

Swarthmore College

cout << 1/0 << endl;

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24

a73310, 66 lines

2 Math Hints

Templates (1)

int cur = \sim li[0];

FOR(i,2,sz(li)) {

cur = ~next;

int dfs(int i) {

int next = addVar();

either(cur, ~li[i]);

either(~li[i], next);

vi val, comp, z; int time = 0;

if (low == val[i]) do {
 x = z.back(); z.pop_back();

comp[x] = low;

for(int e : qr[i]) if (!comp[e])

low = min(low, val[e] ?: dfs(e));

int low = val[i] = ++time, x; z.push_back(i);

either(cur, next);

either(cur, ~li[1]);

2sat.cpp

/* kactl 2-SAT Solver * Negated variables are represented by bit-inversions (\) $texttt\{\ tilde\{\}x\}\}$. * Usage: * TwoSat ts(number of boolean variables); * ts.either(0, \tilde3); // Var 0 is true or var 3 is * ts.setValue(2); // Var 2 is true * $ts.atMostOne(\{0, \forall tilde1, 2\}); // \le 1 \text{ of } vars 0, \forall tilde1$ and 2 are true * ts.solve(); // Returns true iff it is solvable * ts.values[0..N-1] holds the assigned values to the vars struct TwoSat { int N; vector<vi> gr; vi values; // 0 = false, 1 = true TwoSat(int n = 0) : N(n), qr(2*n) {} int addVar() { // (optional) gr.emplace_back(); gr.emplace_back(); return N++; void either(int f, int j) { $f = \max(2 * f, -1 - 2 * f);$ $j = \max(2*j, -1-2*j);$ gr[f].push_back(j^1); gr[j].push_back(f^1); void setValue(int x) { either(x, x); } void atMostOne(const vi& li) { // (optional) if (sz(li) <= 1) return;

```
if (values[x>>1] == -1)
        values[x>>1] = x&1;
    } while (x != i);
    return val[i] = low;
  bool solve() {
    values.assign(N, -1);
    val.assign(2*N, 0); comp = val;
    FOR(i, 0, 2*N) if (!comp[i]) dfs(i);
    FOR(i, 0, N) if (comp[2*i] == comp[2*i+1]) return 0;
    return 1;
};
3dHull.cpp
"Point3D.h"
                                                   de4d99, 49 lines
typedef Point3D<double> P3;
struct PR {
  void ins(int x) { (a == -1 ? a : b) = x; }
  void rem(int x) { (a == x ? a : b) = -1; }
 int cnt() { return (a != -1) + (b != -1); }
 int a, b;
};
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vector<P3>& A) {
 assert (sz(A) >= 4);
 vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
#define E(x,y) E[f.x][f.y]
 vector<F> FS;
  auto mf = [\&] (int i, int j, int k, int l) {
   P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
   if (q.dot(A[1]) > q.dot(A[i]))
      q = q * -1;
   F f{q, i, j, k};
    E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
    FS.push_back(f);
  FOR(i, 0, 4) rep(j, i+1, 4) rep(k, j+1, 4)
   mf(i, j, k, 6 - i - j - k);
  FOR(i, 4, sz(A)) {
   FOR(j,0,sz(FS)) {
      F f = FS[j];
      if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
       E(a,b).rem(f.c);
        E(a,c).rem(f.b);
        E(b,c).rem(f.a);
        swap(FS[j--], FS.back());
        FS.pop_back();
    int nw = sz(FS);
    FOR(i,0,nw) {
      F f = FS[i];
\#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i, f.
      C(a, b, c); C(a, c, b); C(b, c, a);
 trav(it, FS) if ((A[it.b] - A[it.a]).cross(
   A[it.c] - A[it.a]).dot(it.q) \ll 0) swap(it.c, it.b);
```

```
return FS;
};
AdvancedHash.cpp
                                                  395fc0, 86 lines
11 base1[MX], base2[MX];
int base;
11 b1Inv, b2Inv;
11 baseExp(ll power, ll prime) {
       if (power == 0) {
           return 1:
       } else {
           11 cur = baseExp(power / 2, prime); cur = cur *
                cur; cur = cur % prime;
           if (power % 2 == 1) cur = cur * base;
           cur = cur % prime;
            return cur;
   }
struct hsh
   11 p1, p2;
   ll val1, val2;
   vl val1s, val2s;
   vl nums;
   hsh() {
        p1 = 1000000007;
        p2 = 1000000009;
       val1 = 0;
        val2 = 0;
       val1s.pb(0); val2s.pb(0);
   hsh(string S) {
       p1 = 1000000007;
        p2 = 1000000009;
        val1 = 0;
       val2 = 0;
        val1s.pb(0); val2s.pb(0);
        FOR(i, sz(S)) {
           push_back(S[i] - 'a' + 1);
    void push_back(ll v) {
       v++;
        val1 *= base;
        val1 += v;
        val1 %= p1;
        val2 *= base;
        val2 += v:
       val2 %= p2;
    valls.pb(vall);
    val2s.pb(val2);
        nums.pb(v);
   void pop_back() {
```

```
//warning: the extent to which this has been tested
             is limited lol
       11 v = nums[sz(nums)-1]; nums.pop_back();
       val1s.pop_back(); val2s.pop_back();
       val1 += p1 - v; val1 *= b1Inv;
       val1 %= p1;
        val2 += p2 - v; val2 *= b2Inv; val2 %= p2;
   11 get(int L, int R) {
       11 A = (val1s[R+1] - (val1s[L] * base1[R-L+1]) % p1
             + p1) % p1;
       11 B = (val2s[R+1] - (val2s[L] * base2[R-L+1]) % p2
            + p2) % p2;
        return A * p2 + B;
};
void prepHash() {
    base = uniform_int_distribution<int>(1000, MOD-2)(rng);
    base1[0] = 1; base2[0] = 1;
   blInv = baseExp(1000000005, 1000000007);
   b2Inv = baseExp(1000000007, 1000000009);
    FOR(i, 1, MX) {
       base1[i] = (base1[i-1] * base) % 1000000007;
       base2[i] = (base2[i-1] * base) % 1000000009;
```

AhoCorasick.cpp

```
f84242, 78 lines
struct AhoCorasick {
    static const int K = 26;
    struct Vertex {
        int next[K];
       bool leaf = false;
       int p = -1;
        char pch;
       int link = -1;
       int go[K];
       11 \text{ value} = -1;
       11 \text{ num} = 0;
        Vertex(int p=-1, char ch='\$') : p(p), pch(ch) {
            fill(begin(next), end(next), -1);
            fill(begin(go), end(go), -1);
    };
    vector<Vertex> t;
    void init() {
        Vertex v; t.pb(v);
    void add_string(string const& s, int count) {
        int v = 0;
        for (char ch : s) {
            int c = ch - 'a';
            if (t[v].next[c] == -1) {
                t[v].next[c] = t.size();
                t.emplace_back(v, ch);
```

```
v = t[v].next[c];
        t[v].leaf = true;
        t[v].num = count;
    ll dfs(int v) {
        if (t[v].value != -1) {
            return t[v].value;
        ll ans = t[v].num;
        ans += dfs(get_link(v)); // is this right?
        return t[v].value = ans;
    void compute() {
        t[0].value = 0;
        FOR(i, 1, sz(t)) {
            dfs(i);
    int get_link(int v) {
        if (t[v].link == -1) {
            if (v == 0 \mid | t[v].p == 0)
                t[v].link = 0;
                t[v].link = go(get_link(t[v].p), t[v].pch);
        return t[v].link;
    int go(int v, char ch) {
        int c = ch - 'a';
        if (t[v].go[c] == -1) {
            if (t[v].next[c] != -1)
                t[v].qo[c] = t[v].next[c];
                t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch
                     );
        return t[v].go[c];
};
```

Angle.cpp

0f0602, 35 lines

```
struct Angle {
  int x, y;
  Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
  Angle operator-(Angle b) const { return {x-b.x, y-b.y, t};
  int half() const {
   assert(x || v);
    return y < 0 \mid | (y == 0 \&\& x < 0);
  Angle t90() const { return \{-y, x, t + (half() \&\& x \ge 0)\}
  Angle t180() const { return \{-x, -y, t + half()\}; }
  Angle t360() const { return \{x, y, t + 1\}; }
bool operator<(Angle a, Angle b) {
  // add a.dist2() and b.dist2() to also compare distances
  return make_tuple(a.t, a.half(), a.y * (ll)b.x) <</pre>
```

```
make_tuple(b.t, b.half(), a.x * (ll)b.y);
// Given two points, this calculates the smallest angle
// them, i.e., the angle that covers the defined line
    seament.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
 if (b < a) swap(a, b);</pre>
 return (b < a.t180() ?
         make_pair(a, b) : make_pair(b, a.t360()));
Angle operator+(Angle a, Angle b) { // point a + vector b
 Angle r(a.x + b.x, a.y + b.y, a.t);
 if (a.t180() < r) r.t--;
 return r.t180() < a ? r.t360() : r;
Angle angleDiff(Angle a, Angle b) { // angle b - angle a}
 int tu = b.t - a.t; a.t = b.t;
 return \{a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)
```

ArticulationPoints.cpp

23f413, 37 lines

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited:
vector<int> tin, low;
int timer;
void dfs(int v, int p = -1) {
   visited[v] = true;
   tin[v] = low[v] = timer++;
    int children=0;
    for (int to : adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
        } else {
            dfs(to, v);
           low[v] = min(low[v], low[to]);
           if (low[to] >= tin[v] && p!=-1)
                IS CUTPOINT (v);
            ++children:
    if(p == -1 && children > 1)
        IS CUTPOINT (v);
void find_cutpoints() {
   timer = 0;
    visited.assign(n, false);
   tin.assign(n, -1);
   low.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs (i);
```

```
Basis.cpp
                                                   075e70, 30 lines
const int D; //length of masks
11 basis[D]; // basis[i] keeps the mask of the vector whose
    f value is i
int bs = 0; //basis\ size
void insertVector(ll mask) {
  for (int i = 0; i < D; i++) {
    if ((mask & 111 << i) == 0) continue;
    if (!basis[i]) {
     basis[i] = mask;
      ++bs;
      return;
    mask ^= basis[i];
bool inSpan(ll mask) {
  for (int i = 0; i < D; i++) {
    if ((mask & 111 << i) == 0) continue;
    if (!basis[i]) {
           return true;
    mask ^= basis[i];
    return false:
```

Berlekamp-Massev.cpp

```
vector<int> berlekamp_massey(const vector<int> &a) {
 vector<int> v, last; //v is the answer, 0-based, p is
      the module
 int k = -1, delta = 0;
 for (int i = 0; i < (int)a.size(); i++) {
   int tmp = 0;
    for (int j = 0; j < (int)v.size(); <math>j++)
     tmp = (tmp + (long long)a[i - j - 1] * v[j]) % p;
   if (a[i] == tmp) continue;
    if (k < 0) {
     k = i;
     delta = (a[i] - tmp + p) % p;
     v = vector < int > (i + 1);
      continue;
    vector<int> u = v;
    int val = (long long)(a[i] - tmp + p) * power(delta, p -
    if (v.size() < last.size() + i - k) v.resize(last.size()</pre>
         + i - k);
    (v[i - k - 1] += val) %= p;
    for (int j = 0; j < (int) last.size(); <math>j++) {
```

```
v[i - k + j] = (v[i - k + j] - (long long)val * last[j
         ]) % p;
    if (v[i - k + j] < 0) v[i - k + j] += p;
}

if ((int)u.size() - i < (int)last.size() - k) {
    last = u;
        k = i;
    delta = a[i] - tmp;
    if (delta < 0) delta += p;
}

for (auto &x : v) x = (p - x) % p;
v.insert(v.begin(), 1);

return v; // $\forall i, \sum_{f} = 0} ^ m a_{f} - j} v_{j} =
    0$
}</pre>
```

BinSearchSegtree.cpp

853b41, 54 lines

```
const ll identity = 0;
const 11 SZ = 131072;
11 sum[2*SZ], lazy[2*SZ];
ll combine(ll A, ll B) {
   return A+B:
ll combineUpd(ll A, ll B) {
   return A+B;
void push(int index, 11 L, 11 R) {
   sum[index] = combineUpd(sum[index], lazy[index]);
   if (L != R) lazy[2*index] = combineUpd(lazy[2*index],
        lazy[index]), lazy[2*index+1] = combineUpd(lazy[2*
        index+1], lazy[index]);
   lazv[index] = identity;
void pull(int index) {
   sum[index] = combine(sum[2*index], sum[2*index+1]);
bool checkCondition(int index) {
   //FILL THIS IN
11 query (int lo = 0, int hi = SZ-1, int index = 1, 11 L = 0,
     11 R = SZ-1) { //returns first node satisfying con
   push (index, L, R);
   if (lo > R || L > hi) return -1;
   bool condition = checkCondition(index);
        return (condition ? L : -1);
   int M = (L+R) / 2;
   if (checkCondition(2*index)) {
        return query(lo, hi, 2*index, L, M);
   return query(lo, hi, 2*index+1, M+1, R);
```

BipartiteMatching.cpp

7a6c4b, 24 lines

```
//Storing the graph
vector<int> g[maxn];
//Storing whether we have visited a node
bool vis[maxn];
//Storing the vertex matched to
int match[maxn];
bool hungarian(int u){
 for (int i = 0; i < q[u].size(); ++i) {
   int v = q[u][i];
   if (!vis[v]){
     vis[v] = true;
     if (!match[v] || hungarian(match[v])){
       match[u] = v; match[v] = u; return true;
 return false;
//in main: call hungarian for each vertex on one side
for (int i = 1;i <= n1;++i) {
   memset (vis, false, sizeof vis);
   if (hungarian(i)) ans++; //if we can match i
```

BipartiteMatchingWithWeights.cpp

c21d5f, 74 lines

```
while (head <= tail) {
      int u = q[head++];
      visx[u] = true;
      for (int v = 1; v \le n; ++v) {
       if (!visy[v]){
          if (a[u] + b[v] == g[u][v])
            visy[v] = true; pre[v] = u;
            if (!fy[v]){
              augment(v); return;
            g[++tail] = fy[v];continue;
          if (slack[v] > a[u] + b[v] - g[u][v]){
            slack[v] = a[u] + b[v] - q[u][v];
            pre[v] = u;
    11 d = inf;
    for (int i = 1; i \le n; ++i) {
     if (!visy[i]) d = min(d, slack[i]);
    for (int i = 1; i \le n; ++i) {
      if (visx[i]) a[i] -= d;
      if (visy[i]) b[i] += d;
      else slack[i] -= d;
    for (int v = 1; v \le n; ++v) {
      if (!visv[v] && !slack[v]) {
        visy[v] = true;
        if (!fv[v]) {
         augment (v);
          return;
        q[++tail] = fy[v];
11 km(){
  for (int i = 1; i \le n; ++i) {
   a[i] = -inf;
   b[i] = 0;
   for (int j = 1; j \le n; ++j) a[i] = max(a[i], g[i][j]);
 memset(fx, 0, sizeof fx);
 memset(fy, 0, sizeof fy);
  for (int i = 1; i \le n; ++i) bfs(i);
  11 \text{ ans} = 0;
  for (int i = 1; i \le n; ++i) ans += a[i] + b[i];
  //vertex i on left is matched to g2[i][fx[i]] * fx[i]
  //g2[a][b]=1 iff exists edge ab
 return ans;
```

BIT.cpp

992848, 14 lines

```
const int maxn = 30005;
int n, bit[maxn];
void add(int i, int x) {
   for (; i <= n; i += i & (-i))
       bit[i] += x;</pre>
```

```
int sum(int i) {
   int r = 0;
   for (; i; i -= i & (-i)) {
      r += bit[i];
   }
   return r;
}
```

Bridge.cpp

a44485, 33 lines

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, low;
int timer;
void dfs(int v, int p = -1) {
   visited[v] = true;
   tin[v] = low[v] = timer++;
    for (int to : adj[v]) {
       if (to == p) continue;
        if (visited[tol) {
           low[v] = min(low[v], tin[to]);
       } else {
            dfs(to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] > tin[v])
                IS_BRIDGE(v, to);
void find_bridges() {
   timer = 0:
   visited.assign(n, false);
   tin.assign(n, -1);
   low.assign(n, -1);
   for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs(i);
```

