

# **CSC4200/5200 – COMPUTER NETWORKING**

**Instructor: Susmit Shannigrahi**

**SPANNING TREE**

**sshannigrahi@tnitech.edu**

**GTA: dereddick42@students.tnitech.edu**



# Exam

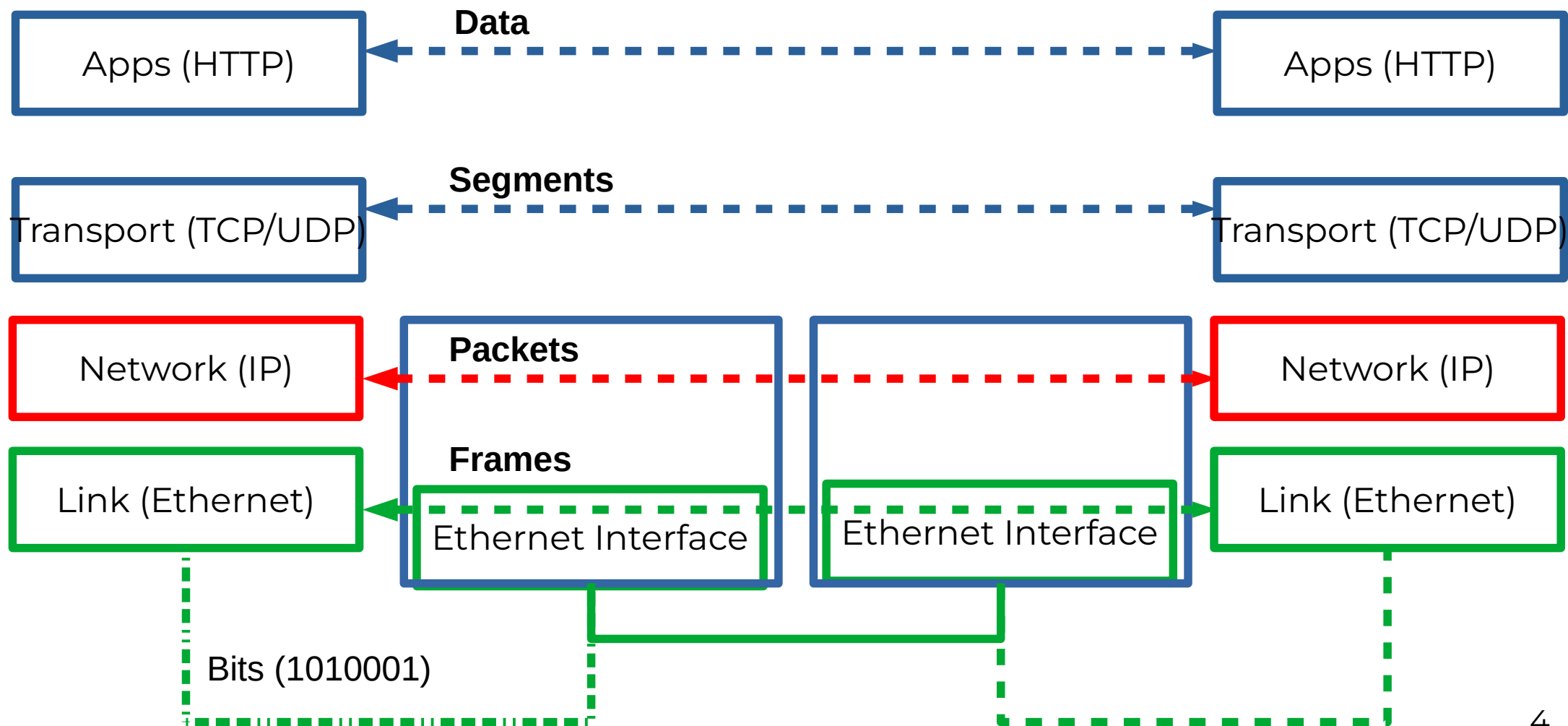
---

- Sept 30th
  - 55 minutes – 11:15AM-12:10PM
  - **If you have a conflict, let me know NOW!**
  - Location – iLearn
- Open book – but you may not have time to look things up.
- Only from the book and lecture notes, no programming questions

# Project groups

---

- Let me know by 09/25, Friday.
  - Groups of 3.
- If I don't hear from you by Friday, you will be assigned to a random group!



