Formal Theory of Communication Topology in Concurrent ML

Thomas Logan July 15, 2018

1 Mathematical Artifacts

```
1
    type thread_id
2
    val spawn : (unit -> unit) -> thread_id
3
    type 'a chan
4
5
    val channel : unit -> 'a chan
    val recv : 'a chan -> 'a
    val send : ('a chan * 'a) -> unit
1
2
    signature SERV = sig
3
      type serv
4
      val make : unit -> serv
      val call : serv * int -> int
5
6
    end
7
8
    structure Serv : SERV = struct
9
      datatype serv = S of (int * int chan) chan
10
11
      fun make () = let
12
        val reqChn = channel ()
13
        fun loop state = let
14
          val (v, replCh) = recv reqChn in
15
           send (replCh, state);
16
          loop v end in
17
         spawn (fn () => loop 0);
18
         S reqChn end
19
20
      fun call (server, v) = let
21
        val S reqChn = server
22
        val replChn = channel () in
23
        send (reqCh, (v, replCh));
24
        recv replChn end end
25
26
1
2
    type 'a event
    val sync : 'a event -> 'a
3
4
    val recvEvt : 'a chan -> 'a event
    val sendEvt : 'a chan * 'a -> unit event
    val choose : 'a event * 'a event -> 'a event
6
7
8
    fun send (ch, v) = sync (sendEvt (ch, v))
9
    fun recv v = sync (recvEvt v)
10
    val thenEvt : 'a event * ('a -> 'b event) -> 'b event
11
```

```
12
13
1
    val server = Serv.make ()
2
    val _ = spawn (fn () => Serv.call (server, 35))
3
    val _ = spawn (fn () =>
      Serv.call (server, 12);
      Serv.call (server, 13))
5
    val _ = spawn (fn () => Serv.call (server, 81))
6
7
    val \_ = spawn (fn () => Serv.call (server, 44))
8
1
    structure Serv : SERV = struct
2
       datatype serv = S of (int * int chan) chan
3
      fun make () = let
4
5
        val reqChn = FanIn.channel()
6
7
8
        fun loop state = let
9
           val (v, replCh) = FanIn.recv reqChn in
10
           OneShot.send (replCh, state);
11
           loop\ v\ end\ in
12
13
         spawn (fn () => loop 0);
14
         S reqChn end
15
16
       fun call (server, v) = let
17
        val S reqChn = server
18
        val replChn = OneShot.channel () in
19
        FanIn.send (reqCh, (v, replCh));
20
         OneShot.recv replChn end
21
22
       end
23
1
    let
      val w = 4
3
      val x = ref 1
4
      val y = ref 2
5
      val z = (!x + 1) + (!y + 2) + (w - 3)
6
      val w = 1 in
7
      y := 0;
       (!y + 2) - (!x + 1) * (w - 3) end
8
1
    let
2
      val x = 1
      val y = 2
3
      val z = ref (4 * 73)
```

```
5
        val x = 4 in
 6
        z := 1;
 7
        x * !z end
 8
1
 2
     let
 3
        val f = fn x \Rightarrow x 1
        val g = fn y \Rightarrow y + 2
 4
        val h = fn z \Rightarrow z + 3 in
 5
        (f g) + (f h) end
 6
1
2
     datatype 'a list = Nil | Cons 'a ('a list)
3
     inductive sorted : :
 5
        ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow
        'a list \Rightarrow bool where
 6
        Nil : sorted P Nil |
 7
        Single : sorted P (Cons x Nil) |
 8
        Cons :
9
10
          P x y \Longrightarrow
11
          sorted P (Cons y ys) \Longrightarrow
12
          sorted P (Cons x (Cons y ys))
13
1
     datatype nat = Z | S nat
2
     inductive lte : : nat \Rightarrow nat \Rightarrow bool where
3
        Eq : lte n n |
4
        Lt : lte n1 n2 \Longrightarrow lte n1 (S n2)
5
 6
7
     theorem "
 8
        sorted lte
9
          (Cons (Z) (Cons (S Z)
10
             (Cons (S Z) (Cons
11
               (S (S (S Z))) Nil)))"
12
        apply (rule Cons)
13
        apply (rule Lt)
14
        apply (rule Eq)
15
        apply (rule Cons)
16
        apply (rule Eq)
17
        apply (rule Cons)
        apply (rule Lt)
18
        apply (rule Lt)
apply (rule Eq)
19
20
21
        apply (rule Single)
22
        done
23
```

```
1
     definition True : : bool where
2
       True \equiv ((\lambdax : :bool. x) = (\lambdax. x))
3
4
     definition False : : bool where
5
       False \equiv (\forallP. P)
6
7
1
     signature CHAN = sig
       type 'a chan
2
3
       val channel : unit -> 'a chan
       val send : 'a chan * 'a -> unit
4
       val recv : 'a chan -> 'a
5
 6
       end
1
2
     structure ManyToManyChan : CHAN = struct
3
       type message_queue = 'a option ref queue
4
5
       datatype 'a chan_content =
6
         Send of (condition * 'a) queue |
7
         Recv of (condition * 'a option ref) queue |
8
         Inac
9
10
       datatype 'a chan =
11
         Chn of 'a chan_content ref * mutex_lock
12
13
       fun channel () = Chn (ref Inac, mutexLock ())
14
       fun send (Chn (conRef, lock)) m =
15
16
         acquire lock;
17
         (case !conRef of
18
           Recv q \Rightarrow let
19
             val (recvCond, mopRef) = dequeue q in
20
             mopRef := Some m;
21
             if (isEmpty q) then conRef := Inac else ();
22
             release lock; signal recvCond; () end |
23
           Send q \Rightarrow let
24
             val sendCond = condition () in
25
             enqueue (q, (sendCond, m));
26
             release lock; wait sendCond; () end |
