

Lecture 7

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1 The Johnson-Lindenstrauss Lemma

The Johnson-Lindenstrauss lemma maps a metric space into a subset of Euclidean space of relatively low dimension, and it follows nicely from Bernstein-type inequalities. Suppose we have a collection $\mathcal{C} = \{\mathbf{x}_i\}_{i=1}^m \subseteq \mathbb{R}^n$ all distinct. Then, $(\mathcal{C}, \|\cdot\|_2)$ is a metric space. In order to compress signal, work in lower dimension, etc., we want to map this space into \mathbb{R}^k where $k \ll n$ so that the metric structure is not distorted too much.

The idea is to use a random $k \times n$ matrix for the dimension reduction. Denote this linear map by $\Phi : \mathbb{R}^n \rightarrow \mathbb{R}^k$. We would like

$$1 - \epsilon \leq \frac{\|\Phi(\mathbf{x}) - \Phi(\mathbf{y})\|_2}{\|\mathbf{x} - \mathbf{y}\|_2} \leq 1 + \epsilon \quad (1)$$

for all $\mathbf{x}, \mathbf{y} \in \mathcal{C}$. To summarize, we have parameters ϵ, k, n, m , and the question is what restriction on these parameters is necessary to find such a mapping Φ .

Lemma 1.1: Johnson-Lindenstrauss

There exists a universal constant $c > 0$ with the following property. Let $k \geq c\epsilon^{-2} \log(m)$ for $\epsilon \in (0, \frac{1}{2}]$. Then, there exists a linear mapping Φ satisfying condition (1). In particular, if \mathbf{A} is a $k \times n$ matrix with iid entries of zero mean, unit variance, and constant subgaussian moment, then $\Phi = \frac{1}{\sqrt{k}} \mathbf{A}$ works with high probability.

The proof idea is given by the corollaries from the last lecture. For every $\mathbf{x}, \mathbf{y} \in \mathcal{C}$, the vector

$$\mathbf{q} = \frac{\mathbf{A}(\mathbf{x} - \mathbf{y})}{\|\mathbf{x} - \mathbf{y}\|_2} \quad (2)$$

is isotropic, centered, and subgaussian with iid subgaussian components. This implies that

$$\Pr\left[\left|\|\mathbf{q}\|_2 - \sqrt{k}\right| \geq t\right] \leq 2\exp(-\tilde{c}t^2) \quad (3)$$

for all $t > 0$, where $\tilde{c} > 0$ may only depend on the subgaussian moment of the entries. If t is taken as an appropriate constant multiple of $\sqrt{\log(m)}$, then we get a quantity much less than m^{-2} on the right-hand side. So we can take a union bound over pairs of \mathbf{x}, \mathbf{y} to obtain the result.

An interesting example is that if \mathcal{C} is the standard basis in \mathbb{R}^n , applying Φ gives us m vectors

in $\mathbb{R}^{c \log(n)}$ such that $\|\Phi(\mathbf{x}_i)\| \in [0.99, 1.01]$ for all i and $|\langle \Phi(\mathbf{x}_i), \Phi(\mathbf{x}_j) \rangle| \leq 0.01$. So, we have just mapped exponentially many vectors into the space in such a way that they are almost unit and almost pairwise orthogonal.

2 Concentration of Sparse Linear Combinations

Theorem 2.1: Chernoff

Let b_1, b_2, \dots, b_n be iid $\text{Ber}(p)$ random variables. Then,

$$\Pr \left[\sum_{i=1}^n b_i - pn \geq tpn \right] \leq \left(\frac{\exp(t)}{(1+t)^{(1+t)}} \right)^{pn} \quad (4)$$

for $t > 0$. For the other tail,

$$\Pr \left[\sum_{i=1}^n b_i - pn \leq -tpn \right] \leq \left(\frac{\exp(-t)}{(1-t)^{(1-t)}} \right)^{pn} \quad (5)$$

for $t \in (0, 1)$.