CentroidDecomp.cpp

b8fb48 54

```
struct CentroidDecomposition {
  vector<set<int>> tree; // it's not vector<vector<int>>!
  vector<int> dad;
  vector<int> sub;
   vector<int> dep;

CentroidDecomposition(vector<set<int>> &tree) : tree(tree)
   {
  int n = tree.size();
  dad.resize(n);
  sub.resize(n);
  build(0, -1);
}

void build(int u, int p) {
  int n = dfs(u, p); // find the size of each subtree
```

```
int centroid = dfs(u, p, n); // find the centroid
       if (p == -1) {
            dep[centroid] = 0;
        } else {
            dep[centroid] = dep[p] + 1;
    if (p == -1) p = centroid; // dad of root is the root
        itself
    dad[centroid] = p;
    // for each tree resulting from the removal of the
        while (!tree[centroid].empty()) {
            int v = *(tree[centroid].begin());
      tree[centroid].erase(v); // remove the edge to
      tree[v].erase(centroid); // the component from the
          tree
     build(v, centroid);
 int dfs(int u, int p) {
   sub[u] = 1;
   for (auto v : tree[u])
     if (v != p) sub[u] += dfs(v, u);
    return sub[u]:
 int dfs(int u, int p, int n) {
   for (auto v : tree[u])
     if (v != p \text{ and } sub[v] > n/2) \text{ return } dfs(v, u, n);
   return u;
 int operator[](int i) {
   return dad[i];
};
```

ChordalGraph.cpp

78d46d, 44 lines

```
//Maximum Cardinality Search
while (cur) {
 p[cur] = h[nww];
 rnk[p[cur]] = cur;
 h[nww] = nxt[h[nww]];
 lst[h[nww]] = 0;
 lst[p[cur]] = nxt[p[cur]] = 0;
 tf[p[cur]] = true;
 for (vector<int>::iterator it = G[p[cur]].begin(); it != G
      [p[curll.end();
      it++)
   if (!tf[*it]) {
     if (h[deg[*it]] == *it) h[deg[*it]] = nxt[*it];
     nxt[lst[*it]] = nxt[*it];
     lst[nxt[*it]] = lst[*it];
     lst[*it] = nxt[*it] = 0;
     dea[*it]++;
     nxt[*it] = h[deg[*it]];
     lst[h[deg[*it]]] = *it;
     h[deq[*it]] = *it;
```

```
cur--;
 if (h[nww + 1]) nww++;
 while (nww && !h[nww]) nww--;
//Checking if a sequence is a perfect elimination ordering
iud = true;
for (int i = 1; i <= n; i++) {
 cur = 0:
  for (vector<int>::iterator it = G[p[i]].begin(); it != G[p
       [i]].end(); it++)
    if (rnk[p[i]] < rnk[*it]) {</pre>
      s[++cur] = *it;
      if (rnk[s[cur]] < rnk[s[1]]) swap(s[1], s[cur]);</pre>
  for (int j = 2; j <= cur; j++)
    if (!st[s[1]].count(s[j])) {
      jud = false;
      break;
if (!jud)
 printf("Imperfect\n");
 printf("Perfect\n");
```

CircleIntersection.cpp

CircleLine.cpp

CirclePolygonIntersection.cpp

```
"../../content/geometry/Point.h" 3e5102, 19 lines
typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
   auto tri = [&] (P p, P q) {
     auto r2 = r * r / 2;
     P d = q - p;
```

CircleTangents.cpp

Circumcircle.cpp

ClosestPair.cpp

```
ConvexHull.cpp
```

da5fe0, 44 lines

```
struct pt {
   ld x, y;
bool cmp(pt a, pt b) {
   return a.x < b.x || (a.x == b.x && a.y < b.y);</pre>
bool cw(pt a, pt b, pt c) {
   return a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y) < 0;
bool ccw(pt a, pt b, pt c) {
   return a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y) > 0;
void convex hull(vector<pt>& a) {
   if (a.size() == 1)
       return:
    sort(a.begin(), a.end(), &cmp);
   pt p1 = a[0], p2 = a.back();
   vector<pt> up, down;
   up.push back(p1);
   down.push back(p1);
    for (int i = 1; i < (int)a.size(); i++) {</pre>
        if (i == a.size() - 1 || cw(p1, a[i], p2)) {
            while (up.size() \ge 2 \&\& !cw(up[up.size()-2], up
                [up.size()-1], a[i]))
                up.pop_back();
            up.push_back(a[i]);
       if (i == a.size() - 1 || ccw(p1, a[i], p2)) {
            while(down.size() >= 2 && !ccw(down[down.size()
                -2], down[down.size()-1], a[i]))
                down.pop_back();
            down.push_back(a[i]);
   a.clear():
    for (int i = 0; i < (int)up.size(); i++)</pre>
        a.push_back(up[i]);
    for (int i = down.size() - 2; i > 0; i--)
       a.push_back(down[i]);
```

ConvexHullTrick.cpp

875d66, 26 line

```
//This represents those relying on j only;i.e intercept
inline ll y_axis(int j) {
    return dp[j] + a * s[j] * s[j] - b * s[j];
}
//This represents those relying on both i and j; i.e slope
inline ll x_axis(int j) {
    return s[j];
}
inline ld getSlope(int j,int k) {
    ld y = y_axis(k) - y_axis(j);
    ld x = x_axis(k) - x_axis(j);
    return y / x;
```

CRT.cpp

118857, 32 lines

```
//each is x \mod p_i = a_i
11 p[maxn], a[maxn];
//for quickmult see pollard rho
ll exgcd(ll x, ll v, ll & a, ll & b) {
    if (y == 0) {
        a = 1; b = 0; return x;
    11 d = exgcd(y, x%y, a, b);
    11 temp = a; a = b; b = temp - (x / y) * b;
    return d:
int first nontrivial = 0;
ll current p :
11 sol = 0; //this is the solution
for (int i = 1; i \le n; i++) {
 if (p[i] != 1) {
    first nontrivial = i;
    current_p = p[i]; sol = a[i];
    break:
for (int i = first_nontrivial+1;i <= n;i++) {</pre>
 11 x, y;
 if (p[i] == 1) continue;
 11 d = exgcd(current_p, p[i], x, y);
 ll r = ((a[i] - sol) % p[i] + p[i]) % p[i];
 ll temp = quickmult(x, r / d,p[i] / d);
  sol = sol + current_p * temp;
  current_p = current_p / d * p[i];
  sol = (sol % current_p + current_p) % current_p;
```

Cutpoints.cpp

23f413, 37 lines

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph

vector<bool> visited;
vector<int> tin, low;
int timer;

void dfs(int v, int p = -1) {
    visited[v] = true;
    tin[v] = low[v] = timer++;
```

```
int children=0;
    for (int to : adj[v]) {
       if (to == p) continue;
        if (visited[to]) {
            low[v] = min(low[v], tin[to]);
           dfs(to, v);
            low[v] = min(low[v], low[to]);
            if (low[to] >= tin[v] && p!=-1)
                IS_CUTPOINT(v);
            ++children;
   if(p == -1 \&\& children > 1)
        IS_CUTPOINT(v);
void find_cutpoints() {
   timer = 0;
   visited.assign(n, false);
   tin.assign(n, -1);
   low.assign(n, -1);
   for (int i = 0; i < n; ++i) {
       if (!visited[i])
            dfs (i);
```

Delaunay Triangulation.cpp

DFSMatching.cpp

0aa9b2, 37 lines

```
* Author: Lukas Polacek
 * Date: 2009-10-28
 * License: CC0
 * Source:
 * Description: Simple bipartite matching algorithm. Graph
      $q$ should be a list
 * of neighbors of the left partition, and $btoa$ should be
     a vector full of
 * -1's of the same size as the right partition. Returns the
 * the matching. $btoa[i]$ will be the match for vertex $i$
     on the right side.
 * or $-1$ if it's not matched.
 * Time: O(VE)
 * Usage: vi btoa(m, -1); dfsMatching(q, btoa);
 * Status: works
bool find(int j, vector<vi>& g, vi& btoa, vi& vis) {
 if (btoa[j] == -1) return 1;
```

```
vis[j] = 1; int di = btoa[j];
for (int e : g[di])
   if (!vis[e] && find(e, g, btoa, vis)) {
      btoa[e] = di;
      return 1;
   }
  return 0;
}
int dfsMatching(vector<vi>& g, vi& btoa) {
   vi vis;
  FOR(i,0,sz(g)) {
      vis.assign(sz(btoa), 0);
      for (int j : g[i])
      if (find(j, g, btoa, vis)) {
       btoa[j] = i;
        break;
      }
}
return sz(btoa) - (int)count(all(btoa), -1);
}
```

Dinic.cpp

b8d428, 83 lines

```
//from https://cp-algorithms.com/graph/dinic.html
//Complexity: O(E*V^2)
struct FlowEdge {
  int v, u;
  long long cap, flow = 0;
  FlowEdge(int v, int u, long long cap) : v(v), u(u), cap(
       cap) {}
};
struct Dinic {
  const long long flow inf = 1e18;
  vector<FlowEdge> edges;
  vector<vector<int>> adj;
  int n, m = 0;
  int s, t;
  vector<int> level, ptr;
  queue<int> q;
  Dinic(int n, int s, int t): n(n), s(s), t(t) {
      adj.resize(n);
      level.resize(n);
      ptr.resize(n);
  void add_edge(int v, int u, long long cap) {
      edges.emplace_back(v, u, cap);
      edges.emplace back(u, v, 0);
      adj[v].push_back(m);
      adj[u].push_back(m + 1);
      m += 2:
  bool bfs() {
      while (!a.emptv()) {
          int v = q.front();
          q.pop();
          for (int id : adj[v]) {
               if (edges[id].cap - edges[id].flow < 1)</pre>
                   continue:
               if (level[edges[id].u] != -1)
                  continue;
               level[edges[id].u] = level[v] + 1;
```

```
q.push(edges[id].u);
       return level[t] != -1;
  long long dfs(int v, long long pushed) {
      if (pushed == 0)
           return 0;
      if (v == t)
           return pushed;
       for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid
           int id = adj[v][cid];
           int u = edges[id].u;
           if (level[v] + 1 != level[u] || edges[id].cap -
               edges[id].flow < 1)
               continue:
           long long tr = dfs(u, min(pushed, edges[id].cap -
                 edges[id].flow));
           if (tr == 0)
              continue;
           edges[id].flow += tr;
           edges[id ^ 1].flow -= tr;
           return tr;
       return 0;
  long long flow() {
      long long f = 0;
      while (true) {
           fill(level.begin(), level.end(), -1);
           level[s] = 0;
           q.push(s);
           if (!bfs())
               break:
           fill(ptr.begin(), ptr.end(), 0);
           while (long long pushed = dfs(s, flow_inf)) {
               f += pushed;
       return f:
};
```

DSU.cpp

89f1c6, 19 lines

```
int parent[MX], si[MX];
void init(int N) {
   FOR(i, N) parent[i] = i, si[i] = 0;
int get(int x) {
    if (parent[x] != x) parent[x] = get(parent[x]);
    return parent[x]:
void unify(int x, int y) {
    x = get(x); v = get(v);
    if (x == y) return;
   if (si[x] < si[y]) swap(x, y);
   if (si[x] == si[y]) si[x]++;
   parent[y] = x;
```

```
EulerPath.cpp
                                                  274951, 25 lines
int N. M:
vector<vpi> graph(MX); //{ed, edNum}
vector<vpi::iterator> its(MX);
vector<bool> used(MX);
vpi eulerPath(int r) {
    FOR(i, N) its[i] = begin(graph[i]);
    FOR(i, M) used[i] = false;
    vpi ans, s\{\{r, -1\}\};
    int lst = -1;
    while (sz(s)) {
        int x = s.back().f; auto &it = its[x], en = end(
             graph[x]);
        while (it != en && used[it->s]) it++;
        if (it == en) {
            if (lst != -1 && lst != x) return {};
            ans.pb(s.back()); s.pop_back(); if (sz(s)) lst =
                 s.back().f;
        } else {
            s.pb(*it);
            used[it->s] = 1;
    if (sz(ans) != M+1) return {};
    return ans;
FastDelaunay.cpp
"Point.h"
                                                   1c46ca, 88 lines
typedef Point<11> P;
typedef struct Quad* Q;
typedef __int128_t lll; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
struct Ouad {
 bool mark; Q o, rot; P p;
 P F() { return r()->p; }
 Q r() { return rot->rot; }
 O prev() { return rot->o->rot; }
 Q next() { return r()->prev(); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
 111 p2 = p.dist2(), A = a.dist2()-p2,
      B = b.dist2()-p2, C = c.dist2()-p2;
  return p.cross(a,b) *C + p.cross(b,c) *A + p.cross(c,a) *B >
O makeEdge(P orig, P dest) {
  Q q[] = \{new Quad\{0,0,0,oriq\}, new Quad\{0,0,0,arb\},
           new Ouad{0,0,0,dest}, new Ouad{0,0,0,arb}};
```

q[i] -> o = q[-i & 3], q[i] -> rot = q[(i+1) & 3];

swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);

return *q;

void splice(0 a, 0 b) {

Q connect(Q a, Q b) {

```
splice(q, a->next());
  splice(q->r(), b);
 return q;
pair<Q,Q> rec(const vector<P>& s) {
 if (sz(s) \le 3) {
   Q = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back())
   if (sz(s) == 2) return { a, a->r() };
   splice(a->r(), b);
   auto side = s[0].cross(s[1], s[2]);
   Q c = side ? connect(b, a) : 0;
   return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
\#define H(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (e->F().cross(H(base)) > 0)
 O A, B, ra, rb;
 int half = sz(s) / 2;
 tie(ra, A) = rec({all(s) - half});
 tie(B, rb) = rec({sz(s) - half + all(s)});
 while ((B->p.cross(H(A)) < 0 \&& (A = A->next()))
         (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
 Q base = connect(B->r(), A);
  if (A->p == ra->p) ra = base->r();
 if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
   while (circ(e->dir->F(), H(base), e->F())) {
     Q t = e->dir; \
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
     e = t; \
  for (;;) {
   DEL(LC, base->r(), o); DEL(RC, base, prev());
   if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
     base = connect(RC, base->r());
   else
     base = connect(base->r(), LC->r());
 return { ra, rb };
vector<P> triangulate(vector<P> pts) {
 sort(all(pts)); assert(unique(all(pts)) == pts.end());
 if (sz(pts) < 2) return {};
 Q e = rec(pts).first;
 vector < Q > q = \{e\};
  int qi = 0;
  while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p)
    ; \
 g.push back(c\rightarrow r()); c = c\rightarrow next(); } while (c != e); }
 ADD; pts.clear();
 while (qi < sz(q)) if (!(e = q[qi++])->mark) ADD;
 return pts;
FastHashTable.cpp
```

Q = makeEdge(a->F(), b->p);

<ext/pb_ds/assoc_container.hpp> 8f2fd7, 2 lines using namespace __gnu_pbds;

FFT GeneralMatching GeometrySnippets

```
gp hash table<int, int> table;
```

FFT.cpp

89f37e, 61 lines

```
using cd = complex<double>;
const double PI = acos(-1);
int reverse(int num, int lq n) {
    int res = 0;
    for (int i = 0; i < lq n; i++) {
       if (num & (1 << i))
           res |= 1 << (lg_n - 1 - i);
    return res;
void fft(vector<cd> & a, bool invert) {
    int n = a.size();
    int lq_n = 0;
    while ((1 \ll lq n) \ll n)
       lg_n++;
    for (int i = 0; i < n; i++) {
       if (i < reverse(i, lq_n))</pre>
            swap(a[i], a[reverse(i, lg_n)]);
    for (int len = 2; len <= n; len <<= 1) {
       double ang = 2 * PI / len * (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len) {
            cd w(1);
            for (int j = 0; j < len / 2; j++) {
               cd u = a[i+j], v = a[i+j+len/2] * w;
               a[i+j] = u + v;
               a[i+j+len/2] = u - v;
                w \neq wlen;
        }
    if (invert) {
        for (cd & x : a)
           x /= n;
vector<11> multiply(vector<11> const& a, vector<11> const& b
    vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end()
       );
    int n = 1:
    while (n < a.size() + b.size())</pre>
       n <<= 1;
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for (int i = 0; i < n; i++)
       fa[i] *= fb[i];
    fft(fa, true);
    vector<ll> result(n);
    for (int i = 0; i < n; i++)
        result[i] = round(fa[i].real());
```

