Heaps



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Announcements

Please fill out the teacher evaluation form, thank you!

- Smartphone: www.hunter.cuny.edu/mobilete
- Computer: www.hunter.cuny.edu/te

Announcements

Focus group rescheduled, please see Blackboard announcement to register

Fill out paper survey to get credit for the workshops!

<u>Announcements</u>

- Final Exam MONDAY May 20, 9-11am
- Review Tuesday May 14
 - Student Led
 - Come with questions

Heap

A Heap is a complete binary tree that is either

- Empty or
- Its root contains a value ≥ (or ≤) both of its children and has heaps as subtrees

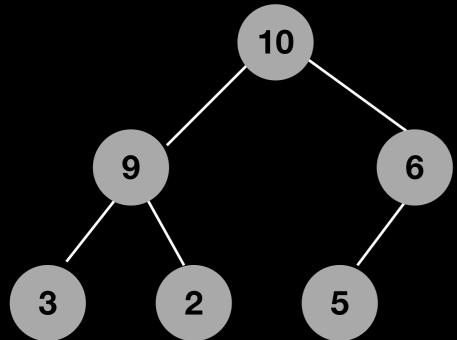
Heap

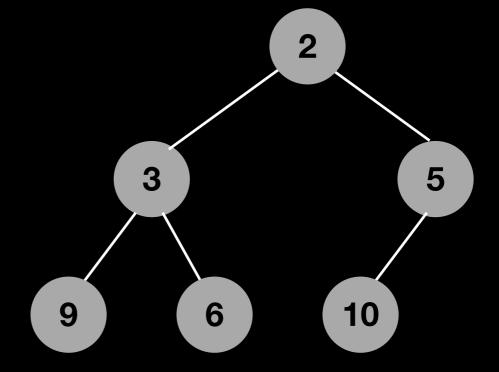
A special binary tree:

- Ordered in a <u>weaker sense</u>
- Always a **complete** binary tree

MaxHeap



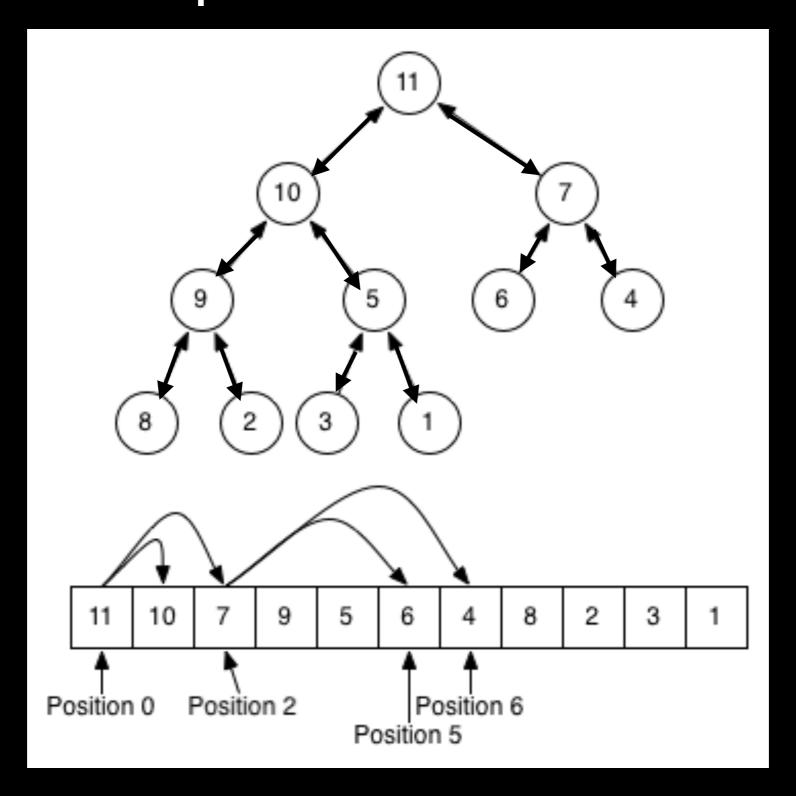


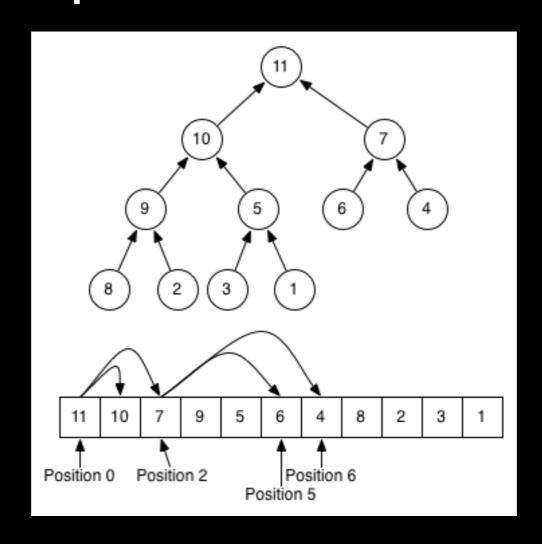


How would you implement it???

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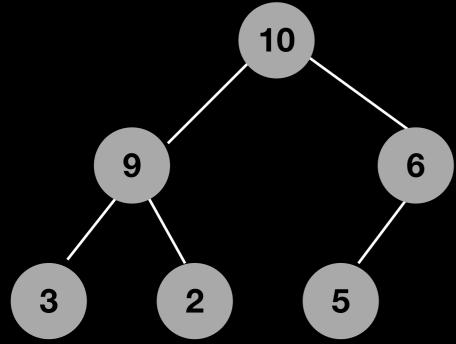
Insight: it is always complete





```
root_ = items_[0]
items_[i] left_child = items_[2 * i + 1]
items_[i] right_child = items_[ 2 * i + 2]
items_[i] parent = items_[(i-1)//2]
```

MaxHeap



Priority Queue

10 ? ? 5

Retrieve

Can only retrieve max/min item

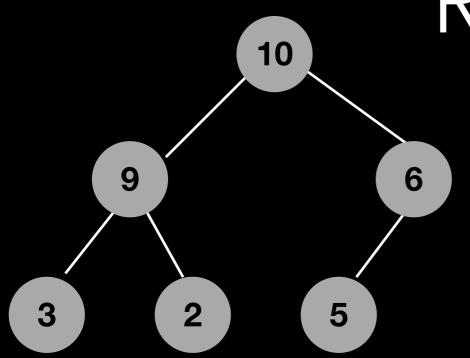
Stored at root

O(1)

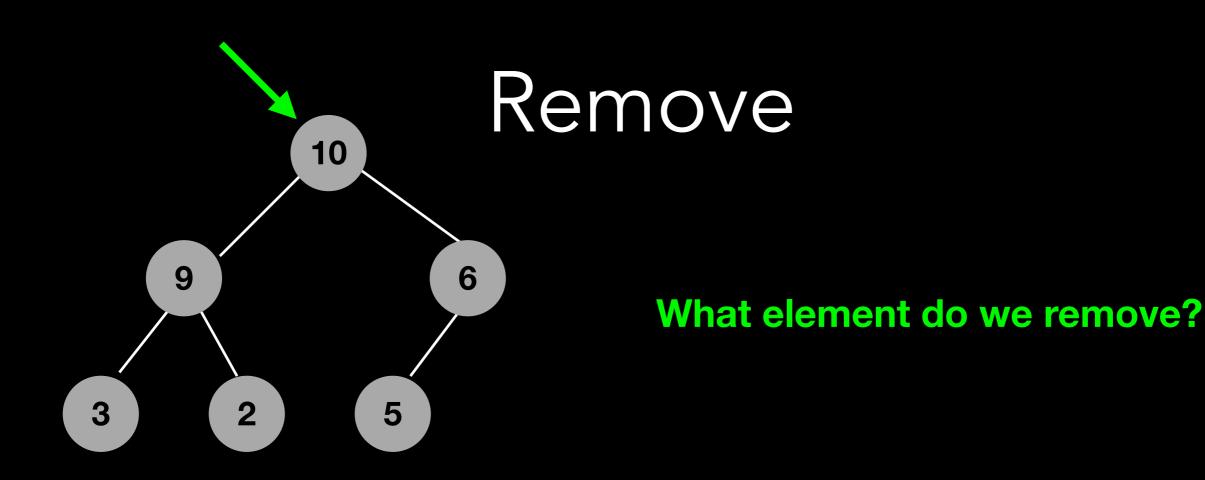
Remove max/min item (the root)

Must retain Heap

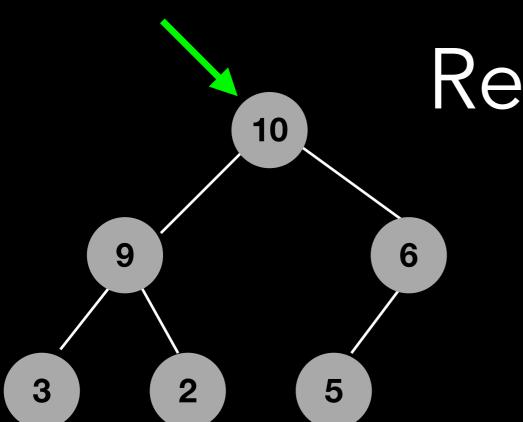
- Heap ordering property
- Complete



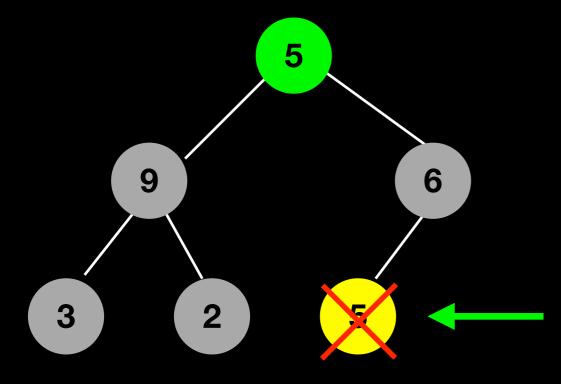
What element do we remove?



