Introduction To SQL

CS 564- Fall 2018

ANNOUNCEMENTS

- Enroll in Piazza!
- PS #1 will be posted tomorrow (due next Sunday)
- Group formation:
 - send an email to Ankur with:
 - 3 x (Student IDs + emails)
 - only one person from every team!

WHAT IS THIS LECTURE ABOUT

- The Relational Model
- SQL: Basics
 - creating a table
 - primary keys
- SQL: Single-table queries
 - SELECT-FROM-WHERE structure
 - DISTINCT/ORDER BY/LIMIT
- SQL: Multi-table queries
 - foreign keys
 - joins

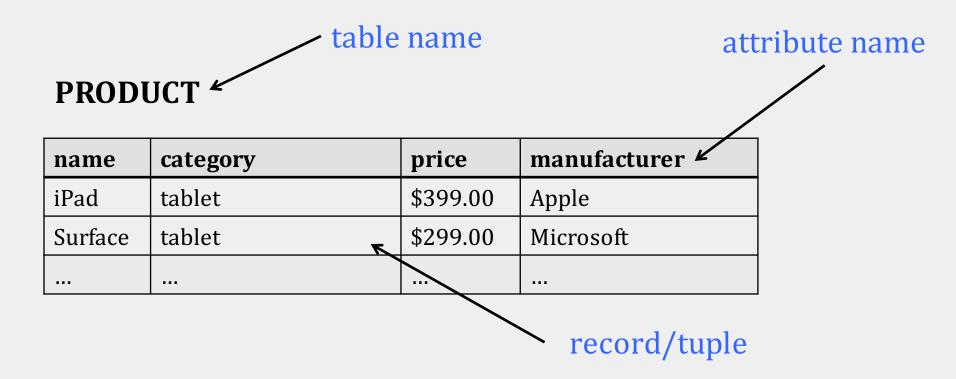
RELATIONAL MODEL

RELATIONAL MODEL

- first proposed by Codd in 1969
- has just a single concept: relation
- the world is represented as a collection of tables
- well-suited for efficient manipulations on computers

RELATION

The data is stored in **tables** (or relations)



DOMAINS

- Each attribute has an atomic type called domain
- A <u>domain</u> specifies the set of values allowed
- Examples:
 - integer
 - string
 - real

```
PRODUCT(name: string, category: string, price: real, manufacturer: string)
```

SCHEMA

The **schema** of a *relation*:

- relation name + attribute names
- Product(name, price, category, manufacturer)
- In practice we add the domain for each attribute

The **schema** of a *database*:

a collection of relation schemas

INSTANCE

The **instance** of a *relation*:

a set of tuples or records

The **instance** of a *database*:

a collection of relation instances

EXAMPLE

PRODUCT(name: string,

category: string,

price: real,

manufacturer: string)

instance

schema

name	category	price 🗸	manufacturer
iPad	tablet	\$399.00	Apple
Surface	tablet	\$299.00	Microsoft

SCHEMA VS INSTANCE

- Analogy with programming languages:
 - − schema ~ type
 - instance ∼ value
- Important distinction
 - schema: stable over long periods of time
 - instance: changes constantly, as data is inserted/updated/deleted

SQL: BASICS

WHAT IS SQL?

- The most widely used database language
- Used to query and manipulate data
- SQL stands for <u>S</u>tructured <u>Q</u>uery <u>L</u>anguage
 - many SQL standards: SQL-92, SQL:1999, SQL:2011
 - vendors support different subsets
 - we will discuss the common functionality

CREATING A TABLE

```
table name
       CREATE TABLE Author(
          →authorid INTEGER PRIMARY KEY,
attributes
          firstname CHAR(20),
                                      atomic types
           lastname CHAR(30),
       );
```

PRIMARY KEYS

A **primary key** is a minimal subset of attributes that is a unique identifier of tuples in a relation

- A key is an implicit constraint on which tuples can be in the relation
- In SQL we specify that an attribute is the primary key with the keyword PRIMARY KEY

UNIQUE KEYS

 We can also define a unique key: a subset of attributes that uniquely defines a row:

```
CREATE TABLE Author(
    authorid INTEGER UNIQUE,
    firstname CHAR(20));
```

There can be only one primary key, but many unique keys!

