

**IBM Research**

# Sieving for closest lattice vectors (with preprocessing)

Thijs Laarhoven

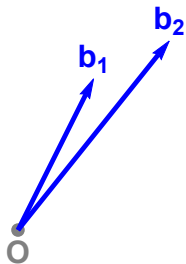
`mail@thijs.com`  
`http://www.thijs.com/`

Lorentz Center 2016, Leiden, The Netherlands  
(August 24, 2016)

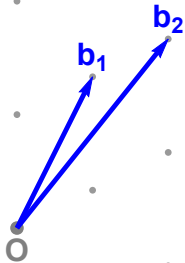
# Lattices



# Lattices

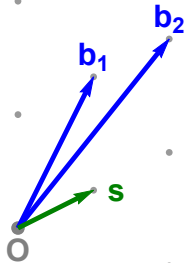


# Lattices



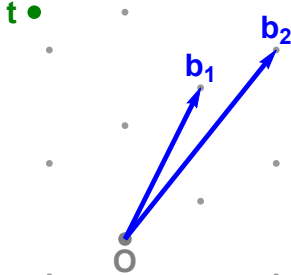
# Lattices

## Shortest Vector Problem (SVP)



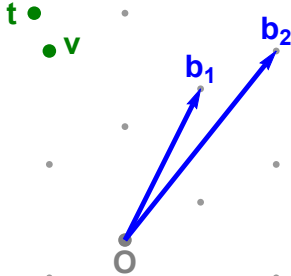
# Lattices

Closest Vector Problem (CVP)



# Lattices

## Closest Vector Problem (CVP)





# Outline

Sieving for SVP

Sieving for CVP

Sieving for CVPP

Conclusion





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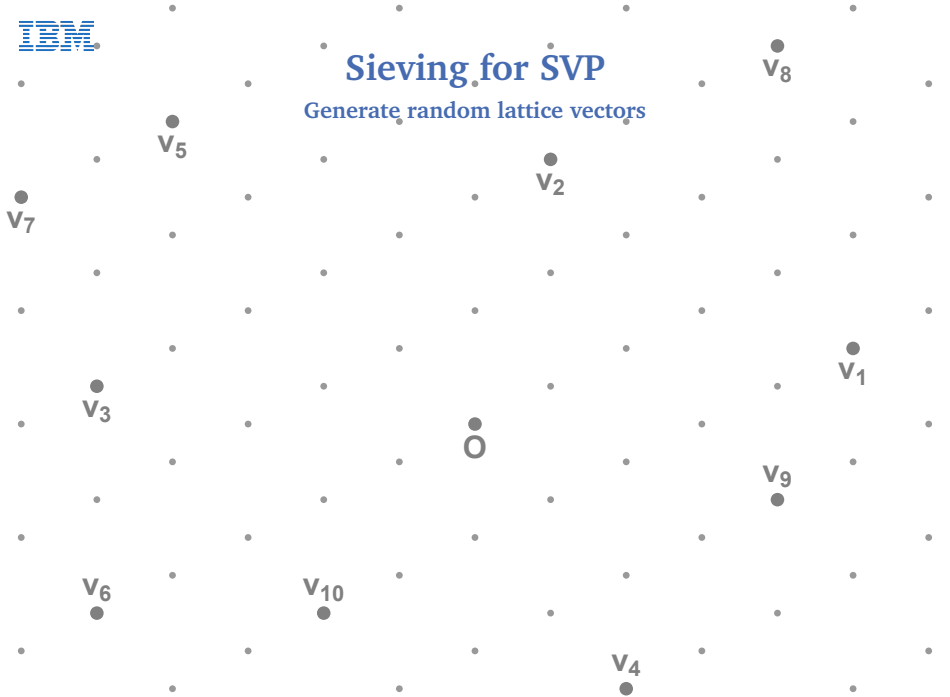
# Sieving for SVP

