## Formal Theory of Communication Topology in Concurrent ML

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## 1 Mathematical Artifacts

```
f(x) = x^2
```

```
1
      type thread_id
2
      val spawn: (unit -> unit) -> thread_id
3
4
      type 'a chan
5
      val channel : unit -> 'a chan
      val recv : 'a chan -> 'a
6
      val send : ('a chan * 'a) -> unit
7
8
1
2
      signature SERV = sig
3
        type serv
        val make : unit -> serv
4
5
        val call : serv * int -> int
6
7
8
      structure Serv : SERV = struct
9
         datatype serv = S of (int * int chan) chan
10
         fun make () = let
11
12
           val reqChn = channel ()
13
           fun loop state = let
14
             val (v, replCh) = recv reqChn in
15
             send (replCh, state);
16
             loop v end in
17
           spawn (fn () => loop 0);
18
           S reqChn end
19
20
         fun call (server, v) = let
21
           val S reqChn = server
22
           val replChn = channel () in
           send (reqCh, (v, replCh));
23
24
           recv replChn end end
25
26
1
2
      type 'a event
3
      val sync : 'a event -> 'a
4
      val recvEvt : 'a chan -> 'a event
5
      val sendEvt : 'a chan * 'a -> unit event
6
      val choose: 'a event * 'a event -> 'a event
7
8
      fun send (ch, v) = sync (sendEvt (ch, v))
      fun recv v = sync (recvEvt v)
9
10
```

```
11
      val thenEvt: 'a event * ('a -> 'b event) -> 'b event
12
13
1
      val server = Serv.make ()
2
      val _ = spawn (fn () => Serv.call (server, 35))
3
      val _ = spawn (fn () =>
4
        Serv.call (server, 12);
5
        Serv.call (server, 13))
      val _ = spawn (fn () => Serv.call (server, 81))
6
      val _ = spawn (fn () => Serv.call (server, 44))
7
8
1
      structure Serv : SERV = struct
2
        datatype serv = S of (int * int chan) chan
3
4
        fun make () = let
5
6
           val reqChn = FanIn.channel()
7
8
           fun loop state = let
9
             val (v, replCh) = FanIn.recv reqChn in
10
             OneShot.send (replCh, state);
             loop v end in
11
12
13
           spawn (fn () => loop 0);
14
           S reqChn end
15
16
         fun call (server, v) = let
17
           val S reqChn = server
18
           val replChn = OneShot.channel () in
19
           FanIn.send (reqCh, (v, replCh));
20
           OneShot.recv replChn end
21
22
         end
23
1
      let
2
        val w = 4
3
        val x = ref 1
4
        val y = ref 2
5
        val z = (!x + 1) + (!y + 2) + (w - 3)
6
        val w = 1 in
7
        y := 0;
         (!y + 2) - (!x + 1) * (w - 3) end
9
1
      let
2
        val x = 1
3
        val y = 2
```

```
val z = ref (4 * 73)
          val x = 4 in
 5
 6
          z := 1;
 7
          x * !z end
 8
 1
 2
        let
 3
          val f = fn x \Rightarrow x 1
          val g = fn y \Rightarrow y + 2
 4
          val h = fn z \Rightarrow z + 3 in
 5
 6
          (f g) + (f h) end
 7
1
 2
        datatype 'a list = Nil | Cons 'a ('a list)
 3
 4
        inductive sorted ::
 5
          ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow
          'a list \Rightarrow bool where
 6
 7
          Nil : sorted P Nil |
          Single : sorted P (Cons x Nil) |
 8
9
          Cons :
10
            P x y \Longrightarrow
11
            sorted P (Cons y ys) \Longrightarrow
12
            sorted P (Cons x (Cons y ys))
13
1
        datatype nat = Z \mid S nat
 2
 3
        inductive lte :: nat \Rightarrow nat \Rightarrow bool where
4
          Eq : lte n n |
          Lt : lte n1 n2 \Longrightarrow lte n1 (S n2)
5
 6
        theorem "
 7
          sorted lte
 8
             (Cons (Z) (Cons (S Z)
9
10
               (Cons (S Z) (Cons
                 (S (S (S Z))) Nil)))"
11
12
          apply (rule Cons)
13
          apply (rule Lt)
14
          apply (rule Eq)
15
          apply (rule Cons)
16