NULL VALUES

- tuples in SQL relations can have NULL as a value for one or more attributes
- The meaning depends on context:
 - missing value: e.g. we know that Greece has some population, but we don't know what it is
 - inapplicable: e.g. the value of attribute spouse for an unmarried person

NULL VALUES

When creating a table in SQL, we can assert that a particular attribute takes no **NULL** values

```
CREATE TABLE Author(
          authorid INTEGER PRIMARY KEY,
          firstname CHAR(20) NOT NULL,
          lastname CHAR(30)
     );
```

POPULATING A TABLE

• To insert a single tuple:

```
INSERT INTO <relation>
VALUES (<list of values>);
```

 We may add to the relation name a list of attributes (if we forget the order)

```
INSERT INTO Author
VALUES(001, 'Dan', 'Brown');
```

SQL: SINGLE-TABLE QUERIES

BASIC SQL QUERY

optional **SELECT** [DISTINCT] attributes one or more tables **FROM** WHERE conditions on the tables conditions of the form: Attr1 **op** Attr2

EXAMPLE

What is the population of USA?

SELECT Population

FROM Country

WHERE Code = 'USA';

SELECTION: filters the tuples of the relation

SEMANTICS

- 1. Think of a *tuple variable* ranging over each tuple of the relation mentioned in **FROM**
- 2. Check if the current tuple satisfies the WHERE clause
- 3. If so, compute the attributes or expressions of the **SELECT** clause using this tuple

* IN SELECT CLAUSES

When there is one relation in the **FROM** clause, * in the **SELECT** clause stands for "all attributes of this relation"

```
SELECT *
FROM City
WHERE Population >= '1000000'
AND CountryCode = 'USA';
```

RENAMING ATTRIBUTES

If we want the output schema to have different attribute names, we can use **AS** < new name > to rename an attribute

```
SELECT Name AS LargeUSACity
FROM City
WHERE Population >= '1000000'
AND CountryCode = 'USA';
```

ARITHMETIC EXPRESSIONS

We can use any arithmetic expression (that makes sense) in the **SELECT** clause

```
SELECT Name,
  (Population/ 1000000) AS PopulationInMillion
FROM City
WHERE Population >= '1000000';
```

WHAT CAN WE USE IN WHERE CLAUSES?

- attribute names of the relations that appear in the FROM clause
- comparison operators: =, <>, <, >, <=, >=
- arithmetic operations (+, -, /, *)
- AND, OR, NOT to combine conditions
- operations on strings (e.g. concatenation)
- pattern matching: s LIKE p
- special functions for comparing dates and times

PATTERN MATCHING

s **LIKE** p: pattern matching on strings

- % = any sequence of characters
- _ = any single character

```
SELECT Name, GovernmentForm
FROM Country
WHERE GovernmentForm LIKE '%Monarchy%';
```

USING DISTINCT

- The default semantics of SQL is **bag** semantics (duplicate tuples are allowed in the output)
- The use of **DISTINCT** in the **SELECT** clause removes all duplicate tuples in the result, and returns a **set**

SELECT DISTINCT GovernmentForm **FROM** Country;

ORDER BY

The use of **ORDER BY** orders the tuples by the attribute we specify in decreasing (**DESC**) or increasing (**ASC**) order

```
SELECT Name, Population
FROM City
WHERE Population >= '1000000'
ORDER BY Population DESC;
```

LIMIT

- The use of **LIMIT** < number > limits the output to be only the specified number of tuples
- It can be used with ORDER BY to get the maximum or minimum value of an attribute!

```
SELECT Name, Population
FROM City
ORDER BY Population DESC
LIMIT 2;
```

SQL: MULTI-TABLE QUERIES

FOREIGN KEYS

Suppose that we want to create a table Book, and make sure that the author of the book exists in the table Author

```
CREATE TABLE Book(
   bookid INTEGER PRIMARY KEY,
   title TEXT,
   authorid INTEGER,
   FOREIGN KEY (authorid) REFERENCES
   Author(authorid));
```

FOREIGN KEYS

• Use the keyword **REFERENCES**, as:

```
FOREIGN KEY ( < list of attributes > )
REFERENCES < relation > ( < attributes > )
```

 Referenced attributes must be declared PRIMARY KEY or UNIQUE

ENFORCING FK CONSTRAINTS

If there is a foreign-key constraint from attributes of relation \mathbf{R} to the primary key of relation \mathbf{S} , two violations are possible:

- 1. An insert or update to *R* introduces values not found in *S*
- 2. A deletion or update to *S* causes some tuples of *R* to dangle

There are 3 ways to enforce foreign key constraints!