```
return result;
```

```
GeneralMatching.cpp
                                                 04bdb9, 78 lines
//belong is a DSU; unit = union
int n, match[N], next[N], mark[N], vis[N], Q[N];
std::vector<int> e[N];
int rear;
int LCA(int x, int y) {
 static int t = 0; t++;
 while (true) {
   if (x != -1) {
     x = findb(x);
     if (vis[x] == t) return x;
     vis[x] = t;
     if (match[x] != -1) x = next[match[x]];
     else x = -1:
   std::swap(x, y);
void group(int a, int p){
   while (a != p) {
     int b = match[a], c = next[b];
     if (findb(c) != p) next[c] = b;
     if (mark[b] == 2) mark[Q[rear++] = b] = 1;
     if (mark[c] == 2) mark[Q[rear++] = c] = 1;
     unit(a, b); unit(b, c);
void aug(int s){
 for (int i = 0; i < n; i++)
   next[i] = -1, belong[i] = i, mark[i] = 0, vis[i] = -1;
 mark[s] = 1;
 Q[0] = s; rear = 1;
 for (int front = 0; match[s] == -1 && front < rear; front
   int x = Q[front];
   for (int i = 0; i < (int)e[x].size(); i++){
     int v = e[x][i];
      if (match[x] == y) continue;
      if (findb(x) == findb(y)) continue;
      if (mark[y] == 2) continue;
      if (mark[y] == 1) {
       int r = LCA(x, y);
        if (findb(x) != r) next[x] = y;
       if (findb(y) != r) next[y] = x;
       group(x, r);
       group(y, r);
      else if (match[y] == -1) {
       next[v] = x;
       for (int u = y; u != -1; ) {
         int v = next[u];
         int mv = match[v];
         match[v] = u, match[u] = v; u = mv;
       break;
      else{
```

```
next[y] = x;
       mark[Q[rear++] = match[y]] = 1;
       mark[y] = 2;
//the\ graph\ is\ stored\ as\ e[N]\ and\ g[N]
for (int i = 0; i < n; i++) match[i] = -1;
for (int i = 0; i < n; i++) if (match[i] == -1) aug(i);
int tot = 0;
for (int i = 0; i < n; i++) {
if (match[i] != -1) tot++;
//matched\ pairs = tot/2
printf("%d\n", tot/2);
for (int i = 0; i < n; i++) {
   printf("%d ", match[i] + 1);
```

f647cc, 86 lines

```
GeometrySnippets.cpp
They are from KACTL, KTH's Team Reference Document
//check point in convex hull
typedef Point<11> P;
bool inHull(const vector<P>& 1, P p, bool strict = true) {
 int a = 1, b = sz(1) - 1, r = !strict;
 if (sz(1) < 3) return r && onSegment(1[0], 1.back(), p);
 if (sideOf(1[0], 1[a], 1[b]) > 0) swap(a, b);
 if (sideOf(1[0], 1[a], p) >= r || sideOf(1[0], 1[b], p) <=
      -r)
    return false;
  while (abs(a - b) > 1) {
   int c = (a + b) / 2;
    (sideOf(1[0], 1[c], p) > 0 ? b : a) = c;
  return sqn(l[a].cross(l[b], p)) < r;</pre>
//center of mass of polygon
typedef Point < double > P;
P polygonCenter(const vector<P>& v) {
 P res(0, 0); double A = 0;
 for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
   res = res + (v[i] + v[j]) * v[j].cross(v[i]);
   A += v[i].cross(v[i]);
 return res / A / 3;
// Returns a vector with the vertices of a polygon with
    everything to the left
// of the line going from s to e cut away.
vpedef Point<double> P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
 vector<P> res;
 FOR(i, 0, sz(polv)) {
   P cur = poly[i], prev = i ? poly[i-1] : poly.back();
   bool side = s.cross(e, cur) < 0;</pre>
   if (side != (s.cross(e, prev) < 0))</pre>
      res.push_back(lineInter(s, e, cur, prev).second);
```

HalfPlaneIntersection HopcroftKarp

```
if (side)
      res.push_back(cur);
  return res;
//volumn of polyhedron, with face outwards
template<class V, class L>
double signed_poly_volume(const V& p, const L& trilist) {
 double v = 0;
 trav(i, trilist) v += p[i.a].cross(p[i.b]).dot(p[i.c]);
 return v / 6;
//intersection of two lines
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
 auto d = (e1 - s1).cross(e2 - s2);
 if (d == 0) // if parallel
   return \{-(s1.cross(e1, s2) == 0), P(0, 0)\};
  auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
 return \{1, (s1 * p + e1 * q) / d\};
// Returns where p is as seen from s towards e. 1/0/-1
      $\Leftrightarrow$ left/on line/right.
template<class P>
int sideOf(P s, P e, P p) { return sqn(s.cross(e, p)); }
template<class P>
int sideOf(const P& s, const P& e, const P& p, double eps) {
 auto a = (e-s).cross(p-s);
 double l = (e-s).dist()*eps;
 return (a > 1) - (a < -1);
//f is longitude, t is latitude
double sphericalDistance(double f1, double t1,
 double f2, double t2, double radius) {
 double dx = \sin(t2) \cdot \cos(f2) - \sin(t1) \cdot \cos(f1);
 double dy = sin(t2) * sin(f2) - sin(t1) * sin(f1);
 double dz = cos(t2) - cos(t1);
 double d = sgrt(dx*dx + dy*dy + dz*dz);
  return radius *2 *asin(d/2);
```

HalfPlaneIntersection.cpp

80d545, 86 lines

```
// Comparator for sorting.
    bool operator < (const Halfplane& e) const {
        return angle < e.angle;
    // Intersection point of the lines of two half-planes.
         It is assumed they're never parallel.
    friend Point inter(const Halfplane& s, const Halfplane&
        long double alpha = cross((t.p - s.p), t.pq) / cross
             (s.pq, t.pq);
        return s.p + (s.pq * alpha);
};
// Actual algorithm
vector<Point> hp_intersect(vector<Halfplane>& H) {
    Point box[4] = { // Bounding box in CCW order
        Point(inf, inf),
        Point(-inf, inf),
        Point (-inf, -inf),
        Point(inf, -inf)
    for (int i = 0; i<4; i++) { // Add bounding box half-
        Halfplane aux(box[i], box[(i+1) % 4]);
        H.push_back(aux);
    // Sort by angle and start algorithm
    sort(H.begin(), H.end());
    deque<Halfplane> dq;
    int len = 0;
    for(int i = 0; i < int(H.size()); i++) {</pre>
        // Remove from the back of the deque while last half
             -plane is redundant
        while (len > 1 && H[i].out(inter(dg[len-1], dg[len
            -21))) {
            dq.pop_back();
            --len;
        // Remove from the front of the deque while first
             half-plane is redundant
        while (len > 1 && H[i].out(inter(dq[0], dq[1]))) {
            dq.pop_front();
            --len;
        // Special case check: Parallel half-planes
        if (len > 0 && fabsl(cross(H[i].pg, dg[len-1].pg)) <
              eps) {
            // Opposite parallel half-planes that ended up
                 checked against each other.
            if (dot(H[i].pq, dq[len-1].pq) < 0.0)</pre>
                return vector<Point>();
            // Same direction half-plane: keep only the
                 leftmost half-plane.
            if (H[i].out(dq[len-1].p)) {
                dq.pop_back();
                --len;
            else continue:
        // Add new half-plane
        dq.push_back(H[i]);
        ++len;
```

```
// Final cleanup: Check half-planes at the front against
      the back and vice-versa
while (len > 2 && dq[0].out(inter(dq[len-1], dq[len-2]))
    dq.pop_back();
    --len:
while (len > 2 && dq[len-1].out(inter(dq[0], dq[1]))) {
    dq.pop_front();
    --len:
// Report empty intersection if necessary
if (len < 3) return vector<Point>();
// Reconstruct the convex polygon from the remaining
     half-planes.
vector<Point> ret(len);
for(int i = 0; i+1 < len; i++) {
    ret[i] = inter(dq[i], dq[i+1]);
ret.back() = inter(dq[len-1], dq[0]);
return ret;
```

HopcroftKarp.cpp

Time: $\mathcal{O}\left(\sqrt{V}E\right)$

// OK

Description: Fast bipartite matching algorithm. Graph g should be a list of neighbors of the left partition, and btoa should be a vector full of -1's of the same size as the right partition. Returns the size of the matching. btoa[i] will be the match for vertex i on the right side, or -1 if it's not matched.

```
Usage: vi btoa(m, -1); hopcroftKarp(g, btoa);
```

```
2ff797, 42 lines
bool dfs(int a, int L, vector<vi>& g, vi& btoa, vi& A, vi& B
 if (A[a] != L) return 0;
 A[a] = -1;
  for (int b : g[a]) if (B[b] == L + 1) {
   B[b] = 0;
   if (btoa[b] == -1 \mid | dfs(btoa[b], L + 1, g, btoa, A, B))
     return btoa[b] = a, 1;
 return 0;
int hopcroftKarp(vector<vi>& g, vi& btoa) {
 int res = 0;
 vi A(g.size()), B(btoa.size()), cur, next;
  for (;;) {
    fill(all(A), 0);
   fill(all(B), 0);
    cur.clear();
    for (int a : btoa) if (a != -1) A[a] = -1;
   FOR(a, 0, sz(q)) if (A[a] == 0) cur.push back(a);
    for (int lay = 1;; lay++) {
     bool islast = 0;
     next.clear();
      for (int a : cur) for (int b : g[a]) {
       if (btoa[b] == -1) {
         B[b] = lav;
         islast = 1;
        else if (btoa[b] != a && !B[b]) {
         B[b] = lay;
```

```
next.push_back(btoa[b]);
  if (islast) break;
 if (next.empty()) return res;
 for (int a : next) A[a] = lay;
 cur.swap(next);
FOR(a, 0, sz(g))
  res += dfs(a, 0, q, btoa, A, B);
```

HullDiameter.cpp

0c6e60, 12 lines

```
typedef Point<11> P;
array<P, 2> hullDiameter(vector<P> S) {
 int n = sz(S), j = n < 2 ? 0 : 1;
 pair<11, array<P, 2>> res({0, {S[0], S[0]}});
 FOR(i,0,j)
    for (;; j = (j + 1) % n) {
      res = \max(\text{res}, \{(S[i] - S[j]).dist2(), \{S[i], S[j]\}\});
      if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >=
          0)
        break;
 return res.second;
```

InsidePolygon.cpp

```
"Point.h", "OnSegment.h", "SegmentDistance.h"
                                                    fd40d6, 11 lines
template<class P>
bool inPolygon(vector<P> &p, P a, bool strict = true) {
 int cnt = 0, n = sz(p);
 FOR(i,0,n) {
   P q = p[(i + 1) % n];
   if (onSegment(p[i], q, a)) return !strict;
    //or: if (segDist(p[i], q, a) \le eps) return ! strict;
    cnt ^= ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) >
         0;
  return cnt;
```

InsidePolygonFast.cpp

96e771, 52 lines

```
bool lexComp(const pt &1, const pt &r) {
    return 1.x < r.x || (1.x == r.x && 1.y < r.y);
int sgn(long long val) { return val > 0 ? 1 : (val == 0 ? 0
    : -1); }
vector<pt> seq;
pt translation;
bool pointInTriangle(pt a, pt b, pt c, pt point) {
    long long s1 = abs(a.cross(b, c));
    long long s2 = abs(point.cross(a, b)) + abs(point.cross(
         b, c)) + abs(point.cross(c, a));
    return s1 == s2;
void prepare(vector<pt> &points) {
    n = points.size();
    int pos = 0;
    for (int i = 1; i < n; i++) {
        if (lexComp(points[i], points[pos]))
```

```
rotate(points.begin(), points.begin() + pos, points.end
   n--;
   seq.resize(n);
   for (int i = 0; i < n; i++)
        seq[i] = points[i + 1] - points[0];
   translation = points[0];
bool pointInConvexPolygon(pt point) {
   point = point - translation;
   if (seq[0].cross(point) != 1 &&
            sqn(seq[0].cross(point)) != sqn(seq[0].cross(seq
                [n - 1])))
        return false:
   if (seq[n - 1].cross(point) != 0 &&
            sqn(seq[n-1].cross(point)) != sqn(seq[n-1].
                cross(seq[0])))
        return false;
    if (seq[0].cross(point) == 0)
        return seq[0].sqrLen() >= point.sqrLen();
   int 1 = 0, r = n - 1;
   while (r - 1 > 1) {
        int mid = (1 + r) / 2;
        int pos = mid;
        if (seq[pos].cross(point) >= 0)
           1 = mid;
            r = mid;
   int pos = 1;
    return pointInTriangle(seq[pos], seq[pos + 1], pt(0, 0),
```

IntervalContainer.cpp

Description: Add and remove intervals from a set of disjoint intervals. Will merge the added interval with any overlapping intervals in the set when adding. Intervals are [inclusive, exclusive). Time: $\mathcal{O}(\log N)$

```
set<pii>::iterator addInterval(set<pii>& is, int L, int R) {
 if (L == R) return is.end();
 auto it = is.lower_bound({L, R}), before = it;
 while (it != is.end() && it->first <= R) {</pre>
   R = max(R, it->second);
   before = it = is.erase(it);
 if (it != is.begin() && (--it)->second >= L) {
   L = min(L, it->first);
   R = max(R, it->second);
   is.erase(it);
 return is.insert(before, {L,R});
void removeInterval(set<pii>& is, int L, int R) {
 if (L == R) return;
 auto it = addInterval(is, L, R);
 auto r2 = it->second;
```

```
if (it->first == L) is.erase(it);
 else (int&)it->second = L;
 if (R != r2) is.emplace(R, r2);
KdTree.cop
Description: KD-tree (2d, can be extended to 3d)
                                                  bac5b0, 63 lines
typedef long long T;
typedef Point<T> P;
const T INF = numeric_limits<T>::max();
bool on_x(const P& a, const P& b) { return a.x < b.x; }
bool on_y(const P& a, const P& b) { return a.y < b.y; }
struct Node {
 P pt; // if this is a leaf, the single point in it
 T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; // bounds
 Node *first = 0, *second = 0;
 T distance (const P& p) { // min squared distance to a
   T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
   T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
   return (P(x,y) - p).dist2();
 Node (vector<P>&& vp) : pt(vp[0]) {
   for (P p : vp) {
     x0 = min(x0, p.x); x1 = max(x1, p.x);
     y0 = min(y0, p.y); y1 = max(y1, p.y);
   if (vp.size() > 1) {
     // split on x if width >= height (not ideal...)
     sort(all(vp), x1 - x0 >= y1 - y0 ? on_x : on_y);
      // divide by taking half the array for each child (not
      // best performance with many duplicates in the middle
      int half = sz(vp)/2;
     first = new Node({vp.begin(), vp.begin() + half});
     second = new Node({vp.begin() + half, vp.end()});
};
struct KDTree {
 Node* root;
 KDTree(const vector<P>& vp) : root(new Node({all(vp)})) {}
 pair<T, P> search(Node *node, const P& p) {
   if (!node->first) {
      // uncomment if we should not find the point itself:
      // if (p = node \rightarrow pt) return \{INF, P()\};
     return make_pair((p - node->pt).dist2(), node->pt);
   Node *f = node->first, *s = node->second;
   T bfirst = f->distance(p), bsec = s->distance(p);
   if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);
   // search closest side first, other side if needed
   auto best = search(f, p);
   if (bsec < best.first)</pre>
```

best = min(best, search(s, p));

return best;

KMP KnuthOptimization LazySegtree LCA

for (int $k = m[1][r - 1]; k \le m[1 + 1][r]; ++k$)

m[1][r] = k; // Update Decision Point

if (f[1][r] > f[1][k] + f[k + 1][r] + w(1, r)) {

 $/*Class2: 2D DP, f_{\{i,j\}} = min_{\{k \le j\}} \{f_{\{i-1,k\}}\} + w(k,j)$

f[1][r] = f[1][k] + f[k + 1][r] + w(1, r); //Update

// Enumerate Intervals of Length Len

vector<long long> dp_before(n), dp_cur(n);

pair<long long, int> best = {INF, -1};

//Call compute for each possible i.

// compute $dp_cur[l]$, ... $dp_cur[r]$ (inclusive)

void compute(int 1, int r, int opt1, int optr) {

for (int k = optl; k <= min(mid, optr); k++) {</pre>

best = min(best, {dp_before[k] + C(k, mid), k});

f[1][r] = INF;

Where $1 \le i \le n$, $1 \le i \le m$

long long C(int i, int j);

dp_cur[mid] = best.first;

compute(1, mid - 1, opt1, opt);

compute(mid + 1, r, opt, optr);

int opt = best.second;

if (1 > r) return; int mid = (1 + r) >> 1;