27
           Inac => let
28
             val sendCond = condition () in
29
             conRef := Send (queue [(sendCond, m)]);
30
             release lock; wait sendCond; () end)
31
32
       fun recv (Chn (conRef, lock)) =
33
         acquire lock;
34
         (case !conRef of
35
           Send q \Rightarrow let
```

```
36
             val (sendCond, m) = dequeue q in
37
             if (isEmpty q) then
38
               conRef := Inac
39
             else
40
               ();
41
             release lock; signal sendCond; m end |
42
           Recv q => let
43
             val recvCond = condition ()
44
             val mopRef = ref None in
45
             enqueue (q, (recvCond, mopRef));
             release lock; wait recvCond;
46
             valOf (!mopRef) end |
47
           Inac => let
48
49
             val recvCond = condition ()
50
             val mopRef = ref None in
51
             conRef := Recv (queue [(recvCond, mopRef)]);
52
             release lock; wait recvCond;
53
             valOf (!mopRef) end)
54
55
       end
56
57
1
2
       structure FanOutChan : CHAN = struct
3
4
       datatype 'a chan_content =
         Send of condition * 'a |
5
6
         Recv of (condition * 'a option ref) queue \mid
7
         Inac
8
9
       datatype 'a chan =
10
         Chn of 'a chan_content ref * mutex_lock
11
12
       fun channel () = Chn (ref Inac, mutexLock ())
13
       fun send (Chn (conRef, lock)) m = let
14
15
         val sendCond = condition () in
16
         case cas (conRef, Inac, Send (sendCond, m)) of
17
           Inac => (* conRef already set *)
             wait sendCond; () |
18
19
           Recv q =>
20
           (* the current thread is
21
             * the only one that updates from this state *)
22
             acquire lock;
23
             (let
24
               val (recvCond, mopRef) = dequeue q in
25
               mopRef := Some m;
               if (isEmpty q) then conRef := Inac else ();
26
27
               release lock; signal (recvCond);
```

```
28
               () end) |
29
           Send _ => raise NeverHappens end
30
31
       fun recv (Chn (conRef, lock)) =
32
         acquire lock;
33
         (case !conRef of
34
           Inac => let
35
             val recvCond = condition ()
36
             val mopRef = ref None in
37
             conRef := Recv (queue [(recvCond, mopRef)]);
38
             release lock; wait recvCond;
39
             valOf (!mopRef) end |
40
           Recv q => let
41
             val recvCond = condition ()
42
             val mopRef = ref None in
43
             enqueue (q, (recvCond, mopRef));
44
             release lock; wait recvCond;
45
             valOf (!mopRef) end |
46
           Send (sendCond, m) =>
47
             conRef := Inac;
48
             release lock;
49
             signal sendCond;
50
             m end)
51
52
       end
53
    structure FanInChan : CHAN = struct
1
2
3
    datatype 'a chan_content =
4
       Send of (condition * 'a) queue |
5
       Recv of condition * 'a option ref |
6
       Inac
7
8
    datatype 'a chan =
       Chn of 'a chan_content ref * mutex_lock
9
10
11
    fun channel () = Chn (ref Inac, mutexLock ())
12
13
    fun send (Chn (conRef, lock)) m =
14
       acquire lock;
15
       case !conRef of
16
       Recv (recvCond, mopRef) =>
17
         mopRef := Some m; conRef := Inac;
18
         release lock; signal recvCond;
19
         () |
20
       Send q => let
21
         val sendCond = condition () in
22
         enqueue (q, (sendCond, m));
23
         release lock; wait sendCond;
```

```
24
         () end |
25
       Inac => let
26
         val sendCond = condition () in
27
         conRef := Send (queue [(sendCond, m)])
28
         release lock; wait sendCond; () end
29
30
    fun recv (Chn (conRef, lock)) = let
31
       val recvCond = condition ()
32
       val mopRef = ref None in
33
       case cas (conRef, Inac, Recv (recvCond, mopRef)) of
         Inac => (* conRef already set *)
34
           wait recvCond; valOf (!mopRef) |
35
36
         Send q \Rightarrow
37
           (* the current thread is the only one
38
           -* that updates the state from this state *)
39
           acquire lock;
40
           (let
41
             val (sendCond, m) = dequeue q in
42
             if (isEmpty q) then conRef := Inac else ();
43
             release lock; signal sendCond; m end) |
44
         Recv _ => raise NeverHappens end end
45
46
2 \text{ structure OneToOneChan} : CHAN = struct
3
    datatype 'a chan_content =
4
5
       Send of condition * 'a |
6
       Recv of condition * 'a option ref |
7
       Inac
8
9
    datatype 'a chan = Chn of 'a chan_content ref
10
11
    fun channel () = Chn (ref Inac)
12
     fun send (Chn conRef) m = let
13
14
       val sendCond = condition () in
15
       case cas (conRef, Inac, Send (sendCond, m)) of
16
         Inac =>
17
           (* conRef already set to Send *)
18
           wait sendCond; () |
19
         Recv (recvCond, mopRef) =>
20