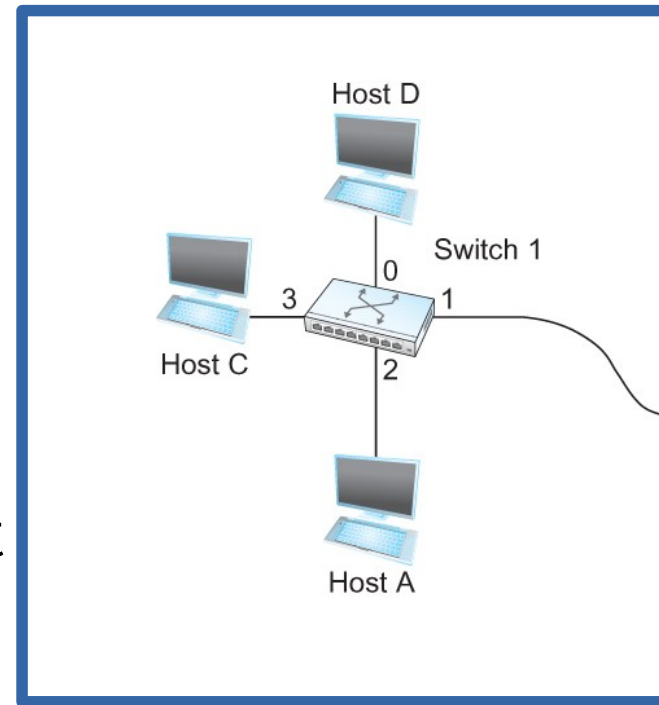
# So far...

---

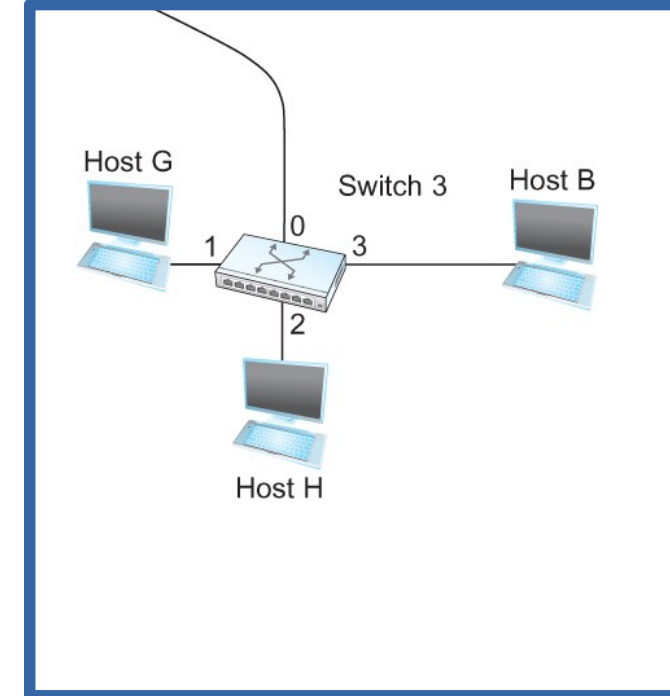
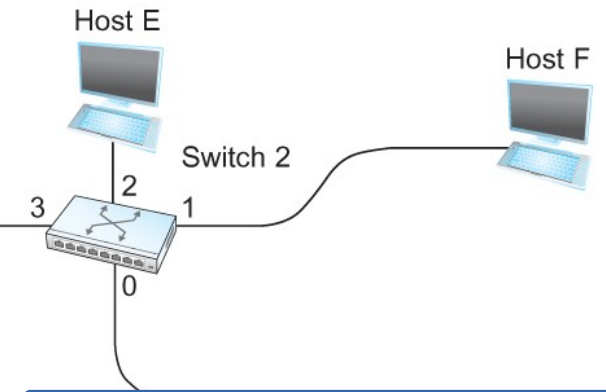
- we saw how to build a local network
- How do we interconnect different types of networks to build a large global network?