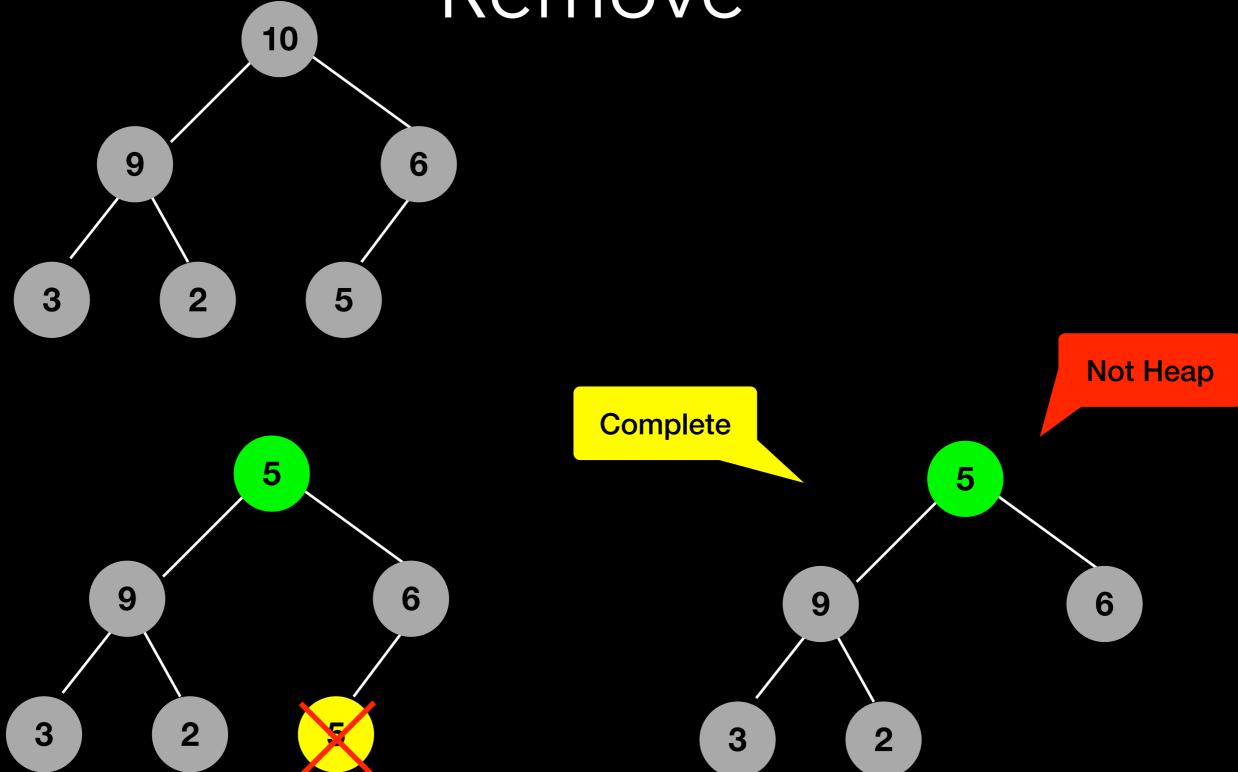
What node do we remove?



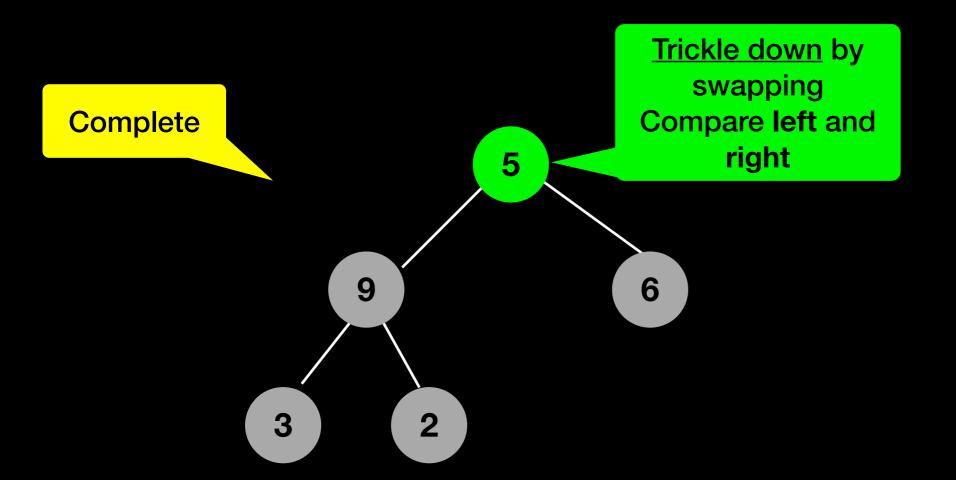
What element do we remove?



