

## § 2.1 Affine & Convex Sets

### § 2.1.1 Line & Line Segments

Def<sup>n</sup>: Line := For  $x_1 \neq x_2, \in \mathbb{R}^n$ ,  $y = \theta x_1 + (1-\theta)x_2$  with  $\theta \in \mathbb{R}$  is a line.



Def<sup>n</sup>: Line segment := For  $x_1 \neq x_2, \in \mathbb{R}^n$ ,  $y = \theta x_1 + (1-\theta)x_2$  with  $0 < \theta < 1$  is a line segment.

Remark:  $y = x_2 + \theta(x_1 - x_2)$  is also equivalent.  
Start from  $x_2$  & move along the right direction.

### § 2.1.2 Affine sets

Def<sup>n</sup>: Affine set := set  $C$  s.t. a line through any 2 points in  $C$ , lies inside  $C$ . (redundant)

= For  $x_1, x_2 \in C$ ,  $\theta x_1 + (1-\theta)x_2 \in C \quad \forall \theta \in \mathbb{R}$ .

Def<sup>n</sup>: Affine combination := given  $x_1, x_2, \dots, x_k$ , the point  $\theta_1 x_1 + \theta_2 x_2 + \dots + \theta_k x_k$  s.t.  $\sum \theta_i = 1$

Remark: From induction,  $\Rightarrow$  an affine set has all affine combinations.

Claim:  $V = C - x_0$  is a subspace, where  $C$  is an affine set,  $x_0 \in C$ .

Proof: Let  $v_1, v_2 \in V$  &  $\alpha, \beta \in \mathbb{R}$ . Now  $v_1 + x_0 \in C$  &  $v_2 + x_0 \in C$ .  
Also,  $\alpha v_1 + \beta v_2 + x_0 = \alpha(v_1 + x_0) + \beta(v_2 + x_0) + (1-\alpha-\beta)x_0 \in C$   
from the remark. This means  $\alpha v_1 + \beta v_2 \in V$  since  $(\alpha v_1 + \beta v_2) + x_0 \in C$ .

$$\text{[viz. } \underbrace{\alpha v_1 + \beta v_2}_{\in V} = C - x_0]$$

Concl<sup>n</sup>: Affine set  $C$  can be written as:  $C = V + x_0$  (a vector space + an offset).

Def<sup>n</sup>: Dimension of an affine set  $C$  := Dimension of the subspace  $V = C - x_0$

NB: The subspace  $V$  doesn't depend on the choice of  $x_0$ .

E.g.: Soln. of linear eq<sup>n</sup>s is an affine set:  $C = \{x \mid Ax = b\}$ .

Let  $x' = \theta x_1 + (1-\theta)x_2$ , &  $x_1, x_2$  be solns. Then

$$Ax' = \theta Ax_1 + (1-\theta)Ax_2 = (\theta + 1 - \theta)b = b \Rightarrow x' \text{ is also a soln.}$$

$\Rightarrow C$  is an affine set.

: Every affine set can be expressed as a soln. of linear eq<sup>n</sup>s. (§#1)

Def<sup>n</sup>: Affine Hull := set of all affine combinations of some set  $S \subseteq \mathbb{R}^n$ .  
 $\therefore \text{aff}(S) = \{\theta_1 x_1 + \dots + \theta_k x_k \mid \sum \theta_i = 1, x_i \in S\}$

NB:  $\text{aff}(S) \subseteq C$  where  $C$  is affine &  $S \subseteq C$ .  $\text{aff}(S)$  is the smallest affine set with  $S$  in it.

### § 2.1.3 Affine dimension & relative interior

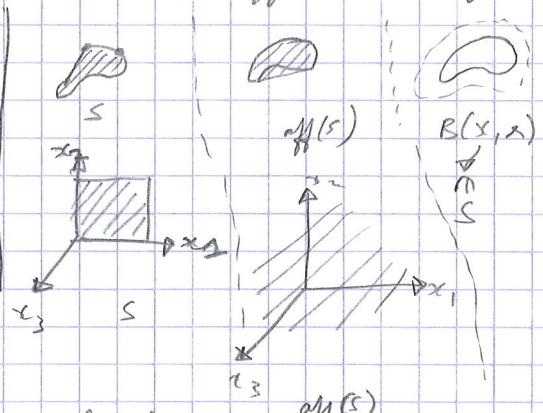
Def<sup>n</sup>: Affine dim := Dimension of  $\text{aff}(S)$  for any set  $S$ .