# Switching

- Switch
  - A mechanism to interconnect links to form a large network
- Forward **frames**
- Separate the collision domains
- Filter packets between LANs
- Connects two or more LAN segments - **Bridging**



**LAN 1**  
**Collision domain 1**



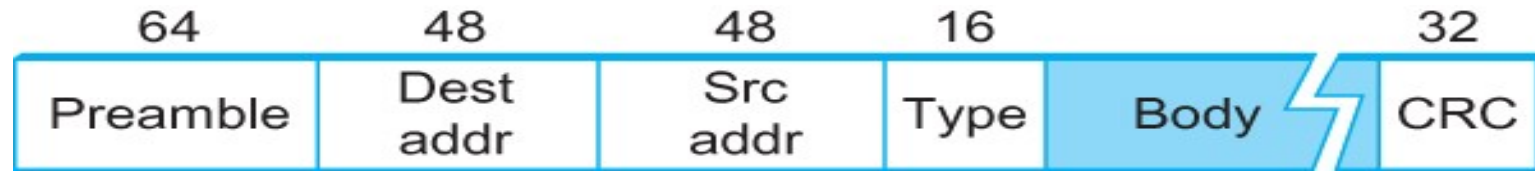
**LAN 2**  
**Collision domain 2**

# Switches are self learning!

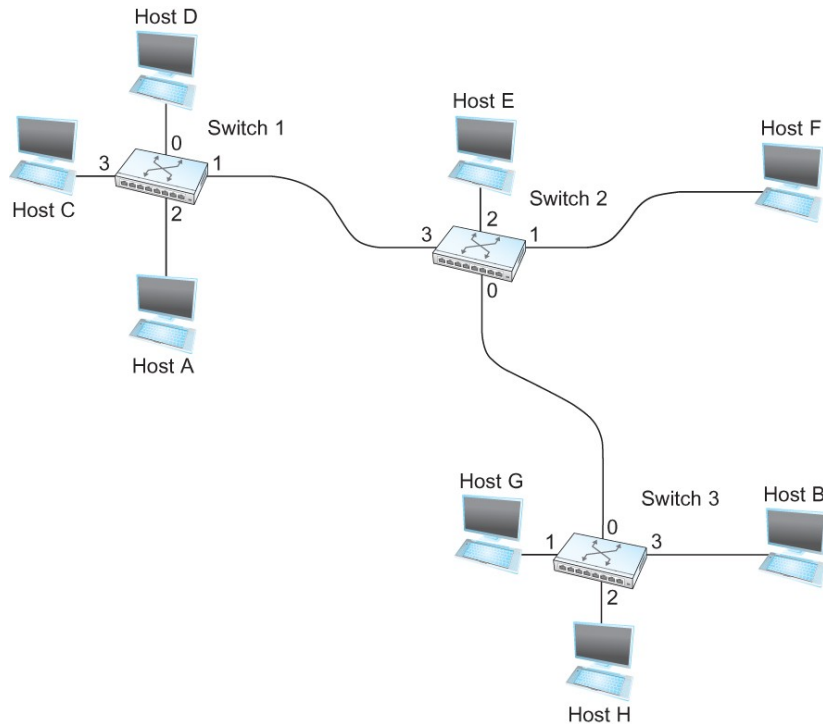
---

- Inspect the source MAC address
  - **What is a mac address?**
- Associate mac address and incoming interface
- Store this association for later use, (for some time)
  - aging-timer

