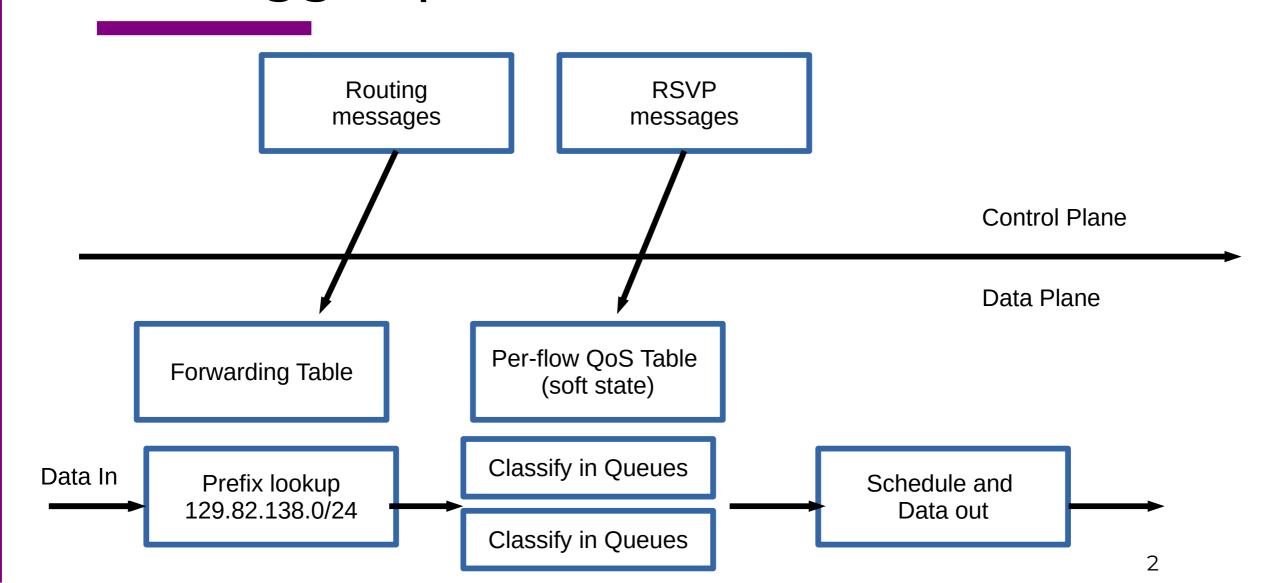
CSC7970 - NEXT-GENERATION NETWORKING

RESERVATION PROTOCOLS AND P2P

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The bigger picture - IntServ

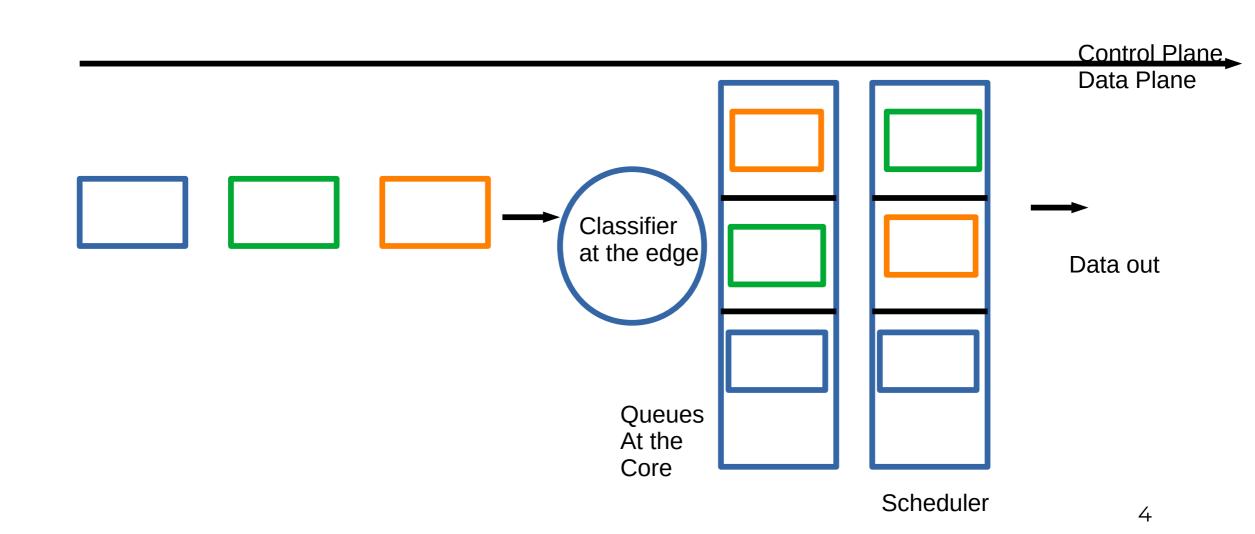


RSVP RESV messages Path messages Data Source Router Router Requester

Everyone needs to agree Per-flow – lots of state Should we create admission control?