```
// find nearest point to a point, and its squared distance
  // (requires an arbitrary operator< for Point)
 pair<T, P> nearest(const P& p) {
    return search(root, p);
};
KMP.cpp
                                                    e6241c, 37 lines
 nxt[0] = 0;
  for (int i = 1;i < length;i++) {</pre>
    while (k > 0 \&\& str[k] != str[i]) {
     k = nxt[k - 1];
    if (str[k] == str[i]) {
     k++;
```

LazySegtree.cpp

a5b59d, 38 lines

int n;

8ed0ff, 45 lines

```
const 11 identity = 0;
const 11 SZ = 131072;
11 sum[2*SZ], lazy[2*SZ];
ll combine(ll A, ll B) {
   return A+B;
ll combineUpd(ll A, ll B) {
    return A+B;
void push(int index, ll L, ll R) {
    sum[index] = combineUpd(sum[index], lazy[index]);
    if (L != R) lazy[2*index] = combineUpd(lazy[2*index],
        lazy[index]), lazy[2*index+1] = combineUpd(lazy[2*
        index+1], lazy[index]);
    lazy[index] = identity;
void pull(int index) {
    sum[index] = combine(sum[2*index], sum[2*index+1]);
11 query (int lo, int hi, int index = 1, 11 L = 0, 11 R = SZ
    push (index, L, R);
    if (lo > R | | L > hi) return identity;
    if (lo <= L && R <= hi) return sum[index];</pre>
    int M = (L+R) / 2;
```

```
return combine (query (lo, hi, 2*index, L, M), query (lo,
        hi, 2*index+1, M+1, R));
void update(int lo, int hi, ll increase, int index = 1, ll L
     = 0, 11 R = SZ-1)  {
   push(index, L, R);
   if (hi < L || R < lo) return;
   if (lo <= L && R <= hi) {
       lazy[index] = increase;
       push(index, L, R);
        return;
   int M = (L+R) / 2;
   update(lo, hi, increase, 2*index, L, M); update(lo, hi,
        increase, 2*index+1, M+1, R);
   pull(index);
LCA.cpp
                                                   041ec2, 55 lines
const int L; /\!/\!SET THIS TO CEIL(LOG(MX\_N))
int N;
int anc[MX][L];
int depth[MX];
int parent[MX];
vector<vi> graph (MX);
int LCA(int a, int b) {
   if (depth[a] < depth[b]) {</pre>
       int c = b;
       b = a;
        a = c;
   int dist = depth[a] - depth[b];
   while (dist > 0) {
       FOR(i, L) {
            if (dist & 1 << i) {
                a = anc[a][i];
                dist -= 1 << i;
   }
   if (a == b) return a;
   FORd(j, L) {
        if (anc[a][i] != -1 \&\& anc[a][i] != anc[b][i]) {
            a = anc[a][j]; b = anc[b][j];
    return parent[a];
void parDFS(int v, int p, int d) {
   parent[v] = p; depth[v] = d;
    FOR(i, sz(graph[v])) {
        int nxt = graph[v][i];
       if (nxt == p) continue;
       parDFS(nxt, v, d+1);
```

```
int nxt[maxn]; //next array
/*Core KMP @param nxt: the resulting "next" array*/
void build_nxt(char * str,int * nxt, int length) {
    nxt[i] = k;
/*Matching the string with the pattern
@return number of occurences of pattern string in the
    original string
int match (char * str, char * pattern, int length_str, int
    length pattern) {
 int total = 0;
 int p = 0;
  for (int i = 0;i < length_str;i++) {</pre>
    while (p > 0 && pattern[p] != str[i]) {
     p = nxt[p - 1];
    if (pattern[p] == str[i]) {
      p++;
    if (p == length_pattern) {
     total++;
     p = nxt[p - 1];
  return total;
//build nxt for pattern
```

 $/*Class1 : Interval DP: f_{\{l,r\}} = min_{\{k=l\}}^{r-1} f_{\{l,k\}+f_{-1}\}}$

(1) For any $l \le l' \le r' \le r$, we have $w(l', r') \le w(l, r)$.

(2) (The important one): For any $l1 \le l2 \le r1 \le r2$, we

for (int len = 2; len <= n; ++len) // Enumerate Interval</pre>

weights w(l,r) satisfying the following inequality:

 $w(l1,r1) + w(l2,r2) \le w(l1,r2) + w(l2,r1).$

for (int l = 1, r = len; r <= n; ++1, ++r) {

build_nxt(pattern, nxt, len_pattern);

KnuthOptimization.cpp

 $\{k+1,r\} + w(l,r)$

void preprocess() {

parDFS(0, -1, 0);

FOR(i, N) FOR(j, L) anc[i][j] = -1;

FOR(i, N) anc[i][0] = parent[i];

LCT LeftistTree LinearProgramming

```
FOR(j, 1, L) {
       FOR(i, N) {
            if (anc[i][j-1] != -1) {
                anc[i][j] = anc[anc[i][j-1]][j-1];
LCT.cpp
                                                  1809e5, 116 lines
struct rec
    int ls, rs, p; //ls = left \ son; \ rs = right \ son; \ p =
        parent
    uint siz; //siz = size of the subtree
    uint key, sum; //sum: sum of weights in the subtree
    uint mult, add; //two lazy tags
    bool rev; //denote whether this segment has been
         reverted
};
rec splay[maxn];
void clear(){
    splay[0].p = splay[0].ls = splay[0].rs = splay[0].rev =
         splay[0].key = splay[0].sum = 0;
    splay[0].siz = 0;
 void update(int x){
    clear():
    splay[x].sum = splay[splay[x].ls].sum + splay[splay[x].
         rs].sum + splay[x].key;
    splay[x].sum %= modi;
    splay[x].siz = splay[splay[x].ls].siz + splay[splay[x].
         rs].siz + 1;
    splay[x].siz %= modi;
void zig(int x){
    int y = splay[x].p, z = splay[y].p;
    if (y == splay[z].ls) splay[z].ls = x;
    else if (y == splay[z].rs) splay[z].rs = x;
    splav[x].p = z;
    // Switch is and rs for zag.
    if (splay[x].rs) splay[splay[x].rs].p = y;
    splay[y].ls = splay[x].rs;
    splay[x].rs = y;
    splav[v].p = x;
    update(y);
bool is root(int x){
    return x != splay[splay[x].p].ls && x != splay[splay[x].
void rev(int x){
    if (!x)return;
    swap(splay[x].ls, splay[x].rs);
    splav[x].rev ^= true;
void pushdown(int x){
    if (splay[x].rev) {
        rev(splay[x].ls);
        rev(splay[x].rs);
```

```
splav[x].rev = false;
    //Todo: Push lazy tags here.
void set_root(int x){
    static int q[maxn];
    static int top;
    int i:
    for (i = x; !is_root(i); i = splay[i].p) {
        q[++top] = i;
    q[++top] = i;
    while (top) {
        pushdown(q[top--]);
    while (!is_root(x)){
        int v = splav[x].p;
        if (is_root(y)){
            if (x == splay[y].ls) zig(x); else zag(x);
        else{
            int z = splay[y].p;
            if (y == splay[z].ls){
                if (x == splay[y].ls) zig(y), zig(x);
                else zag(x), zig(x);
            else{
                if (x == splay[y].rs) zag(y), zag(x);
                else ziq(x), zaq(x);
    update(x);
//this is a special operation on LCT
void access(int x)
    for (int t = 0; x; t = x, x = splay[x].p){
        set root(x);
        splay[x].rs = t;
        update(x);
//we will make x be the new root of the tree it belongs to
void makeroot(int x) {access(x); set_root(x); rev(x);}
void split(int x, int y) {makeroot(x);access(y);set_root(y);}
//link vertex x and vertex y
void link(int x, int y) {makeroot(x); makeroot(y); splay[x].p =
      y;}
//cut the edge between x and y
void cut(int x, int y){
    split(x, y);
    splay[y].ls = splay[x].p = 0;
    update(y);
//find the root; x connected with y IFF findroot(x) =
    findroot(y)
int findroot(int x){
    access(x);
    set_root(x);
    while (splay[x].ls) {
        pushdown(x);
        x = splay[x].ls;
```

```
set root(x);
    return x;
//Adding\ edge\ between\ u\ and\ v:\ link(u,\ v);
//Removing edge between u and v: cut(u1, v1);
//Adding vertices on route between u and v by c :
/* split(u, v);
   calc(v, 1, c);*/
//Query the sum on route from u to v: split(u1, v1) print(
     splay[v1].sum);
LeftistTree.cpp
                                                    7d92ca, 21 lines
struct node{
   node *1, *r;
    //key is the priority
    int key, id;
    //distanct to the leftist child - it is used to maintain
          the properties of the lefitst tree
    int dist:
    int rdist(){return (r==NULL)?0:r->dist;}
    int ldist(){return (l==NULL)?0:1->dist;}
node* merge(node*1,node*r)
    if (1 == NULL) return r;
   if (r == NULL) return 1;
    //we want to make sure the root has the smallest key
    if (1->key > r->key) swap(1,r);
   1->r = merge(1->r,r);
    //maintain the properties of the leftist tree
    if (1->ldist() < 1->rdist()) swap(1->1,1->r);
   1->dist = 1->rdist()+1;
    return 1:
LinearProgramming.cpp
                                                    220aa9, 96 lines
/*solving linear programming with simplex algorithm.
Problem: Given the i-th condition to be \sum_{j=1}^n a_{ij}
     x_i \setminus leq b_i
x_{-j} \setminus geq \ 0 \ (m \ conditions \ in \ total)
Find X-j that maximizes \sum_{j=1}^n c_{jx-j*}
const double eps = 1e-8;
const double inf = 1e100;
double a[maxn][maxn], b[maxn], c[maxn], ans[maxn];
int id[maxn << 1];</pre>
int n, m; //n = num \ of \ variables; m = constraints
void pivot(int base, int base2) {
 swap(id[base + n], id[base2]);
 for (int j = 0; j \le n; ++j) {
   if (j != base2) {
      a[base][i] /= a[base][base2];
 a[base][base2] = 1 / a[base][base2];
 for (int i = 0; i \le m; ++i) {
   if (i != base && fabs(a[i][base2]) > eps){
      for (int j = 0; j \le n; ++j) {
        if (i != base2)
          a[i][j] -= a[base][j] * a[i][base2];
      a[i][base2] *= -a[base][base2];
```

900055, 29 lines

```
bool init(){
  for (int i = 1; i \le n + m; ++i)
   id[i] = i;
  while (true) {
    int x = 0, y = 0;
    double mini = 0;
    for (int i = 1;i <= m;++i) {
      if (a[i][0] < mini - eps){</pre>
       x = i;
        mini = a[i][0];
    if (!x) break;
    for (int j = 1; j \le n; ++j) {
     if (a[x][j] < -eps) y = j;
    if (!y) return false;
    pivot(x, y);
  return true;
bool simplex() {
  while (true) {
    int x = 0, y = 0;
    double mini = inf;
    for (int j = 1; j \le n; ++j) {
      if (a[0][j] > eps){
        y = j;
        break:
    if (!y) break;
    for (int i = 1; i \le m; ++i) {
      if (a[i][y] > eps && a[i][0] / a[i][y] < mini - eps){</pre>
       x = i;
        mini = a[i][0] / a[i][y];
    if (!x) return false;
    pivot(x, y);
  return true;
//main program
for (int i = 1; i \le n; ++i) a[0][i] = c[i];
for (int i = 1; i \le m; ++i) a[i][0] = b[i];
if (!init()){
 printf("Infeasible\n"); return 0;
else if (!simplex()){
 printf("Unbounded\n"); return 0;
else{
 printf("%.101f\n", -a[0][0]);
  for (int i = 1; i \le m; ++i) {
   if (id[n + i] \le n) ans[id[n + i]] = a[i][0];
  if (t) {
      for (int i = 1; i \le n; ++i) printf("%.101f ", ans[i]);
```

```
printf("\n");
 return 0;
LinearTransformation.cpp
Description:
Apply the linear transformation (translation, rotation and
soalin which takes line p0-p1 to line q0-q1 to point r.
                                                    03a306, 6 lines
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1,
   const P& q0, const P& q1, const P& r) {
 P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
 return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2
LineDistance.cpp
"Point.h"
                                                    f6bf6b, 7 lines
Returns the signed distance between point p and the line
     containing points a and b. Positive value on left side
     and negative on right as seen from a towards b. a=b
     gives nan. P is supposed to be Point<T> or Point3D<T>
     where T is e.g. double or long long. It uses products in
     intermediate steps so watch out for overflow if using
     int or long long. Using Point3D will always give a non-
     negative distance. For Point3D, call .dist on the result
     of the cross product.
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
 return (double) (b-a).cross(p-a)/(b-a).dist();
LineHullIntersection.cpp
Time: \mathcal{O}(N + Q \log n)
"Point.h"
                                                    f78f76, 39 lines
typedef array<P, 2> Line;
#define cmp(i, j) sqn(dir.perp().cross(poly[(i)%n]-poly[(j)%n
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
int extrVertex(vector<P>& poly, P dir) {
 int n = sz(polv), lo = 0, hi = n;
 if (extr(0)) return 0;
 while (lo + 1 < hi) {
   int m = (lo + hi) / 2;
   if (extr(m)) return m;
    int 1s = cmp(1o + 1, 1o), ms = cmp(m + 1, m);
    (ls < ms \mid | (ls == ms \&\& ls == cmp(lo, m)) ? hi : lo) =
```

#define cmpL(i) sqn(line[0].cross(poly[i], line[1]))

int endA = extrVertex(poly, (line[0] - line[1]).perp());

int endB = extrVertex(poly, (line[1] - line[0]).perp());

array<int, 2> lineHull(Line line, vector<P> poly) {

return lo;

```
if (cmpL(endA) < 0 \mid | cmpL(endB) > 0)
   return {-1, -1};
 array<int, 2> res;
 FOR(i, 0, 2) {
   int lo = endB, hi = endA, n = sz(poly);
   while ((lo + 1) % n != hi) {
     int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(endB) ? lo : hi) = m;
   res[i] = (lo + !cmpL(hi)) % n;
   swap (endA, endB);
 if (res[0] == res[1]) return {res[0], -1};
 if (!cmpL(res[0]) && !cmpL(res[1]))
   switch ((res[0] - res[1] + sz(poly) + 1) % sz(poly)) {
     case 0: return {res[0], res[0]};
     case 2: return {res[1], res[1]};
 return res;
LineContainer.cpp
                                                   8ec1c7, 34 lines
/* Author: KACTL Line Container
* Description: Container where you can add lines of the
     form kx+m, and query maximum values at points x.
   Useful for dynamic programming (''convex hull trick'').
* Time: O(\log N)*/
struct Line {
 mutable 11 k, m, p;
 bool operator<(const Line& o) const { return k < o.k; }</pre>
 bool operator<(ll x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>>> {
 // (for doubles, use inf = 1/.0, div(a,b) = a/b)
  static const 11 inf = LLONG MAX;
 ll div(ll a, ll b) { // floored division
   return a / b - ((a ^ b) < 0 && a % b); }
 bool isect(iterator x, iterator v) {
   if (v == end()) return x \rightarrow p = inf, 0;
   if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
   else x->p = div(y->m - x->m, x->k - y->k);
   return x->p >= v->p;
 void add(ll k, ll m) {
   auto z = insert(\{k, m, 0\}), y = z++, x = y;
   while (isect(v, z)) z = erase(z);
   if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y))
   while ((y = x) != begin() \&\& (--x)->p >= y->p)
     isect(x, erase(y));
 ll query(ll x) {
   assert(!empty());
   auto 1 = *lower_bound(x);
   return 1.k * x + 1.m;
};
```

Manacher.cpp

using namespace std;

<set>, <cstdio>, <cstring>, <algorithm>

Matrix MinCostMaxFlow MinimumEnclosingCircle

```
//you need twice the length of original string here, since
    you need space for adding $
const int maxn = 2100000;
char str[maxn], str2[maxn];
//p[i] is just "how long you can extend from i in both ways"
int p[maxn];
void manacher(int n)
   int mx = 0;
    int id = 0;
    for (int i = 1;i <= n;++i) {
       if (mx >= i) {
            p[i] = min(mx - i, p[2 * id - i]);
        else{
            p[i] = 0;
        for (; str[i + p[i] + 1] == str[i - p[i] - 1];) {
           p[i]++;
        if (p[i] + i > mx) {
            id = i;
            mx = p[i] + i;
```

Matrix.cpp

67a0be, 21 lines