# Switching Table



- To decide how to forward a packet, a switch consults a *forwarding table*



Destination, Port

---

|   |   |
|---|---|
| A | 3 |
| B | 0 |
| C | 3 |
| D | 3 |
| E | 2 |
| F | 1 |
| G | 0 |
| H | 0 |

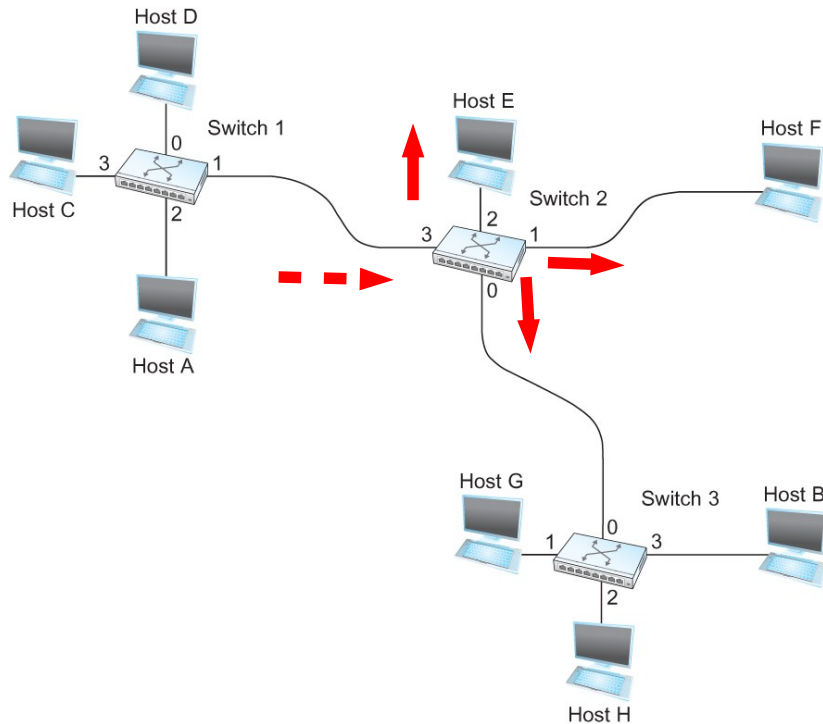
**Forwarding Table for  
Switch 2**



# Switching Table

- Unknown destination → send out on all Interfaces **(flooding)**

- **Skip the incoming interface**



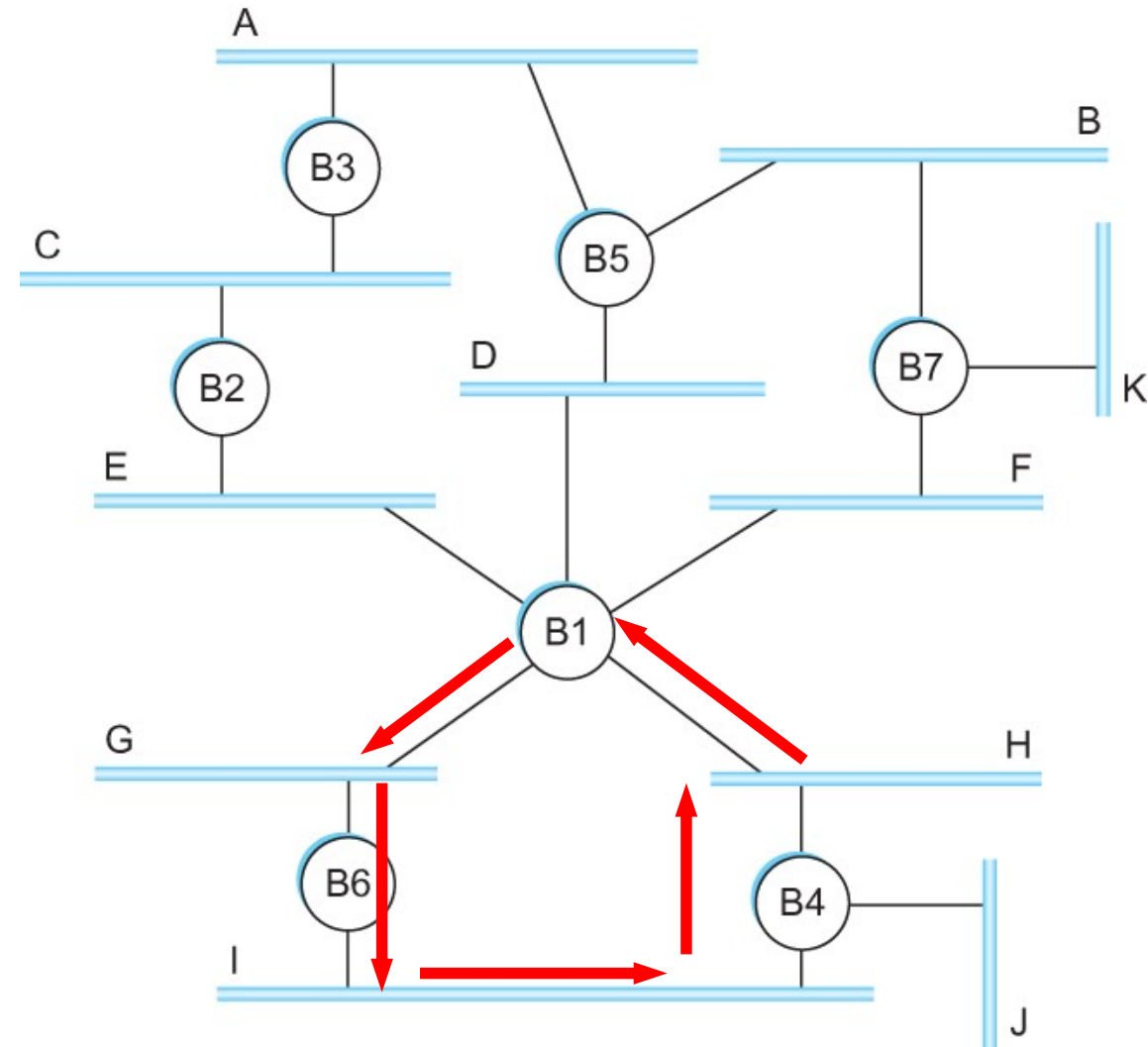
Destination, Port

---

|   |   |
|---|---|
| A | 3 |
| B | 0 |
| C | 3 |
| D | 3 |
| E | 2 |