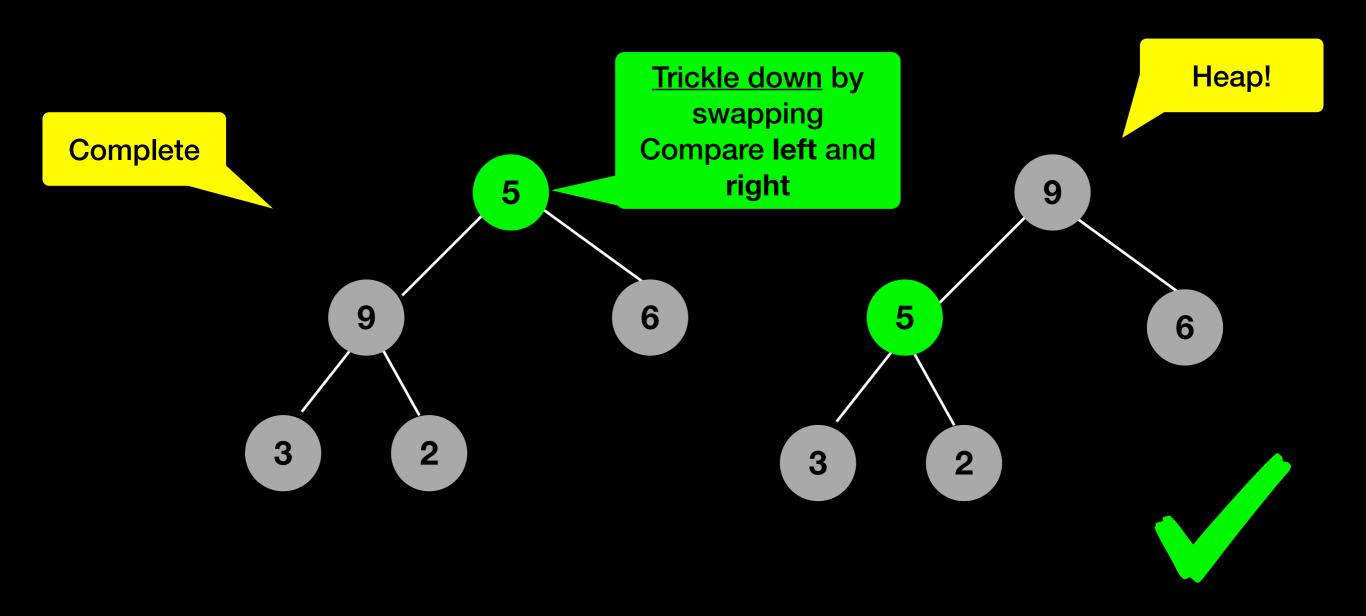
Remove this node form complete tree

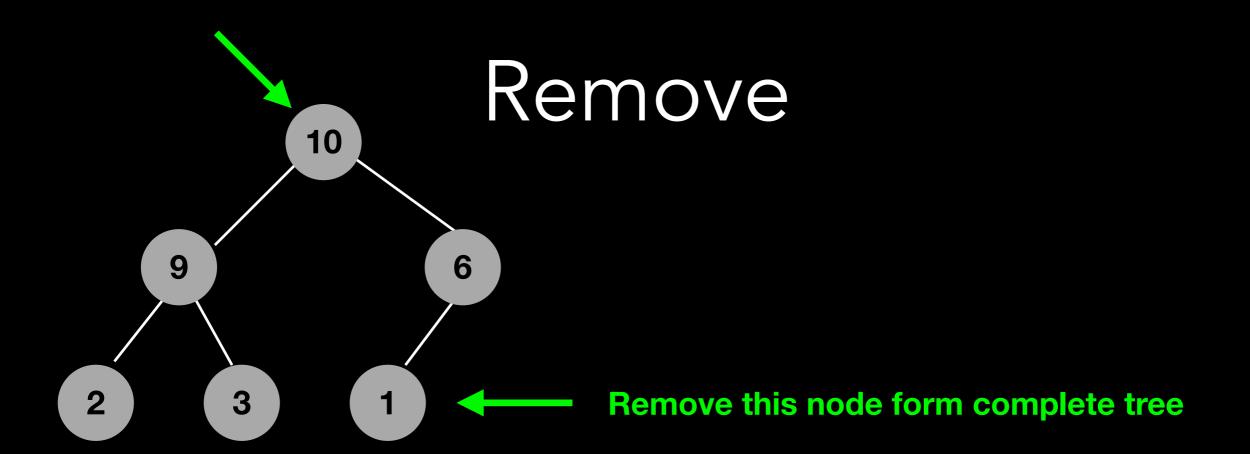


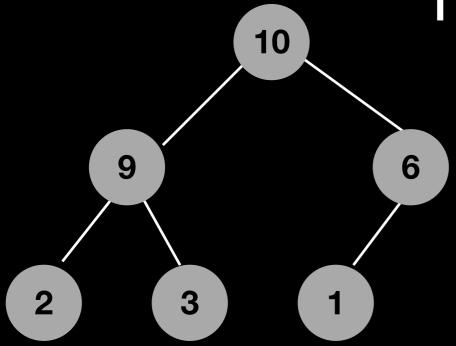
heapRebuild

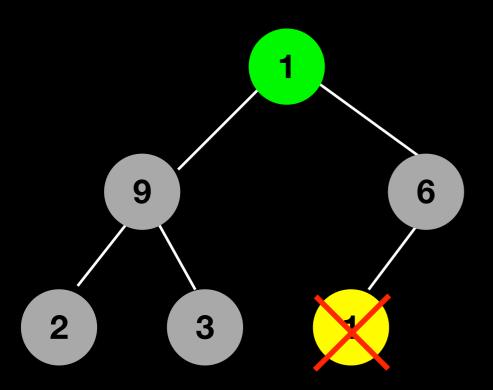


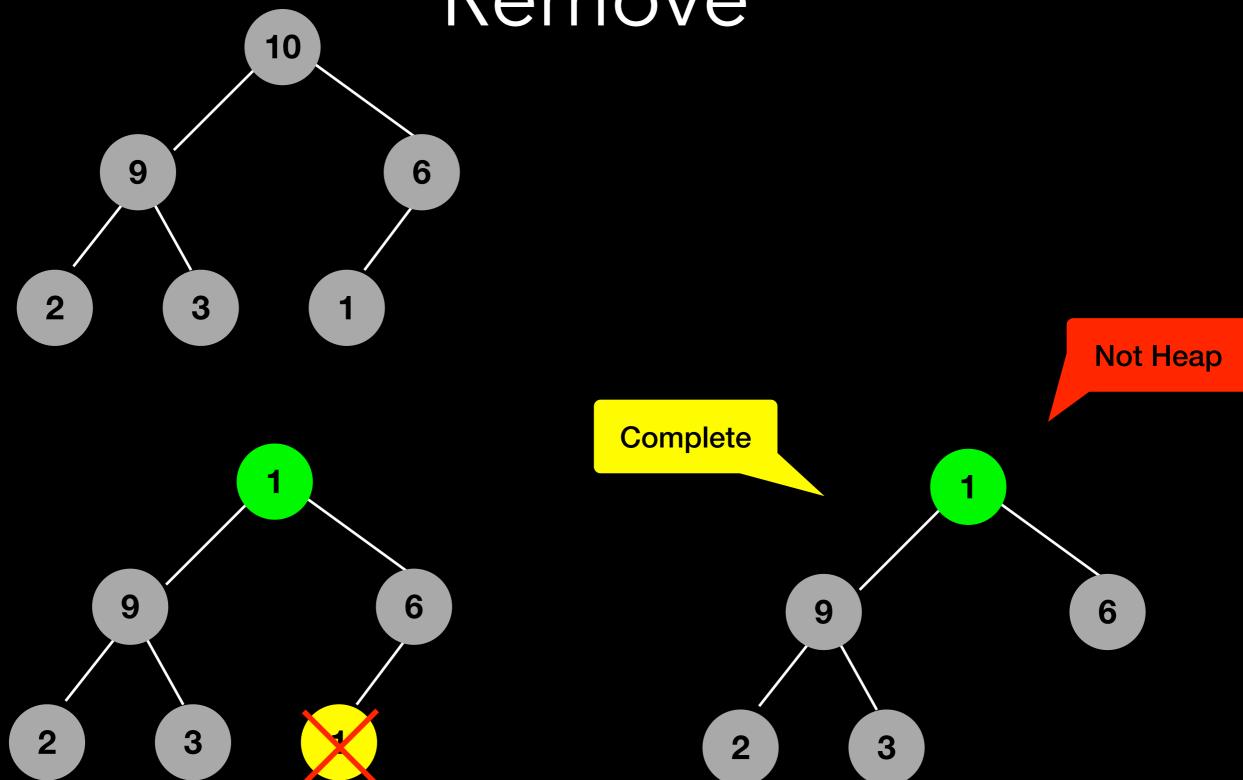
heapRebuild

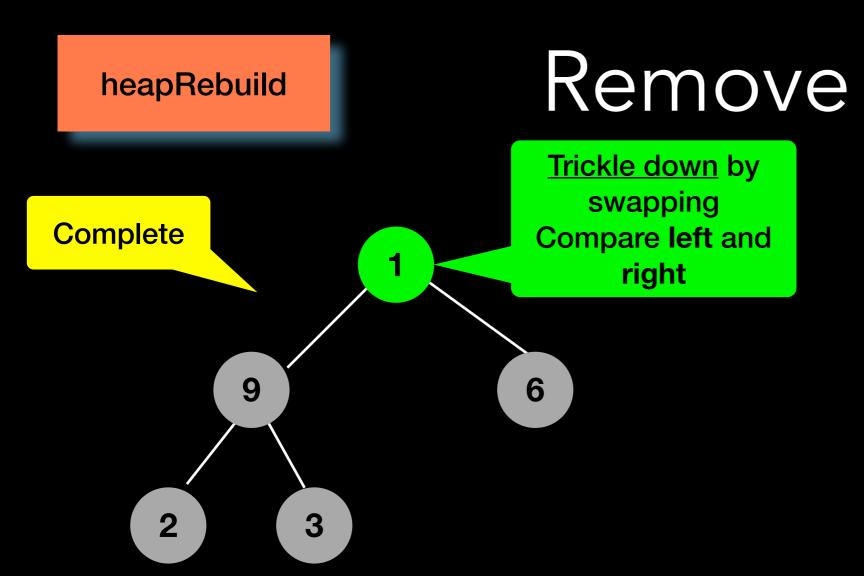






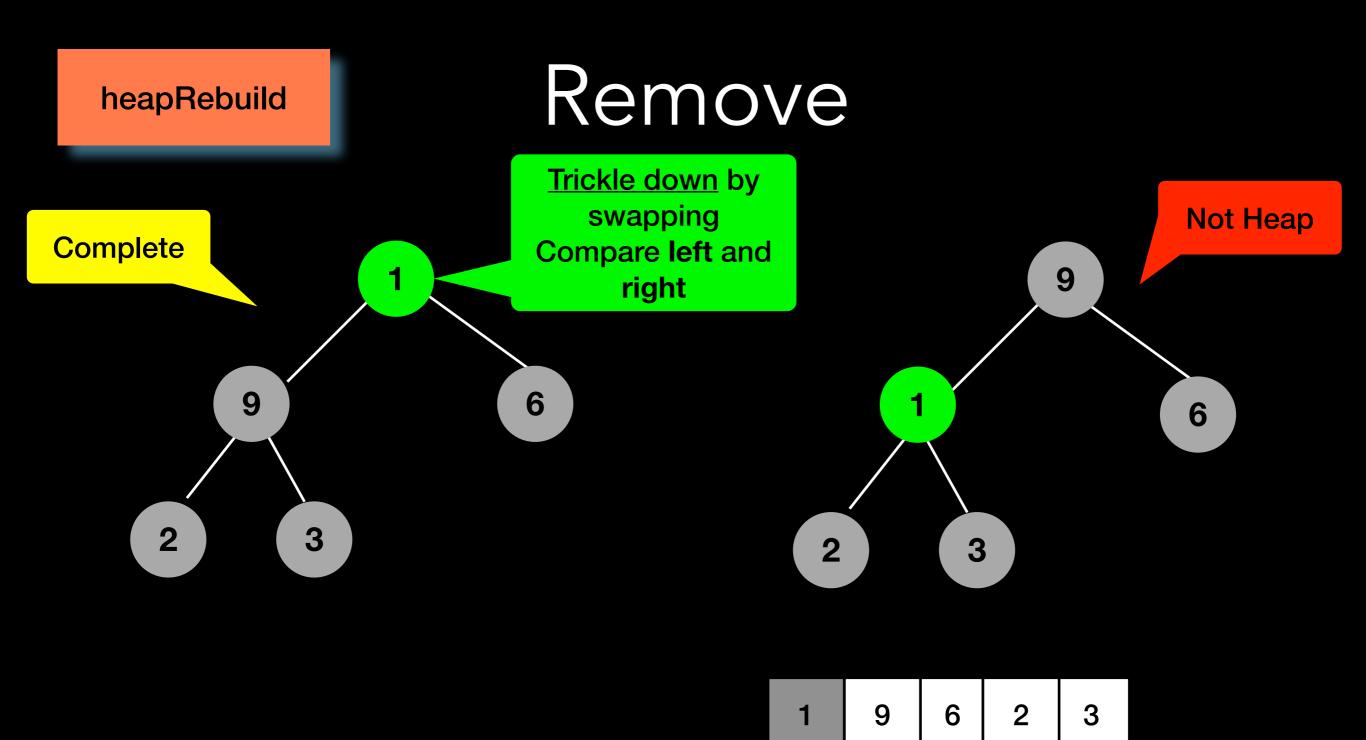




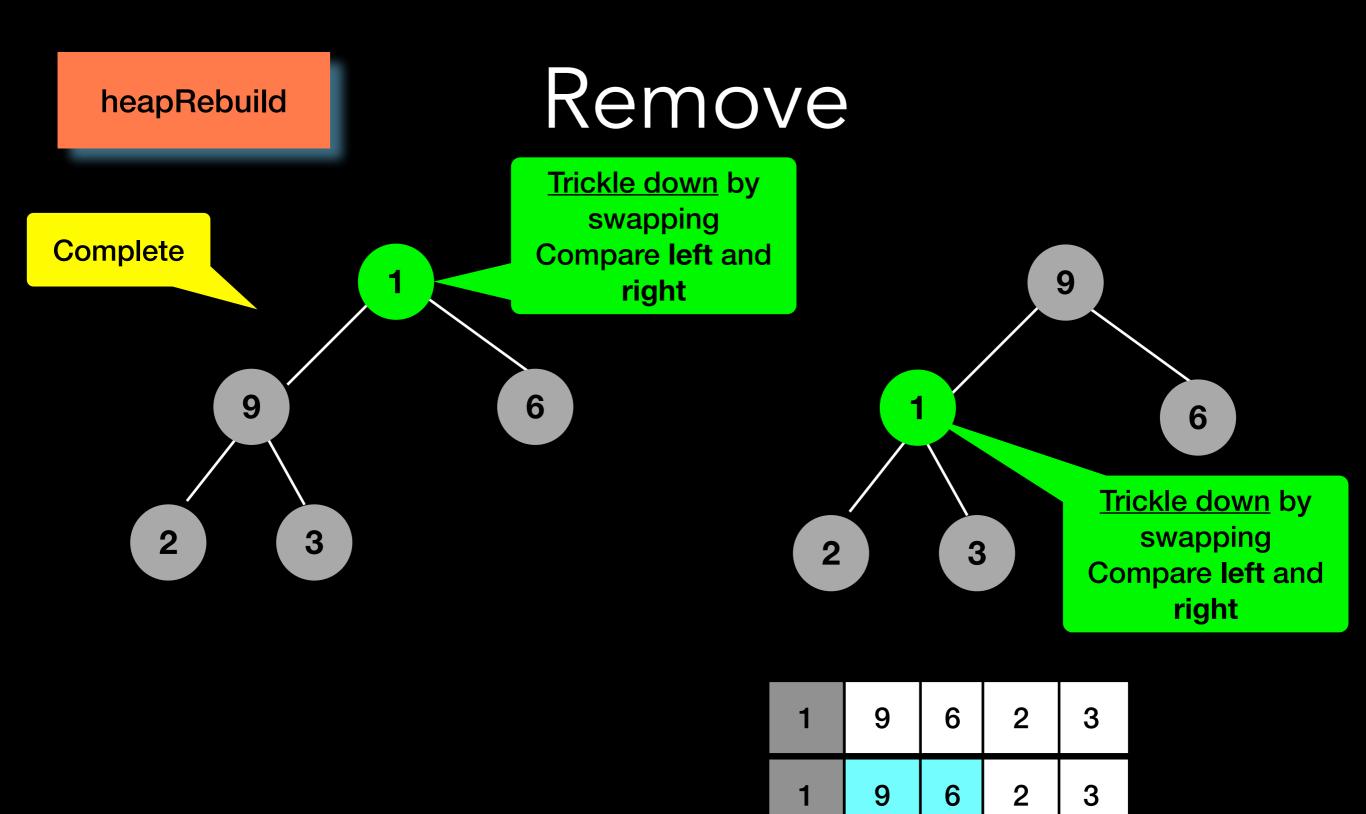


```
1 9 6 2 3
```

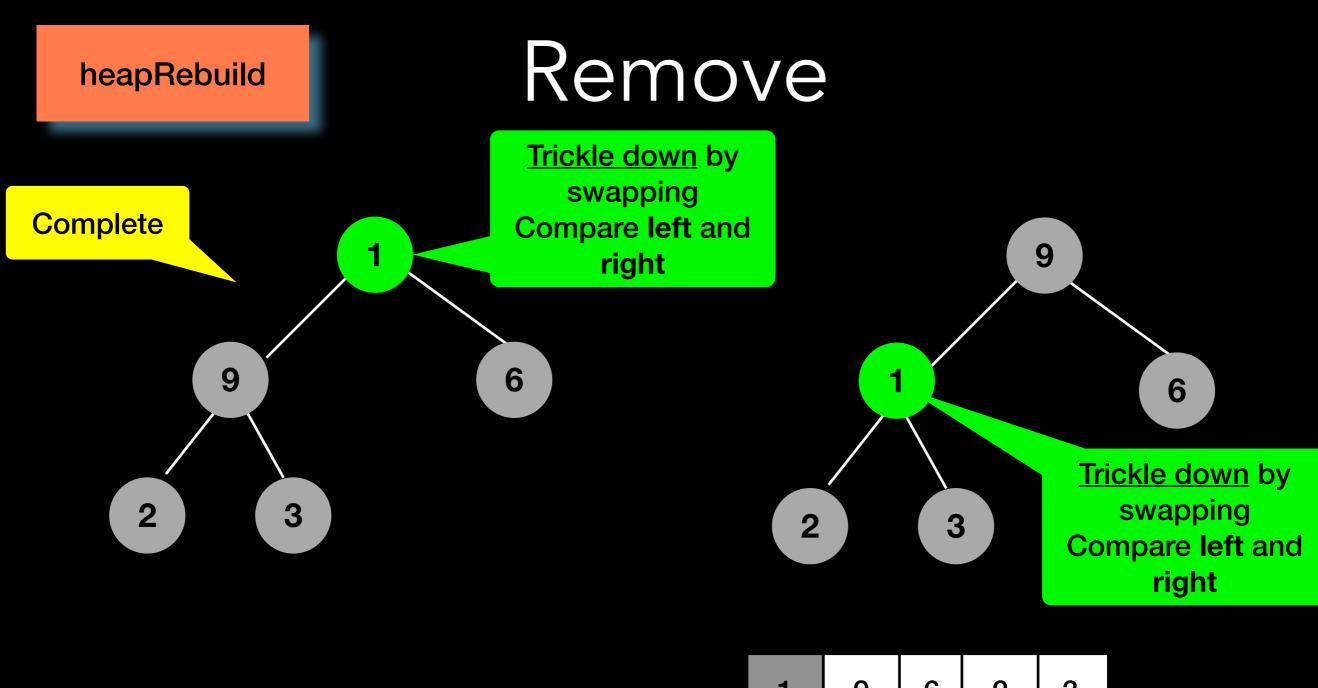
```
items_[i] left_child = items_[2 * i + 1]
items_[i] right_child = items_[ 2 * i + 2]
```



items_[i]	<pre>left_child = items_[2</pre>	*	i	+ 1]
<pre>items_[i]</pre>	right_child = items_[2	*	i + 2]