Generate random lattice vectors



# Sieving for SVP

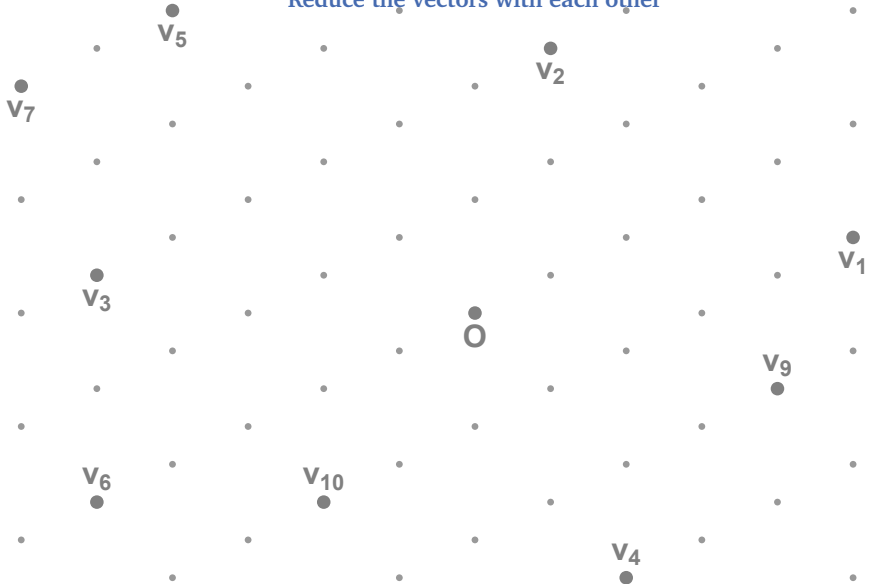
Generate random lattice vectors





# Sieving for SVP

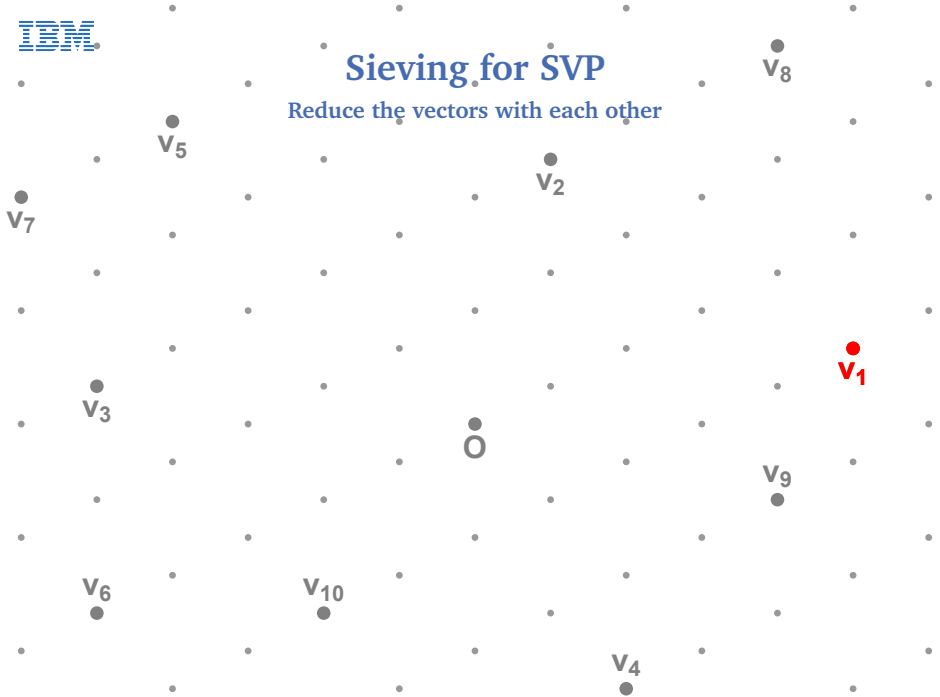
Reduce the vectors with each other





# Sieving for SVP

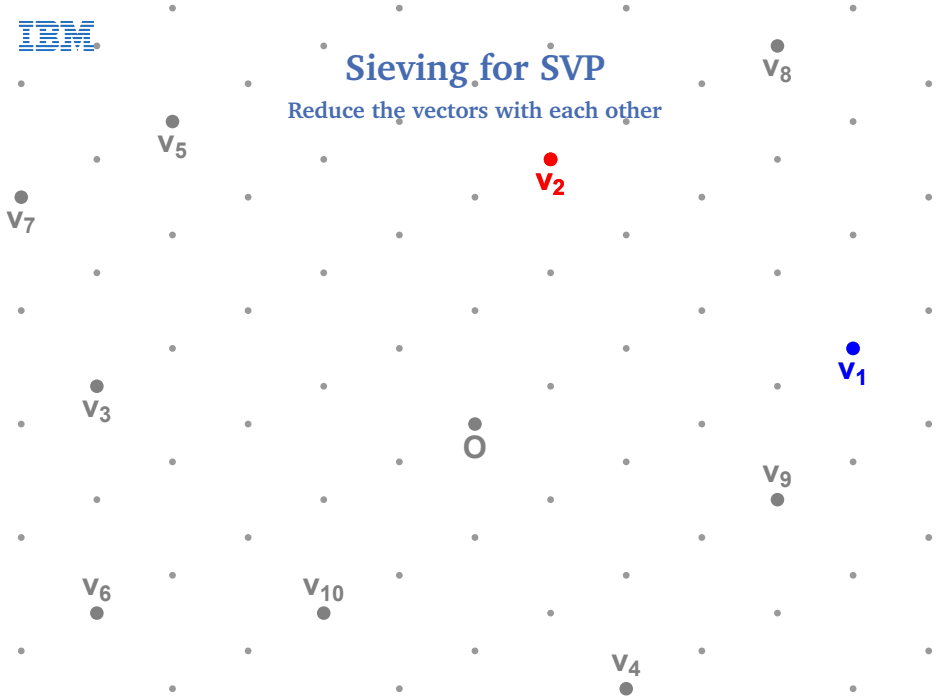
Reduce the vectors with each other





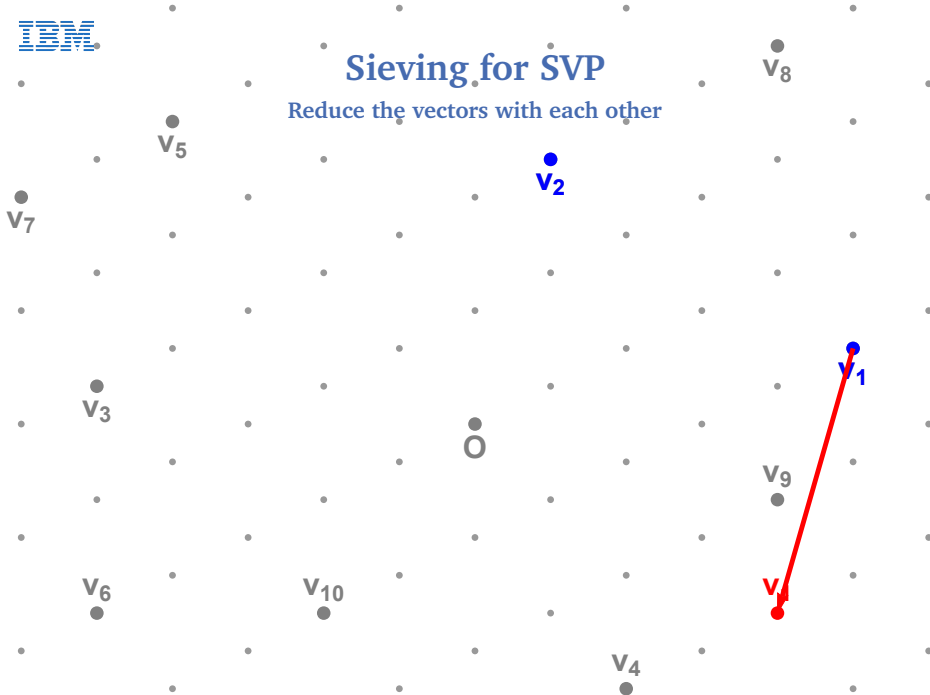
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Reduce the vectors with each other



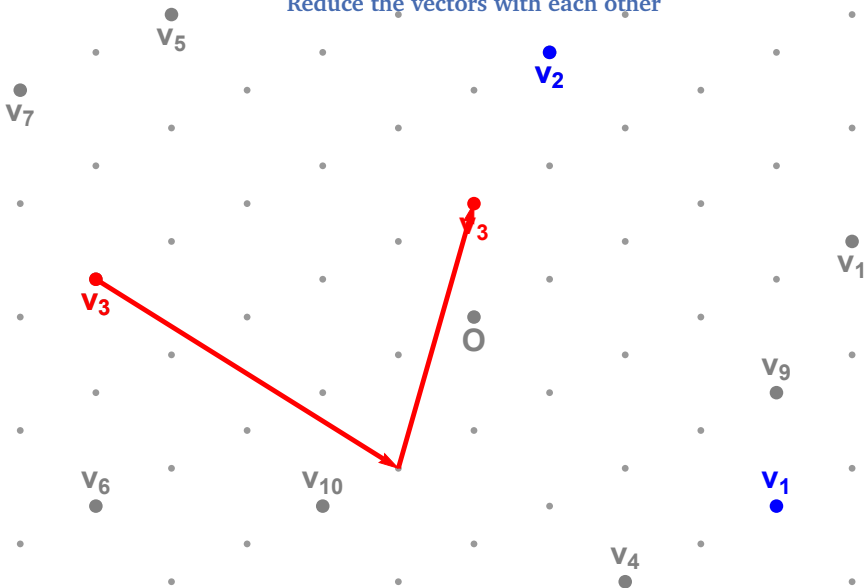
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# Sieving for SVP

