1. Remember: Core STM Concepts and Concurrency Issues

Define TVar and Basic STM Operations

```
import Control.Concurrent.MVar
import Control.Concurrent
-- Define a transactional variable (TVar)
data TVar a = TVar (MVar a)
-- Function to create a new TVar
newTVar :: a -> IO (TVar a)
newTVar val = TVar <$> newMVar val
-- Function to read from a TVar
readTVar :: TVar a -> IO a
readTVar (TVar mvar) = do
                val <- readMVar mvar</pre>
                threadDelay 5000000 -- Add a delay of 5 second
                return val
-- Function to write to a TVar
writeTVar :: TVar a -> a -> IO ()
writeTVar (TVar mvar) val = modifyMVar mvar (\ -> return val)
```

Identifying concurrency issues: This simple version of TVar does not yet solve concurrency problems like race conditions or deadlocks. It uses MVar for thread safety but does not handle transaction rollbacks or retries.

Atomicity, Consistency, Isolation: The next steps will aim to build the transaction mechanism to achieve atomicity and isolation by running multiple operations as a single unit (the core of STM).

As an experiment, we can proceed as follows:

- 1. Define thread variable ty with a default value.
- 2. Read and double check the value
- 3. Have a readTVar func running in background
- 4. Attempt to writeTVar during the background execution
- 5. Read values to double check.

As a result, everything plays out as expected, the function reads the value before the value gets overwritten.

2. Understand: STM for Preventing Concurrency Issues

Next, I implement a very basic version of STM to group operations on TVars within an atomic block. This will help prevent inconsistent states when multiple threads access shared resources concurrently.

Implement a Simple STM Monad

```
-- Function to write to a TVar
writeTVar :: TVar a -> a -> IO ()
writeTVar (TVar mvar) val = modifyMVar_ mvar (\_ -> return val)
-- Stage 2 update
-- Define a simple STM monad
newtype STM a = STM { runSTM :: IO a }
-- The atomically function to execute STM actions
atomically :: STM a -> IO a
atomically (STM action) = do
   -- Wrap action in exception handling to simulate a retry mechanism
   result <- try action
   case result of
     Right val -> return val
     Left ( :: SomeException) -> atomically (STM action) -- Retry on
failure
-- Functions to lift TVar operations into STM
readTVarSTM :: TVar a -> STM a
readTVarSTM tvar = STM $ readTVar tvar
writeTVarSTM :: TVar a -> a -> STM ()
writeTVarSTM tvar val = STM $ writeTVar tvar val
```

Here, I introduced the STM monad, which wraps IO operations on TVars. The atomically function executes STM transactions, and in case of an exception (such as a conflict in a real STM system), it retries the operation.

Concurrency issues and STM's retry mechanism: We simulate the retry mechanism
here by catching exceptions, but a real implementation would involve detecting conflicts
between transactions.

We can repeat the same experiment as mentioned in the previous stage which yields the same expected result.

```
ghci> :1 main2.hs
                                    ( main2.hs, interpreted )
[1 of 2] Compiling Main
Ok, one module loaded.
ghci> tv <- newTVar 0
ghci> tv <- newTVar 1
ghci> atomically (readTVarSTM tvar) >>= print
<interactive>:53:25: error: Variable not in scope: tvar :: TVar a0
ghci> atomically (readTVarSTM tv) >>= print
ghci> forkIO (atomically (readTVarSTM tv) >>= print)
ThreadId 1566
ghci> atomically (writeTVarSTM tv 100)
ghci> 1
ghci>
ghci>
ghci> atomically (readTVarSTM tv) >>= print
100
```

3. Apply: Implement a Job Scheduler Using the Custom STM

Next, I built a simple job queue using the STM and TVar implementations which will process jobs concurrently while using STM to ensure safe access to shared data.

```
type JobQueue = TVar [String]

-- Submit a job to the queue
submitJobSTM :: JobQueue -> String -> STM ()
submitJobSTM queue job = do
    jobs <- readTVarSTM queue
    writeTVarSTM queue (jobs ++ [job])

-- Process a job from the queue
processJobSTM :: JobQueue -> STM (Maybe String)
processJobSTM queue = do
    jobs <- readTVarSTM queue
    case jobs of
    [] -> return Nothing
```