```
items_[i] left_child = items_[2 * i + 1]
items_[i] right_child = items_[ 2 * i + 2]
```

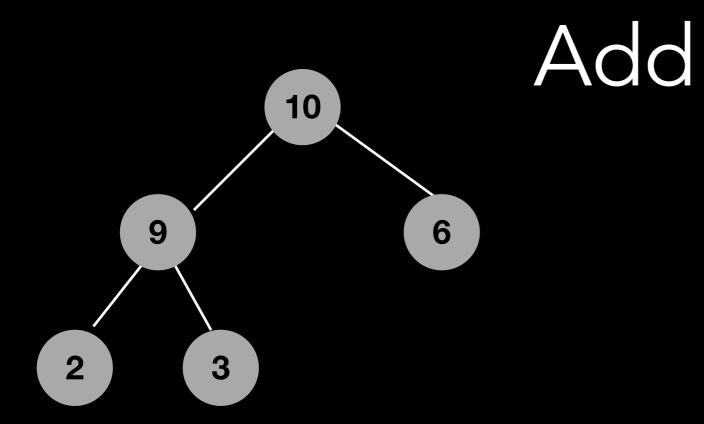


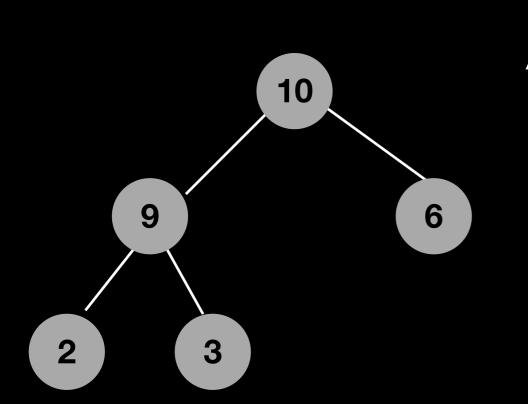
1	9	6	2	3
1	9	6	2	3
9	1	6	2	3

```
items_[i] left_child = items_[2 * i + 1]
items_[i] right_child = items_[ 2 * i + 2]
```

Remove heapRebuild Trickle down by swapping Complete Compare left and 9 right Trickle down by swapping Compare left and right 9 Heap! 9 6 2 3 9 6 2 9 6 3 9 3 6 2 29

Remove heapRebuild Trickle down by swapping Complete Compare left and 9 right Trickle down by swapping Compare left and right 9 Heap! 9 6 2 3 9 6 6 9 3 O(logn) 9 3 6 30

