

A: Query Containment

1. (a) There is a homomorphism from q' to q : $y \rightarrow x, z \rightarrow z, t \rightarrow z, w \rightarrow y$

But there is no homomorphism from q to q' . Hence, $q \subseteq q'$

(b) There is a homomorphism from q' to q : $x \rightarrow x, y \rightarrow y, z \rightarrow z, u \rightarrow x, v \rightarrow y$

But there is no homomorphism q to q' . Hence, $q \subseteq q'$

(c) There is no homomorphism from q to q' nor q' to q . Hence, None

2. Given $q(x): -R(x, z_1)$, we want to prove $S_K \equiv q$

1° There is a homomorphism from S_K to q : $x \rightarrow x, z_n \rightarrow z_1$ ($n \in [1, K]$)

Hence, $q \subseteq S_K$

2° There is a homomorphism from q to S_K : $x \rightarrow x, z_1 \rightarrow z_1$

Hence, $S_K \subseteq q$

In conclusion, $S_K \equiv q$ and q is the minimal CQ since it only has one atom.

3. We want $q \subset q_2$ and $q \neq q_2$, q could only be the three cases:

1° $q = R(x, y)R(z) \dots$ } \Rightarrow there is no homomorphism from q to q_1

2° $q = R(x, y)R(z, z) \dots$

3° $q = R(x, y)R(y, t) \Rightarrow$ implying $q_1 = q$ which violates $q_1 \neq q$

Thus, such q does not exist.

4. Let $q_1 = q = R(x, 10), x < 10$
 $q_2 = q = R(x, 10), x < 10 \Rightarrow \underline{q = q_1 \cup q_2} \quad (1)$

Let $q'_1 = R(y, 10)R(9, y), y < 10$
 $q'_2 = R(y, 10)R(9, y), y = 10 \Rightarrow \underline{q' = q'_1 \cup q'_2} \quad (2)$

There is a homomorphism from q_1 to q'_1 : $x \rightarrow y \Rightarrow \underline{q'_1 \subseteq q_1} \quad (3)$

There is a homomorphism from q_2 to q'_2 : $x \rightarrow 9 \Rightarrow \underline{q'_2 \subseteq q_2} \quad (4)$

By (1), (2), (3), (4), $q' \subseteq q$

B. Query Complexity

1. (a) Given $|q|$ is the size of query, $|I|$ is the size of dataset.

(b) Since we can join the variables and atoms one by one $\Rightarrow \underline{O(|q| \cdot |I|)}$
the combined time complexity is polynomial

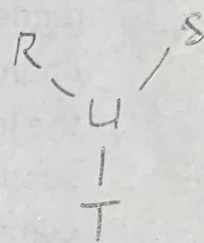
2. We use GYO to check

(a) choose ear R , remaining $\{\{y, t_2, z\}, \{z, t_3, x\}, \{x, y, z\}\}$

choose ear S , remaining $\{\{z, t_3, x\}, \{x, y, z\}\}$

choose ear T , remaining $\{\{x, y, z\}\}$

choose ear U , remaining $\{\} \Rightarrow \underline{\text{acyclic}}$



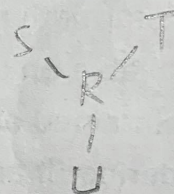
(b) Not acyclic since we can not choose ears for reduction

(c) choose ear S , remaining $\{\{x, y, z\}, \{y, z\}, \{z, x\}\}$

choose ear T , remaining $\{\{x, y, z\}, \{z, x\}\}$

choose ear U , remaining $\{\{x, y, z\}\}$

choose ear R , remaining $\{\} \Rightarrow \underline{\text{acyclic}}$



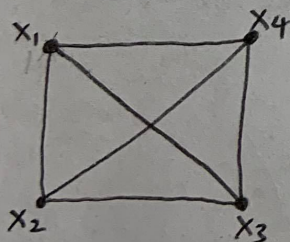
3.

(a) Choose $X(v_i) = \{x, z_i, z_i, z_{i+1}\}, i \in [1, k] \Rightarrow ghw = 3$

(b) Choose $X(v_i) = \{x_i, x_i, x_{i+1}\}$ or $\{x_i, z_i, z_{i+1}\}, i \in [1, k] \Rightarrow ghw = 2$

4. When n atoms serve the edges of a complete graph

ex. $n = 6, q = R(x_1, x_2) R(x_1, x_3) R(x_1, x_4) R(x_2, x_3) R(x_2, x_4) R(x_3, x_4)$



then there is only one way to choose X as $X(v) = \{x_1, x_2, x_3, x_4\}$
then $ghw = \# \text{vertices} - 1 = 4 - 1 = 3$

Generally, x vertices and n edges to form a complete graph

$$\frac{x(x-1)}{2} = n \Rightarrow x = \frac{1 + \sqrt{1+8n}}{2}$$

$$\underline{\underline{\text{Thus, } ghw = \frac{1 + \sqrt{1+8n}}{2} - 1}}$$