

# Efficient Convex Relaxations for Streaming PCA

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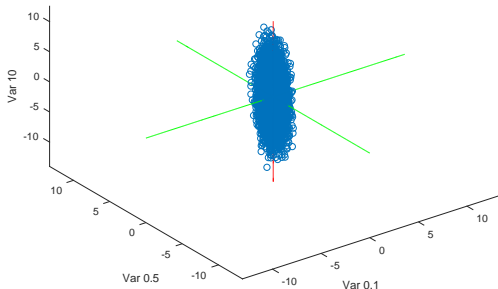
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# Outline

- 1 Introduction to the empirical PCA problem
- 2 The stochastic optimization perspective
- 3 Convex relaxations to PCA
- 4 Controlling the rank of iterates
- 5 Two convex optimization algorithms
- 6 Some insights into the proofs
- 7 Empirical results

# PCA as a geometric problem

- Given a data matrix  $X \in \mathbb{R}^{d \times T}$
- Output a  $U \in \mathbb{R}^{d \times k}$ ,  $U^T U = I_k$  to minimize reconstruction error:  
$$\frac{1}{T} \|X - UU^T X\|_F^2$$



# PCA as a geometric problem

## Optimization problem

$$\begin{aligned} \underset{U \in \mathbb{R}^{d \times k}}{\text{minimize:}} \quad & \frac{1}{T} \|X - UU^\top X\|_F^2 \\ \text{subject to:} \quad & U^\top U = I_k \end{aligned}$$

# PCA as a geometric problem

## Optimization problem

$$\begin{aligned} \underset{U \in \mathbb{R}^{d \times k}}{\text{minimize:}} \quad & \frac{1}{T} \|X - UU^T X\|_F^2 \\ \text{subject to:} \quad & U^T U = I_k \end{aligned}$$

Quickly rewriting the objective:

$$\begin{aligned} \frac{1}{T} \|X - UU^T X\|_F^2 &= \frac{1}{T} \{ \|X\|_F^2 + \|UU^T X\|_F^2 - 2\text{Tr}(X^T UU^T X) \} \\ &= \frac{1}{T} \{ \|X\|_F^2 + \text{Tr}((UU^T X)^T UU^T X) - 2\text{Tr}(X^T UU^T X) \} \\ &= \frac{1}{T} \{ \|X\|_F^2 + \text{Tr}(X^T UU^T UU^T X) - 2\text{Tr}(X^T UU^T X) \} \\ &= \frac{1}{T} \{ \|X\|_F^2 - \text{Tr}(U^T X X^T U) \} \end{aligned}$$

## Equivalent optimization problem

$$\begin{aligned} & \underset{U \in \mathbb{R}^{d \times k}}{\text{maximize}} && \frac{1}{T} \text{Tr} \left( U^\top X X^\top U \right) \\ & \text{subject to} && U^\top U = I_k \end{aligned}$$

# PCA as variance maximization

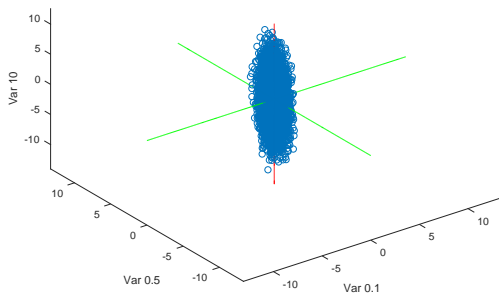
## Equivalent optimization problem

$$\begin{aligned} & \underset{U \in \mathbb{R}^{d \times k}}{\text{maximize}} && \frac{1}{T} \text{Tr} \left( U^\top X X^\top U \right) \\ & \text{subject to} && U^\top U = I_k \end{aligned}$$

## Solution to optimization problem

Optimal solution is given by a set of  $k$  eigenvectors associated with the top  $k$  eigenvalues of  $\frac{1}{T} X X^\top$

# The stochastic optimization perspective



Each  $\mathbf{x}_t \sim \mathcal{N}\left(\mathbf{0}, \begin{pmatrix} 0.1 & 0 & 0 \\ 0 & 0.5 & 0 \\ 0 & 0 & 10 \end{pmatrix}\right)$  independently.