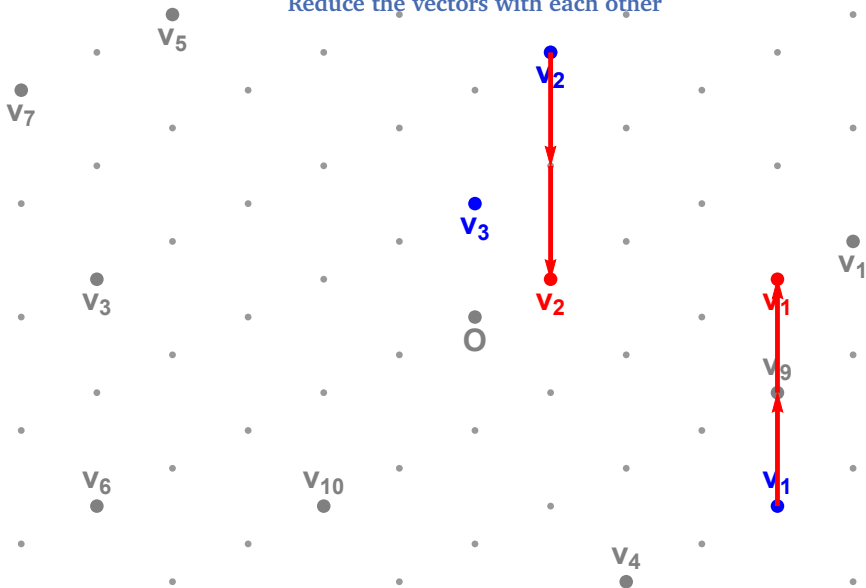
Reduce the vectors with each other





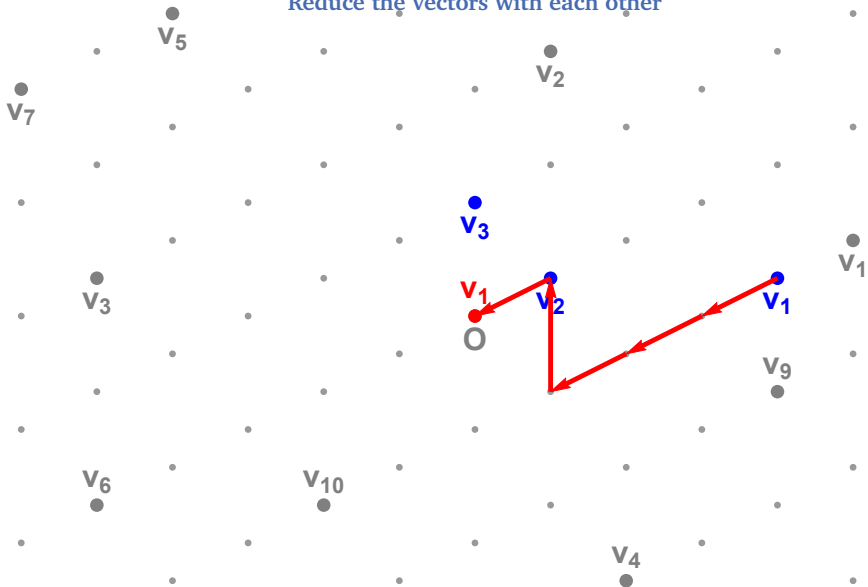
# Sieving for SVP

Reduce the vectors with each other



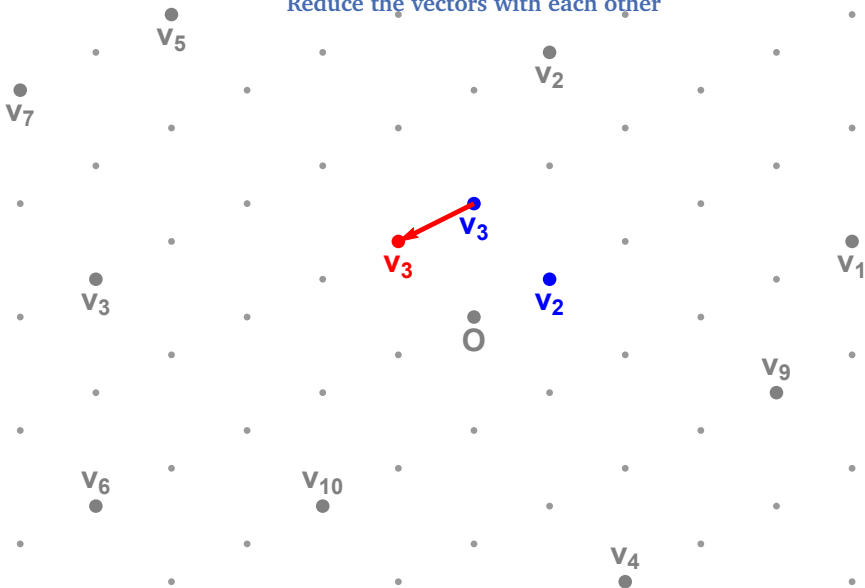
# Sieving for SVP

Reduce the vectors with each other



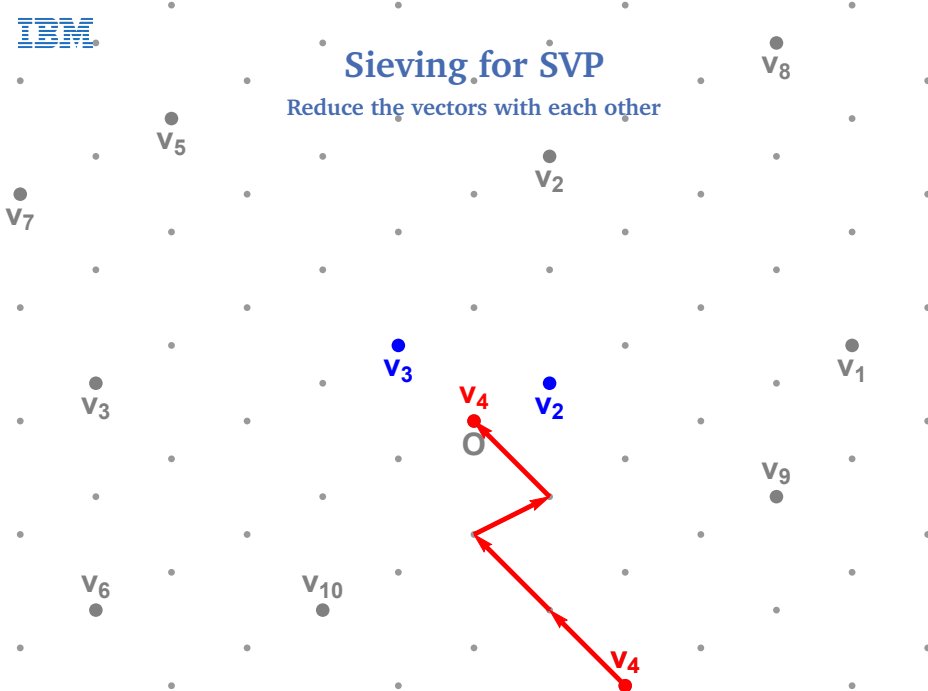
# Sieving for SVP

Reduce the vectors with each other



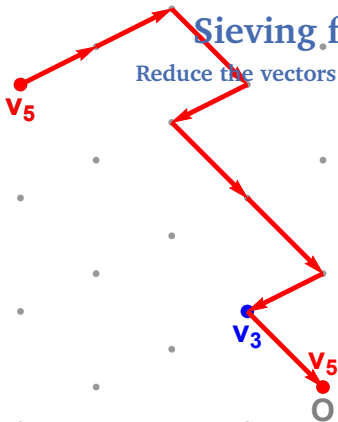
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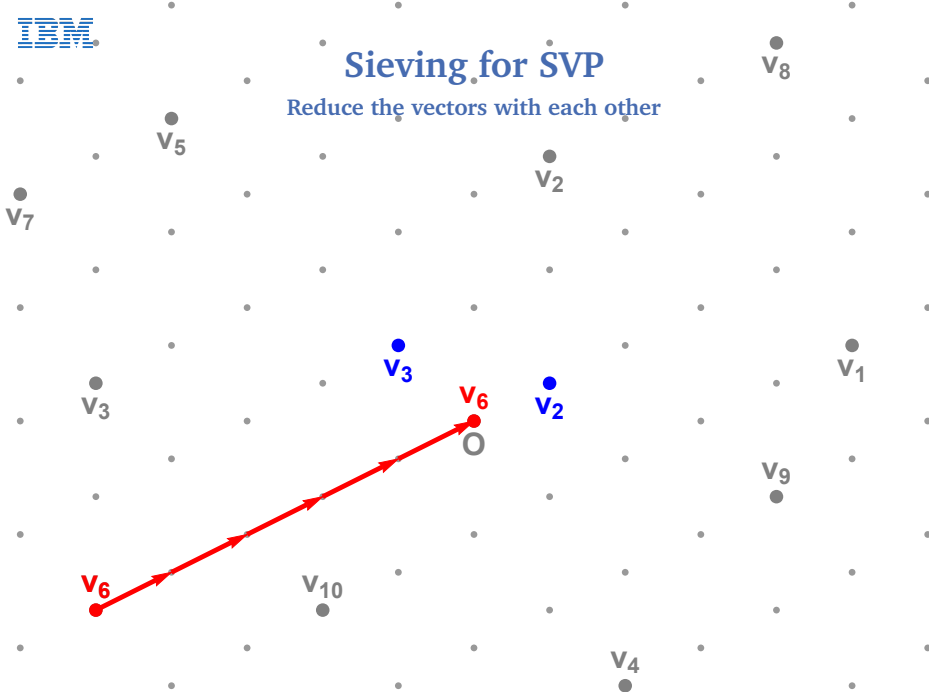
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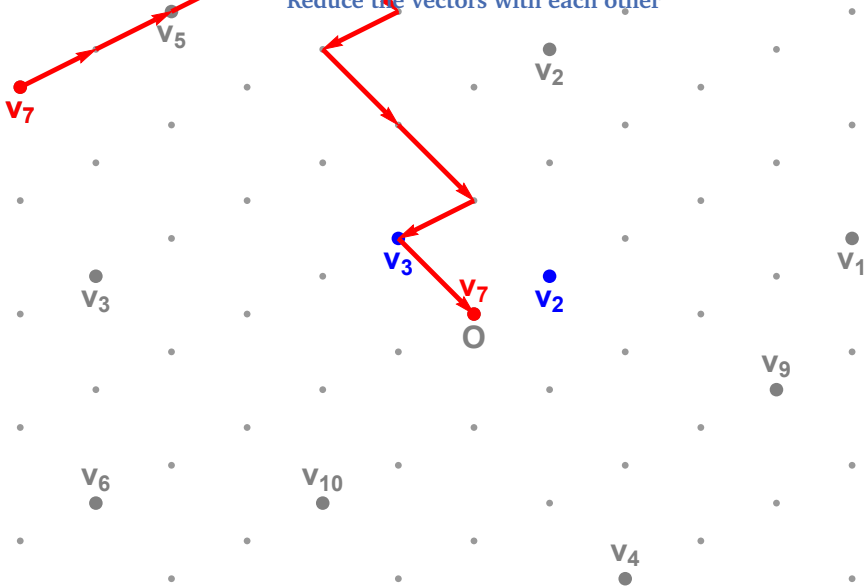
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Reduce the vectors with each other



