Formal Theory of Communication Topology in Concurrent ML

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1 Mathematical Artifacts

```
type thread_id
val spawn: (unit -> unit) -> thread_id

type 'a chan
val channel : unit -> 'a chan
val recv : 'a chan -> 'a
val send : ('a chan * 'a) -> unit
```

```
signature SERV = sig
  type serv
  val make : unit -> serv
  val call : serv * int -> int
end

structure Serv : SERV = struct
  datatype serv = S of (int * int chan) chan

  fun make () = let
    val reqChn = channel ()
    fun loop state = let
      val (v, replCh) = recv reqChn in
      send (replCh, state);
      loop v end in
    spawn (fn () => loop 0);
    S reqChn end

  fun call (server, v) = let
    val S reqChn = server
    val replChn = channel () in
    send (reqCh, (v, replCh));
    recv replChn end end
```

```
type 'a event
val sync : 'a event -> 'a
val recvEvt : 'a chan -> 'a event
val sendEvt : 'a chan * 'a -> unit event
val choose: 'a event * 'a event -> 'a event

fun send (ch, v) = sync (sendEvt (ch, v))
fun recv v = sync (recvEvt v)

val thenEvt: 'a event * ('a -> 'b event) -> 'b event
```

```
12
13
```

```
1   val server = Serv.make ()
2   val _ = spawn (fn () => Serv.call (server, 35))
3   val _ = spawn (fn () =>
4     Serv.call (server, 12);
5     Serv.call (server, 13))
6   val _ = spawn (fn () => Serv.call (server, 81))
7   val _ = spawn (fn () => Serv.call (server, 44))
8
```

```
1   structure Serv : SERV = struct
2     datatype serv = S of (int * int chan) chan
3
4     fun make () = let
5
6       val reqChn = FanIn.channel()
7
8       fun loop state = let
9         val (v, replCh) = FanIn.recv reqChn in
10        OneShot.send (replCh, state);
11        loop v end in
12
13      spawn (fn () => loop 0);
14      S reqChn end
15
16    fun call (server, v) = let
17      val S reqChn = server
18      val replChn = OneShot.channel () in
19      FanIn.send (reqCh, (v, replCh));
20      OneShot.recv replChn end
21
22    end
23
```

```
1   let
2     val w = 4
3     val x = ref 1
4     val y = ref 2
5     val z = (!x + 1) + (!y + 2) + (w - 3)
6     val w = 1 in
7     y := 0;
8     (!y + 2) - (!x + 1) * (w - 3) end
9
```

```
1   let
2     val x = 1
3     val y = 2
4     val z = ref (4 * 73)
```

```
    val x = 4 in
    z := 1;
    x * !z end
```

```
let
  val f = fn x => x 1
  val g = fn y => y + 2
  val h = fn z => z + 3 in
  (f g) + (f h) end
```

```
datatype 'a list = Nil | Cons 'a ('a list)

inductive sorted ::
  ('a ⇒ 'a ⇒ bool) ⇒
  'a list ⇒ bool where
  Nil : sorted P Nil |
  Single : sorted P (Cons x Nil) |
  Cons :
    P x y ⟹
    sorted P (Cons y ys) ⟹
    sorted P (Cons x (Cons y ys))
```

```
datatype nat = Z | S nat

inductive lte :: nat ⇒ nat ⇒ bool where
  Eq : lte n n |
  Lt : lte n1 n2 ⟹ lte n1 (S n2)

theorem "
  sorted lte
    (Cons (Z) (Cons (S Z)
      (Cons (S Z) (Cons
        (S (S (S Z))) Nil))))"
  apply (rule Cons)
  apply (rule Lt)
  apply (rule Eq)
  apply (rule Cons)
  apply (rule Eq)
  apply (rule Cons)
  apply (rule Lt)
  apply (rule Lt)
  apply (rule Eq)
  apply (rule Single)
  done
```

```
definition True :: bool where
  True ≡ ((λx::bool. x) = (λx. x))

