CPSC 335 Project 4: Hashing

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Introduction

In this project you will design, implement, and analyze one algorithm for the hashing problem. The algorithm is called Cuckoo Hashing, presented in the class.

For this problem, you will design and implement one algorithm in C/C++/Java/Python, test it on various inputs and complete a hash table with a given input. No algorithm analysis is needed for this project.

The Cuckoo Hashing Algorithm

There are several versions of cuckoo hashing. The version we learned in class is the simplest, where there are two hash functions, and thus only two places where any given item could be stored in the table. Let us consider the set of keys to be printable ASCII strings of length no more than 80. Let us consider the hash table size to be 17.

If key is the string representing the key, then let keysize be the size of the string and key[i] be the ASCII code of the $(i+1)^{th}$ character in the string key read from left to right: $key = key_0 key_1...key_{kevsize-1}$

Java uses the following hash function on strings:

$$val = 31^{keysize-1} \cdot key_0 + 31^{keysize-2} \cdot key_1 + ... + 31^1 \cdot key_{keysize-2} + key_{keysize-1}$$

Let us consider two hash functions, f_1 and f_2 . Function f_2 will compute the hash value using Java's hash function formula, while the function f_1 computes a different hash value using a different hash function. Function f_1 computes first a large number then it brings the result into the proper range using the formulas below:

$$val = \sum_{i=0}^{keysize-1} key[i] \cdot 31^{i}$$

$$f_1 = val \% tablesize$$
if $f_1 < 0$ then $f_1 = f_1 +$ tablesize

Function f₂ computes first a large number then it brings the result into the proper range using the formulas below:

$$val = \sum_{i=0}^{keysize-1} key[keysize - i - 1] \cdot 31^{i}$$

$$f_2 = val \% tablesize$$
if $f_2 < 0$ then $f_2 = f_2 +$ tablesize

Both functions f_1 and f_2 compute first a large number then it brings the result into the proper range 0..tablesize-1. But we bring the intermediate results into the proper range after each calculation, we do not need to wait until we compute the final result. Also, we can ring the power term 31^{index} into the proper range before multiplying it with key_{index}

You need to insert the strings below (also given in the input file in6.txt) into the hash table provided next. Please put an empty line at the end of the file.

Algorithm Engineering California State University Fullerton College of Engineering and Computer Science **Department of Computer Science Dynamic Programming Monge Properties String Matching Matrix Searching Optimal Tree Construction Online algorithms** emphasis on **Server Problem** Some related problem **Self-Stabilization** One of the greatest mysteries in science **Quantum Nature of Universe** In physics and are known Cuckoo hashing is fun

into the hash table (next page) using f_1 for the first table and f_2 for the second table. Show the result of the insertion in the table shown on next page.

<u>Hint</u>: consider a two-dimensional table of strings t, where t[0] is T1 and t[1] is T2. Consider a variable index that oscillates between 0 and 1 as it would have oscillated between T1 and T2. In C++, the value of index could be changed using the tertiary operator: index = index? 0:1. Depending on the value of index, either apply hash function f_1 (index == 0) or f_2 (index == 1).

	Table T1	Table T2
[0]		
[1]		

[2]	
[3]	
[4]	
[5]	
[6]	
[7]	
[8]	
[9]	
[10]	
[11]	
[12]	
[13]	
[14]	
[15]	
[16]	

What to do

- 1. Write clear pseudocode for the algorithm.
- 2. Type these notes (electronically or on paper) and submit it as a PDF report.
- 3. Implement your algorithm in C/C++/Java/Python. You may use the templates provided at the end of this file.
- 4. Compile and execute the program.
- 5. Complete the table using the strings from the file **in6.txt** as the input and insert it into the PDF report.
- 6. Create a file with the output of the program for an input value and submit it together with the program. Note, the output can be redirected to a file (for easy printing). For example, the following command line will create an output file in Linux-based operating system called alout.txt by re-directing the output from the screen (display) to the file alout.txt:

 $K:\cpscs335> a.out > a4out.txt$

Sample outputs for the Cuckoo Hashing Algorithm:

```
Example #1:
K:\202> ast4
CPSC 335-x – Programming Assignment #4: Cuckoo Hashing algorithm
Input the file name (no spaces)!
in4 31.txt
CPSC 335-x - Programming Assignment #4: Cuckoo Hashing algorithm
Input the file name (no spaces)!
in4 31.txt
String < Algorithm Engineering> will be placed at t[11][0]
String < California > will be placed at t[16][0]
String <State University> will be placed at t[5][0]
String <Fullerton> will be placed at t[15][0]
String < College of Engineering and Computer Science > will be placed at t[10][0]
String < Department of Computer Science > will be placed at t[5][0] replacing < State University >
String <State University> will be placed at t[7][1]
String < Dynamic Programming > will be placed at t[3][0]
String <Monge Properties> will be placed at t[9][0]
String <String Matching> will be placed at t[16][0] replacing <California>
String < California > will be placed at t[2][1]
String <Matrix Searching> will be placed at t[5][0] replacing <Department of Computer Science>
String < Department of Computer Science > will be placed at t[12][1]
String < Optimal Tree Construction > will be placed at t[5][0] replacing < Matrix Searching >
String <Matrix Searching> will be placed at t[11][1]
String <Online algorithms> will be placed at t[0][0]
String <emphasis on> will be placed at t[15][0] replacing <Fullerton>
String <Fullerton> will be placed at t[3][1]
String <Server Problem> will be placed at t[9][0] replacing <Monge Properties>
String < Monge Properties > will be placed at t[2][1] replacing < California >
```

```
String < California > will be placed at t[16][0] replacing < String Matching >
String <String Matching> will be placed at t[16][1]
Example #2:
K:\202> ast4
CPSC 335-x – Programming Assignment #4: Cuckoo Hashing algorithm
Input the file name (no spaces)!
in5 31.txt
String < Algorithm Engineering> will be placed at t[11][0]
String < California > will be placed at t[16][0]
String <State University> will be placed at t[5][0]
String <Fullerton> will be placed at t[15][0]
String < College of Engineering and Computer Science > will be placed at t[10][0]
String < Department of Computer Science > will be placed at t[5][0] replacing < State University >
String <State University> will be placed at t[7][1]
String < Dynamic Programming > will be placed at t[3][0]
String < Monge Properties > will be placed at t[9][0]
String <String Matching> will be placed at t[16][0] replacing <California>
String < California > will be placed at t[2][1]
String < Matrix Searching > will be placed at t[5][0] replacing < Department of Computer Science >
String < Department of Computer Science > will be placed at t[12][1]
String < Optimal Tree Construction > will be placed at t[5][0] replacing < Matrix Searching >
String <Matrix Searching> will be placed at t[11][1]
String <Online algorithms> will be placed at t[0][0]
String <emphasis on> will be placed at t[15][0] replacing <Fullerton>
String <Fullerton> will be placed at t[3][1]
String <Server Problem> will be placed at t[9][0] replacing <Monge Properties>
String < Monge Properties> will be placed at t[2][1] replacing < California>
String <California> will be placed at t[16][0] replacing <String Matching>
String <String Matching> will be placed at t[16][1]
String <Some related problem> will be placed at t[11][0] replacing <Algorithm Engineering>
String < Algorithm Engineering > will be placed at t[2][1] replacing < Monge Properties >
String < Monge Properties > will be placed at t[9][0] replacing < Server Problem >
String <Server Problem> will be placed at t[4][1]
String <Self-Stabilization> will be placed at t[2][0]
String <One of the greatest> will be placed at t[6][0]
```

