

Sieving for shortest lattice vectors using fast search algorithms

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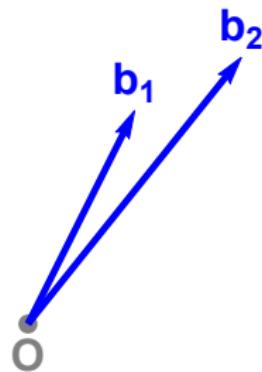
Lattices

What is a lattice?



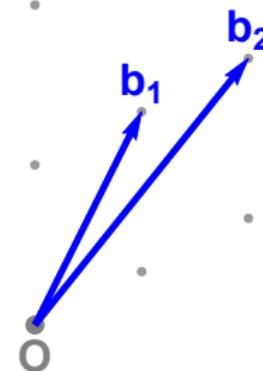
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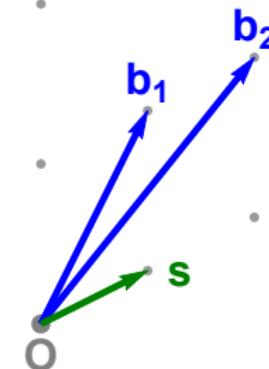
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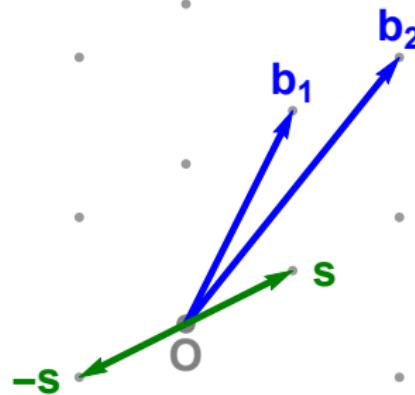
Lattices

Shortest Vector Problem (SVP)



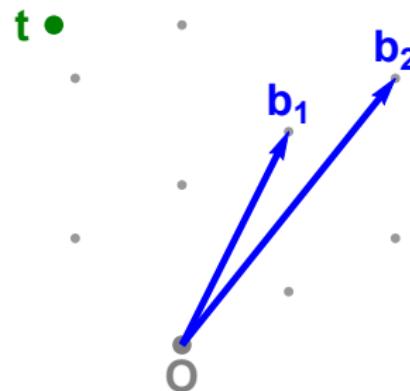
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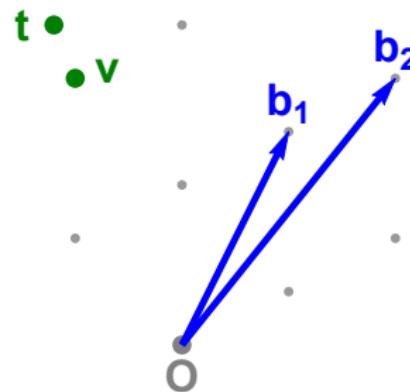
Lattices

Closest Vector Problem (CVP)



Lattices

Closest Vector Problem (CVP)



Lattices

Applications

- “Constructive cryptography”: Lattice-based cryptosystems
 - ▶ Based on hard lattice problems (SVP, CVP, LWE, SIS)
 - ▶ NTRU cryptosystem [[HPS98, ...](#)]
 - ▶ Fully Homomorphic Encryption [[Gen09, ...](#)]
 - ▶ Candidate for post-quantum cryptography

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 - ▶ Attack RSA with Coppersmith’s method [Cop97]
 - ▶ Attack lattice-based cryptosystems [Ngu99, JJ00]

Lattices

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How hard are hard lattice problems such as SVP?

Nguyen-Vidick sieve

O

Nguyen-Vidick sieve

1. Sample a list L of random lattice vectors

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Nguyen-Vidick sieve

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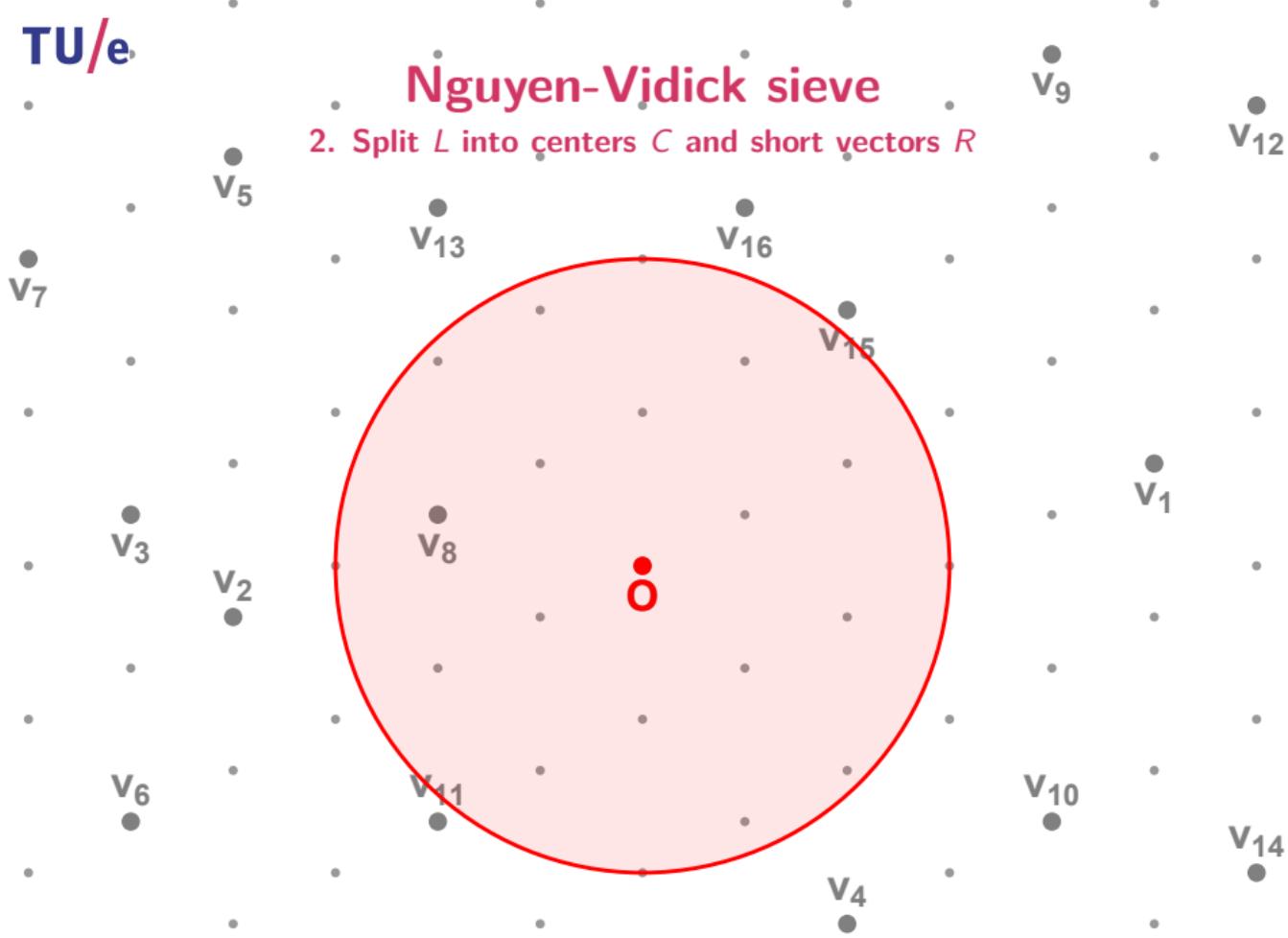
Nguyen-Vidick sieve

