

Efficient Convex Relaxations for Streaming PCA

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1. Streaming PCA in Stochastic Setting

- Given a sequence of vectors $(\mathbf{x}_t)_{t=1}^\infty$ i.i.d. $\mathbf{x}_t \sim \mathcal{D}$.
- Minimize reconstruction error $\mathbb{E}_{\mathbf{x} \sim \mathcal{D}} \|\mathbf{x} - \mathbf{P}\mathbf{x}\|^2$, where $\mathbf{P} \in \mathcal{P}_k = \{\mathbf{P} : \mathbf{P}^2 = \mathbf{P}, \mathbf{P}^\top = \mathbf{P}, \text{rank}(\mathbf{P}) = k\}$.

- Equivalently solve:

$$\begin{aligned} & \underset{\mathbf{P} \in \mathbb{R}^{d \times d}}{\text{maximize}} && \mathbb{E}_{\mathbf{x} \sim \mathcal{D}} \langle \mathbf{P}, \mathbf{x}\mathbf{x}^\top \rangle \\ & \text{subject to} && \mathbf{P}^2 = \mathbf{P}, \mathbf{P}^\top = \mathbf{P}, \text{rank}(\mathbf{P}) = k \end{aligned} \quad (1)$$

- Usually solved by Oja's algorithm:

$$\begin{aligned} \mathbf{U}_{t+1} &\leftarrow (\mathbf{I} + \eta_t \mathbf{x}_t \mathbf{x}_t^\top) \mathbf{U}_t, \mathbf{U}_{t+1} \leftarrow \text{Orth}(\mathbf{U}_{t+1}) \\ \mathbf{P}_{t+1} &\leftarrow \mathbf{U}_{t+1} \mathbf{U}_{t+1}^\top \end{aligned}$$

2. A Convex Relaxation

- Convexify $\mathcal{P}_k \rightarrow \mathcal{C} = \{\mathbf{P} : \text{Tr}(\mathbf{P}) = k, 0 \preceq \mathbf{P} \preceq \mathbf{I}, \mathbf{P} = \mathbf{P}^\top\}$.
- Can now solve the following convex optimization problem:

$$\begin{aligned} & \underset{\mathbf{P} \in \mathbb{R}^{d \times d}}{\text{maximize}} && \mathbb{E}_{\mathbf{x} \sim \mathcal{D}} \langle \mathbf{P}, \mathbf{x}\mathbf{x}^\top \rangle \\ & \text{subject to} && \mathbf{P} \in \mathcal{C} \end{aligned} \quad (2)$$

- Can also add regularization to achieve distribution dependent guarantees:

$$\begin{aligned} & \underset{\mathbf{P} \in \mathbb{R}^{d \times d}}{\text{maximize}} && \mathbb{E}_{\mathbf{x} \sim \mathcal{D}} \langle \mathbf{P}, \mathbf{x}\mathbf{x}^\top \rangle - \frac{\lambda}{2} \|\mathbf{P}\|_F^2 \\ & \text{subject to} && \mathbf{P} \in \mathcal{C} \end{aligned} \quad (3)$$

3. Prior Work Guarantees

- Let \mathbf{P}^* be the optimal solution to Problem 1, $\mathbf{C} = \mathbb{E}[\mathbf{x}\mathbf{x}^\top]$, and $\Delta(\mathbf{C}) = \lambda_k(\mathbf{C}) - \lambda_{k+1}(\mathbf{C})$.
- Projected Stochastic Gradient Descent on Problem 2 (MSG) is:

$$\mathbf{P}_{t+1} = \Pi_{\mathcal{C}}(\mathbf{P}_t + \eta_t \mathbf{x}_t \mathbf{x}_t^\top)$$
- Projected Stochastic Gradient Descent on Problem 3 (RMSG) is:

$$\mathbf{P}_{t+1} = \Pi_{\mathcal{C}}((1 - \lambda \eta_t) \mathbf{P}_t + \eta_t \mathbf{x}_t \mathbf{x}_t^\top)$$
- MSG and RMSG guarantee that $\mathbb{E} \langle \mathbf{P}^* - \mathbf{P}_t, \mathbf{C} \rangle$ is in $\tilde{O}(1/\sqrt{t})$ and $\tilde{O}(1/(\Delta(\mathbf{C})t))$ [2].
- RMSG is statistically optimal, however, MSG and RMSG can take up to $\Omega(d^3)$ computations per iteration!