           (* the current thread is the only one
21
           -* that accesses conRef for this state *)
22
           mopRef := Some m; conRef := Inac;
23
           signal recvCond; () |
24
         Send _ => raise NeverHappens end end
25
26
```

```
27
    fun recv (Chn conRef) = let
28
       val recvCond = condition ();
29
      val mopRef = ref None in
30
       case cas (conRef, Inac, Recv (recvCond, mopRef)) of
31
         Inac => (* conRef already set to Recv*)
32
           wait recvCond; valOf (!mopRef) |
33
         Send (sendCond, m) =>
34
           (* the current thread is the only one
35
           -* that accesses conRef for this state *)
36
           conRef := Inac; signal sendCond; m |
37
         Recv _ => raise NeverHappens end end
38
39
    end
40
1
    structure OneShotChan : CHAN = struct
2
3
    datatype 'a chan_content =
       Send of condition * 'a |
4
5
      Recv of condition \ast 'a option ref |
6
      Inac
7
8
    datatype 'a chan = Chn of 'a chan_content ref * mutex_lock
9
10
    fun channel () = Chn (ref Inac, lock ())
11
12
    fun send (Chn (conRef, lock)) m = let
13
      val sendCond = condition () in
14
       case (conRef, Inac, Send (sendCond, m)) of
15
         Inac =>
16
           (* conRef already set to Send*)
17
           wait sendCond; () |
18
         Recv (recvCond, mopRef) =>
19
           mopRef := Some m; signal recvCond;
20
21
         Send _ => raise NeverHappens end end
22
23
24
    fun recv (Chn (conRef, lock)) = let
25
      val recvCond = condition ()
26
       val mopRef = ref None in
27
       case (conRef, Inac, Recv (recvCond, mopRef)) of
28
        Inac =>
29
           (* conRef already set to Recv*)
30
           wait recvCond; valOf (!mopRef) |
31
         Send (sendCond, m) =>
32
           acquire lock; signal sendCond;
33
           (* never relases lock;
34
           -* blocks others forever *)
           m |
35
```

```
Recv _ =>
36
37
           acquire lock;
38
           (* never able to acquire lock;
39
           -* blocked forever *)
40
          raise NeverHappens end end
41
42
    end
43
1 structure OneShotToOneChan : CHAN = struct
3
    datatype 'a chan =
      Chn of condition * condition * 'a option ref
4
5
6
    fun channel () =
7
      Chn (condition (), condition (), ref None)
8
9
    fun send (Chn (sendCond, recvCond, mopRef)) m =
10
      mopRef := Some m; signal recvCond;
11
      wait sendCond; ()
12
    fun recv (Chn (sendCond, recvCond, mopRef)) =
13
14
      wait recvCond; signal sendCond;
      valOf (!mopRef)
15
16
17
    end
18
```

2 Syntax

```
2
     datatype var = Var string
 3
 4
     datatype exp =
       Let var boundexp exp |
 5
 6
       Rslt var
 7
     boundexp =
 8
9
       Unt |
10
       MkChn |
11
       Prim prim |
12
       Spwn exp |
13
       Sync var |
14
       Fst var |
       Snd var |
15
16
       Case var var exp var exp |
17
       App var var and
18
19
     prim =
```

```
20 SendEvt var var |
21 RecvEvt var |
22 Pair var var |
23 Lft var |
24 Rht var |
25 Abs var var ex
```

3 Dynamic Semantics

```
1
     datatype ctrl_label =
 2
        LNxt var | LSpwn var | LCall var | LRtn var
 3
 4
     type_synonym ctrl_path = (ctrl_label list)
 5
 6
     datatype chan = Chn ctrl_path var
 7
 8
     datatype val =
9
        {\tt VUnt \ | \ VChn \ chan \ | \ VClsr \ prim \ (var \rightharpoonup val)}
10
11
     datatype ctn = Ctn var exp (var \rightarrow val)
12
13
     datatype state = Stt exp (var \rightarrow val) (ctn list)
14
15
 1
 2
     inductive seq_step : :
 3
        bind * (var \rightarrow val)) \Rightarrow val \Rightarrow bool where
 4
        LetUnt :
5
          seq_step (Unt, env) VUnt |
 6
        LetPrim :
 7
          seq_step (Prim p, env) (VClsr p env) |
 8
        LetFst :
 9
          env xp = Some (VClsr (Pair x1 x2) envp) \Longrightarrow
10
          envp x1 = Some v \Longrightarrow
11
          seq_step (Fst xp, env) v |
12
        LetSnd :
13
          env xp = Some (VClsr (Pair x1 x2) envp) \Longrightarrow
14
          envp x2 = Some v \Longrightarrow
15
          seq_step (Snd xp, env) v
16
17
18
 1
 2
 3
     inductive seq_step_up : :
        bind * (var \rightharpoonup val)) \Rightarrow exp * val_env \Rightarrow bool where
 4
        LetCaseLft :
```

```
6
           env xs = Some (VClsr (Lft xl') envl) \Longrightarrow
 7
           envl xl' = Some vl \Longrightarrow
8
           seq_step_up
9
              (Case xs xl el xr er, env)
10
              (el, env(xl \mapsto vl)) |
11
         {\tt LetCaseRht} \; : \;
12
           env xs = Some (VClsr (Rht xr') envr) \Longrightarrow
13
           envr xr' = Some vr \Longrightarrow
14
           seq_step_up
15
              (Case xs xl el xr er, env)
16
              (er, env(xr \mapsto vr)) |
17
         LetApp :
           env f = Some (VClsr (Abs fp xp el) envl) \Longrightarrow
18
19
           env xa = Some va \Longrightarrow
20
           seq_step_up
21
              (App f xa, env)
22
              (el, envl(
23
                 fp \mapsto (VClsr (Abs fp xp el) envl),