definition False :: bool where
  False ≡ (∀P. P)
```

```
signature CHAN = sig
  type 'a chan
  val channel: unit -> 'a chan
  val send: 'a chan * 'a -> unit
  val recv: 'a chan -> 'a
  end
```

```
structure ManyToManyChan : CHAN = struct
  type message_queue = 'a option ref queue

  datatype 'a chan_content =
    Send of (condition * 'a) queue |
    Recv of (condition * 'a option ref) queue |
    Inac

  datatype 'a chan =
    Chn of 'a chan_content ref * mutex_lock

  fun channel () = Chn (ref Inac, mutexLock ())

  fun send (Chn (conRef, lock)) m =
    acquire lock;
    (case !conRef of
      Recv q => let
        val (recvCond, mopRef) = dequeue q in
        mopRef := Some m;
        if (isEmpty q) then conRef := Inac else ();
        release lock; signal recvCond; () end |
      Send q => let
        val sendCond = condition () in
        enqueue (q, (sendCond, m));
        release lock; wait sendCond; () end |
      Inac => let
        val sendCond = condition () in
        conRef := Send (queue [(sendCond, m)]);
        release lock; wait sendCond; () end)

  fun recv (Chn (conRef, lock)) =
    acquire lock;
    (case !conRef of
      Send q => let
```

```
36          val (sendCond, m) = dequeue q in
37          if (isEmpty q) then
38            conRef := Inac
39          else
40            ();
41          release lock; signal sendCond; m end |
42        Recv q => let
43          val recvCond = condition ()
44          val mopRef = ref None in
45          enqueue (q, (recvCond, mopRef));
46          release lock; wait recvCond;
47          valOf (!mopRef) end |
48        Inac => let
49          val recvCond = condition ()
50          val mopRef = ref None in
51          conRef := Recv (queue [(recvCond, mopRef)]);
52          release lock; wait recvCond;
53          valOf (!mopRef) end)
54
55    end
56
57
```

```
1
2     structure FanOutChan : CHAN = struct
3
4     datatype 'a chan_content =
5       Send of condition * 'a |
6       Recv of (condition * 'a option ref) queue  |
7       Inac
8
9     datatype 'a chan =
10      Chn of 'a chan_content ref * mutex_lock
11
12    fun channel () = Chn (ref Inac, mutexLock ())
13
14    fun send (Chn (conRef, lock)) m = let
15      val sendCond = condition () in
16      case cas (conRef, Inac, Send (sendCond, m)) of
17        Inac => (* conRef already set *)
18          wait sendCond; () |
19        Recv q =>
20        (* the current thread is
21          * the only one that updates from this state *)
22          acquire lock;
23          (let
24            val (recvCond, mopRef) = dequeue q in
25            mopRef := Some m;
26            if (isEmpty q) then conRef := Inac else ();
27            release lock; signal (recvCond);
```

```
28              () end) |
29          Send _ => raise NeverHappens end
30
31    fun recv (Chn (conRef, lock)) =
32      acquire lock;
33      (case !conRef of
34        Inac => let
35          val recvCond = condition ()
36          val mopRef = ref None in
37          conRef := Recv (queue [(recvCond, mopRef)]);
38          release lock; wait recvCond;
39          valOf (!mopRef) end |
40        Recv q => let
41          val recvCond = condition ()