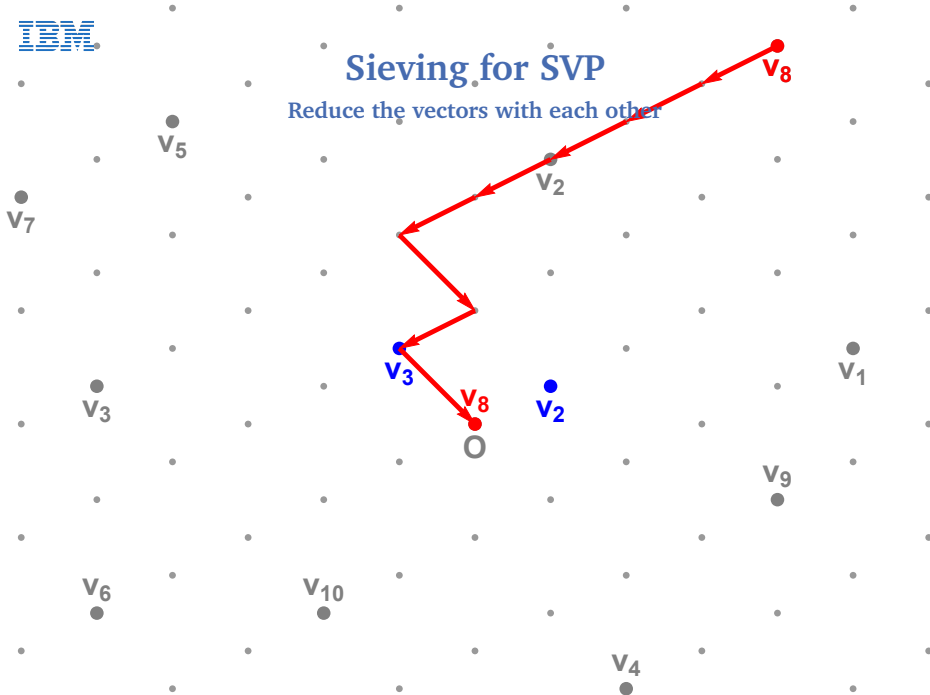
# Sieving for SVP

Reduce the vectors with each other



# Sieving for SVP

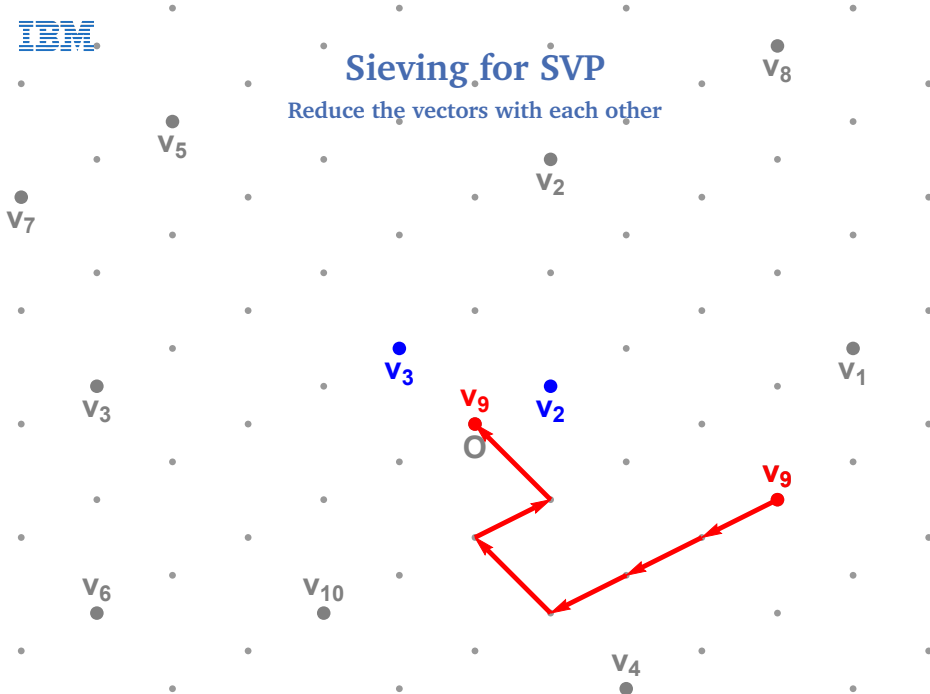
Reduce the vectors with each other





# Sieving for SVP

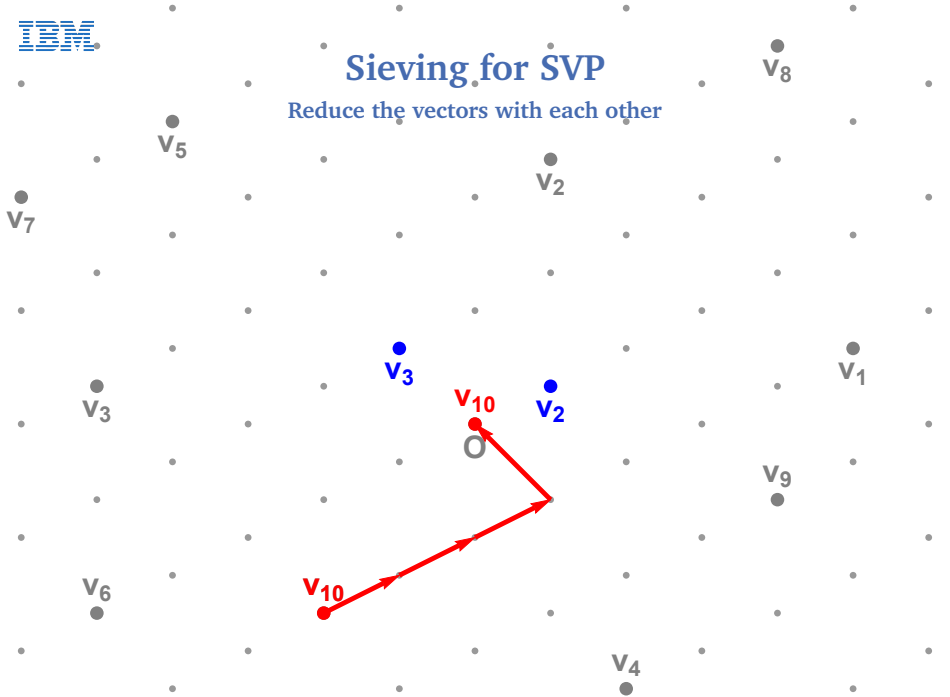
Reduce the vectors with each other





# Sieving for SVP

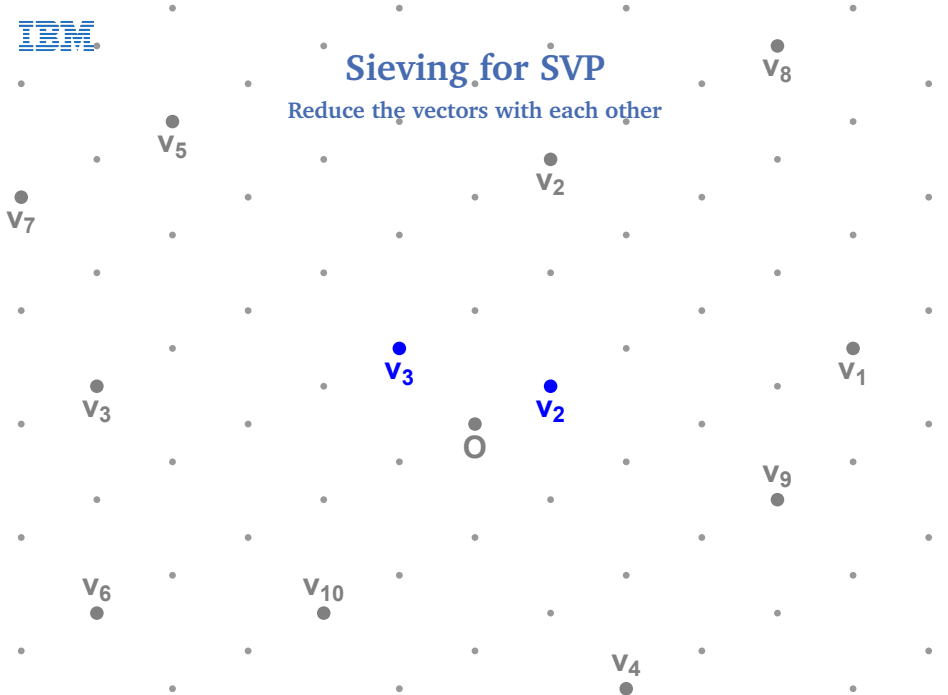
Reduce the vectors with each other





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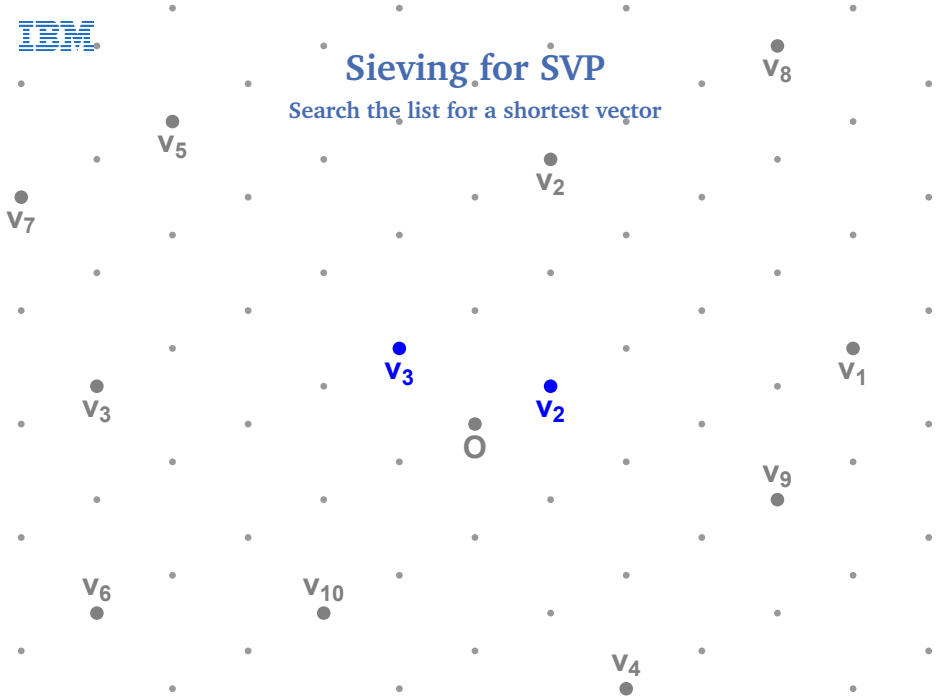
Reduce the vectors with each other





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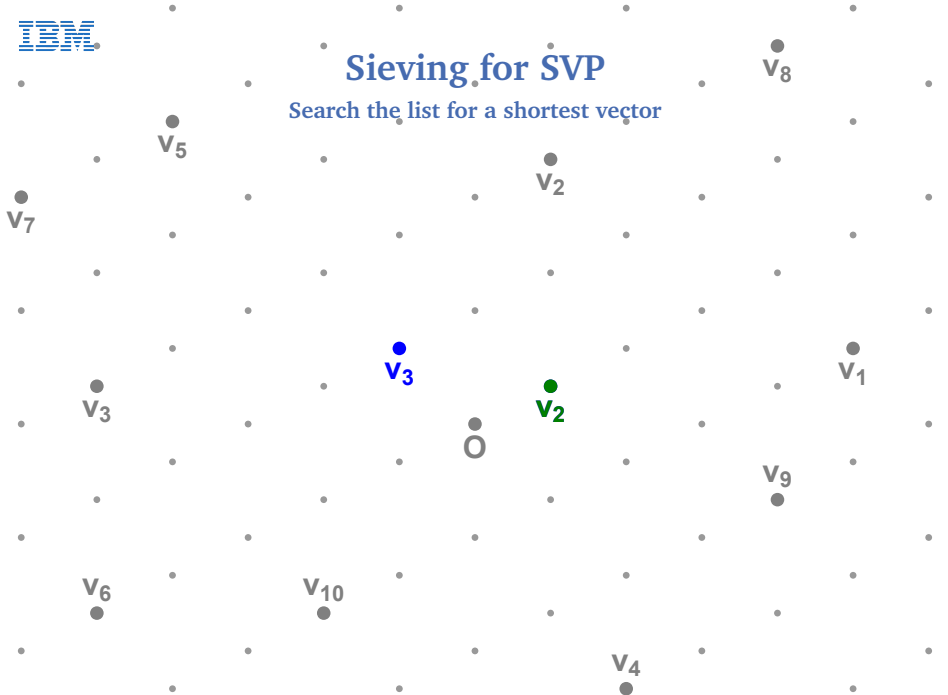
Search the list for a shortest vector





