

# Protocol

Alice  $|H\rangle_x \leftarrow$  choice for  $x$

$|H\rangle_r \leftarrow$  choice for  $r$

$\frac{1}{2}I_H \leftarrow$  Alice's part of  $|GHZ\rangle$

$|0\rangle_S \leftarrow$  message to Bob

$|0\rangle_A \leftarrow$  Alice's GHZ outcome, to be determined

$S_0 =$  Alice's state after measuring to get a  $+$  creating  $S$ .

$\text{Tr}_S(S_0) =$  Alice's state after sending  $S$  to Bob.

Bob now has a purification of  $\text{Tr}_S(S_0)$ , so purifying the post-measured state, and  $S$ .

$S_1 = S_0 + G$  space. So  $\text{Tr}_G(S_1) = \text{Tr}_S(S_0)$

First check: Alice measures  $S_1$  to see if  $x \oplus y = 1$ .

$$\max \langle \Pi_{11} S_1 \rangle$$

$$\text{Tr}_G(S_1) = \text{Tr}_S(S_0)$$

This should give 0.75 (or less, perhaps).

# Full protocol SDP

Alice appends  $\frac{1}{2}I_y$  to  $S_1$  so Bob has a copy/  
a purification.

Bob responds w/  $d$  (in space  $D$ )  $d = b \oplus c$  here.

Alice checks if  $a \oplus d = xy(x \oplus y \oplus 1) \oplus 1$

SDP:  $\max \langle \Pi_2, S_2 \rangle$  checks if  $x \oplus y = 1$   
and  $a \oplus d = \dots$

$$\text{Tr}_D(S_2) = S_1 \otimes \frac{1}{2}I_y$$

$$\text{Tr}_G(S_1) = \text{Tr}_S(S_0)$$

$$S_0 \in \text{Pos}(H \otimes R \otimes A \otimes S)$$

$$S_1 \in \text{Pos}(H \otimes R \otimes A \otimes G)$$

$$S_2 \in \text{Pos}(H \otimes R \otimes A \otimes D)$$

$$\sigma_0 = \frac{1}{2}I_H \otimes |HX\rangle\langle X| \otimes |HX\rangle\langle R| \otimes |0\rangle\langle 0|_A \otimes |0\rangle\langle 0|_S$$

$$\sigma_1 = \underset{\uparrow \text{GHZ}}{U} \left( \frac{1}{2}I_H \otimes |HX\rangle\langle X| \otimes |HX\rangle\langle R| \otimes |0\rangle\langle 0|_A \right) U^\dagger \otimes |0\rangle\langle 0|_S$$

$$\sigma_2 = \text{creates } S \text{ from } a, x, r.$$

↑ This is  $S_0$ , broken up into steps.

## Continuity argument

Totally made-up lemma

Keep doing GHSZ  $n$  times.

$$\text{Then } \|S - S_{\text{actual}}\| \leq f(n)$$

↑  
as defined  
above

↑  
approximation

↑  
goes to 0 as  $n \rightarrow \infty$

## Lemma

$\alpha$  = SDP value from above.

$\alpha_{\text{actual}}$  = SDP value using  $S_{\text{actual}}$  instead of  $S$ .

$$|\alpha - \alpha_{\text{actual}}| \rightarrow 0 \text{ as } n \rightarrow \infty.$$

Proof: ① Take the dual

② SDP magic.

③ Profit.

□

Concern: Since in this particular protocol implementation, Bob does not separate the "b & c boxes" there is not really much "G#Z" happening. I can't remember if this should be an issue or not. (In any case,  $\frac{3}{4}$  is an upper bound.)

Question: If Bob sends back boxes B & C to Alice, do we get a nice SDP still?  
We might need NPA at that point.