NB: Affine dim of unit circle is 2 unlike most def<sup>n</sup>s.

Consider:  $\dim(\text{aff}(S) + \mathbb{R}^n)$  in but  $S \subseteq \mathbb{R}^n$ .

Def<sup>n</sup>: Relative interior :=  $\text{relint}(S) = \{x \in S \mid B(x, \epsilon) \cap \text{aff}(S) \subseteq S \text{ for } \epsilon > 0\}$

→ Not clear, eg 2.2

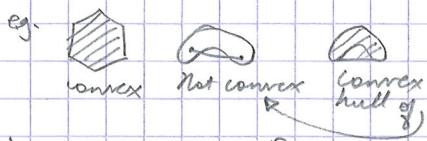


### § 2.1.4 Convex sets

Def<sup>n</sup>: Convex: If  $x_1, x_2 \in C$ , for  $0 \leq \theta \leq 1$ ,  $\theta x_1 + (1-\theta)x_2 \in C$ .

NB: Every affine set is convex.

Def<sup>n</sup>: Convex combination :=  $\sum \theta_i x_i$ , s.t.  $\sum \theta_i = 1 \wedge \theta_i \geq 0 \forall i$ .



Claim: Set is convex  $\Leftrightarrow$  contains every convex combination. (if time permits)

Def<sup>n</sup>: Convex hull := (implicit)  $\text{conv}(S) = \{\sum \theta_i x_i \mid x_i \in S, \sum \theta_i = 1, \theta_i \geq 0 \forall i\}$

Remark: Smallest convex set containing  $S$  is  $\text{conv}(S)$

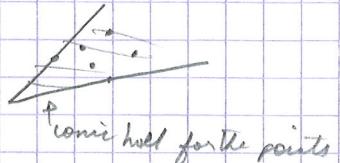
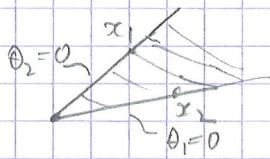
Generalisations: (a) Infinite sums:  $\theta_i \geq 0, \sum \theta_i = 1$ ;  $\sum \theta_i x_i \in S$  if convex.

(b) Integrals:  $p: \mathbb{R}^n \rightarrow \mathbb{R}$ ,  $p(x) > 0 \forall x \in S$ ,  $\int_S p(x) dx = 1$ ;  $\int_S p(x) x dx \in S$

→ Not clear.

(c) Probabilities:

### § 2.1.5 Cones

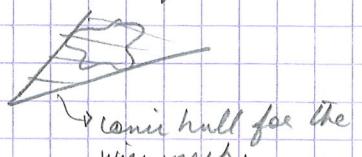


Def<sup>n</sup>: Cone: A set  $S$  is a cone if  $\forall x \in S, \theta \geq 0$  we have  $\theta x \in S$ .

Def<sup>n</sup>: Convex cone: If a set is both convex & a cone.  $\Leftrightarrow \theta_1 x_1 + \theta_2 x_2 \in S \text{ if } \theta_i \geq 0$

Def<sup>n</sup>: Conic combination := (non-neg. linear combination)  $\theta_1 x_1 + \dots + \theta_k x_k$  for  $\theta_i \geq 0$ .

Def<sup>n</sup>: Conic Hull := (implicit)



## § 2.2 Some Important Examples

- \* Affine ( $\Rightarrow$  convex):
  - (a)  $\emptyset$  (null set)
  - (b)  $\{x_0\}$  (singleton)
  - (c)  $\mathbb{R}^n$
- \* (a) line — affine
- (b) line through the origin — subspace & also convex

- \* line segment — convex (not affine unless a point)
- \* ray,  $\{x_0 + \theta v \mid \theta \geq 0\}$ ,  $v \neq 0$ ; convex (not affine)
- \* convex cone if  $x_0 = 0$ .

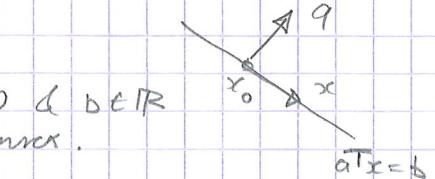
- \* subspace — affine, convex & a convex cone.

### § 2.2.1 Hyperplanes & halfspaces

Def<sup>n</sup>: Hyperplane :=  $\{x \mid a^T x = b\}$  where  $a \in \mathbb{R}^n$ ,  $a \neq 0$  &  $b \in \mathbb{R}$

Remarks: It is a soln. of a set of linear eq's. Thus it's convex.