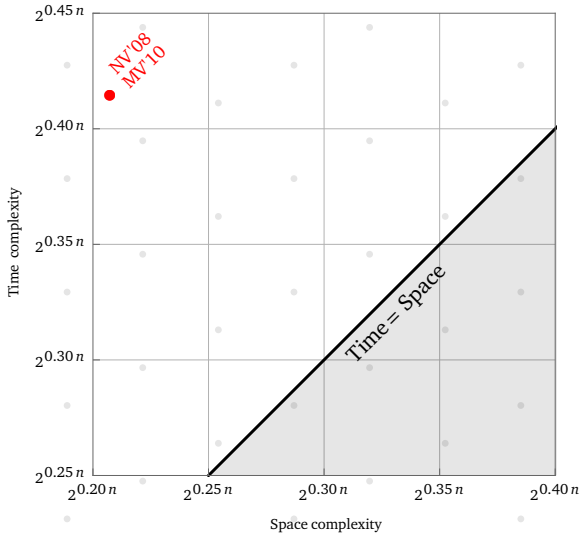
# Sieving for SVP

Search the list for a shortest vector



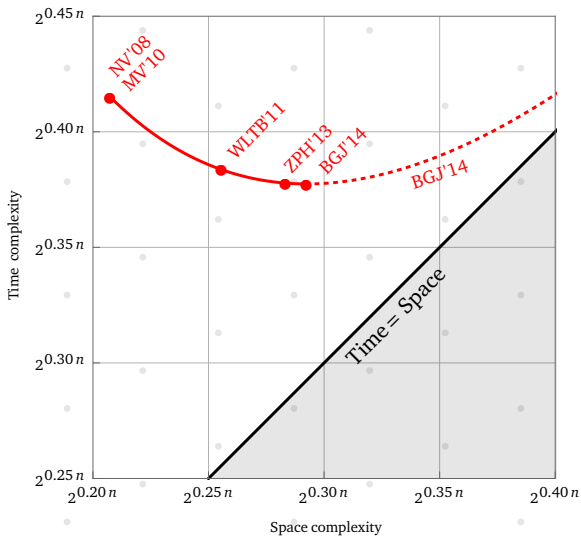
# Sieving for SVP

The GaussSieve and Nguyen-Vidick sieve



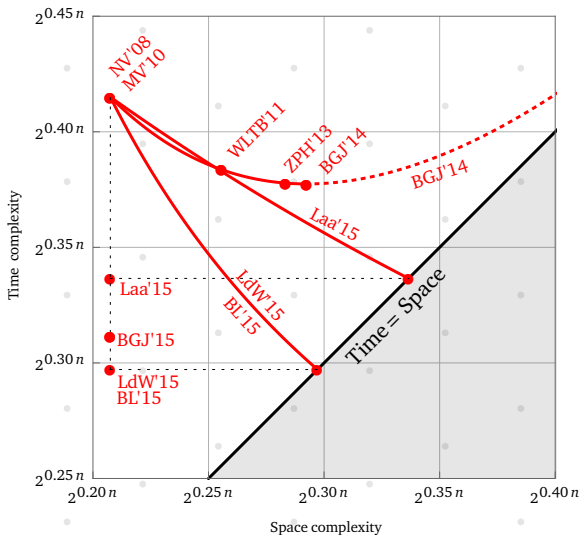
# Sieving for SVP

## Leveled sieving approaches



# Sieving for SVP

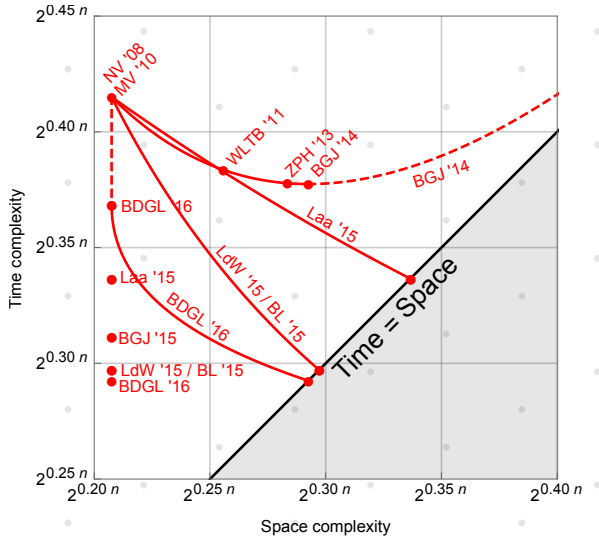
## Locality-Sensitive Hashing (LSH)





# Sieving for SVP

## Locality-Sensitive Filters (LSF)





# Outline

Sieving for SVP

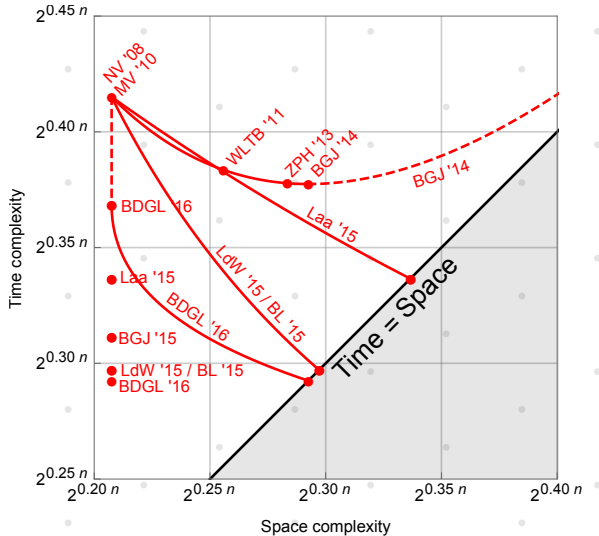
Sieving for CVP

Sieving for CVPP

Conclusion

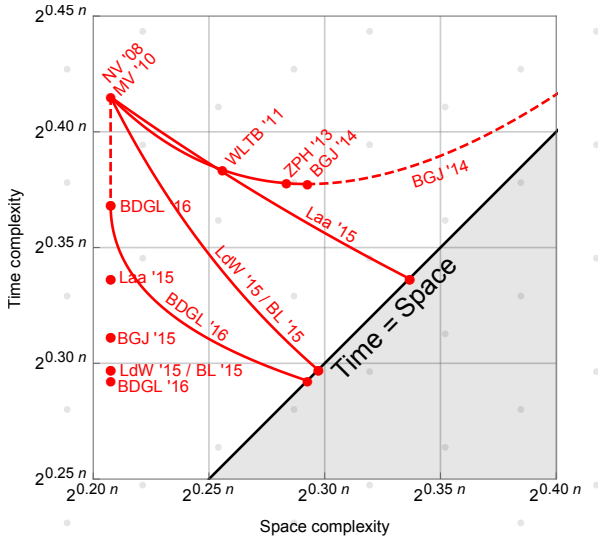
# Sieving for SVP

Space/time trade-offs



# Sieving for CVP

Space/time trade-offs





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# Sieving for CVPP

Run a GaussSieve as preprocessing

# Sieving for CVPP

Run a GaussSieve as preprocessing



$v_2$

$v_1$

$O$

# Sieving for CVPP

Reduce the target vectors with the list





# Sieving for CVPP

Reduce the target vectors with the list



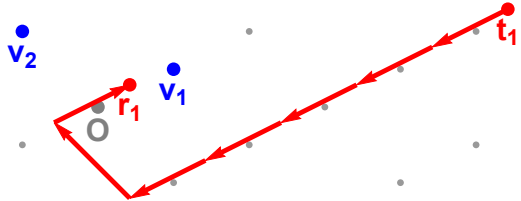
$v_2$

$v_1$

$t_1$

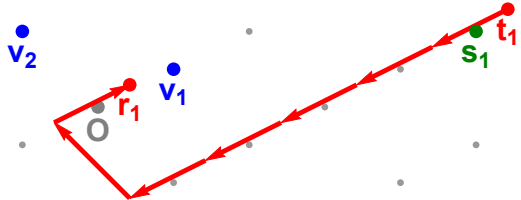
# Sieving for CVPP

Reduce the target vectors with the list



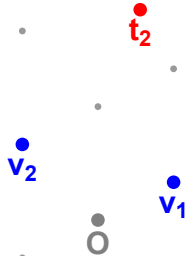
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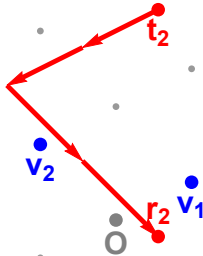
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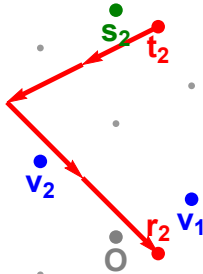
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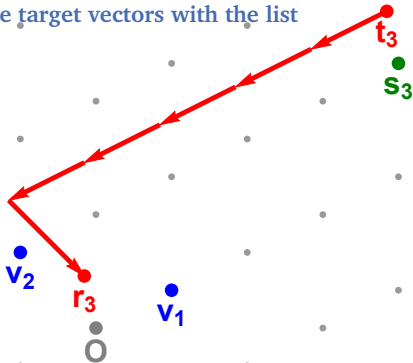
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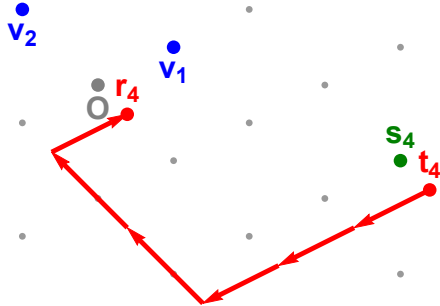
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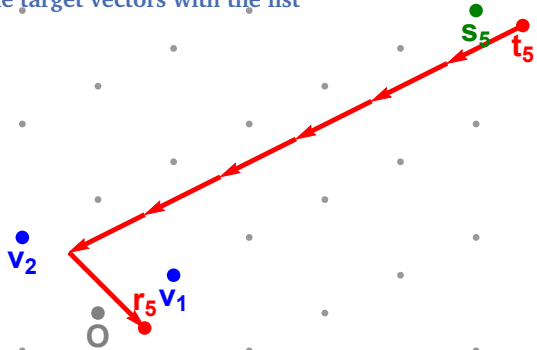
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# Sieving for CVPP