2. Split L into centers C and short vectors R



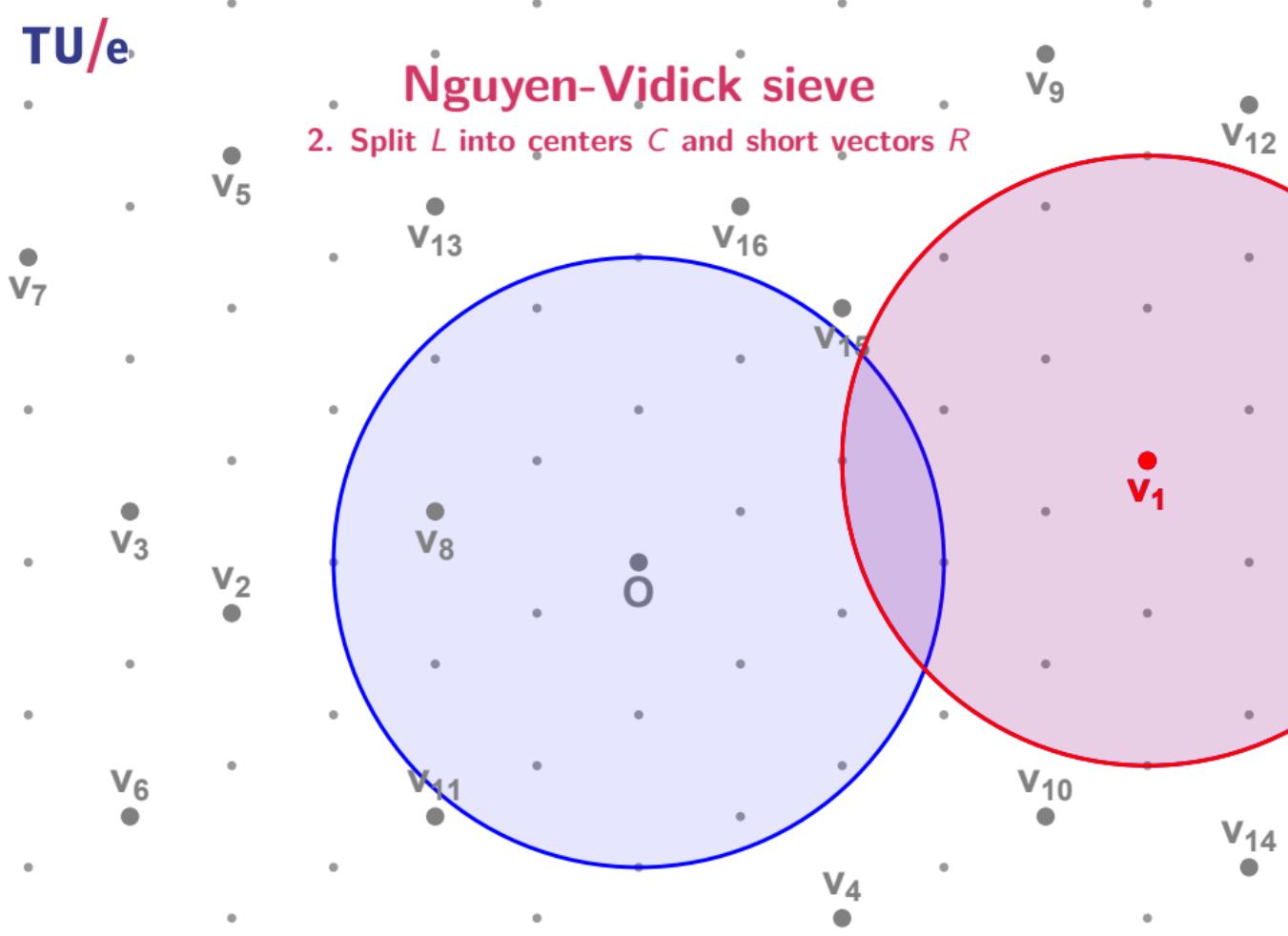
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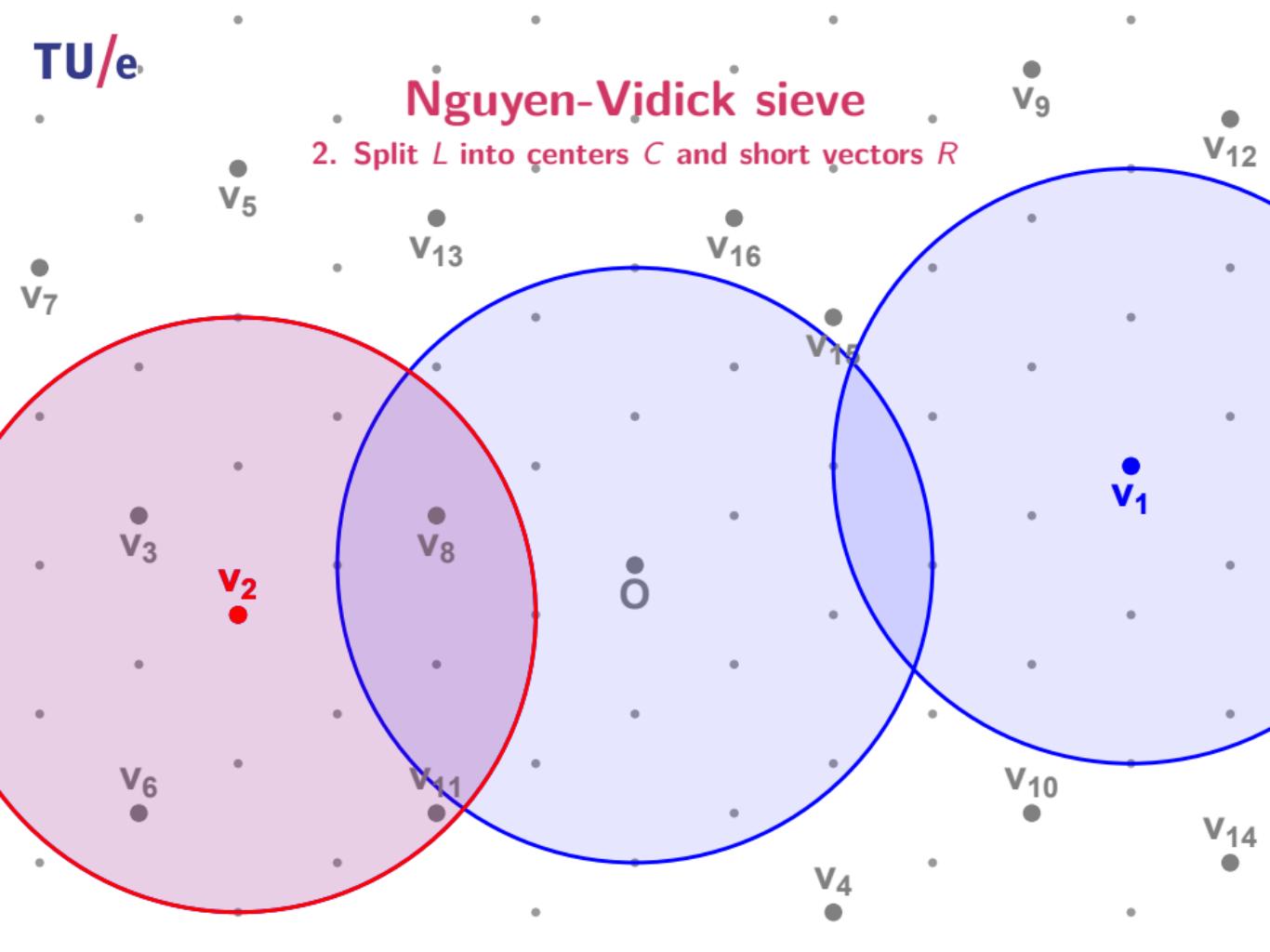
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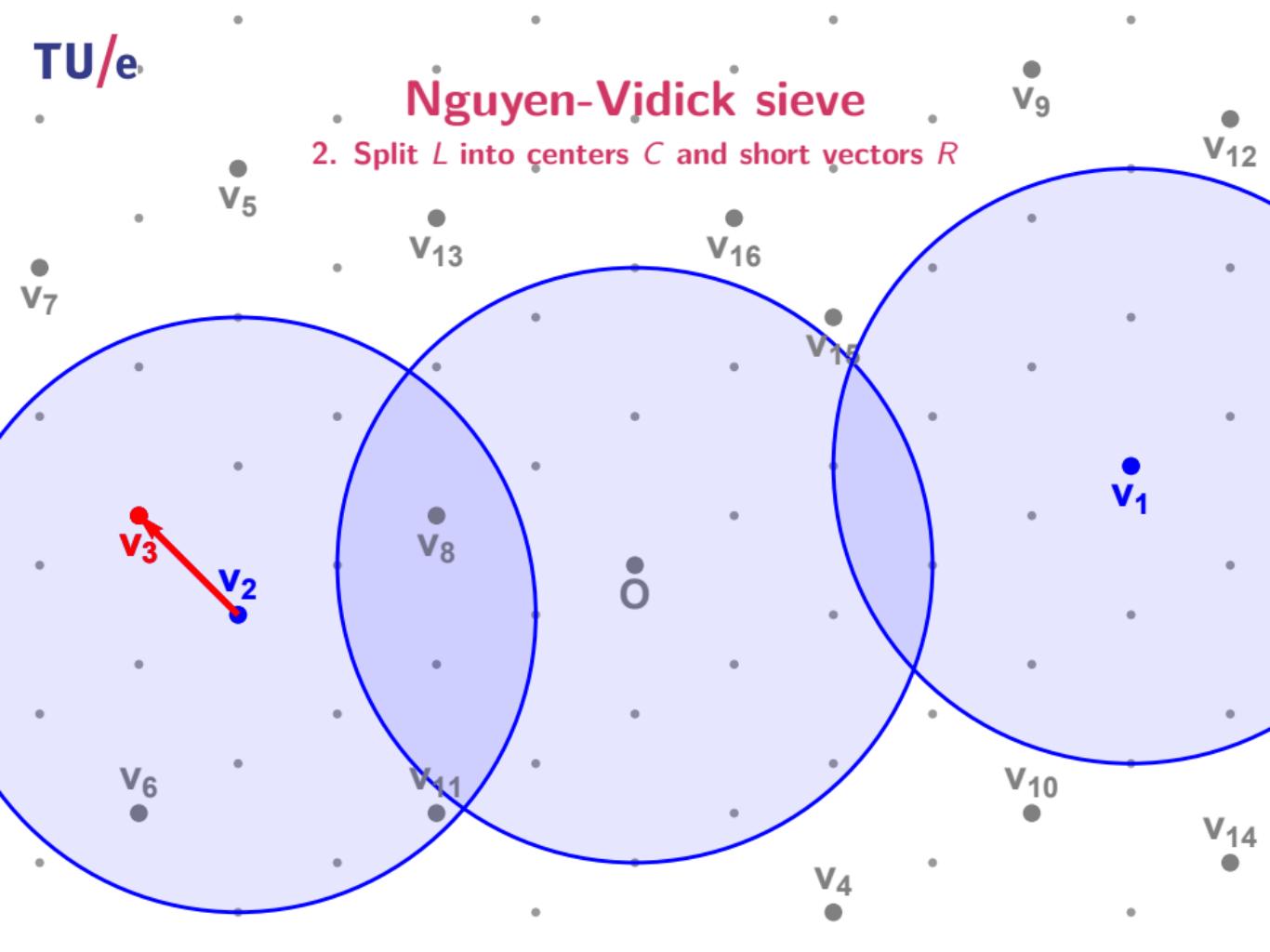
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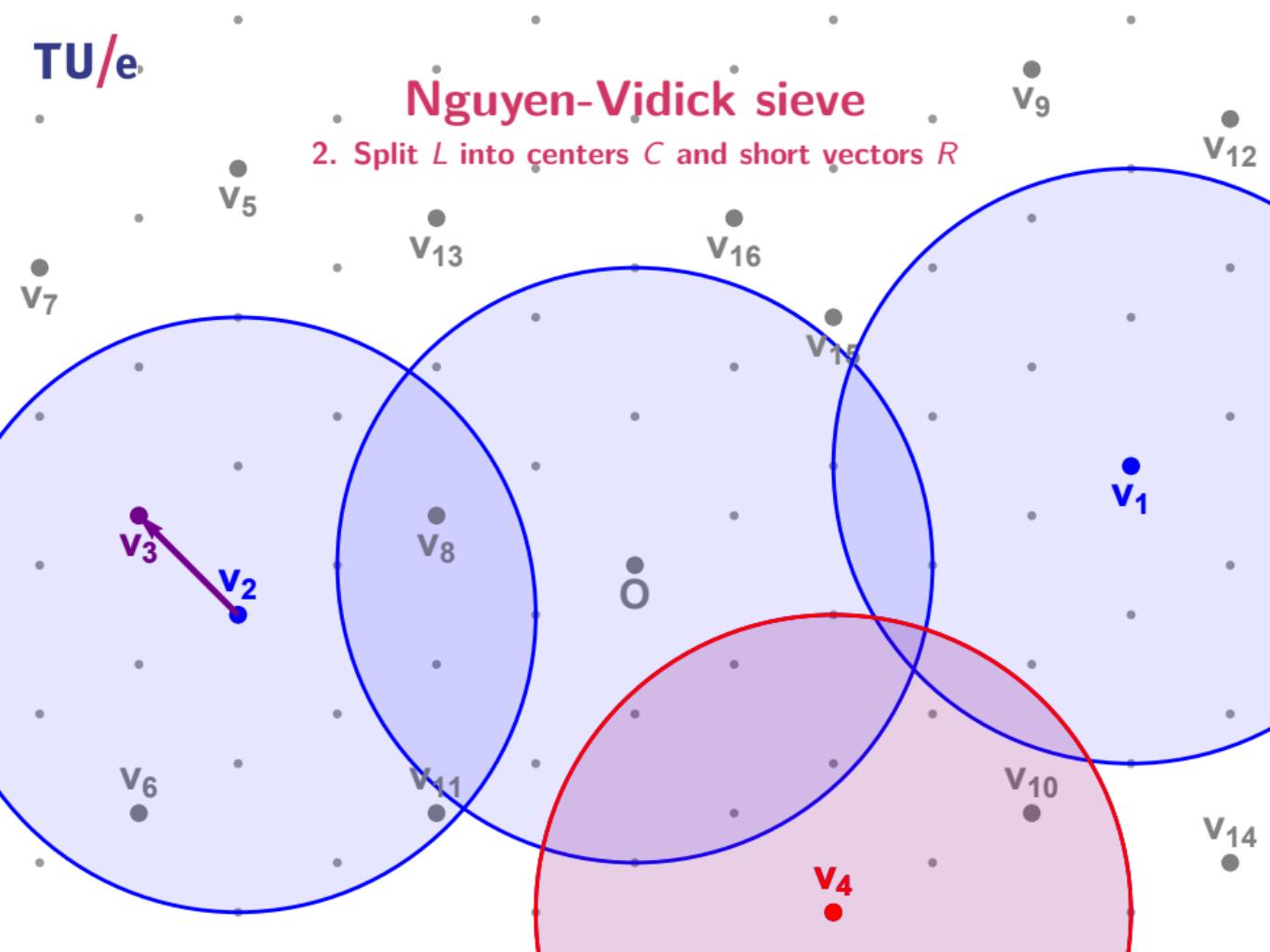
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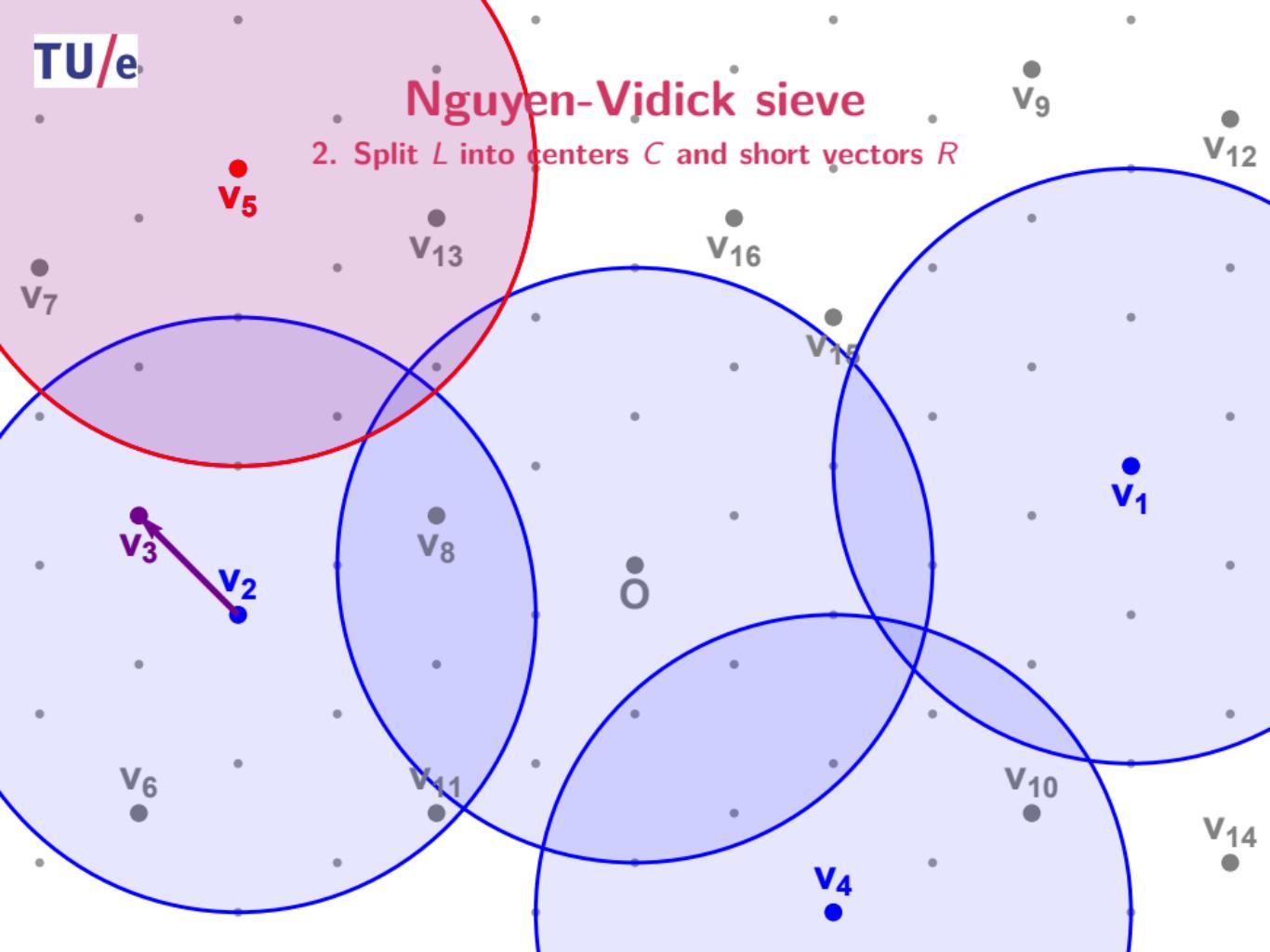
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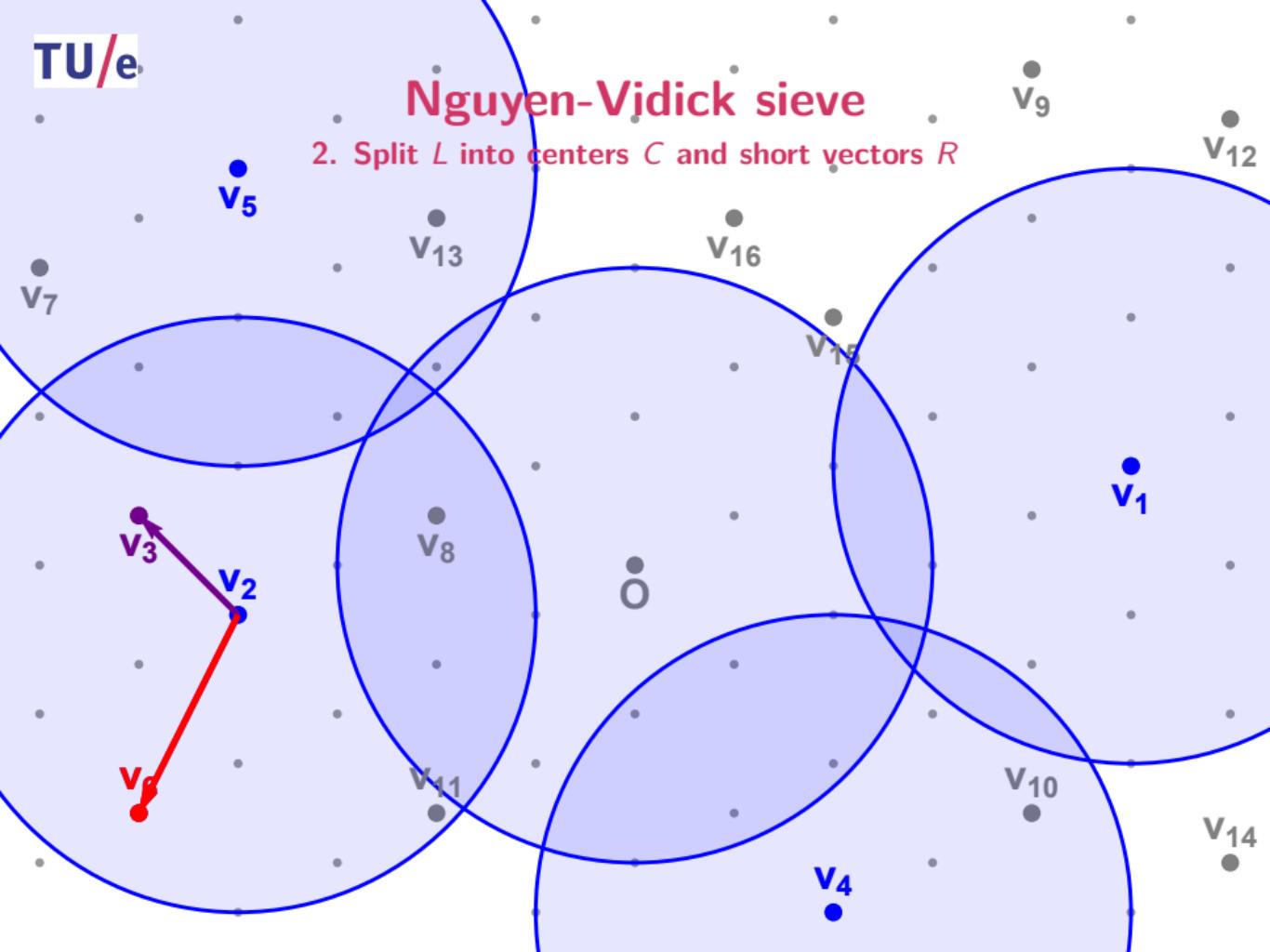
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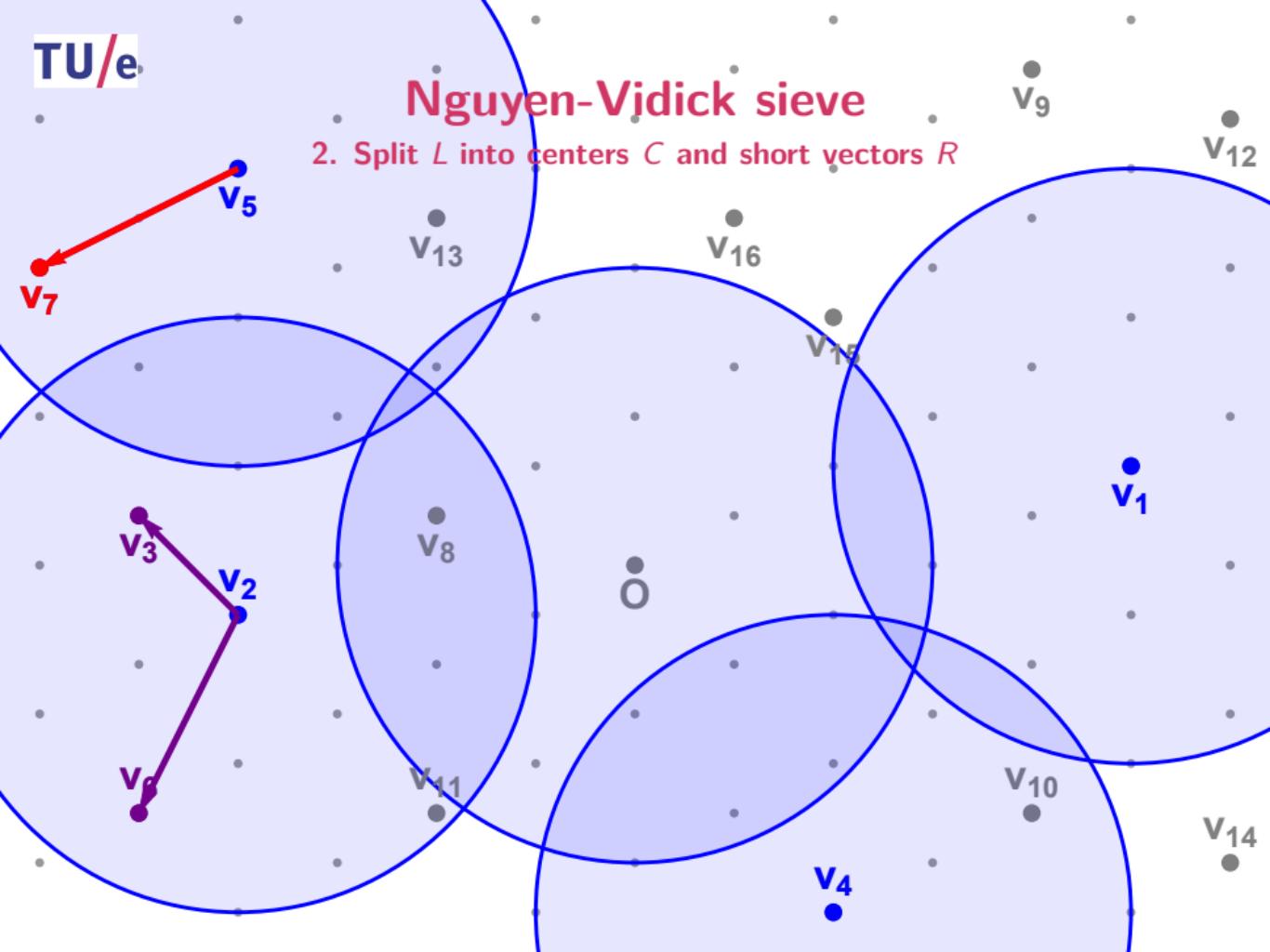
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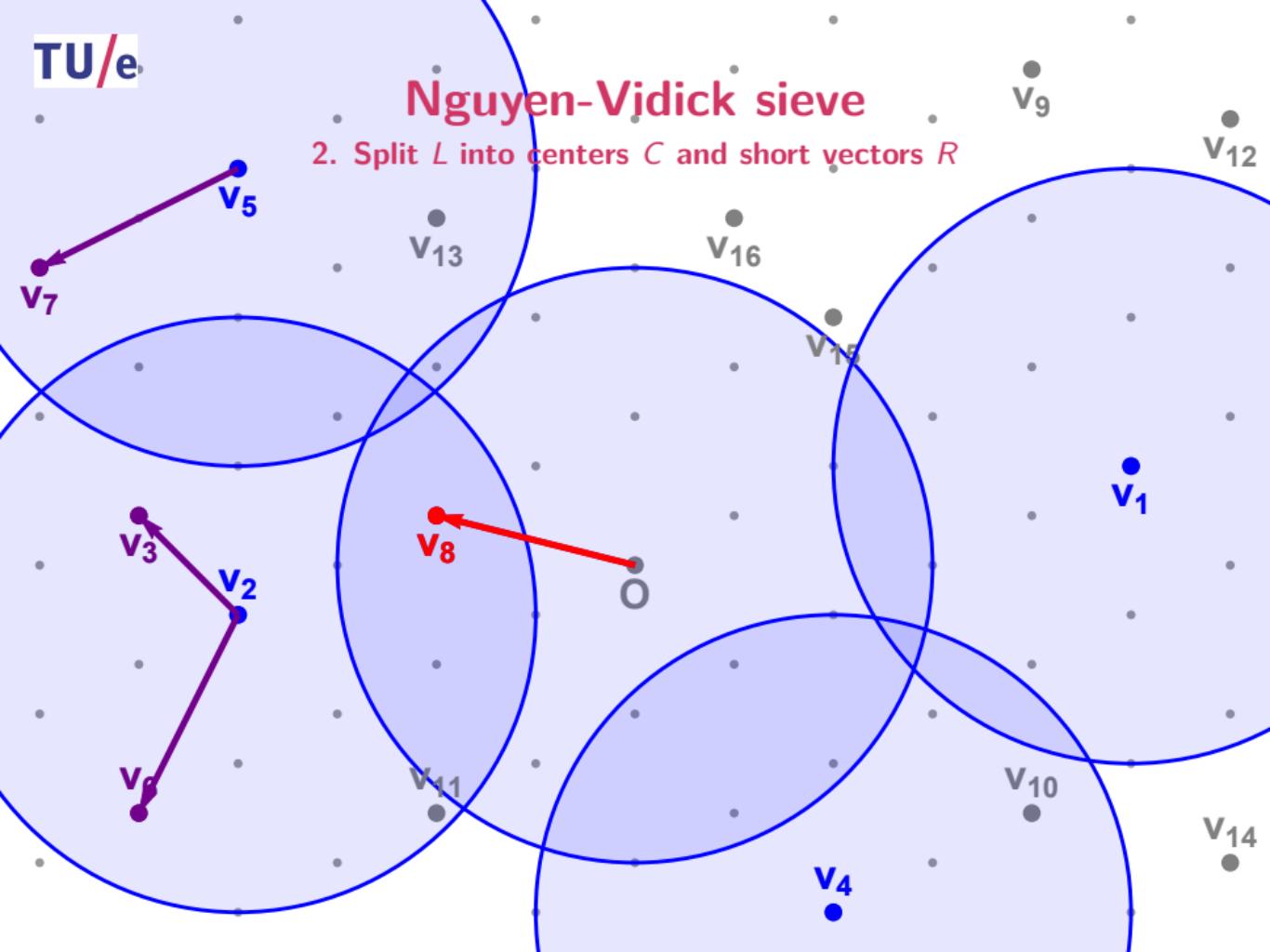
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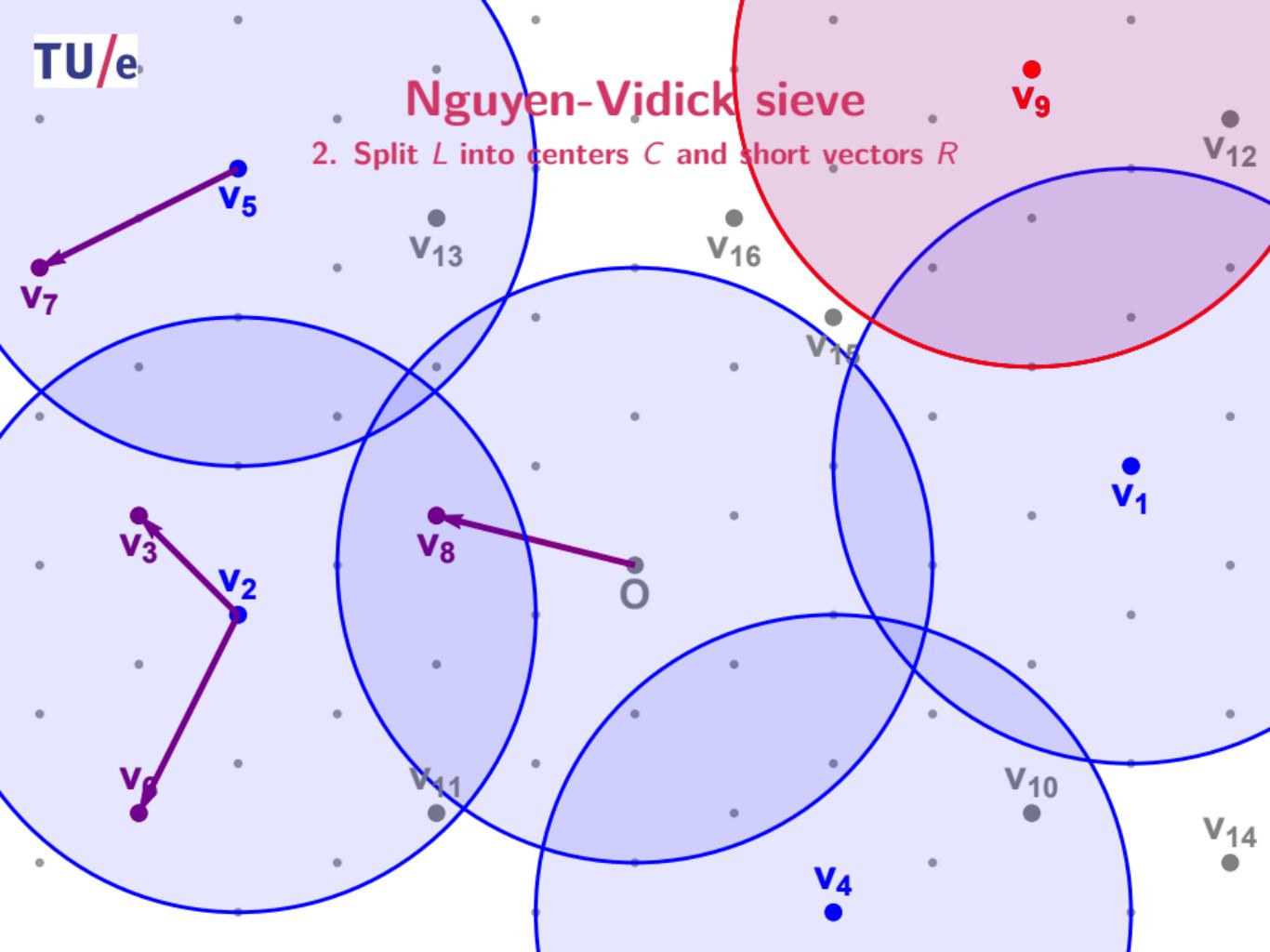
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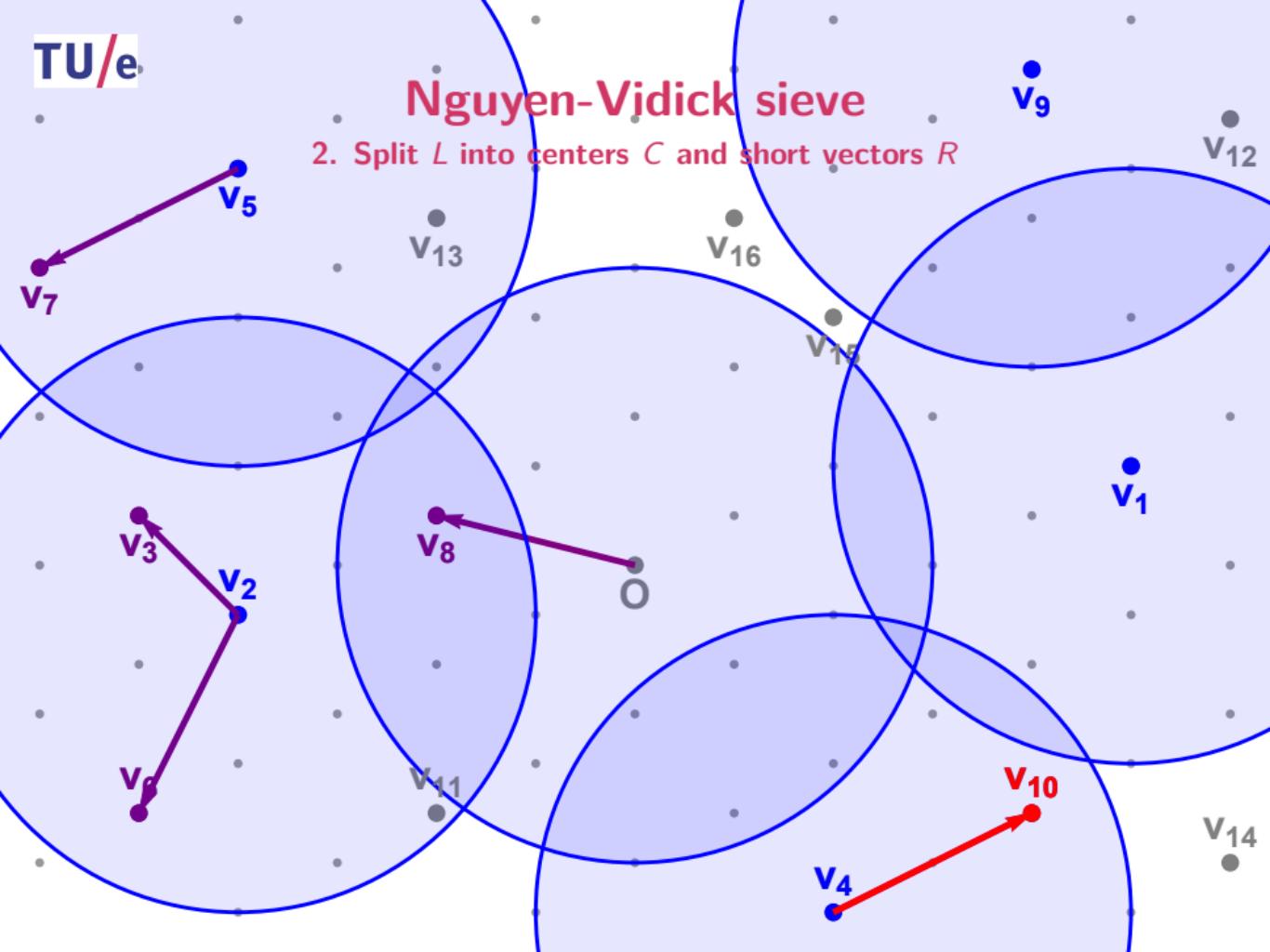
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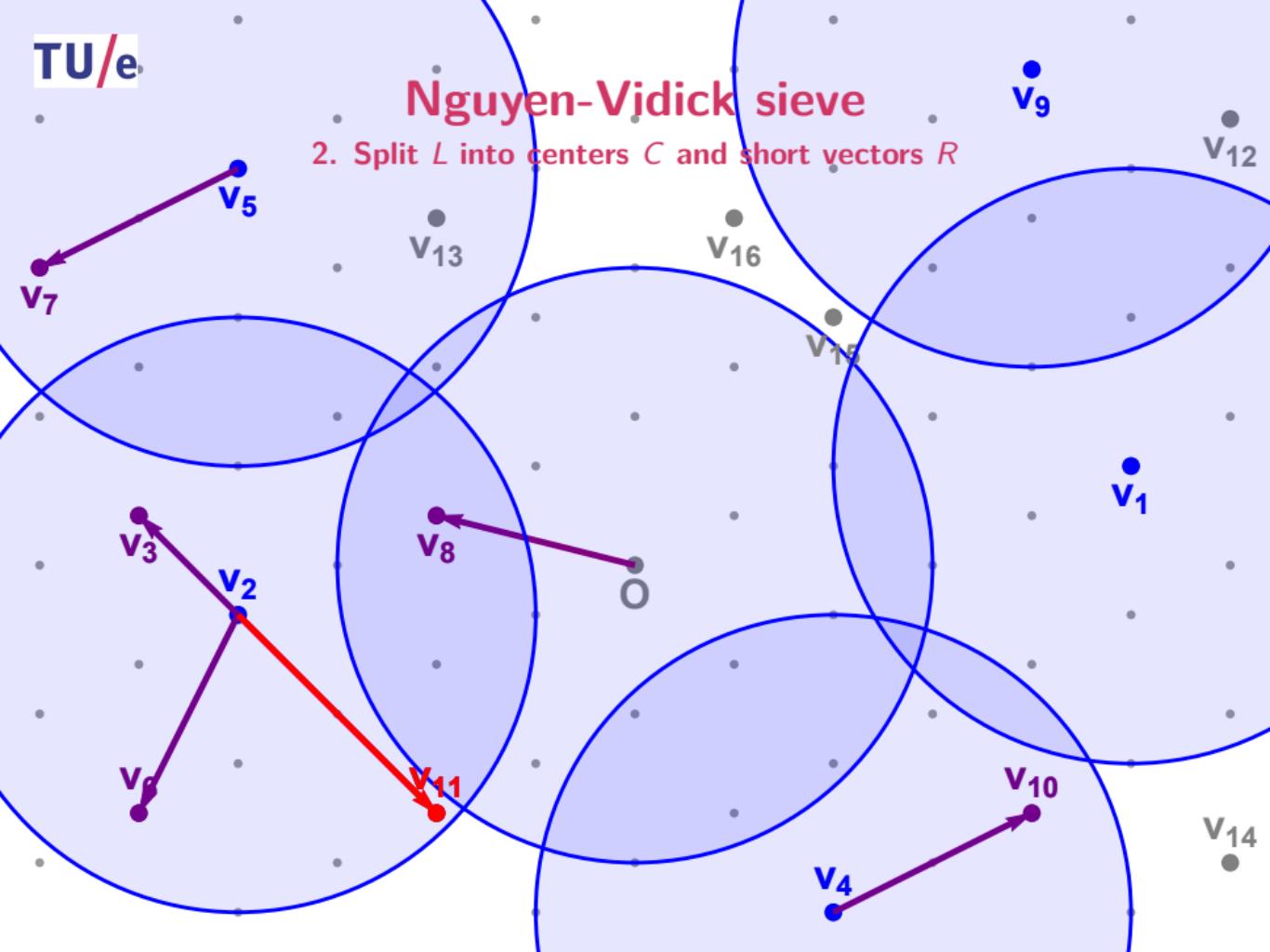
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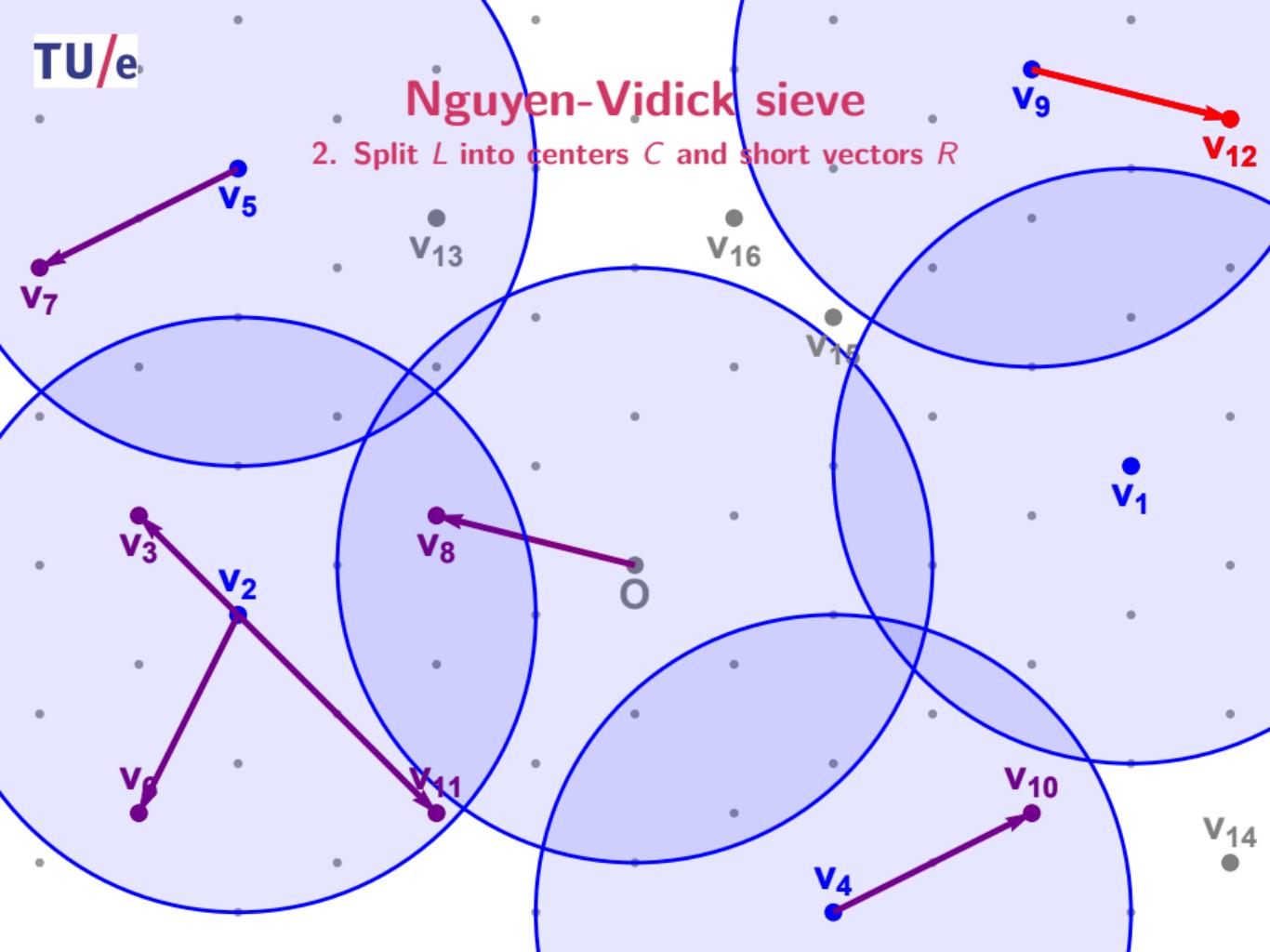
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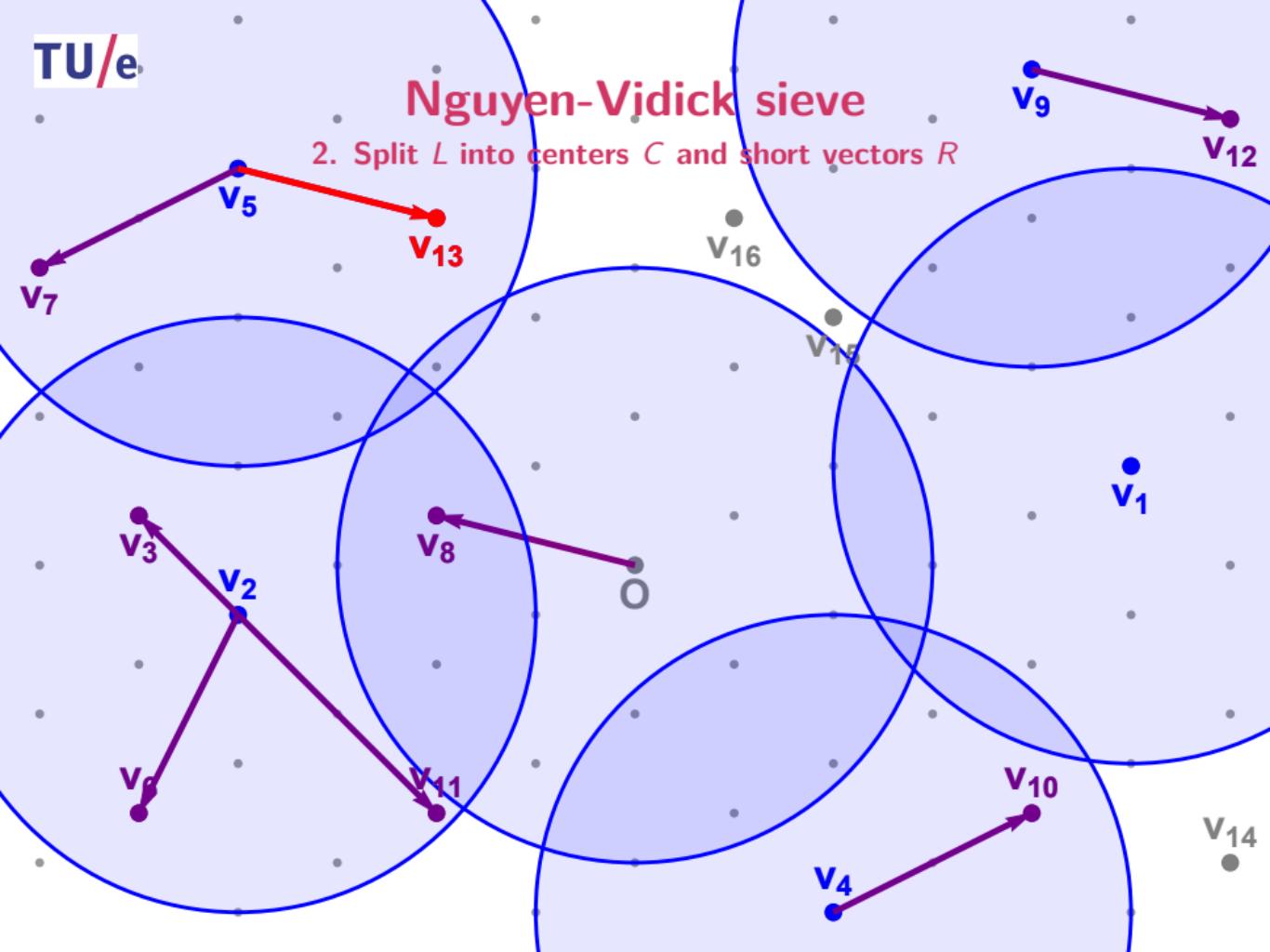
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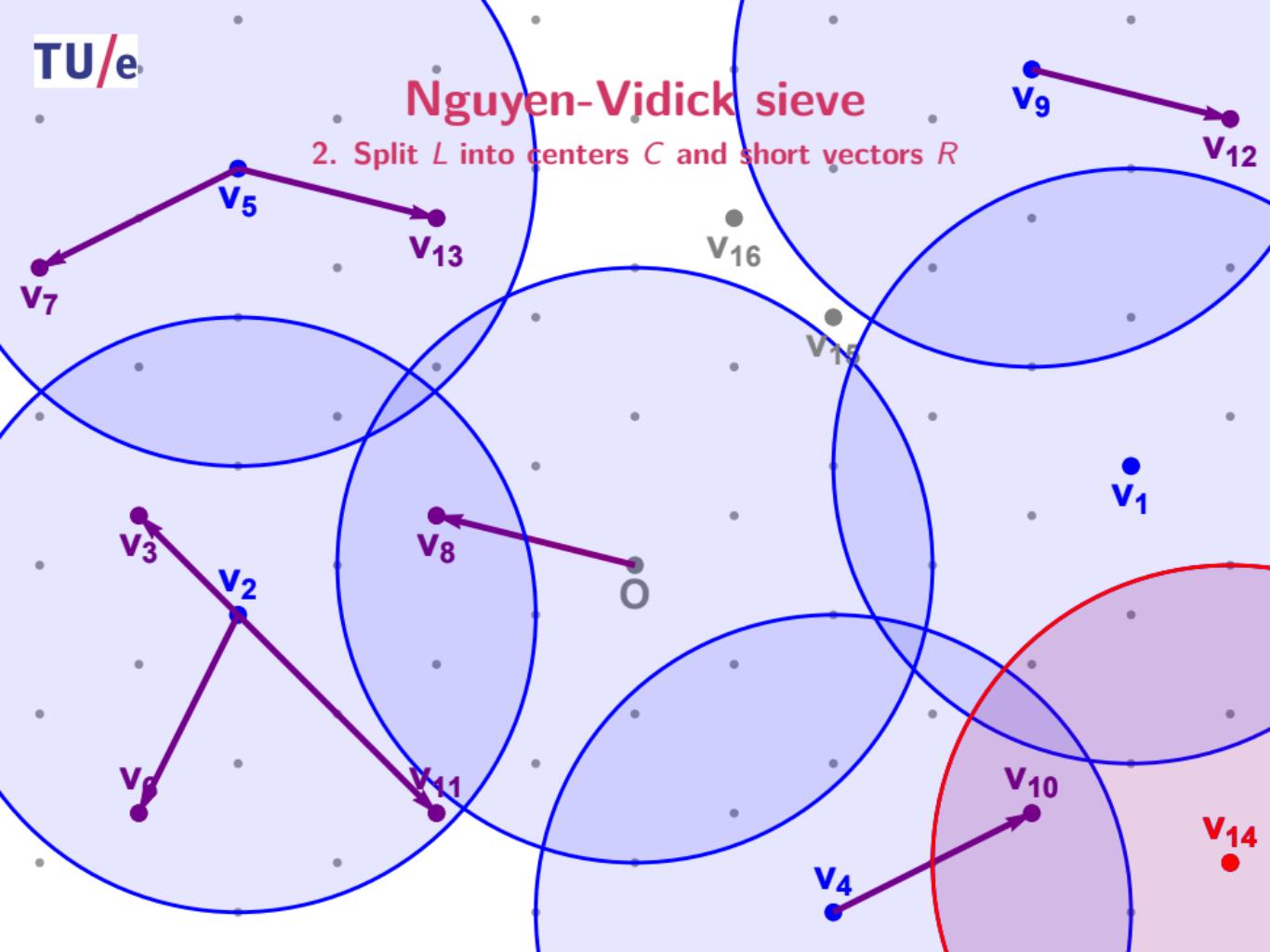
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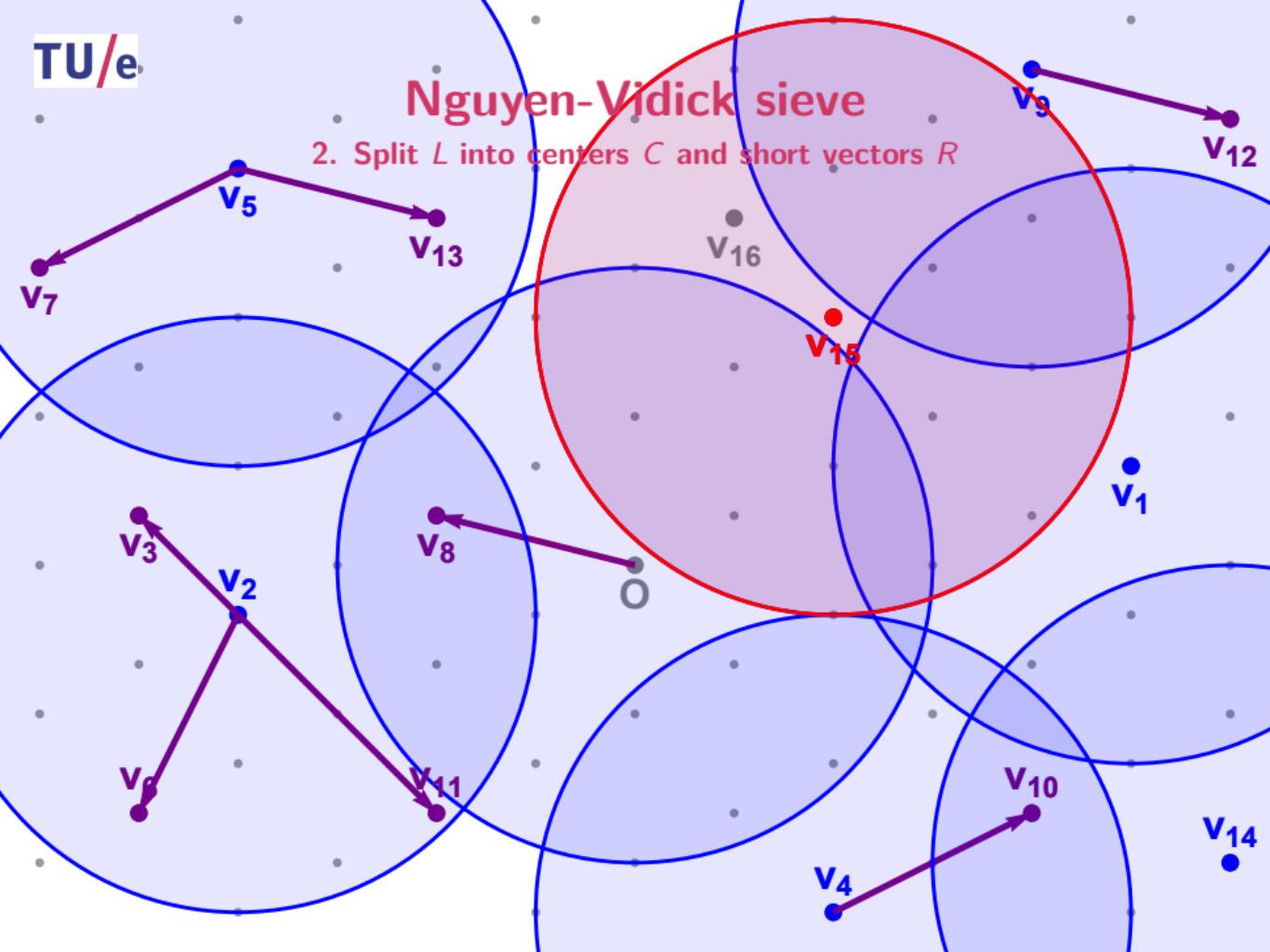
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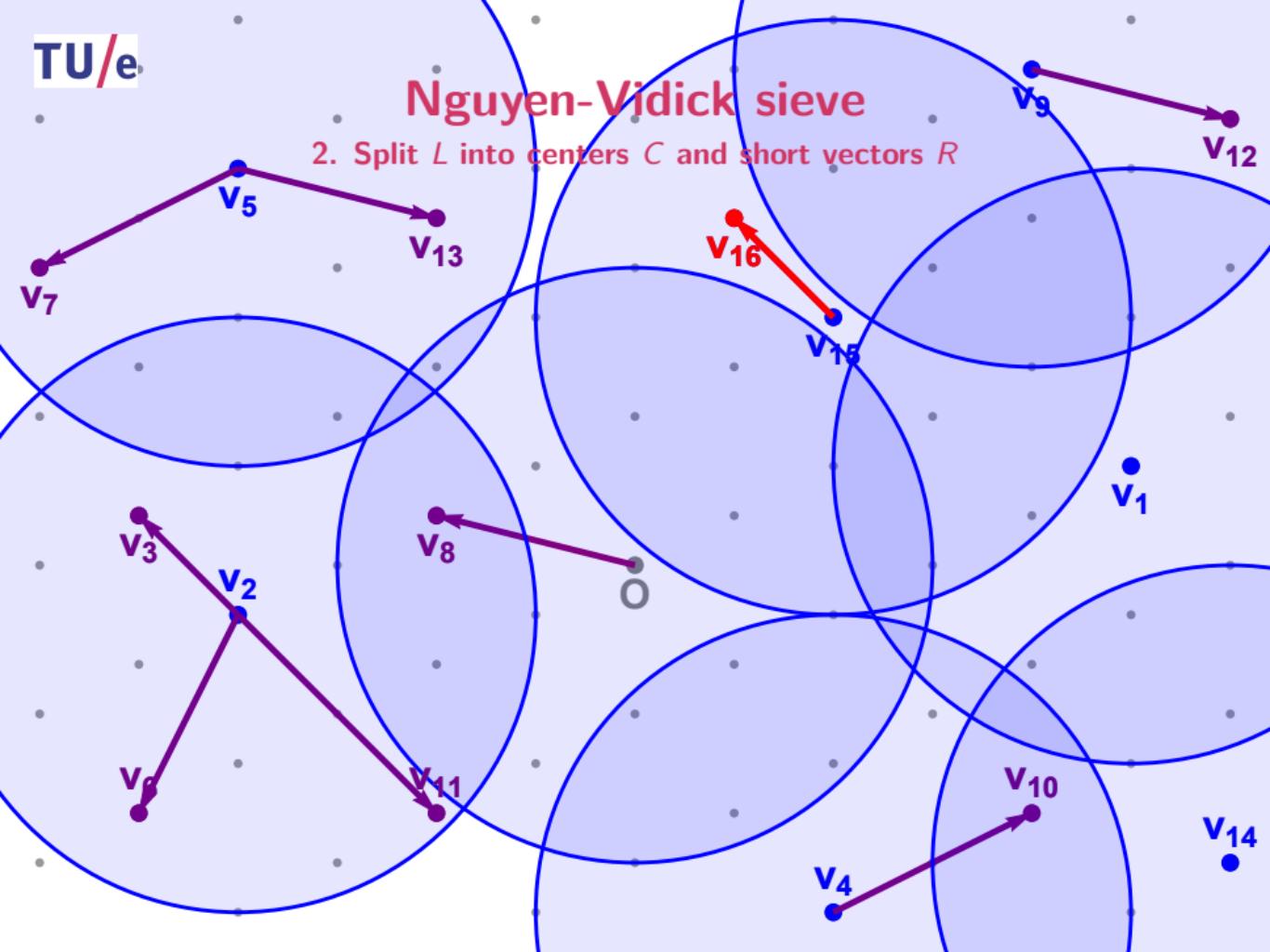
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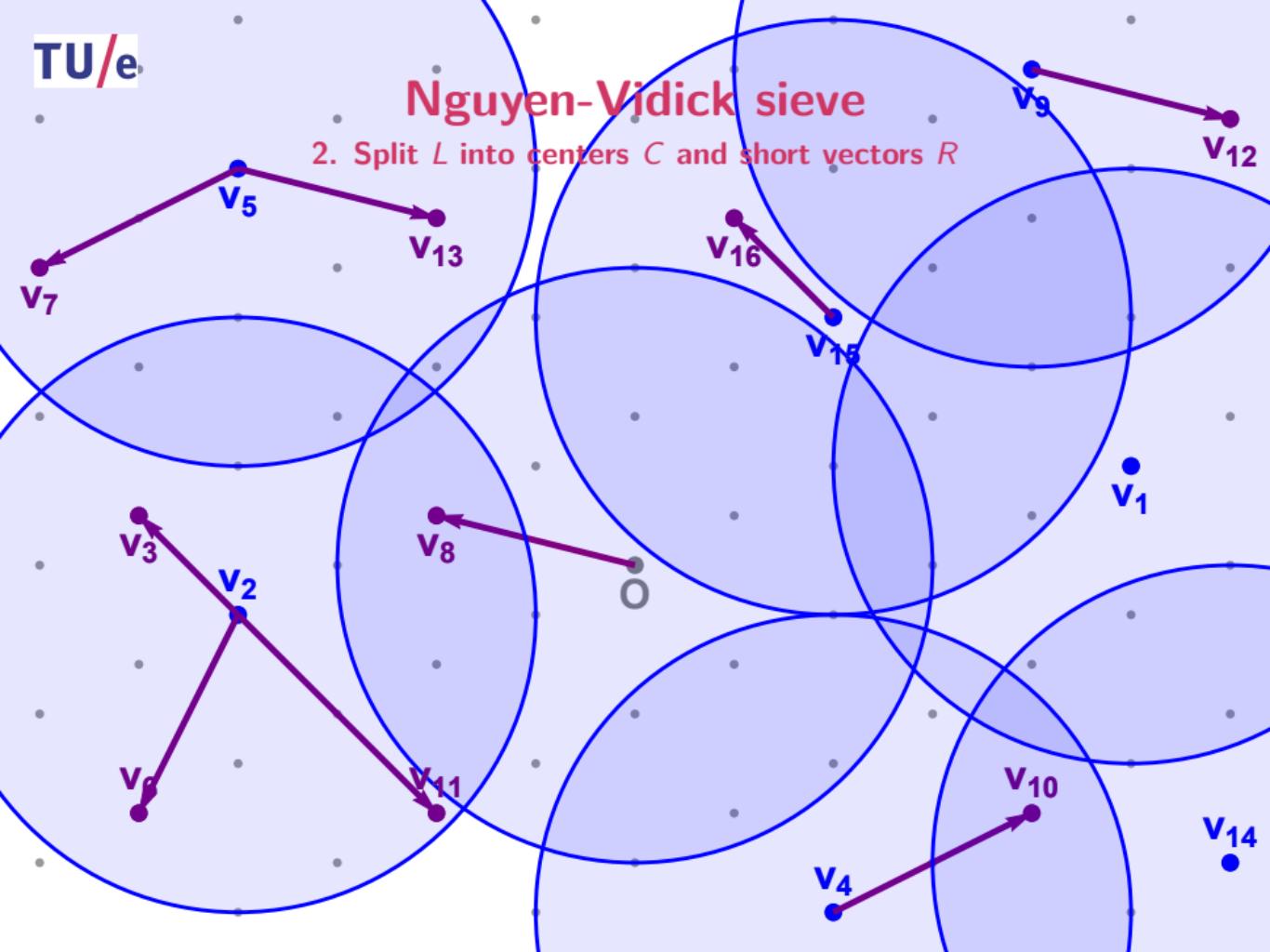
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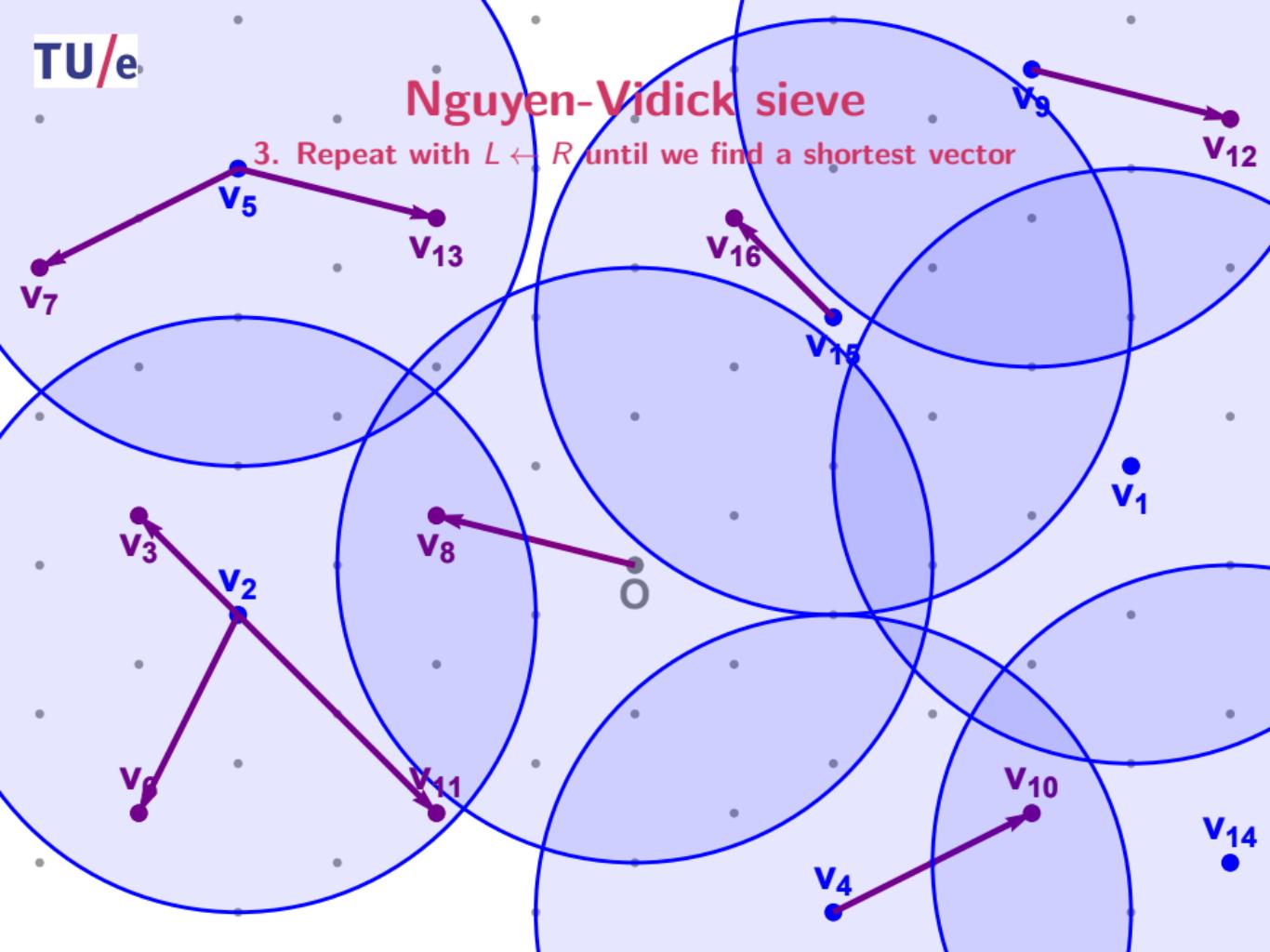
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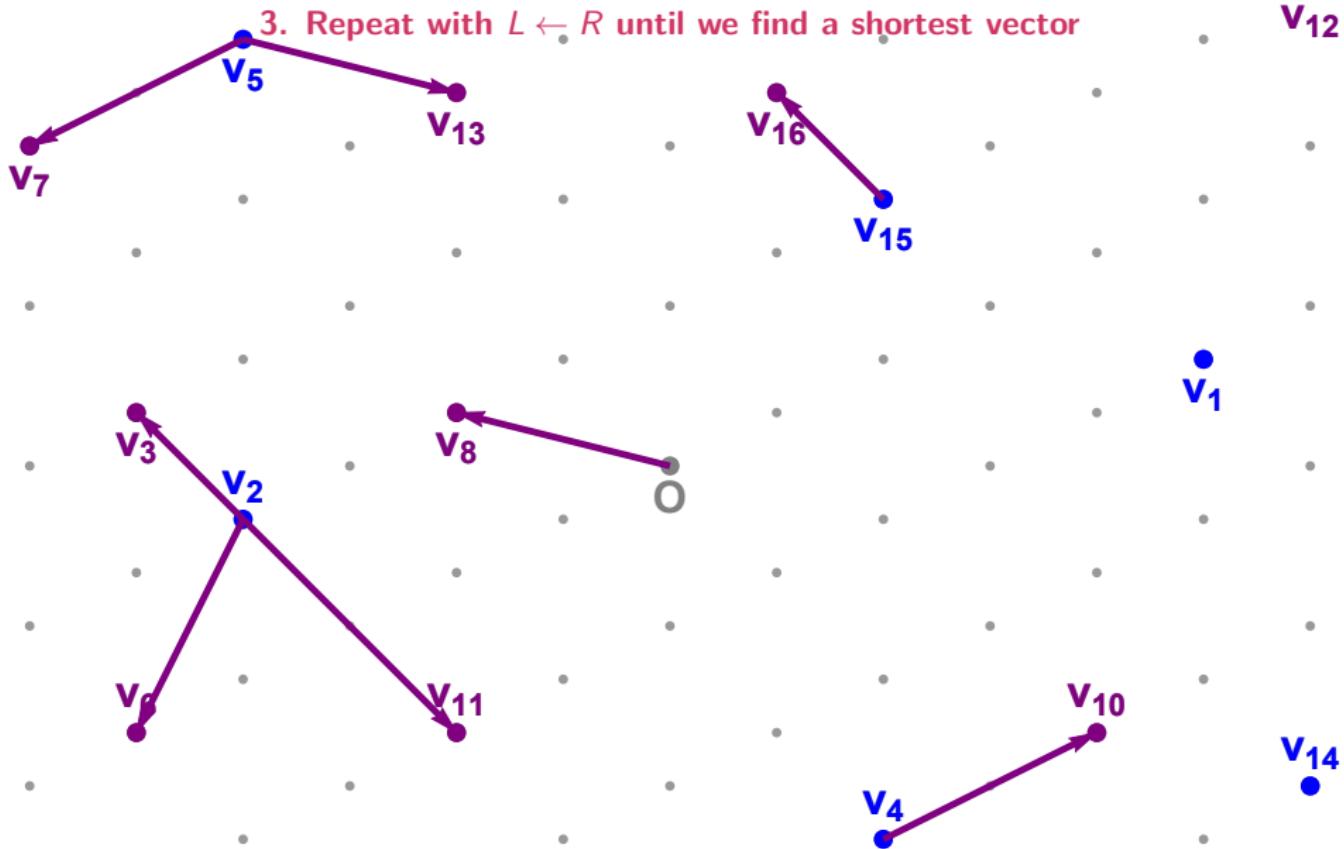
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3. Repeat with $L \leftarrow R$ until we find a shortest vector