# The stochastic optimization perspective

- Assume each column  $\mathbf{x}_t$  of  $\mathbf{X}$  is i.i.d as some probability law  $\mathcal{D}$ .
- Empirical problem is just a proxy for the stochastic optimization problem:

## Stochastic optimization problem

$$\begin{aligned} & \underset{\mathbf{U} \in \mathbb{R}^{d \times k}}{\text{minimize:}} && \mathbb{E}_{\mathbf{x} \sim \mathcal{D}} \|\mathbf{x} - \mathbf{U}\mathbf{U}^\top \mathbf{x}\|_F^2 \\ & \text{subject to:} && \mathbf{U}^\top \mathbf{U} = \mathbf{I}_k \end{aligned}$$

# The stochastic optimization perspective

## Minimizing reconstruction

$$\begin{aligned} & \underset{U \in \mathbb{R}^{d \times k}}{\text{minimize:}} && \mathbb{E}_{x \sim \mathcal{D}} \|x - UU^\top x\|_F^2 \\ & \text{subject to:} && U^\top U = I_k \end{aligned}$$



## Variance maximization

$$\begin{aligned} & \underset{U \in \mathbb{R}^{d \times k}}{\text{maximize}} && \mathbb{E}_{x \sim \mathcal{D}} \text{Tr} \left( U^\top x x^\top U \right) \\ & \text{subject to} && U^\top U = I_k \end{aligned} \tag{1}$$

# Oja's algorithm

- One possible way to solve Problem 1 is to use Stochastic Gradient Descent (SGD)
- Gradient for objective in Problem 1 is  $2\mathbb{E}_{\mathbf{x} \sim \mathcal{D}}[\mathbf{x}\mathbf{x}^\top]\mathbf{U}$
- Since we do not have direct access to distribution  $\mathcal{D}$ , we use an unbiased estimator of the gradient based on a sample  $\mathbf{x}_t \sim \mathcal{D}$  given by  $\mathbf{x}_t\mathbf{x}_t^\top\mathbf{U}$

# Oja's algorithm

SGD on Problem 1 is also known as Stochastic Power Method (Oja's algorithm) Allen-Zhu and Li [2017]

$$U_t \leftarrow (I + \eta_t x_t x_t^\top) U_{t-1}, U_t = \text{Orth}(U_t); \quad (2)$$

## Convergence guarantee (informal)

After  $T$  iterations of Oja's algorithm w.p.  $1 - \delta$  it holds that

$$\|U_*^\top U_T\|_F^2 \leq k - \tilde{O}\left(\frac{1}{\Delta(C)^2 T}\right),$$

where  $C = \mathbb{E}_{x \sim \mathcal{D}}[xx^\top]$  and  $\Delta(C)$  is the eigengap at the  $k$ -th eigenvalue.

$$\Delta(C) := \lambda_k(C) - \lambda_{k+1}(C)$$

# Convex relaxations to PCA

## Maximizing variance formulation

$$\begin{aligned} & \underset{U \in \mathbb{R}^{d \times k}}{\text{maximize}} && \mathbb{E}_{x \sim \mathcal{D}} \text{Tr} \left( U U^\top x x^\top \right) \\ & \text{subject to} && U^\top U = I_k \end{aligned}$$

The above is equivalent to:

## Maximizing variance formulation

$$\begin{aligned} & \underset{P \in \mathbb{R}^{d \times d}}{\text{maximize}} && \mathbb{E}_{x \sim \mathcal{D}} \text{Tr} \left( P x x^\top \right) \\ & \text{subject to} && P^2 = P, P^\top = P, \text{rank}(P) = k \end{aligned}$$

# Convex relaxations to PCA (continued)

The convex hull of  $\{P \in \mathbb{R}^{d \times d} : P^2 = P, P^\top = P, \text{rank}(P) = k\}$  is  $\{P \in \mathbb{R}^{d \times d} : \text{Tr}(P) \leq k, 0 \preceq P \preceq I, P^\top = P\}$

