OSY.SSI[2018][2]

ACCESS DENIED

In the news...

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- ... "It's unfortunate, but each camera will need to be updated manually by users"

In the last episode...

 $\ \cdot \ CIA \cdot 50/50 \cdot \times$

Insert video here.

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What could go wrong?

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Access control

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Formal models: The Bell-LaPadula ACM

The goal of access control is to provide

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Access control implementations must be **NEAT**

- ► Non-bypassable
- Evaluable
- Always invoked
- Tamper-proof





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Example: chmod

So far so good

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Not obvious... More on that in a minute!

AC in practice? CBAC vs ACL

- ► Capability-based (CBAC): You are given a token that provides access (think key).
- ► Access control lists (ACL): Access is granted by your presence on a list (think VIP Party!)

Whichever flavour you fancy most, they both rely on

- ► A certain notion of *identity*
- ► A form of *authority* in control

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What is identity?

When Metaphysics meets Science

Strong version: what characterises an individual?

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Henceforth,

 $\overline{\mathsf{Identity}} \Leftrightarrow \mathsf{Having} \ \mathsf{the} \ \mathsf{Secret}$

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Examples ? Realistic examples?

In short, secrets seem to require... access control themselves.

Aside: Keep secrets secret!



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$$Pr[Auth_{X,Z}(Y) \mid Auth_{X,Y}(X)] = negl$$

Examples?

How to achieve non-transferability?

Any idea?

Zero-knowledge proofs

Fact:

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An example from my magical hat

Exercise: Can you think of another one?

Important distinctions

So we have:

- ► A **secret** (which is secret)
- An authentication mechanism (which is public)
- ► An access control policy (which may be public or not, but better be)

These three things are different and (in principle) independent.

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REMEMBER THIS: Separation of policy and mechanism

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Formal models

The goal of a formal model is to prove security properties.

This becomes necessary as soon as the system becomes large.

Bell and LaPadula designed the first provable AC model.

The archetypical Access Control

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- ► Key idea: "Good" state + "valid" operation ⇒ "Good" state.

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Question: Does it work?

Some limitations...

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Do you see why?

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Dilemma: expressive AC vs. correct AC.

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Dilemma: expressive AC vs. correct AC.

BLP crime: confusing policy and mechanism

Information flow?

Question: can information flow be blocked?

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More about that in a later lecture!

P *clap* A *clap* U *clap* S *clap* E

Send me CVs • Ask questions • Brush your teeth