Proof: Of the first inequality. Choose $\lambda > 0$ (we will optimize over λ later). Using MGF,

$$\Pr \left[\sum_{i=1}^n b_i - pn \geq tpn \right] = \Pr \left[\exp \left(\lambda \sum_{i=1}^n b_i - \lambda pn \right) \geq \exp(\lambda tpn) \right]. \quad (6)$$

Applying Markov's inequality, this is at most

$$\frac{\mathbb{E}[\exp(\lambda \sum_{i=1}^n b_i - \lambda pn)]}{\exp(\lambda tpn)} = \frac{(\mathbb{E}[\exp(\lambda b_1 - \lambda p)])^n}{\exp(\lambda tpn)}. \quad (7)$$

Since the variables are Bernoulli we can easily and directly compute the MGF. So the above is equal to

$$\frac{((1-p)\exp(-\lambda p) + p\exp(\lambda - \lambda p))^n}{\exp(\lambda tpn)} = \frac{\exp(-\lambda pn)}{\exp(\lambda tpn)} (1 - p + p\exp(\lambda))^n. \quad (8)$$

If we take $\lambda = \ln(1+t)$ and use $1+s \leq \exp(s)$, this simplifies to

$$\frac{(1+pt)^n}{\exp(tpn \ln(1+t) + pn \ln(1+t))} \leq \frac{\exp(pn \ln(1+t))}{(1+t)^{tpn+pn}} = \left(\frac{\exp(t)}{(1+t)^{(1+t)}} \right)^{pn}. \quad (9)$$

■

Corollary 2.1

For a universal constant $c > 0$,

$$\Pr \left[\left| \sum_{i=1}^n b_i - pn \right| \geq t \right] \leq 2 \exp \left(-\frac{ct^2}{pn} \right). \quad (10)$$

Note that this Chernoff-based bound is stronger than Hoeffding's inequality in the regime $p \rightarrow 0$. If we apply the version of Hoeffding for variables supported on bounded intervals, we would get

$$\Pr \left[\left| \sum_{i=1}^n b_i - pn \right| \geq t \right] \leq 2 \exp \left(-\frac{ct^2}{n} \right). \quad (11)$$

This is similar for the other version of Hoeffding and the general Bernstein inequality.

Proof: For every $t \in (0, pn]$ we have by Chernoff

$$\Pr \left[\left| \sum_{i=1}^n b_i - pn \right| \geq t \right] \leq \Pr \left[\sum_{i=1}^n b_i - pn \geq t \right] + \Pr \left[\sum_{i=1}^n b_i - pn \leq -t \right] \quad (12)$$

$$\leq \left(\frac{\exp\left(\frac{t}{pn}\right)}{\left(1 + \frac{t}{pn}\right)^{1 + \frac{t}{pn}}} \right)^{pn} + \left(\frac{\exp\left(-\frac{t}{pn}\right)}{\left(1 - \frac{t}{pn}\right)^{1 - \frac{t}{pn}}} \right)^{pn} \quad (13)$$

$$= \exp \left(t - pn \left(1 + \frac{t}{pn} \right) \ln \left(1 + \frac{t}{pn} \right) \right) + \exp \left(-t - pn \left(1 - \frac{t}{pn} \right) \ln \left(1 - \frac{t}{pn} \right) \right). \quad (14)$$

We now apply the fact from calculus that for all t such that $\frac{t}{pn} \in (0, 1]$, we have

$$\frac{t}{pn} - \left(1 + \frac{t}{pn} \right) \ln \left(1 + \frac{t}{pn} \right) \leq -\frac{1}{4} \frac{t^2}{p^2 n^2}. \quad (15)$$

This can be checked directly by taking derivatives. So, the first term in (14) is at most

$$\exp \left(-\frac{pn}{4} \frac{t^2}{p^2 n^2} \right) = \exp \left(-\frac{1}{4} \frac{t^2}{pn} \right). \quad (16)$$

Similar ideas work for the second term in (14), and the corollary follows. ■

Lemma 2.1: Sparse Bernstein Inequality

Let b_1, b_2, \dots, b_n be iid $\text{Ber}(p)$ random variables and $a_1, a_2, \dots, a_n \in \mathbb{R}^n$ such that $\|\mathbf{a}\|_2 = 1$. Then, for all $t > 0$ we have

$$\Pr \left[\sum_{i=1}^n a_i b_i - p \sum_{i=1}^n a_i \geq tp \right] \leq \exp \left(-cp \min \left(t^2, \|\mathbf{a}\|_\infty^{-1} t \right) \right). \quad (17)$$

and

$$\Pr \left[\sum_{i=1}^n a_i b_i - p \sum_{i=1}^n a_i \leq -tp \right] \leq \exp \left(-cp \min \left(t^2, \|\mathbf{a}\|_\infty^{-1} t \right) \right). \quad (18)$$

This statement can be viewed as a generalization of the last corollary (it is exactly the scenario with all $a_i = 1/\sqrt{n}$).