42          val mopRef = ref None in
43          enqueue (q, (recvCond, mopRef));
44          release lock; wait recvCond;
45          valOf (!mopRef) end |
46        Send (sendCond, m) =>
47          conRef := Inac;
48          release lock;
49          signal sendCond;
50          m end)
51
52    end
53
```

```
1   structure FanInChan : CHAN = struct
2
3   datatype 'a chan_content =
4     Send of (condition * 'a) queue |
5     Recv of condition * 'a option ref |
6     Inac
7
8   datatype 'a chan =
9     Chn of 'a chan_content ref * mutex_lock
10
11  fun channel () = Chn (ref Inac, mutexLock ())
12
13  fun send (Chn (conRef, lock)) m =
14    acquire lock;
15    case !conRef of
16    Recv (recvCond, mopRef) =>
17      mopRef := Some m; conRef := Inac;
18      release lock; signal recvCond;
19      () |
20    Send q => let
21      val sendCond = condition () in
22      enqueue (q, (sendCond, m));
23      release lock; wait sendCond;
```

```
      () end |
    Inac => let
      val sendCond = condition () in
      conRef := Send (queue [(sendCond, m)])
      release lock; wait sendCond; () end

  fun recv (Chn (conRef, lock)) = let
    val recvCond = condition ()
    val mopRef = ref None in
    case cas (conRef, Inac, Recv (recvCond, mopRef)) of
      Inac => (* conRef already set *)
        wait recvCond; valOf (!mopRef) |
      Send q =>
        (* the current thread is the only one
        -* that updates the state from this state *)
        acquire lock;
        (let
          val (sendCond, m) = dequeue q in
          if (isEmpty q) then conRef := Inac else ();
          release lock; signal sendCond; m end) |
      Recv _ => raise NeverHappens end end
```

```
structure OneToOneChan : CHAN = struct

  datatype 'a chan_content =
    Send of condition * 'a |
    Recv of condition * 'a option ref |
    Inac

  datatype 'a chan = Chn of 'a chan_content ref

  fun channel () = Chn (ref Inac)

  fun send (Chn conRef) m = let
    val sendCond = condition () in
    case cas (conRef, Inac, Send (sendCond, m)) of
      Inac =>
        (* conRef already set to Send *)
        wait sendCond; () |
      Recv (recvCond, mopRef) =>
        (* the current thread is the only one
        -* that accesses conRef for this state *)
        mopRef := Some m; conRef := Inac;
        signal recvCond; () |
      Send _ => raise NeverHappens end end
```

```
27  fun recv (Chn conRef) = let
28    val recvCond = condition ();
29    val mopRef = ref None in
30    case cas (conRef, Inac, Recv (recvCond, mopRef)) of
31      Inac => (* conRef already set to Recv*)
32        wait recvCond; valOf (!mopRef) |
33      Send (sendCond, m) =>
34        (* the current thread is the only one
35        -* that accesses conRef for this state *)
36        conRef := Inac; signal sendCond; m |
37      Recv _ => raise NeverHappens end end
38
39  end
40
```

```
 1  structure OneShotChan : CHAN = struct
 2
 3  datatype 'a chan_content =
 4    Send of condition * 'a |
 5    Recv of condition * 'a option ref |
 6    Inac
 7
 8  datatype 'a chan = Chn of 'a chan_content ref * mutex_lock
 9
10  fun channel () = Chn (ref Inac, lock ())
11