          apply (rule Eq)
17
          apply (rule Cons)
18
          apply (rule Lt)
19
          apply (rule Lt)
20
          apply (rule Eq)
21
          apply (rule Single)
```

```
22
         done
23
1
2
       definition True :: bool where
3
         True \equiv ((\lambda x::bool. x) = (\lambda x. x))
4
       definition False :: bool where
5
         False \equiv (\forallP. P)
6
7
8
1
2
       signature CHAN = sig
3
         type 'a chan
         val channel: unit \rightarrow 'a chan
4
5
         val send: 'a chan * 'a -> unit
 6
         val recv: 'a chan -> 'a
 7
         end
1
2
       structure ManyToManyChan : CHAN = struct
3
         type message_queue = 'a option ref queue
4
5
         datatype 'a chan_content =
6
           Send of (condition * 'a) queue |
7
           Recv of (condition \ast 'a option ref) queue |
8
           Inac
9
10
         datatype 'a chan =
11
           Chn of 'a chan_content ref * mutex_lock
12
         fun channel () = Chn (ref Inac, mutexLock ())
13
14
15
         fun send (Chn (conRef, lock)) m =
16
           acquire lock;
17
           (case !conRef of
18
             Recv q => let
19
                val (recvCond, mopRef) = dequeue q in
20
                mopRef := Some m;
21
                if (isEmpty q) then conRef := Inac else ();
22
                release lock; signal recvCond; () end |
23
             Send q => let
24
                val sendCond = condition () in
25
                enqueue (q, (sendCond, m));
26
                release lock; wait sendCond; () end |
27
             Inac => let
28
                val sendCond = condition () in
29
                conRef := Send (queue [(sendCond, m)]);
```

```
30
               release lock; wait sendCond; () end)
31
32
         fun recv (Chn (conRef, lock)) =
33
           acquire lock;
34
           (case !conRef of
35
             Send q \Rightarrow let
36
               val (sendCond, m) = dequeue q in
37
               if (isEmpty q) then
                 conRef := Inac
38
39
               else
40
                 ();
               release lock; signal sendCond; m end |
41
42
             Recv q => let
43
               val recvCond = condition ()
44
               val mopRef = ref None in
45
               enqueue (q, (recvCond, mopRef));
46
               release lock; wait recvCond;
47
               valOf (!mopRef) end |
48
             Inac => let
49
               val recvCond = condition ()
               val mopRef = ref None in
50
51
               conRef := Recv (queue [(recvCond, mopRef)]);
52
               release lock; wait recvCond;
               valOf (!mopRef) end)
53
54
55
         end
56
57
1
2
         structure FanOutChan : CHAN = struct
3
4
         datatype 'a chan_content =
           Send of condition * 'a |
5
6
           Recv of (condition * 'a option ref) queue |
7
           Inac
8
9
         datatype 'a chan =
10
           Chn of 'a chan_content ref * mutex_lock
11
12
         fun channel () = Chn (ref Inac, mutexLock ())
13
14
         fun send (Chn (conRef, lock)) m = let
15
           val sendCond = condition () in
16
           case cas (conRef, Inac, Send (sendCond, m)) of
17
             Inac => (* conRef already set *)
18
               wait sendCond; () |
19
             Recv q =>
20
             (* the current thread is
21
              * the only one that updates from this state *)
```

```
22
               acquire lock;
23
               (let
24
                 val (recvCond, mopRef) = dequeue q in
25
                 mopRef := Some m;
26
                 if (isEmpty q) then conRef := Inac else ();
27
                 release lock; signal (recvCond);
28
                 () end) |
29
             Send _ => raise NeverHappens end
30
31
         fun recv (Chn (conRef, lock)) =
32
           acquire lock;
33
           (case !conRef of
             Inac => let
34
