## Formal Theory of Communication Topology in Concurrent ML

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## 1 Mathematical Artifacts

```
f(x) = x^2
```

```
1
      type thread_id
2
      val spawn: (unit -> unit) -> thread_id
3
4
      type 'a chan
5
      val channel : unit -> 'a chan
      val recv : 'a chan -> 'a
6
      val send : ('a chan * 'a) -> unit
7
8
1
2
      signature SERV = sig
3
        type serv
        val make : unit -> serv
4
5
        val call : serv * int -> int
6
7
8
      structure Serv : SERV = struct
9
         datatype serv = S of (int * int chan) chan
10
         fun make () = let
11
12
           val reqChn = channel ()
13
           fun loop state = let
14
             val (v, replCh) = recv reqChn in
15
             send (replCh, state);
16
             loop v end in
17
           spawn (fn () => loop 0);
18
           S reqChn end
19
20
         fun call (server, v) = let
21
           val S reqChn = server
22
           val replChn = channel () in
           send (reqCh, (v, replCh));
23
24
           recv replChn end end
25
26
1
2
      type 'a event
3
      val sync : 'a event -> 'a
4
      val recvEvt : 'a chan -> 'a event
5
      val sendEvt : 'a chan * 'a -> unit event
6
      val choose: 'a event * 'a event -> 'a event
7
8
      fun send (ch, v) = sync (sendEvt (ch, v))
      fun recv v = sync (recvEvt v)
9
10
```

```
11
      val thenEvt: 'a event * ('a -> 'b event) -> 'b event
12
13
1
      val server = Serv.make ()
2
      val _ = spawn (fn () => Serv.call (server, 35))
3
      val _ = spawn (fn () =>
4
        Serv.call (server, 12);
5
        Serv.call (server, 13))
      val _ = spawn (fn () => Serv.call (server, 81))
6
      val _ = spawn (fn () => Serv.call (server, 44))
7
8
1
      structure Serv : SERV = struct
2
        datatype serv = S of (int * int chan) chan
3
4
        fun make () = let
5
6
           val reqChn = FanIn.channel()
7
8
           fun loop state = let
9
             val (v, replCh) = FanIn.recv reqChn in
10
             OneShot.send (replCh, state);
             loop v end in
11
12
13
           spawn (fn () => loop 0);
14
           S reqChn end
15
16
         fun call (server, v) = let
17
           val S reqChn = server
18
           val replChn = OneShot.channel () in
19
           FanIn.send (reqCh, (v, replCh));
20
           OneShot.recv replChn end
21
22
         end
23
1
      let
2
        val w = 4
3
        val x = ref 1
4
        val y = ref 2
5
        val z = (!x + 1) + (!y + 2) + (w - 3)
6
        val w = 1 in
7
        y := 0;
         (!y + 2) - (!x + 1) * (w - 3) end
9
1
      let
2
        val x = 1
3
        val y = 2
```

```
val z = ref (4 * 73)
          val x = 4 in
 5
 6
          z := 1;
 7
          x * !z end
 8
 1
 2
        let
 3
          val f = fn x \Rightarrow x 1
          val g = fn y \Rightarrow y + 2
 4
          val h = fn z \Rightarrow z + 3 in
 5
 6
          (f g) + (f h) end
 7
1
 2
        datatype 'a list = Nil | Cons 'a ('a list)
 3
 4
        inductive sorted ::
 5
          ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow
          'a list \Rightarrow bool where
 6
 7
          Nil : sorted P Nil |
          Single : sorted P (Cons x Nil) |
 8
9
          Cons :
10
            P x y \Longrightarrow
11
            sorted P (Cons y ys) \Longrightarrow
12
            sorted P (Cons x (Cons y ys))
13
1
        datatype nat = Z \mid S nat
 2
 3
        inductive lte :: nat \Rightarrow nat \Rightarrow bool where
4
          Eq : lte n n |
          Lt : lte n1 n2 \Longrightarrow lte n1 (S n2)
5
 6
        theorem "
 7
          sorted lte
 8
             (Cons (Z) (Cons (S Z)
9
10
               (Cons (S Z) (Cons
                 (S (S (S Z))) Nil)))"
11
12
          apply (rule Cons)
13
          apply (rule Lt)
14
          apply (rule Eq)
15