12  fun send (Chn (conRef, lock)) m = let
13    val sendCond = condition () in
14    case (conRef, Inac, Send (sendCond, m)) of
15      Inac =>
16        (* conRef already set to Send*)
17        wait sendCond; () |
18      Recv (recvCond, mopRef) =>
19        mopRef := Some m; signal recvCond;
20        () |
21      Send _ => raise NeverHappens end end
22
23
24  fun recv (Chn (conRef, lock)) = let
25    val recvCond = condition ()
26    val mopRef = ref None in
27    case (conRef, Inac, Recv (recvCond, mopRef)) of
28      Inac =>
29        (* conRef already set to Recv*)
30        wait recvCond; valOf (!mopRef) |
31      Send (sendCond, m) =>
32        acquire lock; signal sendCond;
33        (* never relases lock;
34        -* blocks others forever *)
35        m |
```

```
      Recv _ =>
        acquire lock;
        (* never able to acquire lock;
        -* blocked forever *)
        raise NeverHappens end end

  end
```

```
structure OneShotToOneChan : CHAN = struct

  datatype 'a chan =
    Chn of condition * condition * 'a option ref

  fun channel () =
    Chn (condition (), condition (), ref None)

  fun send (Chn (sendCond, recvCond, mopRef)) m =
    mopRef := Some m; signal recvCond;
    wait sendCond; ()

  fun recv (Chn (sendCond, recvCond, mopRef)) =
    wait recvCond; signal sendCond;
    valOf (!mopRef)

  end
```

2 Syntax

```

  datatype var = Var string

  datatype exp =
    Let var boundexp exp |
    Rslt var

  boundexp =
    Unt |
    MkChn |
    Prim prim |
    Spwn exp |
    Sync var |
    Fst var |
    Snd var |
    Case var var exp var exp |
    App var var and

  prim =
```

```
    SendEvt var var |
    RecvEvt var |
    Pair var var |
    Lft var |
    Rht var |
    Abs var var ex
```

3 Dynamic Semantics

```
datatype ctrl_label =
  LNxt var | LSpwn var | LCall var | LRtn var

type_synonym ctrl_path = (ctrl_label list)

datatype chan = Chn ctrl_path var

datatype val =
  VUnt | VChn chan | VClsr prim (var ⇀ val)

datatype ctn = Ctn var exp (var ⇀ val)

datatype state = Stt exp (var ⇀ val) (ctn list)
```

```
inductive seq_step ::
  bind * (var ⇀ val)) ⇒ val ⇒ bool where
  LetUnt:
    seq_step (Unt, env) VUnt |
  LetPrim:
    seq_step (Prim p, env) (VClsr p env) |
  LetFst:
    env xp = Some (VClsr (Pair x1 x2) envp) ⟹
    envp x1 = Some v ⟹
    seq_step (Fst xp, env) v |
  LetSnd:
    env xp = Some (VClsr (Pair x1 x2) envp) ⟹
    envp x2 = Some v ⟹
    seq_step (Snd xp, env) v
```

```
inductive seq_step_up ::
  bind * (var ⇀ val)) ⇒ exp * val_env ⇒ bool where
  LetCaseLft:
```

```
      env xs = Some (VClsr (Lft xl') envl) ⟹
      envl xl' = Some vl ⟹
      seq_step_up
        (Case xs xl el xr er, env)
        (el, env(xl ↦ vl)) |
    LetCaseRht:
      env xs = Some (VClsr (Rht xr') envr) ⟹
      envr xr' = Some vr ⟹
      seq_step_up
        (Case xs xl el xr er, env)
        (er, env(xr ↦ vr)) |
    LetApp:
      env f = Some (VClsr (Abs fp xp el) envl) ⟹
      env xa = Some va  ⟹
      seq_step_up
        (App f xa, env)
        (el, envl(
          fp ↦ (VClsr (Abs fp xp el) envl),
          xp ↦ va))
```

```

type_synonym cmmn_set = (ctrl_path * chan * ctrl_path) set

