

## Unconstrained Optimization

### Taylor's Theorem

We can use this theorem to approximate functions when  $f$  is of class  $C^2$ , (**page 14**)

$$f(\mathbf{x} + \mathbf{p}) = f(\mathbf{x}) + \nabla f(\mathbf{x})^T \mathbf{p} + \frac{1}{2} \mathbf{p}^T \nabla^2 f(\mathbf{x}) \mathbf{p} + O(\|\mathbf{p}\|^3)$$

### Q/Quotient Convergence

A sequence  $\{x_k\}$  which converges to  $x^*$ , is said to converge (**page 619**)

- Linearly with rate of convergence  $c$ . if  $r = 1$  and  $c \in (0, 1)$ .
- Superlinearly if  $r = 1$  and  $c = 0$ .

- Quadratically if  $r = 2$  and  $c$  is finite.

$$c = \lim_{k \rightarrow \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|^r}$$

### Types of minima

A point  $x^*$  is a **global** minima of  $f$  if,

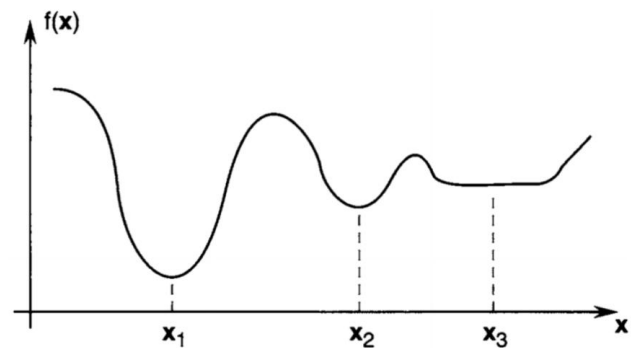
$$f(x^*) \leq f(x) \quad \forall x \in \mathbb{R}^n$$

A point  $x^*$  is a **strict/strong local** minima of  $f$  if  $\exists \mathbb{N}$  such that,

$$f(x^*) < f(x) \quad \forall x \in \mathbb{N}$$

A point  $x^*$  is a **weak local** minima of  $f$  if  $\exists \mathbb{N}$  such that,

$$f(x^*) \leq f(x) \quad \forall x \in \mathbb{N}$$



$x_1$ : strict global minimizer;  $x_2$ : strict local minimizer;  $x_3$ : local minimizer

In both cases,  $\mathbb{N}$  is a neighbourhood of  $x^*$ , that is any open set containing  $x^*$ . (**page 12**)

### Optimality conditions

If  $f$  is continuously differentiable in an open neighbourhood of  $x^*$  and if  $x^*$  is also a weak local minima then the **first order necessary** condition is just  $\nabla f(\mathbf{x}^*) = \mathbf{0}$ .

The **second order necessary** conditions are  $\nabla f(\mathbf{x}^*) = \mathbf{0}$  and  $\nabla^2 f(x^*)$  is positive **semi** definite.

The **second order sufficient** conditions for  $x^*$  to be a strong local minima is  $\nabla f(\mathbf{x}^*) = \mathbf{0}$  and  $\nabla^2 f(x^*)$  is positive definite. (page 15)

## Line search and Trust region Overview

These are the two broad strategies used to move from  $x_k$  to the next iterate  $x_{k+1}$ ,

1. Line Search : Choose a direction  $p_k$ , solve  $\min_{\alpha > 0} f(x_k + \alpha p_k)$  to get the step length  $\alpha$ .  
Now  $x_{k+1} = x_k + \alpha p_k$
2. Trust Region : Choose a trust radius  $\Delta_k$ . Approximate  $f$  around  $x_k$  as a quadratic function,  
 $m_k(x_k + p) = f_k + p^T \nabla f_k + \frac{1}{2} p^T B_k p$ . Now we find  $x_{k+1} = x_k + p$  by solving  $\min_p m_k(x_k + p)$  for  $p$  where  $|x_{k+1} - x_k| < \Delta_k$ . If this new iterate is not satisfactory we shrink  $\Delta_k$  and repeat the steps.