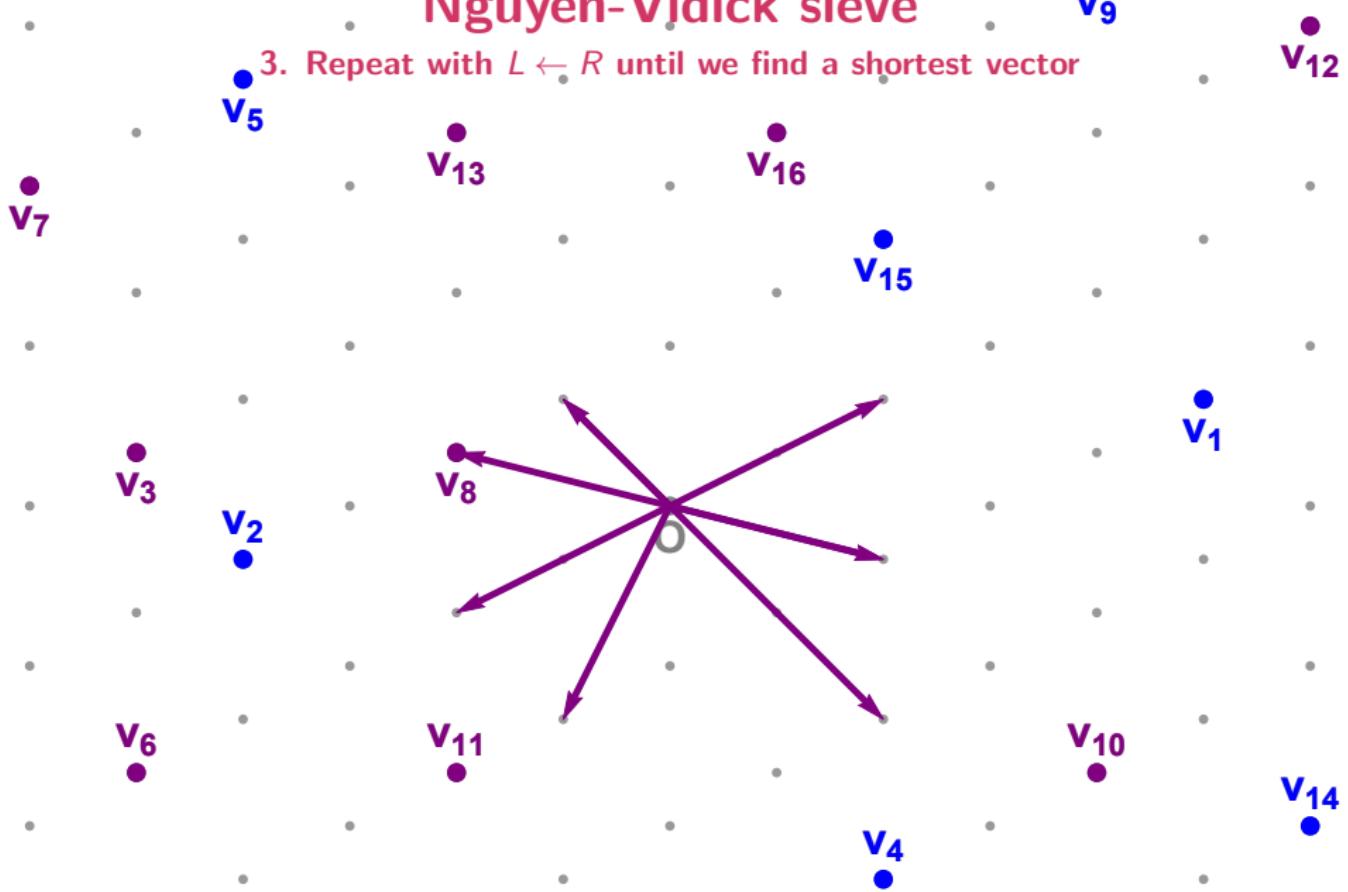
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Nguyen-Vidick sieve

Overview



Nguyen-Vidick sieve

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The Nguyen-Vidick sieve runs in time $(4/3)^n$ and space $\sqrt{4/3}^n$.



Nguyen-Vidick sieve

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The Nguyen-Vidick sieve runs in time $2^{0.415n}$ and space $2^{0.208n}$.



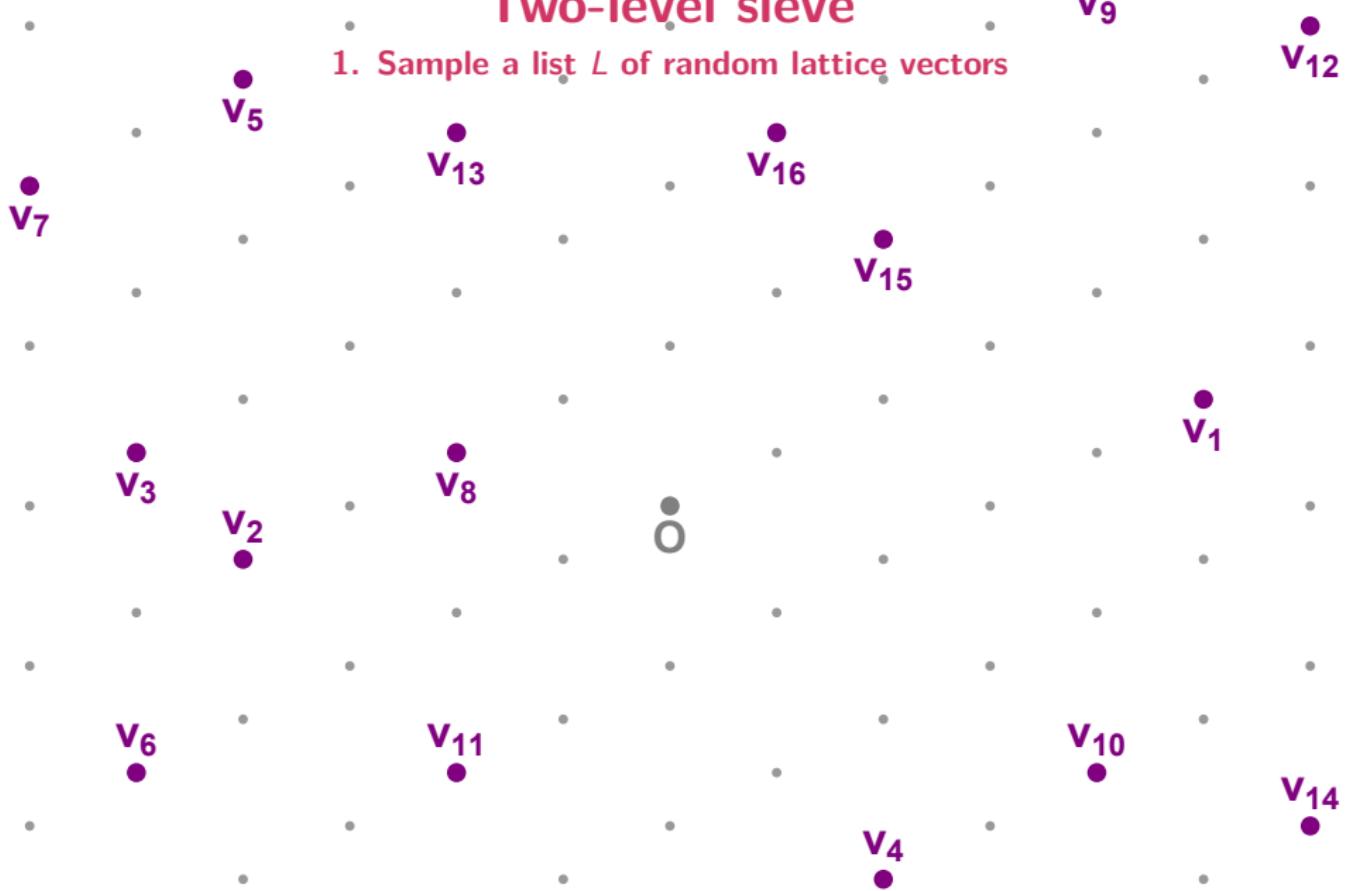
Two-level sieve

1. Sample a list L of random lattice vectors



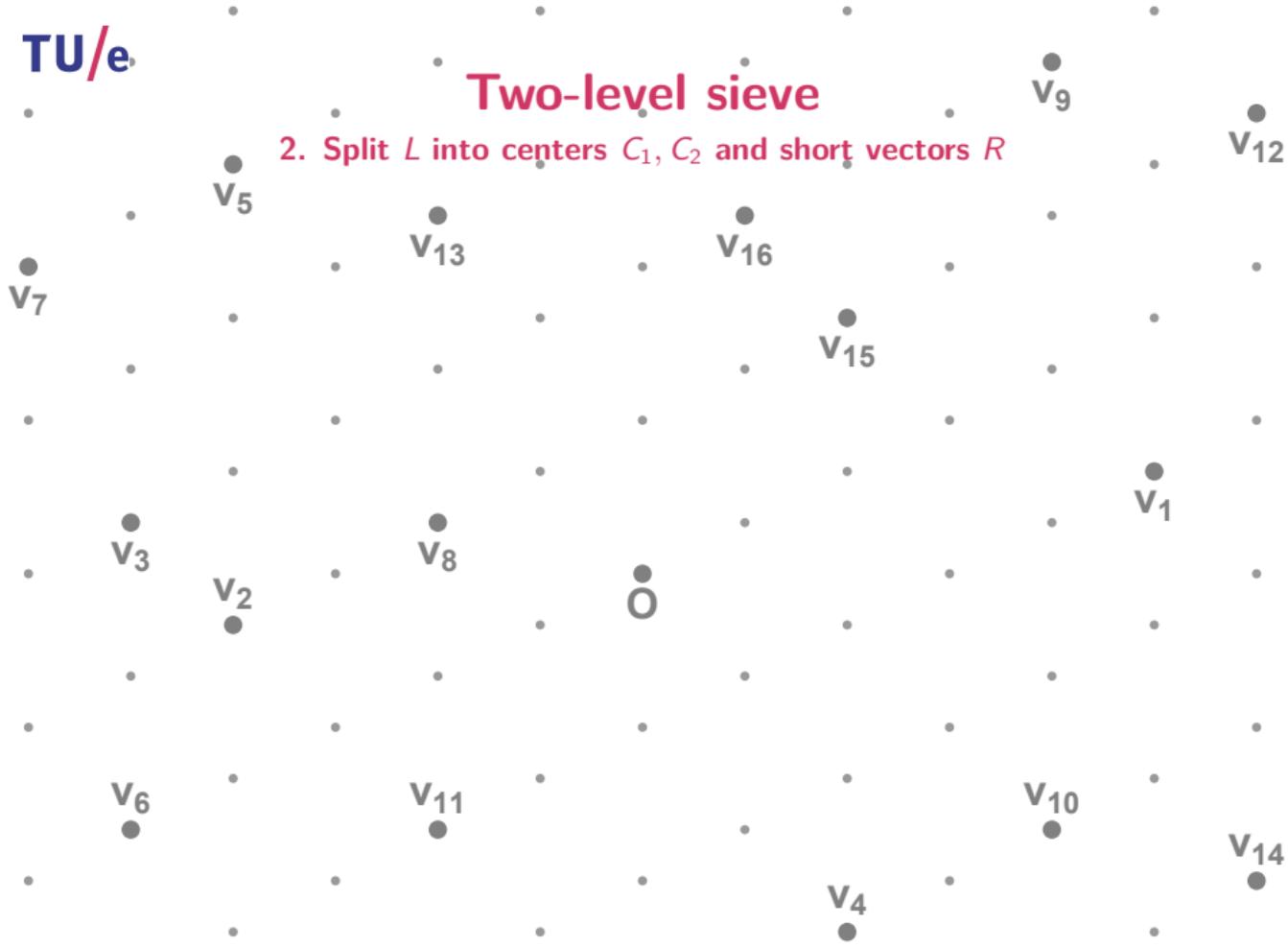
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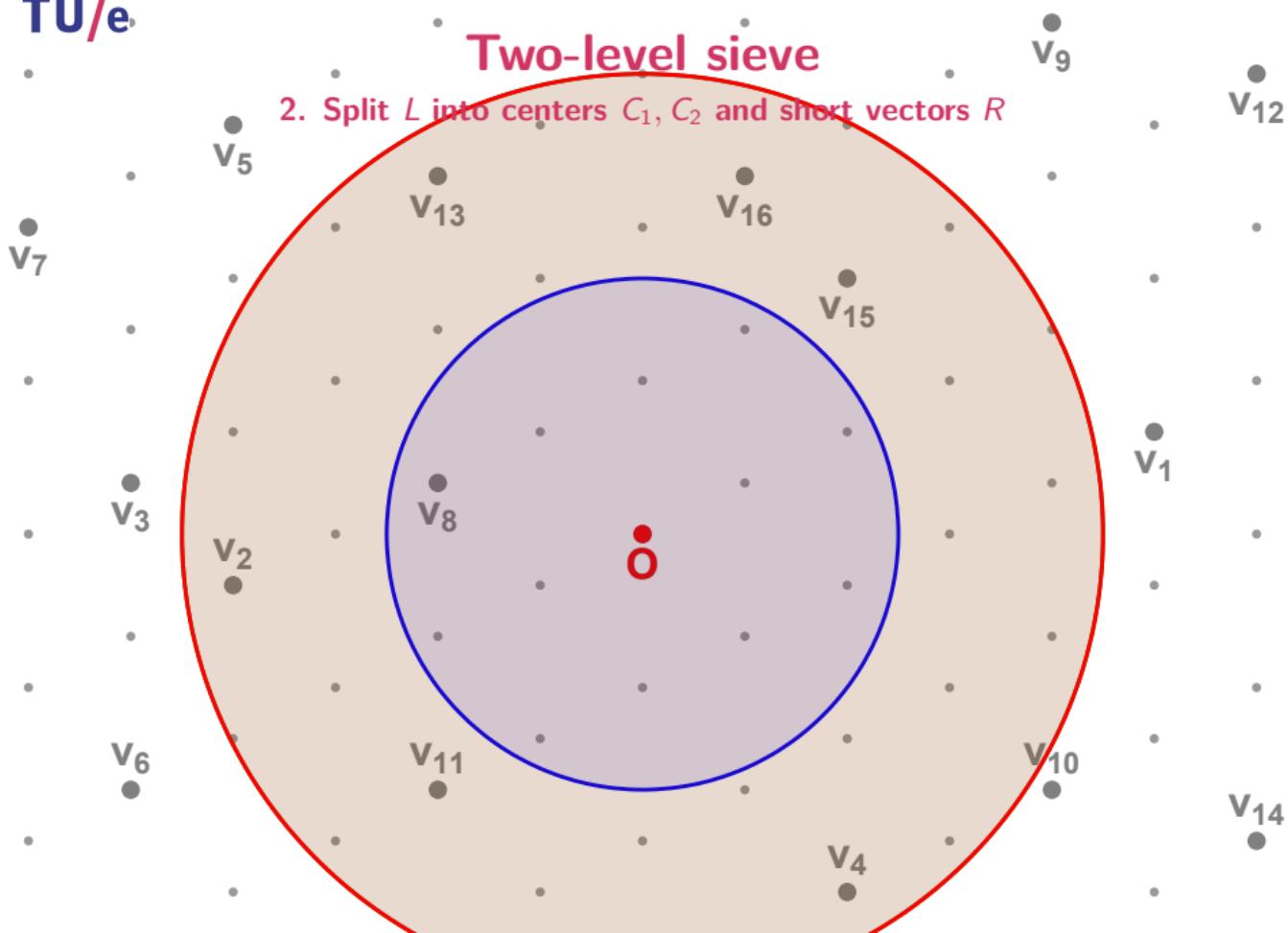
Two-level sieve

2. Split L into centers C_1, C_2 and short vectors R



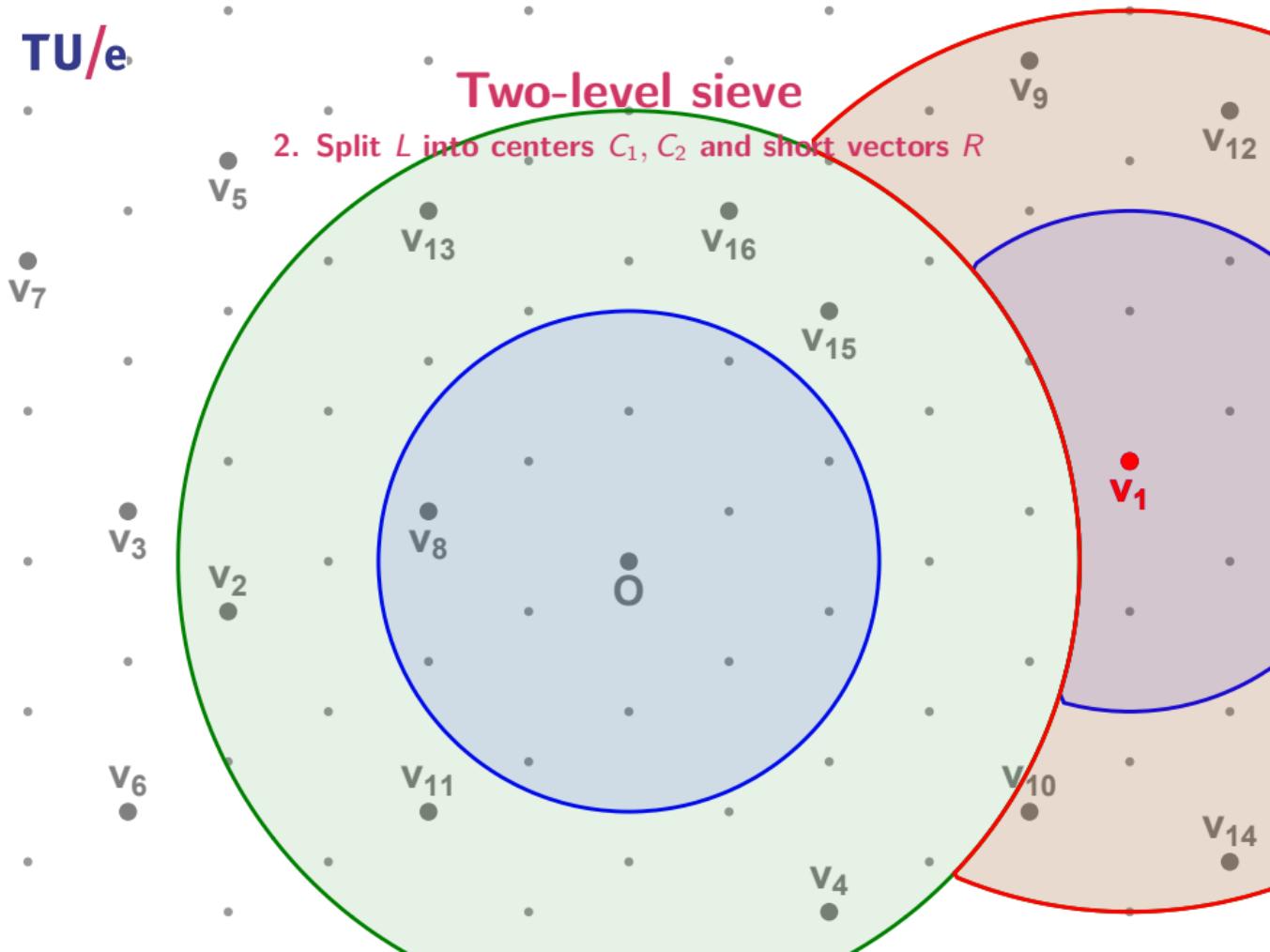
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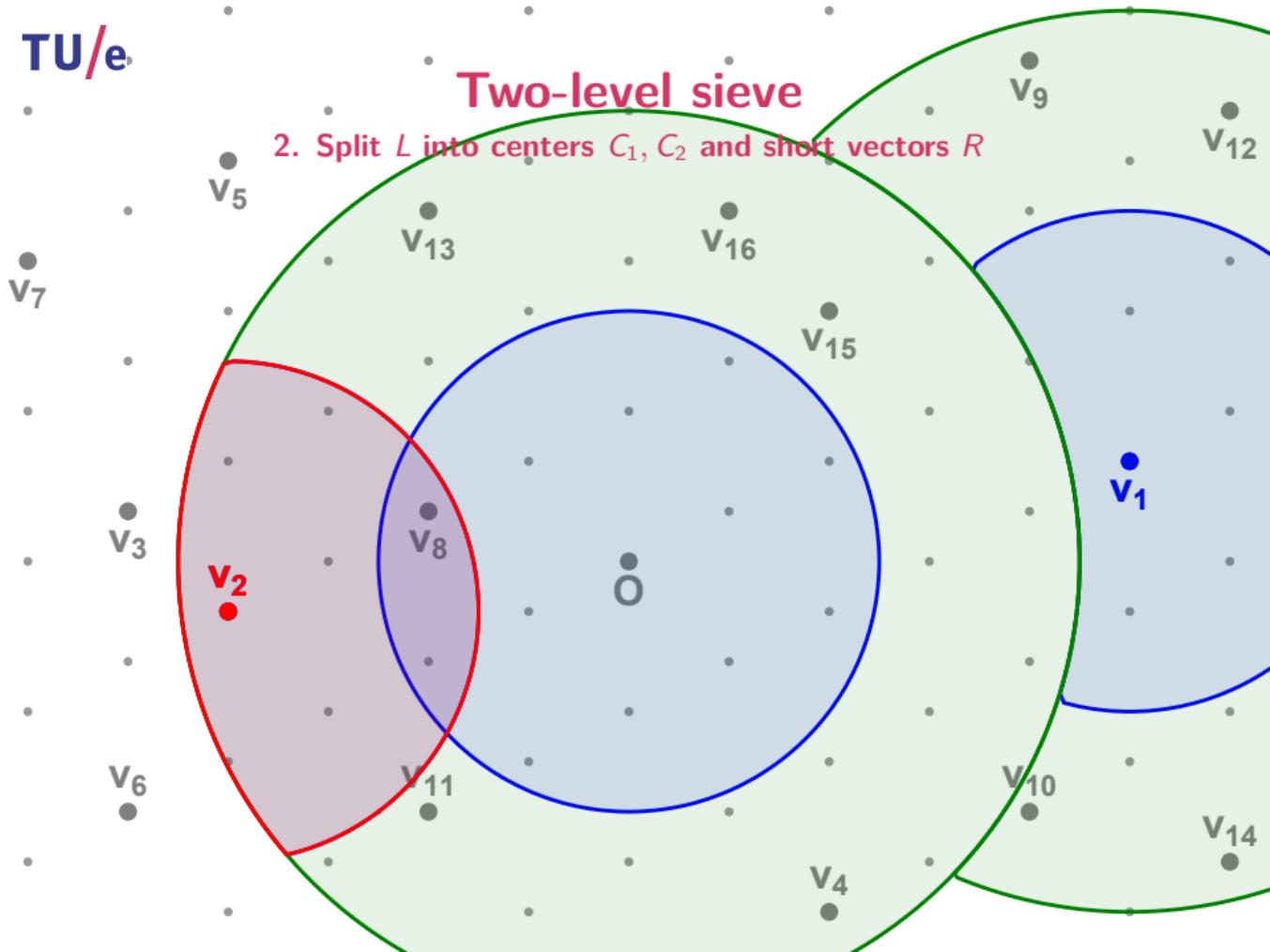
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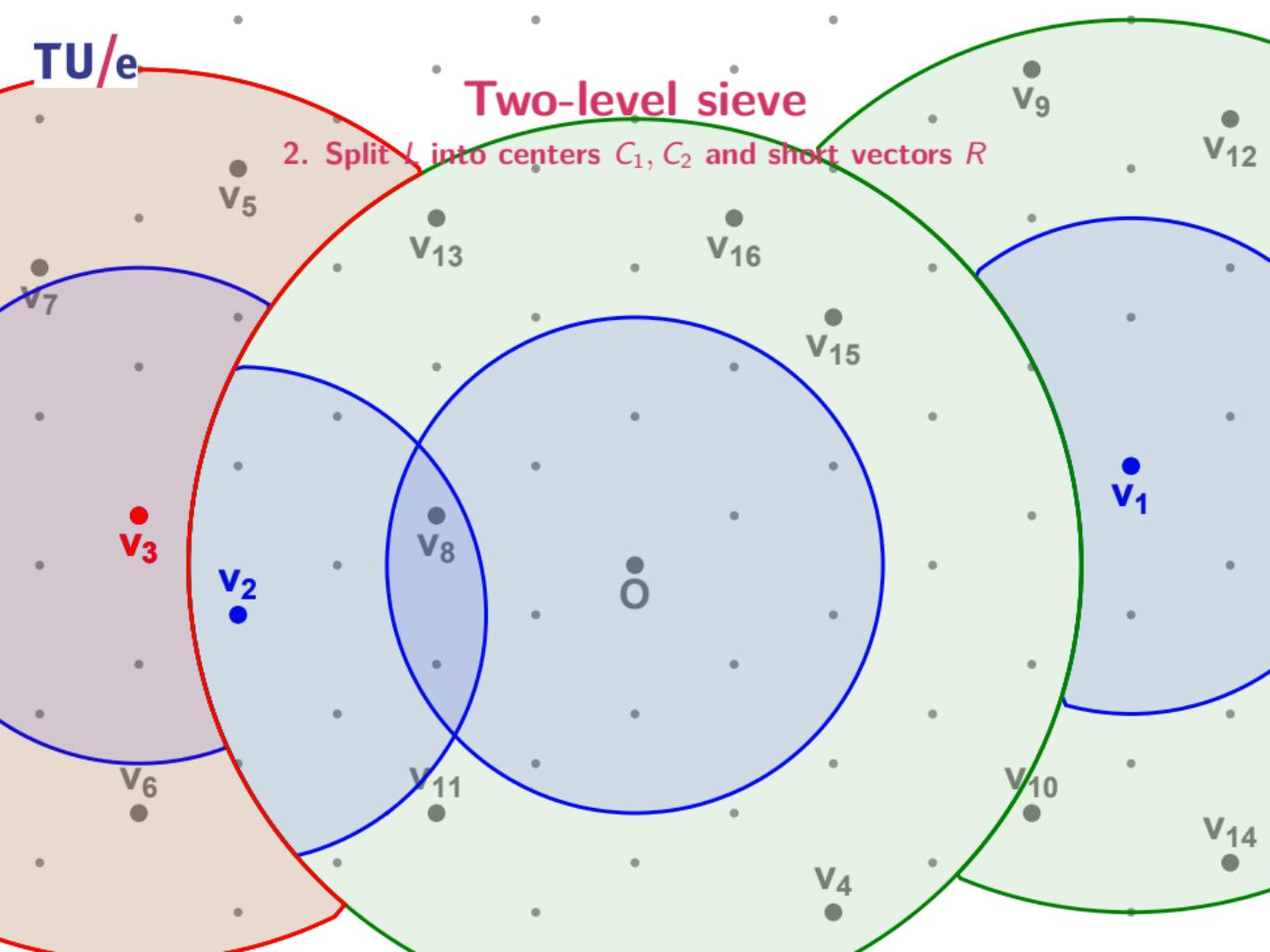
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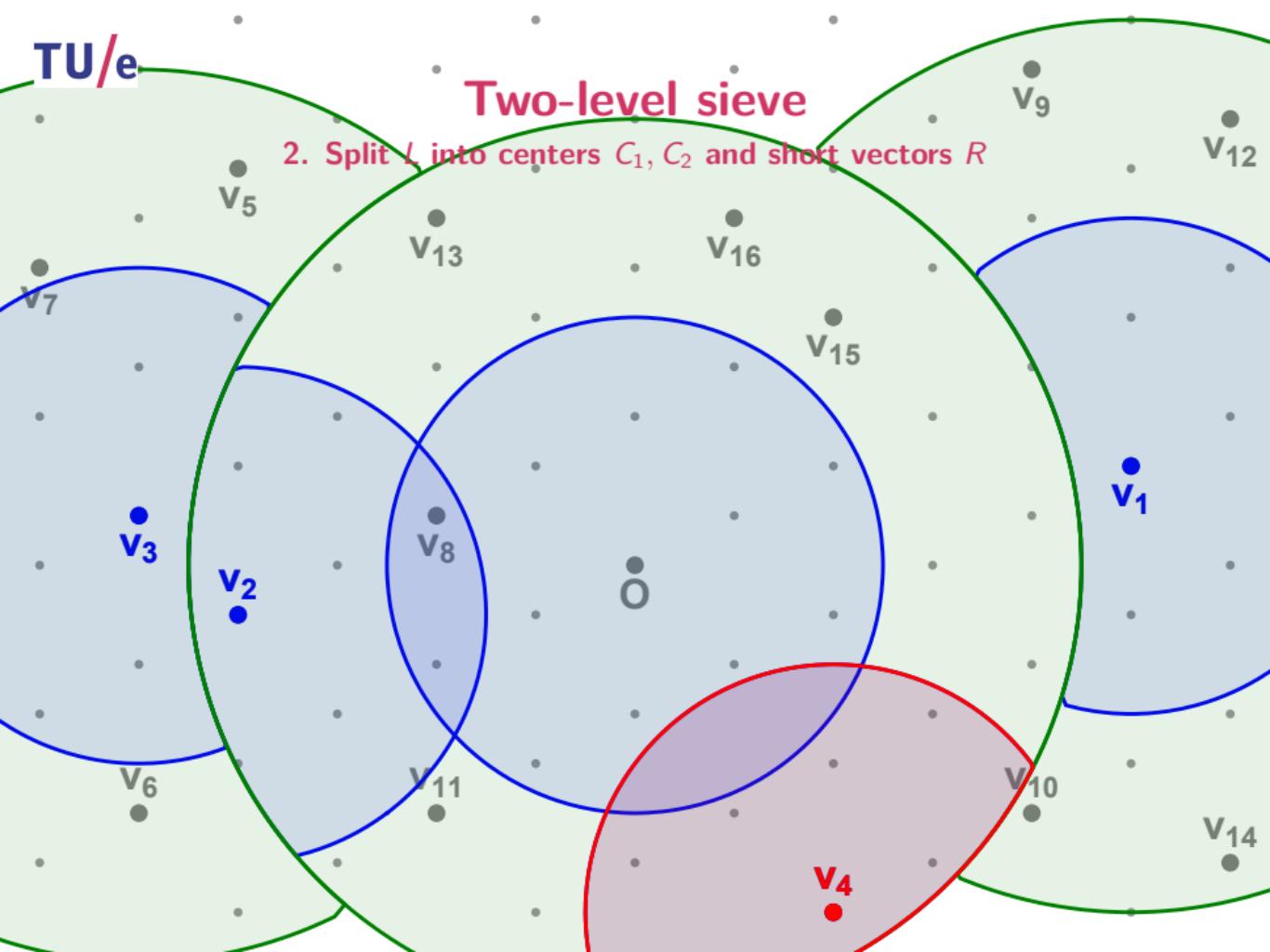
Two-level sieve