```
vector<vl> matMul(vector<vl>& A, vector<vl> &B) {
    int R = sz(A), C = sz(B[0]), I = sz(B);
    vector<vl> res(R, vector<vi>(C, 0));
    FOR(i, R) {
       FOR(j, I) {
           FOR(k, C) {
                res[i][k] += A[i][j] + B[j][k];
    return res;
vector<vi> exp(vector<vi> M, ll P) {
    vector < vi > res(sz(M), vi(sz(M[0]), 0));
    FOR(i, sz(M)) {
       res[i][i] = 1;
    for (; P; p \neq 2; M = matMul(M, M)) if (P & 1) R =
        matMul(R, M);
    return res;
```

MinCostMaxFlow.cpp

Description: Min-cost max-flow. cap[i][j] != cap[j][i] is allowed; double edges are not. If costs can be negative, call setpi before maxflow, but note that negative cost cycles are not supported. To obtain the actual flow, look at positive values only.

Time: Approximately $\mathcal{O}(E^2)$

c3f497, 81 lines

```
#include <bits/extc++.h>
const 11 INF = numeric limits<11>::max() / 4;
typedef vector<11> VL;
```

```
struct MCMF {
  int N;
  vector<vi> ed, red;
  vector<VL> cap, flow, cost;
  vi seen;
  VL dist, pi;
  vector<pii> par;
  MCMF (int N) .
   N(N), ed(N), red(N), cap(N, VL(N)), flow(cap), cost(cap)
    seen(N), dist(N), pi(N), par(N) {}
  void addEdge(int from, int to, ll cap, ll cost) {
    this->cap[from][to] = cap;
    this->cost[from][to] = cost;
    ed[from].push_back(to);
    red[to].push_back(from);
  void path(int s) {
    fill(all(seen), 0);
    fill(all(dist), INF);
    dist[s] = 0; ll di;
    __gnu_pbds::priority_queue<pair<11, int>> q;
    vector<decltype(q)::point_iterator> its(N);
    q.push({0, s});
    auto relax = [&](int i, ll cap, ll cost, int dir) {
      11 val = di - pi[i] + cost;
      if (cap && val < dist[i]) {
        dist[i] = val;
        par[i] = {s, dir};
        if (its[i] == q.end()) its[i] = q.push({-dist[i], i})
        else q.modify(its[i], {-dist[i], i});
    };
    while (!q.empty()) {
      s = q.top().second; q.pop();
      seen[s] = 1; di = dist[s] + pi[s];
      for (int i : ed[s]) if (!seen[i])
        relax(i, cap[s][i] - flow[s][i], cost[s][i], 1);
      for (int i : red[s]) if (!seen[i])
        relax(i, flow[i][s], -cost[i][s], 0);
   FOR(i, 0, N) pi[i] = min(pi[i] + dist[i], INF);
  pair<11, 11> maxflow(int s, int t) {
   11 \text{ totflow} = 0, \text{ totcost} = 0;
    while (path(s), seen[t]) {
      11 fl = INF;
      for (int p,r,x = t; tie(p,r) = par[x], x != s; x = p)
        fl = min(fl, r ? cap[p][x] - flow[p][x] : flow[x][p]
             ]);
      totflow += fl;
      for (int p,r,x = t; tie(p,r) = par[x], x != s; x = p)
        if (r) flow[p][x] += fl;
        else flow[x][p] -= fl;
```

```
FOR(i, 0, N) rep(j, 0, N) totcost += cost[i][j] * flow[i][j
        1:
   return {totflow, totcost};
  // If some costs can be negative, call this before maxflow
 void setpi(int s) { // (otherwise, leave this out)
   fill(all(pi), INF); pi[s] = 0;
   int it = N, ch = 1; 11 v;
   while (ch-- && it--)
     FOR(i, 0, N) if (pi[i] != INF)
        for (int to : ed[i]) if (cap[i][to])
         if ((v = pi[i] + cost[i][to]) < pi[to])</pre>
           pi[to] = v, ch = 1;
    assert(it >= 0); // negative cost cycle
};
```

Minimum Enclosing Circle.cpp

6d8e96, 44 lines

```
point a[maxn];
void getCenter2(point a, point b, point & c)
   c.x = (a.x+b.x)/2;
   c.y = (a.y+b.y)/2;
void getCenter3(point a, point b, point c, point &d)
   double a1 = b.x-a.x, b1 = b.y-a.y, c1 = (a1*a1+b1*b1)/2;
   double a2 = c.x-a.x, b2 = c.y-a.y, c2 = (a2*a2+b2*b2)/2;
   double de = (a1 * b2 - b1 * a2);
   d.x = a.x + (c1 * b2 - c2 * b1)/de;
   d.y = a.y + (a1 * c2 - a2 * c1)/de;
    //randomP_shuffle before using
    radius = 0:
   center = a[0];
    for (int i = 1; i < n; ++i)
        if (!isIn(a[i], center, radius))
            radius = 0;
            center = a[i];
            for (int j = 0; j < i; ++j)
                if (!isIn(a[j], center, radius))
                    getCenter2(a[i], a[j], center);
                    radius = dis(a[i],center);
                    for (int k = 0; k < j; ++k) if (!isIn(a[k
                         ], center, radius))
                        getCenter3(a[i], a[j], a[k], center)
                        radius = dis(a[k],center);
    //printf("\%.2lf \n\%.2lf \mw.2lf \mw.2lf \n", radius, center.x, center.y
    printf("%.31f\n", radius);
```

```
return 0;
}
```

MinimumVertexCover.cpp

Description: Finds a minimum vertex cover in a bipartite graph. The size is the same as the size of a maximum matching, and the complement is a maximum independent set.

```
"DFSMatching.h"
                                                  8c9aed, 20 lines
vi cover(vector<vi>& q, int n, int m) {
 vi match(m, -1);
 int res = dfsMatching(q, match);
 vector<bool> lfound(n, true), seen(m);
 for (int it : match) if (it != -1) lfound[it] = false;
 vi a, cover;
 FOR(i,0,n) if (lfound[i]) q.push_back(i);
 while (!a.emptv()) {
   int i = q.back(); q.pop_back();
   lfound[i] = 1;
    for (int e : q[i]) if (!seen[e] && match[e] != -1) {
     seen[e] = true;
     q.push_back(match[e]);
 FOR(i,0,n) if (!lfound[i]) cover.push_back(i);
 FOR(i,0,m) if (seen[i]) cover.push_back(n+i);
 assert(sz(cover) == res);
 return cover:
```

MinkowskiSum.cpp

```
18ebb1, 31 lines
void reorder_polygon(vector<pt> & P) {
    size t pos = 0;
    for(size_t i = 1; i < P.size(); i++) {</pre>
        if(P[i].y < P[pos].y \mid | (P[i].y == P[pos].y && P[i].
            x < P[pos].x)
            pos = i;
    rotate(P.begin(), P.begin() + pos, P.end());
vector<pt> minkowski(vector<pt> P, vector<pt> Q) {
    // the first vertex must be the lowest
    reorder polygon(P);
    reorder_polygon(Q);
    // we must ensure cyclic indexing
    P.push back(P[0]);
    P.push_back(P[1]);
    Q.push_back(Q[0]);
    Q.push_back(Q[1]);
    // main part
    vector<pt> result;
    size_t i = 0, j = 0;
    while (i < P.size() - 2 | | j < Q.size() - 2) {
        result.push_back(P[i] + Q[j]);
        auto cross = (P[i + 1] - P[i]).cross(Q[j + 1] - Q[j
            ]);
        if(cross >= 0)
            ++i;
        if(cross <= 0)
            ++j;
    return result;
```

```
Mo.cpp
void remove(idx); // TODO: remove value at idx from data
     structure
void add(idx);
                   // TODO: add value at idx from data
    structure
int get_answer(); // TODO: extract the current answer of
     the data structure
int block_size;
struct Query {
   int 1, r, idx:
   bool operator<(Query other) const
        return make_pair(l / block_size, r) <</pre>
               make_pair(other.l / block_size, other.r);
};
vector<int> mo_s_algorithm(vector<Query> queries) {
   vector<int> answers(queries.size());
   sort(queries.begin(), queries.end());
   // TODO: initialize data structure
   int cur_1 = 0;
   int cur_r = -1;
    // invariant: data structure will always reflect the
        range [cur_l, cur_r]
    for (Query q : queries) {
        while (cur_1 > q.1) {
            cur_1--;
            add(cur 1);
        while (cur_r < q.r) {
            cur_r++;
            add(cur r);
        while (cur_l < q.l) {</pre>
            remove(cur_l);
            cur_1++;
        while (cur_r > q.r) {
            remove(cur_r);
            cur_r--;
        answers[q.idx] = get_answer();
   return answers;
NTT.cpp
const 11 MOD = (119 << 23) + 1, root = 62; // = 998244353
// For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 <<
// and 483 \ll 21 (same root). The last two are > 10^9.
typedef vector<ll> vl;
void ntt(vl &a) {
 int n = sz(a), L = 31 - builtin clz(n);
 static vl rt(2, 1);
 for (static int k = 2, s = 2; k < n; k *= 2, s++) {
   rt.resize(n);
   ||z|| = \{1, modExp(root, MOD >> s)\};
   FOR(i,k,2*k) rt[i] = rt[i / 2] * z[i & 1] % MOD;
```

```
vi rev(n);
 FOR(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
 FOR(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
  for (int k = 1; k < n; k *= 2)
    for (int i = 0; i < n; i += 2 * k) FOR(j,k) {
     ll z = rt[j + k] * a[i + j + k] % MOD, &ai = a[i + j];
     a[i + j + k] = ai - z + (z > ai ? MOD: 0);
     ai += (ai + z >= MOD ? z - MOD : z);
vl conv(const vl &a, const vl &b) {
 if (a.empty() || b.empty()) return {};
 int s = sz(a) + sz(b) - 1, B = 32 - \underline{builtin\_clz(s)}, n =
      1 << B;
  int inv = modExp(n, MOD - 2);
  vl L(a), R(b), out(n);
 L.resize(n), R.resize(n);
 ntt(L), ntt(R);
 FOR(i,n) out[-i \& (n-1)] = (11)L[i] * R[i] % MOD * inv %
 ntt(out);
  return {out.begin(), out.begin() + s};
orderedset.cpp
<ext/pb.ds/assoc_container.hpp>
                                                   647225, 4 lines
using namespace __gnu_pbds;
template<typename T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
PlanarPointLocation.cpp
                                                 2512ba, 199 lines
* CP algorithm point point_location
* This implementation assumes that the subdivision is
     correctly stored inside a DCEL
  and the outer face is numbered -1.
  For each query a pair (1,i) is returned if point is
       strictly inside face number i.
  (0,i) returned if point lies on the edge number i.
bool edge_cmp(Edge* edge1, Edge* edge2)
   const pt a = edge1->1, b = edge1->r;
   const pt c = edge2->1, d = edge2->r;
   int val = sgn(a.cross(b, c)) + sgn(a.cross(b, d));
   if (val != 0)
        return val > 0;
    val = sgn(c.cross(d, a)) + sgn(c.cross(d, b));
    return val < 0;
enum EventType { DEL = 2, ADD = 3, GET = 1, VERT = 0 };
struct Event {
   EventType type;
    int pos;
   bool operator<(const Event& event) const { return type <
         event.type; }
```

};

```
vector<Edge*> sweepline(vector<Edge*> planar, vector<pt>
    queries)
    using pt_type = decltype(pt::x);
    // collect all x-coordinates
    auto s =
        set<pt_type, std::function<bool(const pt_type&,
            const pt_type&)>>(lt);
    for (pt p : queries)
        s.insert(p.x);
    for (Edge* e : planar) {
       s.insert(e->1.x);
        s.insert(e->r.x);
    // map all x-coordinates to ids
    int cid = 0:
    auto id =
        map<pt_type, int, std::function<bool(const pt_type&,</pre>
              const pt_type&)>>(
            lt);
    for (auto x : s)
        id[x] = cid++;
    // create events
    auto t = set<Edge*, decltype(*edge_cmp)>(edge_cmp);
    auto vert_cmp = [](const pair<pt_type, int>& 1,
                       const pair<pt_type, int>& r) {
        if (!eq(l.first, r.first))
            return lt(l.first, r.first);
        return 1.second < r.second;
    auto vert = set<pair<pt_type, int>, decltype(vert_cmp)>(
         vert_cmp);
    vector<vector<Event>> events(cid);
    for (int i = 0; i < (int) queries.size(); i++) {
        int x = id[queries[i].x];
        events[x].push_back(Event{GET, i});
    for (int i = 0; i < (int)planar.size(); <math>i++) {
        int lx = id[planar[i]->1.x], rx = id[planar[i]->r.x
            1:
        if (lx > rx) {
            swap(lx, rx);
            swap(planar[i]->1, planar[i]->r);
        if (lx == rx) {
            events[lx].push_back(Event{VERT, i});
        } else {
            events[lx].push_back(Event{ADD, i});
            events[rx].push back(Event{DEL, i});
    // perform sweep line algorithm
    vector<Edge*> ans(gueries.size(), nullptr);
    for (int x = 0; x < cid; x++) {
        sort(events[x].begin(), events[x].end());
        vert.clear();
        for (Event event : events[x]) {
            if (event.type == DEL) {
               t.erase(planar[event.pos]);
            if (event.type == VERT) {
```

```
vert.insert(make pair(
            min(planar[event.pos]->1.y, planar[event
                .pos]->r.y),
            event.pos));
   if (event.type == ADD) {
        t.insert(planar[event.pos]);
   if (event.type == GET) {
        auto jt = vert.upper_bound(
            make_pair(queries[event.pos].y, planar.
                size()));
        if (jt != vert.begin()) {
            --jt;
            int i = jt->second;
            if (ge(max(planar[i]->1.y, planar[i]->r.
                   queries[event.pos].y)) {
                ans[event.pos] = planar[i];
                continue;
        Edge* e = new Edge;
        e->1 = e->r = queries[event.pos];
        auto it = t.upper_bound(e);
        if (it != t.begin())
            ans[event.pos] = *(--it);
        delete e;
for (Event event : events[x]) {
   if (event.type != GET)
        continue:
   if (ans[event.pos] != nullptr &&
        eg(ans[event.pos]->1.x, ans[event.pos]->r.x)
        continue;
    Edge* e = new Edge;
   e->1 = e->r = queries[event.pos];
   auto it = t.upper_bound(e);
   delete e:
   if (it == t.begin())
        e = nullptr;
        e = *(--it);
   if (ans[event.pos] == nullptr) {
        ans[event.pos] = e;
        continue;
   if (e == nullptr)
        continue:
   if (e == ans[event.pos])
        continue;
   if (id[ans[event.pos]->r.x] == x) {
        if (id[e->1.x] == x) {
            if (gt(e->1.y, ans[event.pos]->r.y))
                ans[event.pos] = e;
   } else {
        ans[event.pos] = e;
```

```
return ans;
struct DCEL {
   struct Edge {
        pt origin;
        Edge* nxt = nullptr:
        Edge* twin = nullptr;
        int face;
   };
    vector<Edge*> body;
vector<pair<int, int>> point_location(DCEL planar, vector<pt
    > queries)
    vector<pair<int, int>> ans(queries.size());
    vector<Edge*> planar2;
    map<intptr_t, int> pos;
    map<intptr_t, int> added_on;
    int n = planar.body.size();
    for (int i = 0; i < n; i++) {
        if (planar.body[i]->face > planar.body[i]->twin->
             face)
            continue;
        Edge * e = new Edge;
        e->1 = planar.body[i]->origin;
        e->r = planar.body[i]->twin->origin;
        added_on[(intptr_t)e] = i;
        pos[(intptr_t)e] =
            lt(planar.body[i]->origin.x, planar.body[i]->
                twin->origin.x)
                ? planar.body[i]->face
                : planar.body[i]->twin->face;
        planar2.push_back(e);
    auto res = sweepline(planar2, queries);
    for (int i = 0; i < (int) queries.size(); i++) {
        if (res[i] == nullptr) {
            ans[i] = make_pair(1, -1);
            continue;
        pt p = queries[i];
        pt 1 = res[i] ->1, r = res[i] ->r;
        if (eq(p.cross(1, r), 0) && le(p.dot(1, r), 0)) {
            ans[i] = make_pair(0, added_on[(intptr_t)res[i
                11);
            continue:
        ans[i] = make_pair(1, pos[(intptr_t)res[i]]);
    for (auto e : planar2)
        delete e:
    return ans:
```

PersistentSegtree.cpp

ed1804, 86 lines

16