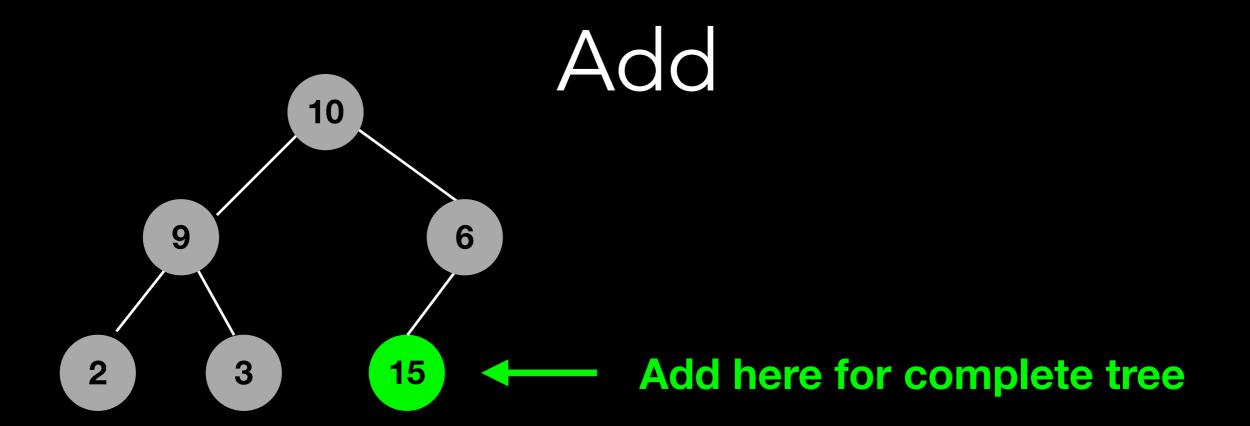




Add

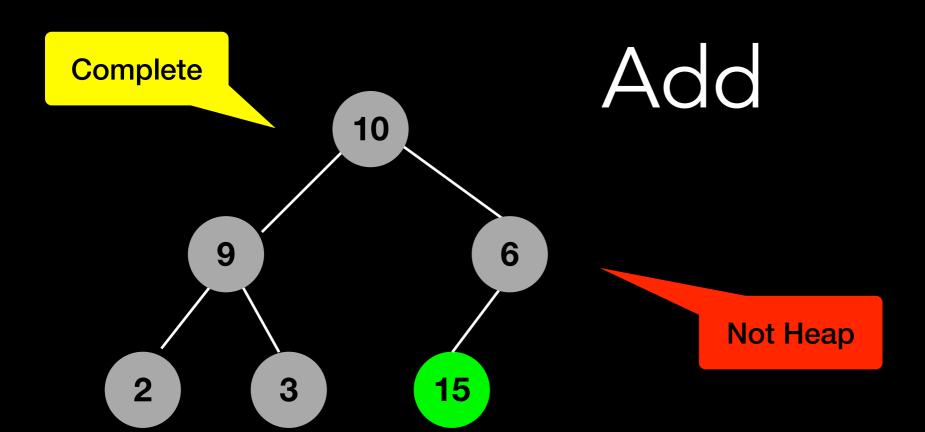
Where do we add?

10 9 6 2 3

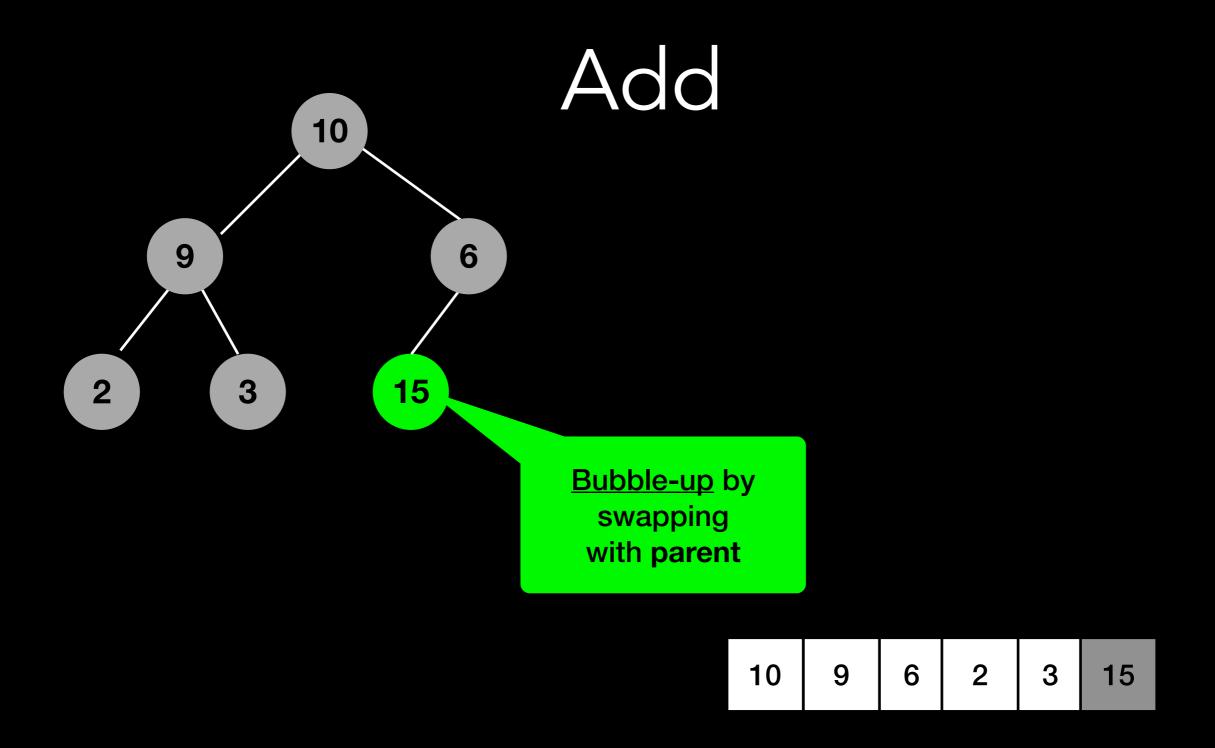


10 9 6 2 3 15

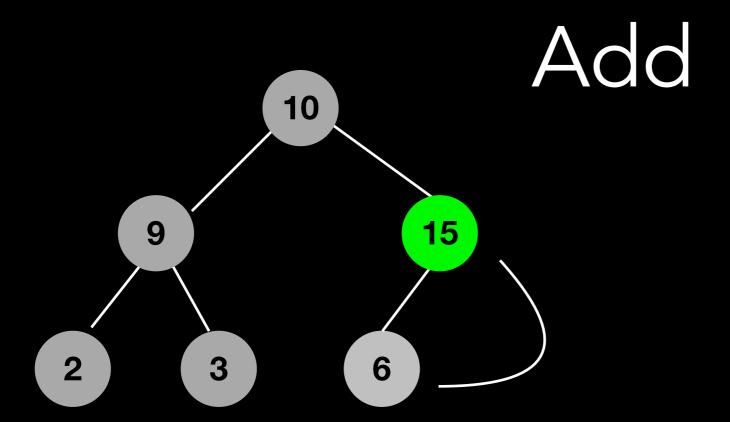
items_[i] left_child = items_[2 * i + 1]
33



10 9 6 2 3 15

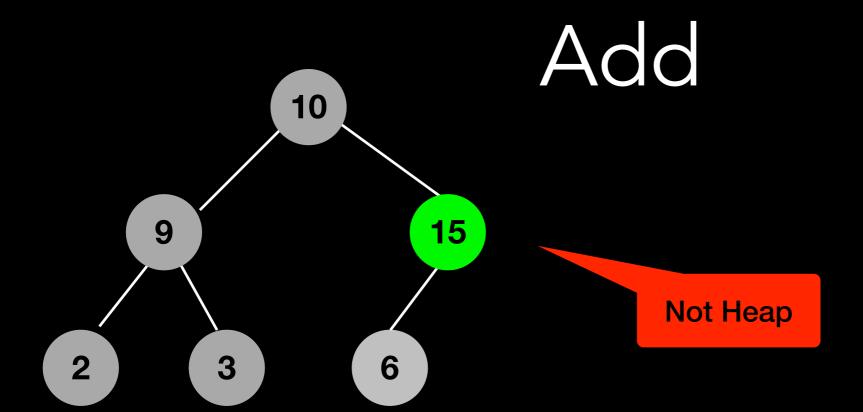


items_[i] parent = items_[(i-1)//2]

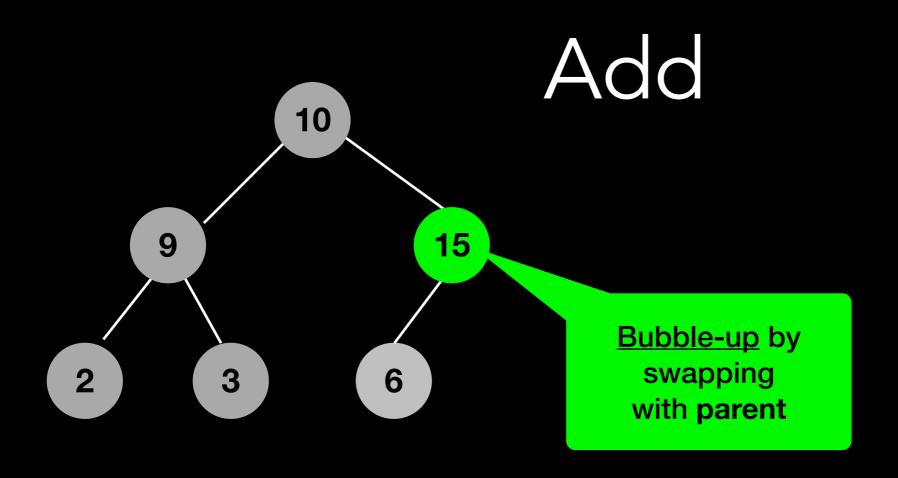


10	9	6	2	3	15
10	9	15	2	3	6

items_[i] parent = items_[(i-1)//2]