2. Split \mathcal{V} into centers C_1, C_2 and short vectors R



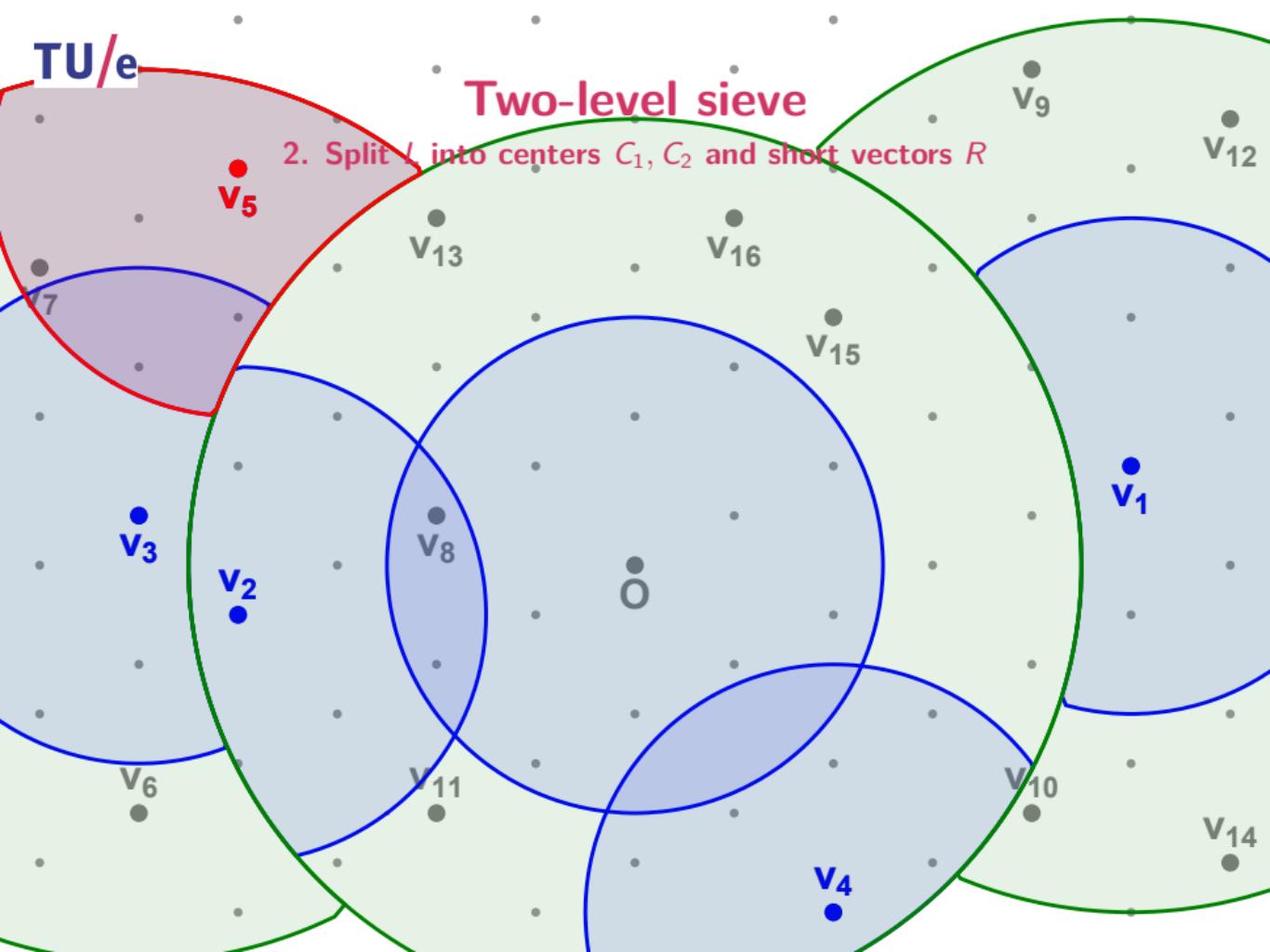
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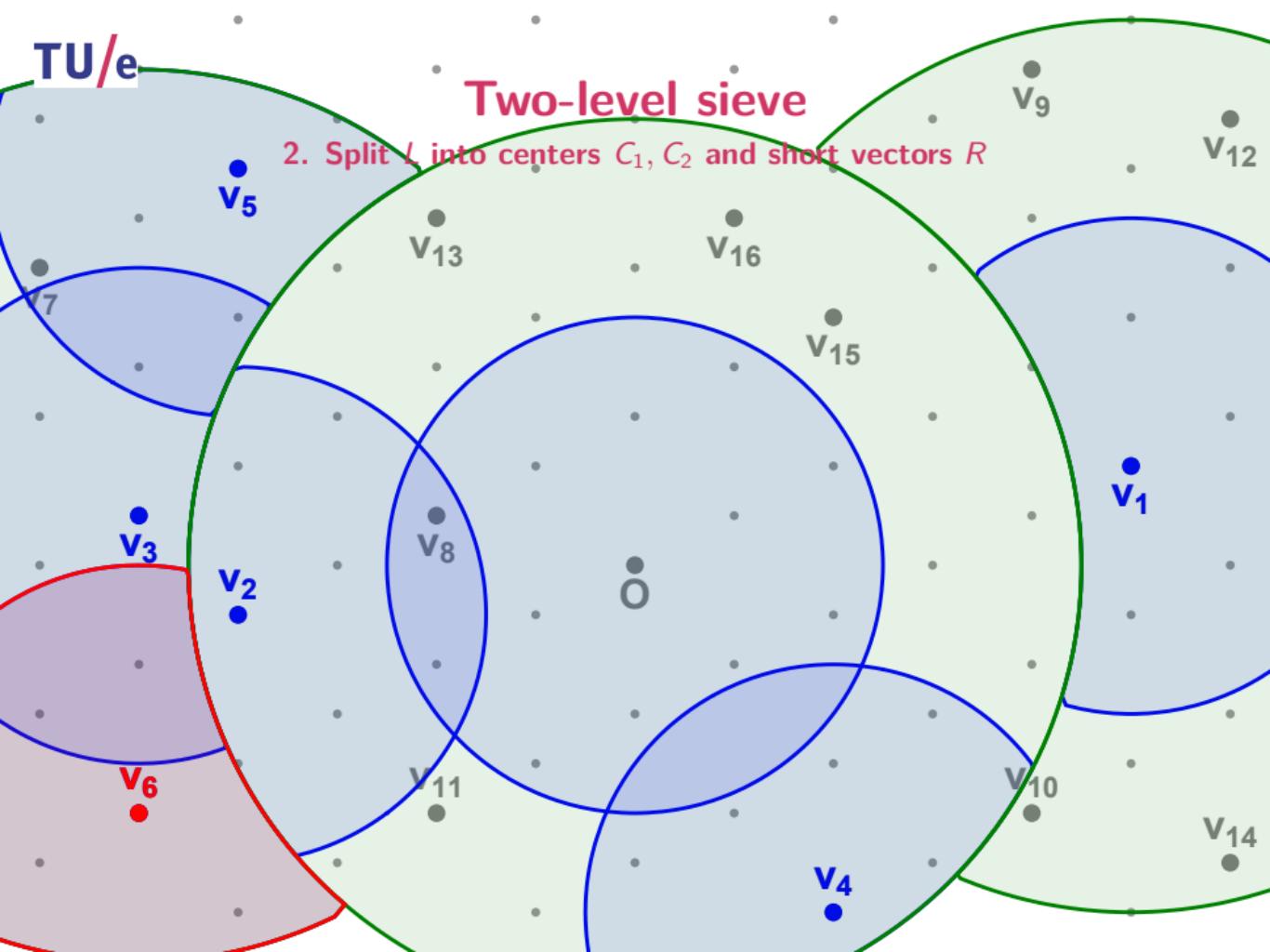
Two-level sieve

2. Split V into centers C_1, C_2 and short vectors R



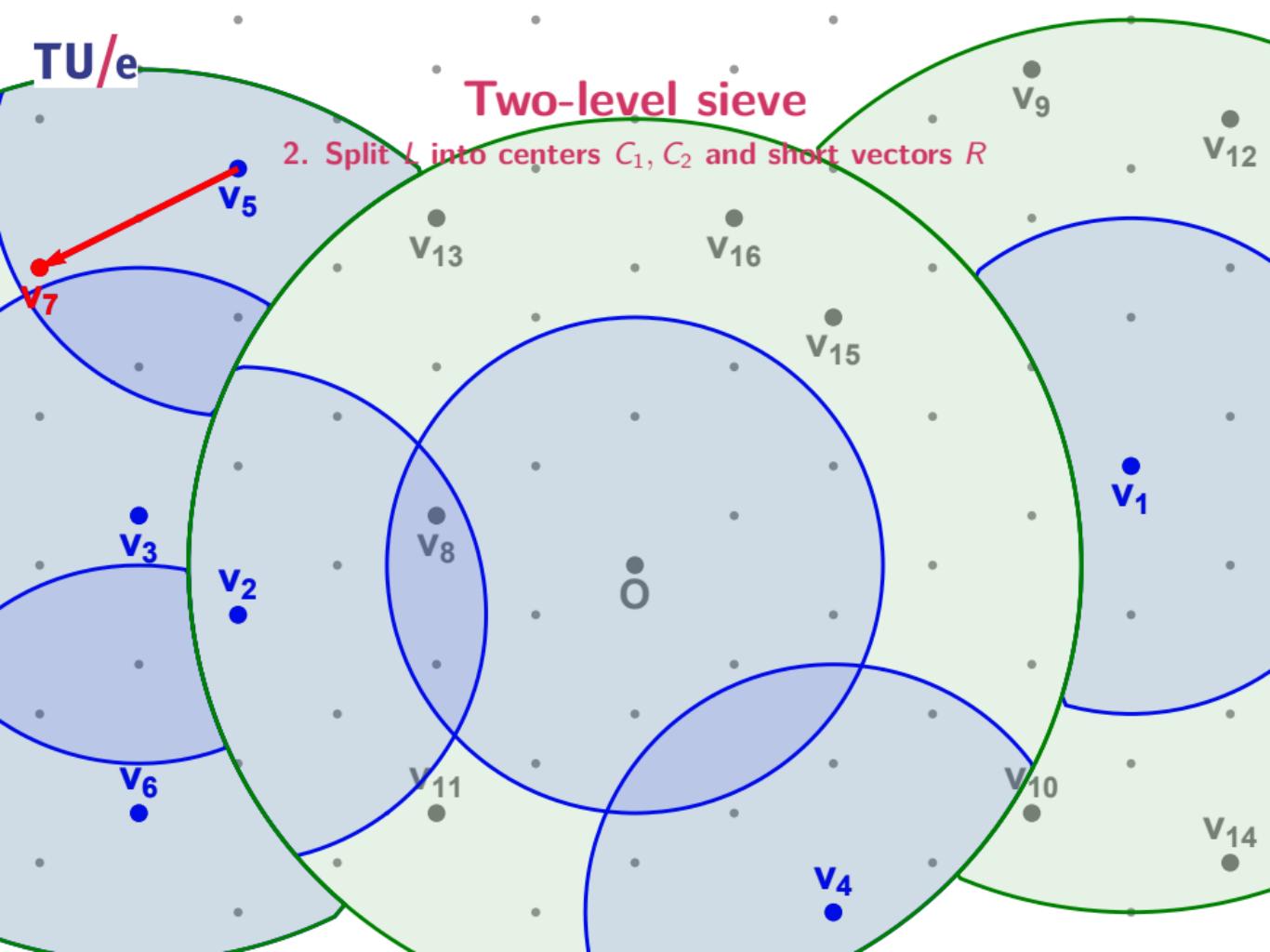
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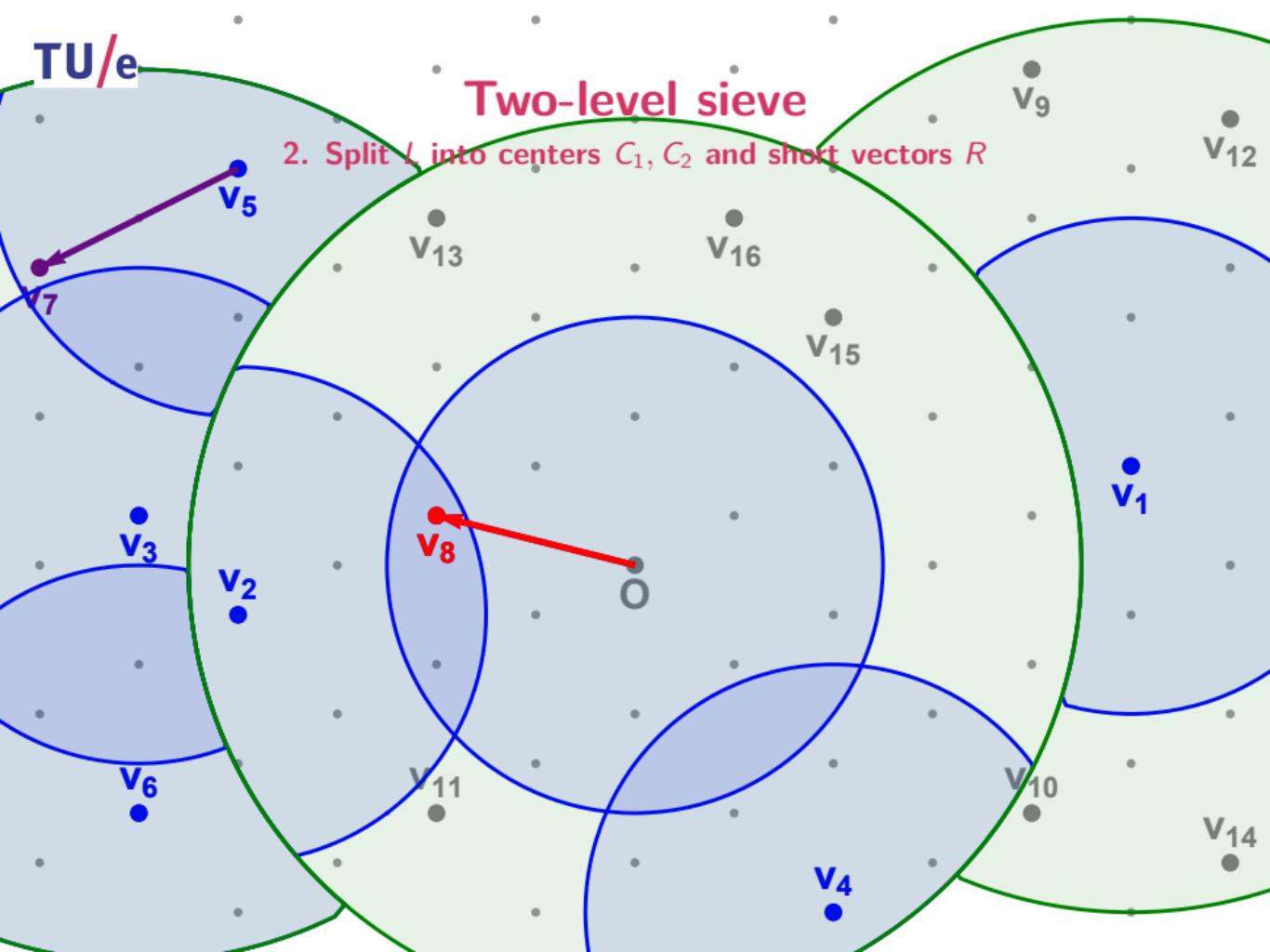
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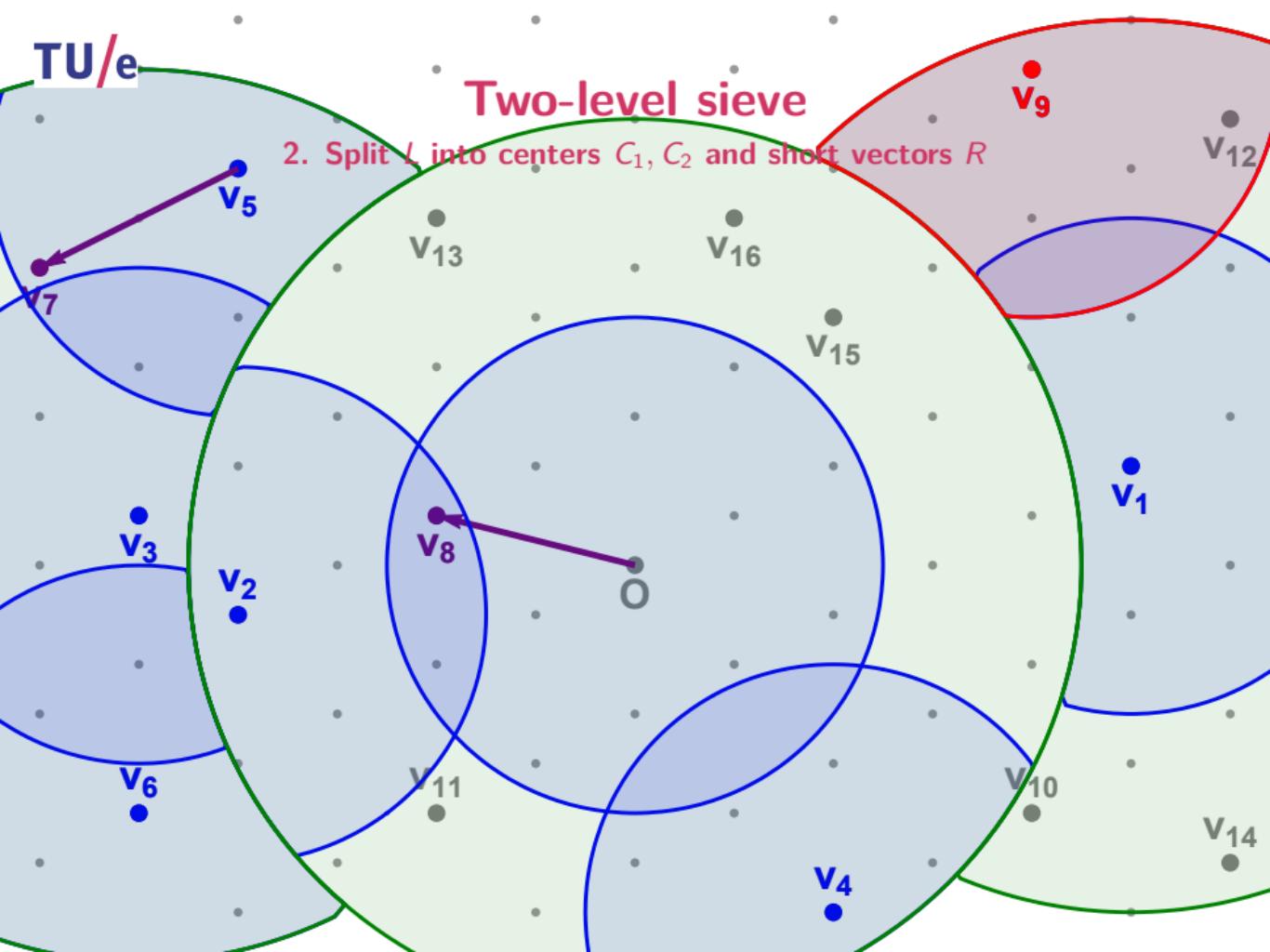
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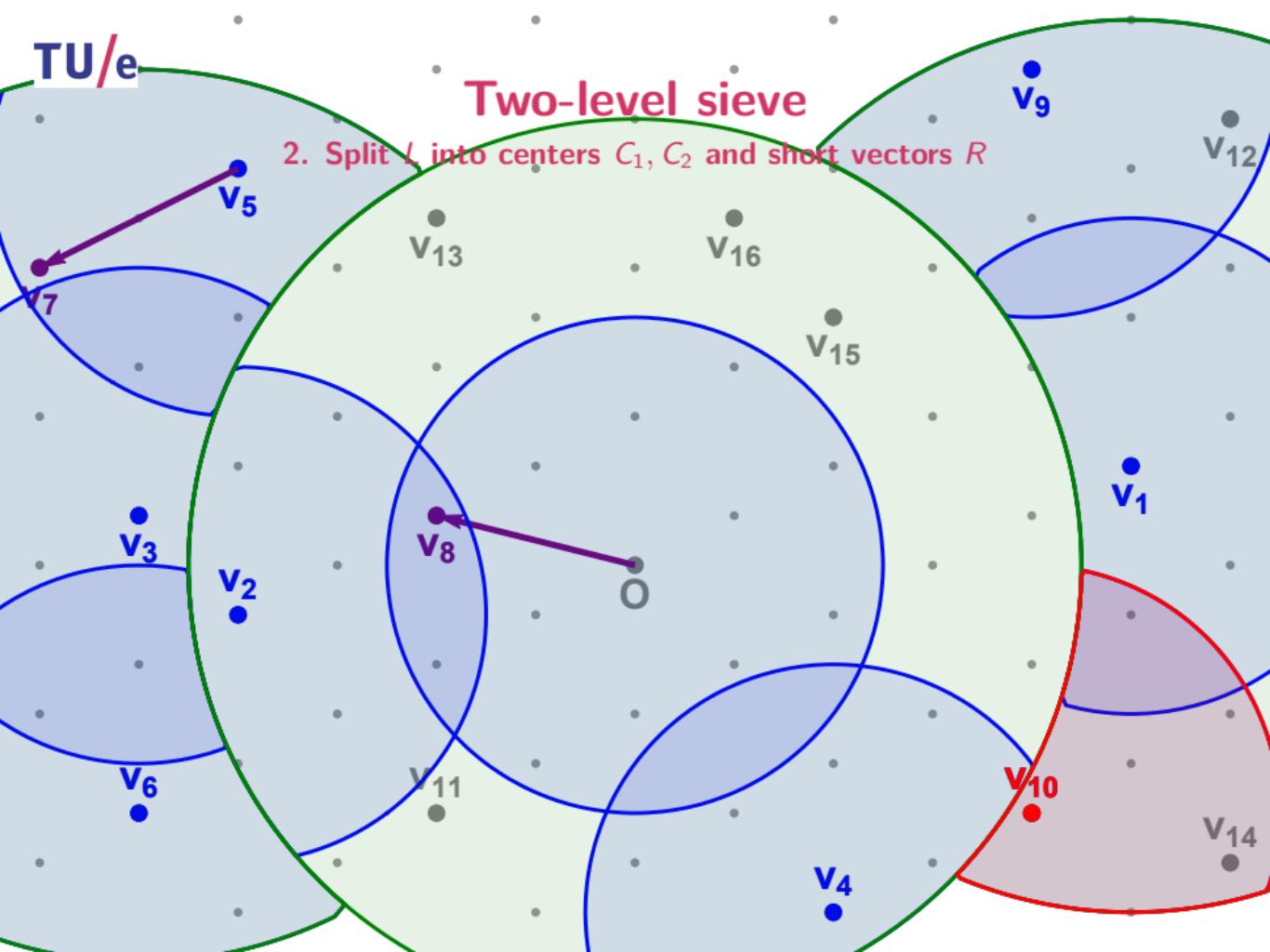
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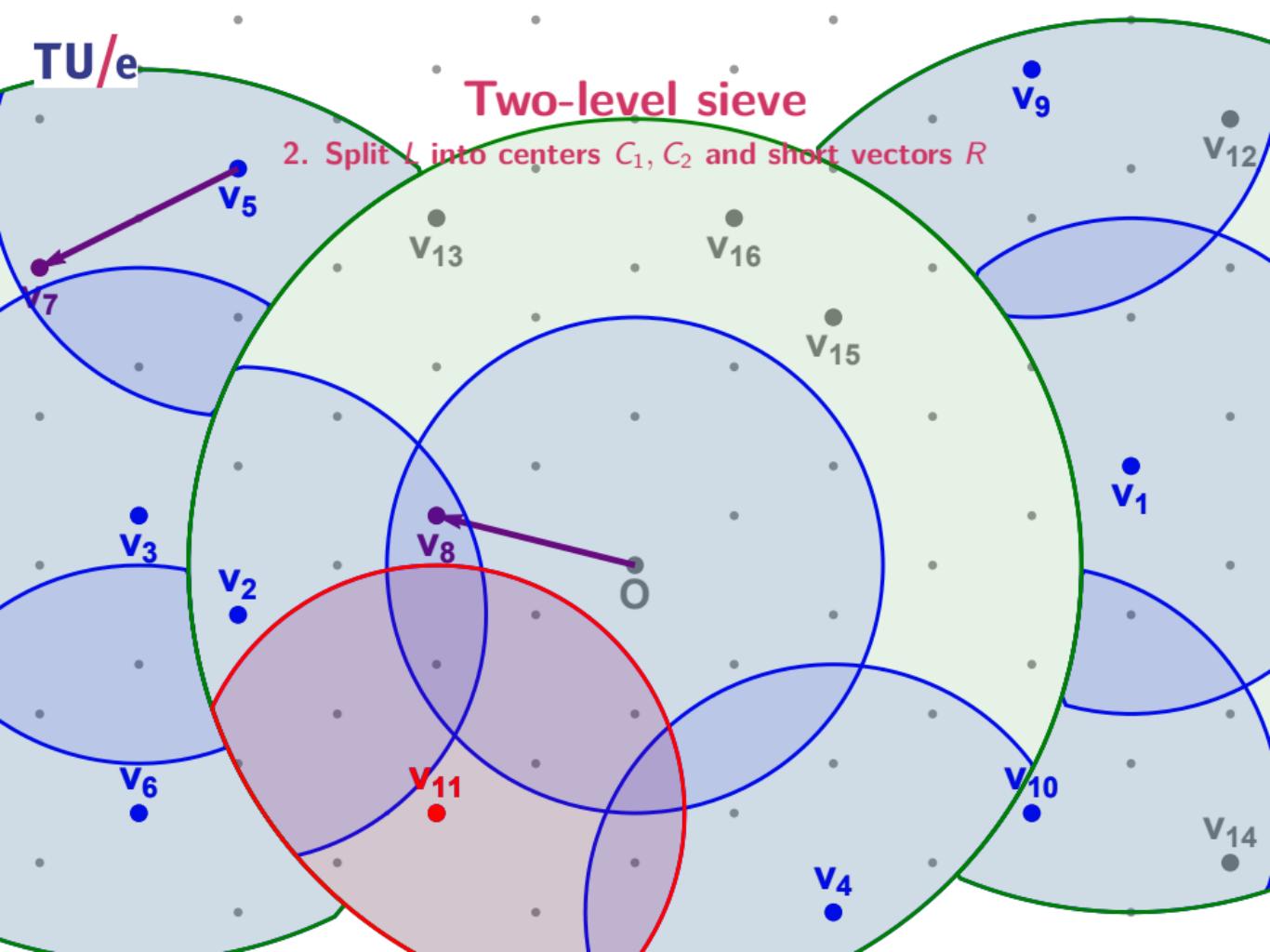
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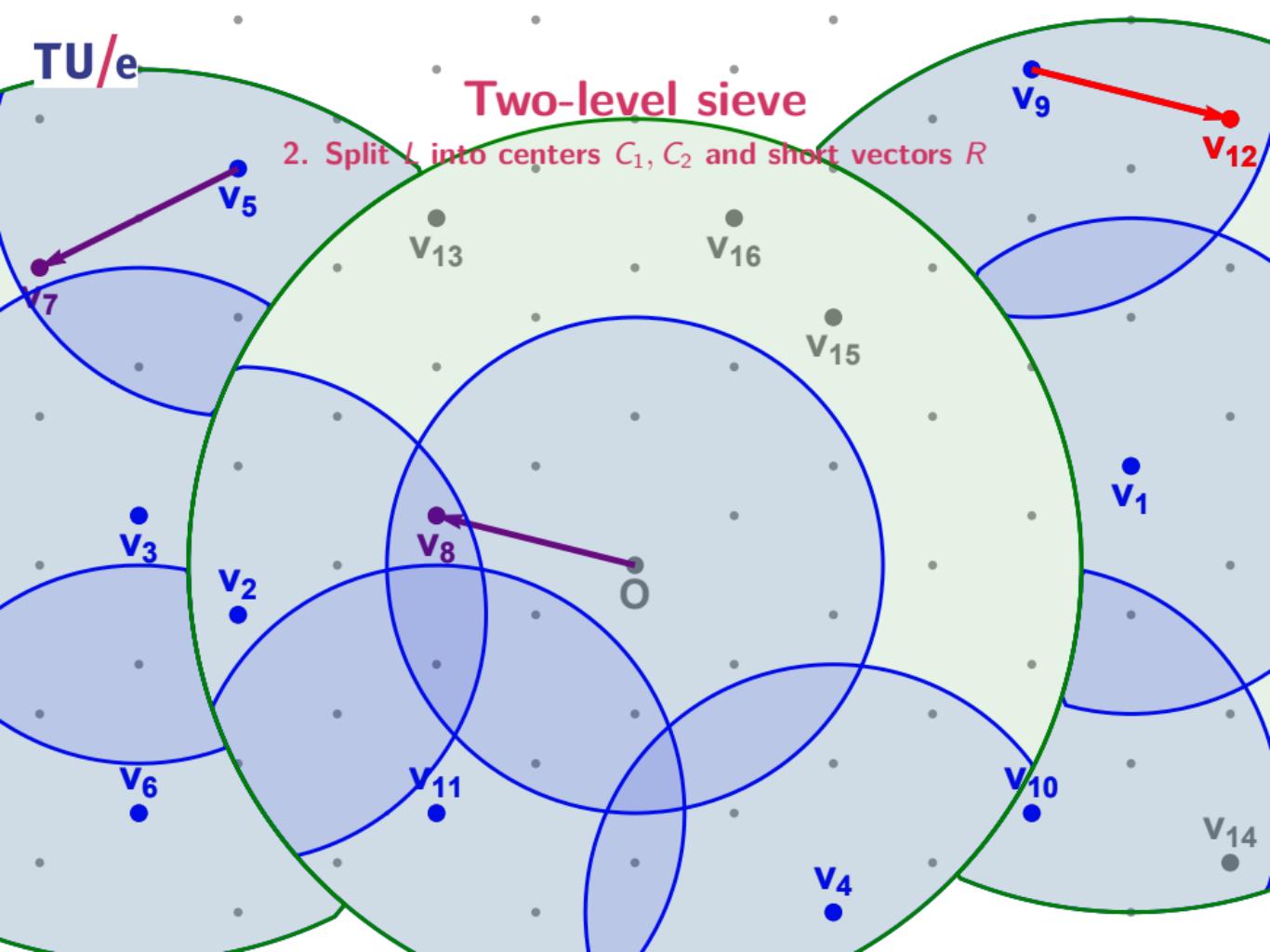
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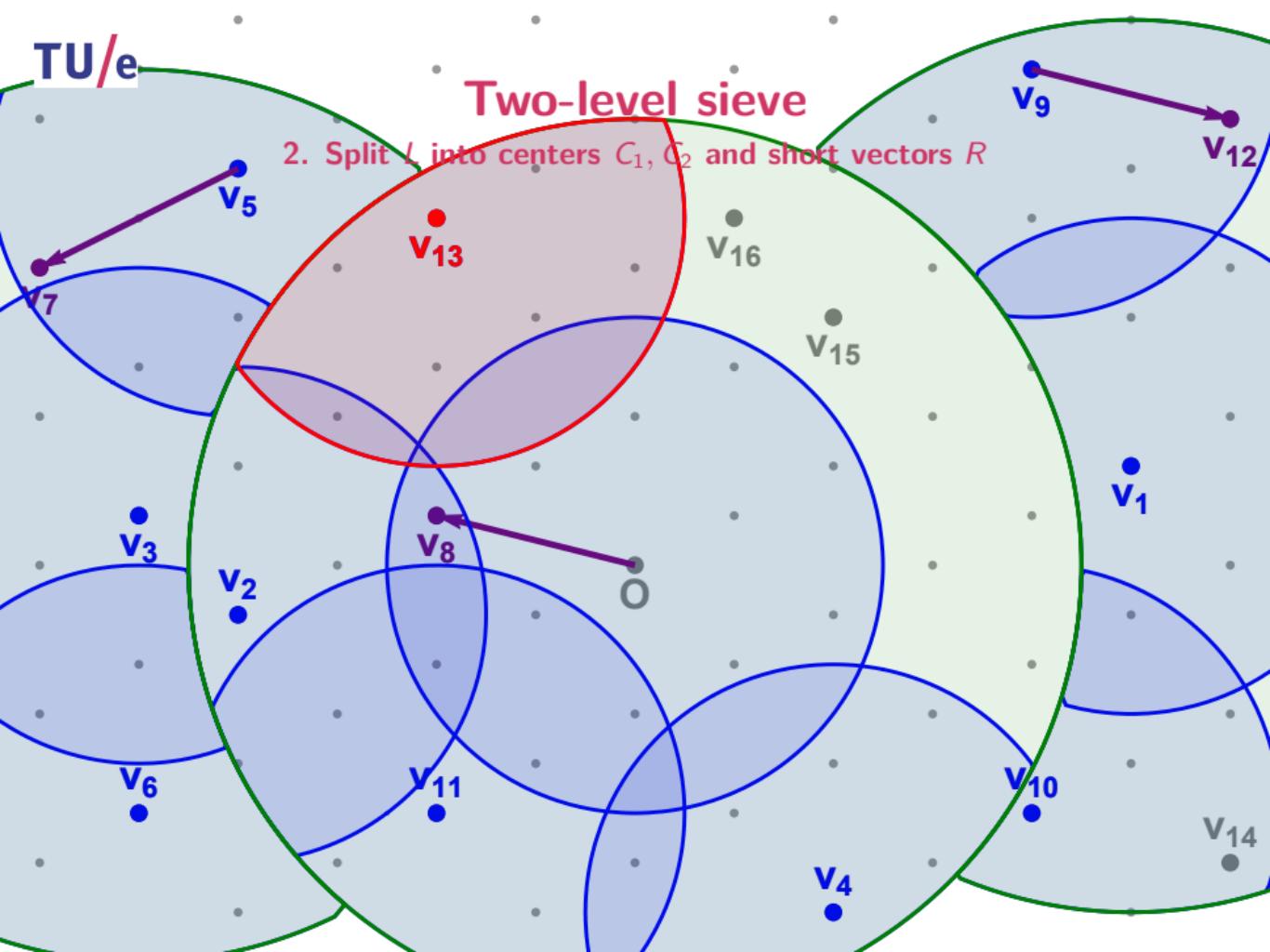
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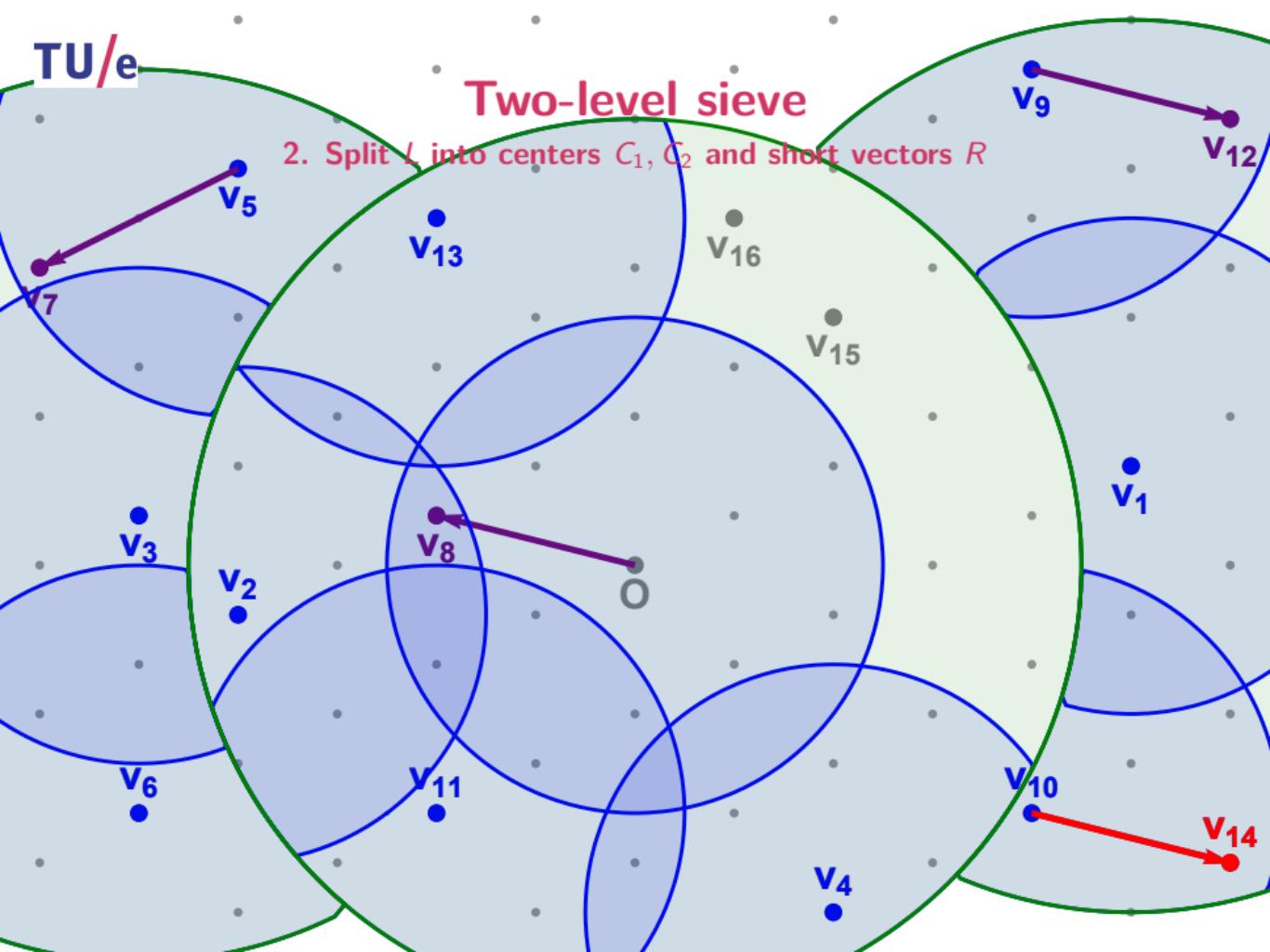
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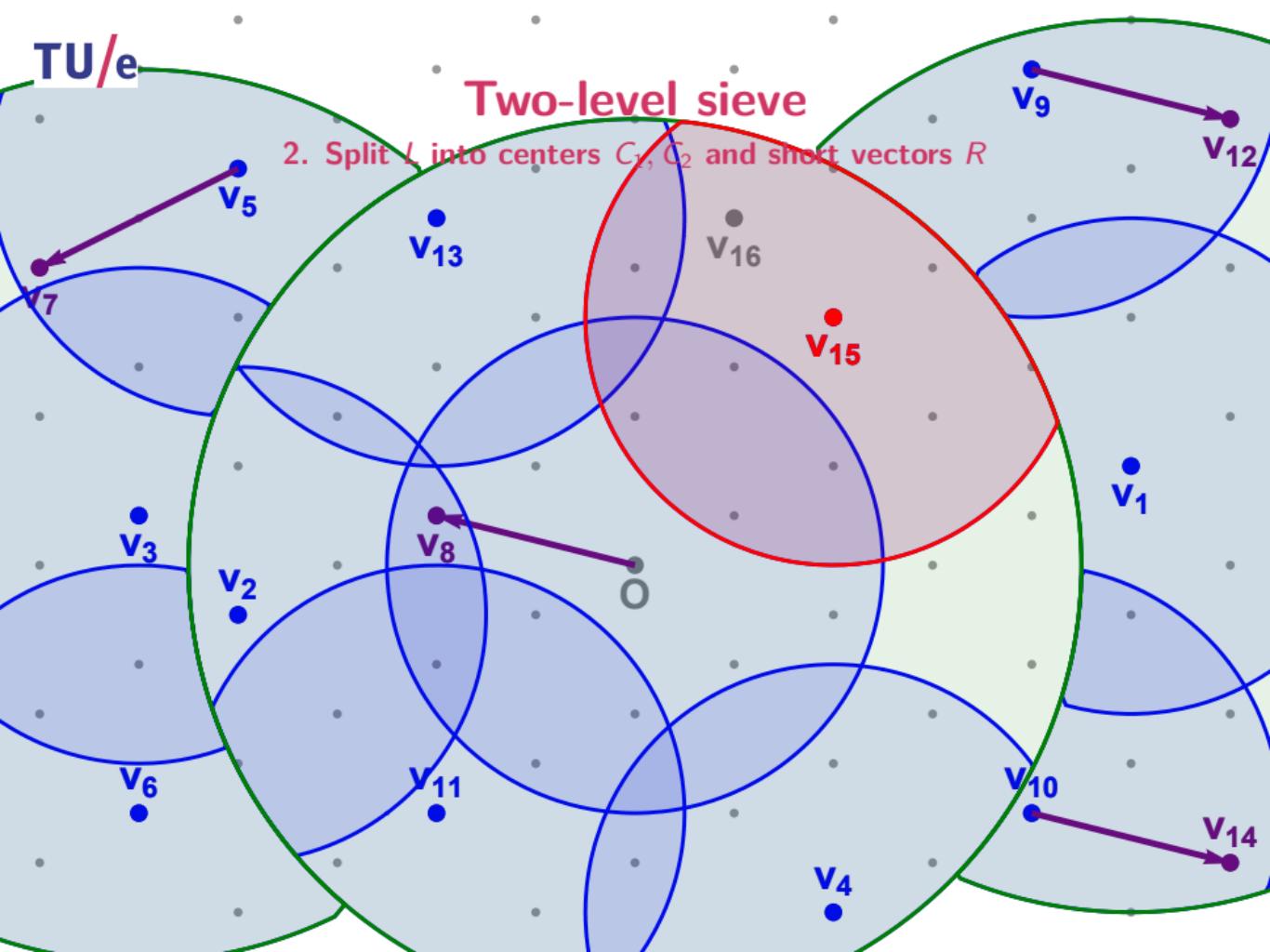
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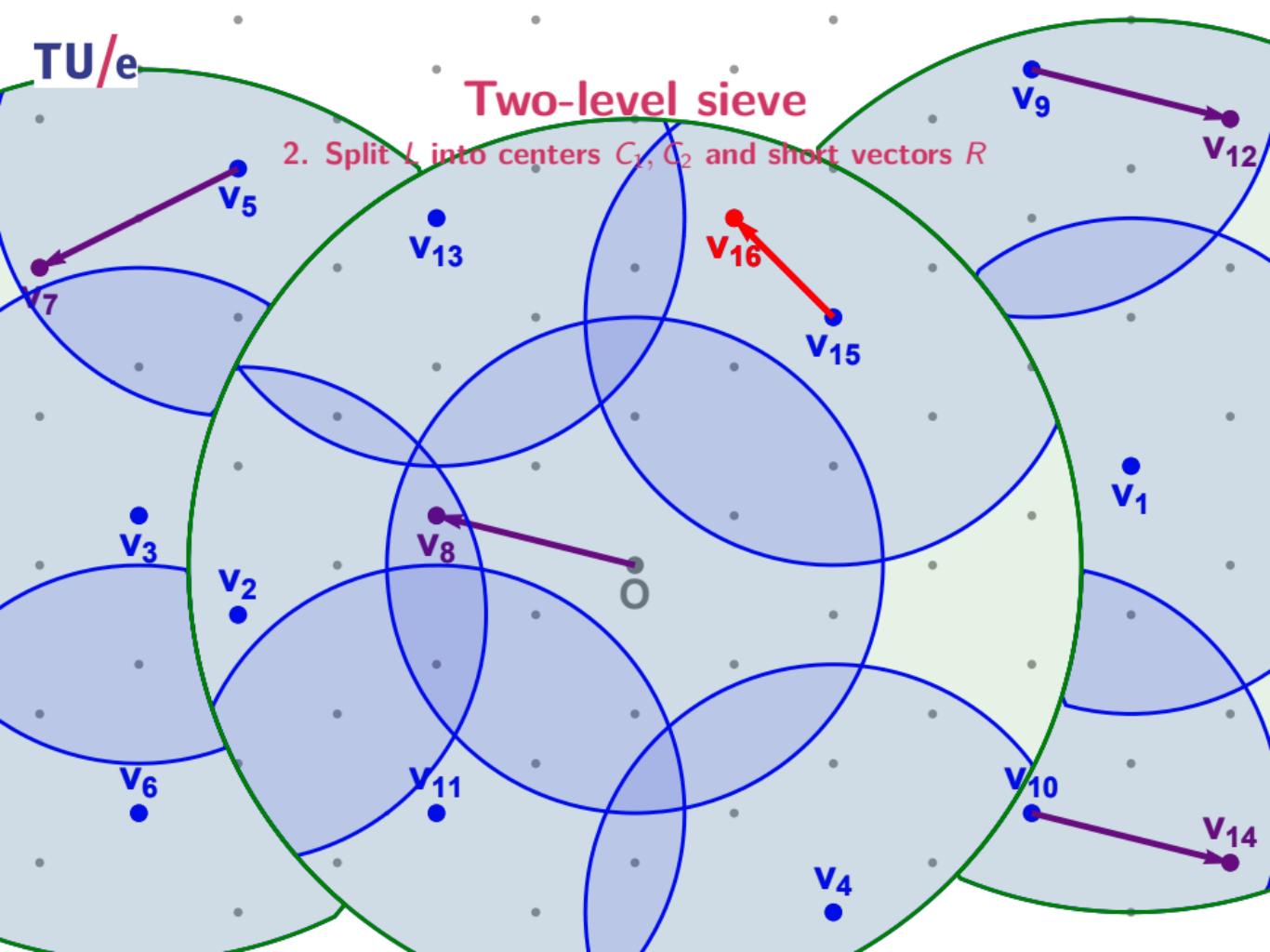
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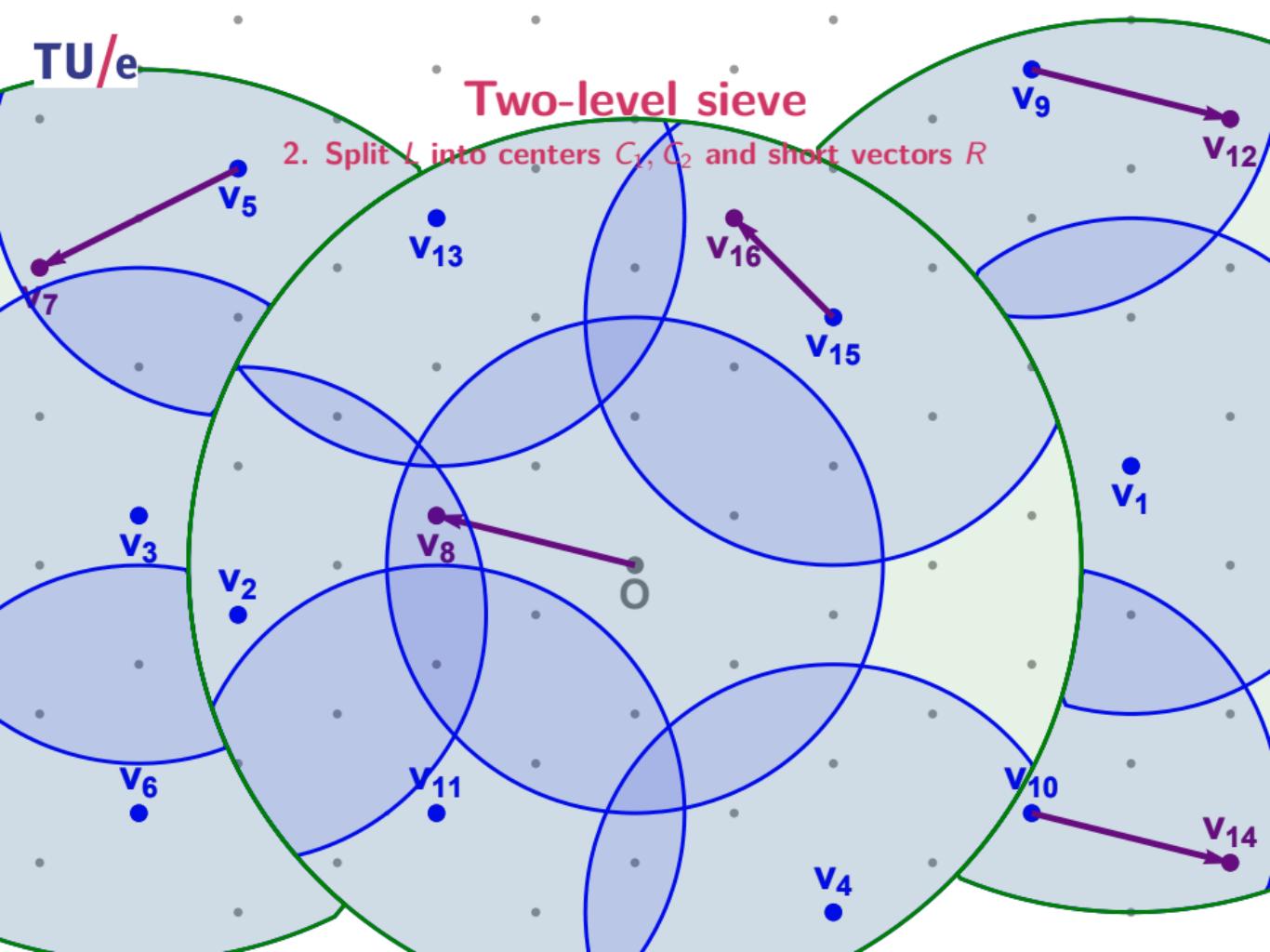
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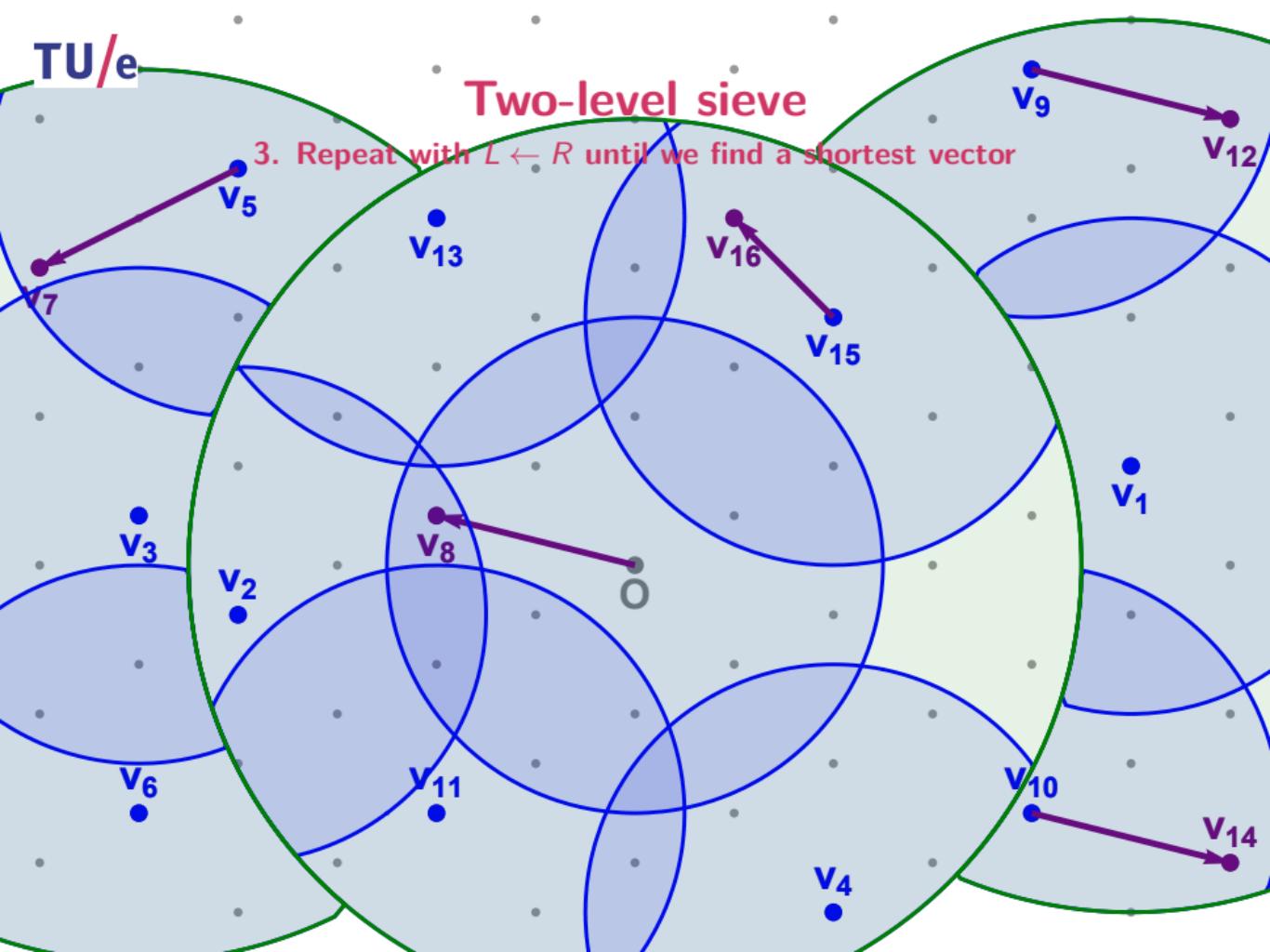
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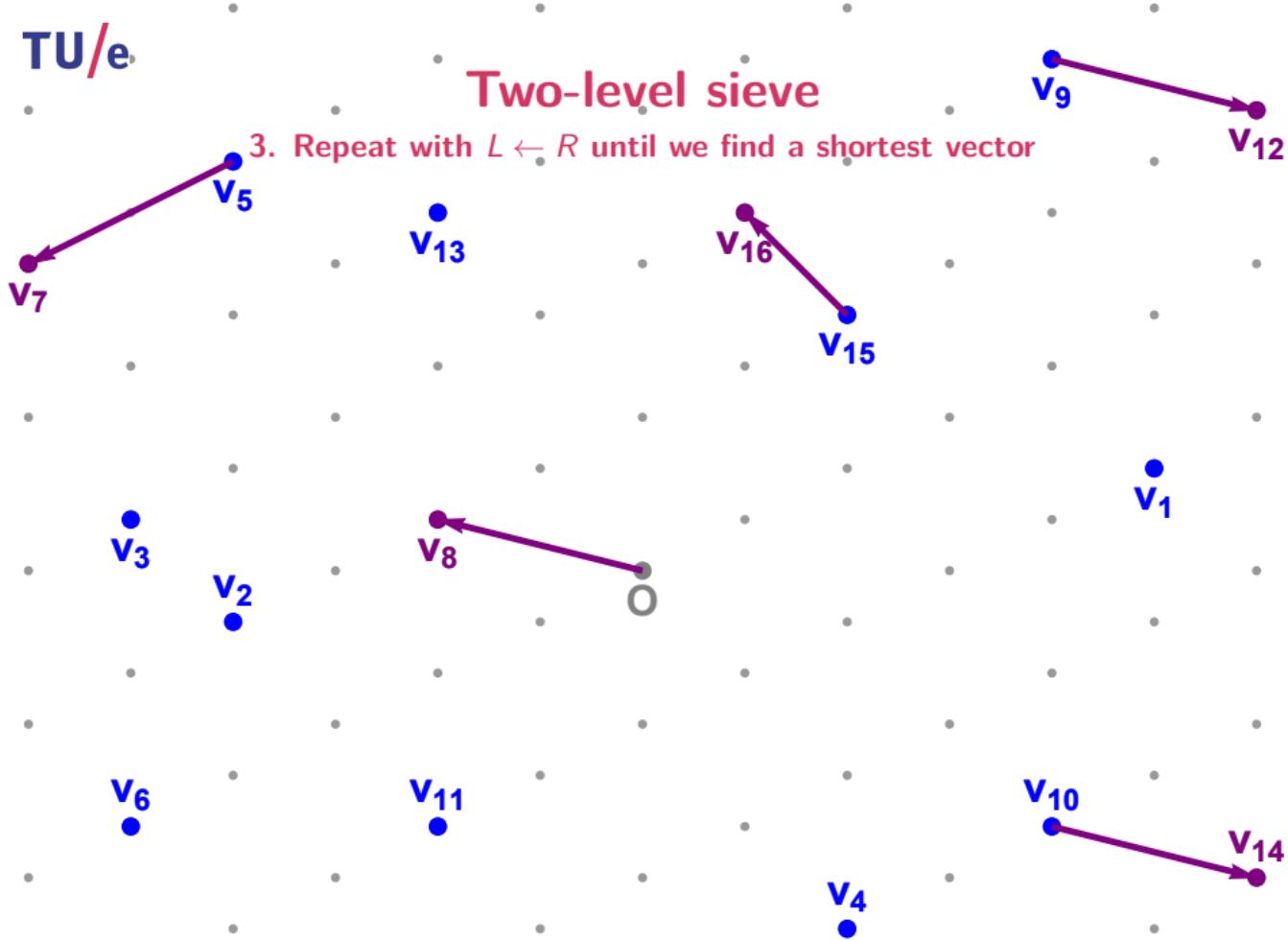
Two-level sieve