Reduce the target vectors with the list



# Sieving for CVPP

Reduce the target vectors with the list

$v_2$

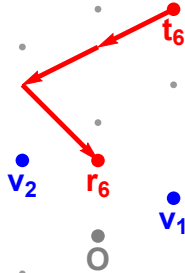
$v_1$

$O$

$t_6$

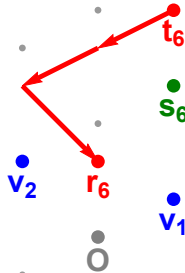
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Reduce the target vectors with the list



# Sieving for CVPP

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# Sieving for CVPP

Relation with the Voronoi cell



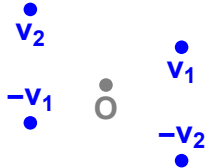
$v_2$

$v_1$

O

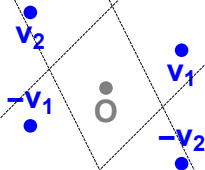
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Relation with the Voronoi cell



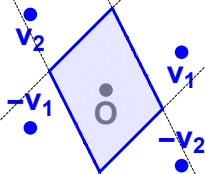
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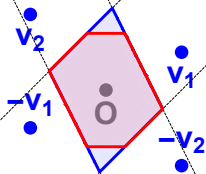
Relation with the Voronoi cell





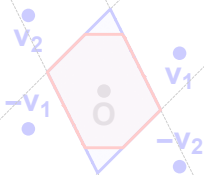
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# Sieving for CVPP

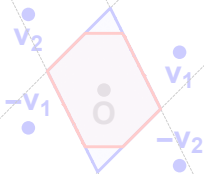
## Overview



# Sieving for CVPP

## Overview

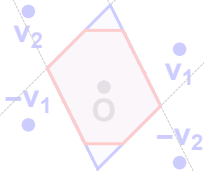
- Blue region: Gauss cell  $\mathcal{G}$



# Sieving for CVPP

## Overview

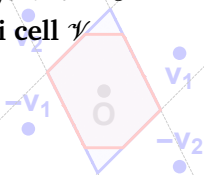
- Blue region: **Gauss cell**  $\mathcal{G}$ 
  - ▶ Defined by  $2^{0.21n+o(n)}$  short lattice vectors
  - ▶ Volume:  $\text{Vol}(\mathcal{G}) = 2^{O(n)} \cdot \det(\mathcal{L})$
  - ▶ Reductions always land in  $\mathcal{G}$



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- Red region: **Voronoi cell**  $\mathcal{V}$



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  - ▶ Volume:  $\text{Vol}(\mathcal{V}) = \det(\mathcal{L})$
  - ▶ Reductions almost never land in  $\mathcal{V}$

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- Problems:

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  - ▶ Volume:  $\text{Vol}(\mathcal{V}) = \det(\mathcal{L})$
  - ▶ Reductions almost never land in  $\mathcal{V}$
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  - ▶ Exponentially small success probability  $\text{Vol}(\mathcal{V})/\text{Vol}(\mathcal{G})$



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  - ▶ Reductions almost never land in  $\mathcal{V}$
- Problems:
  - ▶ Exponentially small success probability  $\text{Vol}(\mathcal{V})/\text{Vol}(\mathcal{G})$
  - ▶ Probability only over randomness of targets



# Sieving for CVPP

Solving the problems

- Idea 1: **Larger lists, weaker reductions**

# Sieving for CVPP

Solving the problems

- Idea 1: **Larger lists, weaker reductions**
  - Problem: Exponentially small success probability

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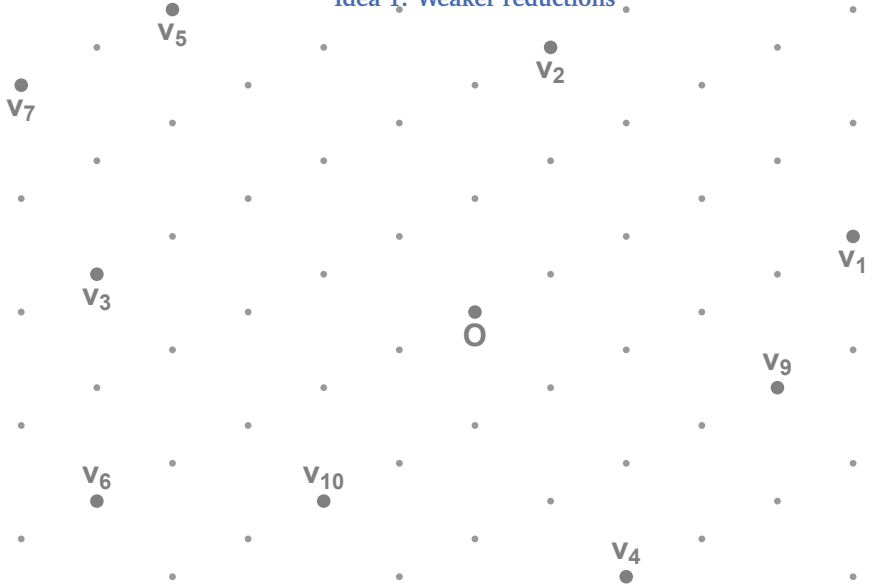
Solving the problems

- Idea 1: **Larger lists, weaker reductions**
  - ▶ Problem: Exponentially small success probability
  - ▶ To guarantee  $\text{Vol}(\mathcal{G}) \approx \text{Vol}(\mathcal{V})$ , need  $2^{n/2+o(n)}$  vectors
  - ▶ Preprocessing: reduce  $\mathbf{v}_1$  with  $\mathbf{v}_2$  iff  $\|\mathbf{v}_1 - \mathbf{v}_2\| \ll \|\mathbf{v}_1\|$



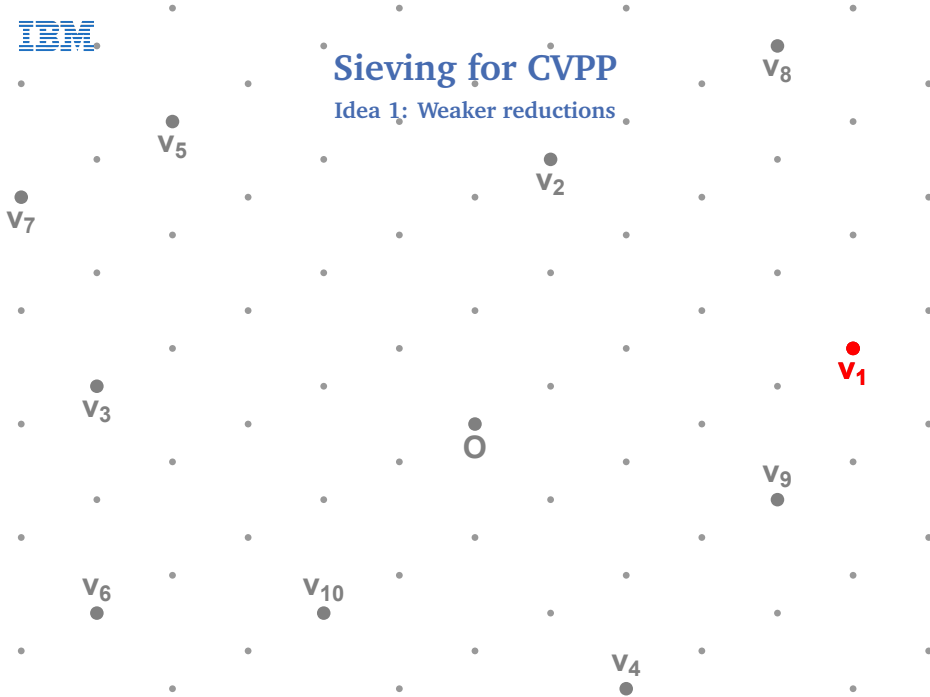
# Sieving for CVPP

Idea 1: Weaker reductions



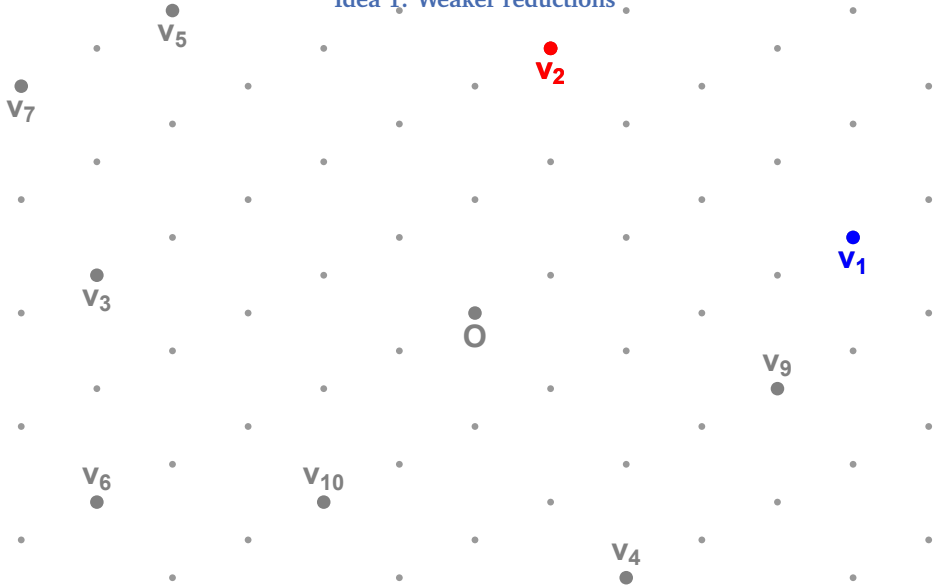
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Idea 1: Weaker reductions