Convex relaxation [Arora et al., 2013]

$$\begin{aligned} & \underset{P \in \mathbb{R}^{d \times d}}{\text{maximize}} && \text{Tr}(PC) \\ & \text{subject to} && \text{Tr}(P) \leq k, 0 \preceq P \preceq I, P^\top = P \end{aligned} \quad (3)$$

Convex relaxation with regularization [Mianjy and Arora, 2018]

$$\begin{aligned} & \underset{P \in \mathbb{R}^{d \times d}}{\text{maximize}} && \text{Tr}(PC) - \frac{\lambda}{2} \|P\|_F^2 \\ & \text{subject to} && \text{Tr}(P) \leq k, 0 \preceq P \preceq I, P^\top = P \end{aligned} \quad (4)$$

# Convex relaxations to PCA (continued)

## Projected SGD for Problem 3 (MSG)

$$P_t \leftarrow \mathcal{P} \left( P_{t-1} + \eta_t \mathbf{x}_t \mathbf{x}_t^\top \right)$$

## Projected SGD for Problem 4 (RMSG)

$$P_t \leftarrow \mathcal{P} \left( (1 - \lambda \eta_t) P_{t-1} + \eta_t \mathbf{x}_t \mathbf{x}_t^\top \right)$$

In the above  $\mathcal{P}(\cdot)$  is the projection onto the convex set of constraints  $\{P : \text{Tr}(P) \leq k, 0 \preceq P \preceq I, P^\top = P\}$ .

# Convex relaxations to PCA (continued)

Running projected SGD on the above problems, comes with the following guarantees:

## Convergence guarantee for MSG (informal)

After  $T$  iterations of MSG, it holds that

$$\mathbb{E}[\langle P_* - P_T, C \rangle] \leq \tilde{O}\left(\frac{1}{\sqrt{T}}\right).$$

## Convergence guarantee for RMSG (informal)

After  $T$  iterations of RMSG, it holds that

$$\mathbb{E}[\langle P_* - P_T, C \rangle] \leq \tilde{O}\left(\frac{1}{\Delta(C)^2 T}\right).$$



# Angle between subspaces and suboptimality in objective

- Oja's guarantee is of the form  $k - \|U_*^\top U_T\|_F^2 \leq \epsilon$
- MSG and RMSG guarantees are of the form  $\langle U_* U_*^\top - U_T U_T^\top, C \rangle \leq \epsilon$
- We have [Mianjy and Arora, 2018]

$$\langle U_* U_*^\top - U_T U_T^\top, C \rangle \leq \lambda_1(C)(k - \|U_*^\top U_T\|_F^2)$$

- No known relation in opposite direction

# Per iteration complexity

- The computational complexity of Oja's algorithm per iteration is  $O(dk^2)$  (can be reduced to  $O(dk)$  if we do not call the (*Orth*) procedure)
- The computational complexity of MSG and RMSG per iteration is  $O(d\text{rank}(P_t)^2)$
- Worst case for MSG and RMSG is  $O(d^3)$

# Per iteration complexity in practice

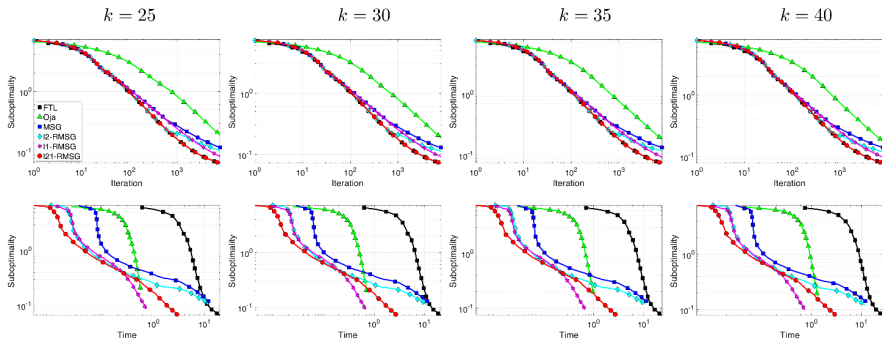


Figure 3: Experiment on MNIST by Mianjy and Arora [2018]