3. Repeat with $L \leftarrow R$ until we find a shortest vector



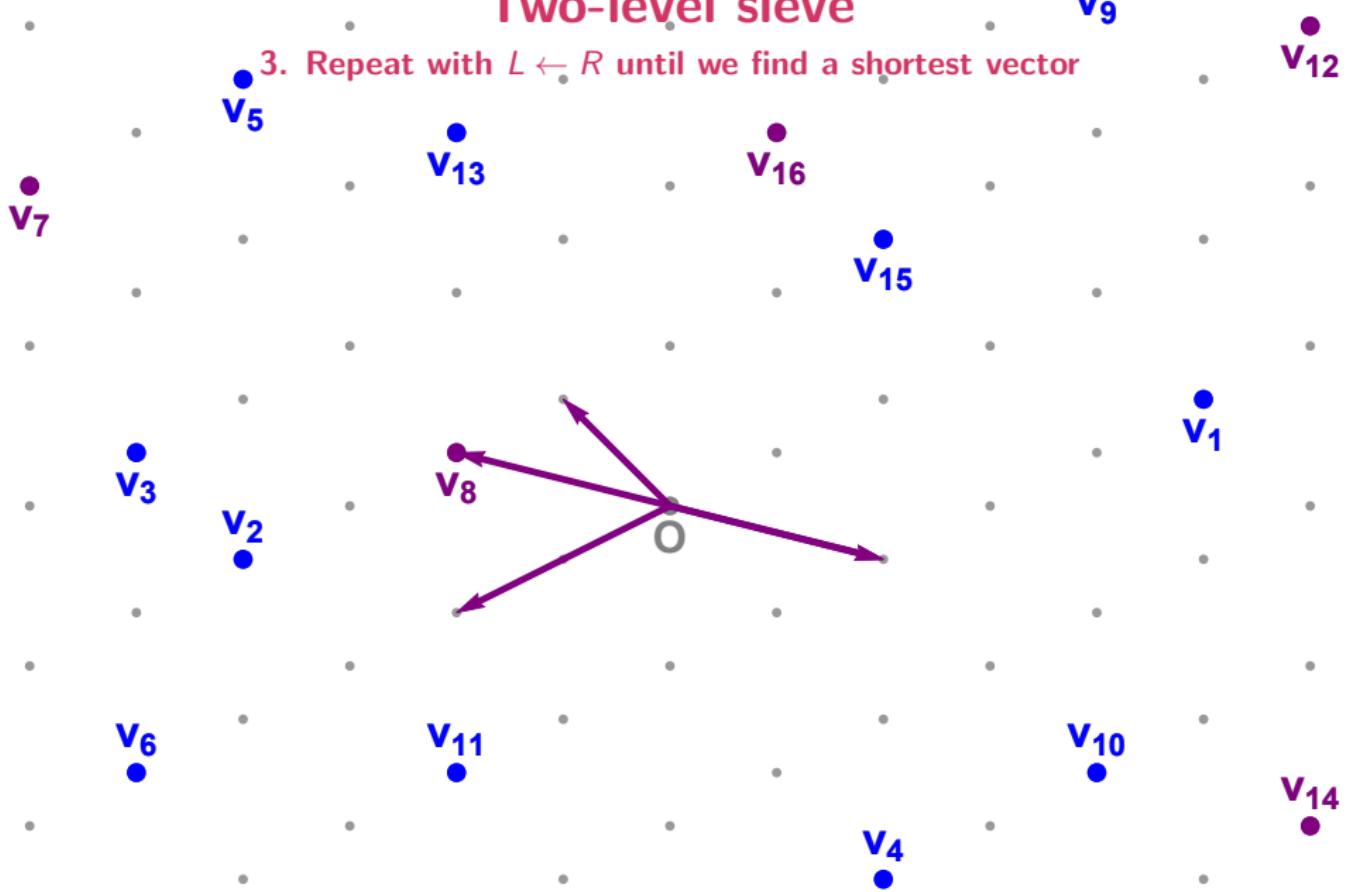
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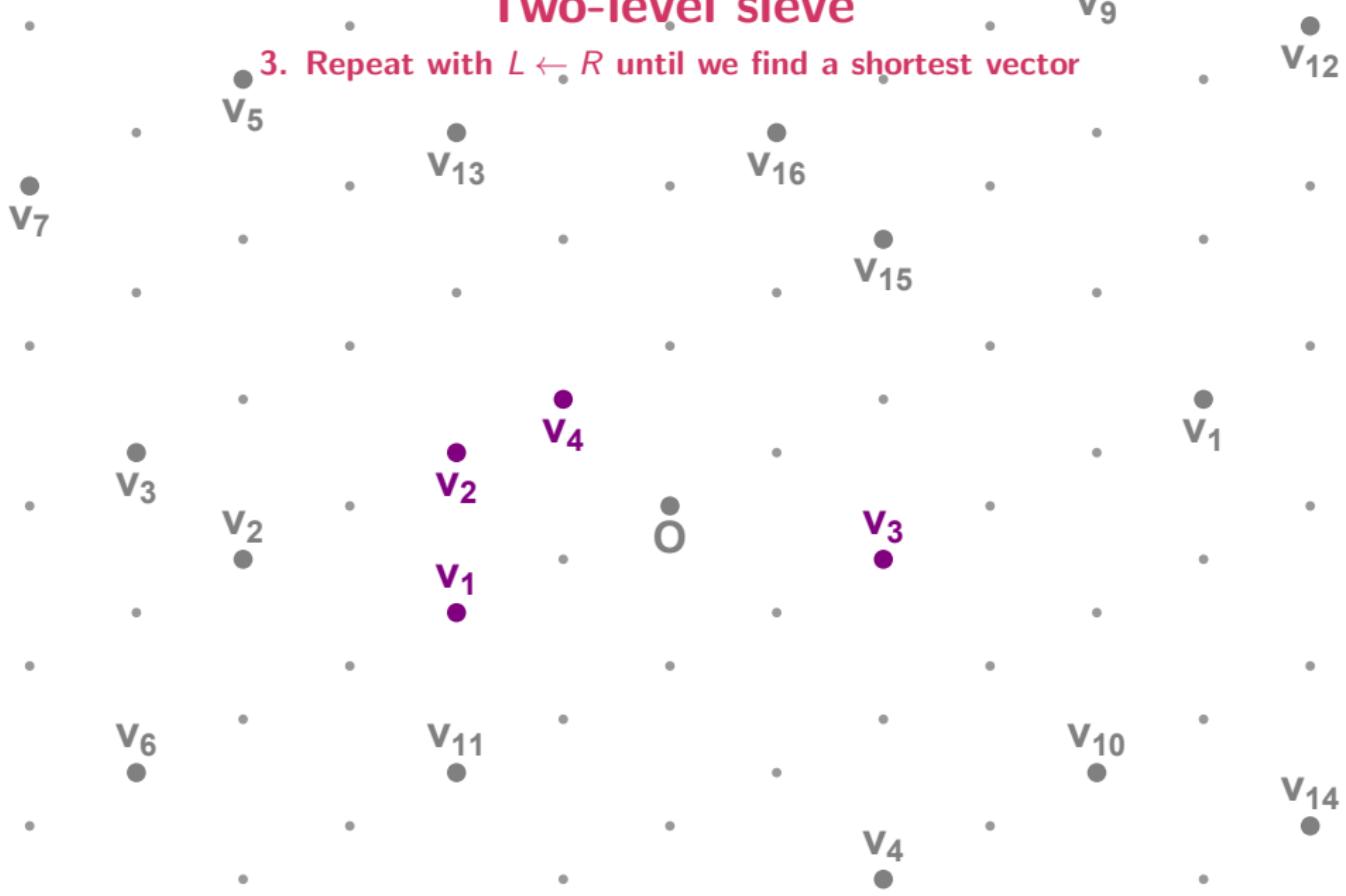
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Two-level sieve

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Multiple levels

Overview



Multiple levels

Overview

Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The one-level sieve runs in time $2^{0.4150n}$ and space $2^{0.2075n}$.

v₃

v₂

v₂

v₁

v₄

v₆

v₁₁

v₁₅

v₃

v₄

v₁₀

v₁₄

v₁

v₅

v₇

v₉

v₁₂

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Heuristic (Wang et al., ASIACCS'11)

The two-level sieve runs in time $2^{0.3836n}$ and space $2^{0.2557n}$.

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Heuristic (Zhang et al., SAC'13)

The three-level sieve runs in time $2^{0.3778n}$ and space $2^{0.2833n}$.

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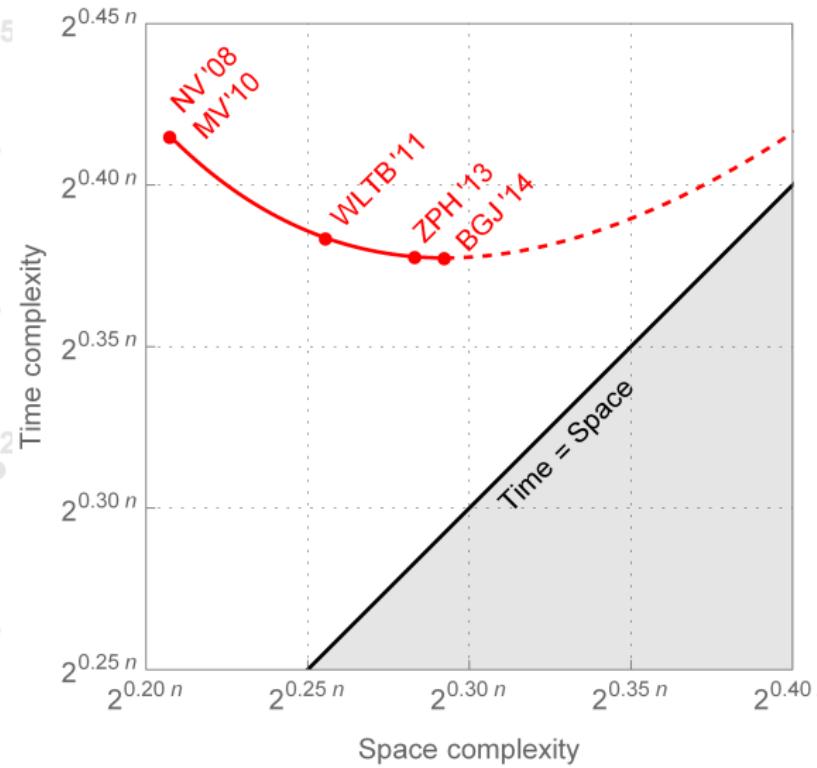
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Conjecture

The four-level sieve runs in time $2^{0.3774n}$ and space $2^{0.2925n}$, and higher-level sieves are not faster than this.

Sieving

Space/time trade-off



Quantum Search

Classical form

Problem: Given a list L of size N , and a function $f : L \rightarrow \{0, 1\}$ such that there is exactly one element $e \in L$ with $f(e) = 1$. Find this element e .

- Classical search: $\Theta(N)$ time
- Quantum search: $\Theta(\sqrt{N})$ time [Gro96]

Quantum Search

General form

Problem: Given a list L of size N , and a function $f : L \rightarrow \{0, 1\}$ such that there are $c = O(1)$ elements $e \in L$ with $f(e) = 1$. Find one such element e .