          apply (rule Cons)
16
          apply (rule Eq)
17
          apply (rule Cons)
18
          apply (rule Lt)
19
          apply (rule Lt)
20
          apply (rule Eq)
21
          apply (rule Single)
```

```
22
         done
23
1
2
       definition True :: bool where
3
         True \equiv ((\lambda x::bool. x) = (\lambda x. x))
4
       definition False :: bool where
5
         False \equiv (\forallP. P)
6
7
8
1
2
       signature CHAN = sig
3
         type 'a chan
         val channel: unit \rightarrow 'a chan
4
5
         val send: 'a chan * 'a -> unit
 6
         val recv: 'a chan -> 'a
 7
         end
1
2
       structure ManyToManyChan : CHAN = struct
3
         type message_queue = 'a option ref queue
4
5
         datatype 'a chan_content =
6
           Send of (condition * 'a) queue |
7
           Recv of (condition \ast 'a option ref) queue |
8
           Inac
9
10
         datatype 'a chan =
11
           Chn of 'a chan_content ref * mutex_lock
12
         fun channel () = Chn (ref Inac, mutexLock ())
13
14
15
         fun send (Chn (conRef, lock)) m =
16
           acquire lock;
17
           (case !conRef of
18
             Recv q => let
19
                val (recvCond, mopRef) = dequeue q in
20
                mopRef := Some m;
21
                if (isEmpty q) then conRef := Inac else ();
22
                release lock; signal recvCond; () end |
23
             Send q => let
24
                val sendCond = condition () in
25
                enqueue (q, (sendCond, m));
26
                release lock; wait sendCond; () end |
27
             Inac => let
28
                val sendCond = condition () in
29
                conRef := Send (queue [(sendCond, m)]);
```

```
30
               release lock; wait sendCond; () end)
31
32
         fun recv (Chn (conRef, lock)) =
33
           acquire lock;
34
           (case !conRef of
35
             Send q \Rightarrow let
36
               val (sendCond, m) = dequeue q in
37
               if (isEmpty q) then
                 conRef := Inac
38
39
               else
40
                 ();
               release lock; signal sendCond; m end |
41
42
             Recv q => let
43
               val recvCond = condition ()
44
               val mopRef = ref None in
45
               enqueue (q, (recvCond, mopRef));
46
               release lock; wait recvCond;
47
               valOf (!mopRef) end |
48
             Inac => let
49
               val recvCond = condition ()
               val mopRef = ref None in
50
51
               conRef := Recv (queue [(recvCond, mopRef)]);
52
               release lock; wait recvCond;
               valOf (!mopRef) end)
53
54
55
         end
56
57
1
2
         structure FanOutChan : CHAN = struct
3
4
         datatype 'a chan_content =
           Send of condition * 'a |
5
6
           Recv of (condition * 'a option ref) queue |
7
           Inac
8
9
         datatype 'a chan =
10
           Chn of 'a chan_content ref * mutex_lock
11
12
         fun channel () = Chn (ref Inac, mutexLock ())
13
14
         fun send (Chn (conRef, lock)) m = let
15
           val sendCond = condition () in
16
           case cas (conRef, Inac, Send (sendCond, m)) of
17
             Inac => (* conRef already set *)
18
               wait sendCond; () |
19
             Recv q =>
20
             (* the current thread is
21
              * the only one that updates from this state *)
```

```
22
               acquire lock;
23
               (let
24
                 val (recvCond, mopRef) = dequeue q in
25
                 mopRef := Some m;
26
                 if (isEmpty q) then conRef := Inac else ();
27
                 release lock; signal (recvCond);
28
                 () end) |
29
             Send _ => raise NeverHappens end
30
31
         fun recv (Chn (conRef, lock)) =
32