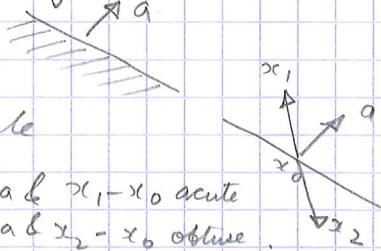
- : Recall the geometric interpretation
- : can write as  $\{x \mid a^T(x - x_0) = 0\}$  ( $a^T x_0 = b$ )
- : also as  $x = x_0 + a^\perp$  where  $a^\perp = \{v \mid a^T v = 0\}$  to satisfy the sameeq?
- : The hyperplane is an offset  $x_0$  plus all vectors orthogonal  $a$  (normal vector).



Def<sup>n</sup>: Halfspace (closed) :=  $\{x \mid a^T x \leq b\}$

Remarks: Halfspaces are convex but not affine.

- : The  $\{x \mid a^T(x - x_0) \leq 0\}$  rep helps say that the angle of  $x - x_0$  must be obtuse (or right).



### § 2.2.2 Euclidian ball & ellipsoids

Def<sup>n</sup>: (Euclidian) ball in  $\mathbb{R}^n$  :=  $B(x_c, r) = \{x \mid \|x - x_c\|_2 \leq r\} = \{x \mid (x - x_c)^T(x - x_c) \leq r^2\}$   
for  $r > 0$ ,  $\|\cdot\|_2$  is the euclidian norm ( $\|u\|_2 = \sqrt{u^T u}$ )

Remarks:  $x_c$  is the centre,  $r$  is the radius.

- : another rep  $B(x_c, r) = \{x_c + ru \mid \|u\|_2 \leq 1\}$

Claim: A euclidian ball is a convex set.

Proof: Let  $\|x_1 - x_c\|_2 \leq r$ ,  $\|x_2 - x_c\|_2 \leq r$ . For  $0 \leq \theta \leq 1$ ,

$$\begin{aligned} \|(\theta x_1 + (1-\theta)x_2) - x_c\|_2 &= \|\theta(x_1 - x_c) + (1-\theta)(x_2 - x_c)\|_2 \\ &\leq \theta \|x_1 - x_c\|_2 + (1-\theta) \|x_2 - x_c\|_2 \\ &\leq r \end{aligned}$$

Def<sup>n</sup>: Ellipsoid :=  $E = \{x \mid (x - x_c)^T P^{-1}(x - x_c) \leq 1\}$  for  $P = P^T \succ 0$ , i.e.  $P$  is symmetric & positive definite.

Claim: Semi-axis lengths of  $E$  are  $\sqrt{\lambda_i}$  where  $\lambda_i$  are eigenvalues of  $P$ .

Proof: Hint:  $P = \lambda^2 I$  yields a ball; explains also the  $P^{-1}$  instead of  $P$ .

Remark: <missing>

### § 2.2.3 Norm balls & norm cones.

+ Def<sup>n</sup>:  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is a norm if  $\star f \geq 0 \wedge x \in \mathbb{R}^n$

$$\star f(x) = 0 \text{ only if } x = 0; x = 0 \Rightarrow f(x) = 0$$

$$\star f(tx) = |t|f(x)$$

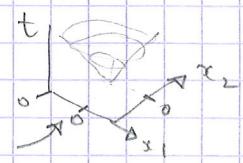
$$\star f(x+y) \leq f(x) + f(y), \forall x, y \in \mathbb{R}$$

+ Def<sup>n</sup>: Norm Ball := Radius  $r$  & centre  $x_c$   $\{x \mid \|x - x_c\| \leq r\}$

Claim: Norm balls are convex  $\Leftrightarrow$

Def<sup>n</sup>: Norm Cone :=  $\{(x, t) \mid \|x\| \leq t\} \subseteq \mathbb{R}^{n+1}$

e.g. Boundary of second-order cone in  $\mathbb{R}^3$   $\{(x, x_2, t) \mid (x_1^2 + x_2^2)^{1/2} \leq t\}$



## § 2.2.4 Polyhedra

Def<sup>n</sup>: polyhedron :=  $P = \{x \mid a_j^T x \leq b_j, j=1 \dots m; c_j^T x = d_j, j=1 \dots p\}$

Remarks: It's an intersection of halfspaces & hyperplanes.