| F | 1 |
| G | 0 |
| H | 0 |

**Forwarding Table for  
Switch 2**

# Loop

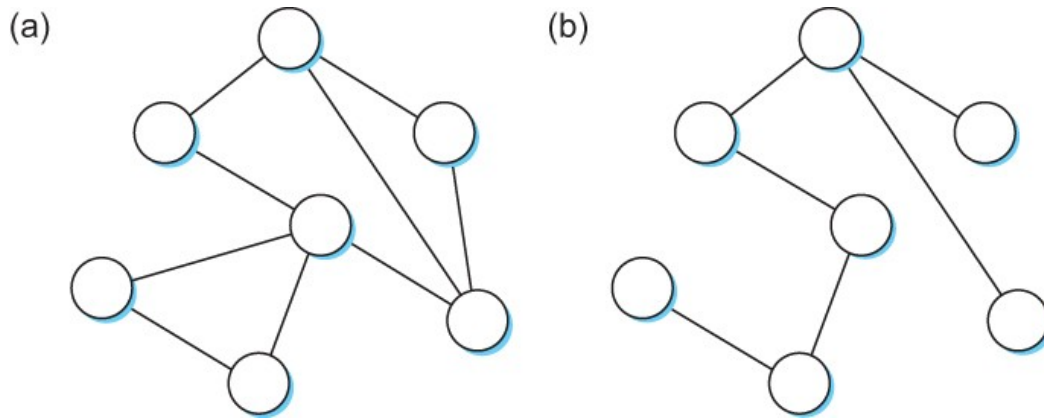


**Spot the loop**  
**Why?**

# Solution? Spanning Tree

Think of the extended LAN as being represented by a graph that possibly has loops (cycles)

- A spanning tree is a sub-graph of this graph that covers all the vertices but contains no cycles
- Spanning tree keeps all the vertices of the original graph but throws out some of the edges



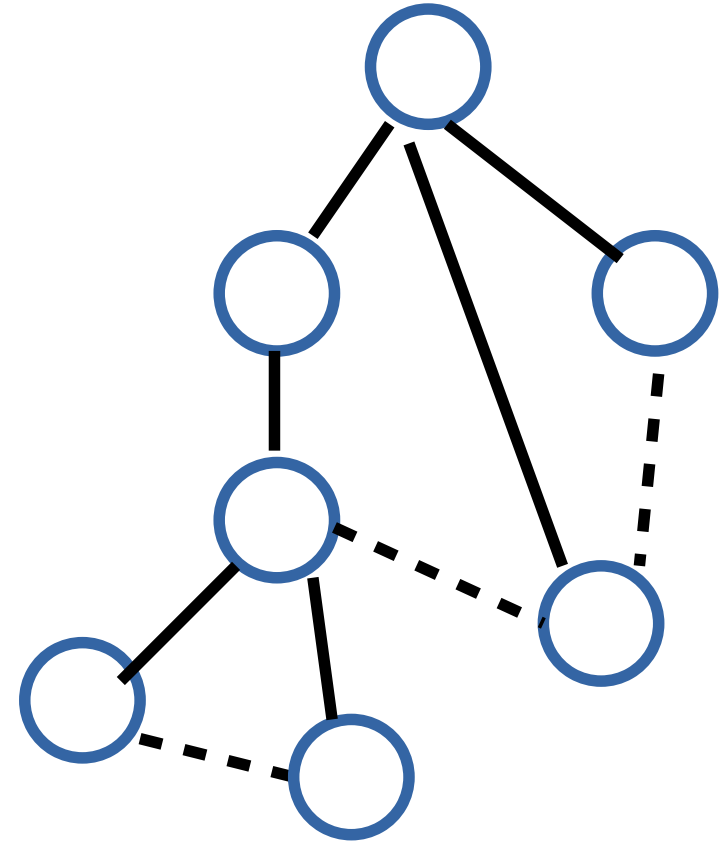
Example of (a) a cyclic graph; (b) a corresponding spanning tree.