Need 5Gbps and 20 ms latency

The bigger picture - DiffServ



QoS

- Different applications can benefit from different handling
 - IP is best effort
 - If you look at the header, it used to have a type-of-service field
 - Not used in production, and later changed to DSCP

												IF	Pv4	Head	der F	orm	at													V	VIK	pe	dia
Offsets	Octet				(0					1							2									3						
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	3 24	1 25	5 20	6 27	7 2	8 2	3	0 31
0	0	Version IHL									DSCP ECN						Total Length																
4	32	Identification													Flags Fragment Offset																		
8	64	Time To Live									Protocol							Header Checksum															
12	96		Source IP Address																														
16	128		Destination IP Address																														
20	160																																
24	192															0=+		/:£		٠,													
28	224															Opt	ons	(II IH	L > 5)													
32	256																																

QoS

- What does DSCP stand for?
 - Differentiated Services Code Point
 - Used for DiffServ, upto the application to decide how to use it
 - You can say things like
 - Expedited Forwarding
 - Class selection
 - and so on..

QoS

```
+----+
| bit0 | bit1 | bit2 | bit3 | bit4 | bit5 | bit6 | bit7 |
|----+
DSCP | ECN |
+----+
```

- PHB
 - 000 000 \rightarrow Default
 - Expedited Forwarding(EF) \rightarrow 101 110
- RFC 2474, 2597, 2598

Drop Priority

- All sent packets get same average delay
 - Why?
- Drop packets early to avoid TCP congestion control
- You will spend a lot of time to recover than dropping a few packets

The Current State of the Internet

- The Internet is becoming more and more centralized
 - Dominated by tech companies that want to control it
 - (Almost) any communication today has to go through a (cloud) server
 - Where it is mined, sold, resold, and used for various things
- A number of reasons to avoid centralization
 - Privacy
 - Dependability
 - Restriction of free speech

Decentralization through P2P

High-level idea

- Why having to go to a server to communication with other users?
- Let's communicate with them directly!
- BitTorrent
 - still dominates Internet traffic
 - Apple/MS/RedHat use them for software distribution

Where BitTorrent is not so great?

- Implementing a data-centric design on top of TCP/IP
 - No way to reason about trust
 - Explicit peer discovery (which is the IP address of each peer?)
 - Have to bind data to a specific peer (aka IP address)
 - Does not do well with "flash-crowd" scenarios
 - Peers have no clue how "far" or close another peer is
 - Problems with that?

Flash crowd scenarios





How does BitTorrent deal with flash crowds? (self.Bittorrent)

submitted 1 year ago by Mathbound314

When I tried to leech a popular torrent the day it came out, the seeder to leeches ratio was extremely low. It would have taken a week to torrent it. How does BitTorrent deal with this? Do they have a special protocol for flashcrowds?

1 comment share save hide report

BitTorrent flashcrowds only occur in very small fractions (0.3-2%) of the swarms but that they can affect over ten million users.

B. Zhang, A. Iosup, J. Pouwelse and D. Epema, "Identifying, analyzing, and modeling flashcrowds in BitTorrent," 2011 IEEE International Conference on Peer-to-Peer Computing, Kyoto, 2011, pp. 240-249.

How can we find torrent-files?

- Usually through a website
 - Users go online, search for what they want, and download the torrent-file

- Do we need to download the torrent-file from a website?
 - Not necessarily...

- Magnet links as an alternative
 - Download a torrent-file directly from another peer

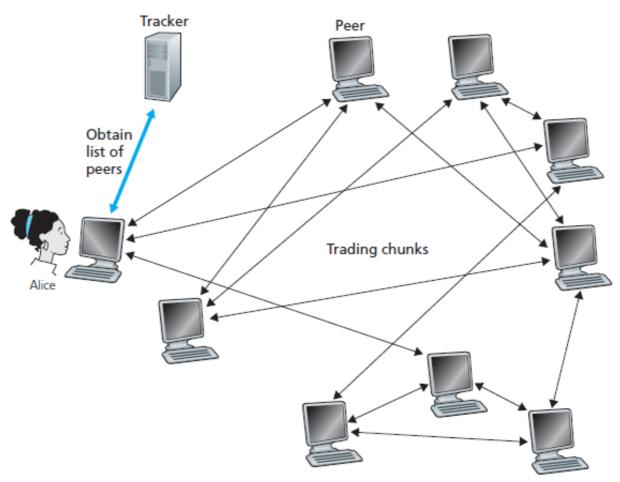
Peer Discovery

- Discovery of available peers to download a torrent from?
 - Operate on top of an IP network
 - We need to know the IP address of peers that have the data we want

- 2 ways to do it
 - Through a tracker centralized approach
 - Through a Distributed Hash Table (DHT) distributed approach

What happens if the tracker fails?