10	9	6	2	3	15
10	9	15	2	3	6

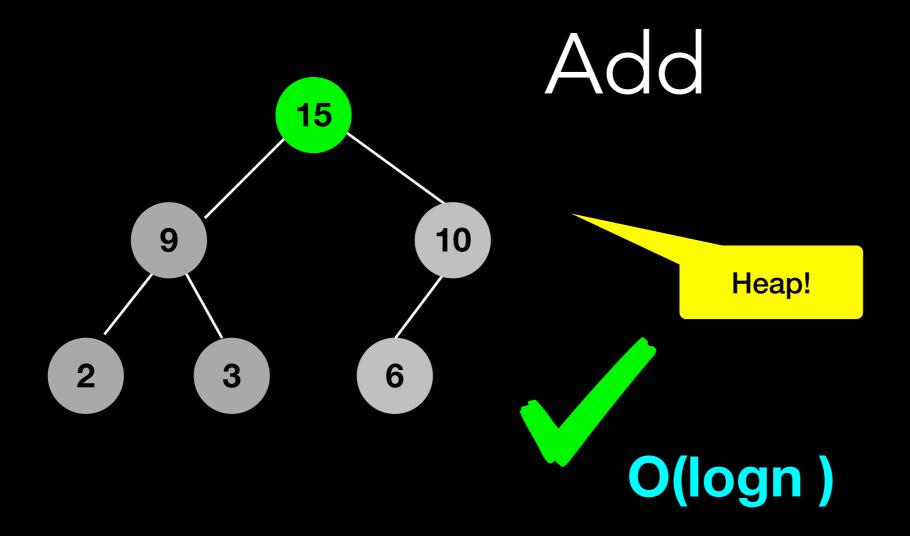


10	9	6	2	3	15
10	9	15	2	3	6

9 10

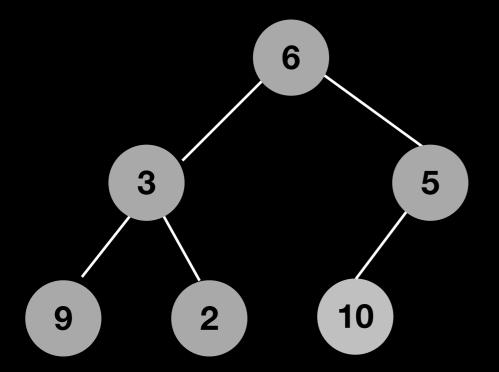
Λ	
Δ	

10	9	6	2	3	15
10	9	15	2	3	6
15	9	10	2	3	6



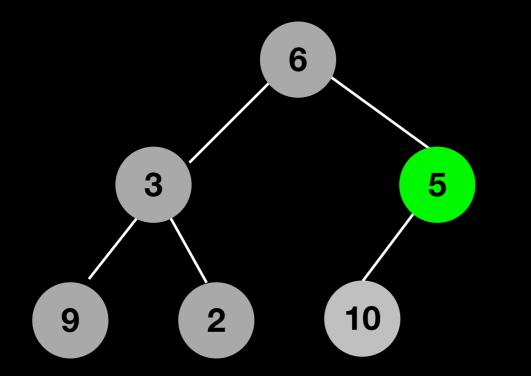
10	9	6	2	3	15
10	9	15	2	3	6
15	9	10	2	3	6

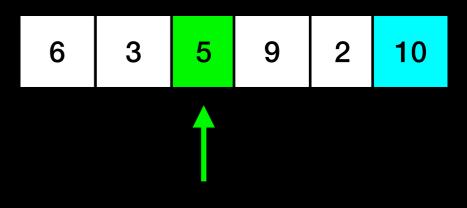
6 3 5 9 2 10



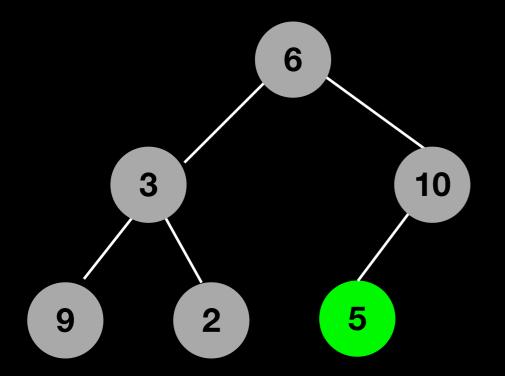
```
6 3 5 9 2 10
```

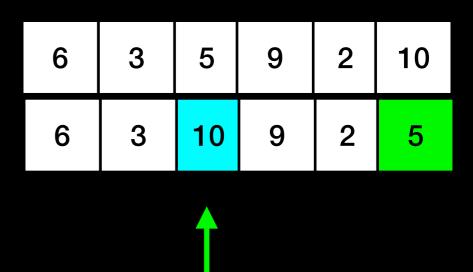
```
for(int i=(itemCount//2)-1; i >=0; i--)
{
    heapRebuild(i);
}
```



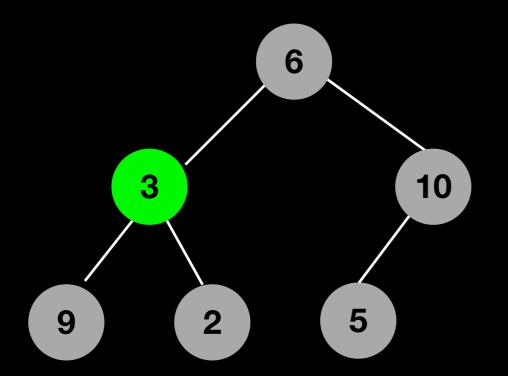


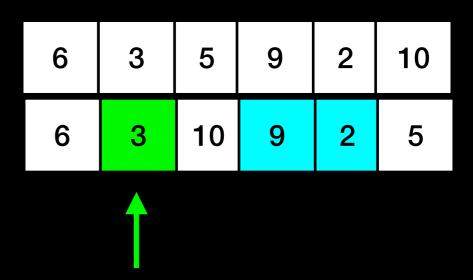
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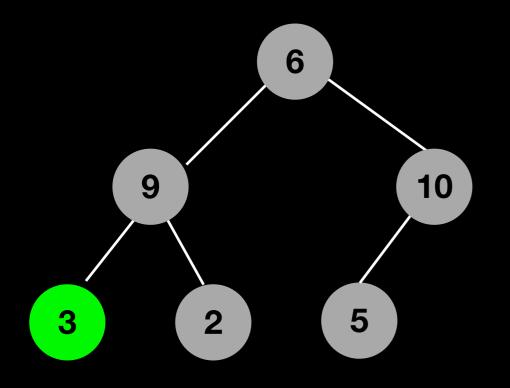


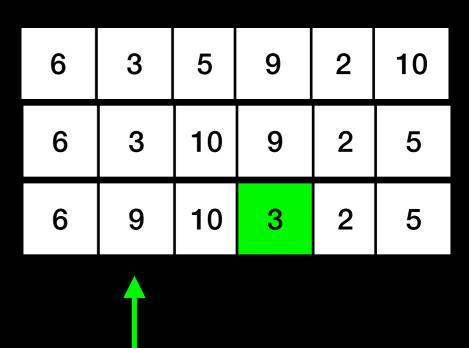
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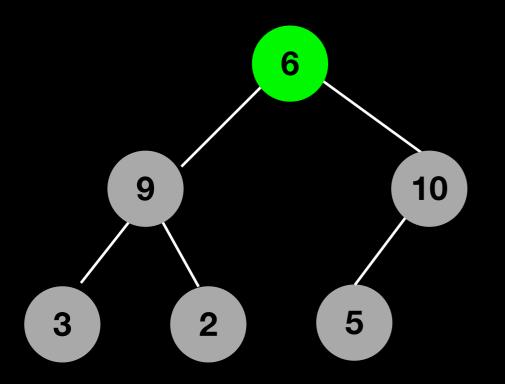


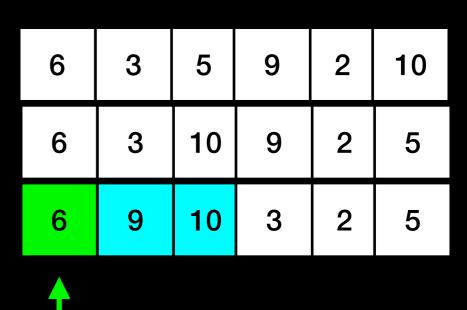
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```



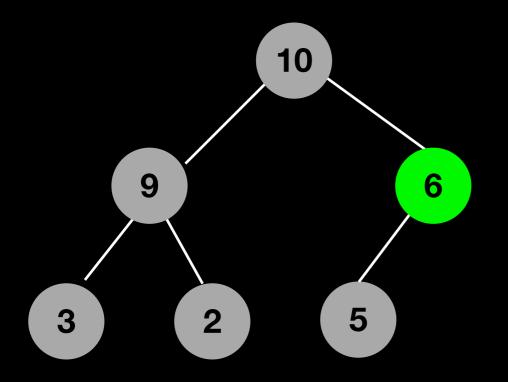


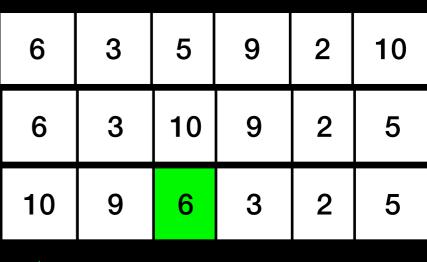
```
for(int i=(itemCount/2)-1; i >=0; i--)
{
    heapRebuild(i);
}
```