ACTION 1: REJECT

- The insertion/deletion/update query is rejected and not executed in the DBMS
- This is the default action if a foreign key constraint is declared

ACTION 2: CASCADE UPDATE

When a tuple referenced is *updated*, the update **propagates** to the tuples that reference it

```
CREATE TABLE Book(
   bookid INTEGER PRIMARY KEY,
   title TEXT,
   authorid INTEGER,
   FOREIGN KEY (authorid) REFERENCES
   Author(authorid)
   ON UPDATE CASCADE);
```

ACTION 2: CASCADE DELETE

When a tuple referenced is *deleted*, the deletion **propagates** to the tuples that reference it

```
CREATE TABLE Book(
   bookid INTEGER PRIMARY KEY,
   title TEXT,
   authorid INTEGER,
   FOREIGN KEY (authorid) REFERENCES
   Author(authorid)
   ON DELETE CASCADE);
```

ACTION 3: SET NULL

 When a delete/update occurs, the values that reference the deleted tuple are set to NULL

```
CREATE TABLE Book(
   bookid INTEGER PRIMARY KEY,
   title TEXT,
   authorid INTEGER,
   FOREIGN KEY (authorid) REFERENCES
   Author(authorid)
   ON UPDATE SET NULL);
```

WHAT SHOULD WE CHOOSE?

 When we declare a foreign key, we may choose policies SET NULL or CASCADE independently for deletions and updates

ON [UPDATE, DELETE] [SET NULL, CASCADE]

• Otherwise, the default policy (reject) is used

MULTIPLE RELATIONS

- We often want to combine data from more than one relation
- We can address several relations in one query by listing them all in the FROM clause
- If two attributes from different relations have the same name, we can distinguish them by writing <relation>.<attribute>

EXAMPLE

What is the name of countries that speak Greek?

```
FROM Country, CountryLanguage
WHERE Code = CountryCode
AND Language = 'Greek';
```

This is **BAD** style!!

EXAMPLE: GOOD STYLE

```
SELECT Country.Name
FROM Country, CountryLanguage
WHERE Country.Code=CountryLanguage.CountryCode
AND CountryLanguage.Language = 'Greek';
```

```
FROM Country C, CountryLanguage L
WHERE C.Code = L.CountryCode
AND L.Language = 'Greek';
```

VARIABLES

Variables are necessary when we want to use two copies of the same relation in the **FROM** clause

```
FROM Country C, CountryLanguage L1,
CountryLanguage L2
WHERE C.Code = L1.CountryCode
   AND C.Code = L2.CountryCode
   AND L1.Language = 'Greek'
   AND L2.Language = 'English';
```

SEMANTICS: SELECT-FROM-WHERE

- 1. Start with the cross product of all the relations in the **FROM** clause
- 2. Apply the conditions from the WHERE clause
- Project onto the list of attributes and expressions in the SELECT clause
- 4. If **DISTINCT** is specified, eliminate duplicate rows

SEMANTICS OF SQL: EXAMPLE

SELECT R.D FROM R, S WHERE R.A = S.B AND S.C = 'e';

A	D
1	a
2	b
2	С

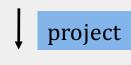
cross product

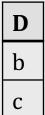
В	С
1	d
2	e

A	D	В	С
1	a	1	d
1	a	2	e
2	b	1	d
2	b	2	e
2	С	1	d
2	С	2	e

select

A	D	В	С
2	b	2	e
2	С	2	e





SEMANTICS OF SQL: NESTED LOOP

```
SELECT a_1, a_2, ..., a_k
          R_1 AS X_1, R_2 AS X_2, ..., R_n AS X_n
FROM
WHERE Conditions
answer := {}
for x_1 in R_1 do
   for x_2 in R_2 do
        for x_n in R_n do
            if Conditions
                then answer := answer \cup \{(a_1,...,a_k)\}
return answer
```

SEMANTICS OF SQL

- The query processor will almost never evaluate the query this way
- SQL is a declarative language
- The DBMS figures out the most efficient way to compute it (we will discuss this later in the course when we talk about query optimization)