

Conjugate Gradient - 1

1 Co-ordinate Descent Method

This method attempts to minimize the objective function, by minimizing the function along one (axis) direction at a time. For a function of the form,

$$f(X) = 4X_1^2 + X_2^2$$

this method tries to minimize the function along X_1 first and then along X_2 (i.e. along the axes), eventually reaching the minima (0,0) in exactly two iterations. However for a function of the form,

$$g(X) = 4X_1^2 + X_2^2 - 2X_1X_2$$

the same method takes multiple iterations to converge to the minimum.

- The reason for multiple iterations for $g(X)$ to converge can be attributed to the fact that the elliptically shaped contours of $g(X)$ do not have their axes lined up along co-ordinate axes, X_1 and X_2 .
- Whenever the Hessian(H) is a diagonal matrix (as in case of $f(X)$, not for $g(X)$), co-ordinate descent method converges to the solution in exactly N iterations (N being the dimension of X).

It is also necessary to note that this method would fail for non-smooth functions as it might get stuck at a non-stationary point.

However, the method attracts substantial interest due its inherent advantages of high computational efficiency, scalability, quicker convergence and ease of implementation.

2 Linear Conjugate Gradient Method

The Linear Conjugate Gradient Method is an iterative method to solve the system of equations $Ax = b$, (A is symmetric and positive definite) which can be equivalently formulated as a minimization problem of the quadratic form:

$$\phi(x) = \frac{1}{2}x^T Ax - b^T x$$

Hence, the method can well be extended to solve optimization problems.

3 Conjugate Gradient Method

Let us consider a minimization problem of the type

$$\phi(x) = \frac{1}{2}x^T Hx + c^T x, \quad (1)$$

where H is assumed to be symmetric and positive definite. Let d^0, d^1, \dots, d^{n-1} be n such linearly independent vectors (directions) and x be the iterative solution for the system. We have,

$$x = x^0 + \sum_{i=0}^{n-1} \alpha^i d^i, \quad (2)$$

Upon back substitution to the original problem, we get

$$\phi(\alpha) = \frac{1}{2}\alpha^T D^T H D \alpha + c^T D \alpha, \quad (3)$$

where $D = [d^0 d^1 \dots d^{n-1}]$ and $\alpha = [\alpha^0 \alpha^1 \dots \alpha^{n-1}]$.

Let $H' = D^T H D$, then H' is a matrix with elements as $d^{iT} H d^j \forall i, j$.

- For H' to be a diagonal matrix, $d^{iT} H d^j = 0 \forall i \neq j$. This condition ensures that $\phi(\alpha)$ is minimized.

This reduces to condition of H -orthogonality, i.e. the vectors (d^i) must be orthogonal to each other. Such set of vectors (d^i) , which satisfy the relation $d^{iT} H d^j = 0 \forall i \neq j$, are referred to as H -conjugate vectors.

Proof of Existence: Let H be a symmetric matrix with n orthogonal eigen-vectors, if v_1, v_2 be two such distinct eigen vectors, then $v_1^T v_2 = 0$ and $H v_1 = \lambda_1 v_1$, a multiplication by v_2^T results in $v_2^T H v_1 = 0$. This suffices the proof of existence of H -conjugate vectors, for symmetric Hessian.

Gradient, Error and Residuals

- Residual(r_k) is defined as $-\nabla \phi(x) = -c - Hx$ (from 1).
- Error(e_k) = $x_k - x^*$, x^* is the true solution.

Let x_k be an iterative solution to a system, with α_k and d_k being the step length and direction at the k^{th} iterate, then $x_{k+1} = x_k + \alpha_k d_k$. Hence, $f(x_{k+1}) = f(x_k + \alpha_k d_k)$. Differentiating w.r.t α_k , we get $\nabla f(x_{k+1})^T d_k = 0$ (since, $f(x_{k+1})$ is some constant), equivalently,

$$g_{k+1}^T d_k = 0,$$

where $g_{k+1} = \nabla f(x_{k+1})$. Also, $g_k = H e_k$. Hence, $(H e_{k+1})^T d_k = 0$, which is equivalent to

$$e_{k+1}^T H^T d_k = 0, \quad (4)$$

which means error in $(k+1)^{th}$ iteration is H -conjugate to $d^k \forall k \in [0, n-1]$.