24
                 xp \mapsto va))
25
26
1
 2
 3
      type_synonym cmmn_set = (ctrl_path * chan * ctrl_path) set
 5
      6
7
      inductive leaf : :
 8
         {\tt trace\_pool} \ \Rightarrow \ {\tt ctrl\_path} \ \Rightarrow \ {\tt bool} \ \ {\tt where}
9
         intro :
10
           \mathtt{trpl}\ \mathtt{pi}\ \neq\ \mathtt{None}\ \Longrightarrow
11
            (\nexists pi' . trpl pi' \neq None \land strict_prefix pi pi') \Longrightarrow
12
           leaf trpl pi
13
14
 1
 2
      inductive concur_step : :
 3
        {\tt trace\_pool} \; * \; {\tt cmmn\_set} \; \Rightarrow \;
         {\tt trace\_pool} \;\; * \;\; {\tt cmmn\_set} \;\; \Rightarrow \;\;
4
5
        bool where
 6
         Seq_Sttep_Down:
 7
           leaf trpl pi \Longrightarrow
 8
           trpl pi = Some
9
              (Stt (Rslt x) env
10
                 ((Ctn xk ek envk) # k)) \Longrightarrow
11
           env x = Some v \Longrightarrow
12
           concur_step
13
              (trpl, ys)
```

```
14
              (trpl(pi @ [LRtn xk] \mapsto
15
                 (Stt ek (envk(xk \mapsto v)) k)), ys) |
16
        Seq_Step :
17
           \texttt{leaf trpl pi} \implies
18
           trpl pi = Some
19
              (Stt (Let x b e) env k) \Longrightarrow
20
           seq\_step (b, env) v\Longrightarrow
21
           concur_step
22
              (trpl, ys)
23
              (trpl(pi 0 [LNxt x] \mapsto
24
                (Stt e (env(x \mapsto v)) k), ys) |
25
        Seq_Step_Up :
26
           \texttt{leaf trpl pi} \Longrightarrow
27
           trpl pi = Some
28
              (Stt (Let x b e) env k) \Longrightarrow
29
           seq\_step\_up (b, env) (e', env') \Longrightarrow
30
           concur_step
31
              (trpl, ys)
32
              (trpl(pi @ [LCall x] \mapsto
33
                (Stt e' env'
34
                   ((Ctn x e env) # k))), ys) |
35
        LetMkCh :
36
           \texttt{leaf trpl pi} \implies
37
           trpl pi = Some (Stt (Let x MkChn e) env k) \Longrightarrow
38
           concur_step
39
              (trpl, ys)
40
              (trpl(pi @ [LNxt x] \mapsto
41
                 (Stt e (env(x \mapsto (VChn (Chn pi x)))) k)), ys) |
42
        LetSpwn :
43
           \texttt{leaf trpl pi} \implies
44
           trpl pi = Some
45
              (Stt (Let x (Spwn ec) e) env k) \Longrightarrow
46
           concur_step
47
              (trpl, ys)
48
              (trpl(
49
                pi @ [LNxt x] \mapsto
50
                     (St e (env(x \mapsto VUnt)) k),
51
                pi @ [LSpwn x] \mapsto
52
                     (St ec env []), ys) |
53
        LetSync :
54
           \texttt{leaf trpl pis} \Longrightarrow
55
           trpl pis = Some
56
              (Stt (Let xs (Sync xse) es) envs ks) \Longrightarrow
57
           envs xse = Some
58
              (VClsr (SendEvt xsc xm) envse) \Longrightarrow
59
           leaf trpl pir \Longrightarrow
           trpl pir = Some
60
61
              (Stt (Let xr (Sync xre) er) envr kr) \Longrightarrow
62
           envr xre = Some
63
              (VClsr (RecvEvt xrc) envre) \Longrightarrow
```

```
64
            envse xsc = Some (VChn c) \Longrightarrow
65
            envre xrc = Some (VChn c) \Longrightarrow
66
            envse xm = Some vm \Longrightarrow
67
            concur_step
68
               (trpl, ys)
69
               (trpl(
70
                 pis @ [LNxt xs] \mapsto
71
                    (Stt es (envs(xs \mapsto VUnt)) ks),
72
                 pir 0 [LNxt xr] \mapsto
73
                    (Stt er (envr(xr \mapsto vm)) kr)),
74
                 ys \cup \{(pis, c, pir)\})
75
76
      inductive star : ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow 'a \Rightarrow 'a \Rightarrow bool
1
 2
         for r where
 3
         refl : star r x x |
         step : r x y \Longrightarrow star r y z \Longrightarrow star r x z
```

4 Dynamic Communication

```
1
      inductive is_send_path : :
 2
         {\tt trace\_pool} \; \Rightarrow \; {\tt chan} \; \Rightarrow \;
 3
         control_path \Rightarrow bool where
 4
         intro :
 5
            trpl piy = Some
               (Stt (Let xy (Sync xe) en) env k) \Longrightarrow
 7
            env xe = Some
               (VClsr (SendEvt xsc xm) enve) \Longrightarrow
 8
9
            enve xsc = Some (VChn c) \Longrightarrow
10
            is_send_path trpl c piy
11
12
      inductive is_recv_path : :
13
         \texttt{trace\_pool} \ \Rightarrow \ \texttt{chan} \ \Rightarrow
14
         control_path \Rightarrow bool where
15
         intro :
16
            trpl piy = Some
17
               (Stt (Let xy (Sync xe) en) env k) \Longrightarrow
18
            env xe = Some
19
               (VClsr (RecvEvt xrc) enve) \Longrightarrow
20
            enve xrc = Some (VChn c) \Longrightarrow
21
            is_recv_path trpl c piy
22
23
1
 2
      inductive every_two : :
         ('a \Rightarrow bool) \Rightarrow
```

```
('a \Rightarrow 'a \Rightarrow bool) \Rightarrow
 4
 5
          bool where
 6
          intro : (\forall pi1 pi2 .