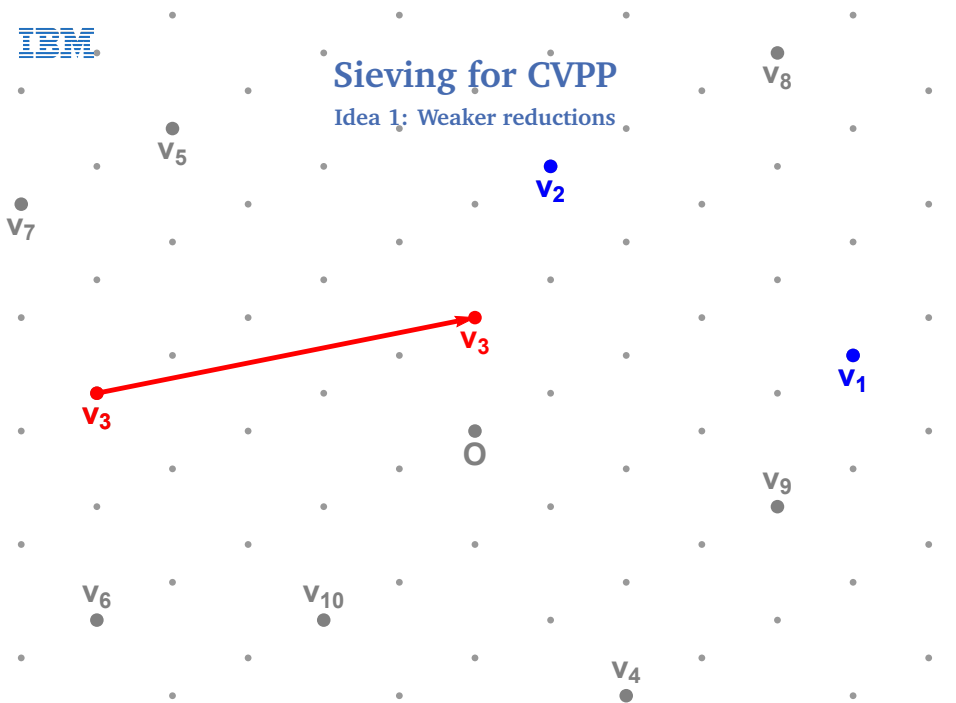
# Sieving for CVPP

Idea 1: Weaker reductions



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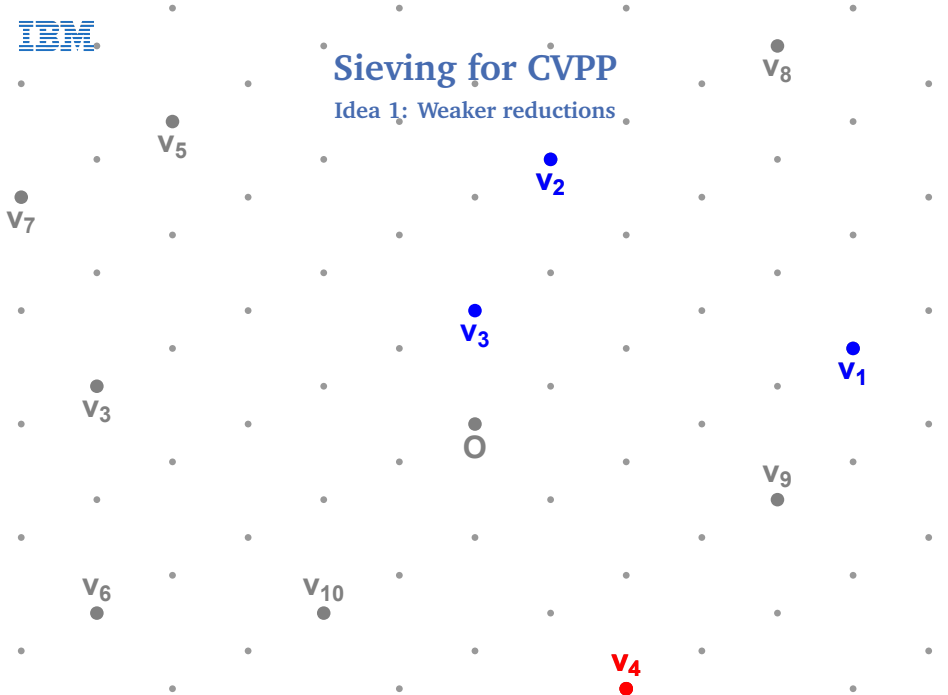
Idea 1: Weaker reductions





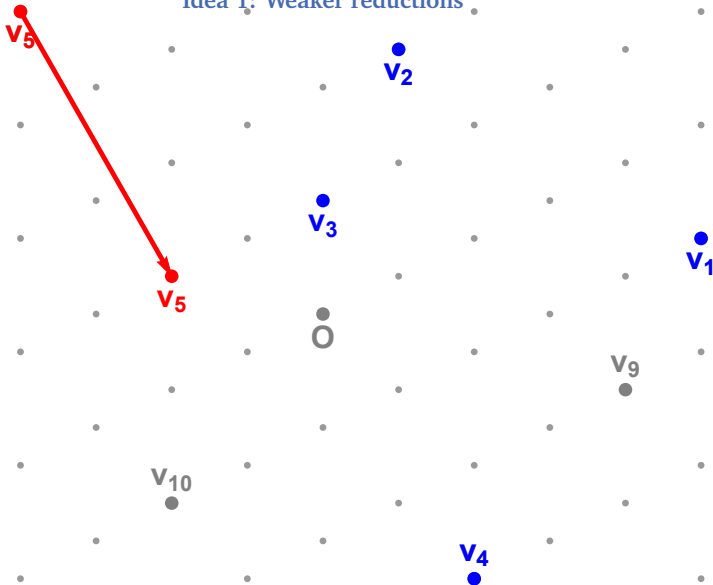
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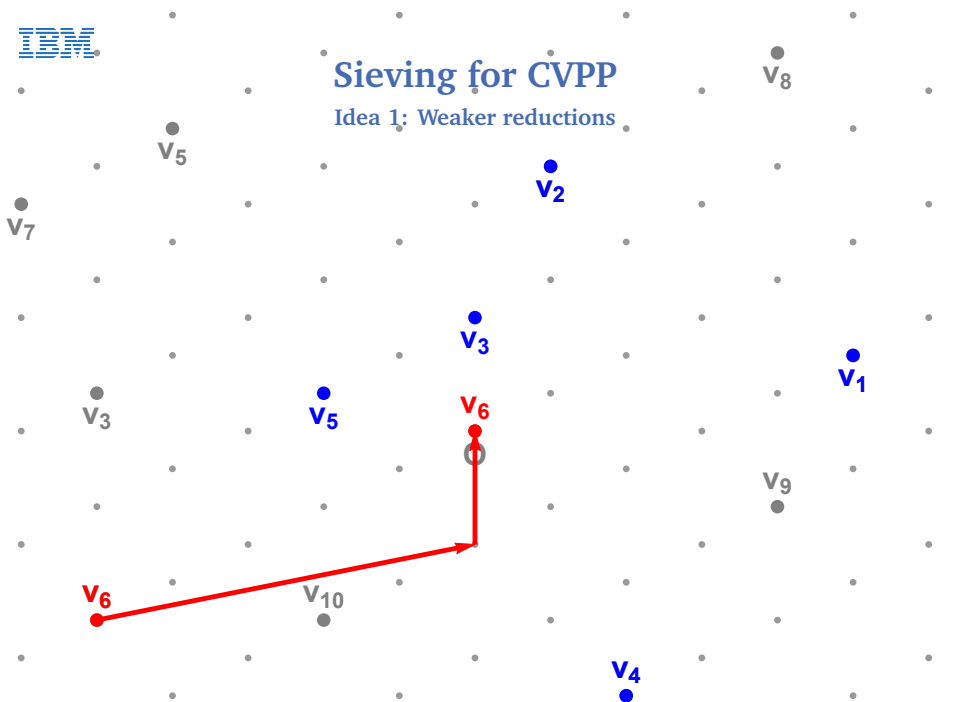
# Sieving for CVPP