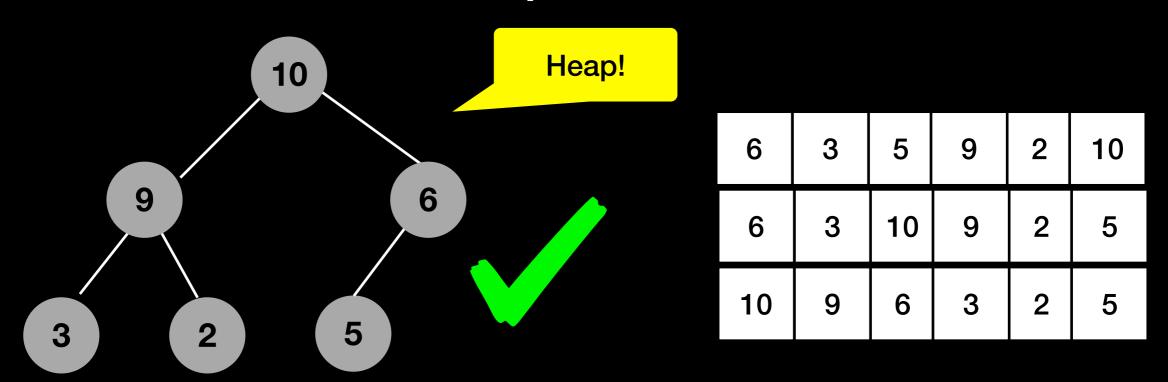
```
for(int i=(itemCount/2)-1; i >=0; i--)
{
    heapRebuild(i);
}
```







```
for(int i=(itemCount/2)-1; i >=0; i--)
{
   heapRebuild(i);
}
```

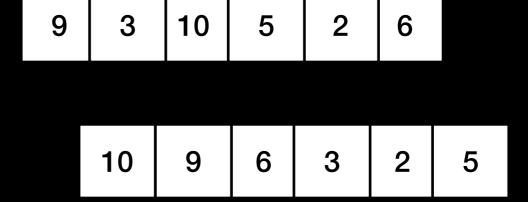


n/2 heapRebuild

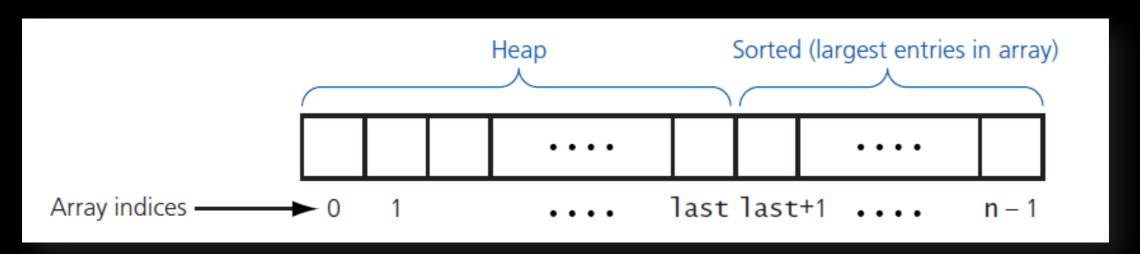
```
for(int i=(itemCount/2)-1; i >=0; i---)
{
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}
```

Given an unsorted array:

- heapCreate
- last = n 2
- repeat:



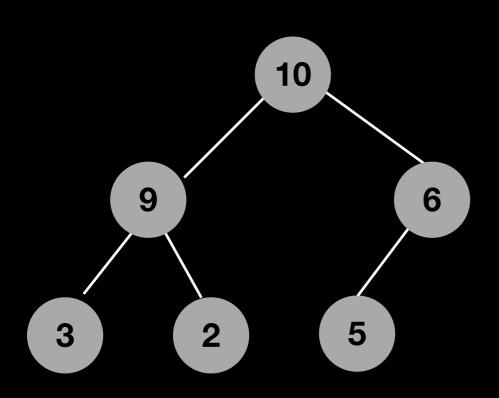
- swap items[0] with items[last+1]
- last--
- heapRebuild



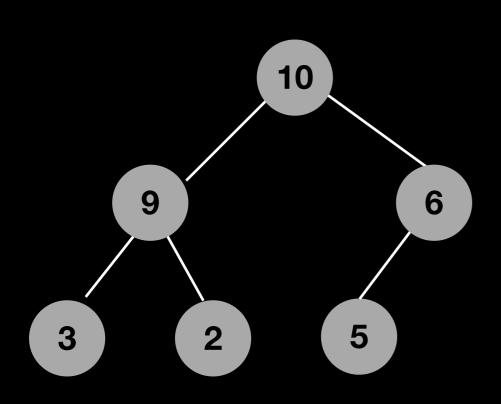
Given an unsorted array:

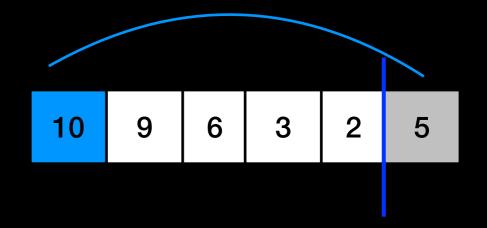
- heapCreate
- last = n 2
- repeat:
 - swap items[0] with items[last+1]
 - last--
 - heapRebuild