$\nabla f_k$  is the gradient of  $f$  evaluated at  $x_k$ .  $\nabla^2 f_k$  is the hessian of  $f$  evaluated at  $x_k$ .  $B_k$  is an approximation of  $\nabla^2 f_k$ . (page 19)

## Line search directions

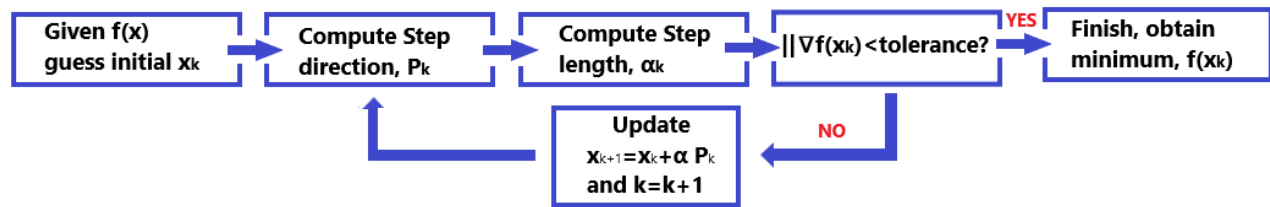
1. Steepest Direction :  $p_k = \frac{-\nabla f_k}{\|\nabla f_k\|}$ , we get the most rapid decrease in this direction from  $f(x_k)$ .
2. Newton Direction :  $p_k = -(\nabla^2 f_k)^{-1} \nabla f_k$ , the exact solution to the quadratic approximation of  $f$ . It is a descent direction only if the hessian is positive definite.
3. Quasi Newton Direction :  $p_k = B_k^{-1} \nabla f_k$ , where we avoid the costly hessian computation by using either Symmetric-rank-one (SR1) or BFGS iterative approximation. We impose symmetry and secant equation conditions,  $B_k^T = B_k$  and  $B_{k+1}(x_{k+1} - x_k) = \nabla f_{k+1} - \nabla f_k$ .
4. Non Linear Conjugate Gradient directions :  $p_k = -\nabla f_k + \beta_k p_{k-1}$ , here  $\beta_k$  is a scalar that ensures  $p_k$  and  $p_{k-1}$  is conjugate. (page 21)

## Wolfe Conditions

The armijo condition (to ensure sufficient decrease) and the curvature condition (to ensure the steps are not too small) are together known as the Wolfe conditions. (page 33)

1. Armijo condition :  $f(x_k + \alpha p_k) \leq f(x_k) + c_1 \alpha \nabla f_k^T p_k$  : this enforces sufficient decrease in  $f_k$  every iteration.  $c_1 \in (0, 1)$ , usually  $10^{-4}$
2. Curvature condition :  $\nabla f(x_k + \alpha_k p_k)^T p_k \geq c_2 \nabla f_k^T p_k$  : here  $c_2 \in (c_1, 1)$  (usually 0.9) and this ensures the new slope  $\phi'(\alpha)$  is less negative than  $c_2 \times \phi'(0)$ . Where  $\phi(\alpha) = f(x_k + \alpha p_k)$

## Algorithm flow chart of line search methods



## Convergence of line search methods

To analyze convergence, we use the angle between our descent direction and the steepest descent direction.  
(Page-37)

$$\cos \theta_k = \frac{-\nabla f_k^T p_k}{\|\nabla f_k\| \|p_k\|}$$

## Zoutendijk Theorem:

### (Theorem-3.2)

Consider any iteration of the form  $x_{k+1} = x_k + \alpha_k p_k$ ,

- where  $p_k$  is a descent direction
- $\alpha_k$  satisfies the Wolfe conditions
- $f$  is bounded below in  $\mathbb{R}^n$
- $f$  is continuously differentiable in an open set  $N$  containing the level set  $L = \{x : f(x) \leq f(x_0)\}$ ,
- Assume also that the gradient  $\nabla f$  is Lipschitz continuous on  $N$ .

