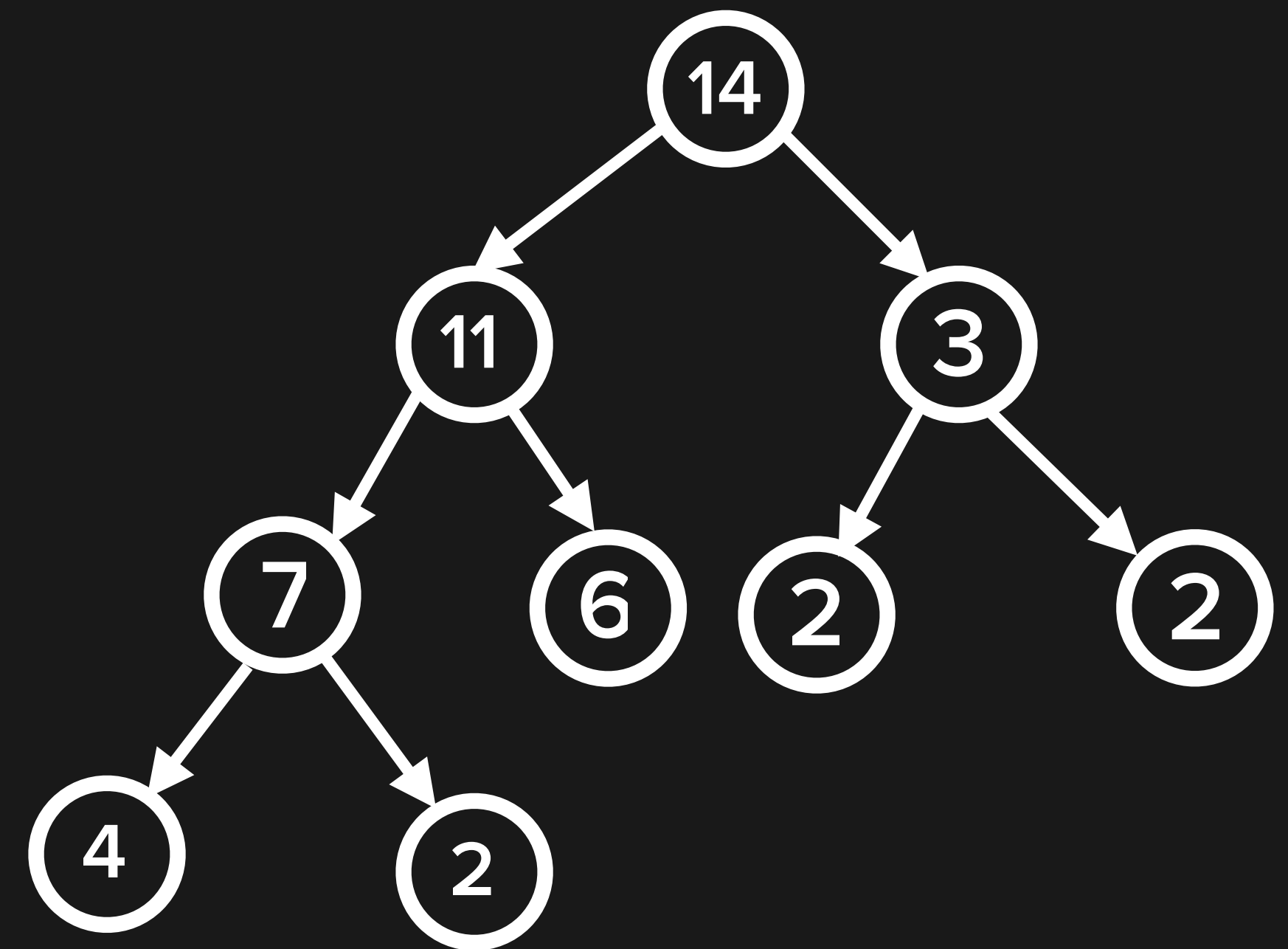
# OTHER METHODS

*peek* (aka *find-min* or *find-max*)

returns the root value

*size* (aka *count* or *length*)

returns number of elements



*max-heap*

# HEAPIFY

Input is an array (usually unsorted, unordered)