```
//Define the node of a persistent segment tree
struct node{
   int 1, r, sum;
//the persistent segment tree. Warning: Check memory limit
    before using persistent segment tree!
node tree[maxn*32];
```

```
//Storing the root of versions of segment tree
int head[maxn];
//allocate next position. You can implement in a way that
    support garbage collection.
int nextPos(){
    static int ct; return ++ct;
//Building the first version of our segmetn tree
void build(int cur,int 1,int r){
 tree[cur].sum = 0;
 tree[cur].l = nextPos();
 tree[cur].r = nextPos();
 if (1 == r) {
   tree[cur].1 = tree[cur].r = 0;
 else{
    int mid = (1+r) >> 1;
   build(tree[cur].1,1,mid);
   build(tree[cur].r,mid+1,r);
//This function is: currently we are at node cur, which is a
     node in the latest version of segment tree
//we want to make modifications based on some past segment
    tree, and the corresponding node in the last version is
    at last
//we want to add 1 at position key
void modify(int cur,int last,int key,int l,int r){
  //this is creating the node for our latest version
 tree[curl.sum = tree[last].sum;
 tree[cur].1 = nextPos();
 tree[curl.r = nextPos();
 if (1 == r) {
    //base case:add on current version of our segment tree
    tree[cur].sum++;
   tree[cur].1 = tree[cur].r = 0;
 else{
    int mid = (1+r) >> 1;
    if (key <= mid) {</pre>
      //we are going to modify in the left part, so we can
           reuse the right child
     tree[cur].r = tree[last].r;
      modifv(tree[cur].1, tree[last].1, key,1, mid);
      //update information for the current version of
           seament tree
     tree[cur].sum++;
    else
     tree[cur].l = tree[last].l;
     modify(tree[cur].r, tree[last].r, key,mid+1, r);
     tree[cur].sum++;
int query(int cur,int last,int l,int r,int k){
 if (1 == r) return 1;
 int mid = (1+r) >> 1;
 //notice the subtraction here - we want too see the
       dfiffernce between today's version and old versions.
 int ct = tree[tree[cur].1].sum - tree[tree[last].1].sum ;
```

```
//if there are to many larger than mid, the k-th element
       should be in the left
 if (ct >= k) {
   return query(tree[cur].1, tree[last].1, 1, mid, k);
  //otherwise, the k-th element should be in the right
 elset
   return query(tree[cur].r, tree[last].r, mid+1, r, k-ct);
//Build segment tree to support gueres k-th element in a
    subinterval
void build(int n) {
 for (int i = 0; i \le n; ++i) {
   head[i] = nextPos();
 build(head[0], 1, n);
 for (int i = 1; i <= n; ++i) {
   modify(head[i], head[i-1],c[i], 1, n);
/*Query the k-th element in [l,r]:
printf("%d\n", a[query(head[r], head[l-1], 1, n, k)].key); */
```

Point3d.cpp

8058ae, 32 lines

```
template<class T> struct Point3D {
  typedef Point3D P:
  typedef const P& R;
  T x, y, z;
  explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z) {
  bool operator<(R p) const {</pre>
   return tie(x, y, z) < tie(p.x, p.y, p.z); }
  bool operator==(R p) const {
   return tie(x, y, z) == tie(p.x, p.y, p.z); }
  P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z); }
  P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
  P operator*(T d) const { return P(x*d, y*d, z*d); }
  P operator/(T d) const { return P(x/d, y/d, z/d); }
  T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
  P cross(R p) const {
    return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
  T dist2() const { return x*x + y*y + z*z; }
  double dist() const { return sqrt((double)dist2()); }
  //Azimuthal angle (longitude) to x-axis in interval [-pi,
  double phi() const { return atan2(y, x); }
  //Zenith angle (latitude) to the z-axis in interval [0, pi
  double theta() const { return atan2(sqrt(x*x+y*y),z); }
  P unit() const { return *this/(T)dist(); } //makes dist()
  //returns unit vector normal to *this and p
  P normal(P p) const { return cross(p).unit(); }
  //returns point rotated 'angle' radians ccw around axis
  P rotate (double angle, P axis) const {
    double s = sin(angle), c = cos(angle); P u = axis.unit()
    return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
};
```

```
Point.cpp
```

17

```
template <class T> int sgn(T x) { return (x > 0) - (x < 0);
template<class T>
struct Point {
 typedef Point P;
 T x, y;
  explicit Point (T x=0, T y=0) : x(x), y(y) {}
 bool operator (P p) const \{ return tie(x,y) < tie(p.x,p.y) \}
  bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y)
      ; }
  P operator+(P p) const { return P(x+p.x, y+p.y); }
  P operator-(P p) const { return P(x-p.x, y-p.y); }
  P operator*(T d) const { return P(x*d, y*d); }
  P operator/(T d) const { return P(x/d, y/d); }
  T dot(P p) const { return x*p.x + y*p.y; }
 T cross(P p) const { return x*p.y - y*p.x; }
 T cross(P a, P b) const { return (a-*this).cross(b-*this);
 T dist2() const { return x*x + y*y; }
  double dist() const { return sqrt((double)dist2()); }
  // angle to x-axis in interval [-pi, pi]
  double angle() const { return atan2(y, x); }
  P unit() const { return *this/dist(); } // makes dist()=1
  P perp() const { return P(-y, x); } // rotates +90 degrees
  P normal() const { return perp().unit(); }
  // returns point rotated 'a' radians ccw around the origin
  P rotate(double a) const {
   return P(x*\cos(a)-y*\sin(a),x*\sin(a)+y*\cos(a)); }
  friend ostream& operator << (ostream& os, P p) {
   return os << "(" << p.x << "," << p.y << ")"; }
//Line distance
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
 return (double) (b-a).cross(p-a)/(b-a).dist();
//LineProjectionReflection
template<class P>
P lineProj(P a, P b, P p, bool refl=false) {
 P v = b - a;
 return p - v.perp() * (1+refl) *v.cross(p-a) /v.dist2();
//Point-Segment Distance
typedef Point < double > P;
double segDist(P& s, P& e, P& p) {
 if (s==e) return (p-s).dist();
 auto d = (e-s) . dist2(), t = min(d, max(.0, (p-s) . dot(e-s)));
 return ((p-s)*d-(e-s)*t).dist()/d;
//Segment-Segment Distance
double TwoSegMinDist(P A, P B, P C, P D)
    return min (min (segDist (A, B, C), segDist (A, B, D)),
               min(segDist(C,D,A),segDist(C,D,B)));
//On Seament
template < class P > bool on Segment (P s, P e, P p) {
 return p.cross(s, e) == 0 \&\& (s - p).dot(e - p) <= 0;
```

```
//Segment Intersection
/*If a unique intersection point between the line segments
    going from s1 to e1 and from s2 to e2 exists then it is
    returned.
If no intersection point exists an empty vector is returned.
      If infinitely many exist a vector with 2 elements is
    returned, containing the endpoints of the common line
    seament.
The wrong position will be returned if P is Point<1l> and
    the intersection point does not have integer coordinates
Products of three coordinates are used in intermediate steps
      so watch out for overflow if using int or long long.*/
template<class P> vector<P> segInter(P a, P b, P c, P d) {
 auto oa = c.cross(d, a), ob = c.cross(d, b),
      oc = a.cross(b, c), od = a.cross(b, d);
  // Checks if intersection is single non-endpoint point.
 if (sqn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
   return { (a * ob - b * oa) / (ob - oa) };
  set<P> s:
  if (onSegment(c, d, a)) s.insert(a);
  if (onSegment(c, d, b)) s.insert(b);
  if (onSegment(a, b, c)) s.insert(c);
 if (onSegment(a, b, d)) s.insert(d);
  return {all(s)};
//The notations from here might be different
//Find the symmetric point of point p about line p1p2
Point SymPoint (Point p, Line 1)
    Point result:
    double a=1.p2.x-l.p1.x;
    double b=1.p2.v-1.p1.v;
    double t = ((p.x-1.p1.x)*a+(p.y-1.p1.y)*b)/(a*a+b*b);
    result.x=2*1.p1.x+2*a*t-p.x;
    result.v=2*1.p1.v+2*b*t-p.v;
    return result;
PollardRho.cpp
                                                  d9508f, 82 lines
//Quick\ Multiplication - Calculate\ x\ *\ y\ mod\ modi
     efficiently
```

```
11 d = (p - 1) / (111 << r);
    for (int i = 0; i < 12; i++) {
        if (p == tests[i]) {
            return true;
        11 x = quickpow2(tests[i], d, p);
        for (int j = 1; j <= r; j++) {
            11 y = quickmult(x, x, p);
            if (y == 1 \&\& x != 1 \&\& x != p - 1) return false
            x = v;
        if (x != 1) return false;
    return true;
//We will store factors in a global variable to save time
map<11, int> factors;
11 gcd(11 x, 11 y) {
    if (y == 0) return x;
    return gcd(y, x % y);
l get_next(ll x, ll addi,ll modi){
    return (quickmult(x, x, modi) + addi);
//find a prime factor of n based on the seed, if we cannot
     find it return -1
ll rho_find(ll n, ll seed, ll addi) {
    11 a = seed;
    11 b = get_next(seed, addi, n);
    while (a != b) {
        11 p = gcd(abs(a - b), n);
        if (p > 1) {
            if (p == n) return -1;
            return p;
        a = get next(a, addi, n);
        b = get_next(get_next(b, addi, n), addi, n);
    return -1;
//factorizing n via pollard rho
void pollard rho(ll n) {
    if (n == 1) {
        return;
    if (prime test(n)){
        factors[n]++;
        return:
    11 p = -1;
    while (p == -1) {
        11 addi = rand() % (n - 1) + 1;
        p = rho_find(n, rand() % (n - 1) + 1, addi);
        if (p != -1) {
            pollard_rho(p);
            pollard_rho(n / p);
            return;
```

```
PolygonUnion.cpp
```

```
"Point.h", "sideOf.h"
                                                  b73c5a, 39 lines
/* KACTL Polygon Union
* Description: Calculates the area of the union of $n$
     polygons (not necessarily
* convex). The points within each polygon must be given in
     CCW order.
 * (Epsilon checks may optionally be added to sideOf/sqn,
     but shouldn't be needed.)
 * Time: $O(N^2)$. where $N$ is the total number of points*/
typedef Point < double > P;
double rat(P a, P b) { return sqn(b.x) ? a.x/b.x : a.y/b.y;
double polyUnion(vector<vector<P>>& poly) {
 double ret = 0;
 FOR(i, 0, sz(poly)) rep(v, 0, sz(poly[i])) {
   P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
   vector<pair<double, int>> segs = {{0, 0}, {1, 0}};
   FOR(j, 0, sz(poly)) if (i != j) {
     FOR(u, 0, sz(poly[j])) {
       P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])
            ];
        int sc = sideOf(A, B, C), sd = sideOf(A, B, D);
        if (sc != sd) {
         double sa = C.cross(D, A), sb = C.cross(D, B);
          if (min(sc, sd) < 0)
           segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
        } else if (!sc && !sd && j<i && sgn((B-A).dot(D-C))
          segs.emplace_back(rat(C - A, B - A), 1);
          segs.emplace_back(rat(D - A, B - A), -1);
   sort(all(segs));
    for (auto& s : segs) s.first = min(max(s.first, 0.0),
        1.0);
   double sum = 0;
   int cnt = segs[0].second;
   FOR(j, 1, sz(segs)) {
     if (!cnt) sum += segs[j].first - segs[j - 1].first;
     cnt += seqs[j].second;
   ret += A.cross(B) * sum;
 return ret / 2;
```

Polynomial.cpp

92b62a, 87 lines

```
constexpr int maxn = 262144;

constexpr int mod = 998244353;

using i64 = long long;

using poly_t = int[maxn];

using poly = int *const;

//Find f^{-1} mod x^{n}.

void polyinv(const poly &h, const int n, poly &f) {

/* f = 1 / h = f_{-0} (2 - f_{-0} h) */

static poly_t inv_t;

std::fill(f, f + n + n, 0);
```

PushRelabel QuasiExgcdSum

```
f[0] = fpow(h[0], mod - 2);
  for (int t = 2; t <= n; t <<= 1) {
    const int t2 = t \ll 1;
    std::copy(h, h + t, inv_t);
    std::fill(inv_t + t, inv_t + t2, 0);
    DFT(f, t2);
    DFT(inv t, t2);
    for (int i = 0; i != t2; ++i)
      f[i] = (i64) f[i] * sub(2, (i64) f[i] * inv_t[i] % mod)
    IDFT(f, t2);
    std::fill(f + t, f + t2, 0);
}
//Find h(x) such that h^2(x) = f(x) \mod x^{deg}.
inline void sqrt(int deg, int *f, int *h) {
 if (deg == 1) {
   h[0] = 1;
   return;
  sqrt(deq + 1 >> 1, f, h);
  int len = 1;
 while (len < deg * 2) { // doubling
   len *= 2:
  fill(g, g + len, 0);
  inv(deg, h, g);
  copv(f, f + deq, t);
  fill(t + deg, t + len, 0);
 NTT(t, len, 1);
 NTT(q, len, 1);
 NTT(h, len, 1);
  for (int i = 0; i < len; i++) {
   h[i] = (long long)1 * inv2 *
           ((long long)1 * h[i] % mod + (long long)1 * q[i]
                * t[i] % mod) % mod;
 NTT(h, len, -1);
 fill(h + deg, h + len, 0);
/*This is Fast Walsh Transformation
Goal: Given A, B, compute C_{-i} = \sum_{j=1}^{n} A_{-j}B_{-k}
? is or, and, xor*/
void FWT(int *f.int pd){
    for (int d=1; d<n; d<<=1)</pre>
        for (int m=d<<1, i=0; i<n; i+=m)</pre>
            for (int j=0; j<d; j++) {</pre>
                 int x=f[i+j],y=f[i+j+d];
                if (pd==0) f[i+j+d]=(x+y)%p; // or
                if (pd==1) f[i+j]=(x+y) p; // and
                 if(pd==2) f[i+j]=(x+y) p, f[i+j+d]=(x-y+p) p;
void IFWT(int *f,int pd){
    for (int d=1; d<n; d<<=1)
        for (int m=d<<1, i=0; i<n; i+=m)</pre>
            for (int j=0; j<d; j++) {</pre>
                int x=f[i+j],y=f[i+j+d];
                 if (pd==0) f[i+j+d]=(y-x+p)%p; //OR
                if (pd==1) f[i+j]=(x-y+p) p;; // AND
```

```
if(pd==2) f[i+j]=111*(x+y)*inv*p, f[i+j+d]=1
                      11*(x-y+p)*inv%p; //XOR
void solve or()
    memcpy(a,A,sizeof a);
    memcpy(b,B,sizeof b);
    FWT(a,0); FWT(b,0);
    for(int i=0;i<n;i++)</pre>
        a[i]=111*a[i]*b[i]%p;
    IFWT(a,0);
PushRelabel.cpp
Description: Push-relabel using the highest label selection rule and the
gap heuristic. Quite fast in practice. To obtain the actual flow, look at
positive values only.
Time: \mathcal{O}\left(V^2\sqrt{E}\right)
                                                      be9cf6, 48 lines
struct PushRelabel {
 struct Edge {
    int dest, back;
    11 f, c;
 };
  vector<vector<Edge>> q;
  vector<11> ec;
  vector<Edge*> cur;
  vector<vi> hs; vi H;
  PushRelabel(int n): g(n), ec(n), cur(n), hs(2*n), H(n) {}
  void addEdge(int s, int t, ll cap, ll rcap=0) {
   if (s == t) return;
    g[s].push_back({t, sz(g[t]), 0, cap});
    g[t].push_back({s, sz(g[s])-1, 0, rcap});
  void addFlow(Edge& e, ll f) {
    Edge &back = g[e.dest][e.back];
    if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
    e.f += f; e.c -= f; ec[e.dest] += f;
    back.f -= f; back.c += f; ec[back.dest] -= f;
  11 calc(int s, int t) {
    int v = sz(q); H[s] = v; ec[t] = 1;
    vi co(2*v); co[0] = v-1;
    FOR(i, 0, v) cur[i] = q[i].data();
    for (Edge& e : g[s]) addFlow(e, e.c);
    for (int hi = 0;;) {
      while (hs[hi].empty()) if (!hi--) return -ec[s];
      int u = hs[hi].back(); hs[hi].pop back();
      while (ec[u] > 0) // discharge u
        if (\operatorname{cur}[u] == \operatorname{g}[u].\operatorname{data}() + \operatorname{sz}(\operatorname{g}[u])) {
          H[u] = 1e9;
           for (Edge& e : q[u]) if (e.c && H[u] > H[e.dest
            H[u] = H[e.dest]+1, cur[u] = &e;
          if (++co[H[u]], !--co[hi] && hi < v)
             FOR(i,0,v) if (hi < H[i] && H[i] < v)
               --co[H[i]], H[i] = v + 1;
          hi = H[u];
        } else if (cur[u]->c && H[u] == H[cur[u]->dest]+1)
          addFlow(*cur[u], min(ec[u], cur[u]->c));
```