This job scheduler submits and processes jobs using the custom TVar and STM. It ensures safe concurrent access by grouping readTVar and writeTVar operations into atomic transactions

To test the job queue, I also added a test case to analyze how it works.

```
-- Test case to test concurrency in the job queue
testConcurrentJobProcessing :: IO ()
testConcurrentJobProcessing = do
    -- Step 1: Create an empty job queue
   queue <- newTVar []</pre>
   -- Step 2: Spawn multiple threads to concurrently submit jobs
   let numJobs = 10 -- Total number of jobs to submit
   putStrLn "Submitting jobs concurrently..."
   replicateM 10 $ forkIO $ replicateM (numJobs `div` 10) $ do
        atomically $ submitJobSTM queue "TestJob"
   -- Step 3: Spawn multiple threads to process jobs concurrently
   putStrLn "Processing jobs concurrently..."
   replicateM 5 $ forkIO $ replicateM (numJobs `div` 5) $
processJobsConcurrentlySTM queue
    -- Step 4: Allow time for job submission and processing
   threadDelay 2000000 -- Wait for 2 seconds to allow all jobs to be
processed
    -- Step 5: Check if all jobs were processed
```

```
jobsRemaining <- readTVar queue
putStrLn $ "Remaining jobs: " ++ show jobsRemaining

-- Since we've submitted `numJobs` jobs, the queue should be empty
after all are processed
  when (null jobsRemaining) $
    putStrLn "All jobs processed. PASS"
  when (not (null jobsRemaining)) $
    putStrLn "Some jobs were not processed. FAIL"</pre>
```

And here is the result of it:

```
ghci> testConcurrentJobProcessing
Submitting jobs concurrently...
Submitting jobs concurrently...
Processing jobs concurrently...
Processing jobs concurrently...
PrPPOPPrPrcroroeococscecesesesissssnsisiginin ngngjg g o j jbjojo:obob b:b:T: : e T TsTeTetesesJststotJtJ
bJoJo
obobb
b
Remaining jobs: []
All jobs processed. PASS
ghci>
```

4. Analyze: Implementing and Benchmarking Concurrency

To analyze the performance of my custom implementation, I implemented two versions of the concurrency code, one with custom STM and one with MVar.

The main difference between is how the JobQueue is instantiated

CustomSTM	MVar
type JobQueue = TVar [String]	type JobQueue = MVar [String]

After a couple of experiments, there is no significant difference between the two implementations. As expected, CustomSTM performs worse in time (by 0.2%).

Here are the results of the test (CustomSTM above, MVar below):

```
ghci> :l main4tvar.hs
                                  ( main4tvar.hs, interpreted )
[1 of 2] Compiling Main
main4tvar.hs:49:5: warning: [-Wnoncanonical-monad-instances]
    Noncanonical `return' definition detected
    in the instance declaration for `Monad STM'.
    `return' will eventually be removed in favour of `pure'
    Either remove definition for `return' (recommended) or define as `return = pure'
    See also: https://gitlab.haskell.org/ghc/ghc/-/wikis/proposal/monad-of-no-return
49 l
        return x = STM (return x)
        ^^^^^^
Ok, one module loaded.
ghci> testCustomSTMJobProcessing
Submitting jobs concurrently...
All jobs processed. PASS
Elapsed time: 0.2102979s
ghci> :l main4mvar.hs
[1 of 2] Compiling Main
                                          ( main4mvar.hs, interpreted )
Ok, one module loaded.
ghci> testConcurrentJobProcessing
Submitting jobs concurrently...
All jobs processed. PASS
```

Elapsed time: 0.2098653s