# **ISO** Presentation

Quantum Programming

Tristan Nemoz Supervisor: Mario Berta

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London

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- Not all ratings are known
- Can we predict whether user j will like product k?

## Mathematical description

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Let  $\mathbf{P} \in \{0; 1\}^{m \times n}$ , whose coefficients are known with a given probability. A quantum recommendation system, given  $\hat{\mathbf{P}}$  the incomplete respresentation of  $\mathbf{P}$  and an user j, predicts k such that  $\mathbf{P}_{j,k} = 1$  with high probability.

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Impossible problem without assumptions

#### The low-rank assumption

Let **P** be the unknown preference matrix. For a given precision parameter  $\varepsilon$ , there exists a rank-k matrix **T** such that:

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#### **Problem initialization**

**P** can be sampled into  $\hat{\mathbf{P}}$  such that  $\hat{\mathbf{P}}$  holds the low-rank assumption. Hence,  $\varepsilon$  is assumed to be known.

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Imperial College London

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#### **Problem statement**

How does the real-world implementation of the Quantum Recommendation System algorithm differs from its theoretic implementation?

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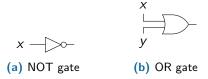
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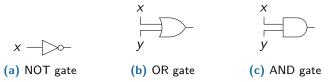


Figure: Classical gates

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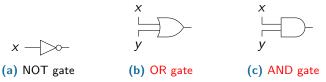


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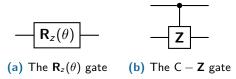
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(a) The  $\mathbf{R}_z(\theta)$  gate

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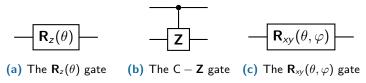


Figure: Quantum gates

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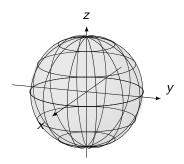
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A qubits-string is called a **quantum register**. It is mathematically described as the tensor product of these qubits, that is:

$$|q_1\rangle |q_2\rangle \cdots |q_n\rangle = \bigotimes_{i=1}^n |q_k\rangle$$
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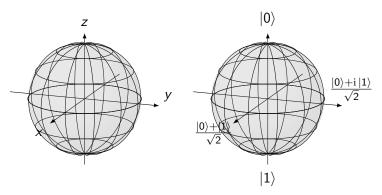


Figure: The Bloch Sphere

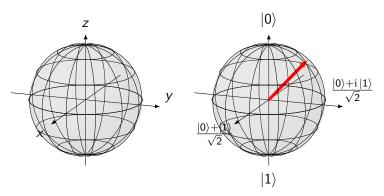


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# Loading a vector as a quantum state

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Loading **x** means creating  $|x\rangle$  from  $|0\rangle^{\otimes \lceil \log_2(n) \rceil}$  with a polylogarithmic number of gates<sup>4</sup>.

<sup>&</sup>lt;sup>4</sup>Prakash, "Quantum Algorithms for Linear Algebra and Machine Learning.

### **QRAM**<sup>5</sup>

A QRAM is a binary tree whose leaves stores the coefficients of a vector  $\mathbf{x}$  and whose nodes stores the sum of its leaves values.

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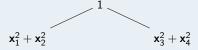


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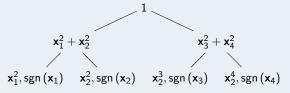


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# **Loading from QRAM**

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- Uses 2<sup>k</sup> rotations gates around the y-axis at level k of the tree whose angle are determined using QRAM
- Goal: parallelize the execution of the rotations using superposition

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### **QRAM** function assumption

It is possible, using QRAM, to design a gate  $L_k$  such that:

$$\mathbf{L}_{k} |k\rangle |0\rangle^{\otimes t} = |k\rangle |\overline{\theta_{k}}\rangle$$

where  $\overline{\theta_k}$  is the best *t*-bits approximation of  $\theta_k$ .

# A simpler problem

#### Parallel execution of rotations

It is possible to get the state  $|\theta\rangle$   $e^{i\theta}$   $|x\rangle$  from the state  $|\theta\rangle$  ( $\alpha$   $|0\rangle + \beta$   $|1\rangle$ )

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#### The z-rotation

Given an angle  $\theta$ , the rotation around the z-axis of the Bloch sphere is given by the gate:

$$\mathbf{R}_z( heta) = egin{pmatrix} 1 & 0 \ 0 & \mathrm{e}^{\mathrm{i}\, heta} \end{pmatrix} \,.$$

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Controlling the z-rotation gates on the target qubit in the previous circuit applies a rotation on the target qubit around the z-axis of angle  $\theta_k$  only if the first quantum register is in state  $|k\rangle$  in time O(t).

Is it possible to transform a *y*-rotation into a *z*-rotation of the same angle?

## **Converting** a *z*-rotation to a *y*-rotation

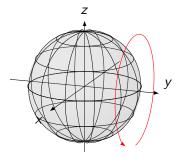


Figure: An exemple of a transformation of a y-rotation into a z-rotation

# **Converting** a *z*-rotation to a *y*-rotation

# Rotate the sphere

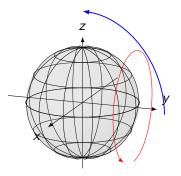
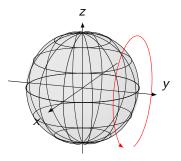


Figure: An exemple of a transformation of a y-rotation into a z-rotation

# **Converting** a *z*-rotation to a *y*-rotation

### Rotate the sphere



**Figure:** An exemple of a transformation of a *y*-rotation into a *z*-rotation

# Errata

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# Conclusion

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### References I

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