 7
               p x1 \longrightarrow
 8
               \texttt{p} \quad \texttt{x2} \quad \longrightarrow \quad
 9
               r x1 x2) \Longrightarrow
10
             every_two p r
11
12
       inductive ordered : : 'a list \Rightarrow 'a list \Rightarrow bool where
13
          left : prefix pi1 pi2 \Longrightarrow ordered pi1 pi2 |
          \texttt{right} \; : \; \texttt{prefix} \; \texttt{pi2} \; \texttt{pi1} \implies \texttt{ordered} \; \texttt{pi1} \; \texttt{pi2}
14
15
16
 1
 2
       inductive one_shot : : trace_pool \Rightarrow chan \Rightarrow bool where
 3
          intro :
             every_two
 5
                (is_send_path trpl c) op= \Longrightarrow
 6
             one_shot trpl c
 7
 8
       inductive fan_out : : trace_pool \Rightarrow chan \Rightarrow bool where
 9
          intro :
10
             every_two
11
                (is_send_path trpl c) ordered \Longrightarrow
12
             fan_out trpl c
13
14
       inductive fan_in : : trace_pool \Rightarrow chan \Rightarrow bool where
15
          intro :
16
             every_two
17
                (is_recv_path trpl c) ordered \Longrightarrow
18
             fan_in trpl c
19
20
       inductive one_to_one : : trace_pool \Rightarrow chan \Rightarrow bool where
21
          intro :
22
             fan_out trpl c \Longrightarrow
23
             fan_in trpl c \Longrightarrow
24
             one_to_one trpl c
25
26
```

5 Static Semantics

```
1
2   datatype abstract_value =
3     AChn var |
4     AUnt |
5     APrim prim
6
```

```
7
      \texttt{type\_synonym} \ \ \texttt{abstract\_env} \ = \ \texttt{var} \ \Rightarrow \ \texttt{abstract\_value} \ \ \texttt{set}
8
9
      fun rslt_var : : exp \Rightarrow var where
10
         rslt_var (Rslt x) = x |
11
         rslt_var (Let _ _ e) = (rslt_var e)
12
13
 1
 2
 3
      inductive may_be_eval_exp : :
         \verb|abstract_env| * \verb|abstract_env| \Rightarrow
 4
 5
         exp \Rightarrow bool where
         Result :
 6
 7
            may_be_eval_exp (V, C) (RESULT x) |
 8
         let_Unt :
 9
            {AUnt} \subseteq V x \Longrightarrow
10
            may_be_eval_exp (V, C) e \Longrightarrow
11
            may_be_eval_exp (V, C) (Let x Unt e) |
12
         let_Chan :
13
            \{\mathtt{AChn}\ \mathtt{x}\}\ \subseteq\ \mathtt{V}\ \mathtt{x}\quad\Longrightarrow\quad
14
            may_be_eval_exp (V, C) e \Longrightarrow
15
            may_be_eval_exp (V, C) (Let x (MkChn) e) |
16
         let_SendEvt :
17
            \texttt{\{APrim (SendEvt xc xm)\}} \subseteq \texttt{V} \texttt{ x} \Longrightarrow
18
            may_be_eval_exp (V, C) e \Longrightarrow
19
            may_be_eval_exp (V, C) (Let x (Prim (SendEvt xc xm)) e
        ) |
20
         let_RecvEvt :
21
            {APrim (RecvEvt xc)} \subseteq V x \Longrightarrow
22
            may_be_eval_exp (V, C) e \Longrightarrow
23
            may_be_eval_exp (V, C) (Let x (Prim (RecvEvt xc)) e) |
24
         let_Pair :
25
            \{\texttt{APrim (Pair x1 x2)}\} \subseteq \texttt{V x} \implies
26
            \verb"may_be_eval_exp" (V, C) e \implies
            may_be_eval_exp (V, C) (Let x (Pair x1 x2) e) |
27
28
         let_Left :
29
            \texttt{\{APrim (Left xp)\}} \subseteq \texttt{V} \texttt{ x} \implies
30
            \verb"may_be_eval_exp" (V, C) e \implies
31
            may_be_eval_exp (V, C) (Let x (Left xp) e) |
32
         let_Right :
33
            {APrim (Right xp)} \subseteq V x \Longrightarrow
34
            may_be_eval_exp (V, C) e \Longrightarrow
35
            may_be_eval_exp (V, C) (Let x (Right xp) e) |
36
         let_Abs :
37
            {APrim (Abs f' x' e')} \subseteq V f' \Longrightarrow
38
            may_be_eval_exp (V, C) e' \Longrightarrow
39
            {APrim (Abs f' x' e')} \subseteq V x \Longrightarrow
40
            may_be_eval_exp (V, C) e \Longrightarrow
            may_be_eval_exp (V, C) (Let x (Abs f' x' e') e) |
41
```

```
42
           let_Spawn :
43
              \texttt{{AUnt}} \subseteq \texttt{V} \texttt{ x} \implies
44
              may_be_eval_exp (V, C) ec \Longrightarrow
45
              \verb"may_be_eval_exp" (V, C) e \implies
46
              may_be_eval_exp (V, C) (Let x (Spwn ec) e) |
47
           let_Sync :
48
              \forall xsc xm xc .