# How do we create a spanning tree?

- Properties: No loops
- How?
  - Selectively flood
  - Distributed algorithm, no coordination!
  - Automatic reconciliation when failure occurs

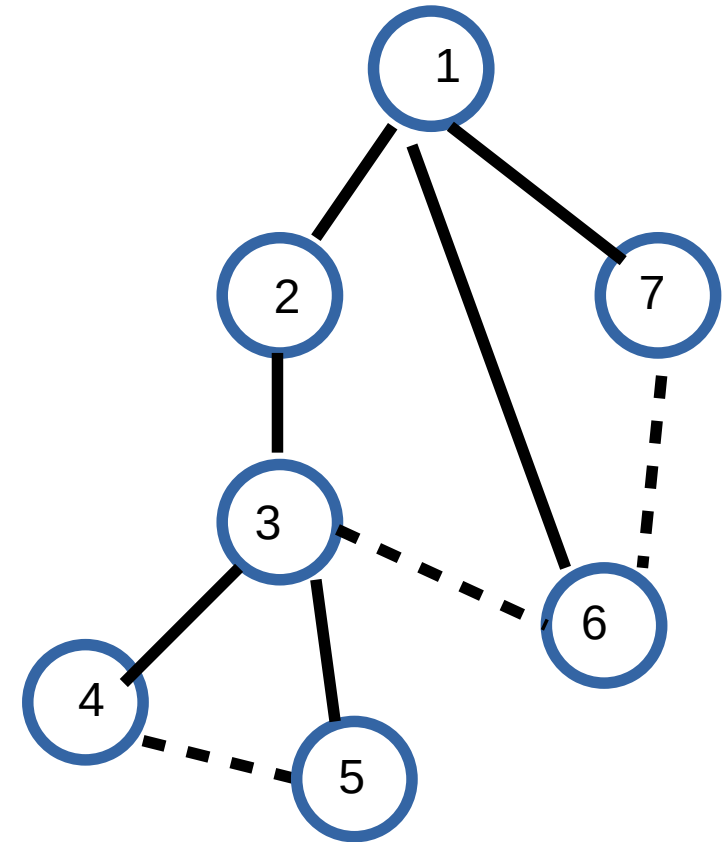
# How do we create a spanning tree?

- Properties: No loops
- How?
  - Selectively flood
  - Distributed algorithm, no coordination!
  - Automatic reconciliation when failure occurs
- Switches elect a root
  - The switch with the smallest identifier
  - Each switch identifies if its interface is on the shortest path from the root
  - Exclude if not
- Send message  $(Y, d, X)$
- From  $x$ , claims  $Y$  is the root, distance is  $d$