Idea 1: Weaker reductions



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Idea 1: Weaker reductions



# Sieving for CVPP

Idea 1: Weaker reductions

$v_7$

$v_5$

$v_8$

$v_2$

$v_3$

$v_1$

$v_5$

$O$

$v_9$

$v_3$

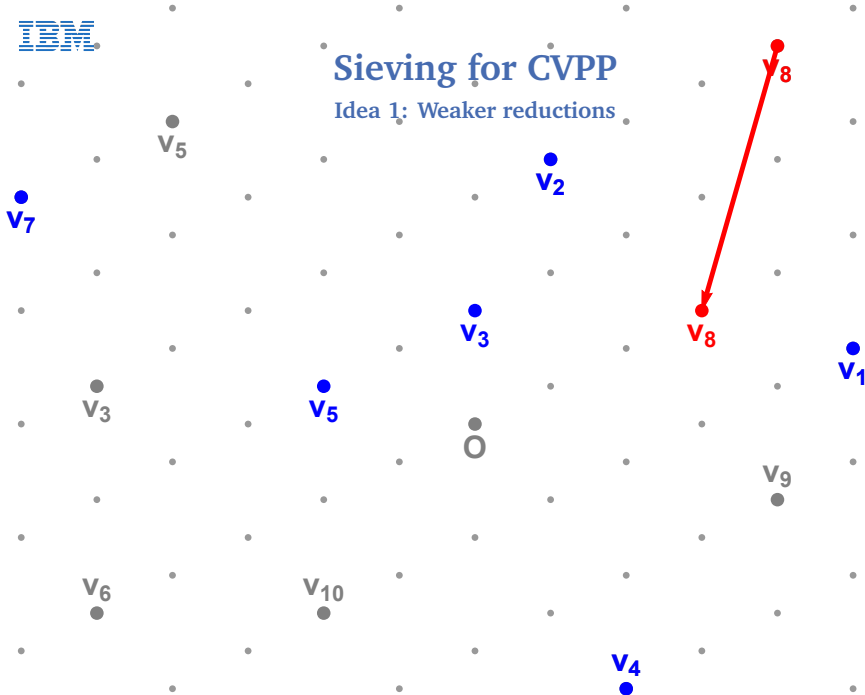
$v_6$

$v_{10}$

$v_4$

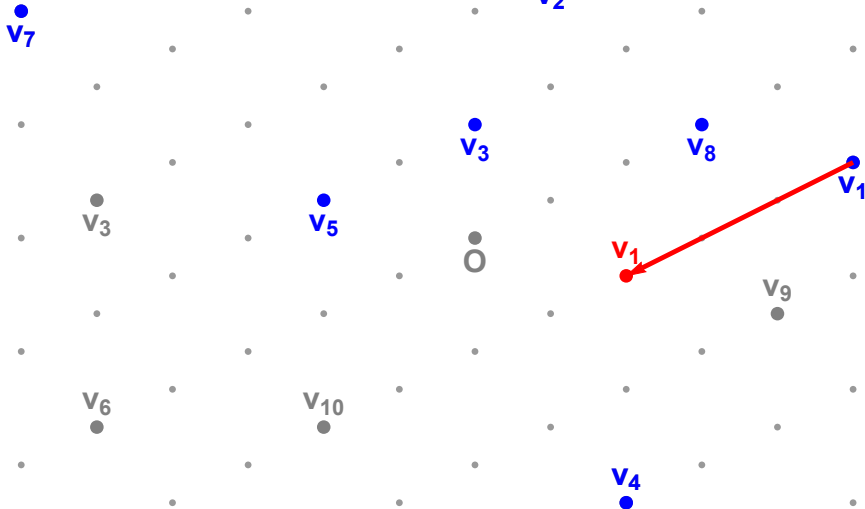
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$v_8$

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$v_4$



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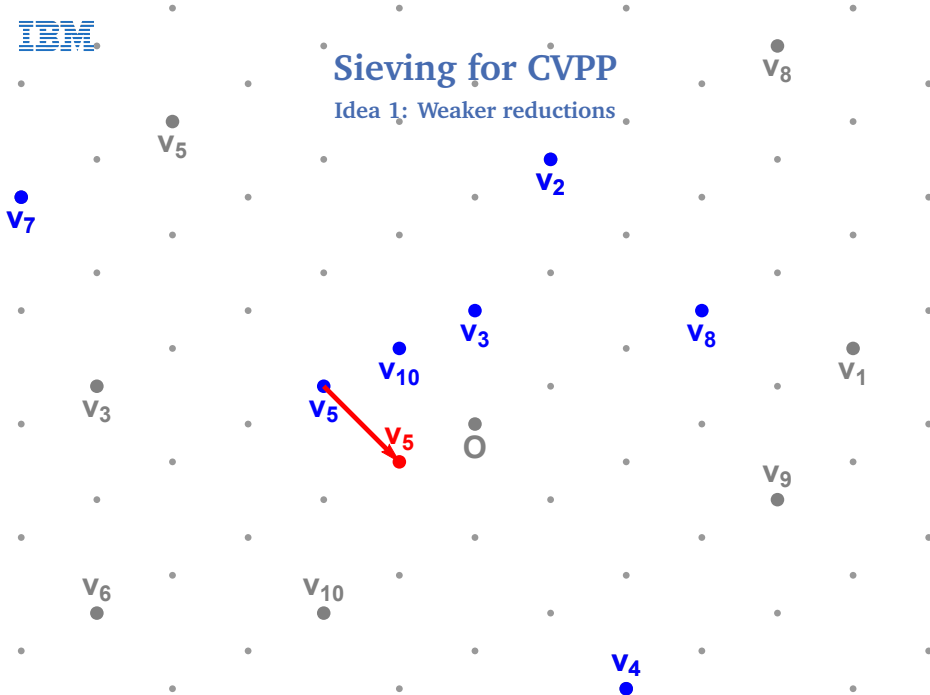
$v_{10}$

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# Sieving for CVPP

## Solving the problems

- Idea 1: **Larger lists, weaker reductions**
  - ▶ Problem: Exponentially small success probability
  - ▶ To guarantee  $\text{Vol}(\mathcal{G}) \approx \text{Vol}(\mathcal{V})$ , need  $2^{n/2+o(n)}$  vectors
  - ▶ Preprocessing: reduce  $\mathbf{v}_1$  with  $\mathbf{v}_2$  iff  $\|\mathbf{v}_1 - \mathbf{v}_2\| \ll \|\mathbf{v}_1\|$

# Sieving for CVPP

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  - ▶ Randomize target  $\mathbf{t}$  before reducing ( $\mathbf{t}' \in_R \mathbf{t} + \mathcal{L}$ )
  - ▶ Randomness now over algorithm, independently of target

# Sieving for CVPP

Idea 2: Rerandomize the target

$v_2$

$v_1$

$0$



# Sieving for CVPP

Idea 2: Rerandomize the target

$v_2$

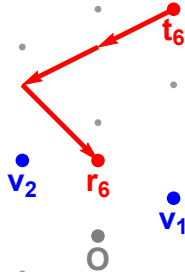
$v_1$

$O$

$t_6$

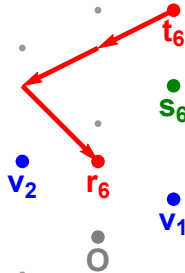
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$O$

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$t_6$

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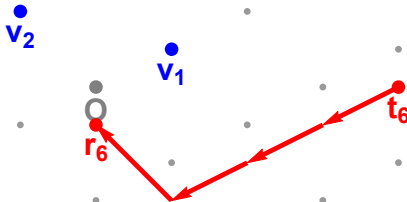
$v_1$

$O$

$t_6$

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$\dot{V}_2$ 
$$\mathbf{v}_1$$

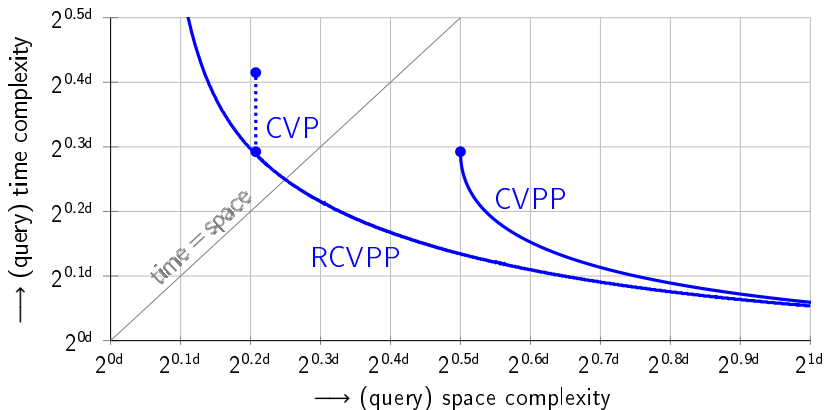
**S<sub>6</sub>**

**r6**

 $t_6$

# Sieving for CVPP

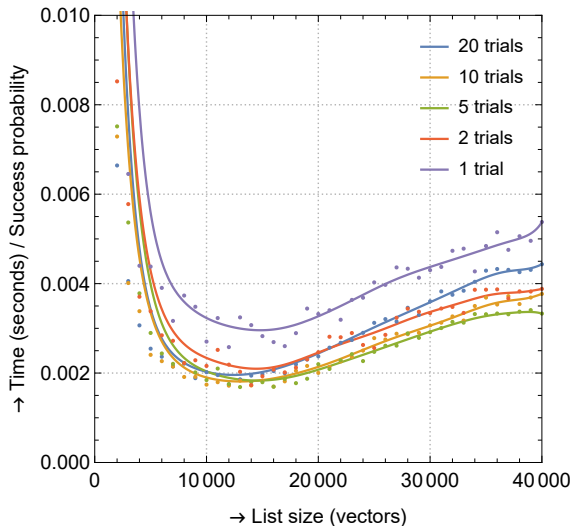
Trade-offs





# Sieving for CVPP

Preliminary experiments ( $n = 50$ , HashSieve)



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  - ▶ Randomized CVPP  $\implies$  embarrassingly parallel

# Questions?