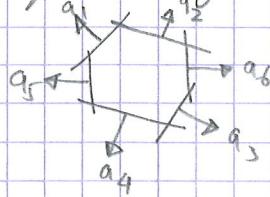
e.g. affine sets (e.g. subspaces, hyperplanes, lines), rays, line segments & halfspaces

Def<sup>n</sup>: polytope := a bounded polyhedron.

Not<sup>n</sup>: polyhedron :=  $P = \{x \mid Ax \leq b, Cx = d\}$  with

$$A = \begin{bmatrix} a_1^T \\ \vdots \\ a_m^T \end{bmatrix}, C = \begin{bmatrix} c_1^T \\ \vdots \\ c_p^T \end{bmatrix}$$

& the symbol  $\leq$  denotes component wise inequality in  $\mathbb{R}^m$ ;



## Simplices

Def<sup>n</sup>: affinely independent :=  $v_0, \dots, v_k$  are so if  $v_1 - v_0, \dots, v_k - v_0$  are linearly independent.

Def<sup>n</sup>: simplex := For  $v_0, \dots, v_k \in \mathbb{R}^n$  affinely independent,  $C = \text{conv}\{v_0, \dots, v_k\}$   
 $= \{\theta_0 v_0 + \dots + \theta_k v_k \mid \theta \geq 0, \mathbf{1}^T \theta = 1\}$

Remark: Affine dimension is  $k$  for the aforesaid.

E.g. 1-dimensional simplex: line segment. 3-dimensional simplex: tetrahedron

2-dimensional simplex: triangle (plus interior)

Def<sup>n</sup>: Unit simplex :=  $n$ -dimensional simplex determined by  $0, e_1, \dots, e_n \in \mathbb{R}^n$   
 $x \geq 0, \mathbf{1}^T x = 1$  (vector).

Def<sup>n</sup>: Probability simplex := <skipped>

Aim: To describe a simplex as a polyhedron; (a)  $x \in C$  iff  $x = \theta_0 v_0 + \dots + \theta_k v_k$  for  $\theta \geq 0$   
(b) equivalently, with  $y = (\theta_1, \dots, \theta_k)$ ,  $B = [v_1 - v_0, \dots, v_k - v_0] \in \mathbb{R}^{(n-1) \times k}$ , we have  
 $x \in C$  iff  $x = v_0 + By$ . (B#2) for  $y \geq 0$  with  $\mathbf{1}^T y \leq 1$

(c)  $B$  has rank  $k$  ::  $v_0, v_1, \dots, v_k$  are affinely independent.  $\Rightarrow \exists$  a non-singular matrix, s.t.  $AB = [A_1, A_2] B = \begin{bmatrix} I \\ 0 \end{bmatrix}$  (row echelon form)

(d)  $A \{x = v_0 + By\}$  yields  $A_1 x = A_1 v_0 + y, A_2 x = A_2 v_0 \Rightarrow x \in C$  iff (i)  $y = A_2 x - A_2 v_0$  granted  
 $y \geq 0, \mathbf{1}^T y \leq 1$   
(ii)  $A_2 x = A_2 v_0$

These are just a set of linear equalities & inequalities in  $y$ . Thus this is a polyhedron.

## Convex Hull Description of Polyhedra

Recall:  $\text{conv}\{v_1, \dots, v_k\} = \{\theta_1 v_1 + \dots + \theta_k v_k \mid \theta \geq 0, \mathbf{1}^T \theta = 1\}$

## § 2.2.5 The positive semi-definite cone.

Def<sup>n</sup>:  $S^n := \{X \in \mathbb{R}^{n \times n} \mid X = X^T\}$  (symmetric matrices)

Remark:  $S^n$  has dimension  $n(n+1)/2$

Def<sup>n</sup>:  $S_+^n := \{X \in S^n \mid X \succeq 0\}$  (symmetric positive definite matrix)

Def<sup>n</sup>:  $S_{++}^n := \{X \in S^n \mid X > 0\}$  (symmetric positive definite matrix)

Remark: Analogous to  $\mathbb{R}_+$  &  $\mathbb{R}_{++}$ .

Recall: Positive semi-definite  $\Rightarrow y^T A x \geq 0$ , then  $A$  is PSD.

Claim:  $S_+^n$  is a convex cone: if  $\theta_1, \theta_2 \geq 0$ , &  $A, B \in S_+^n$ , then  $\theta_1 A + \theta_2 B \in S_+^n$ .

Proof:  $x^T(\theta_1 A + \theta_2 B)x = \theta_1 x^T A x + \theta_2 x^T B x \geq 0$  if of course,  $A \succeq 0$  &  $B \succeq 0$  with  $\theta_1, \theta_2 \geq 0$ .