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Quantum Search

General form

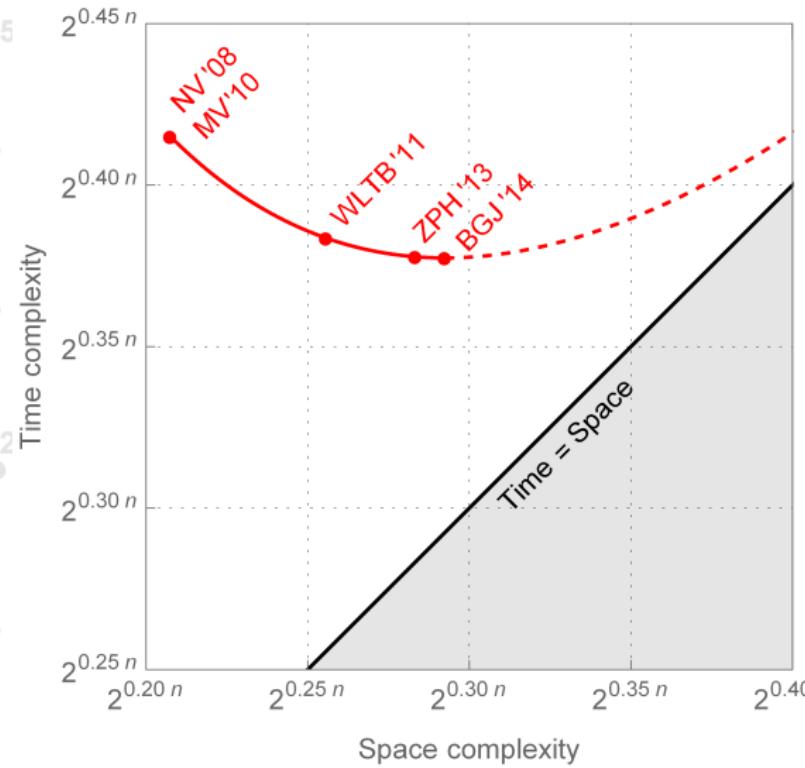
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Potentially speed up search subroutines in sieving

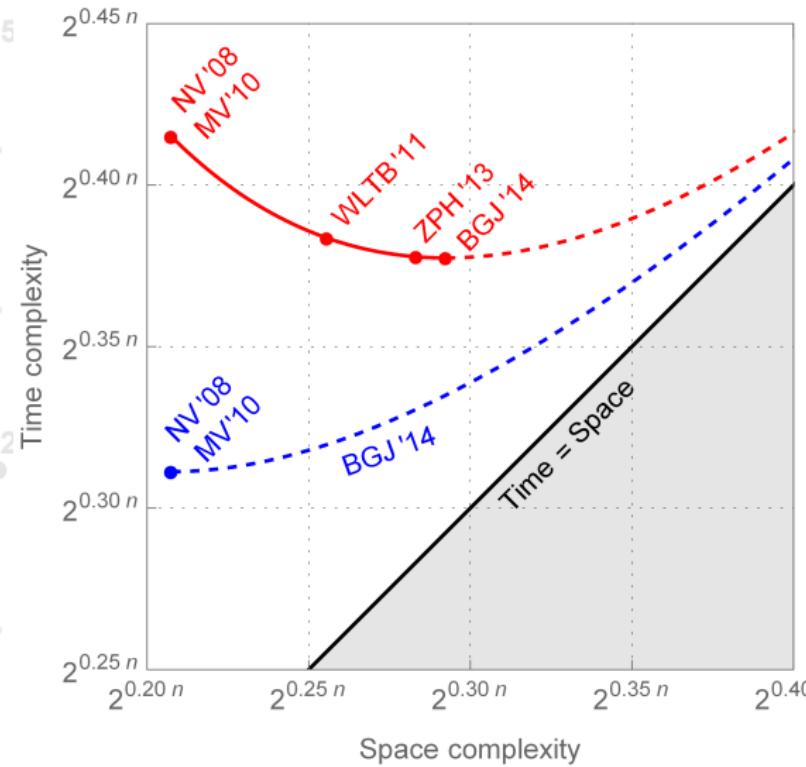
Sieving

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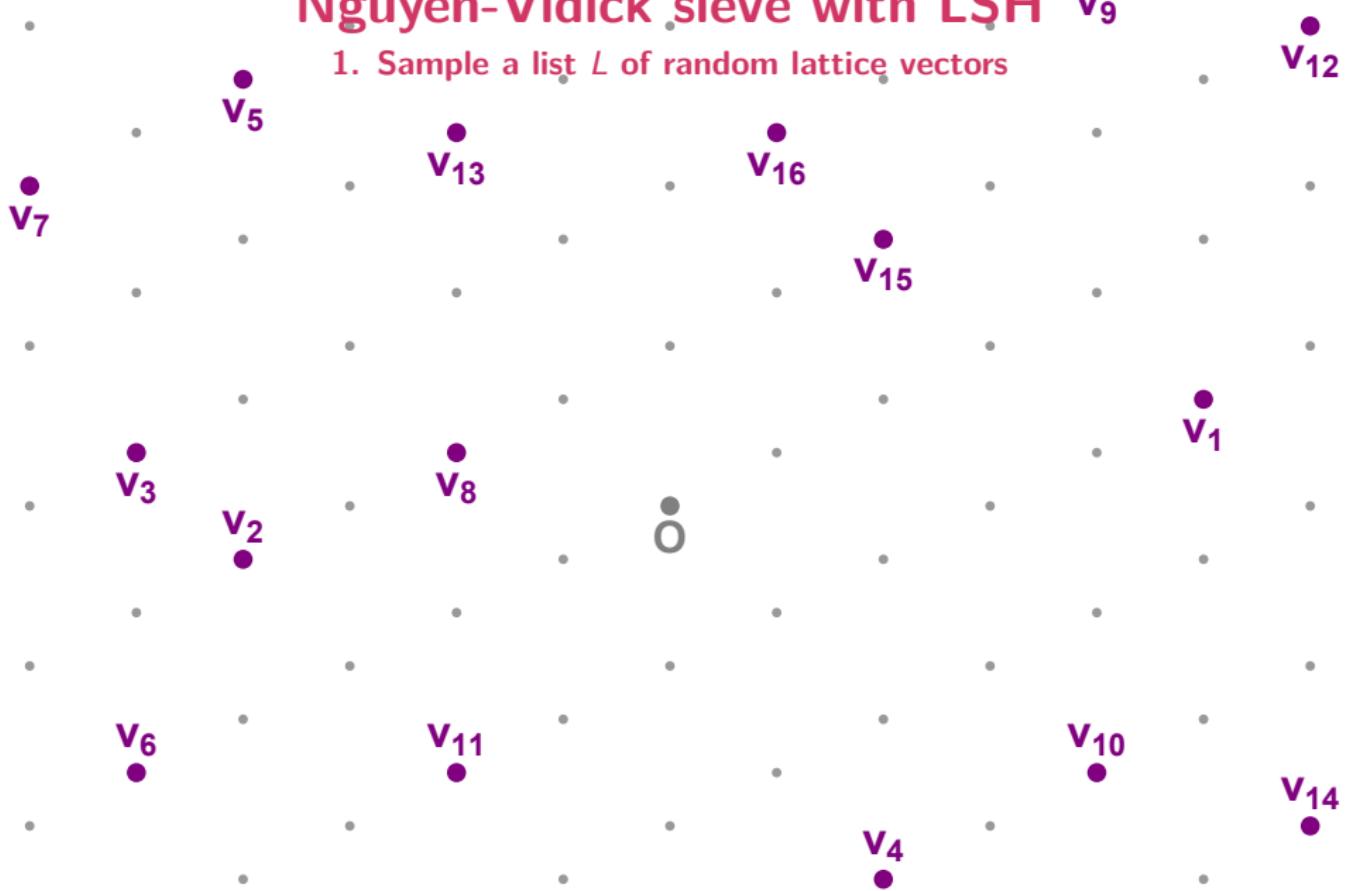
Nguyen-Vidick sieve with LSH

1. Sample a list L of random lattice vectors



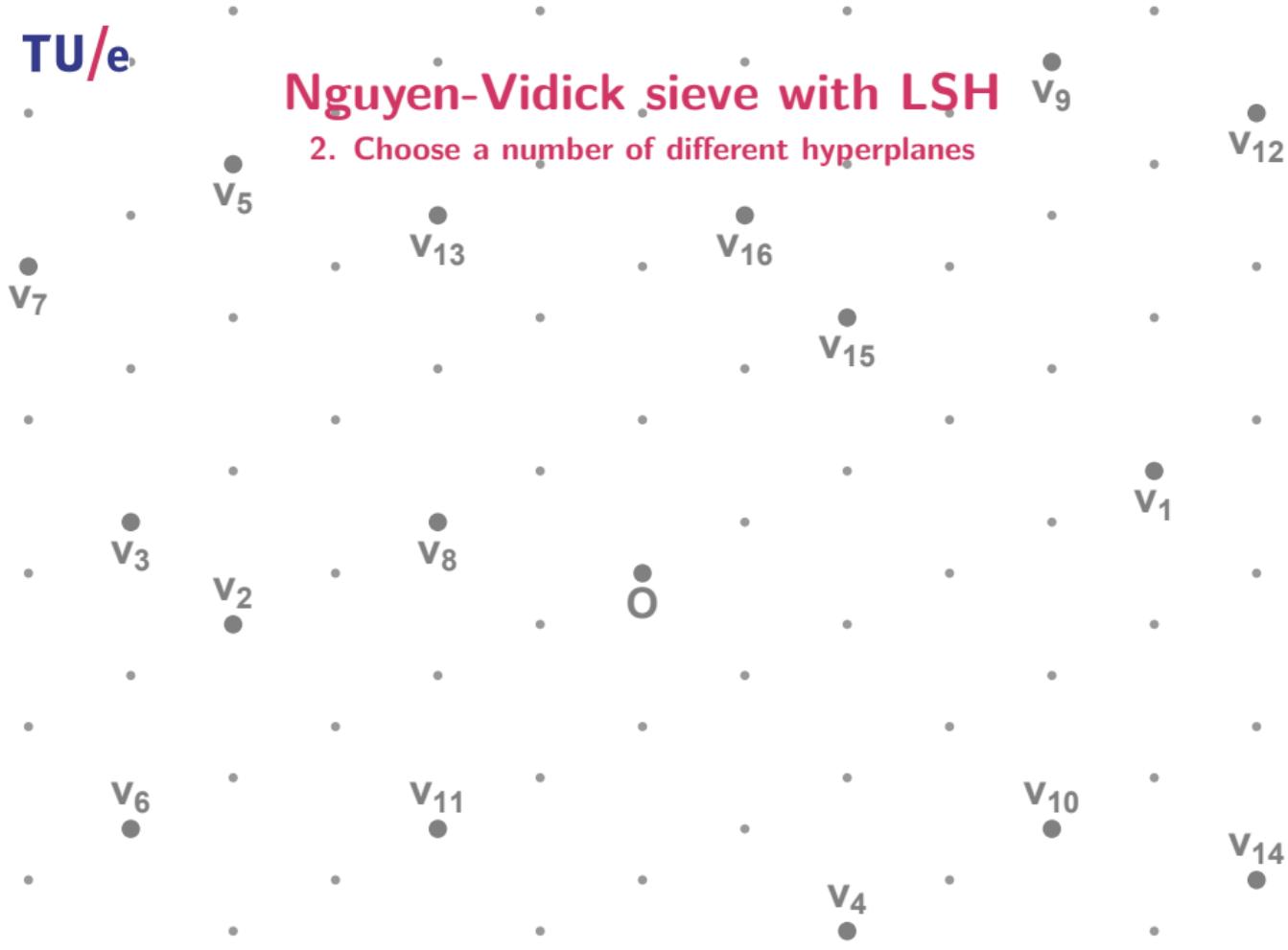
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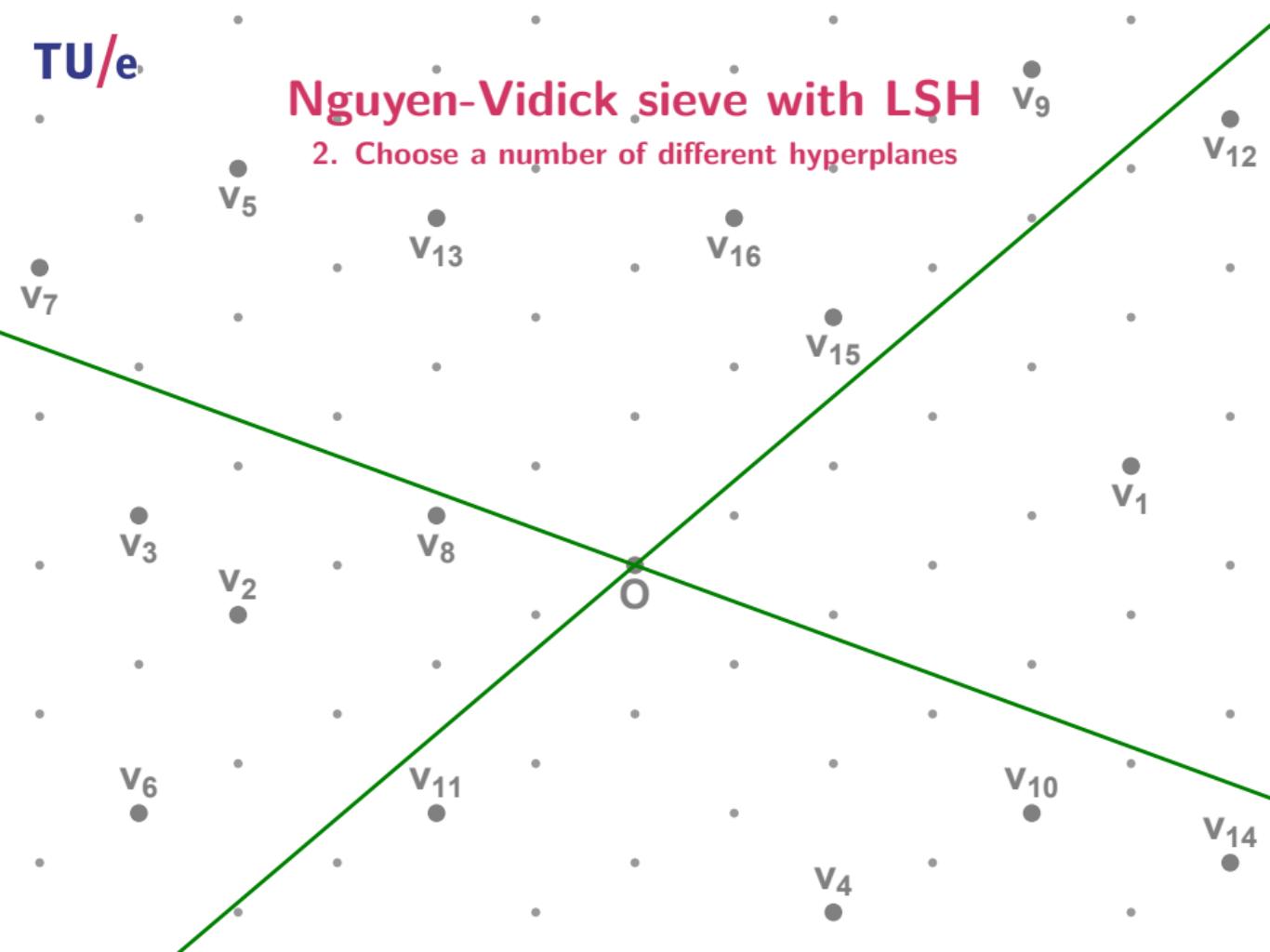
Nguyen-Vidick sieve with LSH

2. Choose a number of different hyperplanes



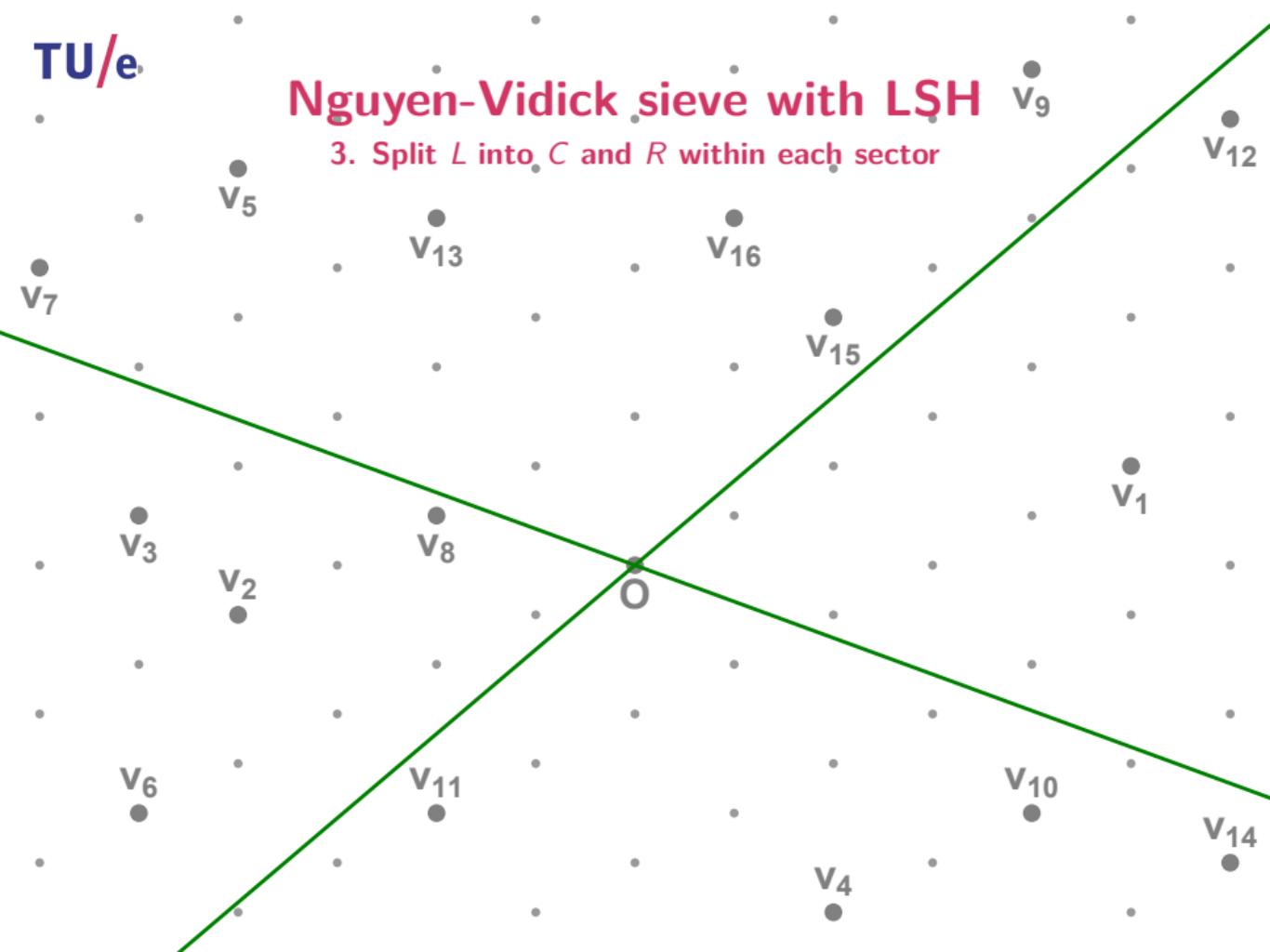
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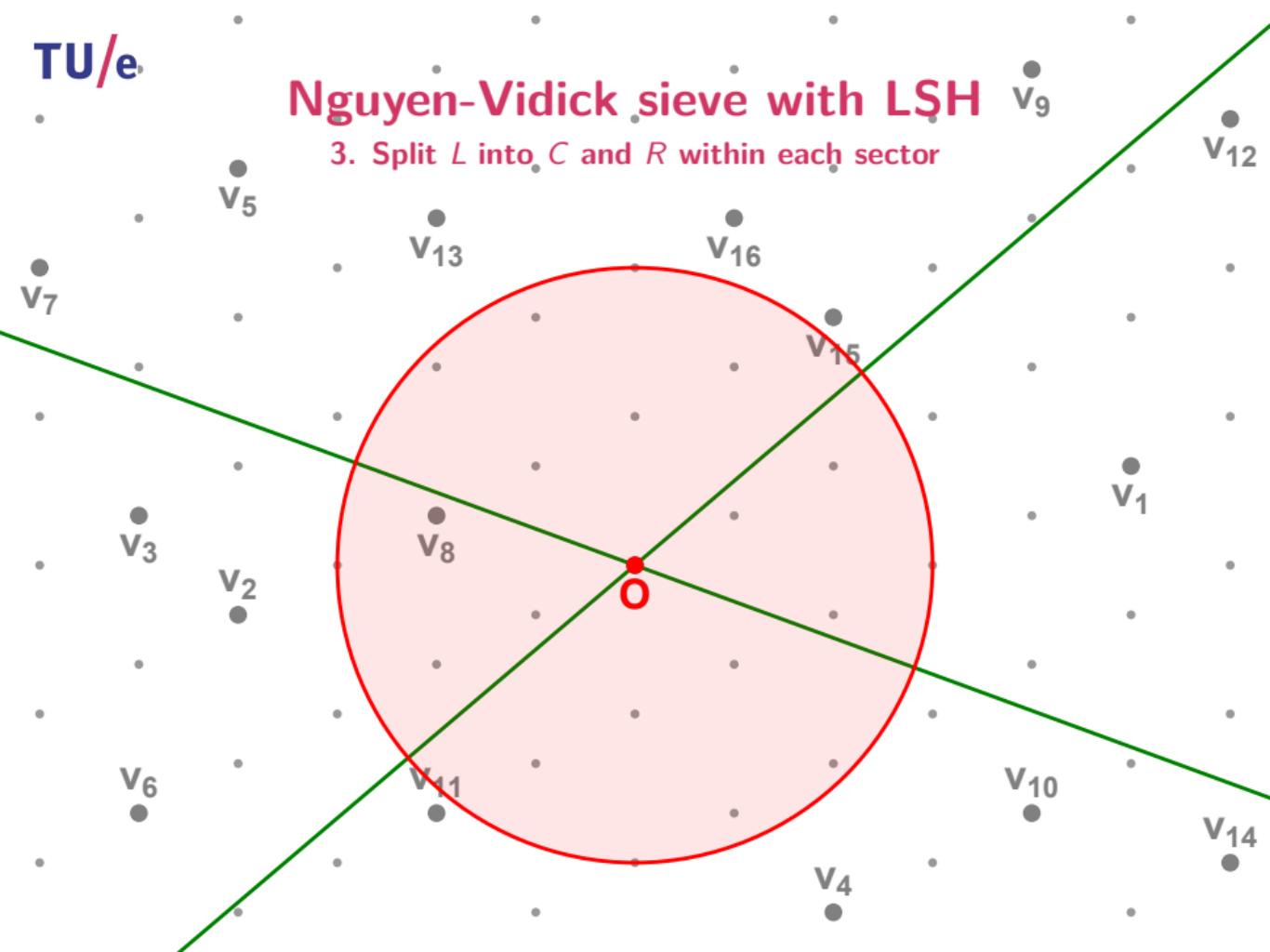
Nguyen-Vidick sieve with LSH