49
                 (APrim (SendEvt xsc xm)) \in V xe \longrightarrow
50
                 \mathtt{AChn} \ \mathtt{xc} \ \in \ \mathtt{V} \ \mathtt{xsc} \ \longrightarrow
51
                 \texttt{{AUnt}} \; \subseteq \; \texttt{V} \; \; \texttt{x} \; \land \; \; \texttt{V} \; \; \texttt{xm} \; \subseteq \; \texttt{C} \; \; \texttt{xc} \; \Longrightarrow \;
              \forall xrc xc
52
                 (APrim (RecvEvt xrc)) \in V xe \longrightarrow
53
54
                 \mathtt{AChn} \ \mathtt{xc} \ \in \ \mathtt{V} \ \mathtt{xrc} \ \longrightarrow
55
                 \texttt{C} \ \texttt{xc} \ \subseteq \ \texttt{V} \ \texttt{x} \implies
56
              may_be_eval_exp (V, C) e \Longrightarrow
57
              may_be_eval_exp (V, C) (Let x (Syync xe) e) |
58
           let_Fst :
59
              \forall x1 x2.
60
                 (APrim (Pair x1 x2)) \in V xp \longrightarrow
61
                 V x1 \subseteq V x \Longrightarrow
62
              may_be_eval_exp (V, C) e \Longrightarrow
63
              may_be_eval_exp (V, C) (Let x (Fst xp) e) |
64
           let_Snd:
65
          \forall x1 x2 .
66
              (APrim (Pair x1 x2) \in V xp \longrightarrow
              \tt V x2 \subseteq \tt V x \Longrightarrow
67
68
          may_be_eval_exp (V, C) e \Longrightarrow
69
           may_be_eval_exp (V, C) (Let x (Snd xp) e) |
70
       let_Case:
71
          \forall x1'.
72
              (APrim (Left xl')) \in V xs \longrightarrow
73
                 V xl' \subseteq V xl \wedge V (rslt_var el) \subseteq V x \wedge
74
                 may_be_eval_exp (V, C) el \Longrightarrow
75
          \forall xr, .
76
              (APrim (Right xr')) \in V xs \longrightarrow
77
                 V xr' \subseteq V xr \wedge V (rslt_var er) \subseteq V x \wedge
78
                 may_be_eval_exp (V, C) er \Longrightarrow
79
                 may_be_eval_exp (V, C) e \Longrightarrow
80
              may_be_eval_exp (V, C) (Let x (Case xs xl el xr er) e)
       let_App :
81
82
          \forall f' x' e' .
83
              (APrim (Abs f' x' e') \in V f \longrightarrow
84
              V xa \subseteq V x' \wedge
85
              V (rslt_var e') \subseteq V x \Longrightarrow
86
          may_be_eval_exp (V, C) e \Longrightarrow
87
          may_be_eval_exp (V, C) (Let x (App f xa) e)
88
89
```

```
1
 2
 3
      \texttt{fun abstract : val} \ \Rightarrow \ \texttt{abstract\_value where}
         abstract VUnt = AUnt |
 4
 5
         abstract VChn (Chn pi x) = AChn x \mid
 6
         abstract VClsr p env = APrim p
 7
 8
 1
 2
 3
 4
      inductive
 5
        may_be_eval_val : :
 6
           \verb|abstract_env| * \verb|abstract_env| \Rightarrow \verb|val| \Rightarrow \verb|bool| and
 7
        may_be_eval_env : :
 8
           \verb|abstract_env| * \verb|abstract_env| \Rightarrow \verb|val_env| \Rightarrow \verb|bool| where|
 9
10
           may_be_eval_val (V, C) VUnt |
11
         Chan:
12
           may_be_eval_val (V, C) VChn c |
13
         SendEvt:
           may_be_eval_env (V, C) env \Longrightarrow
14
           may_be_eval_val (V, C) (VClsr (SendEvt _ _) env) |
15
16
17
           may_be_eval_env (V, C) env \Longrightarrow
18
           may_be_eval_val (V, C) (VClsr (RecvEvt _) env) |
19
        Left :
20
           \verb"may_be_eval_env" (V, C) env \implies
21
           may_be_eval_val (V, C) (VClsr (Left _) env) |
22
23
           may_be_eval_env (V, C) env \Longrightarrow
24
           may_be_eval_val (V, C) (VClsr (Right _) env) |
25
         Abs :
26
           \{(\mathtt{APrim}\ (\mathtt{Abs}\ \mathtt{f}\ \mathtt{x}\ \mathtt{e})\}\ \subseteq\ \mathtt{V}\ \mathtt{f} \implies
27
           may_be_eval_exp (V, C) e \Longrightarrow
           may\_be\_eval\_env (V, C) env \Longrightarrow
28
           may_be_eval_val (V, C) (VClsr (Abs f x e) env) |
29
30
        Pair :
31
           \verb"may_be_eval_env" (V, C) env \implies
32
           may_be_eval_val (V, C) (VClsr (Pair _ _) env) |
33
         intro :
34
           \forall x v .