Example



Credits: https://electronicspost.com/how-bittorrent-works

DHT-based discovery

- How can we make this process fully distributed?
 - Instead of contacting a tracker, we can contact directly peers that may have the torrent we are interested in

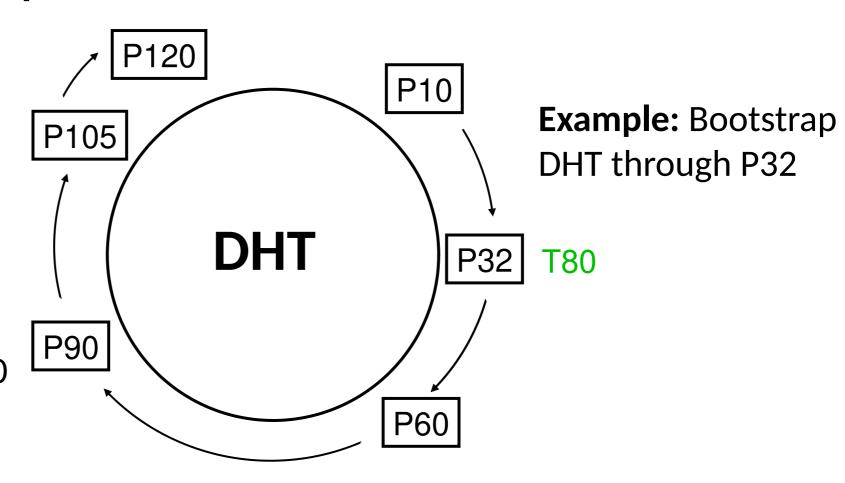
- How to find such a peer?
 - Hash torrent and peer ids (e.g., IP addresses)
 - Map each hash value to a point on the DHT
 - Ask a number of peers on the DHT until you find someone that knows peers that download (have) the torrent

DHT Example

Hash of torrent (e.g., found on a torrent website)

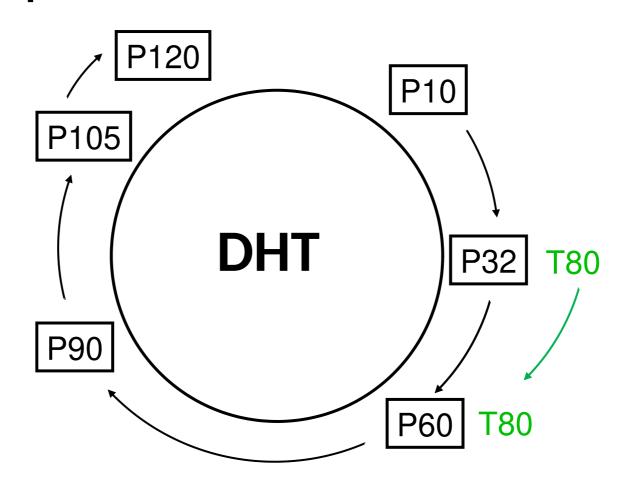
Example:

 $hash(linux_16.04) = T80$

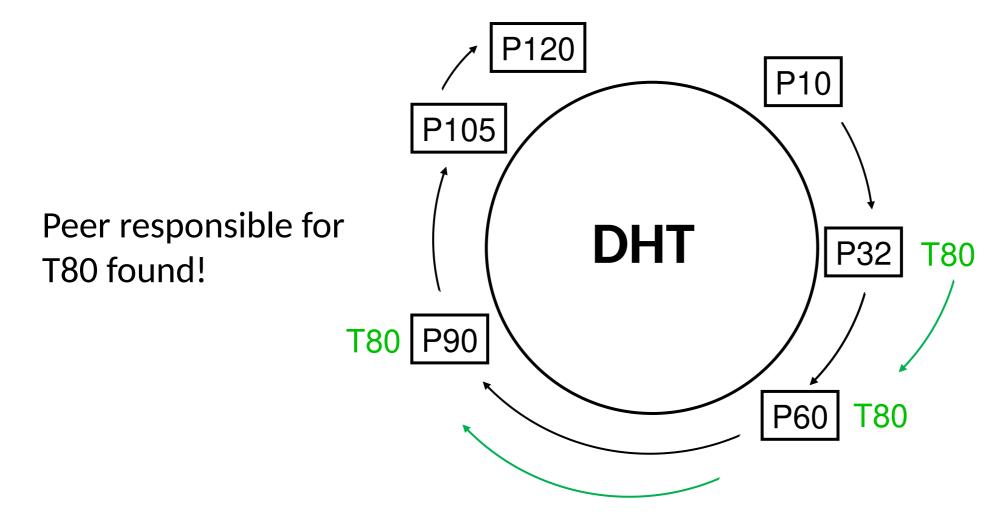


DHT Example

Query successor peer(s) on the DHT until a peer that is responsible for the torrent is found



DHT Example



Discovering what data a peer has

- Peers maintain a bitmap for each torrent
 - Each bit indicates whether a peer has a specific torrent piece (each piece consists of multiple network-layer packets)

- Peers exchange bitmaps with each other
 - Let each other know what data they have

Conclusion

- P2P aims to offer the means to decentralization on the current Internet
 - Peers communicate with each other directly
 - Peers focus on the data to share taking data integrity into account
- BitTorrent is a well-designed system for peer-to-peer communication
 - Hard to trick BitTorrent!
 - Does its best to offer data-centricity on top of a point-to-point network
 - Dominated traffic, slowed down, and picking up again