35
               val recvCond = condition ()
36
               val mopRef = ref None in
37
               conRef := Recv (queue [(recvCond, mopRef)]);
38
               release lock; wait recvCond;
39
               valOf (!mopRef) end |
40
             Recv q => let
41
               val recvCond = condition ()
42
               val mopRef = ref None in
               enqueue (q, (recvCond, mopRef));
43
44
               release lock; wait recvCond;
               valOf (!mopRef) end |
45
             Send (sendCond, m) =>
46
               conRef := Inac;
47
48
               release lock;
49
               signal sendCond;
50
               m end)
51
52
         end
53
1
       structure FanInChan : CHAN = struct
2
3
       datatype 'a chan_content =
4
         Send of (condition * 'a) queue |
5
         Recv of condition \ast 'a option ref |
6
         Inac
7
8
       datatype 'a chan =
9
         Chn of 'a chan_content ref * mutex_lock
10
11
       fun channel () = Chn (ref Inac, mutexLock ())
12
13
       fun send (Chn (conRef, lock)) m =
14
         acquire lock;
15
         case !conRef of
16
         Recv (recvCond, mopRef) =>
17
           mopRef := Some m; conRef := Inac;
```

```
18
           release lock; signal recvCond;
19
           () |
20
         Send q \Rightarrow let
21
           val sendCond = condition () in
22
           enqueue (q, (sendCond, m));
23
           release lock; wait sendCond;
24
           () end |
25
         Inac => let
26
           val sendCond = condition () in
27
           conRef := Send (queue [(sendCond, m)])
           release lock; wait sendCond; () end
28
29
30
       fun recv (Chn (conRef, lock)) = let
31
         val recvCond = condition ()
32
         val mopRef = ref None in
33
         case cas (conRef, Inac, Recv (recvCond, mopRef)) of
34
           Inac => (* conRef already set *)
35
             wait recvCond; valOf (!mopRef) |
36
           Send q \Rightarrow
37
             (* the current thread is the only one
             -* that updates the state from this state *)
39
             acquire lock;
40
             (let
               val (sendCond, m) = dequeue q in
41
               if (isEmpty q) then conRef := Inac else ();
42
               release lock; signal sendCond; m end) |
43
44
           Recv _ => raise NeverHappens end end
45
46
2
    structure OneToOneChan : CHAN = struct
3
4
       datatype 'a chan_content =
         Send of condition * 'a |
5
         Recv of condition * 'a option ref |
6
7
         Inac
8
9
       datatype 'a chan = Chn of 'a chan_content ref
10
       fun channel () = Chn (ref Inac)
11
12
13
       fun send (Chn conRef) m = let
14
         val sendCond = condition () in
15
         case cas (conRef, Inac, Send (sendCond, m)) of
16
           Inac =>
17
             (* conRef already set to Send *)
18
             wait sendCond; () |
19
           Recv (recvCond, mopRef) =>
20
             (* the current thread is the only one
```

```
21
            22
            mopRef := Some m; conRef := Inac;
23
            signal recvCond; () |
24
          Send _ => raise NeverHappens end end
25
26
27
      fun recv (Chn conRef) = let
28
        val recvCond = condition ();
29
        val mopRef = ref None in
30
        case cas (conRef, Inac, Recv (recvCond, mopRef)) of
          Inac => (* conRef already set to Recv*)
31
32
            wait recvCond; valOf (!mopRef) |
33
          Send (sendCond, m) =>
34
            (* the current thread is the only one
35
            36
            conRef := Inac; signal sendCond; m |
37
          Recv _ => raise NeverHappens end end
38
39
      end
40
1
      structure OneShotChan : CHAN = struct
2
3
      datatype 'a chan_content =
        Send of condition * 'a |
4
5
        Recv of condition \ast 'a option ref |
6
        Inac
8
      datatype 'a chan = Chn of 'a chan_content ref *