           acquire lock;
33
           (case !conRef of
             Inac => let
34
35
               val recvCond = condition ()
36
               val mopRef = ref None in
37
               conRef := Recv (queue [(recvCond, mopRef)]);
38
               release lock; wait recvCond;
39
               valOf (!mopRef) end |
40
             Recv q => let
41
               val recvCond = condition ()
42
               val mopRef = ref None in
               enqueue (q, (recvCond, mopRef));
43
44
               release lock; wait recvCond;
               valOf (!mopRef) end |
45
             Send (sendCond, m) =>
46
               conRef := Inac;
47
48
               release lock;
49
               signal sendCond;
50
               m end)
51
52
         end
53
1
       structure FanInChan : CHAN = struct
2
3
       datatype 'a chan_content =
4
         Send of (condition * 'a) queue |
5
         Recv of condition \ast 'a option ref |
6
         Inac
7
8
       datatype 'a chan =
9
         Chn of 'a chan_content ref * mutex_lock
10
11
       fun channel () = Chn (ref Inac, mutexLock ())
12
13
       fun send (Chn (conRef, lock)) m =
14
         acquire lock;
15
         case !conRef of
16
         Recv (recvCond, mopRef) =>
17
           mopRef := Some m; conRef := Inac;
```

```
18
           release lock; signal recvCond;
19
           () |
20
         Send q \Rightarrow let
21
           val sendCond = condition () in
22
           enqueue (q, (sendCond, m));
23
           release lock; wait sendCond;
24
           () end |
25
         Inac => let
26
           val sendCond = condition () in
27
           conRef := Send (queue [(sendCond, m)])
           release lock; wait sendCond; () end
28
29
30
       fun recv (Chn (conRef, lock)) = let
31
         val recvCond = condition ()
32
         val mopRef = ref None in
33
         case cas (conRef, Inac, Recv (recvCond, mopRef)) of
34
           Inac => (* conRef already set *)
35
             wait recvCond; valOf (!mopRef) |
36
           Send q \Rightarrow
37
             (* the current thread is the only one
             -* that updates the state from this state *)
39
             acquire lock;
40
             (let
               val (sendCond, m) = dequeue q in
41
               if (isEmpty q) then conRef := Inac else ();
42
               release lock; signal sendCond; m end) |
43
44
           Recv _ => raise NeverHappens end end
45
46
2
    structure OneToOneChan : CHAN = struct
3
4
       datatype 'a chan_content =
         Send of condition * 'a |
5
         Recv of condition * 'a option ref |
6
7
         Inac
8
9
       datatype 'a chan = Chn of 'a chan_content ref
10
       fun channel () = Chn (ref Inac)
11
12
13
       fun send (Chn conRef) m = let
14
         val sendCond = condition () in
15
         case cas (conRef, Inac, Send (sendCond, m)) of
16
           Inac =>
17
             (* conRef already set to Send *)
18
             wait sendCond; () |
19
           Recv (recvCond, mopRef) =>
20
             (* the current thread is the only one
```

```
21
            22
            mopRef := Some m; conRef := Inac;
23
            signal recvCond; () |
24
          Send _ => raise NeverHappens end end
25
26
27
      fun recv (Chn conRef) = let
28
        val recvCond = condition ();
29
        val mopRef = ref None in
30
        case cas (conRef, Inac, Recv (recvCond, mopRef)) of
          Inac => (* conRef already set to Recv*)
31
32
            wait recvCond; valOf (!mopRef) |
33
          Send (sendCond, m) =>
34
            (* the current thread is the only one
35
            36
            conRef := Inac; signal sendCond; m |
37
          Recv _ => raise NeverHappens end end
38
39
      end
40
1
      structure OneShotChan : CHAN = struct
2
3
      datatype 'a chan_content =
        Send of condition * 'a |
4
5