# Controlling the rank of $P_t$

- If distribution is well-behaved rank of  $P_t$  stays in  $O(k)$
- We get a similar computational complexity to Oja's algorithm but experiments suggest MSG and RMSG have better convergence properties
- Can we formalize the above statements through theoretical results?

# Controlling the rank of $P_t$

- If distribution is well-behaved rank of  $P_t$  stays in  $O(k)$
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- Can we formalize the above statements through theoretical results?

Yes, if we make slight tweaks to MSG and RMSG!

**Input:** Stream of data  $\{x_{t_n}\}$ , parameters  $\Delta(C)$ , probability of failure  $\delta$ , number of components  $k$

**Output:**  $P_T$

```

1:  $n = \log(d/\delta) \frac{128k \log(1/\delta)}{\Delta(C)^5}$ 
2:  $P_1 = \text{Top-k}(\frac{1}{n} \sum_{l=1}^n x_{0_l} x_{0_l}^\top)$ 
3:                                     %%  $\{x_{0_l}\}_{l=1}^n$  is the warm-start mini-batch

4:  $n = \log\left(\frac{Td}{\delta}\right) \frac{8(k+1)}{\Delta(C)^2}$ 
5: for  $t = 1, \dots, T-1$  do
6:    $\eta_t = \frac{1}{\frac{\Delta(C)}{2} \left( t + \frac{128 \log\left(\frac{1}{\delta}\right)}{\Delta(C)^3} \right)}$ 
7:    $C_t \leftarrow \frac{1}{n} \sum_{l=1}^n x_{t_l} x_{t_l}^\top$ 
8:                                     %%  $\{x_{t_l}\}_{l=1}^n$  is the mini-batch for the  $t^{\text{th}}$  epoch

9:    $P_{t+1/2} \leftarrow \left(1 - \frac{\Delta(C)}{2} \eta_t\right) P_t + \eta_t C_t$ 
10:   $P_{t+1} = \mathcal{P}(P_{t+1/2})$ 
11: end for
```

## Theorem

*There exists an algorithm, solving Problem 4, which after  $T$  iterations, with probability at least  $1 - 3e\delta$ , returns a sequence of iterates  $\{P_t\}_{t=1}^T$ , such that for all  $t \leq T$*

$$\langle P^* - P_t, C \rangle \leq \frac{32 \log(1/\delta)}{\Delta(C)^2 \left(t + \frac{1}{\gamma} - 1\right)},$$

*where  $\gamma = \frac{\Delta(C)^3}{128 \log(1/\delta)}$ . Further, for all  $t \leq T$  it holds that  $P_t$  is a rank- $k$  projection matrix and the per-iteration computational complexity of the algorithm is bounded by  $\tilde{O}\left(\frac{dk^2}{\Delta(C)^2} + dk^2\right)$ .*

# Formal guarantee 1 (continued)

- Total computational complexity for  $\epsilon$ -suboptimality of MB-RMSG is  $\tilde{O}(\frac{dk^2}{\Delta(C)^4\epsilon})$
- Total computational complexity for Oja's algorithm to reach  $\epsilon$ -suboptimality is  $\tilde{O}(\frac{dk}{\Delta(C)^2\epsilon})$

We can do better in expectation!