      mutex_lock
9
10
      fun channel () = Chn (ref Inac, lock ())
11
12
      fun send (Chn (conRef, lock)) m = let
13
        val sendCond = condition () in
        case (conRef, Inac, Send (sendCond, m)) of
14
15
          Inac =>
16
            (* conRef already set to Send*)
17
            wait sendCond; () |
18
          Recv (recvCond, mopRef) =>
19
            mopRef := Some m; signal recvCond;
20
            () |
21
          Send _ => raise NeverHappens end end
22
23
24
      fun recv (Chn (conRef, lock)) = let
25
        val recvCond = condition ()
26
        val mopRef = ref None in
27
        case (conRef, Inac, Recv (recvCond, mopRef)) of
28
          Inac =>
```

```
29
             (* conRef already set to Recv*)
30
             wait recvCond; valOf (!mopRef) |
31
           Send (sendCond, m) =>
32
             acquire lock; signal sendCond;
33
             (* never relases lock;
34
             -* blocks others forever *)
35
             m |
36
           Recv _ =>
37
             acquire lock;
38
             (* never able to acquire lock;
             -* blocked forever *)
39
40
             raise NeverHappens end end
41
42
       end
43
1
    structure OneShotToOneChan : CHAN = struct
2
3
       datatype 'a chan =
4
         Chn of condition * condition * 'a option ref
5
6
      fun channel () =
7
         Chn (condition (), condition (), ref None)
8
9
       fun send (Chn (sendCond, recvCond, mopRef)) m =
10
         mopRef := Some m; signal recvCond;
11
         wait sendCond; ()
12
13
       fun recv (Chn (sendCond, recvCond, mopRef)) =
14
         wait recvCond; signal sendCond;
15
         valOf (!mopRef)
16
17
       end
18
1
2
       datatype var = Var string
3
4
       datatype exp =
5
        Let var boundexp exp |
6
         Rslt var
7
8
      boundexp =
9
         Unit |
10
         MkChn |
11
         Prim prim |
12
         Spwn exp |
         Sync var |
13
         Fst var |
14
         Snd var |
15
```

```
16
          Case var var exp var exp |
17
          App var var and
18
19
        prim =
20
          SendEvt var var |
21
          RecvEvt var |
22
          Pair var var |
23
          Lft var |
24
          Rht var |
25
          Abs var var ex
26
27
1
        datatype ctrl_label =
 2
          LNxt var | LSpwn var | LCall var | LRtn var
 3
        type_synonym ctrl_path = (ctrl_label list)
 4
 5
 6
        datatype chan = Chn ctrl_path var
 7
 8
        datatype val =
9
          {\tt VUnit \ | \ VChn \ chan \ | \ VClsr \ prim \ (var \rightharpoonup val)}
10
11
        datatype ctn = Ctn var exp (var \rightarrow val)
12
13
        datatype state = Stt exp (var \rightarrow val) (ctn list)
14
15
 1
 2
        inductive seq_step ::
 3
          bind * (var \rightharpoonup val)) \Rightarrow val \Rightarrow bool where
 4
          LetUnit:
 5
             seq_step (Unit, env) VUnit |
 6
          LetPrim:
 7
             seq_step (Prim p, env) (VClsr p env) |
 8
          LetFst:
9
             env xp = Some (VClsr (Pair x1 x2) envp) \Longrightarrow
10
             envp x1 = Some v \Longrightarrow
11
             seq_step (Fst xp, env) v |
12
          LetSnd:
13
             env xp = Some (VClsr (Pair x1 x2) envp) \Longrightarrow
14
             envp x2 = Some v \Longrightarrow
15
             seq_step (Snd xp, env) v
16
17
        inductive seq_step_up ::
18
          bind * (var \rightharpoonup val)) \Rightarrow exp * val_env \Rightarrow bool where
19
          {\tt LetCaseLft:}
             env xs = Some (VClsr (Lft xl') envl) \Longrightarrow
20
21
             envl xl' = Some vl \Longrightarrow
```

```
22
             seq_step_up
23
                (Case xs xl el xr er, env)
24
                (el, env(xl \mapsto vl)) |
25
           LetCaseRht:
26
             env xs = Some (VClsr (Rht xr') envr) \Longrightarrow
27
             envr xr' = Some vr \Longrightarrow