type_synonym trace_pool = ctrl_path ⇀ state

inductive leaf ::
  trace_pool ⇒ ctrl_path ⇒ bool where
  intro:
    trpl pi ≠ None ⟹
    (∄ pi' . trpl pi' ≠ None ∧ strict_prefix pi pi') ⟹
    leaf trpl pi


```

```

inductive concur_step ::
  trace_pool * cmmn_set ⇒
  trace_pool * cmmn_set ⇒
  bool where
  Seq_Sttep_Down:
    leaf trpl pi ⟹
    trpl pi = Some
      (Stt (Rslt x) env
        ((Ctn xk ek envk) # k)) ⟹
    env x = Some v ⟹
    concur_step
      (trpl, ys)
```

```
14          (trpl(pi @ [LRtn xk] ↦
15            (Stt ek (envk(xk ↦ v)) k)), ys) |
16      Seq_Step:
17        leaf trpl pi ⟹
18        trpl pi = Some
19          (Stt (Let x  b e) env k) ⟹
20        seq_step (b, env) v⟹
21        concur_step
22          (trpl, ys)
23          (trpl(pi @ [LNxt x] ↦
24            (Stt e (env(x ↦ v)) k), ys) |
25      Seq_Step_Up:
26        leaf trpl pi ⟹
27        trpl pi = Some
28          (Stt (Let x b e) env k) ⟹
29        seq_step_up (b, env) (e', env') ⟹
30        concur_step
31          (trpl, ys)
32          (trpl(pi @ [LCall x] ↦
33            (Stt e' env'
34              ((Ctn x e env) # k))), ys) |
35      LetMkCh:
36        leaf trpl pi ⟹
37        trpl pi = Some (Stt (Let x MkChn e) env k) ⟹
38        concur_step
39          (trpl, ys)
40          (trpl(pi @ [LNxt x] ↦
41            (Stt e (env(x ↦ (VChn (Chn pi x)))) k)), ys) |
42      LetSpwn:
43        leaf trpl pi ⟹
44        trpl pi = Some
45          (Stt (Let x (Spwn ec) e) env k) ⟹
46        concur_step
47          (trpl, ys)
48          (trpl(
49            pi @ [LNxt x] ↦
50                (St e (env(x ↦ VUnt)) k),
51            pi @ [LSpwn x] ↦
52                (St ec env []), ys) |
53      LetSync:
54        leaf trpl pis ⟹
55        trpl pis = Some
56          (Stt (Let xs (Sync xse) es) envs ks) ⟹
57        envs xse = Some
58          (VClsr (SendEvt xsc xm) envse) ⟹
59        leaf trpl pir  ⟹
60        trpl pir = Some
61          (Stt (Let xr (Sync xre) er) envr kr) ⟹
62        envr xre = Some
63          (VClsr (RecvEvt xrc) envre) ⟹
```

```
      envse xsc = Some (VChn c) ⟹
      envre xrc = Some (VChn c) ⟹
      envse xm = Some vm ⟹
      concur_step
        (trpl, ys)
        (trpl(
          pis @ [LNxt xs] ↦
            (Stt es (envs(xs ↦ VUnt)) ks),
          pir @ [LNxt xr] ↦
            (Stt er (envr(xr ↦ vm)) kr)),
          ys ∪ {(pis, c, pir)})


inductive star :: ('a ⇒ 'a ⇒ bool) ⇒ 'a ⇒ 'a ⇒ bool
  for r where
  refl: star r x x |
  step: r x y ⟹ star r y z ⟹ star r x z
```

4 Dynamic Communication

```
inductive is_send_path ::
  trace_pool ⇒ chan ⇒
  control_path ⇒ bool where
  intro:
    trpl piy = Some
      (Stt (Let xy (Sync xe) en) env k) ⟹
    env xe = Some
      (VClsr (SendEvt xsc xm) enve) ⟹
    enve xsc = Some (VChn c) ⟹
    is_send_path trpl c piy

inductive is_recv_path ::
  trace_pool ⇒ chan ⇒
  control_path ⇒ bool where
  intro:
    trpl piy = Some
      (Stt (Let xy (Sync xe) en) env k) ⟹
    env xe = Some
      (VClsr (RecvEvt xrc) enve) ⟹
    enve xrc = Some (VChn c) ⟹
    is_recv_path trpl c piy



inductive every_two ::
  ('a ⇒ bool) ⇒
```

```
    ('a ⇒ 'a ⇒ bool) ⇒
    bool where
    intro: (∀ pi1 pi2 .
