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MODULE *WoundWait*

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EXTENDS *TLC, FiniteSets, Integers, Sequences*

CONSTANT *txs, locks, actives*  
VARIABLE *tStat, lStat, lOwn, waitList*

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$Vars \triangleq \langle tStat, lStat, lOwn, waitList \rangle$

$TypeOK \triangleq$   
 $\wedge \forall t \in txs : tStat[t] \in \{ \text{"F"}, \text{"I"} \}$   
 $\wedge \forall l1 \in locks : lStat[l1] \in \{ \text{"F"}, \text{"I"} \}$   
 $\wedge \forall l2 \in locks : lOwn[l2] \in txs \cup \{0\}$   
 $\wedge waitList \subseteq txs$

$Init \triangleq$   
 $\wedge tStat = [t \in txs \mapsto \text{"F"}]$   
 $\wedge lStat = [l \in locks \mapsto \text{"F"}]$   
 $\wedge lOwn = [l \in locks \mapsto 0]$   
 $\wedge waitList = \{ \}$

$Min(L) \triangleq \text{CHOOSE } t \in L : \forall ts \in L : t \leq ts$

$SignalWait \triangleq$   
IF  $waitList \neq \{ \}$   
THEN  $waitList' = waitList \setminus \{ Min(waitList) \}$   
ELSE UNCHANGED  $waitList$

$Wait(t) \triangleq \wedge waitList' = waitList \cup \{ t \}$

$CanAcquire(t, l) \triangleq$   
 $\vee \wedge lStat[l] = \text{"F"}$   
 $\vee \wedge lStat[l] = \text{"I"}$   
 $\wedge lOwn[l] = t$

$AcquireLock(t, l) \triangleq$   
 $\vee \wedge CanAcquire(t, l)$   
 $\wedge lStat' = [lStat \text{ EXCEPT } ![l] = \text{"I"}]$   
 $\wedge lOwn' = [lOwn \text{ EXCEPT } ![l] = t]$   
 $\wedge tStat' = [tStat \text{ EXCEPT } ![t] = \text{"I"}]$   
 $\wedge \text{UNCHANGED } waitList$   
 $\vee \wedge lStat[l] = \text{"I"}$   
 $\wedge lOwn[l] \neq t$   
 $\wedge \text{IF } (t < lOwn[l])$   
THEN  $\wedge lStat' = [lStat \text{ EXCEPT } ![l] = \text{"I"}]$   
 $\wedge lOwn' = [lOwn \text{ EXCEPT } ![l] = t]$   
 $\wedge tStat' = [tStat \text{ EXCEPT } ![lOwn[l]] = \text{"F"}]$

$$\begin{aligned}
& \wedge tStat' = [tStat \text{ EXCEPT } ![t] = \text{"I"}] \\
& \wedge \text{UNCHANGED } waitList \\
\text{ELSE } & \wedge Wait(t) \\
& \wedge \text{UNCHANGED } \langle tStat, lStat, lOwn \rangle
\end{aligned}$$

$$\begin{aligned}
Commit(t, l) & \triangleq \\
& \wedge lOwn[l] = t \\
& \wedge lOwn' = [lOwn \text{ EXCEPT } ![l] = 0] \\
& \wedge lStat' = [lStat \text{ EXCEPT } ![l] = \text{"F"}] \\
& \wedge tStat' = [tStat \text{ EXCEPT } ![t] = \text{"F"}] \\
& \wedge t' = t + \text{actives} \\
& \wedge SignalWait
\end{aligned}$$

$$\begin{aligned}
Next & \triangleq \\
& \exists t \in txs : \exists l \in locks : \\
& \quad \vee AcquireLock(t, l) \\
& \quad \vee Commit(t, l)
\end{aligned}$$

$$\begin{aligned}
Spec & \triangleq \\
& \wedge Init \\
& \wedge \Box [Next]_{Vars}
\end{aligned}$$

$$\begin{aligned}
FairSpec & \triangleq \\
& Spec \wedge \forall t1 \in txs, l1 \in locks : \\
& \quad \text{WF}_{Vars}(AcquireLock(t1, l1)) \\
& \wedge \forall t2 \in txs, l2 \in locks : \\
& \quad \text{WF}_{Vars}(Commit(t2, l2))
\end{aligned}$$

$$DeadLock \triangleq waitList \neq txs$$

$$\begin{aligned}
Starvation & \triangleq \\
& \wedge \forall t1 \in txs, l1 \in locks : \\
& \quad \Box \Diamond (\langle AcquireLock(t1, l1) \rangle_{Vars}) \\
& \wedge \forall t2 \in txs, l2 \in locks : \\
& \quad \Box \Diamond (\langle Commit(t2, l2) \rangle_{Vars})
\end{aligned}$$