3. Split L into C and R within each sector



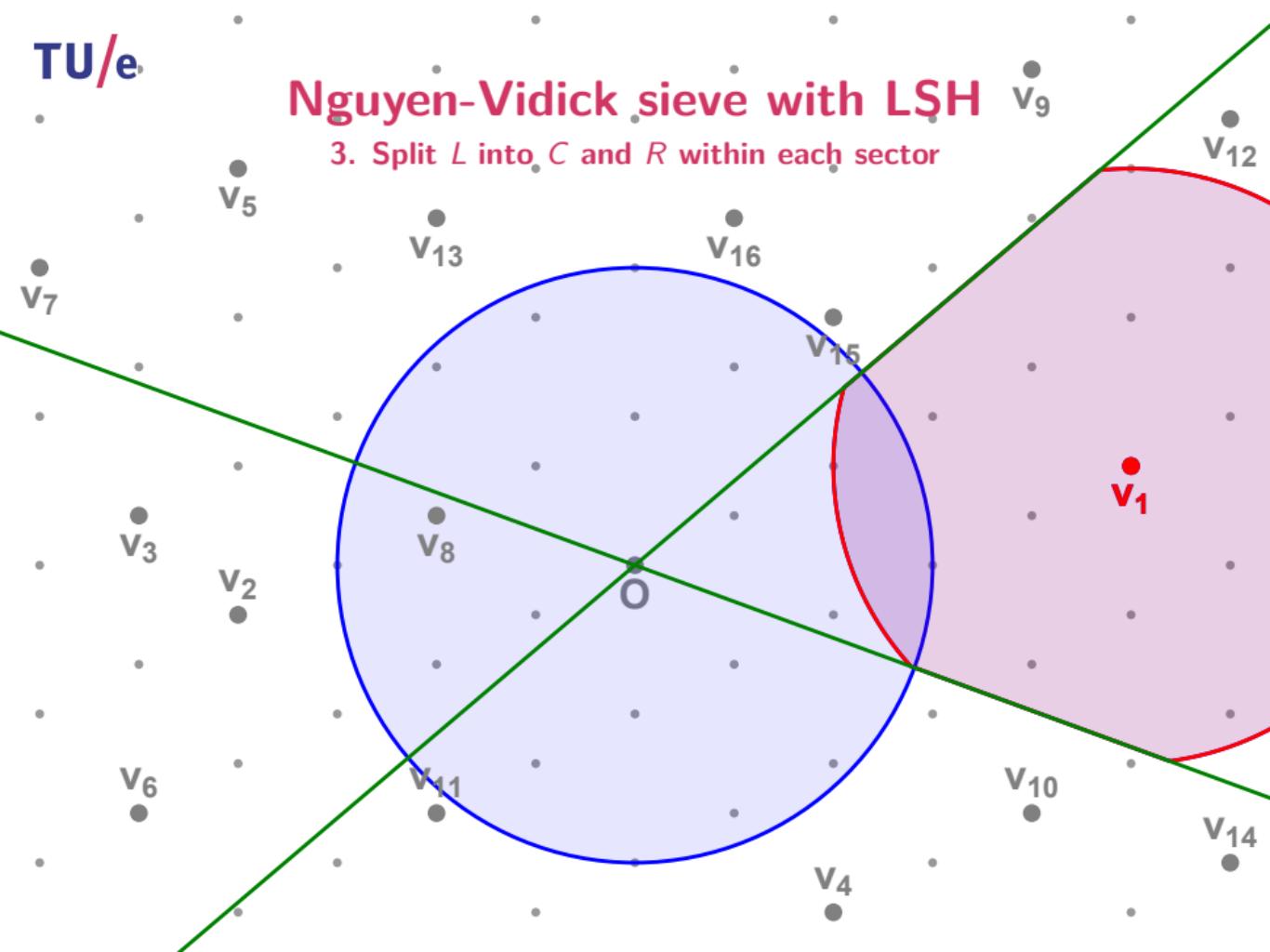
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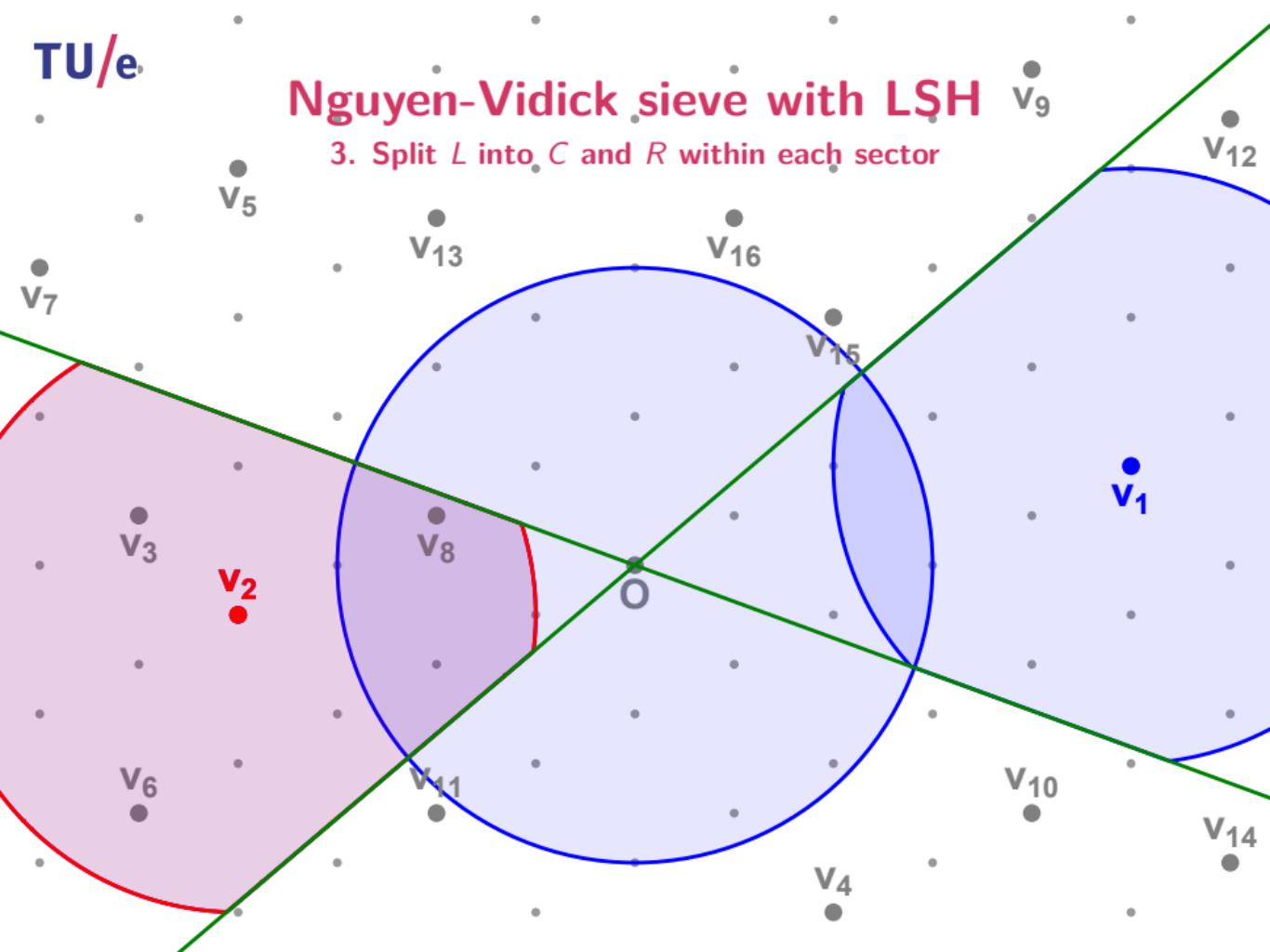
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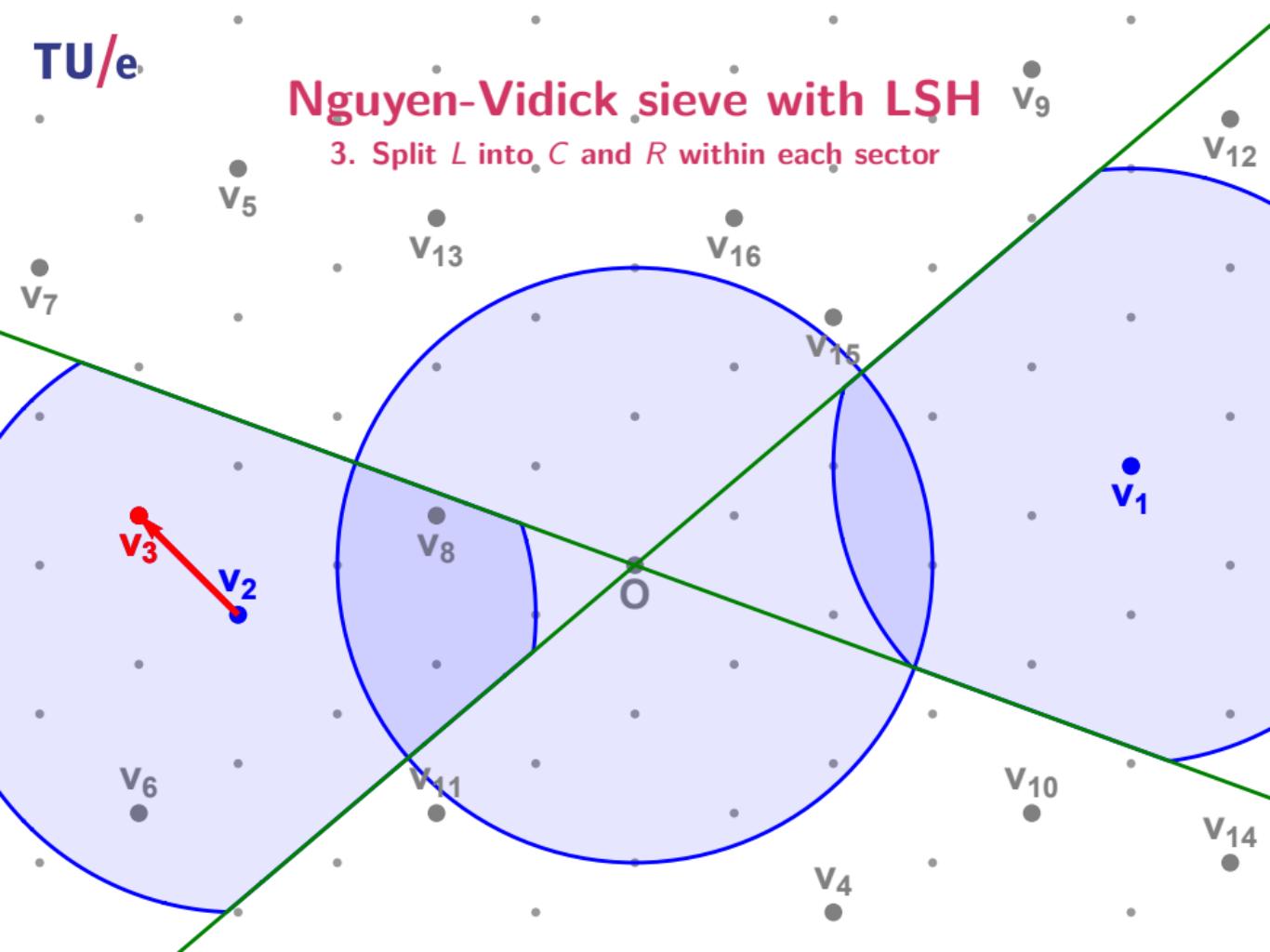
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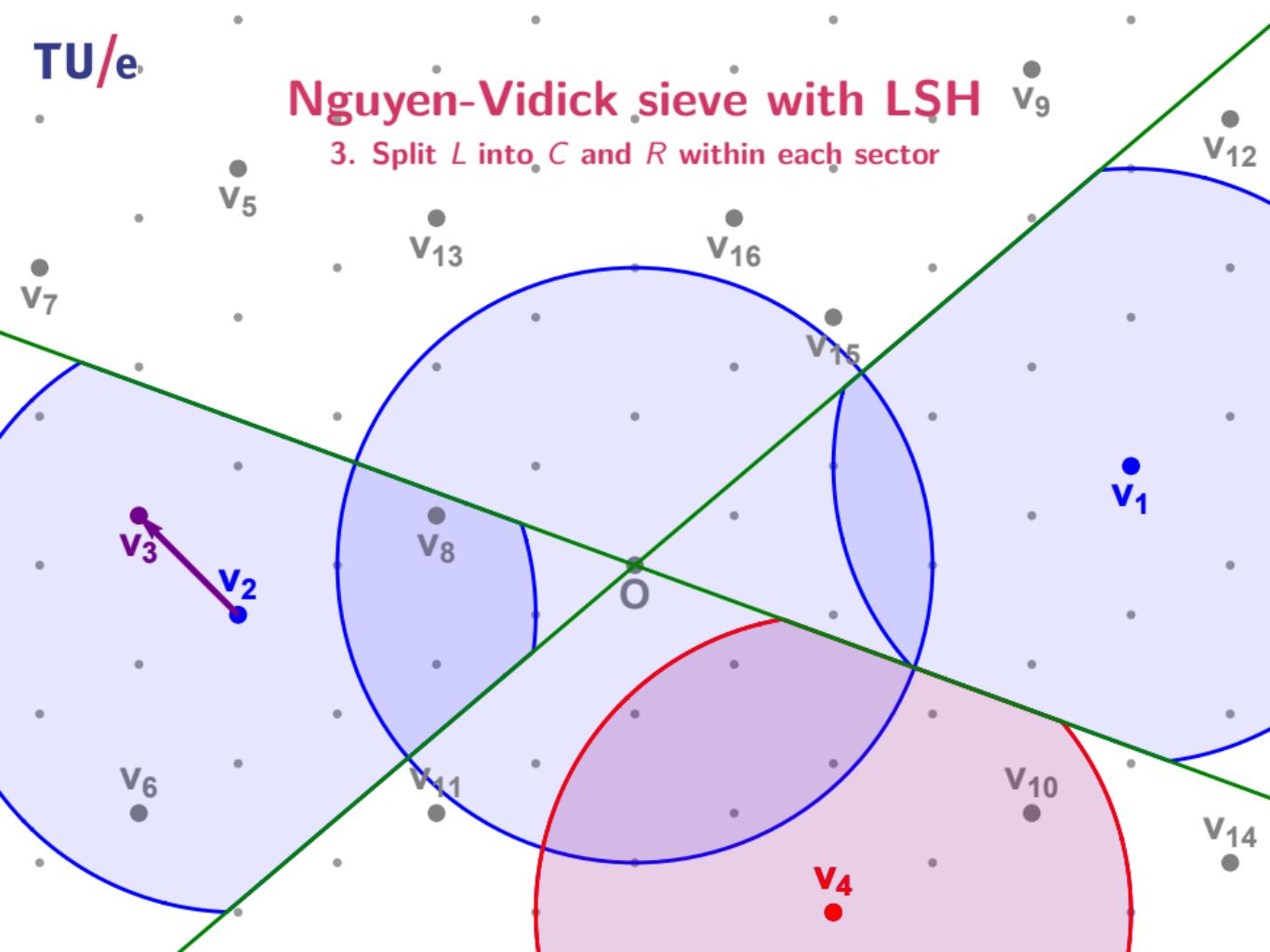
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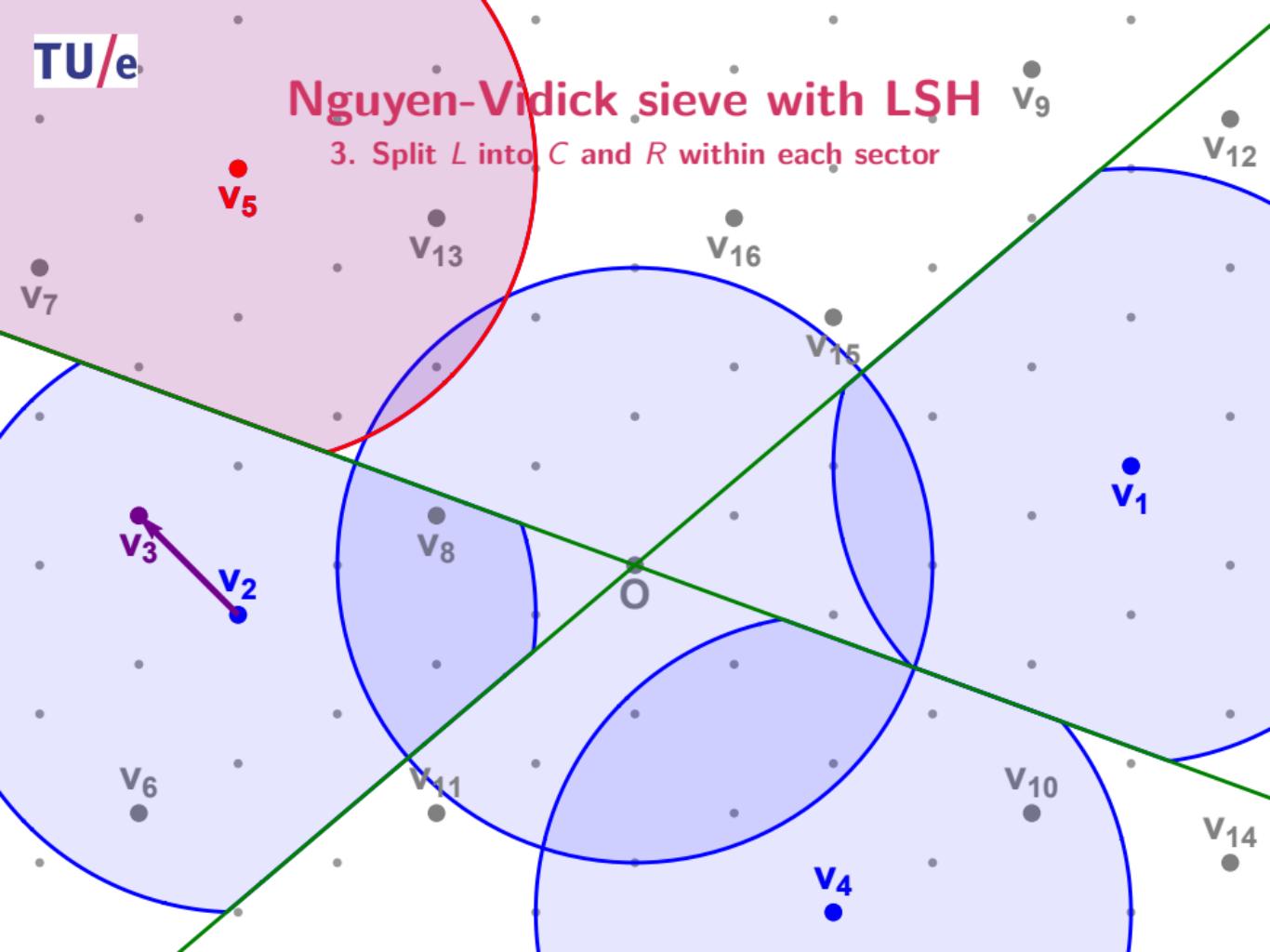
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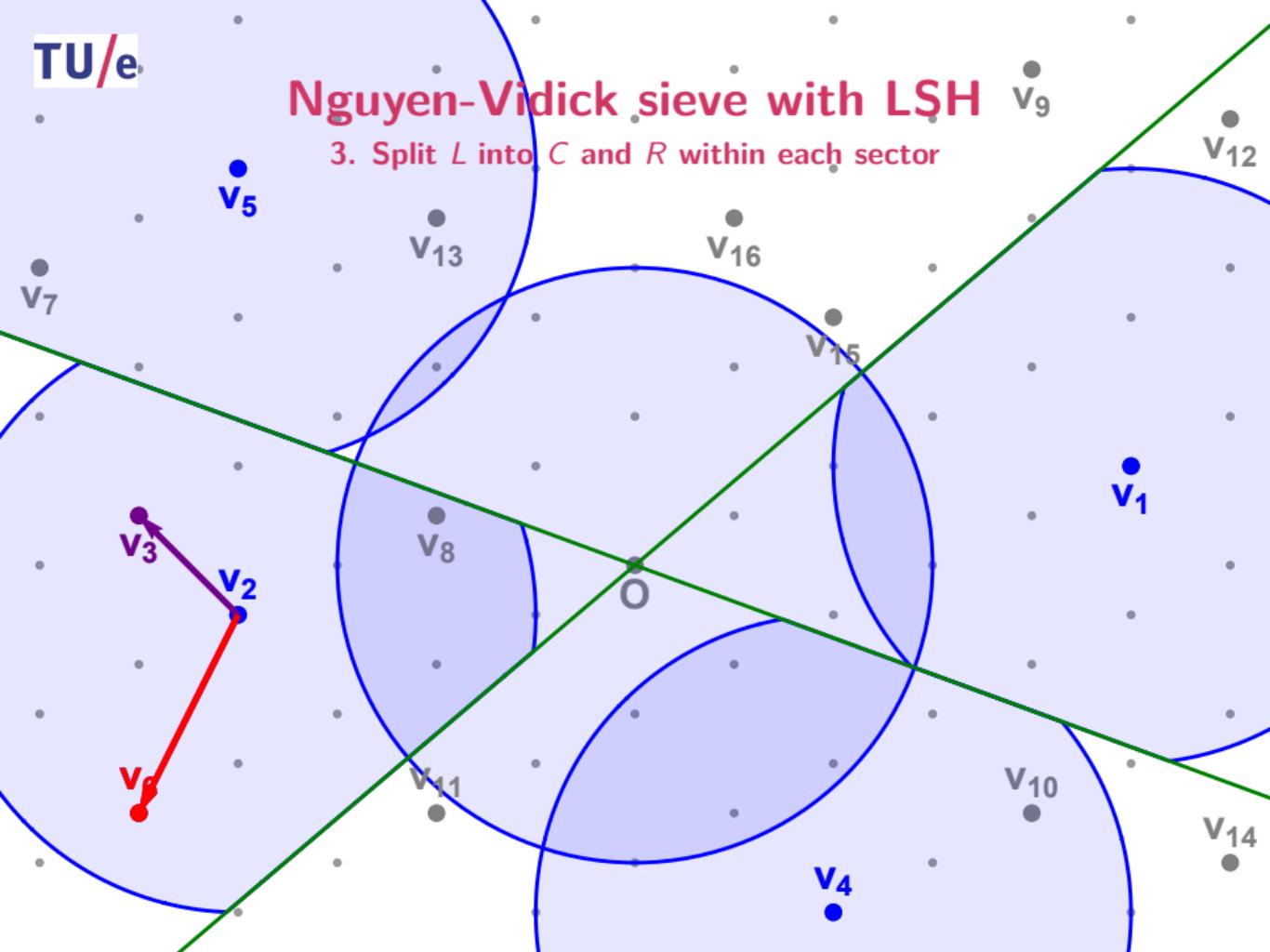
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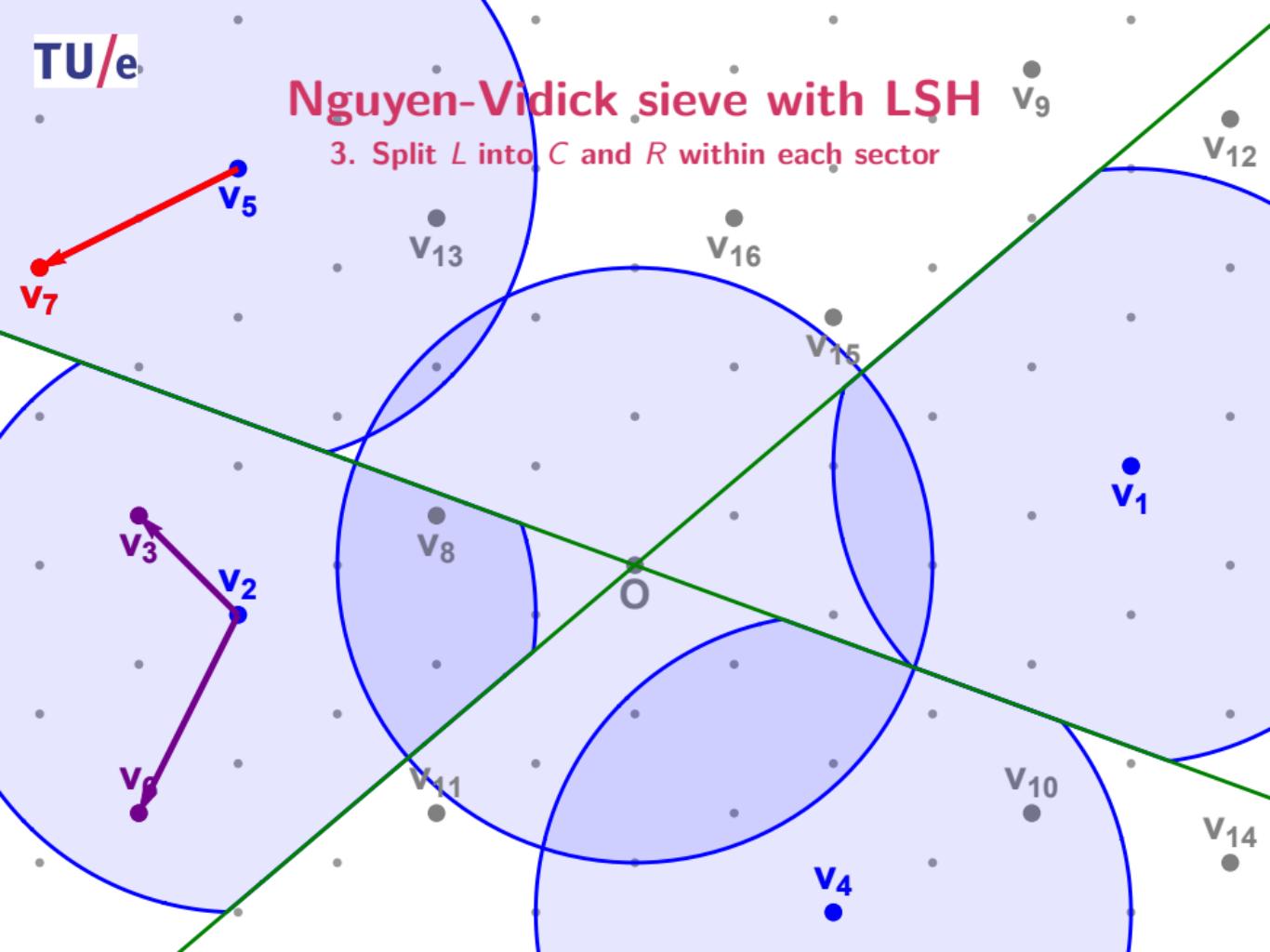
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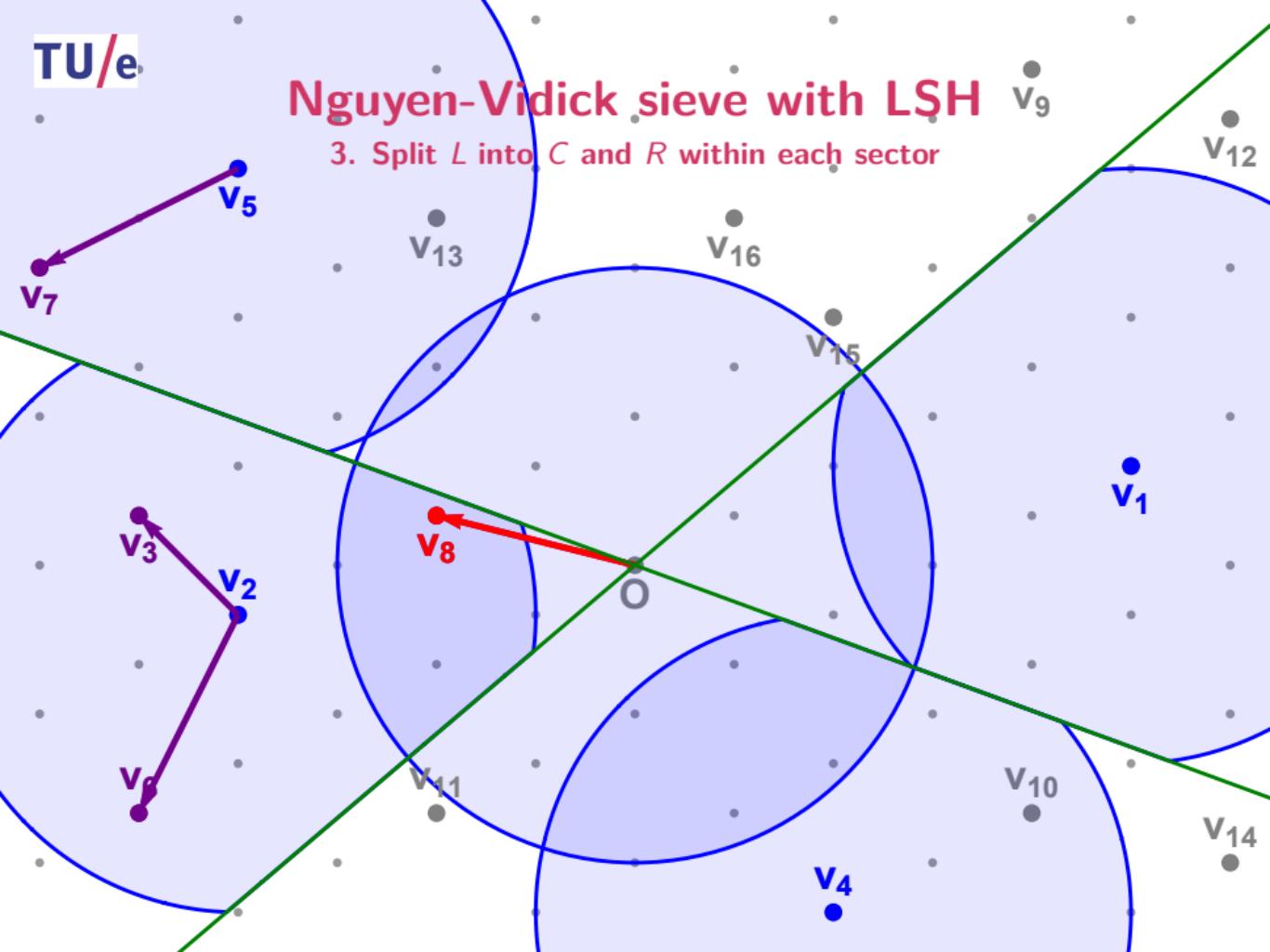
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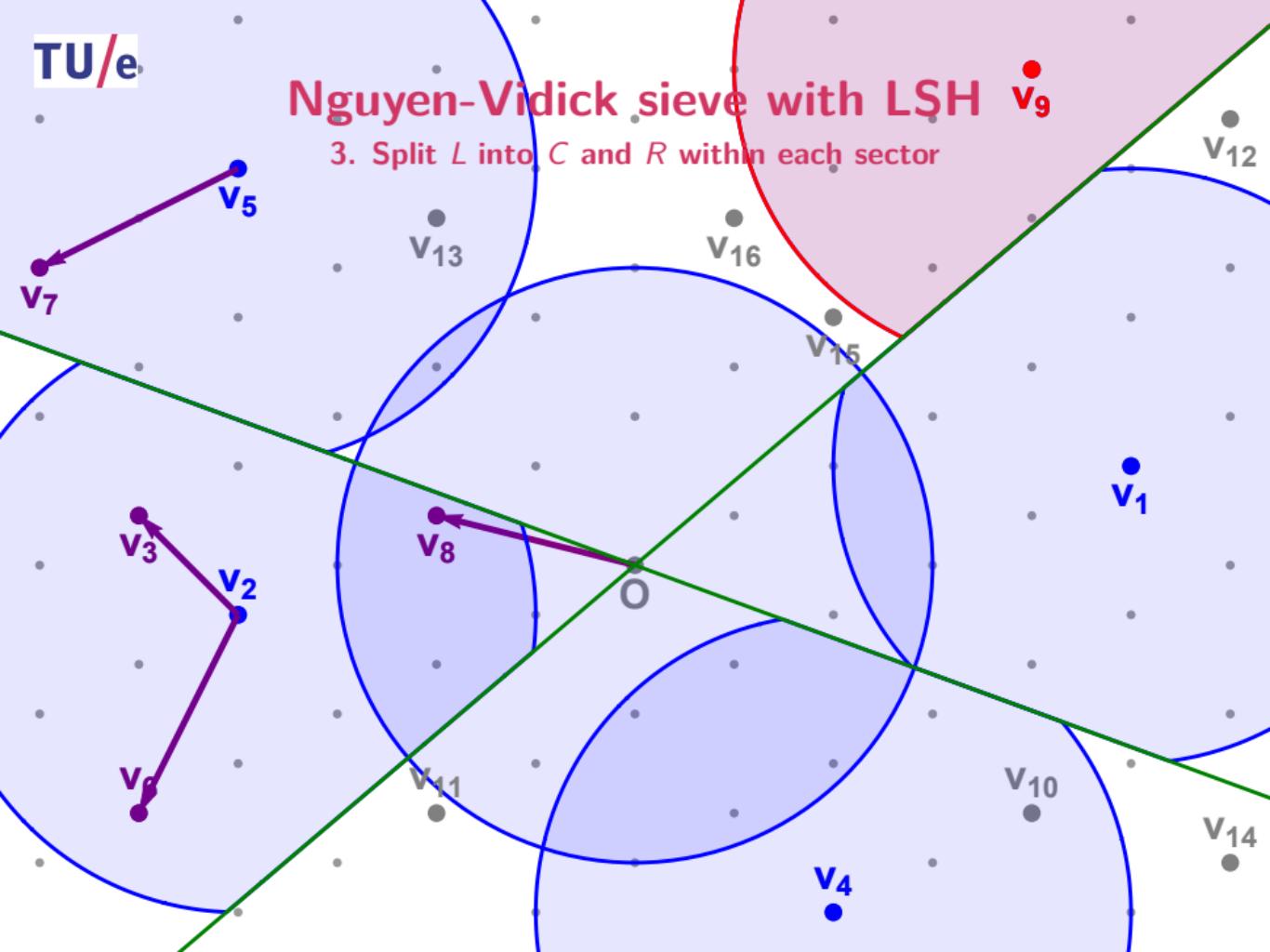
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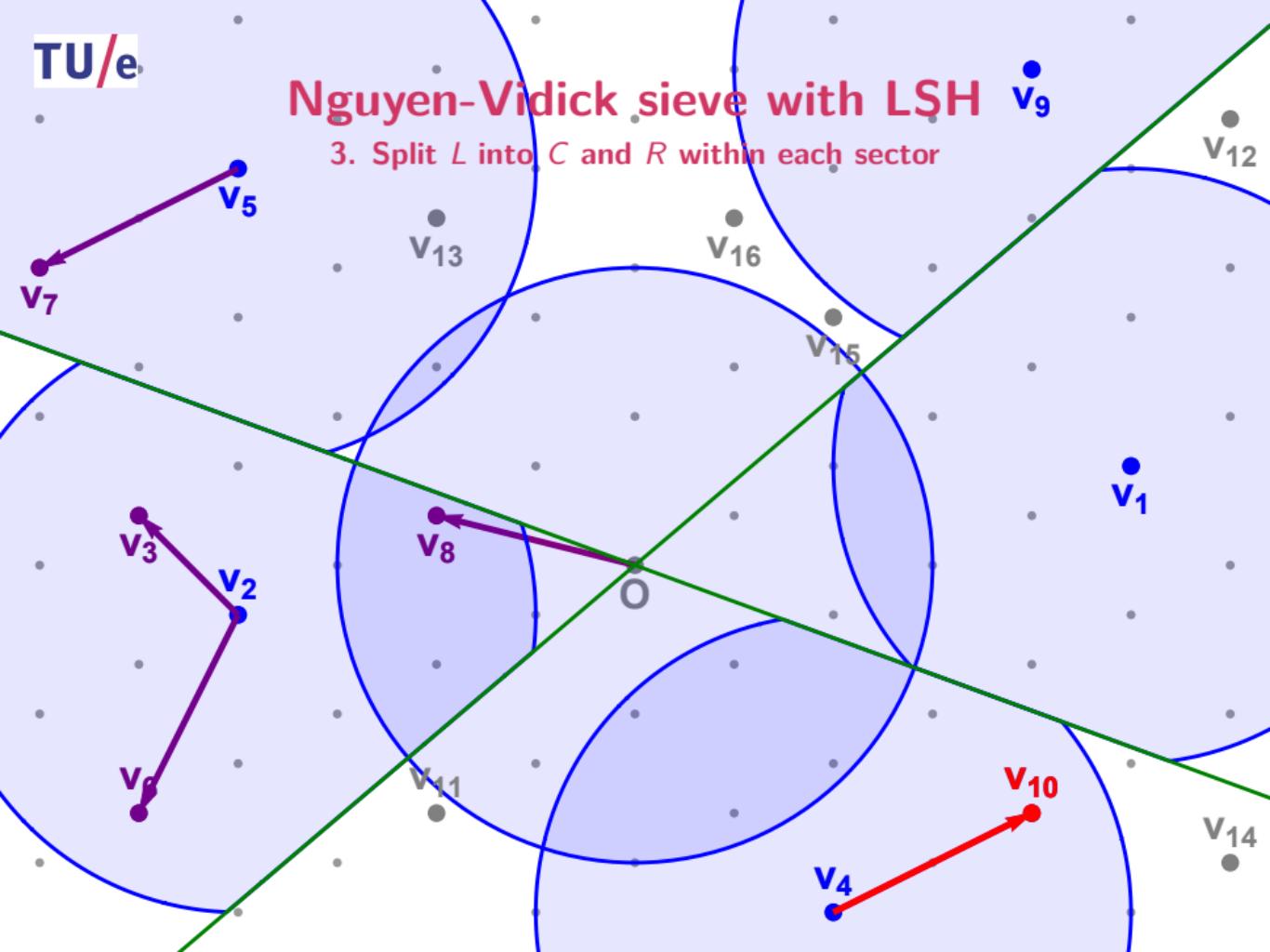
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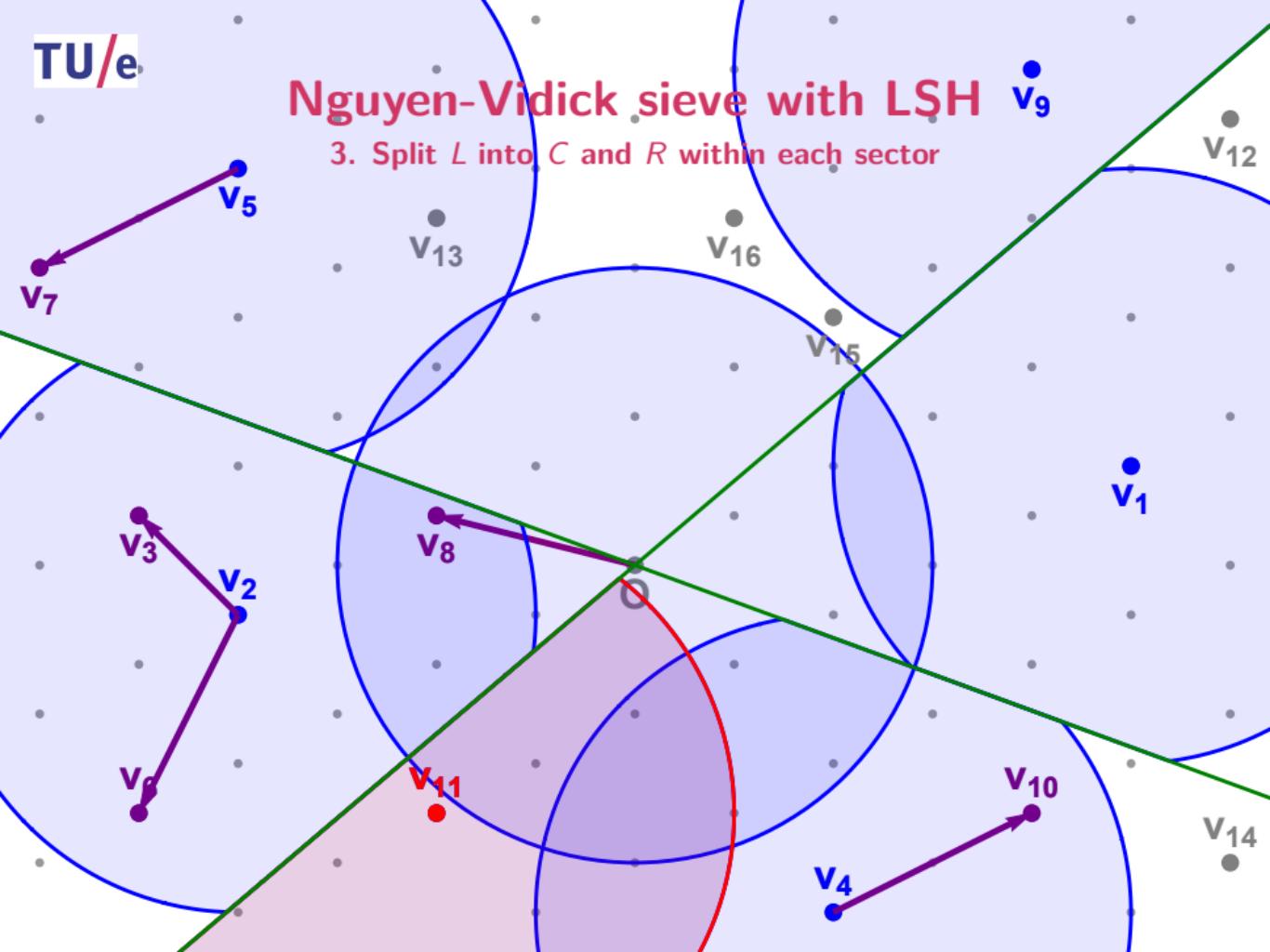
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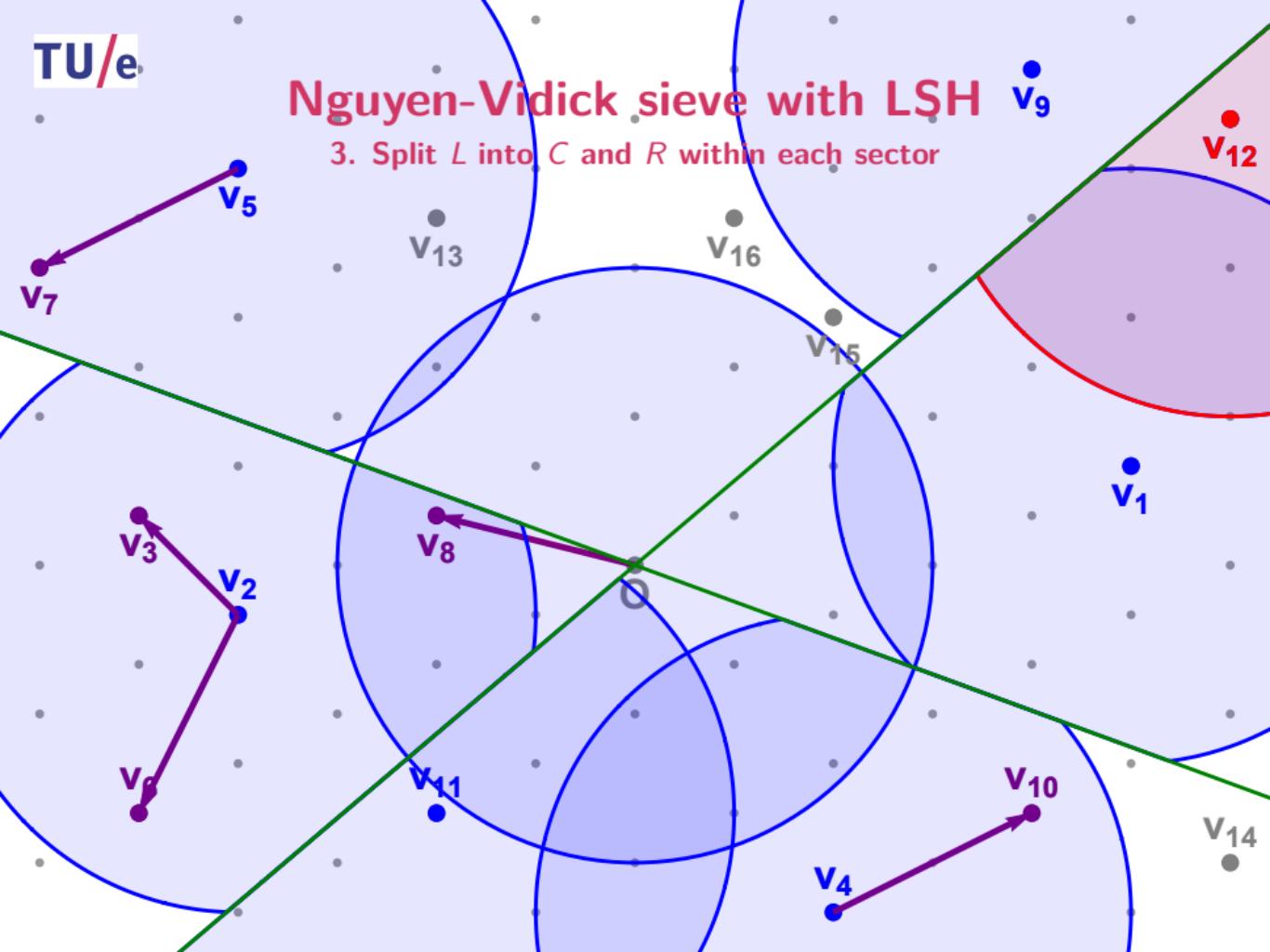
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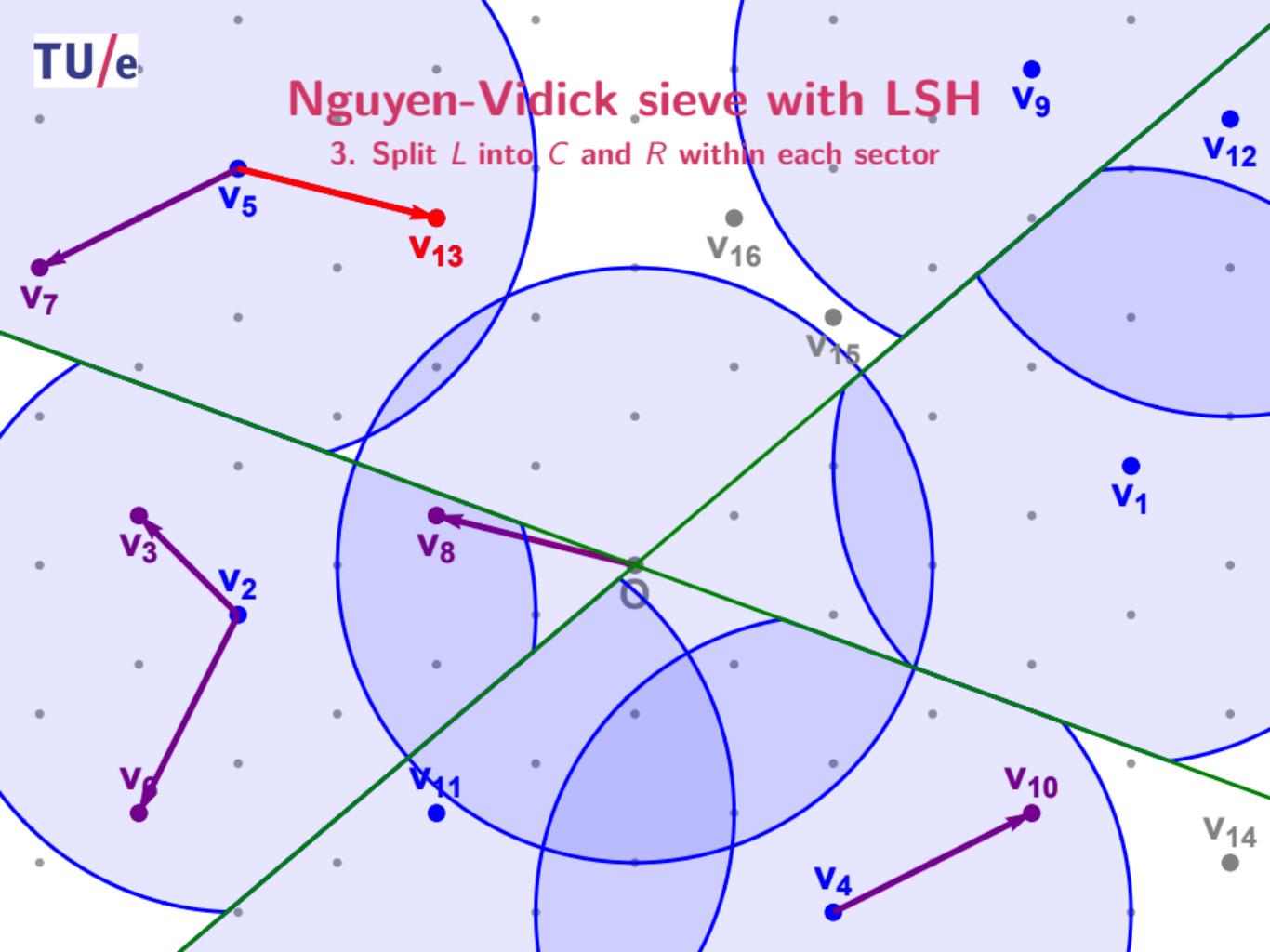
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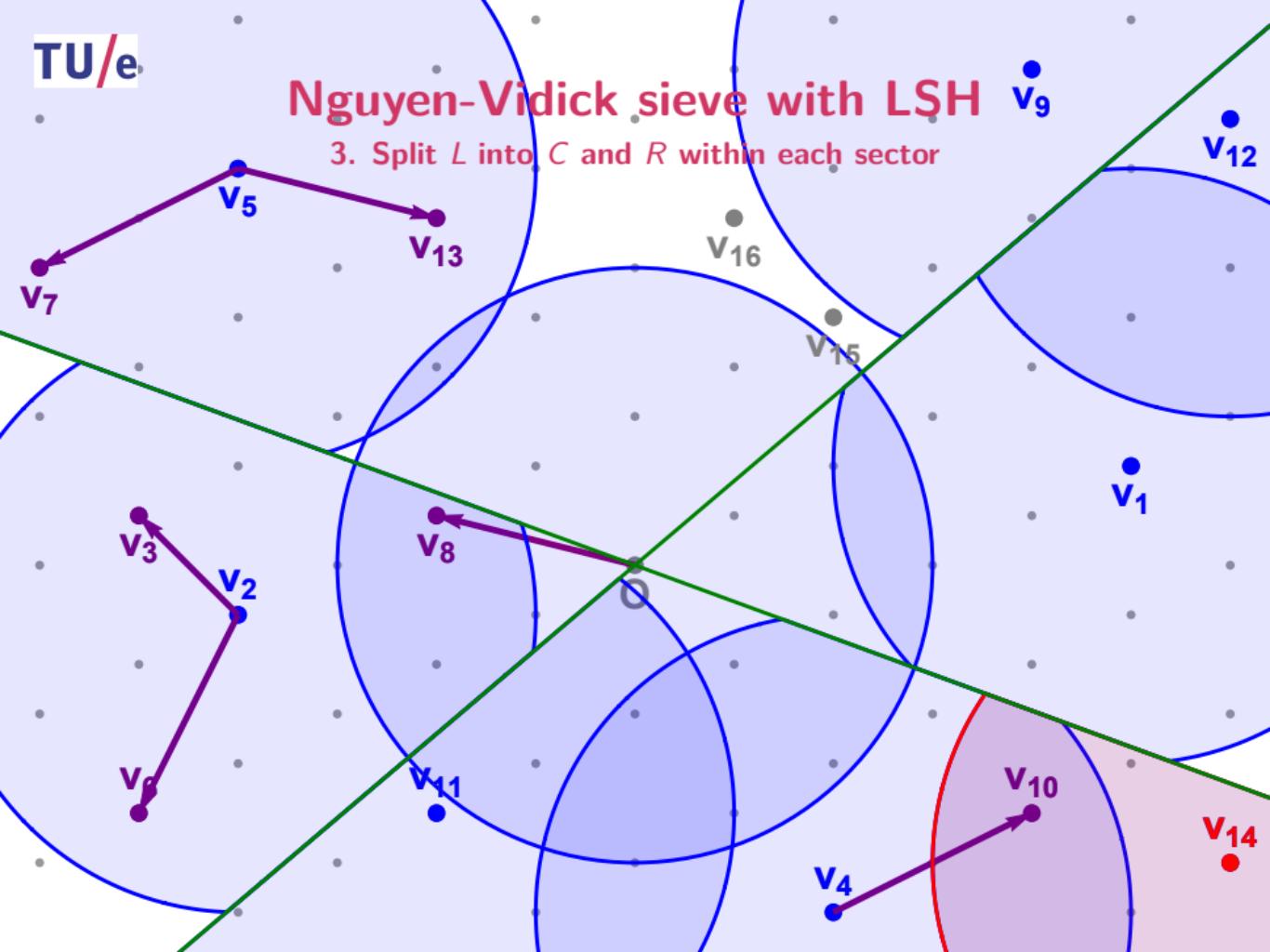
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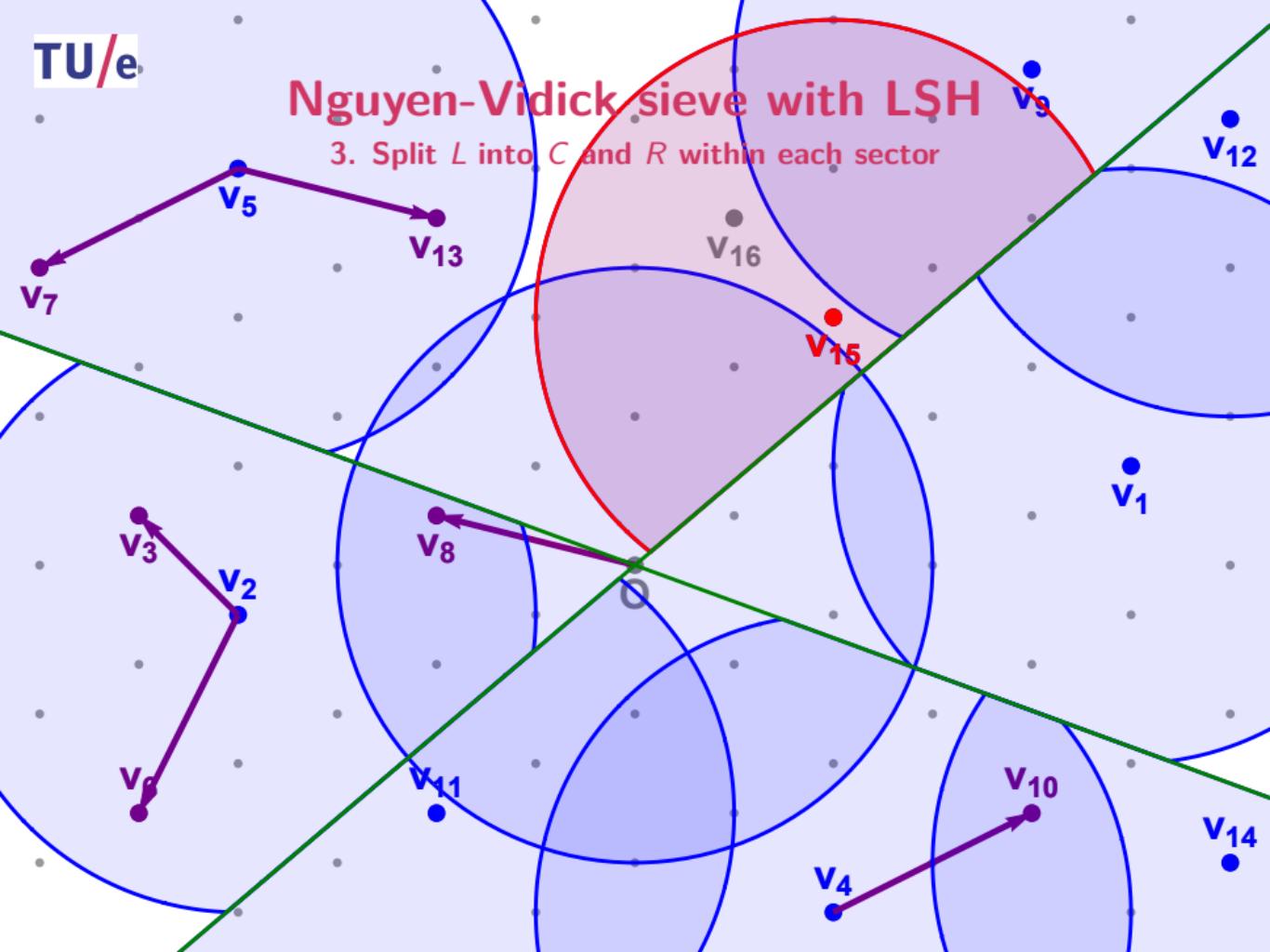
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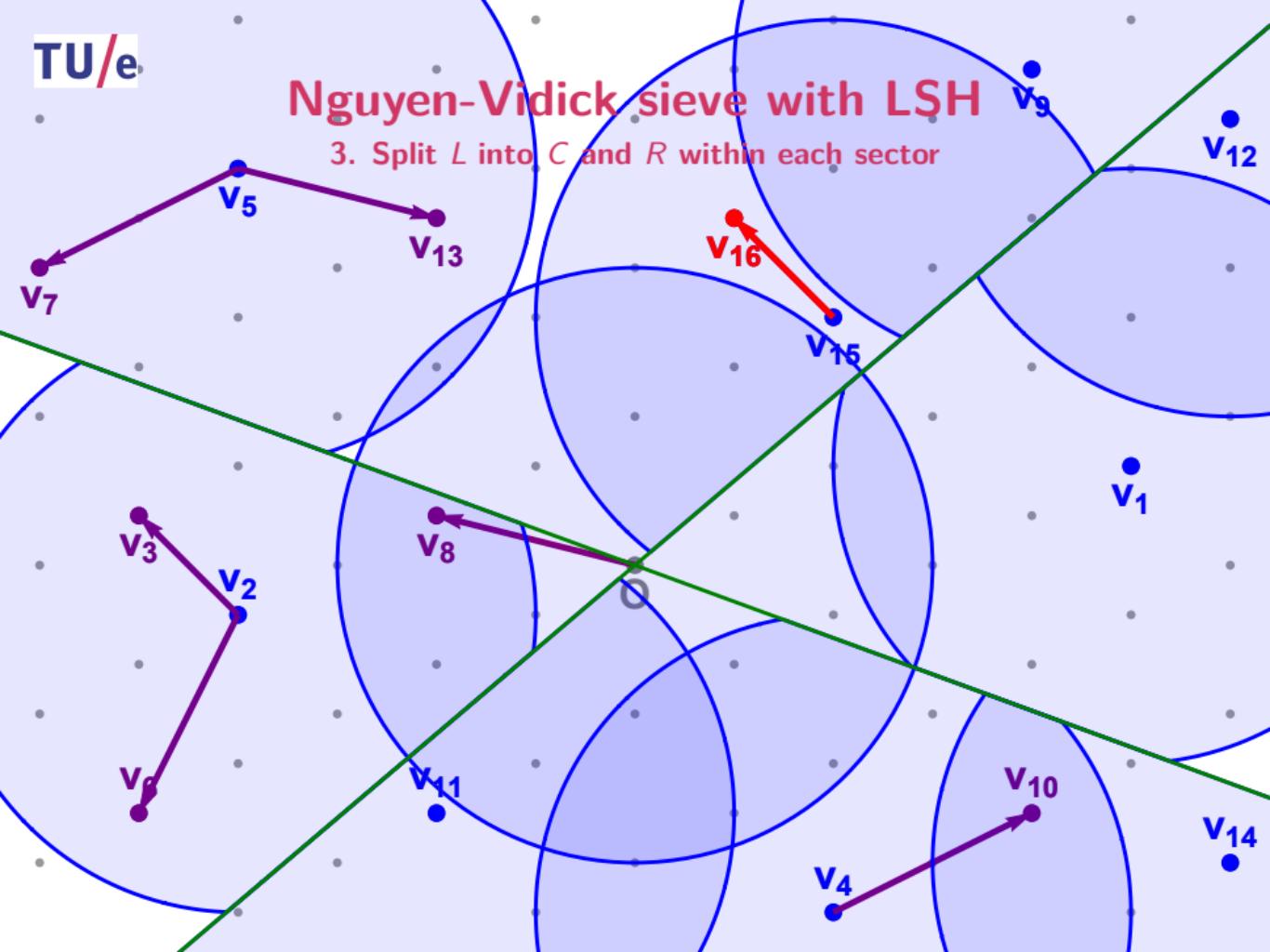
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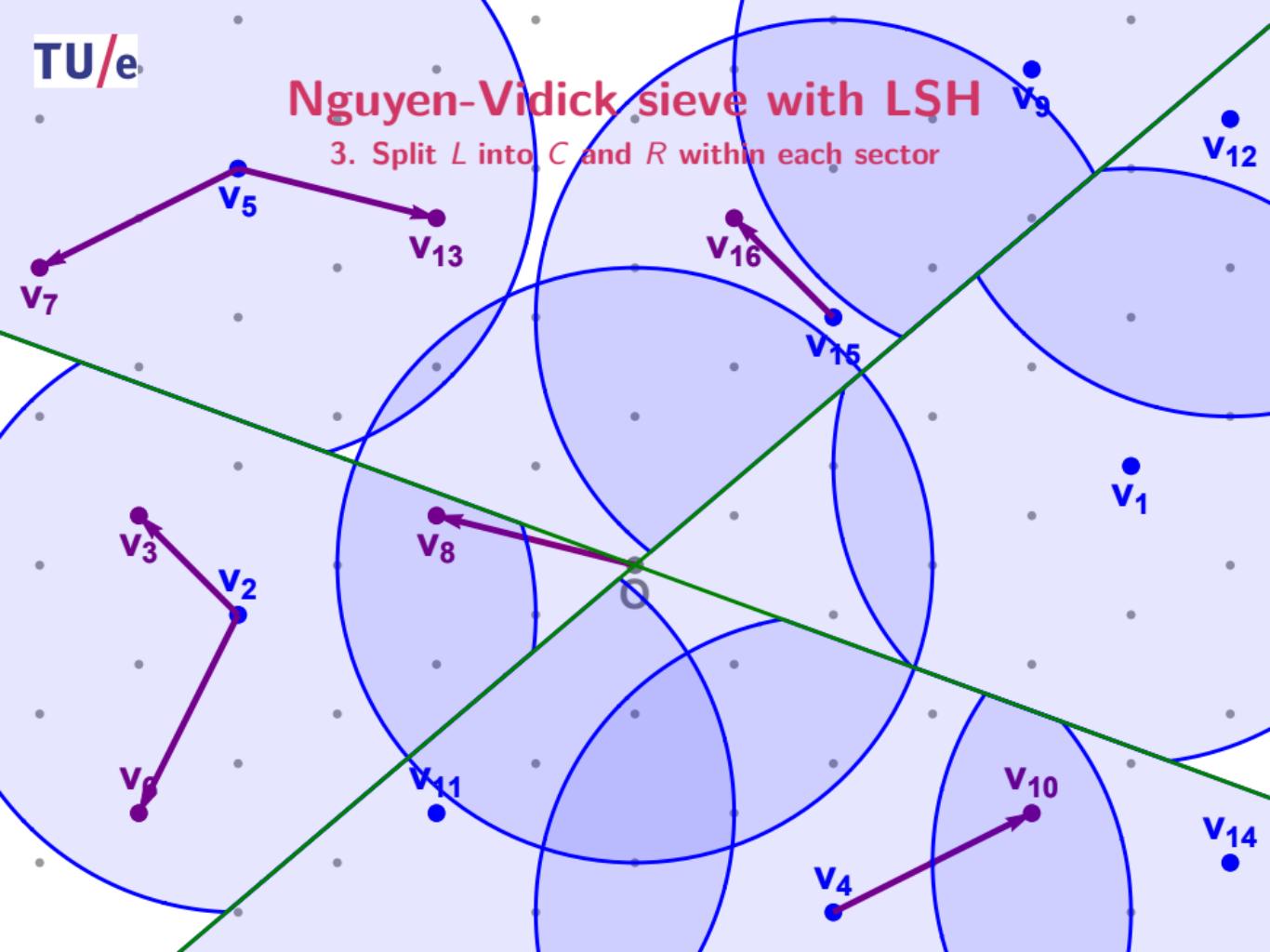
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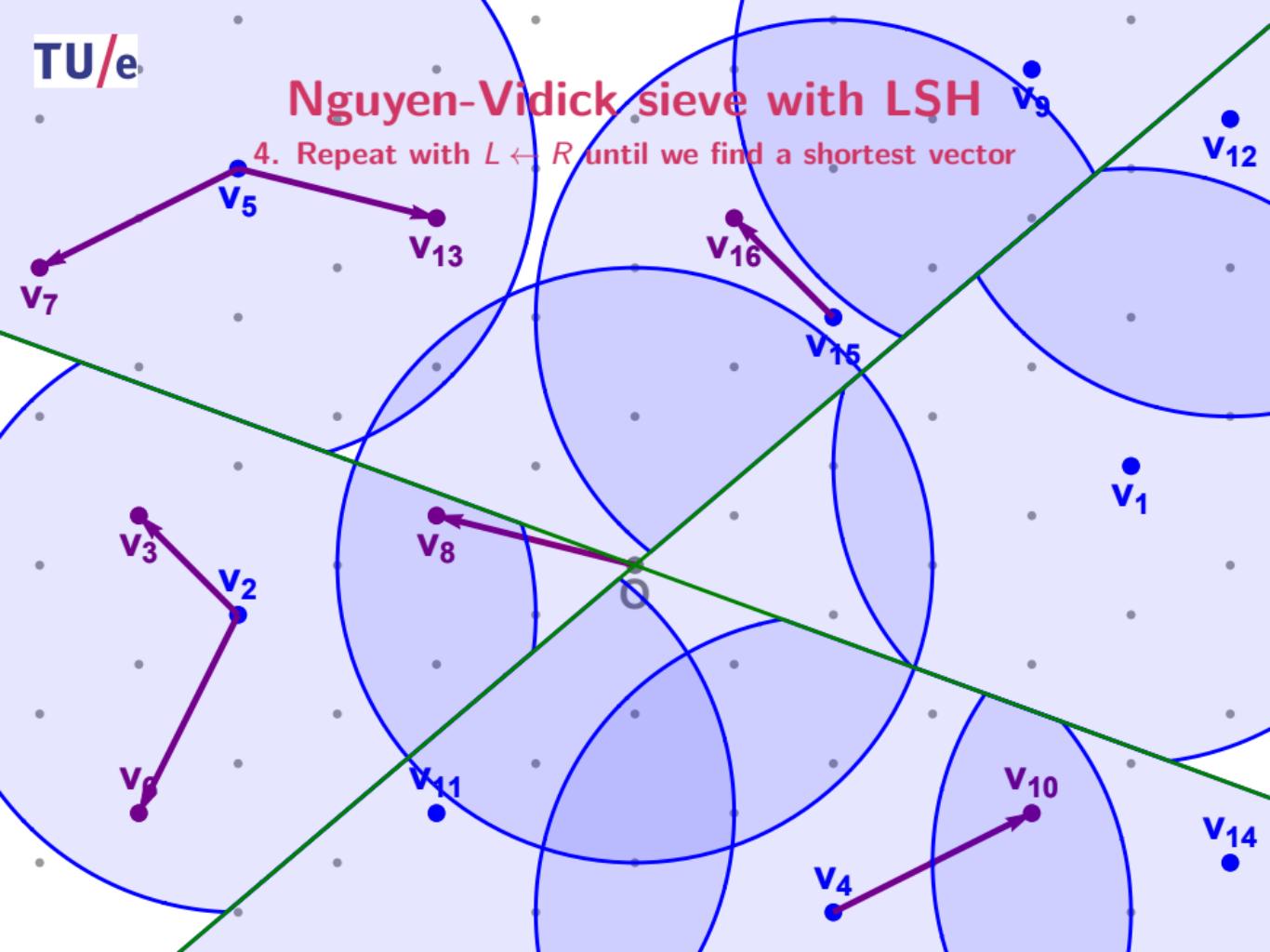
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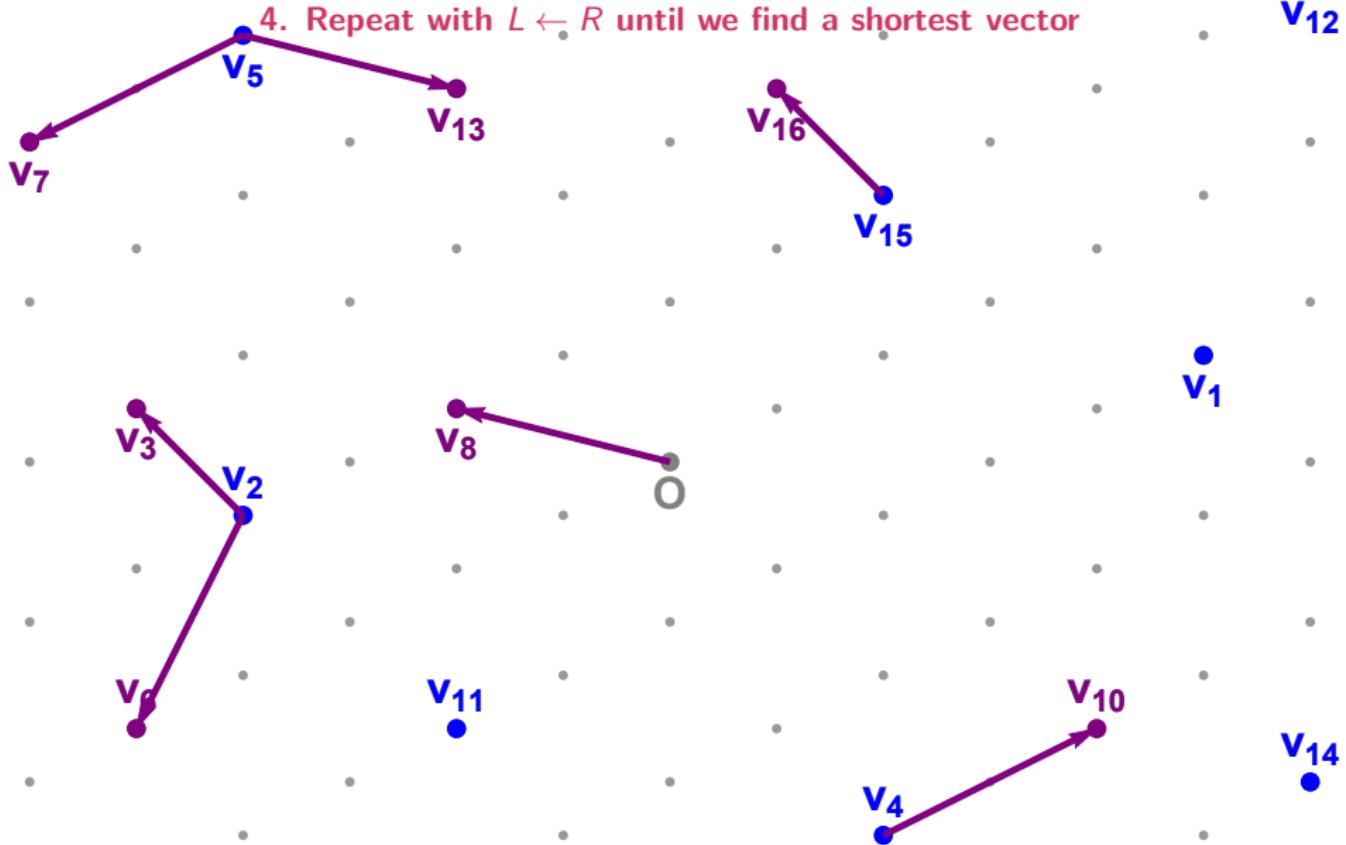