        Recv of condition \ast 'a option ref |
6
        Inac
8
      datatype 'a chan = Chn of 'a chan_content ref *
      mutex_lock
9
10
      fun channel () = Chn (ref Inac, lock ())
11
12
      fun send (Chn (conRef, lock)) m = let
13
        val sendCond = condition () in
        case (conRef, Inac, Send (sendCond, m)) of
14
15
          Inac =>
16
            (* conRef already set to Send*)
17
            wait sendCond; () |
18
          Recv (recvCond, mopRef) =>
19
            mopRef := Some m; signal recvCond;
20
            () |
21
          Send _ => raise NeverHappens end end
22
23
24
      fun recv (Chn (conRef, lock)) = let
25
        val recvCond = condition ()
26
        val mopRef = ref None in
27
        case (conRef, Inac, Recv (recvCond, mopRef)) of
28
          Inac =>
```

```
29
             (* conRef already set to Recv*)
30
             wait recvCond; valOf (!mopRef) |
31
           Send (sendCond, m) =>
32
             acquire lock; signal sendCond;
33
             (* never relases lock;
34
             -* blocks others forever *)
35
             m |
36
           Recv _ =>
37
             acquire lock;
38
             (* never able to acquire lock;
             -* blocked forever *)
39
40
             raise NeverHappens end end
41
42
       end
43
1
    structure OneShotToOneChan : CHAN = struct
2
3
       datatype 'a chan =
4
         Chn of condition * condition * 'a option ref
5
6
      fun channel () =
7
         Chn (condition (), condition (), ref None)
8
9
       fun send (Chn (sendCond, recvCond, mopRef)) m =
10
         mopRef := Some m; signal recvCond;
11
         wait sendCond; ()
12
13
       fun recv (Chn (sendCond, recvCond, mopRef)) =
14
         wait recvCond; signal sendCond;
15
         valOf (!mopRef)
16
17
       end
18
1
2
       datatype var = Var string
3
4
       datatype exp =
5
        Let var boundexp exp |
6
         Rslt var
7
8
      boundexp =
9
         Unit |
10
         MkChn |
11
         Prim prim |
12
         Spwn exp |
         Sync var |
13
         Fst var |
14
         Snd var |
15
```

```
16
          Case var var exp var exp |
17
          App var var and
18
19
        prim =
20
          SendEvt var var |
21
          RecvEvt var |
22
          Pair var var |
23
          Lft var |
24
          Rht var |
25
          Abs var var ex
26
27
1
        datatype ctrl_label =
 2
          LNxt var | LSpwn var | LCall var | LRtn var
 3
        type_synonym ctrl_path = (ctrl_label list)
 4
 5
 6
        datatype chan = Chn ctrl_path var
 7
 8
        datatype val =
9
          {\tt VUnit \ | \ VChn \ chan \ | \ VClsr \ prim \ (var \rightharpoonup val)}
10
11
        datatype ctn = Ctn var exp (var \rightarrow val)
12
13
        datatype state = Stt exp (var \rightarrow val) (ctn list)
14
15
 1
 2
        inductive seq_step ::
 3
          bind * (var \rightharpoonup val)) \Rightarrow val \Rightarrow bool where
 4
          LetUnit:
 5
             seq_step (Unit, env) VUnit |
 6
          LetPrim:
 7
             seq_step (Prim p, env) (VClsr p env) |
 8
          LetFst:
9
             env xp = Some (VClsr (Pair x1 x2) envp) \Longrightarrow
10
             envp x1 = Some v \Longrightarrow
11
             seq_step (Fst xp, env) v |
12
          LetSnd:
13
             env xp = Some (VClsr (Pair x1 x2) envp) \Longrightarrow
14
             envp x2 = Some v \Longrightarrow
15
             seq_step (Snd xp, env) v
16
17
        inductive seq_step_up ::
18
          bind * (var \rightharpoonup val)) \Rightarrow exp * val_env \Rightarrow bool where
19
          {\tt LetCaseLft:}
             env xs = Some (VClsr (Lft xl') envl) \Longrightarrow
20
21
             envl xl' = Some vl \Longrightarrow
```

```
22
             seq_step_up
23
                (Case xs xl el xr er, env)
24
                (el, env(xl \mapsto vl)) |