E.g.  $X = \begin{bmatrix} x & y \\ y & z \end{bmatrix} \in S_+^2 \Leftrightarrow \begin{array}{l} x \geq 0, z \geq 0, xz \geq y^2 \\ \text{#4} \end{array}$

## § 2.3 Operations that preserve convexity

Motivation: These & simple examples as described should help determine/establish convexity of sets.

### § 2.3.1 Intersection

Oct 11, 2016

Claim: Convexity is preserved by intersection; if  $S_1, S_2$  are convex  $\Rightarrow S_1 \cap S_2$  is also convex.

Remark: Extends to infinite sets; if  $S_\alpha$  is convex,  $\bigcap_\alpha S_\alpha$ , ( $\Delta E A$ ) is also convex.

E.g.: (1) A polyhedron - it's convex. It is made of intersection of halfspaces & hyperplanes. Both are convex.

(2) Positive Semi Definite Cones ( $S_+^n$ ) - The can be written as

$$\bigcap \{x \in S^n \mid z^T x \geq 0\}$$

- The set  $\{x \in S^n \mid z^T x \geq 0\}$  is a halfspace (since the constraint is linear) in  $S^n$ . It is hence convex.

- Concl<sup>n</sup>: The positive semi-definite cone is itself convex (is an intersection of convex sets).

(3) Consider set  $S = \{x \in \mathbb{R}^m \mid |p(x)| \leq 1 \text{ for } |t| \leq T/3\}$  for  $p(t) = \sum_{k=1}^m b_k \cos kt$   
< easier example, skipped writing it. :)

Claim: Every closed convex set is an intersection of halfspaces;  $C = \bigcap \{S_i \mid \exists \text{ halfspaces}$

Doubt: Proof? - Mather's argument:  $x, y \in A \cap B \Rightarrow x \in A, y \in B$ . Now for  $0 < \theta < 1$ ,  $C \subseteq B$ ?  
 $\theta x + (1-\theta)y \in A, B$  from convexity of  $A, B$ .  
obviously then  $\theta x + (1-\theta)y \in A \cap B \Rightarrow A \cap B$  is convex.

### § 2.3.2 Affine Functions

Def<sup>n</sup>: Affine function  $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$  is affine if it has the form  $f(x) = Ax + b$  where  $A \in \mathbb{R}^{m \times n}$  &  $b \in \mathbb{R}^m$ .

Claim: Let  $S \subseteq \mathbb{R}^n$  be a convex set & let  $f$  be affine. Then  $f(S)$  is also convex.  
if  $f$  is affine, the inverse image of  $S$  is convex; viz.  $f^{-1}(S) = \{x \mid f(x) \in S\}$  is also convex.

Proof: Let  $x_1, x_2 \in S \Rightarrow \theta x_1 + (1-\theta)x_2 \in S$  ( $\theta \in [0, 1]$ ).

$$f(x_1), f(x_2) \in f(S). \text{ Also } f(\theta x_1 + (1-\theta)x_2) = \theta f(x_1) + (1-\theta)f(x_2)$$

$$\stackrel{\text{f is affine}}{=} \theta f(x_1) + (1-\theta)f(x_2)$$

$$\stackrel{\text{f is affine}}{=} f(\theta x_1 + (1-\theta)x_2) \in f(S) \quad \square$$

: (i) Think  $f^{-1}(B) = A^{-1}B - A^{-1}b$  & then the usual proof can be used)

e.g. (1) Dealing & Translation:  $\alpha S$  &  $S + \alpha$

(2) Projection: If  $S \subseteq \mathbb{R}^m \times \mathbb{R}^n$ , then  $\{x_1 \in \mathbb{R}^m \mid (x_1, x_2) \in S \text{ for some } x_2 \in \mathbb{R}^n\}$

Claim: The sum of 2 convex sets is also convex,  $S_1 + S_2 = \{x+y \mid x \in S_1, y \in S_2\}$  is convex.

Proof: Claim: If  $S_1, S_2$  are convex, then  $\Delta E A$  is their Cartesian product.

Proof: For  $(x_1, x_2) \in S_1 \times S_2$  &  $(y_1, y_2) \in S_1 \times S_2$ ,  $(\theta x_1 + (1-\theta)y_1, \theta x_2 + (1-\theta)y_2) \in S_1 \times S_2$ .  
by convexity of  $S_1$  &  $S_2$ .