        p x1 ⟶
        p x2 ⟶
        r x1 x2) ⟹
      every_two p r

inductive ordered :: 'a list ⇒ 'a list ⇒ bool where
  left: prefix pi1 pi2 ⟹ ordered pi1 pi2 |
  right: prefix pi2 pi1 ⟹ ordered pi1 pi2
```

```

inductive one_shot :: trace_pool ⇒ chan ⇒ bool where
  intro:
    every_two
      (is_send_path trpl c) op= ⟹
    one_shot trpl c

inductive fan_out :: trace_pool ⇒ chan ⇒ bool where
  intro:
    every_two
      (is_send_path trpl c) ordered ⟹
    fan_out trpl c

inductive fan_in :: trace_pool ⇒ chan ⇒ bool where
  intro:
    every_two
      (is_recv_path trpl c) ordered ⟹
    fan_in trpl c

inductive one_to_one :: trace_pool ⇒ chan ⇒ bool where
  intro:
    fan_out trpl c ⟹
    fan_in trpl c ⟹
    one_to_one trpl c
```

5 Static Semantics

```

datatype abstract_value =
  AChn var |
  AUnt |
  APrim prim

```

```
type_synonym abstract_env = var ⇒ abstract_value set

fun rslt_var :: exp ⇒ var where
  rslt_var (Rslt x) = x |
  rslt_var (Let _ _ e) = (rslt_var e)
```

```
inductive may_be_eval_exp ::
  abstract_env * abstract_env ⇒
  exp ⇒ bool where
  Result:
    may_be_eval_exp (V, C) (RESULT x) |
  Let_Unt:
    {AUnt} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x Unt e) |
  Let_Chan:
    {AChn x} ⊆ V x  ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (MkChn) e) |
  Let_SendEvt :
    {APrim (SendEvt xc xm)} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (Prim (SendEvt xc xm)) e
  ) |
  Let_RecvEvt :
    {APrim (RecvEvt xc)} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (Prim (RecvEvt xc)) e) |
  Let_Pair :
    {APrim (Pair x1 x2)} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (Pair x1 x2) e) |
  Let_Left :
    {APrim (Left xp)} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (Left xp) e) |
  Let_Right:
    {APrim (Right xp)} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (Right xp) e) |
  Let_Abs :
    {APrim (Abs f' x' e')} ⊆ V f' ⟹
    may_be_eval_exp (V, C) e' ⟹
    {APrim (Abs f' x' e')} ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_exp (V, C) (Let x (Abs f' x' e') e) |
```

```
42      Let_Spawn:
43        {AUnt} ⊆ V x ⟹
44        may_be_eval_exp (V, C) ec ⟹
45        may_be_eval_exp (V, C) e ⟹
46        may_be_eval_exp (V, C) (Let x (Spwn ec) e) |
47      Let_Sync  :
48        ∀ xsc xm xc .
49          (APrim (SendEvt xsc xm)) ∈ V xe ⟶
50          AChn xc ∈ V xsc ⟶
51          {AUnt} ⊆ V x ∧ V xm ⊆ C xc ⟹
52        ∀ xrc xc .
53          (APrim (RecvEvt xrc)) ∈ V xe ⟶
54          AChn xc ∈ V xrc ⟶
55          C xc ⊆ V x ⟹
56        may_be_eval_exp (V, C) e ⟹
57        may_be_eval_exp (V, C) (Let x (Syync xe) e) |
58      Let_Fst:
59        ∀ x1 x2.
60          (APrim (Pair x1 x2)) ∈ V xp ⟶
61          V x1 ⊆ V x ⟹
62        may_be_eval_exp (V, C) e ⟹
63        may_be_eval_exp (V, C) (Let x (Fst xp) e) |
64      Let_Snd:
65      ∀ x1 x2 .
66        (APrim (Pair x1 x2) ∈ V xp ⟶
67        V x2 ⊆ V x ⟹
68      may_be_eval_exp (V, C) e ⟹
69      may_be_eval_exp (V, C) (Let x (Snd xp) e) |
70    Let_Case:
71      ∀ xl' .
72        (APrim (Left xl')) ∈ V xs ⟶
73          V xl' ⊆ V xl ∧ V (rslt_var el) ⊆ V x ∧
74          may_be_eval_exp (V, C) el ⟹
75      ∀ xr' .
76        (APrim (Right xr')) ∈ V xs ⟶
77          V xr' ⊆ V xr ∧ V (rslt_var er) ⊆ V x ∧
78          may_be_eval_exp (V, C) er ⟹
79          may_be_eval_exp (V, C) e ⟹
80        may_be_eval_exp (V, C) (Let x (Case xs xl el xr er) e)
       |
81    Let_App:
82      ∀ f' x' e' .