Output is an array that satisfies the binary heap ordering property

# HEAPIFY

Start at last parent node

`index = (count - 2) / 2`

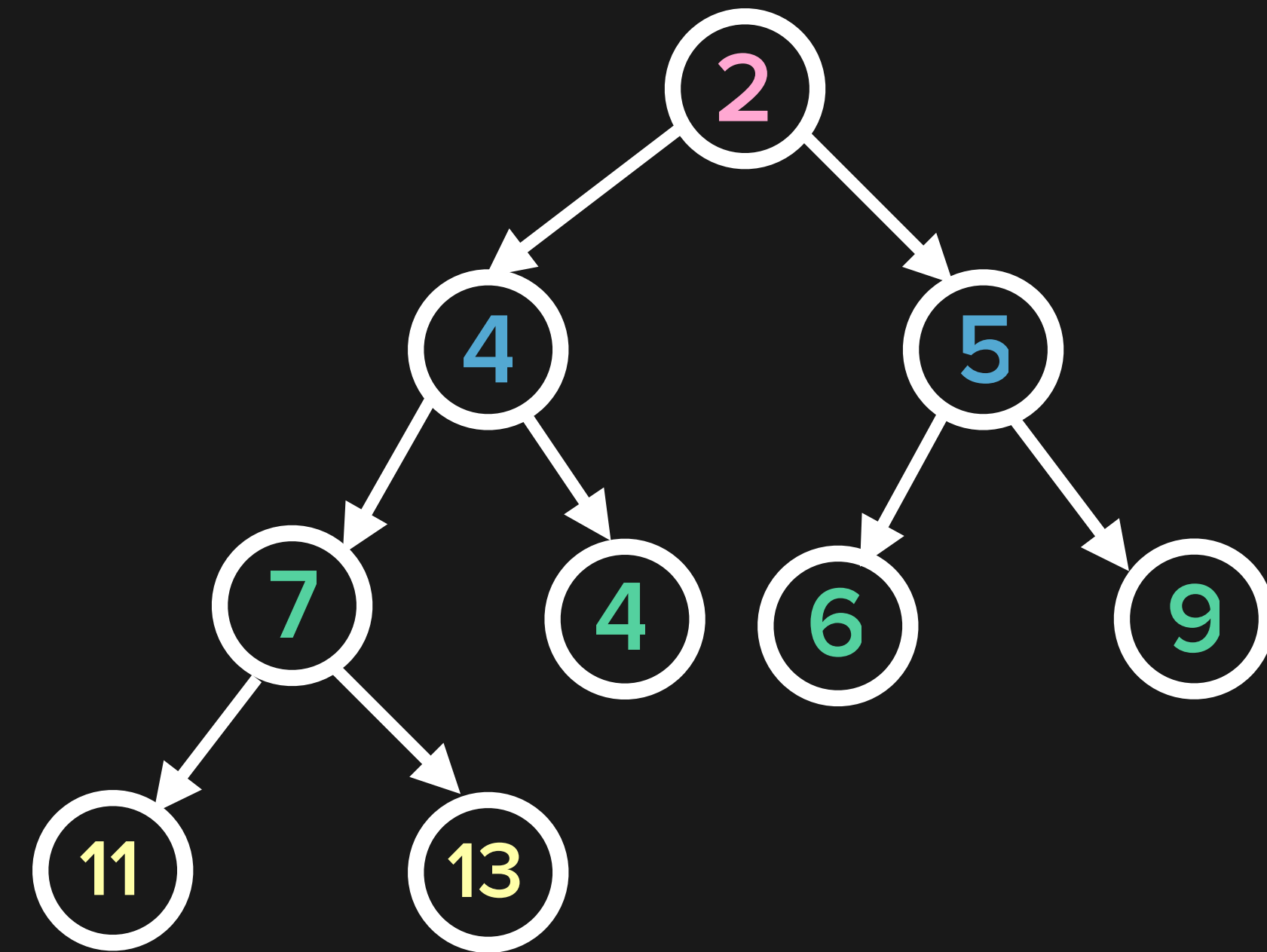
`while index >= 0:`

Sift down element at `index`

`index -= 1`



Start at **index** =  
**(count - 2) / 2**  
because that's the last  
parent node



# HEAPSORT

heapify array

while (count > 0):

Grab **min** or **max** element (*peek*)

delete-min or **delete-max** element

# HEAP RUNTIME

	Average Case	Worst Case
Space	$O(n)$	$O(n)$
Insert	$O(\log n)$	$O(\log n)$
Delete	$O(\log n)$	$O(\log n)$

# HEAPSORT RUNTIME

	Average Case	Worst Case
Space	$O(n)$	$O(n)$
Heapify	$O(n \log n)$	$O(n \log n)$
Heapsort	$O(n \log n)$	$O(n \log n)$