Proof: Of the first inequality. We use the MGF/Markov strategy:

$$\Pr \left[\sum_{i=1}^n a_i b_i - p \sum_{i=1}^n a_i \geq tp \right] = \Pr \left[\exp \left(\lambda \sum_{i=1}^n a_i b_i - \lambda p \sum_{i=1}^n a_i \right) \geq \exp(\lambda tp) \right] \quad (19)$$

$$\leq \frac{\mathbb{E}[\exp(\lambda \sum_{i=1}^n a_i b_i - \lambda p \sum_{i=1}^n a_i)]}{\exp(\lambda tp)} \quad (20)$$

$$= \frac{\prod_{i=1}^n \mathbb{E}[\exp(\lambda a_i b_i - \lambda p a_i)]}{\exp(\lambda tp)}. \quad (21)$$

We can compute the numerator explicitly to find

$$\frac{\prod_{i=1}^n [(1-p) \exp(-\lambda p a_i) + p \exp(\lambda a_i - \lambda p a_i)]}{\exp(\lambda tp)}. \quad (22)$$

Again using $1 + x \leq \exp(x)$,

$$\frac{\prod_{i=1}^n [(1-p) + p \exp(\lambda a_i)]}{\exp(\lambda tp + \lambda p \sum_{i=1}^n a_i)} \leq \frac{\exp(p(\exp(\lambda a_i) - 1))}{\exp(\lambda tp + \lambda p \sum_{i=1}^n a_i)}. \quad (23)$$

This is at most

$$\exp \left(p \sum_{i=1}^n (\exp(\lambda a_i) - 1) - \lambda tp - \lambda p \sum_{i=1}^n a_i \right) = \exp \left(\lambda p \left(\sum_{i=1}^n \frac{\exp(\lambda a_i) - \lambda a_i - 1}{\lambda} - t \right) \right). \quad (24)$$

Note that the difference in the sum is not too large and can be approximated by the square of the number in the exponent. That is, if $\lambda \leq \|\mathbf{a}\|_\infty^{-1}$ so that $|\lambda a_i| \leq 1$ for all i , then we have

$$\sum_{i=1}^n \frac{\exp(\lambda a_i) - \lambda a_i - 1}{\lambda} \leq \tilde{c} \sum_{i=1}^n \frac{\lambda^2 a_i^2}{\lambda} = \tilde{c} \lambda. \quad (25)$$

So (24) is at most

$$\exp(\lambda p(\tilde{c}\lambda - t)) \quad (26)$$

provided that $0 \leq \lambda \leq \|\mathbf{a}\|_\infty^{-1}$. Now we just need the right choice of λ based on t . If $0 \leq t \leq \|\mathbf{a}\|_\infty^{-1}$, we take $\lambda = \frac{t}{2\tilde{c}}$ which gives

$$\exp\left(\frac{t}{2\tilde{c}}p\left(-\frac{t}{2}\right)\right) = \exp\left(-\frac{t^2}{4\tilde{c}}p\right). \quad (27)$$

This is the subgaussian tail. If $t \geq \|\mathbf{a}\|_\infty^{-1}$, we take $\lambda = \frac{1}{2} \|\mathbf{a}\|_\infty^{-1}$ to find

$$\exp\left(\frac{1}{2\tilde{c}} \|\mathbf{a}\|_\infty^{-1} p\left(-\frac{t}{2}\right)\right) = \exp\left(-\frac{t \|\mathbf{a}\|_\infty^{-1} p}{4\tilde{c}}\right). \quad (28)$$

This is the subexponential tail. ■