25
           LetCaseRht:
26
             env xs = Some (VClsr (Rht xr') envr) \Longrightarrow
27
             envr xr' = Some vr \Longrightarrow
28
             seq_step_up
29
                (Case xs xl el xr er, env)
30
                (er, env(xr \mapsto vr)) |
31
           LetApp:
32
             env f = Some (VClsr (Abs fp xp el) envl) \Longrightarrow
33
             env xa = Some va \Longrightarrow
34
             seq_step_up
35
                (App f xa, env)
36
                (el, envl(
37
                  \texttt{fp} \; \mapsto \; \texttt{(VClsr (Abs fp xp el) envl),}
38
                  xp \mapsto va))
39
40
41
        type_synonym cmmn_set = (ctrl_path * chan * ctrl_path)
42
43
        44
45
        inductive leaf ::
46
           trace\_pool \Rightarrow ctrl\_path \Rightarrow bool where
47
           Intro:
48
             \mathtt{trpl}\ \mathtt{pi}\ \neq\ \mathtt{None}\ \Longrightarrow
49
             (∄ pi'. trpl pi' \ne None \land strict_prefix pi pi') \Longrightarrow
50
             leaf trpl pi
51
52
        inductive concur_step ::
53
           trace_pool * cmmn_set ⇒
54
           trace_pool * cmmn_set \Rightarrow
55
           bool where
56
           Seq_Sttep_Down:
57
             leaf trpl pi \Longrightarrow
             trpl pi = Some
58
59
                (Stt (Rslt x) env
60
                  ((Ctn xk ek envk) # k)) \Longrightarrow
61
             env x = Some v \Longrightarrow
62
             concur_step
63
                (trpl, ys)
64
                (trpl(pi @ [LRtn xk] \mapsto
65
                  (Stt ek (envk(xk \mapsto v)) k)), ys) |
66
           Seq_Step:
67
             leaf trpl pi \Longrightarrow
             trpl pi = Some
68
                (Stt (Let x b e) env k) \Longrightarrow
69
70
             seq\_step (b, env) v\Longrightarrow
```

```
71
               concur_step
 72
                  (trpl, ys)
 73
                  (trpl(pi @ [LNxt x] \mapsto
 74
                    (Stt e (env(x \mapsto v)) k), ys) |
 75
            Seq_Step_Up:
 76
               leaf trpl pi \Longrightarrow
 77
               trpl pi = Some
 78
                  (Stt (Let x b e) env k) \Longrightarrow
 79
               seq\_step\_up (b, env) (e', env') \Longrightarrow
 80
               concur_step
                  (trpl, ys)
 81
                  (\texttt{trpl}(\texttt{pi @ [LCall x]} \; \mapsto \;
 82
                    (Stt e' env'
 83
 84
                       ((Ctn x e env) # k))), ys) |
            LetMkCh:
 85
 86
               \texttt{leaf trpl pi} \Longrightarrow
 87
               trpl pi = Some (Stt (Let x MkChn e) env k) \Longrightarrow
 88
               concur_step
 89
                  (trpl, ys)
 90
                  (trpl(pi @ [LNxt x] \mapsto
 91
                    (Stt e (env(x \mapsto (VChn (Chn pi x)))) k)), ys) |
 92
            LetSpwn:
 93
               \texttt{leaf trpl pi} \implies
 94
               trpl pi = Some
 95
                  (Stt (Let x (Spwn ec) e) env k) \Longrightarrow
 96
               concur_step
 97
                  (trpl, ys)
 98
                  (trpl(
 99
                    pi 0 [LNxt x] \mapsto
100
                          (St e (env(x \mapsto VUnit)) k),
101
                    pi @ [LSpwn x] \mapsto
102
                          (St ec env []), ys) |
103
            LetSync:
104
               leaf trpl pis \Longrightarrow
               trpl pis = Some
105
106
                  (Stt (Let xs (Sync xse) es) envs ks) \Longrightarrow
107
               envs xse = Some
108
                  (VClsr (SendEvt xsc xm) envse) \Longrightarrow
109
               leaf trpl pir \Longrightarrow
110
               trpl pir = Some
111
                  (Stt (Let xr (Sync xre) er) envr kr) \Longrightarrow
112
               envr xre = Some
113
                  (VClsr (RecvEvt xrc) envre) \Longrightarrow