35
              env x = Some v \longrightarrow
36
              {abstract v} \subseteq V x \land may_be_eval_val (V, C) v \Longrightarrow
37
           may_be_eval_env (V, C) env
38
39
```

1

```
inductive may_be_eval_stack : :
 3
         abstract_env * abstract_env \Rightarrow
 4
         \verb|abstract_value| \verb|set| \Rightarrow \verb|cont| \verb|list| \Rightarrow \verb|bool| \verb|where||
 5
        Empty :
 6
           may_be_eval_stack (V, C) valset [] |
 7
        Nonempty:
 8
           \mathtt{valset} \ \subseteq \ \mathtt{V} \ \mathtt{x} \implies
 9
           may_be_eval_exp (V, C) e \Longrightarrow
10
           may\_be\_eval\_env (V, C) env \Longrightarrow
11
           \verb"may_be_eval_stack" (V, C) (V (rslt_var e)) k \Longrightarrow
           may_be_eval_stack (V, C) valset ((Ctn x e env) # k)
12
13
14
15
      inductive may_be_eval_state : :
16
         abstract_env * abstract_env \Rightarrow
17
         state \Rightarrow bool where
18
        intro :
19
           may\_be\_eval\_exp (V, C) e \Longrightarrow
20
           may_be_eval_env (V, C) env \Longrightarrow
21
           may_be_eval_stack (V, C) (V (rslt_var e)) k \Longrightarrow
22
           may_be_eval_state (V, C) (Stt e env k)
23
24
      inductive may_be_eval_pool : :
25
         abstract_env * abstract_env ⇒
26
         trace\_pool \Rightarrow bool where
27
        intro :
28
           \forall pi st .
29
              \texttt{trpl pi = Some st} \,\longrightarrow\,
30
             \verb"may_be_eval_state" (V, C) st \Longrightarrow
31
           may_be_eval_pool (V, C) trpl
32
33
 1
 2
      theorem may_be_eval_preserved_under_concur_step : "
        \verb"may_be_eval_pool" (V, C) trpl \Longrightarrow
 3
        concur_step (trpl, ys) (trpl', ys') \Longrightarrow
 4
        may_be_stati_eval_pool (V, C) trpl'"
 5
 6
      proof outline
 7
      qed
 8
9
      theorem may_be_eval_preserved_under_concur_step_star : "
10
        may_be_eval_pool (V, C) trpl \Longrightarrow
11
        star concur_step (trpl, ys) (trpl', ys') \Longrightarrow
12
        may_be_concur_step (V, C) trpl'"
13
      proof outline
14
      qed
15
```

1

```
theorem trace_pool_snapshot_value_not_bound_sound : "
 3
         env x = Some v \Longrightarrow
 4
         trpl pi = Some (Stt e env k) \Longrightarrow
 5
         \verb"may_be_eval_pool" (V, C) trpl \Longrightarrow
 6
         \{\texttt{abstract} \ \texttt{v}\} \ \subseteq \ \texttt{V} \ \texttt{x} \ \texttt{"}
 7
      proof outline
 8
      qed
 9
 1
 2
      theorem trace_pool_always_value_not_bound_sound : "
 3
         env' x = Some v \Longrightarrow
         may_be_eval_pool (V, C) trpl \Longrightarrow
 4
 5
         star concur_step (trpl, ys) (trpl', ys') \Longrightarrow
 6
         trpl' pi = Some (Stt e' env' k') \Longrightarrow
 7
         \{\texttt{abstract} \ \ \texttt{v}\} \ \subseteq \ \texttt{V} \ \ \texttt{x"}
 8
      proof outline
 9
      qed
10
11
      theorem exp_always_value_not_bound_sound : "
12
         env' x = Some v \Longrightarrow
13
         may\_be\_eval\_exp (V, C) e \Longrightarrow
         star concur_step
14
            ([[] \mapsto (Stt e (\lambda _ . None) [])], ys)
15
16
            (trpl', ys') \Longrightarrow
         trpl' pi = Some (Stt e' env' k') \Longrightarrow
17
18
         \{\texttt{abstract} \ \texttt{v}\} \ \subseteq \ \texttt{V} \ \texttt{x"}
19
      proof outline
20
      qed
21
22
1
      inductive is_super_exp : : exp \Rightarrow exp \Rightarrow bool where
 2
         Refl :
 3
            is_super_exp e e |
 4
         let_Spawn_Child
 5
            is\_super\_exp\ ec\ e \Longrightarrow
 6
            is_super_exp (Let x (Spwn ec) en) e |
 7
         let_Case_Left :
 8
            is\_super\_exp\ el\ e \Longrightarrow
 9
            is_super_exp (Let x (case xs xl el xr er) en) e |
10
         let_Case_Right :
11
            is\_super\_exp er e \Longrightarrow
12
            is_super_exp (Let x (case xs xl el xr er) en) e |
13
         let_Abs_Body : "
14
            \verb"is_super_exp" eb e \Longrightarrow
15
            is_super_exp (Let x (Abs f xp eb) en) e |
16
         Let :
17
            is\_super\_exp en e \Longrightarrow