Template for a C/C++ program doing Cuckoo Hashing algorithm:

```
// Assignment 4: Cuckoo Hashing algorithm
// XX YY ( YOU NEED TO COMPLETE YOUR NAME )
// An open addressing method called Cuckoo Hashing
// INPUT: an input file containing strings of characters, one string per line
// OUTPUT: a detailed list of where the strings are inserted.
#include <iostream>
#include <cstring>
#include <stdio.h>
using namespace std;
// cuckoo tables' size
const int tablesize = 17;
// combine the two 1-dimensional table into one 2-dimensional table
char t[tablesize][2][255];
// compute the hash functions
size t f(char*, size t);
// place a string in one of the hash tables
bool place in hash tables (char*);
int main() {
// the strings to be stored in the hash tables
 char s[255]="";
 char null st[] ="";
 size t i, len;
 bool placed;
// clear the tables
 for(i=0; i < tablesize; i++) {
  strepy(t[i][0], null st);
  strcpy(t[i][1], null_st);
 char filename[255] = "";
 // display the header
 cout << endl << "CPSC 335-x - Programming Assignment #4: ";
 cout << "Cuckoo Hashing algorithm" << endl;</pre>
// read the strings from a file
 cout << "Input the file name (no spaces)!" << endl;</pre>
 cin >> filename;
// open the file for reading
 FILE *file = fopen ( filename, "r" );
 if (file!=NULL)
  {
```

```
/* read line by line from the file */
   while (fgets (s, 255, file)!= NULL) {
         // place null character at the end of the line instead of <return>
         len = strlen(s);
         s[len-2]='\0';
         // insert the string in the cuckoo table
         placed = place in hash tables(s);
         // check whether the placement was successful
         if (!placed) {
          cout << "Placement has failed" << endl;</pre>
          return -1;
   fclose (file);
 else
   perror (filename); /* why didn't the file open? */
return 0;
bool place in hash tables (char *s) {
bool placed;
 size t pos;
 int index;
 char temp s[255], temp[255];
strcpy(temp_s, s);
// use a counter to detect loops
 int counter = 0;
// start with table T1
 index = 0;
placed = false;
pos = f(temp s, index);
 while((!placed ) && (counter < 2*tablesize)) {
  if (strcmp(t[pos][index], "") == 0)
   // the entry at index <pos> in the <index> hash table is available so store the string <temp s> there
   cout << "String <" << temp s << "> will be placed at";
   cout << " t[" << pos <<"][" << index << "]" << endl;
   strcpy(t[pos][index], temp s);
   placed = true;
   return placed;
  else {
   // the entry at index <pos> in the <index> hash table is not available so
```

```
// obtain the string stored over there in variable <temp> and store the string <temp s> there
   // now the string <temp> needs to be placed in the other table
   cout << "String <" << temp s << ">> will be placed at" << " t[" << pos;
   cout <<"][" << index << "]" << " replacing <" << t[pos][index] << ">";
   cout << endl;
   // YOU NEED TO WRITE THE CODE TO STORE IN temp THE STRING STORED AT
   // t[pos][index] AND STORE IN t[pos][index] THE STRING temp s
   strcpy(temp s, temp);
   // NOW temp s CONTAINING THE EVICTED STRING NEEDS TO BE STORED
   // IN THE OTHER TABLE
   // WRITE THE CODE TO SET index TO INDICATE THE OTHER TABLE
   // WRITE THE CODE TO CALCULATE IN pos THE HASH VALUE FOR temp s
   counter ++;
return placed;
};
size t f(char *s, size t index) {
// compute the hash functions
// s is the string (the key) to which we apply the hash function
// index indicates which hash function will be used
// index == 0 means the first hash function
// index == 1 means the second hash function
size t po, len;
int i, val, temp;
po = 1;
len = strlen(s);
 if (index == 0) {
  val = s[0];
  val = val % tablesize;
  if (val < 0) val += tablesize;
  if (len == 1)
   return val;
  for (i = 1; i < len; i++)
   temp = s[i];
   po *= 31;
   po = po \% tablesize;
   if (po < 0) po += tablesize;
   val += temp * po;
   val = val % tablesize;
   if (val < 0) val += tablesize;
  return val;
```

```
else {
// YOU NEED TO IMPLEMENT THE STEPS TO CALCULATE THE SECOND
// HASH FUNCTION
val = val % tablesize;
return val;
}
}
```