114
               envse xsc = Some (VChn c) \Longrightarrow
115
               envre xrc = Some (VChn c) \Longrightarrow
116
               envse xm = Some vm \Longrightarrow
117
               concur_step
118
                  (trpl, ys)
119
                  (trpl(
120
                    pis 0 [LNxt xs] \mapsto
```

```
121
                           (Stt es (envs(xs \mapsto VUnit)) ks),
122
                        pir 0 [LNxt xr] \mapsto
123
                           (Stt er (envr(xr \mapsto vm)) kr)),
124
                        ys \cup \{(pis, c, pir)\})
125
126
            inductive star :: ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow 'a \Rightarrow 'a \Rightarrow bool
  1
  2
              for r where
  3
              refl: star r x x |
  4
               \mathtt{step} \colon \mathtt{r} \ \mathtt{x} \ \mathtt{y} \implies \mathtt{star} \ \mathtt{r} \ \mathtt{y} \ \mathtt{z} \implies \mathtt{star} \ \mathtt{r} \ \mathtt{x} \ \mathtt{z}
  5
  1
           inductive is_send_path ::
  2
              \texttt{trace\_pool} \ \Rightarrow \ \texttt{chan} \ \Rightarrow
  3
              {\tt control\_path} \ \Rightarrow \ {\tt bool} \ \ {\tt where}
  4
  5
                  trpl piy = Some
  6
                     (Stt (Let xy (Sync xe) en) env k) \Longrightarrow
  7
                  env xe = Some
  8
                     (VClsr (SendEvt xsc xm) enve) \Longrightarrow
  9
                  enve xsc = Some (VChn c) \Longrightarrow
 10
                  is_send_path trpl c piy
 11
 12
            inductive is_recv_path ::
 13
               {\tt trace\_pool} \; \Rightarrow \; {\tt chan} \; \Rightarrow \;
 14
               {\tt control\_path} \ \Rightarrow \ {\tt bool} \ \ {\tt where}
 15
              intro:
 16
                  trpl piy = Some
 17
                     (Stt (Let xy (Sync xe) en) env k) \Longrightarrow
 18
                  env xe = Some
 19
                     (VClsr (RecvEvt xrc) enve) \Longrightarrow
 20
                  enve xrc = Some (VChan c) \Longrightarrow
 21
                  is_recv_path trpl c piy
 22
 23
            inductive all_two_dyn_paths ::
 24
               (control_path \Rightarrow bool) \Rightarrow
 25
               (control_path \Rightarrow control_path \Rightarrow bool) \Rightarrow
 26
               bool where
               intro: (\forall pi1 pi2 .
 27
 28
                     \texttt{p} \ \texttt{pi1} \ \longrightarrow
 29
                     p pi2 \longrightarrow
 30
                     r pi1 pi2) \Longrightarrow
 31
                  all_two_dyn_paths p r
 32
 33
            inductive ordered :: 'a list \Rightarrow 'a list \Rightarrow bool where
 34
              left: prefix pi1 pi2 \Longrightarrow ordered pi1 pi2 |
 35
               \texttt{right: prefix pi2 pi1} \implies \texttt{ordered pi1 pi2}
 36
 37
            inductive one_shot :: trace_pool \Rightarrow chan \Rightarrow bool where
```

```
38
           intro:
39
              all_two_dyn_paths
40
                 (is_send_path trpl c) op= \Longrightarrow
41
              one_shot trpl c
42
43
         inductive fan_out :: trace_pool \Rightarrow chan \Rightarrow bool where
44
           intro:
45
              all_two_dyn_paths
46
                 (is_send_path trpl c) ordered \Longrightarrow
47
              fan_out trpl c
48
49
         inductive fan_in :: trace_pool \Rightarrow chan \Rightarrow bool where
50
           intro:
51
              all_two_dyn_paths
52
                 (is_recv_path trpl c) ordered \Longrightarrow
53
              fan_in trpl c
54
55
         inductive one_to_one :: trace_pool \Rightarrow chan \Rightarrow bool where
56
           intro:
57
              \texttt{fan\_out trpl c} \implies
58
              \texttt{fan\_in} \ \texttt{trpl} \ \texttt{c} \Longrightarrow
59
              one_to_one trpl c
60
61
```