```
else ++cur[u];
 bool leftOfMinCut(int a) { return H[a] >= sz(g); }
QuasiExgcdSum.cpp
                                                    b1006b, 60 lines
Using Quasi_Exqcd to sum
f(a,b,c,n) = sum_{\{i=0\}}^n \setminus floor_{\{(ai+b)/c\}}
g(a,b,c,n) = sum_{\{i=0\}} \hat{n} \setminus floor i_{\{(ai+b)/c\}}
h(a,b,c,n) = sum_{\{i=0\}} n \left(\frac{1-0}{n} \left(\frac{ai+b}{c}\right)/c\right)^2
all are done under mod p
struct rec{
   11 f, q, h;
//add, sub, quickpow omitted
11 inv2 = quickpow(2, modi-2);
11 inv6 = quickpow(6, modi - 2);
rec solve(ll a, ll b, ll c, ll n) {
 rec ans;
 if (a == 0) {
   ans.f = (b / c) * (n + 1) % modi;
    ans.g = (b / c) * (n + 1) % modi * n % modi * inv2 %
    ans.h = (b / c) * (b / c) % modi * (n+1) % modi;
    return ans;
  ans.f = ans.q = ans.h = 0;
  if (a >= c || b >= c){
    rec temp = solve(a % c, b % c, c, n);
    add(ans.f, (a/c) *n%modi*(n+1)%modi*inv2%modi);
    add(ans.f, (b/c)*(n+1)%modi);
    add(ans.f, temp.f);
    add(ans.g, (a/c) *n%modi*(n+1)%modi*
    ((2*n+1)%modi)%modi*inv6%modi);
    add(ans.q, (b/c)*n%modi*(n+1)%modi*inv2 % modi);
    add(ans.q, temp.q);
    add(ans.h, (a/c) * (a/c) %modi*n%modi*
    (n+1) %modi * ((2*n+1) %modi) %modi * inv6% modi);
    add(ans.h, (b/c)*(b/c)*modi*(n+1)*modi);
    add(ans.h, (a/c)*(b/c)%modi*n%modi*(n+1)%modi);
    add(ans.h, temp.h);
    add(ans.h, 2LL * (a/c)%modi*temp.g%modi);
    add(ans.h, 2LL * (b/c)%modi*temp.f%modi);
    return ans:
  if (a < c && b < c) {
   11 m = (a * n + b) / c;
    rec temp = solve(c, c - b - 1, a, m - 1);
    ans.f = n * m % modi;
    sub(ans.f, temp.f);
    ans.q = n * (n + 1) % modi * m % modi;
    sub(ans.q, temp.f);
    sub(ans.q, temp.h);
    ans.g = ans.g * inv2 % modi;
    ans.h = n * m % modi * (m + 1) % modi;
    sub(ans.h, 2LL * temp.q % modi);
    sub(ans.h, 2LL * temp.f % modi);
    sub(ans.h, ans.f);
```

return ans:

return ans;

QuickPhiSum RootNonTree Sam SCC Schreier-Sims

```
QuickPhiSum.cpp
                                                     024bac, 29 lines
This algorithm concerns efficient evaluation of sum of
    number theoric functions like phi or mu.
We know that using Eulerian sieve, we can only archieve O(n)
      time complexity.
What we are doing is to archieve O(n^{2/3}) time complexity.
The example program shows how to evaluate sum of phi and sum
      of mu efficientyl.
For smaller n (n less than (N^{2}/3)), we use calculate them
For larger n, see getphi and getmu
//See Sieve for more technical details.
//When i is prime, phi[i] = i-1, mu[i] = -1; Otherwise,
     inside the inner loop, let p = primes[j].
//Then it follows phi[p*i] = phi[i] * (p-1); mu[p*i] = -mu[i]
//finally, when i \% p == 0, phi[p*i] = phi[i]*p and mu[p*i]
    = 0:
ll getphi(ll n)
  if (n <= m) return phi[n];</pre>
 if (phi_cheat.find(n) != phi_cheat.end()) return phi_cheat
 ll ans = (ll) n*(n + 1) / 2; //this is <math>\sum_{i=1}^{n} \sum_{i=1}^{n} \sum_{i=1}^{n} t_{i}
       \{d \mid n\} \setminus phi(d)
  //when getting mu, ans = 1
  ll last;
  for (11 i = 2; i \le n; i = last + 1)
    last = n / (n / i);
    ans -= (last-i+1) *getphi(n / i);
 phi_cheat[n] = ans;
 return ans;
```

RootNonTree.cpp

9d4bb6, 38 lines

```
siz[u]++;
void paint(int u,int c){
    if (col[u]) c = col[u];
    else col[u] = c;
    for (int i = 0; i < g[u].size(); i++) {
        int v = q[u][i];
        if (v != p[u]){
            paint(v,c);
//actual blokcing; magic = block size
dfs(1);
if (!tot_cols) {
    cap[++tot\_cols] = 1;
paint(1,tot_cols);
Sam.cpp
<cstdio>, <cstdlib>, <cstring>, <algorithm>
                                                   6196dc, 58 lines
using namespace std;
typedef long long 11;
const int maxn = 510000;
const int sigma = 26;
struct edge
   int v, next;
};
edge g[maxn << 1];
int trie[maxn << 1][siqma], fa[maxn << 1], maxi[maxn << 1],</pre>
    sizia[maxn << 1];
char str[maxn];
int head[maxn << 1];</pre>
int siz,last;
void insert(int u, int v)
    static int id:
    g[++id].v = v;
    g[id].next = head[u];
    head[u] = id;
//This is the core of SAM
void add(int id)
    int p = last;
    int np = last = ++siz;
    sizia[np] = 1;
    maxi[np] = maxi[p] + 1;
    while (p && !trie[p][id]) {
        trie[p][id] = np;
        p = fa[p];
   if (!p){
        fa[np] = 1;
```

```
fa[np] = q;
        else{
            int nq = ++siz;
            maxi[nq] = maxi[p] + 1;
            memcpy(trie[nq], trie[q], sizeof trie[q]);
            fa[nq] = fa[q];
            fa[np] = fa[q] = nq;
            while (trie[p][id] == q) {
                trie[p][id] = nq;
                p = fa[p];
SCC.cpp
                                                   a6b342, 37 lines
vector < vector<int> > q, qr; //q stores graph, gr stores
     graph transposed
vector<bool> used;
vector<int> order, component;
void dfs1 (int v) {
   used[v] = true;
    for (size t i=0; i<q[v].size(); ++i)</pre>
        if (!used[ q[v][i] ])
           dfs1 (q[v][i]);
   order.push_back (v);
void dfs2 (int v) {
   used[v] = true;
    component.push_back (v);
    for (size_t i=0; i<gr[v].size(); ++i)</pre>
        if (!used[ gr[v][i] ])
            dfs2 (gr[v][i]);
void findSCCs() {
   order.clear();
   used.assign (n, false);
    for (int i=0; i<n; ++i)</pre>
        if (!used[i])
           dfs1 (i);
   used.assign (n, false);
   for (int i=0; i<n; ++i) {
        int v = order[n-1-i];
        if (!used[v]) {
            dfs2 (v);
            //SCC FOUND, DO SOMETHING
            component.clear();
    }
Schreier-Sims.cpp
                                                   a288c1, 102 lines
```

// time complexity : $O(n^2 \log^3 |G| + t \log |G|)$

// memory complexity : $O(n^2 \log |G| + tn)$

// t : number of generators

else{

int q = trie[p][id];

if (maxi[p] + 1 == maxi[q]) {

```
//|G|: group size, obviously \leq (n!)
vector<int> inv(vector<int>& p) {
 vector<int> ret(p.size());
 for (int i = 0; i < p.size(); i++) ret[p[i]] = i;
 return ret;
vector<int> operator * (vector<int>& a, vector<int>& b ) {
 vector<int> ret(a.size());
 for (int i = 0; i < a.size(); i++) ret[i] = b[a[i]];
 return ret;
// a group contains all subset products of generators
struct Group {
 int n, m;
 vector<vector<int>> lookup;
 vector<vector<int>>> buckets, ibuckets;
 int yo(vector<int> p, bool add_to_group = 1){
   n = buckets.size();
    for (int i = 0; i < n; i++) {
     int res = lookup[i][p[i]];
     if (res == -1) {
       if (add_to_group) {
         buckets[i].push_back(p);
         ibuckets[i].push_back(inv(p));
         lookup[i][p[i]] = buckets[i].size() - 1;
       return i;
     p = p * ibuckets[i][res];
    return -1;
  ll size() {
    11 ret = 1:
    for (int i = 0; i < n; i++) ret *= buckets[i].size();</pre>
    return ret:
 bool in_group(vector<int> g) { return yo(g, false) == -1;
 Group(vector<vector<int>> &gen, int n){
    n = _n, m = gen.size(); // m permutations of size <math>n, 0
        indexed
    lookup.resize(n);
    buckets.resize(n);
    ibuckets.resize(n);
    for (int i = 0; i < n; i++) {
     lookup[i].resize(n);
     fill(lookup[i].begin(), lookup[i].end(), -1);
    vector<int> id(n):
    for (int i = 0; i < n; i++) id[i] = i;
    for (int i = 0; i < n; i++) {
     buckets[i].push_back(id);
     ibuckets[i].push_back(id);
     lookup[i][i] = 0;
    for (int i = 0; i < m; i++) yo(gen[i]);
    queue<pair<pair<int, int>,pair<int, int>>> q;
    for (int i = 0; i < n; i++) {
      for (int j = i; j < n; j++) {
       for (int k = 0; k < buckets[i].size(); k++) {</pre>
          for (int 1 = 0; 1 < buckets[j].size(); 1++) {</pre>
            q.push({pair<int, int>(i, k), pair<int, int>(j,
                 1) });
```

```
while(!q.empty()) {
      pair<int, int> a = q.front().first;
      pair<int, int> b = q.front().second;
      int res = yo(buckets[a.first][a.second] * buckets[b.
           first][b.second]);
      if (res == -1) continue;
      pair<int, int> cur(res, (int)buckets[res].size() - 1);
      for (int i = 0; i < n; i ++) {
        for (int j = 0; j < (int)buckets[i].size(); ++j){</pre>
          if (i <= res) q.push(make_pair(pair<int, int>(i ,
               j), cur));
          if (res <= i) q.push(make_pair(cur, pair<int, int</pre>
              >(i, j)));
};
int32_t main() {
 ios_base::sync_with_stdio(0);
 cin.tie(0);
 int k, n; cin >> k >> n;
 vector<vector<int>> a;
 while (k--) {
  vector<int> v;
   for (int i = 0; i < n; i++) {
    int x; cin >> x;
    v.push back(x - 1);
  a.push_back(v);
  Group q(a, n);
  cout << q.size() << '\n';
 return 0;
```

Segtree.cpp

c1f7d5, 18 lines

cat input

cat outWrong

cat outSlow

echo "Wrong Output:"

echo "Slow Output:"

```
11 SZ = 262144; //set this to power of two
11* seg = new 11[2*SZ]; //segtree implementation by bqi343's
11 combine(11 a, 11 b) { return max(a, b); }
void update(int p, ll value) {
    for (seq[p += SZ] = value; p > 1; p >>= 1)
        seg[p>>1] = combine(seg[(p|1)^1], seg[p|1]);
ll query(int l, int r) { // sum on interval [l, r]
    11 \text{ resL} = 0, \text{ resR} = 0; \text{ r++;}
    for (1 += SZ, r += SZ; 1 < r; 1 >>= 1, r >>= 1) {
        if (1&1) resL = combine(resL, seg[1++]);
        if (r&1) resR = combine(seq[--r], resR);
    return combine (resL, resR);
```

```
Sieve.cpp
                                                   666866, 16 lines
vi primes, leastFac;
void compPrimes(int N) {
 FOR(i, N) {
   leastFac.pb(0);
 leastFac[0] = 1; leastFac[1] = 1;
  FOR(i, 2, N) {
   if (leastFac[i] == 0) {
     primes.pb(i);
     leastFac[i] = i;
    for (int j = 0; j < sz(primes) && i*primes[j] < N &&
        primes[j] <= leastFac[i]; j++) {</pre>
     leastFac[i*primes[j]] = primes[j];
Simpson.cpp
                                                   88ae08 20 lines
This is a template for solving simpson integration
We are going to integrate \frac{cx+d}{ax+b} over L and R
ld simpson(ld lower, ld upper) {
   ld mid = (lower + upper) / 2;
   return (f(lower) + 4 * f(mid) + f(upper)) * (upper-lower
        ) / 6;
ld simpson_integration(ld lower, ld upper, ld target) {
   ld mid = (lower + upper) / 2;
   ld left_sum = simpson(lower, mid);
   ld right_sum = simpson(mid, upper);
   if (fabs(left_sum + right_sum - target) < eps){</pre>
        return left_sum + right_sum;
    return simpson_integration(lower, mid, left_sum) +
        simpson_integration(mid, upper, right_sum);
//Call: simpson_integration(lower, upper, simpson(lower,
     upper)) << endl;
stress.cpp
                                                   95d57c, 22 lines
#!/usr/bin/env bash
for ((testNum=0;testNum<$4;testNum++))</pre>
    ./$3 > input
    ./$2 < input > outSlow
    ./$1 < input > outWrong
   H1= 'md5sum outWrong'
   H2='md5sum outSlow'
   if ! (cmp -s "outWrong" "outSlow")
        echo "Error found!"
        echo "Input:"
```

SuffixArray SuffixTree Template

```
exit
    fi
done
echo Passed $4 tests
```

SuffixArray.cpp

52732a, 74 lines

```
vector<int> sort_cyclic_shifts(string const& s) {
    int n = s.size();
    const int alphabet = 256;
    vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
    for (int i = 0; i < n; i++)
       cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++)
       cnt[i] += cnt[i-1];
    for (int i = 0; i < n; i++)
       p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int classes = 1;
    for (int i = 1; i < n; i++) {
       if (s[p[i]] != s[p[i-1]])
           classes++;
       c[p[i]] = classes - 1;
    vector<int> pn(n), cn(n);
    for (int h = 0; (1 << h) < n; ++h) {
        for (int i = 0; i < n; i++) {
           pn[i] = p[i] - (1 << h);
            if (pn[i] < 0)
               pn[i] += n;
        fill(cnt.begin(), cnt.begin() + classes, 0);
        for (int i = 0; i < n; i++)
            cnt[c[pn[i]]]++;
        for (int i = 1; i < classes; i++)
            cnt[i] += cnt[i-1];
        for (int i = n-1; i >= 0; i--)
            p[--cnt[c[pn[i]]]] = pn[i];
        cn[p[0]] = 0;
        classes = 1;
        for (int i = 1; i < n; i++) {
            pair < int, int > cur = {c[p[i]], c[(p[i] + (1 << h)]}
            pair < int, int > prev = {c[p[i-1]], c[(p[i-1] + (1))]}
                 << h)) % n]};
            if (cur != prev)
                ++classes;
            cn[p[i]] = classes - 1;
        c.swap(cn);
    return p;
vector<int> suffix array construction(string s) {
    vector<int> sorted shifts = sort cyclic shifts(s);
    sorted shifts.erase(sorted shifts.begin());
    return sorted_shifts;
vector<int> lcp construction(string const& s, vector<int>
    const& p) {
    int n = s.size();
    vector<int> rank(n, 0);
```

```
for (int i = 0; i < n; i++)
    rank[p[i]] = i;
int k = 0;
vector<int> lcp(n-1, 0);
for (int i = 0; i < n; i++) {
    if (rank[i] == n - 1) {
        k = 0;
        continue;
    int j = p[rank[i] + 1];
    while (i + k < n \&\& j + k < n \&\& s[i+k] == s[j+k])
    lcp[rank[i]] = k;
    if (k)
        k--;
return lcp;
```

SuffixTree.cpp

Description: Ukkonen's algorithm for online suffix tree construction. Each node contains indices [l, r) into the string, and a list of child nodes. Suffixes are given by traversals of this tree, joining [l, r) substrings. The root is 0 (has l = -1, r = 0), non-existent children are -1. To get a complete tree, append a dummy symbol – otherwise it may contain an incomplete path (still useful for substring matching, though).

Time: $\mathcal{O}(26N)$ 07bb84, 50 lines