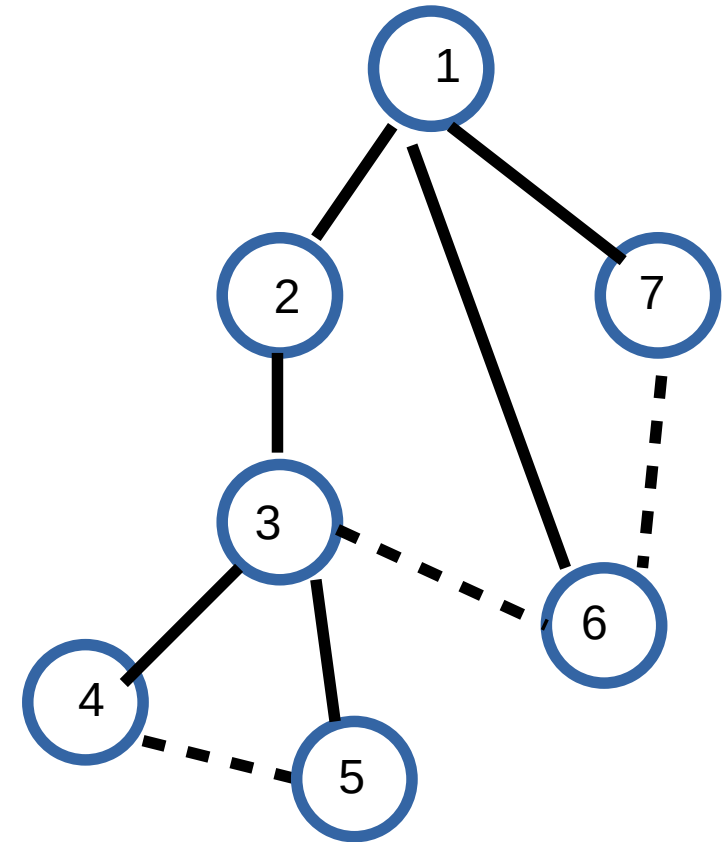
# How do we create a spanning tree?

- **Message (Y, d, X) - (to, distance, from)**
- 4 thinks it's the root
- Sends (4, 0, 4) to 3 and 5
- Receives (3,0,3) from 3
  - Sets it to as the root since  $3 < 4$
- Receives (3,1,5) from 5
  - Sees that this is a longer path to 3
  - 2 hops vs direct path (1 hop)
  - Removes 4-5 link from the tree



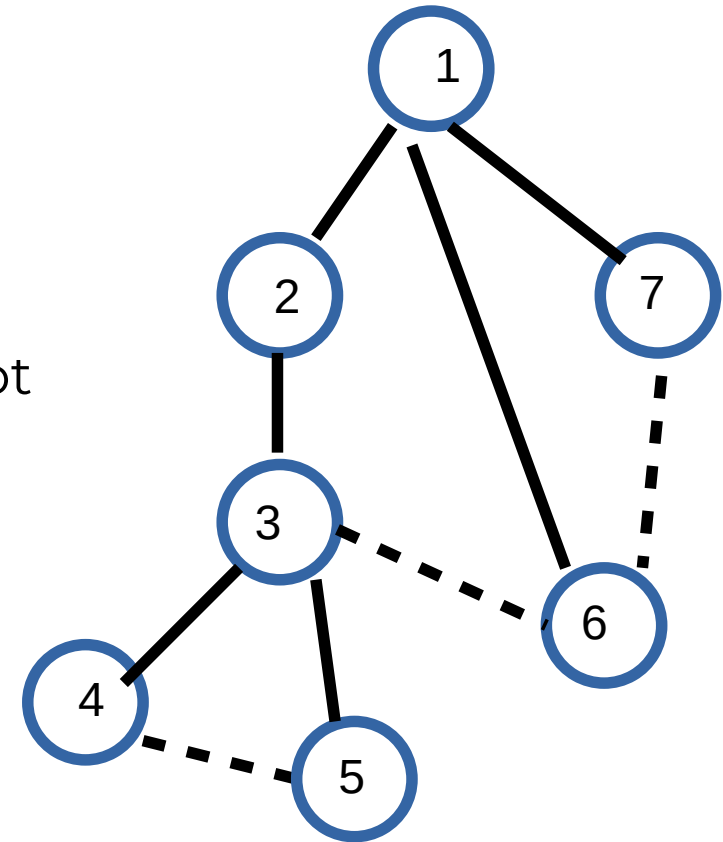
# What does 4 do when it hears from 2?