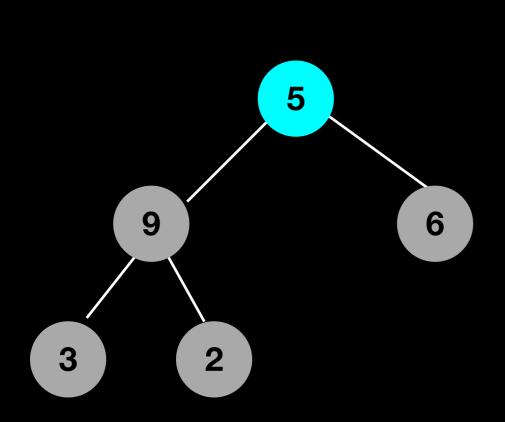
Array indices — 0 1 last last+1 n-1

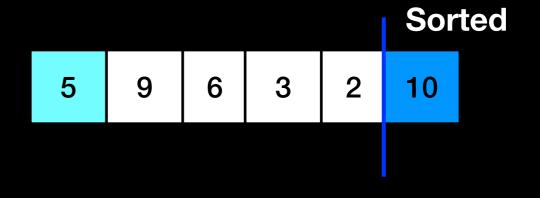




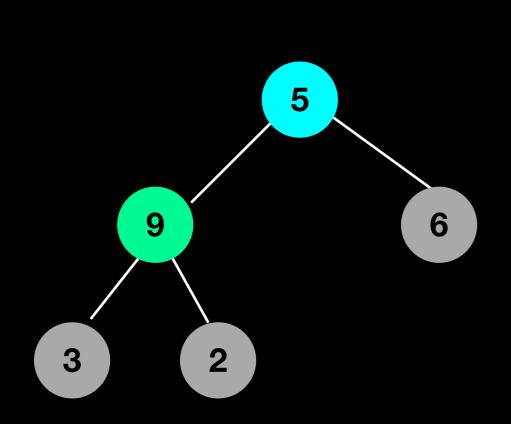


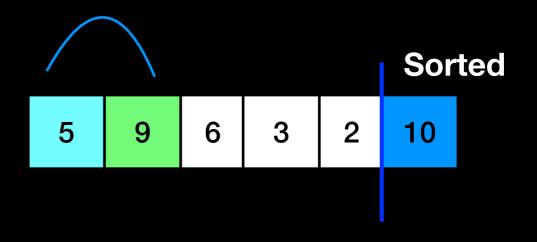


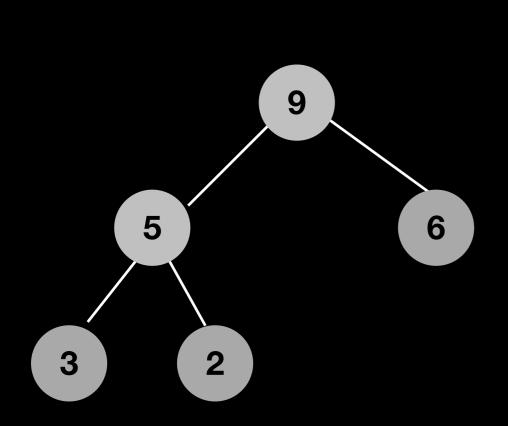


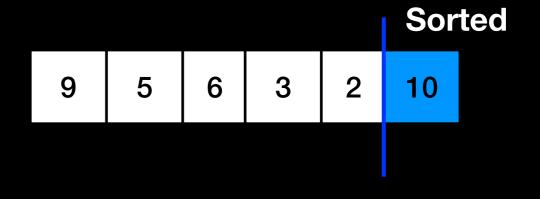


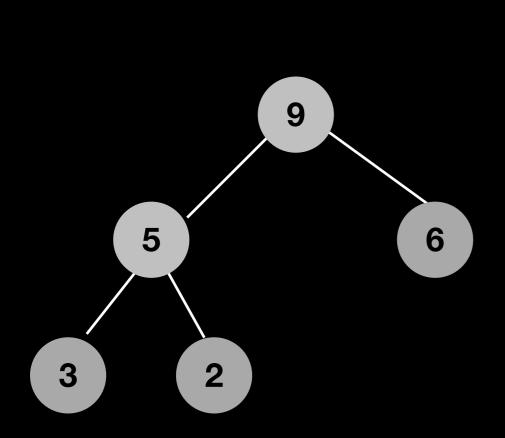
heapRebuild

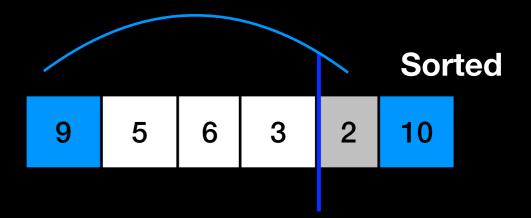




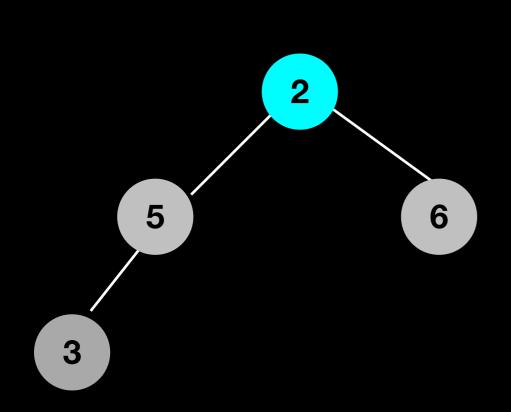


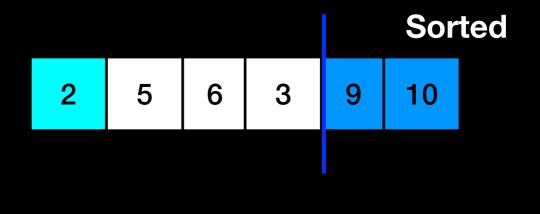




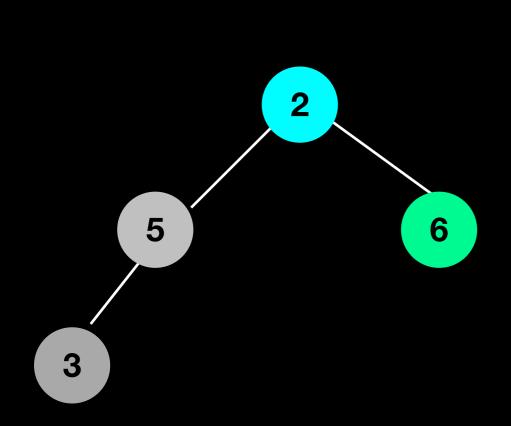


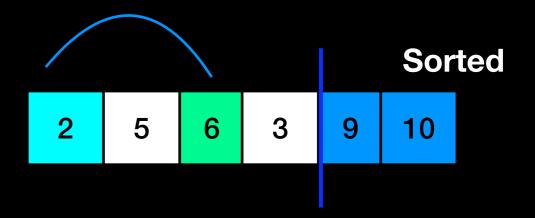
heapRebuild

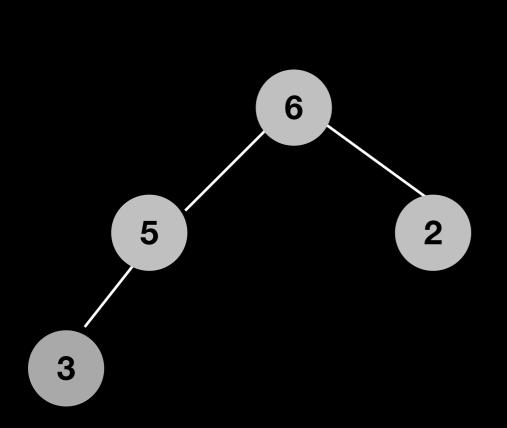


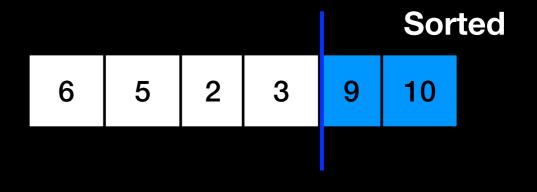


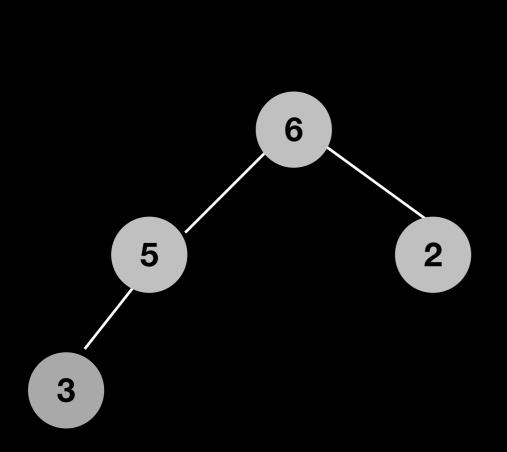
heapRebuild

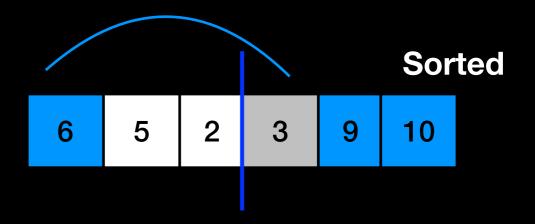




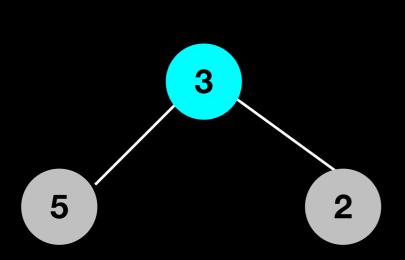


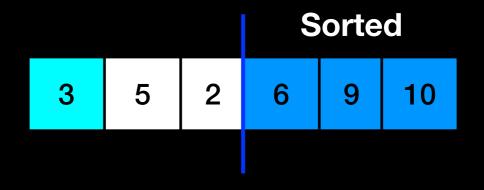




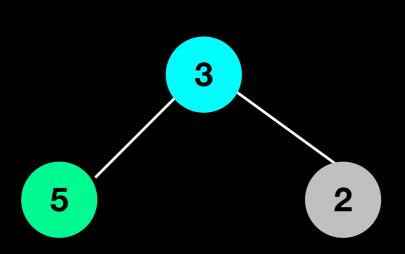


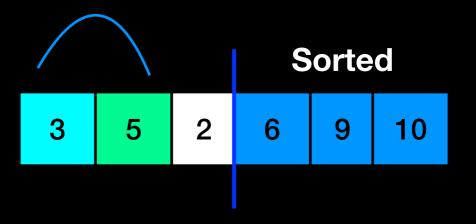
heapRebuild

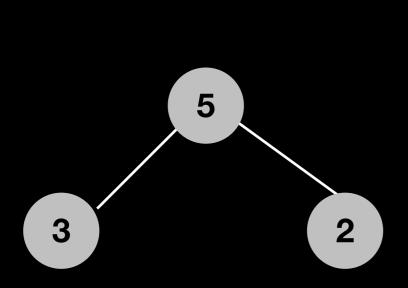


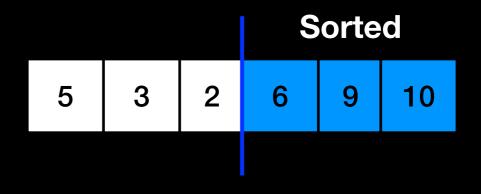


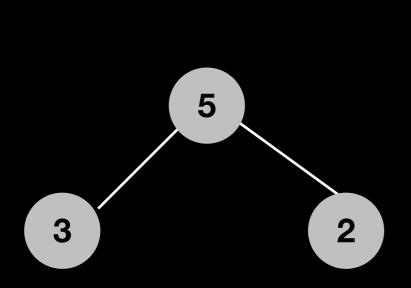
heapRebuild

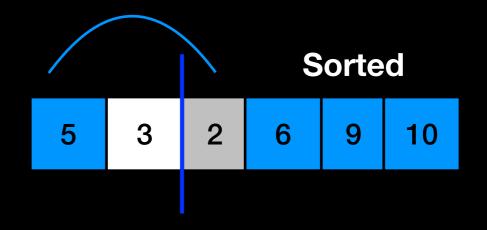




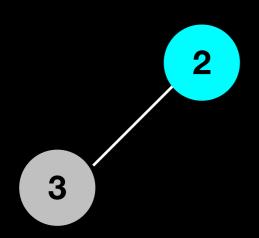


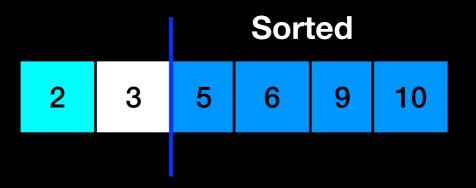




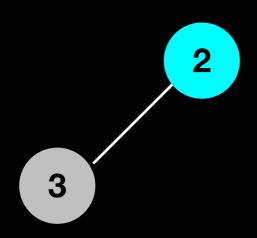


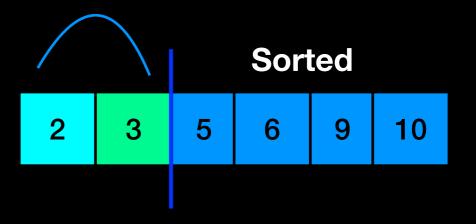
heapRebuild

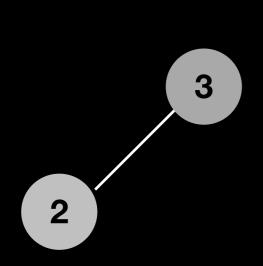


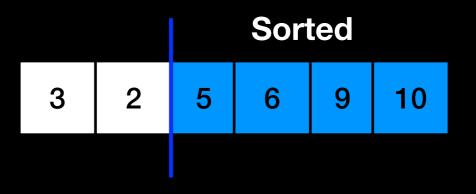


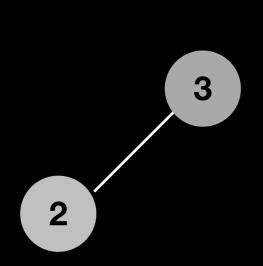
heapRebuild

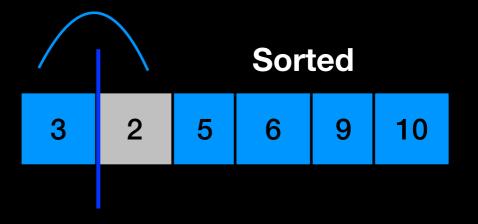














Sorted



Heapsort Analysis

- 1. heapCreate -> O(n logn)
- 2. heapRebuild -> O(logn) repeated for each of the n sorted items

$$O(n \log n) + O(n \log n) = O(n \log n)$$

Like MergeSort but no extra space needed!