28
             seq_step_up
29
                (Case xs xl el xr er, env)
30
                (er, env(xr \mapsto vr)) |
31
           LetApp:
32
             env f = Some (VClsr (Abs fp xp el) envl) \Longrightarrow
33
             env xa = Some va \Longrightarrow
34
             seq_step_up
35
                (App f xa, env)
36
                (el, envl(
37
                  \texttt{fp} \; \mapsto \; \texttt{(VClsr (Abs fp xp el) envl),}
38
                  xp \mapsto va))
39
40
41
        type_synonym cmmn_set = (ctrl_path * chan * ctrl_path)
42
43
        44
45
        inductive leaf ::
46
           trace\_pool \Rightarrow ctrl\_path \Rightarrow bool where
47
           Intro:
48
             \mathtt{trpl}\ \mathtt{pi}\ \neq\ \mathtt{None}\ \Longrightarrow
49
             (∄ pi'. trpl pi' \ne None \land strict_prefix pi pi') \Longrightarrow
50
             leaf trpl pi
51
52
        inductive concur_step ::
53
           trace_pool * cmmn_set ⇒
54
           trace_pool * cmmn_set \Rightarrow
55
           bool where
56
           Seq_Sttep_Down:
             \texttt{leaf trpl pi} \implies
57
             trpl pi = Some (
58
59
               Stt (Rslt x) env ((Ctn xk ek envk) # k)) \Longrightarrow
60
             env x = Some v \Longrightarrow
61
             concur_step
                (trpl, ys)
62
63
                (trpl(pi @ [LRtn xk] \mapsto
64
                  (Stt ek (envk(xk \mapsto v)) k)), ys) |
65
           Seq_Step:
66
             leaf trpl pi \Longrightarrow
67
             trpl pi = Some (Stt (Let x b e) env k) \Longrightarrow
             seq\_step (b, env) v\Longrightarrow
68
69
             concur_step
70
                (trpl, ys)
```

```
71
                 (trpl(pi @ [LNxt x] \mapsto (Stt e (env(x \mapsto v)) k), ys)
 72
           Seq_Step_Up:
 73
              \texttt{leaf trpl pi} \implies
 74
              trpl pi = Some (Stt (Let x b e) env k) \Longrightarrow
 75
              seq\_step\_up (b, env) (e', env') \Longrightarrow
 76
              concur_step
 77
                 (trpl, ys)
 78
                 (\texttt{trpl}(\texttt{pi @ [LCall x]} \; \mapsto \;
 79
                   (Stt e' env' ((Ctn x e env) # k))), ys) |
 80
            LetMkCh:
 81
              leaf trpl pi \Longrightarrow
 82
              trpl pi = Some (Stt (Let x MkChn e) env k) \Longrightarrow
 83
              concur_step
                 (trpl, ys)
 84
 85
                 (trpl(pi @ [LNxt x] \mapsto
 86
                   (Stt e (env(x \mapsto (VChn (Chn pi x)))) k)), ys) |
 87
            LetSpwn:
 88
              leaf trpl pi \Longrightarrow
 89
              trpl pi = Some (Stt (Let x (Spwn ec) e) env k) \Longrightarrow
 90
              concur_step
 91
                 (trpl, ys)
 92
                 (trpl(
 93
                   pi @ [LNxt x] \mapsto (St e (env(x \mapsto VUnit)) k),
 94
                   pi @ [LSpwn x] \mapsto (St ec env []), ys) |
 95
            LetSync:
 96
              leaf trpl pis \Longrightarrow
 97
              trpl pis = Some (Stt (Let xs (Sync xse) es) envs ks)
 98
              envs xse = Some (VClsr (SendEvt xsc xm) envse) \Longrightarrow
 99
              leaf trpl pir \implies
100
              trpl pir = Some (Stt (Let xr (Sync xre) er) envr kr)
101
              envr xre = Some (VClsr (RecvEvt xrc) envre) \Longrightarrow
102
              envse xsc = Some (VChn c) \Longrightarrow
              envre xrc = Some (VChn c) \Longrightarrow
103
              envse xm = Some vm \Longrightarrow
104
105
              concur_step
106
                 (trpl, ys)
107
                 (trpl(
108
                   pis @ [LNxt xs] \mapsto
109
                      (Stt es (envs(xs \mapsto VUnit)) ks),
110
                   pir 0 [LNxt xr] \mapsto
111
                      (Stt er (envr(xr \mapsto vm)) kr)),
112
                   ys \cup \{(pis, c, pir)\})
113
114
```