Nguyen-Vidick sieve with LSH

4. Repeat with $L \leftarrow R$ until we find a shortest vector

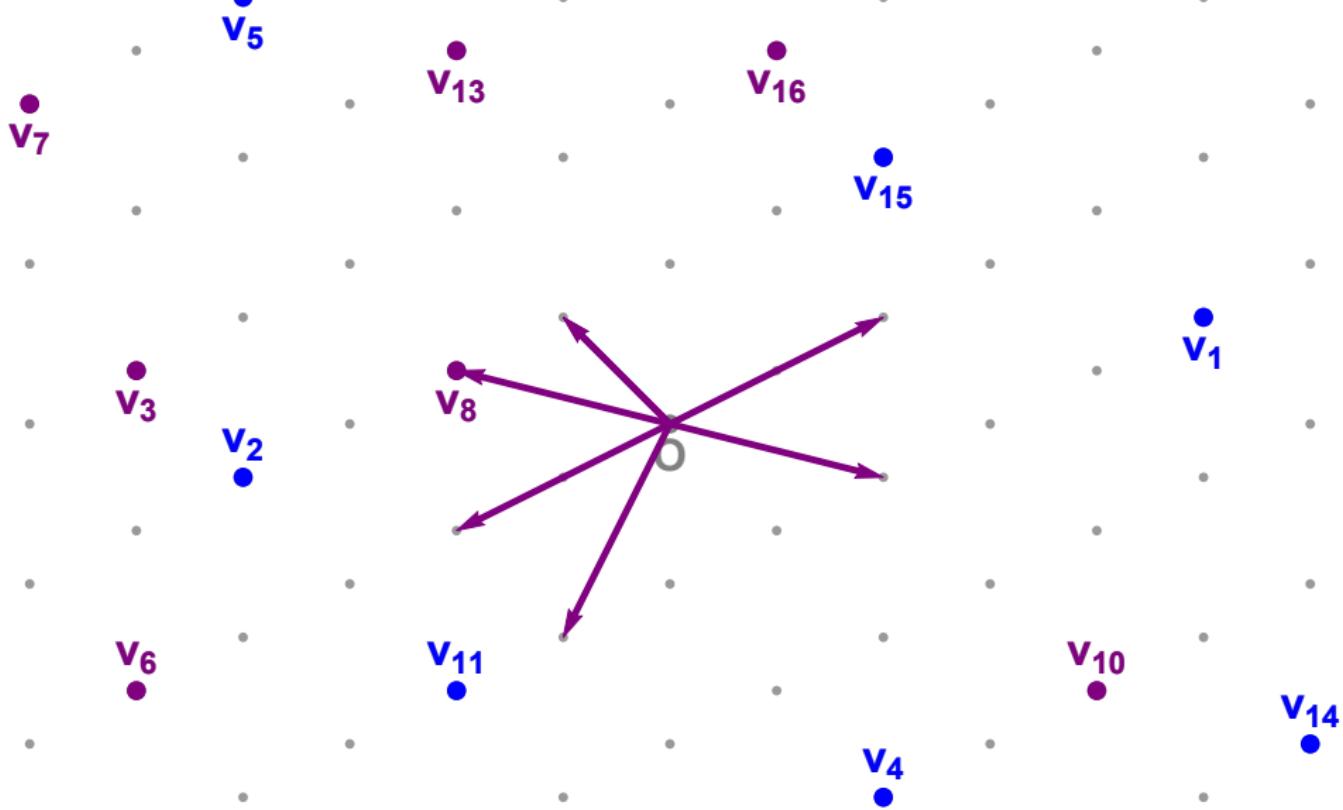


Nguyen-Vidick sieve with LSH

 v_9 v_{12} 4. Repeat with $L \leftarrow R$ until we find a shortest vector

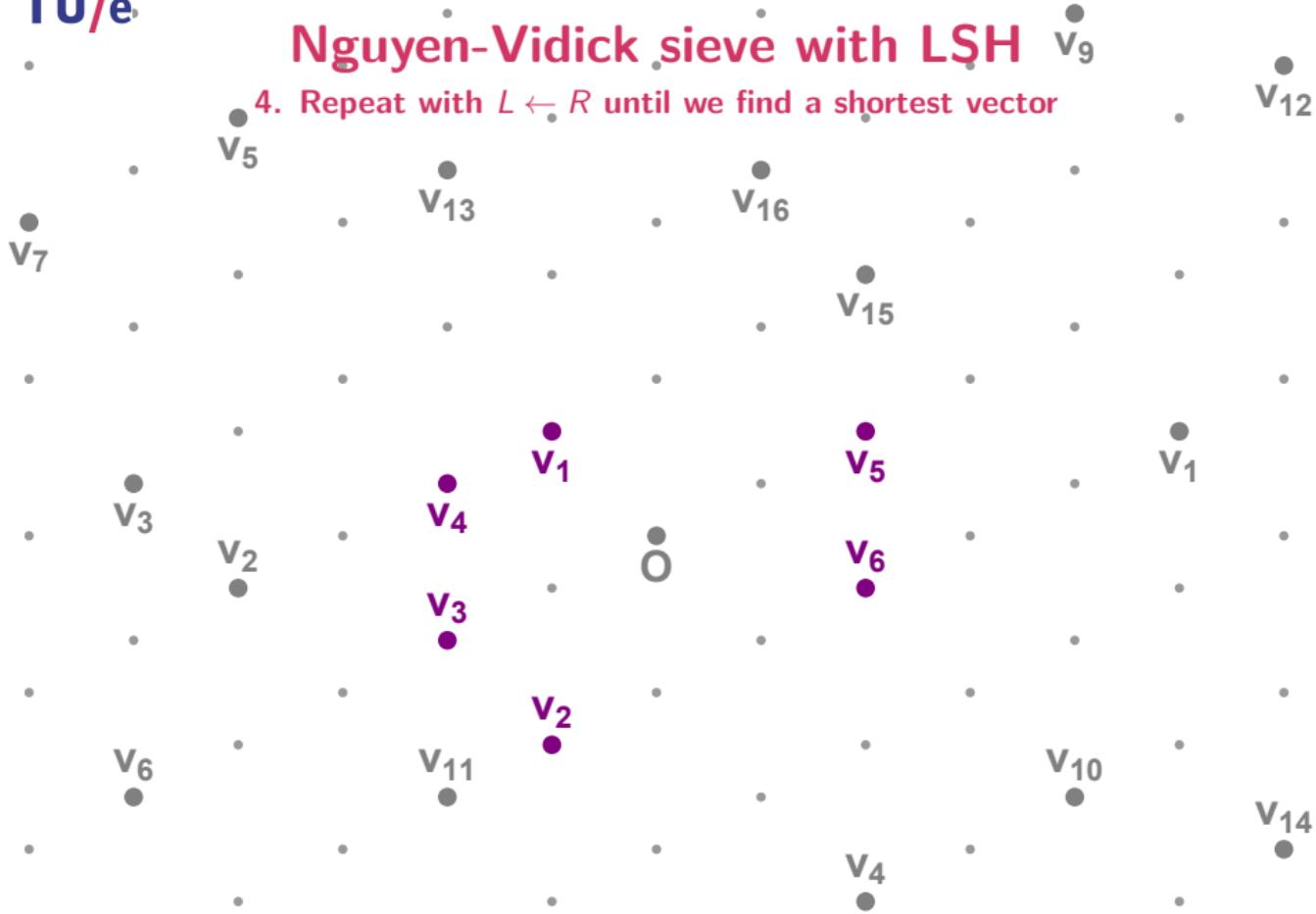
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Nguyen-Vidick sieve with LSH

Overview



Nguyen-Vidick sieve with LSH

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- Two parameters to tune
 - ▶ $k = O(n)$: Number of hyperplanes, leading to 2^k regions
 - ▶ $t = 2^{O(n)}$: Number of different, independent “hash tables”

Nguyen-Vidick sieve with LSH

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Nguyen-Vidick sieve with LSH

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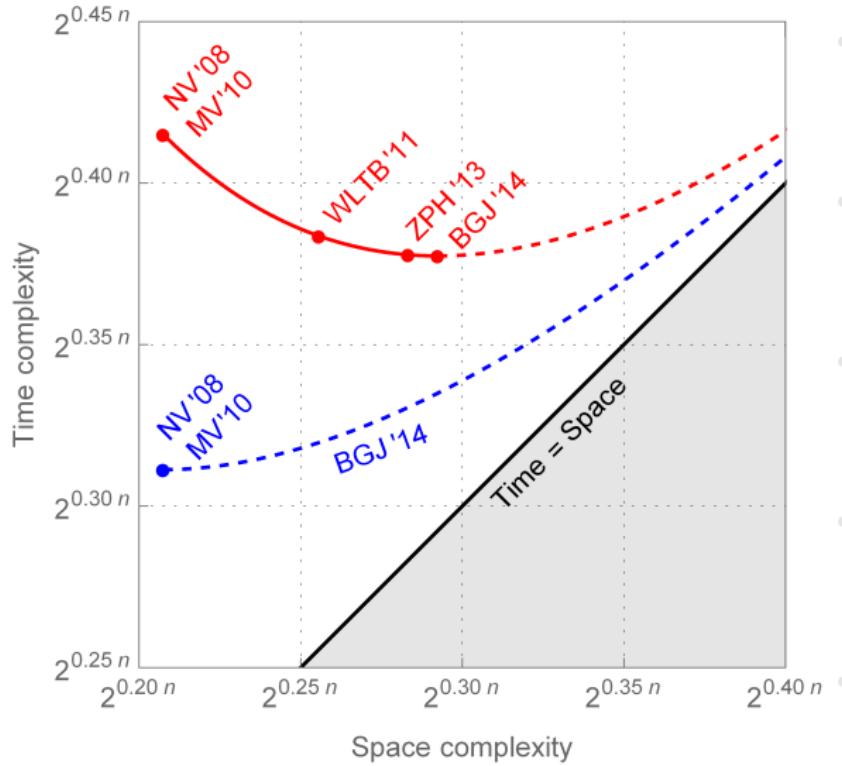
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Heuristic

Sieving with LSH runs in time $2^{0.34n+o(n)}$ and space $2^{0.34n+o(n)}$.

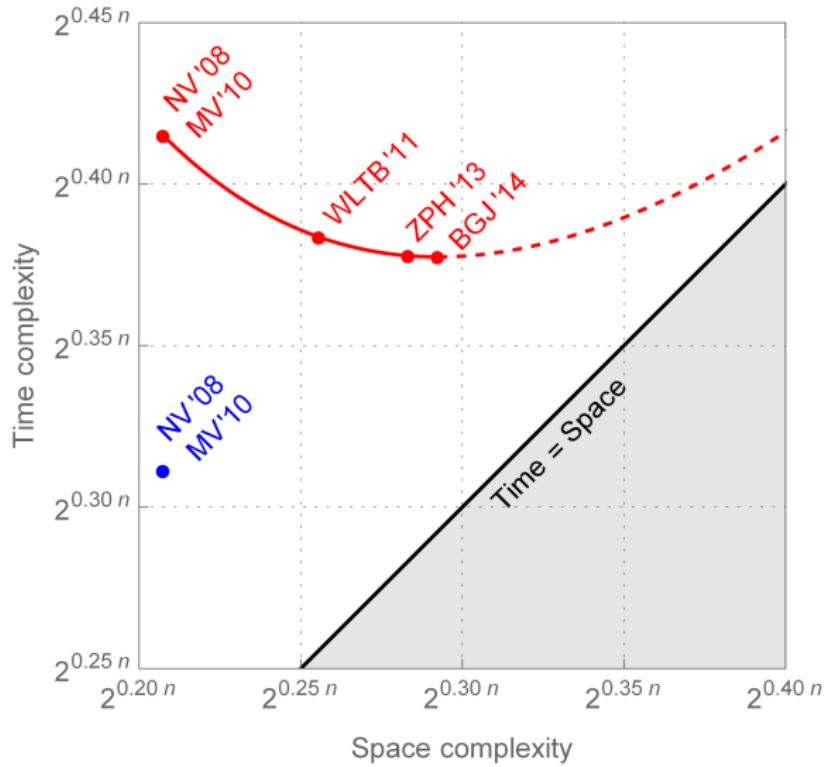
Sieving

Space/time trade-off



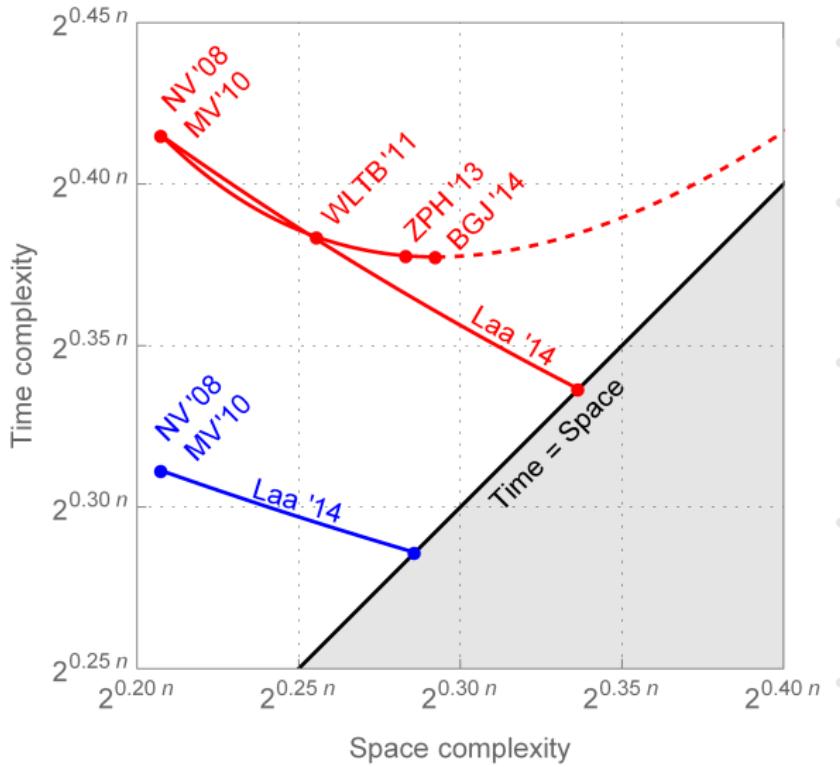
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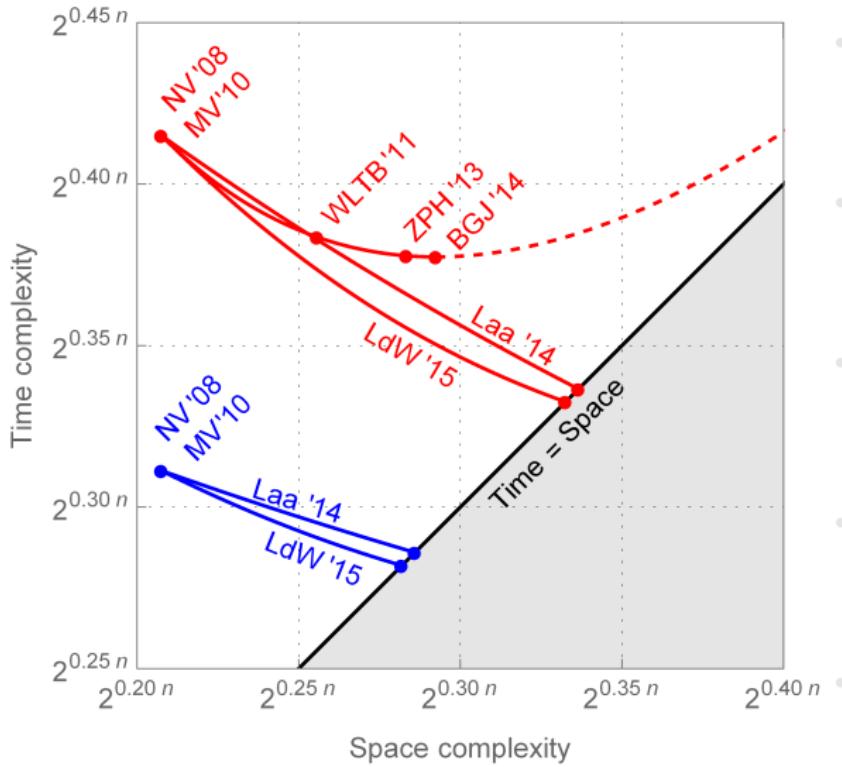
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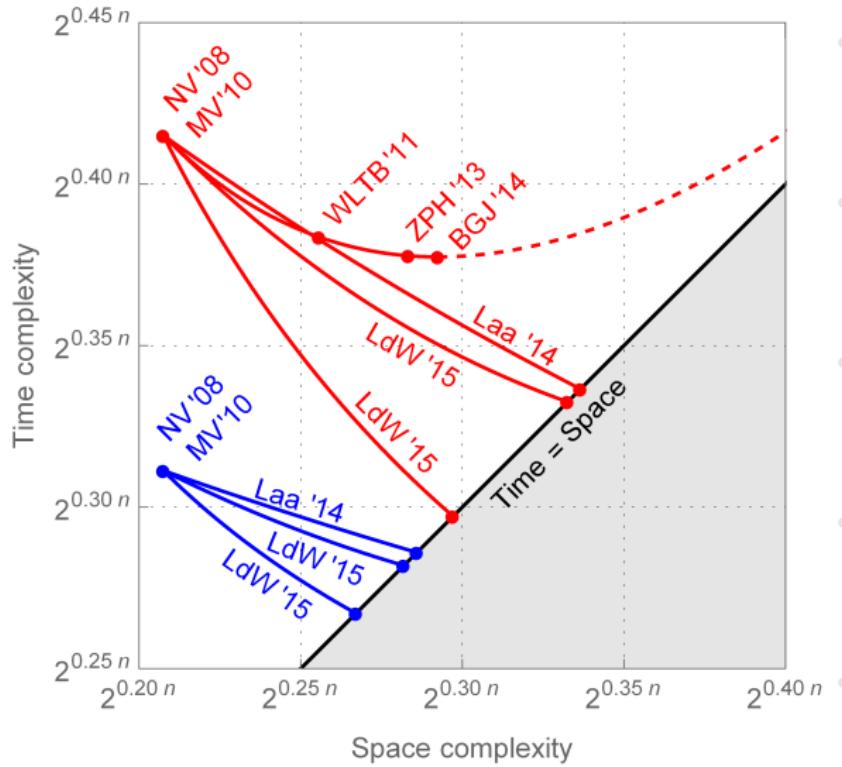
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Questions