- **Message (Y, d, X) - (to, distance, from)**
- 2 hears (1, 0, 1) from 1
- 2 sends (1, 1, 2) to 3
- 3 sends (1, 2, 3) to 5 and 4
- 4 receives (1, 2, 3) from 3
- 4 receives (1, 3, 5) from 5
- Sets 1 as root (id=1 is < id=4)
- Prunes the 4-5 path since it is 4 hops compared to 3 hops via 3



# Failure and Downsides

- Even after the system has stabilized, the root continues to send messages periodically
  - Other bridges continue to forward these messages
- When a bridge fails, the downstream bridges will not receive the configuration messages
  - After waiting a specified period of time, they will once again claim to be the root and the algorithm starts again
- No load balancing





# Virtual LAN (VLANs)



- LANs are on the same Ethernet segments
- Does not scale very well – too many wires
- How can we put multiple people in different locations on the same Ethernet segment (LAN)?
- How do we create multiple LANs over the same wire?

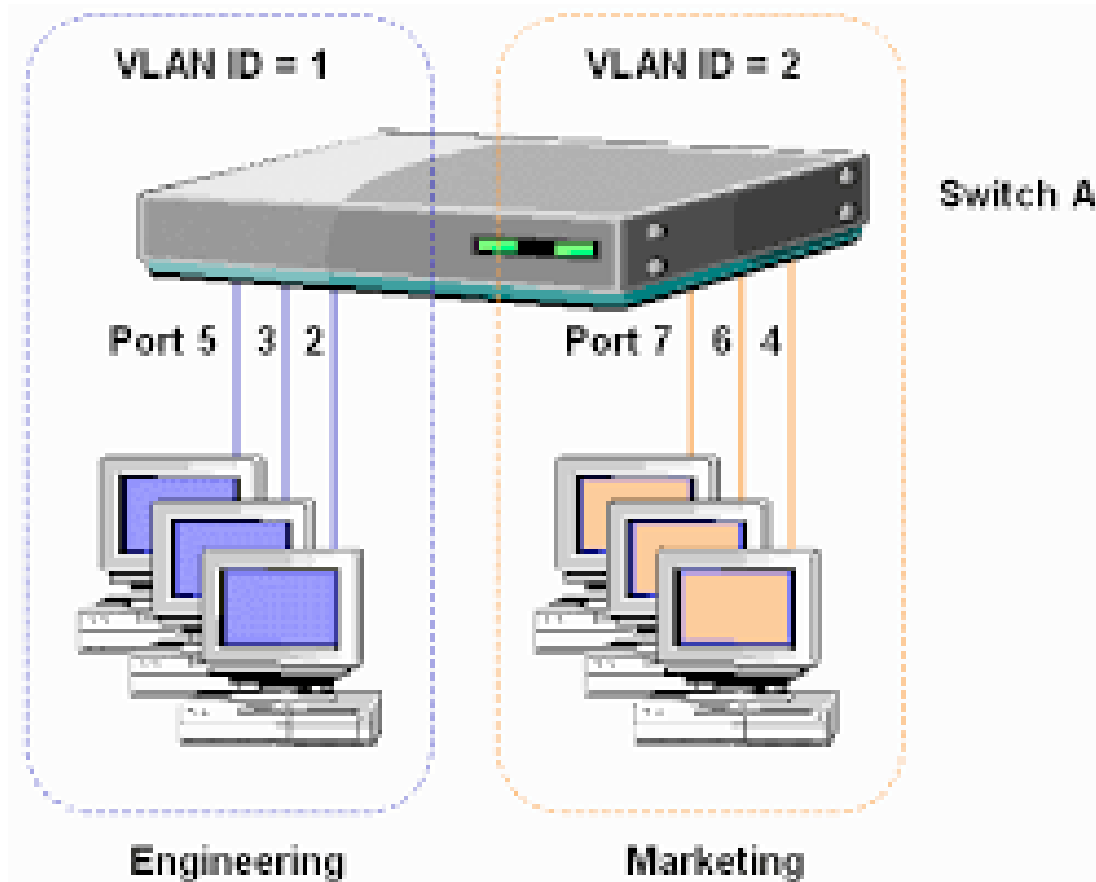
# Why separate at all?

- LANs are on the same Ethernet segments! Security.
- Isolation – sensitive traffic vs normal traffic
- Containment of traffic – your for loop broke the internet
- How do we create multiple LANs over the same wire?



# VLANs

---



- Switches specify which VLAN is accessible over which interface
- Each interface can have a VLAN color
- Each Mac address can have a interface color
- Add VLAN tag to the Ethernet header



# Link Layer Recap – All this for a cat picture