### Theorem

Let  $\mathcal{A}$  be the event that for all  $t \in [T]$  it holds that  $\|C_t - C\| \leq \frac{\Delta(C)}{8(k+1)}$  and  $P_t$  is a rank- $k$  projection matrix. Then Algorithm 22 guarantees that  $\mathcal{A}$  occurs with probability at least  $1 - \delta$  and that

$$\mathbb{E}[\langle P^* - P_T, C \rangle | \mathcal{A}] \leq \tilde{O} \left( \frac{\Delta(C)}{T} + \min(\Delta(C) \times d, 1) \frac{1}{kT} \right).$$

Above theorem implies that the total computational complexity for achieving  $\epsilon$ -suboptimality is  $\tilde{O} \left( \frac{dk^2}{\epsilon \Delta(C)^2} \times \min(d\Delta(C), 1) \right)$ , which is only a factor of  $k$  away from Oja's algorithm whenever the gap is large, and actually improves by a factor of  $1/\Delta(C)$  over Oja's in the case when  $\Delta(C) \in o(1/kd)$ !

## Lemma

Let  $P_t$  be rank  $k$  and suppose  $\|C - C_t\| \leq \beta$ . Then, a sufficient condition for  $P_{t+1}$  to be rank  $k$  is

$$\langle P_t, C \rangle \geq \langle P^*, C \rangle - \frac{\Delta(C)}{2} + \frac{\lambda}{2} + \beta(k+1). \quad (5)$$

- Projection works by shifting all eigenvalues of  $P_t + \eta_t C_t$  and then clipping them between 0 and 1
- Let  $\lambda_k(P_t + \eta_t C_t) = 1 + \lambda_k$  and  $\lambda_{k+1}(P_t + \eta_t C_t) = \lambda_{k+1}$  for some  $\lambda_k$  and  $\lambda_{k+1}$
- If  $\lambda_k > \lambda_{k+1}$ , then the shift from the projection is larger than  $\lambda_{k+1}$  and thus projection will clip  $\lambda_{k+1}$  to 0.
- We can guarantee that this happens with high probability if  $P_t$  is close enough to  $P^*$  and  $C_t$  is close enough to  $C$ .

# Rest of analysis

- The rest of the analysis requires last iterate SGD guarantees with high probability
- For results in expectation we need to adapt the analysis for smooth SGD
- Both of these are non-trivial to do!

# But does it work?

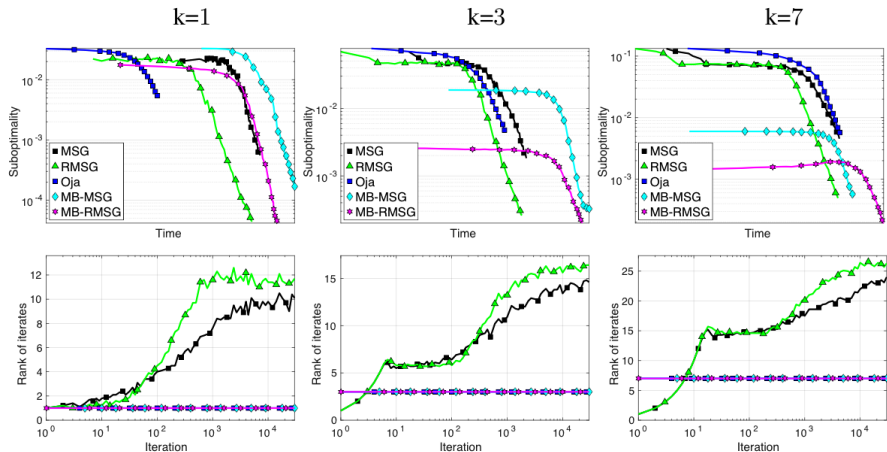


Figure 4: Experiments on MNIST

# Open problems

- We gave an algorithm based on the convex relaxation to PCA which achieves (almost) optimal rates
- But goal was to study MSG and RMSG directly
- We need better tools to control rank of MSG and RMSG iterates

- There are many other works on Streaming PCA, mainly focused on studying Oja's algorithm [De Sa et al., 2014, Hardt and Price, 2014, Balcan et al., 2016, Jain et al., 2016, Shamir, 2016a,b, Allen-Zhu and Li, 2017, Li et al., 2018]
- Other important work in the Online Learning setting is by Warmuth and Kuzmin [2008], Grabowska and Kotłowski [2018], Garber [2018]

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