Then,

$$\sum_{k \geq 0} \cos^2 \theta_k \|\nabla f_k\|^2 < \infty \quad (1)$$

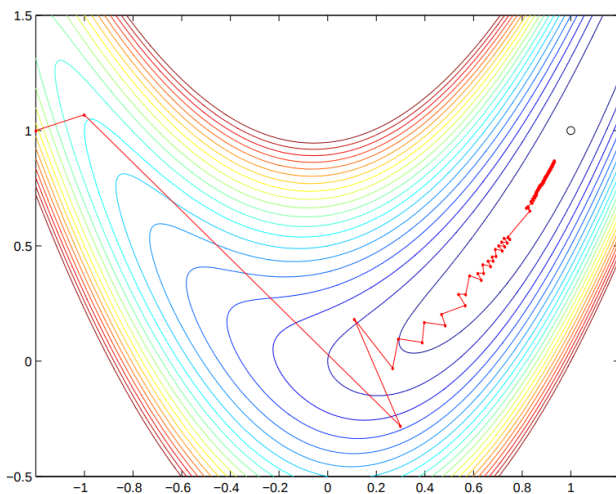
(1) is called **Zoutendijk condition** which implies that

$$\cos^2 \theta_k \|\nabla f_k\|^2 \rightarrow 0 \quad (2)$$

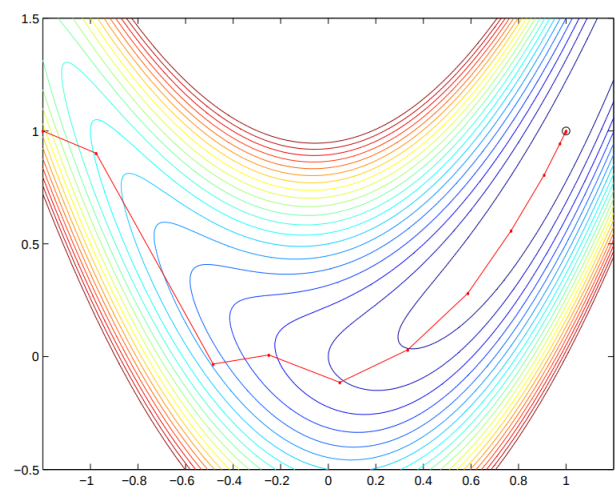
when  $-90^\circ < \theta_k < 90^\circ$  for all  $k$ . It follows immediately from (2) that

$$\lim_{k \rightarrow \infty} \|\nabla f\| = 0 \quad (3)$$

The algorithms which satisfy (3) are called **globally convergent**. Zoutendijk condition shows that the steepest descent method is globally convergent. For other algorithms it describes how far  $p_k$  can deviate from the steepest descent direction and still give rise to a globally convergent iteration.



(a) iterates generated by the generic line search **steepest-descent method**



(b) iterates generated by the Generic Line search **Newton method**

Contours for the objective function  $f(x, y) = 10(y - x^2)^2 + (x - 1)^2$  (**Rosenbrock function**)

## Rate of Convergence

### Convergence of Steepest Descent

- $p_k = \frac{-\nabla f_k}{\|\nabla f_k\|}$
- Globally convergent (converges to a local minimiser from any starting point  $x_0$ ).
- many other methods resort to steepest descent in bad cases
- not scale invariant (changing the inner product on  $\mathbb{R}^n$  changes the notion of gradient!).
- convergence is usually very (very!) slow (linear)
- numerically often not convergent at all

### Convergence of Newton's method

- $p_k = -\nabla^2 f_k^{-1} \nabla f_k$
- convergence is often faster than steepest descent
- may be viewed as “scaled” steepest descent
- If Hessian matrix  $\nabla^2 f_k^{-1}$  is not Positive definite then  $p_k$  is not a descent direction. Two ways for obtaining globally convergent iteration:

- **line search approach**, in which the Hessian is modified, to make it positive definite.
- **trust region approach**, in which Hessian is used to form a quadratic model that is minimized in a ball.

## Step length Selection algorithm

The line search is done in two stages: A **bracketing phase** finds an interval containing desirable step lengths, and a **bisection or interpolation phase** computes a good step length within this interval.

$$\phi(\alpha) = f(x_k + \alpha p_k)$$

If  $f$  is a convex quadratic,  $f(x) = \frac{1}{2}x^T Qx - b^T x$ , its one-dimensional minimizer along the ray  $x_k + \alpha p_k$  can be computed analytically and is given by

$$\alpha_k = \frac{-\nabla f_k^T p_k}{p_k^T Q p_k}$$

For general nonlinear functions, it is necessary to use an iterative procedure. The line search procedure deserves particular attention because it has a major impact on the robustness and efficiency of all nonlinear optimization methods.