18
            is_super_exp (Let x b en) e
```

```
19
20
     inductive is_super_exp_left : : exp \Rightarrow exp \Rightarrow bool where
21
       Refl:
22
          is_super_exp_left e0 e0 |
23
       let_Spawn_Child :
24
          is_super_exp_left e0 (Let x (Spwn ec) en)\Longrightarrow
25
          is_super_exp_left e0 ec |
26
       let_Case_Left :
         is_super_exp_left e0 (Let x (case xs xl el xr er) en)
27
28
          is_super_exp_left e0 el |
29
       let_Case_Right : "
30
         is_super_exp_left e0 (Let x (case xs xl el xr er) en)
31
         is_super_exp_left e0 er |
32
       let_Abs_Body :
33
         is_super_exp_left e0 (Let x (Abs f xp eb) en) \Longrightarrow
34
          is_super_exp_left e0 eb |
35
       Let :
36
          is_super_exp_left e0 (Let x b en) \Longrightarrow
37
          is_super_exp_left e0 en
38
39
     inductive is_super_exp_over_prim : : exp \Rightarrow prim \Rightarrow bool
       where
40
       SendEvt :
41
          is_super_exp_over_prim e0 (SendEvt xC xM) |
42
       RecvEvt :
43
          is_super_exp_over_prim e0 (RecvEvt xC) |
44
       Pair :
45
          is_super_exp_over_prim e0 (Pair x1 x2) |
46
       Left :
47
          is_super_exp_over_prim e0 (Left x) |
48
       Right : "
49
          is_super_exp_over_prim e0 (Right x) |
50
       Abs:
51
          \verb"is_super_exp_left" e0 eb \Longrightarrow
52
          is_super_exp_over_prim e0 (Abs fp xp eb)
53
54
     inductive
55
       is_super_exp_over_env : : exp \Rightarrow env \Rightarrow bool and
56
       is_super_exp_over_val : : exp \Rightarrow val \Rightarrow bool where
57
       VUnt :
58
          is_super_exp_over_val e0 VUnt |
59
       VChn:
60
          is_super_exp_over_val e0 (VChn c) |
61
       VClsr : "
62
          \verb"is_super_exp_over_prim" = 0 p \implies
63
          is_super_exp_over_env e0 env' ⇒
64
          is_super_exp_over_val e0 (VClsr p env') |
65
       intro : "
```

```
66
           \forall x v .
67
              env x = Some v \longrightarrow
68
              \verb"is_super_exp_over_val" = 0 v \implies
69
           is_super_exp_over_env e0 env
70
71
      inductive is_super_exp_over_stack : : exp \Rightarrow cont list \Rightarrow
        bool where
72
        Empty:
73
           is_super_exp_over_stack e0 [] |
74
        Nonempty :
75
           is_super_exp_left e0 ek \Longrightarrow
76
           is\_super\_exp\_over\_env e0 envk \Longrightarrow
77
           is\_super\_exp\_over\_stack e0 k \Longrightarrow
78
           is_super_exp_over_stack e0 ((Ctn xk ek envk) # k)
79
80
     inductive is_super_exp_over_state : : exp \Rightarrow state \Rightarrow bool
        where
81
        intro :
82
           is_super_exp_left e0 e \Longrightarrow
83
           is_super_exp_over_env e0 env \Longrightarrow
84
           is\_super\_exp\_over\_stack e0 k \Longrightarrow
85
           is_super_exp_over_state e0 (Ctn e env k)
86
87
 1
     lemma is_super_exp_trans : "
 2
        \verb"is_super_exp" ez" ey" \Longrightarrow
 3
        \verb"is_super_exp" ey ex \Longrightarrow
 4
        is_super_exp ez ex"
     proof outline
 5
 6
      qed
 7
 8
9
      lemma is_super_exp_over_state_preserved : "
10
        concur_step (trpl, ys) (trpl', ys') \Longrightarrow
11
        \forall pi st.
12
           trpl pi = Some st \longrightarrow
13
           is\_super\_exp\_over\_state e0 st \Longrightarrow
14
        trpl' pi' = Some st' \Longrightarrow
15
        is_super_exp_over_state e0 st'"
16
      proof outline
17
      qed
18
 2
      lemma state_always_exp_not_reachable_sound : "
 3
        star concur_step (trpl0, ys0) (trpl', ys') \Longrightarrow
 4
        trpl0 = [[] \mapsto (Stt e0 (\lambda _ . None) [])] \Longrightarrow
 5
        trpl' pi' = Some st' \Longrightarrow
        is_super_exp_over_state e0 st' "
```

```
7
     proof outline
8
     qed
9
10
     lemma exp_always_exp_not_reachable_sound : "
11
       star concur_step
       ([[] \mapsto (Stt e0 (\lambda _ . None) [])], {})
12
13
          (trpl', H') \Longrightarrow
14
       trpl, pi, = Some (Stt e, env, k,)) \Longrightarrow
15
       is_super_exp e0 e' "
16
     proof outline
17
     qed
18
19
```