```
struct SuffixTree {
 enum { N = 200010, ALPHA = 26 }; // N \sim 2*maxlen+10
 int toi(char c) { return c - 'a'; }
 string a; //v = cur \ node, q = cur \ position
 int t[N][ALPHA],1[N],r[N],p[N],s[N],v=0,q=0,m=2;
 void ukkadd(int i, int c) { suff:
   if (r[v] \le q)  {
     if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
       p[m++]=v; v=s[v]; q=r[v]; goto suff; }
     v=t[v][c]; q=l[v];
   if (q==-1 || c==toi(a[q])) q++; else {
     l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
     p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
     l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
     v=s[p[m]]; q=l[m];
     while (q < r[m]) \{ v = t[v][toi(a[q])]; q + = r[v] - l[v]; \}
     if (q==r[m]) s[m]=v; else s[m]=m+2;
     q=r[v]-(q-r[m]); m+=2; qoto suff;
 SuffixTree(string a) : a(a) {
   fill(r,r+N,sz(a));
   memset(s, 0, sizeof s);
   memset(t, -1, sizeof t);
   fill(t[1],t[1]+ALPHA,0);
   s[0] = 1; 1[0] = 1[1] = -1; r[0] = r[1] = p[0] = p[1] =
   FOR(i,0,sz(a)) ukkadd(i, toi(a[i]));
 // example: find longest common substring (uses ALPHA =
 pii best;
```

```
int lcs(int node, int i1, int i2, int olen) {
   if (1[node] <= i1 && i1 < r[node]) return 1;</pre>
   if (1[node] <= i2 && i2 < r[node]) return 2;</pre>
   int mask = 0, len = node ? olen + (r[node] - 1[node]) :
   FOR(c, 0, ALPHA) if (t[node][c] != -1)
     mask |= lcs(t[node][c], i1, i2, len);
   if (mask == 3)
     best = max(best, {len, r[node] - len});
    return mask;
  static pii LCS(string s, string t) {
   SuffixTree st(s + (char) ('z' + 1) + t + (char) ('z' + 2))
   st.lcs(0, sz(s), sz(s) + 1 + sz(t), 0);
   return st.best;
};
```

```
Template.cpp
"bits/stdc++.h"
                                                   6c68dd, 46 lines
//Randomness stuff, ckmin, ckmax are optional—depends on
    time
using namespace std;
typedef long long 11;
typedef long double ld;
typedef pair<int, int> pi;
typedef pair<11,11> pl;
typedef pair<ld, ld> pd;
typedef vector<int> vi;
typedef vector<ld> vd;
typedef vector<ll> vl;
typedef vector<pi> vpi;
typedef vector<pl> vpl;
#define FOR(i, a, b) for (int i=a; i<(b); i++)
#define FOR(i, a) for (int i=0; i<(a); i++)
#define FORd(i,a,b) for (int i = (b)-1; i \ge a; i--)
#define FORd(i,a) for (int i = (a)-1; i \ge 0; i--)
#define trav(a,x) for (auto& a : x)
#define uid(a, b) uniform_int_distribution<int>(a, b) (rng)
#define sz(x) (int)(x).size()
#define mp make pair
#define pb push_back
#define f first
#define s second
#define lb lower bound
#define ub upper_bound
#define all(x) x.begin(), x.end()
#define ins insert
template < class T > bool ckmin(T& a, const T& b) { return b <
    a ? a = b, 1 : 0; 
template < class T > bool ckmax(T& a, const T& b) { return a <
    b ? a = b, 1 : 0; 
mt19937 rng(chrono::steady clock::now().time since epoch().
    count());
int main() {
```

ios_base::sync_with_stdio(0); cin.tie(0);

#!/usr/bin/env bash

```
return 0;
Treap.cpp
                                                      396fd3, 50 lines
struct node
    node *ch[2]; //ch[0] = left \ child; ch[1] = right \ child;
    int ct,priority,size,key;
    int lsize() {return(ch[0] == NULL)?0:ch[0]->size;}
    int rsize() {return(ch[1] == NULL)?0:ch[1]->size;}
};
typedef node* tree;
void update (tree & o) {//this part depends on the actual info
      to maintain
    o->size = o->ct; o->size += o->lsize(); o->size += o->
         rsize();
void rotate (tree & o, int dir) { //dir = 0: left rotate
    tree temp = o->ch[dir^1]; o->ch[dir^1] = temp->ch[dir];
         temp->ch[dir] = o;
    update(o); update(temp); o = temp;
void insert(tree & o,int kev) {
    if (o == NULL) {
        o = new node;
        o->size = o->ct = 1;o->priority = rand();o->ch[0]=o
              ->ch[1]=NULL; o->key=key;
    else if (kev == o->kev){
        o->ct++;o->size++;
    else{
        int dir = (\text{key}<o->\text{key})?0:1;
        insert (o->ch[dir], key);
        if (o->ch[dir]->priority>o->priority) rotate(o,dir^
             1);
         update(o);
void remove(tree & o.int kev) {
    if (key == o->key) {
        if (o->ct > 1) {
             o->ct--;o->size--;return;
         else if (o->ch[0] ==NULL||o->ch[1] ==NULL) {
             int d = (o->ch[0] == NULL)?0:1;
             tree temp = o; o = o->ch[d^1]; delete temp;
        else{
             int d = (o \rightarrow ch[0] \rightarrow priority > o \rightarrow ch[1] \rightarrow priority)
             rotate(o,d); remove(o,key);
    else{
        int d = (key < o > key) ?0:1;
         remove(o->ch[d],key);
    if (o) update(o);
validate.cpp
                                                      26fb12, 22 lines
```

```
for ((testNum=0;testNum<$4;testNum++))</pre>
    ./$3 > input
    ./$1 < input > out
    cat input out > data
    ./$2 < data > res
    result=$(cat res)
    if [ "${result:0:2}" != "OK" ];
    then
        echo "Error found!"
        echo "Input:"
        cat input
        echo "Output:"
        cat out
        echo "Validator Result:"
        cat res
        avit
   fi
done
echo Passed $4 tests
Vimrc.cpp
                                                        17 lines
source $VIMRUNTIME/defaults.vim
set tabstop=4
set shiftwidth=4
set nocompatible autoindent cindent ruler showcmd incsearch
    number relnumber
set cino+=L0
svntax on
filetype indent on
inoremap {<CR> {<CR>}<Esc>0
inoremap {}
                {}
imap jk
                <Esc>
set belloff=all
" Select region and then type : Hash to hash your selection.
" Useful for verifying that there aren't mistypes.
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space:]' \
\| md5sum \| cut -c-6
Voronoi.cpp
                                                  5f5301, 105 lines
// Source: http://web.mit.edu/~ecprice/acm/acm08/notebook.
    html#file7
#define MAXN 1024
#define INF 1000000
//Voronoi\ diagrams:\ O(N^2*LogN)
//Convex hull: O(N*LogN)
typedef struct {
  int id;
  double x;
  double y;
  double ang:
} chp;
double x[MAXN], v[MAXN]; // Input points
chp inv[2*MAXN]; // Points after inversion (to be given to
     Convex Hull)
int vor[MAXN]; // Set of points in convex hull;
```

```
//starts at lefmost; last same as first!!
PT ans[MAXN][2];
int chpcmp(const void *aa, const void *bb) {
 double a = ((chp *)aa) -> ang;
 double b = ((chp *)bb) -> ang;
 if (a<b) return -1;
 else if (a>b) return 1;
 else return 0; // Might be better to include a
                 // tie-breaker on distance, instead of the
                      cheap hack below
int orient(chp *a, chp *b, chp *c) {
 double s = a->x*(b->y-c->y) + b->x*(c->y-a->y) + c->x*(a->y-a->y)
      y-b->y);
 if (s>0) return 1;
 else if (s<0) return -1;
 else if (a->ang==b->ang && a->ang==c->ang) return -1; //
       Cheap hack
           //for points with same angles
 else return 0;
//the pt argument must have the points with precomputed
    angles (atan2()'s)
//with respect to a point on the inside (e.g. the center of
int convexHull(int n, chp *pt, int *ans) {
 int i, j, st, anses=0;
 gsort(pt, n, sizeof(chp), chpcmp);
  for (i=0; i<n; i++) pt[n+i] = pt[i];
  st = 0:
  for (i=1; i<n; i++) { // Pick leftmost (bottommost)</pre>
                        //point to make sure it's on the
                             convex hull
    if (pt[i].x<pt[st].x || (pt[i].x==pt[st].x && pt[i].y<pt</pre>
        [st].y) st = i;
  ans[anses++] = st;
  for (i=st+1; i<=st+n; i++) {</pre>
    for (j=anses-1; j; j--) {
     if (orient(pt+ans[j-1], pt+ans[j], pt+i)>=0) break;
      // Should change the above to strictly greater.
      // if you don't want points that lie on the side (not
           on a vertex) of the hull
      // If you really want them, you might also put an
           epsilon in orient
   ans[j+1] = i;
   anses = 1+2;
  for (i=0; i < anses; i++) ans[i] = pt[ans[i]].id;</pre>
 return anses;
int main(void) {
 int i, j, jj;
 double tmp;
  scanf("%d", &n);
  for (i=0; i<n; i++) scanf("%lf %lf", &x[i], &y[i]);
  for (i=0; i<n; i++) {
   x[n] = 2*(-INF)-x[i]; y[n] = y[i];
```

x[n+1] = x[i]; y[n+1] = 2*INF-y[i];x[n+2] = 2*INF-x[i]; y[n+2] = y[i];x[n+3] = x[i]; y[n+3] = 2*(-INF)-y[i];for (j=0; j<n+4; j++) if (j!=i) { jj = j - (j>i);inv[jj].id = j;tmp = (x[j]-x[i])*(x[j]-x[i]) + (y[j]-y[i])*(y[j]-y[i])inv[jj].x = (x[j]-x[i])/tmp;inv[jj].y = (y[j]-y[i])/tmp;inv[jj].ang = atan2(inv[jj].y, inv[jj].x); vors = convexHull(n+3, inv, vor); // Build bisectors for (j=0; j<vors; j++) {</pre> ans[j][0].x = (x[i]+x[vor[j]])/2;ans[j][0].y = (y[i]+y[vor[j]])/2;ans[j][1].x = ans[j][0].x - (y[vor[j]]-y[i]);ans[j][1].y = ans[j][0].y + (x[vor[j]]-x[i]);printf("Around (%lf, %lf) \n ", x[i], y[i]); // List all intersections of the bisectors for (j=1; j<vors; j++) {</pre> PT vv; vv = ComputeLineIntersection(ans[j-1][0], ans[j-1][1], ans[j][0], ans[j][1]); printf("%lf, %lf\n", vv.x, vv.y); printf("\n"); return 0;

Z-algorithm.cpp

85d656, 13 lines

```
vector<int> z_function(string s) {
   int n = (int) s.length();
   vector<int> z(n);
   for (int i = 1, 1 = 0, r = 0; i < n; ++i) {
      if (i <= r)
            z[i] = min (r - i + 1, z[i - 1]);
      while (i + z[i] < n && s[z[i]] == s[i + z[i]])
            ++z[i];
   if (i + z[i] - 1 > r)
        l = i, r = i + z[i] - 1;
   }
   return z;
}
```

hash.sh

3 line

Math Hints (2)

Newton's Method: $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$.

Lagrange Multiplier: Let $f: \mathbb{R}^n \to \mathbb{R}$ be the objective function, $g: \mathbb{R}^n \to \mathbb{R}^c$ be the constraints function, let x^* be an optimal solution to the optimization problem such that $Dg(x^*) = c$: maximize f(x) subject to g(x) = 0. There there exists λ such that $Df(x^*) = \lambda^* \top Dg(x^*)$.

Burnside's Lemma: Let G be a finite group acting on set X. Let X^g denote the set of elements in X that are fixed by g. Then the number of orbits is given by $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$.

Linear Time Inverses Modulo p: For $i \ge 2$, $i^{-1} = -\lfloor \frac{p}{i} \rfloor (p \mod i)^{-1}$.

Quadratic Residue: $\left(\frac{a}{p}\right) = a^{(p-1)/2}$. If $x^2 \equiv a(\text{mod}p)$ has a solution, then $\left(\frac{a}{p}\right) = 1$.

 $LGV\ Lemma$: Assume G=(V,E) is a DAG. Let $\omega(P)$ be the product of edge weights on path P. Let $e(u,v):=\sum_{P:u\to v}\omega(P)$ be the sum of $\omega(P)$ for all paths from u to v. The set of sources $A\subseteq V$ and set of sinks $B\subseteq V$. A collection of disjoint paths $A\to B$ consists of n paths S_i such that S_i is a path from A_i to $B_{\sigma(S)_i}$ such that for any $i\neq j$, S_i and S_i does not share a common vertex. Then if we let

$$M = \begin{bmatrix} e(A_1, B_1) & e(A_1, B_2) & \cdots & e(A_1, B_n) \\ \vdots & \vdots & \ddots & \vdots \\ e(A_n, B_1) & e(A_n, B_2) & \cdots & e(A_n, B_n) \end{bmatrix}$$

then det $M = \sum_{S:A\to B} \operatorname{sgn} \sigma(S) \prod_{i=1}^n \omega(S_i)$ where $S:A\to B$ denotes a set of disjoint paths S from A to B.

Network Flow with Lower/Upper Bounds: Suppose the flow must satisfy $b(u,v) \leq f(u,v) \leq c(u,v)$, and have conservation of flows over vertices.

Variant 1: No source/sink (i.e. flow at all vertices must be balanced), check if there is a feasible flow: create a graph G'. For any edge $u \to v$, add an edge with capacity c(u,v) - b(u,v). Now assume initially, at vertex u, the sum of capacities of edges into u minus the sum of capacities out of u is M. If M > 0, add an edge from super source S to u with capacity M. If M < 0, add an edge from u to the super sink T with capacity -M.

Variant 2: Feasible flow with sources and sinks: Add an edge from original sink t to source s with capacity $+\infty$.

Variant 3: Maximum flow with sources and sinks: first, check if there exists a feasible flow. Then augment using source s and sink t (i.e. run Dinic again with source s and sink t).

Variant 4: Minimum flow with sources and sinks: find the feasible flow first, remove the edge from sink to source (the current flow on the edge is the size of the original flow), and then run the maximum flow from the sink t to the source s to see how much we may get rid off from the original flow.

Pick's Theorem: Suppose a polygon has integer coordinates for all of its vertices; let i be the number of integer points that are interior to the polyon, b be the number of integer points on its boundary, the area of the polygon is $A = i + \frac{b}{2} - 1$.

Mobius Transformation/Circle Inversion: Mobius transformations $f: \hat{\mathbb{C}} \to \hat{\mathbb{C}}$ are specified by $f(z) = \frac{az+b}{cz+d}$.

Dilworth's Theorem: For a partially ordered set S, the maximum size of an antichain is equal to the minimum number of chains (i.e. any two elements are comparable) required to cover S.

LP Duality: Suppose the primal linear program is given by maximize $\mathbf{c}^{\top}\mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0$, the dual program is given by minimize $\mathbf{b}^{\top}\mathbf{y}$ subject to $A^{\top}\mathbf{y} \geq \mathbf{c}, \mathbf{y} \geq 0$.

Mobius Inversion: If g, f are arithmetic functions satisfying $g(n) = \sum_{d|n} f(d)$ for $n \ge 1$, then $f(n) = \sum_{d|n} \mu(d)g(\frac{n}{d})$.

Number of Points on Lattice Convex Polygon: A convex lattice polygon with coordinates in [0, N] has at most $O(N^{2/3})$ points.

Green's Theorem: Let C be a positively oriented, smooth, simple closed curve and let D be the region bounded by C. If L and M are functions of (x,y) defined on an open region containing D and having continuous partial derivatives, then

$$\int_{C} (Ldx + Mdy) = \int_{D} (\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y}) dx dy.$$

Polynomial Division: Suppose we are given polynomials f(x), g(x), and we want to write f(x) as f(x) = Q(x)g(x) + R(x). Let $f^R(x) = x^{\deg f}f(\frac{1}{x})$ (i.e. reverse the coefficients of the polynomial). Let $n = \deg f, m = \deg g$. Then

$$f^R(x) \equiv Q^R(x)g^R(x) \mod x^{n-m+1}$$