4. Meta-algorithm and Intuition

Input: Stream of data $\{\mathbf{x}_{t_i}\}$, parameters $\Delta(\mathbf{C})$, probability of failure δ , number of components k

Output: \mathbf{P}_T

- Initialize \mathbf{P}_1 from a warm start
- for** $t = 1, \dots, T-1$ **do**
- $\eta_t = \Theta\left(\frac{1}{\Delta(\mathbf{C})t}\right)$
- Form unbiased mini-batched estimator of gradient: $\mathbf{C}_t = \frac{1}{n} \sum_{l=1}^n \mathbf{x}_{t_l} \mathbf{x}_{t_l}^\top$
- $\mathbf{P}_{t+1/2} \leftarrow (1 - \frac{\Delta(\mathbf{C})}{2} \eta_t) \mathbf{P}_t + \eta_t \mathbf{C}_t$
- $\mathbf{P}_{t+1} = \Pi(\mathbf{P}_{t+1/2})$ (gradient descent upgrade) or equivalently
- $\mathbf{U}_{t+1} = \text{Top-k}\left([\sqrt{1 - \Delta(\mathbf{C})\eta_t/2} \mathbf{U}_t, \sqrt{\eta_t} \mathbf{X}_t]\right)$
- $\mathbf{P}_{t+1} = \mathbf{U}_{t+1} \mathbf{U}_{t+1}^\top$ (efficient update)
- end for**

Lemma 1 Let \mathbf{P}_t be rank k and suppose $\|\mathbf{C} - \mathbf{C}_t\| \leq \epsilon$. Then a sufficient condition for \mathbf{P}_{t+1} to be rank k is

$$\langle \mathbf{P}_t, \mathbf{C} \rangle \geq \langle \mathbf{P}^*, \mathbf{C} \rangle - \frac{\Delta(\mathbf{C})}{4} + \epsilon(k+1). \quad (4)$$

Intuition behind algorithm: Lemma 1 gives a sufficient conditions for \mathbf{P}_{t+1} to be a projection matrix, given that \mathbf{P}_t is a projection matrix. The sufficient condition translates to:

- \mathbf{P}_t is close enough to \mathbf{P}^* in objective and hence the warm start of the algorithm
- \mathbf{C}_t is close enough to \mathbf{C} with high probability which results in the mini-batched stochastic gradients

5. Main results

High Probability Guarantees

The following holds for Algorithm **MB-MSG**: with probability at least $1 - \delta$, for all $t \leq T$

$$\langle \mathbf{P}^* - \mathbf{P}_t, \mathbf{C} \rangle \leq O\left(\frac{k^4 \log(1/\delta) (\log(T))^2}{\sqrt{t + \frac{1}{\gamma}}}\right),$$

where $\gamma = O\left(\frac{\Delta(\mathbf{C})^2}{(k \log(1/\delta))^2}\right)$. Further, it holds that \mathbf{P}_t is a rank- k projection matrix.

The following holds for Algorithm **MB-RMSG**: with probability at least $1 - \delta$, for all $t \leq T$

$$\langle \mathbf{P}^* - \mathbf{P}_t, \mathbf{C} \rangle \leq \frac{32 \log(3e/\delta)}{\Delta(\mathbf{C})^2 \left(t + \frac{1}{\gamma} - 1\right)},$$

where $\gamma = \frac{\Delta(\mathbf{C})^3}{128 \log(1/\delta)}$. Further, it holds that \mathbf{P}_t is a rank- k projection matrix.

Guarantees in Expectation

Let \mathcal{A} be the event that for all $t \in [T]$ it holds that $\|\mathbf{C}_t - \mathbf{C}\| \leq \frac{\Delta(\mathbf{C})}{8(k+1)}$ and \mathbf{P}_t is a rank- k projection matrix. Then Algorithm **MB-RMSG** guarantees that \mathcal{A} occurs with probability at least $1 - \delta$ and that

$$\mathbb{E}[\langle \mathbf{P}^* - \mathbf{P}_T, \mathbf{C} \rangle | \mathcal{A}] \leq \tilde{O}\left(\frac{\Delta(\mathbf{C})}{T} + \min(\Delta(\mathbf{C}) \times d, 1) \frac{1}{kT}\right).$$

Comparison to Oja's Algorithm

(Informal [1]) The following holds for Algorithm **Oja's Algorithm**: with probability at least $1 - \delta$, for all $t \leq T$

$$\langle \mathbf{P}^* - \mathbf{P}_t, \mathbf{C} \rangle \leq \tilde{O}\left(\frac{1}{t \Delta(\mathbf{C})^2}\right)$$

- Total computational complexity for ϵ -suboptimality for **MB-RMSG** is $\tilde{O}(dk^2/(\epsilon \Delta(\mathbf{C})^2) \times \min(d \Delta(\mathbf{C}), 1))$
- Total computational complexity for ϵ -suboptimality for **Oja's Algorithm** is $\tilde{O}(dk/(\epsilon \Delta(\mathbf{C})^2))$

7. REFERENCES

- [1] Allen-Zhu, Z. and Li, Y. First efficient convergence for streaming k-pca: a global, gap-free, and near-optimal rate. *FOCS*, 2017.
- [2] Mianjy, P. and Arora, R. Stochastic PCA with ℓ_2 and ℓ_1 Regularization. *ICML*, 2018.

6. Experimental Results

