1 Let G=(V, E) be the original graph. Suppose there are two distinct MSTs T:=(V, E) and T:=(V, E). The sets (E:-E) and (E:-E) are not empty, so there are at least one e belongs to (E:-E). Since e does not belong to E: adding it to T: areates a cycle. By this appears, the edge ethat has the most weight of this appears, the edge ethat has the most weight of this appearance e e E:-E;) if e'=e, then e'e E: Now both ages are contradicting to the foot that e' is not in any MST. Here, G has exactly are minimum spanning.

2 Given any out in a edge-weighted graph, one of the crossing edges of minimum weight is in the MST of the graph

Hoof:
Let e, f be are of the crossing edges of minimum weight and T be the MST. Suppose that T does not contain e.f. Now consider the graph formed by adding e.f to T. This graph has a aide that contains e.f. and that aide must contain at bost are crossing - say g, which has heaver weight than e. We are get a spanning tree of strickly lower weight by deleting g and adding e or f. contradicting the assumed minimality of T.



