

all the things

Time & Time Complexity

$$T\left(\frac{n}{b}\right) + f(n) \quad \text{Master Theorem}$$

input
of subproblems

by which subproblem is reduced
ch to recombine/split everything back

, if $\log_b a > c$, then

$$T(n) = \Theta(n^{\log_b a} \log n)$$

and $\log_b a \leq c$, then

$$O(n \log n), O(n), O(\log n)$$

(1.) $3T\left(\frac{n}{2}\right) + n^2 \quad c=3$

$$\log_2 3 > 3 \quad \text{not case 3}$$

$$\Omega(n^3) \quad \log_2 3 > c \quad c=1$$

$$\Theta(n^2) \quad \boxed{\text{Case 3}}$$

(2.) $7T\left(\frac{n}{2}\right) + n^2$

$$\log_2 7 > 2 \rightarrow \Theta(n^{\log_2 7}) \quad \boxed{\text{case 1}}$$

(3.) $4T\left(\frac{n}{2}\right) + n^2$

$$\Theta(n^2) \quad \boxed{\text{Case 3}}$$

$$\log_2 4 = 2 \rightarrow \Theta(n^2 \times \log n) \quad \boxed{\text{case 2}}$$

(4.) $\log_4 3 \quad n \log n$

$$\Omega(n') = n \log n$$

$$\log_4 3 < 1 \rightarrow \Theta(n \log n) \quad \boxed{\text{case 3}}$$

(5.) $\log n \quad \log n \quad c=1$

$$\log_2 4 > 1 \rightarrow \Theta(n^2) \quad \boxed{\text{case 1}}$$

all the things

⑥ $T(n-1) + n$

$$\begin{aligned} \frac{n}{b} &= n-1 \\ n &= bn - b \\ (b-1)n &= -b \\ (\cancel{b-1})n &\stackrel{?}{=} -b \\ n &= \frac{-b}{b-1} \\ &\quad \left. \begin{array}{l} \text{Can't with MT} \\ \cancel{n-(b-1)b} \\ \hline b \end{array} \right. \end{aligned}$$

⑦ $4T\left(\frac{n}{2}\right) + n^2 \log n$

$$n^2 \log n = \Theta(n^2) \& \Omega(n^2)$$

$$\Theta(n^2) + O(n)$$

$$\Theta(n^2) \quad p=1$$

$$\Theta(n^2 \log^p n)$$

$$\Theta(n^2 \log^{(k+1)} n)$$

$$\Theta(n^2 \log_2 n)$$

case 2 extended

⑧ $ST\left(\frac{n}{2}\right) + \frac{n}{\log n}$

$$\frac{n}{\log n}$$

$$p=-1$$

all the things

$$\log_2 5 > \log_2 7$$

$$\Theta\left(n^{\log_2 5}\right) \text{ case 1}$$

⑨ $3T\left(\frac{n}{3}\right) + \frac{n}{\log n}$

$$\Theta\left(\frac{n}{\log n}\right) = \Theta(\log^{-1} n)$$

$$O(n) = \frac{n}{\log n}$$

$$\log_3^{-1} = 1 \rightarrow p = -1$$

$n \log \log n$ MT doesn't work

⑩ $2T\left(\frac{n}{2}\right) + C$

$$\Theta(n^1)$$

$$\log_2 < 1 \rightarrow \Theta(f(n)) = \Theta(C)$$

⑪ $T \log n$

$\log_4 1 < 1 \rightarrow C=0$

$T\left(\frac{n}{4}\right) + \log n$

$\Theta(n^{\log_4 1})$

all the things

(12) Doesn't Apply

(13) $P = \log_4 2^{\lfloor \frac{n}{2} \rfloor}$

case 1

$\Theta(n^{\frac{1}{2}})$

$\log^k n$
 $\Theta(n^k \log^p n)$

$k=1$
 $p=1$

$\Theta(n \log n)$

all the things