83        (APrim (Abs f' x' e') ∈ V f ⟶
84        V xa ⊆ V x' ∧
85        V (rslt_var e') ⊆ V x ⟹
86      may_be_eval_exp (V, C) e ⟹
87      may_be_eval_exp (V, C) (Let x (App f xa) e)
88
89
```

```
fun abstract :: val ⇒ abstract_value where
  abstract VUnt = AUnt |
  abstract VChn (Chn pi x) = AChn x |
  abstract VClsr p env = APrim p
```

```
inductive
  may_be_eval_val ::
    abstract_env * abstract_env ⇒ val ⇒ bool and
  may_be_eval_env ::
    abstract_env * abstract_env ⇒ val_env ⇒ bool where
  Unt:
    may_be_eval_val (V, C) VUnt |
  Chan:
    may_be_eval_val (V, C) VChn c |
  SendEvt:
    may_be_eval_env (V, C) env ⟹
    may_be_eval_val (V, C) (VClsr (SendEvt _ _) env) |
  RecvEvt:
    may_be_eval_env (V, C) env ⟹
    may_be_eval_val (V, C) (VClsr (RecvEvt _) env) |
  Left:
    may_be_eval_env (V, C) env ⟹
    may_be_eval_val (V, C) (VClsr (Left _) env) |
  Right:
    may_be_eval_env (V, C) env ⟹
    may_be_eval_val (V, C) (VClsr (Right _) env) |
  Abs:
    {(APrim (Abs f x e)} ⊆ V f ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_env (V, C) env ⟹
    may_be_eval_val (V, C) (VClsr (Abs f x e) env) |
  Pair:
    may_be_eval_env (V, C) env ⟹
    may_be_eval_val (V, C) (VClsr (Pair _ _) env) |
  intro:
    ∀ x v .
      env x = Some v ⟶
      {abstract v} ⊆ V x ∧ may_be_eval_val (V, C) v ⟹
    may_be_eval_env (V, C) env
```

```
inductive may_be_eval_stack ::
  abstract_env * abstract_env ⇒
  abstract_value set ⇒ cont list ⇒ bool where
  Empty:
    may_be_eval_stack (V, C) valset [] |
  Nonempty:
    valset ⊆ V x ⟹
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_env (V, C) env ⟹
    may_be_eval_stack (V, C) (V (rslt_var e)) k ⟹
    may_be_eval_stack (V, C) valset ((Ctn x e env) # k)


inductive may_be_eval_state ::
  abstract_env * abstract_env ⇒
  state ⇒ bool where
  intro:
    may_be_eval_exp (V, C) e ⟹
    may_be_eval_env (V, C) env ⟹
    may_be_eval_stack (V, C) (V (rslt_var e)) k ⟹
    may_be_eval_state (V, C) (Stt e env k)

inductive may_be_eval_pool ::
  abstract_env * abstract_env ⇒
  trace_pool ⇒ bool where
  intro:
    ∀ pi st .
      trpl pi = Some st ⟶
      may_be_eval_state (V, C) st ⟹
    may_be_eval_pool (V, C) trpl
```

```
theorem may_be_eval_preserved_under_concur_step : "
  may_be_eval_pool (V, C) trpl ⟹
  concur_step (trpl, ys) (trpl', ys') ⟹
  may_be_stati_eval_pool (V, C) trpl'"
proof outline
qed

theorem may_be_eval_preserved_under_concur_step_star : "
  may_be_eval_pool (V, C) trpl ⟹
  star concur_step (trpl, ys) (trpl', ys')  ⟹
  may_be_concur_step (V, C) trpl'"
proof outline
qed
```

```
theorem trace_pool_snapshot_not_value_sound : "
  env x = Some v ⟹
  trpl pi = Some (Stt e env k) ⟹
  may_be_eval_pool (V, C) trpl ⟹
  {abstract v} ⊆ V x "
proof outline
qed
```

```
theorem trace_pool_always_not_value_sound : "
  env' x = Some v ⟹
  may_be_eval_pool (V, C) trpl ⟹
  star concur_step (trpl, ys) (trpl', ys') ⟹
  trpl' pi = Some (Stt e' env' k') ⟹
  {abstract v} ⊆ V x"
proof outline
qed

theorem exp_always_not_value_sound : "
  env' x = Some v ⟹
  may_be_eval_exp (V, C) e ⟹
  star concur_step
    ([[] ↦ (Stt e (λ _ . None) [])], ys)
    (trpl', ys') ⟹
  trpl' pi = Some (Stt e' env' k') ⟹
  {abstract v} ⊆ V x
proof outline
qed
```