Show a linear  $f^n$  on  $S_1 \times S_2 \rightarrow f(x_1, x_2) = x_1 + x_2$  (is also affine). Thus, the result follows from the result (claim 1 in § 2.3.2).  $\square$

$\Delta$  of  $^n$ : Partial sum := For  $S_1, S_2 \in \mathbb{R}^n \times \mathbb{R}^m$ ,  $S = \{(x, y_1 + y_2) \mid (x, y_1) \in S_1, (x, y_2) \in S_2\}$   
 Claim: Partial sums of convex sets are convex.  $\square$  (proof)

e.g. (3) Polyhedron; recall:  $\{x \mid Ax \leq b, Cx = d\}$ ; this can be expressed as an inverse image of an affine  $f^n$ , with a domain  $\mathbb{R}^n \setminus \{0\}$ , as follows.

$$\{x \mid Ax \leq b, Cx = d\} = \{x \mid f(x) \in \mathbb{R}^n \times \{0\}\} \text{ where } f(x) = (b - Ax, d - Cx)$$

(4) Sol $^n$  set of a linear matrix inequality: The soln. set of LMI is convex where LMI is  
 def $^n$ : Linear Matrix Inequality:  $\exists A(x) = x_1 A_1 + \dots + x_n A_n \leq B, A_i \in \mathbb{R}^{m \times m}$ .  
 (linear in  $x$ )

Proof: Use  $f(x) = B - A(x)$  (an affine  $f^n$ ) & the positive semidefinite cone  $\mathcal{S}^m$   
 $f: \mathbb{R}^n \rightarrow \mathcal{S}^m$  then it will easily follow.

(5) Hyperbolic cone (Def $^n$ ):  $\{x \mid x^T P x \leq (c^T x)^2, c^T x \geq 0\}$  where  $P \in \mathcal{S}_+^n$  &  $c \in \mathbb{R}^n$   
 Claim: Hyperbolic cone is convex.

Proof: Recall that the second order cone is convex, i.e.  $\{(z, t) \mid z^T z \leq t^2, t \geq 0\}$   
 Use  $f(x) = (P^{1/2}x, c^T x)$  & then it's obvious ( $z = P^{1/2}x, t = c^T x$ )

### § 2.3.3 Linear-fractional & Perspective Functions

Def $^n$ : Perspective  $f^n := P: \mathbb{R}^{n+1} \rightarrow \mathbb{R}^n$  with domain  $\mathbb{R}^n \times \mathbb{R}_{++}$  &  $P(z, t) := \underline{z}$

Remark: The perspective function scales everything to the last component & drops it.

Claim: If  $C (\subseteq \text{dom } P)$  is convex, then  $P(C)$  is also convex.

: Perspective  $f^n$ 's preserve convexity.

Proof: Let  $x = (x_1, x_{n+1}) \succ y = (y_1, y_{n+1}) \in \mathbb{R}^{n+1}$  with  $x_{n+1}, y_{n+1} > 0$ . For  $0 < \theta < 1$ , we have

$$P(\theta x + (1-\theta)y) = \frac{\theta z + (1-\theta)\bar{y}}{\theta x_{n+1} + (1-\theta)y_{n+1}} = \mu P(x) + (1-\mu)P(y) \text{ where } \mu = \frac{\theta x_{n+1}}{\theta x_{n+1} + (1-\theta)y_{n+1}} \text{ (not too hard to see)}$$

Claim:  $\theta$  &  $\mu$  are monotonically related.  $\Leftrightarrow$  for  $\theta=0, \mu=0; \theta=1, \mu=1$ .

So then convexity follows from def $^n$  of it 'cause just use  $\frac{\theta x_{n+1}}{\theta x_{n+1} + (1-\theta)y_{n+1}} > \frac{\theta x_{n+1}}{\theta x_{n+1} + (1-\theta)y_{n+1}}$   
 &  $y$  are chosen from a convex set.

& that the denominator only inc. or dec. with  $\theta$  compared to the rest.

Claim: The inverse image  $P^{-1}$  is also convex. i.e.

for  $C \subseteq \mathbb{R}^n$  convex,  $P^{-1}(C) = \{(x, t) \in \mathbb{R}^{n+1} \mid x/t \in C, t > 0\}$  is convex.

Proof: Suppose:  $(x, t), (y, s) \in P^{-1}(C)$  &  $0 \leq \theta < 1$ . We must show that

$$\theta(x, t) + (1-\theta)(y, s) \in P^{-1}(C), \text{ i.e. } \frac{\theta x + (1-\theta)y}{\theta t + (1-\theta)s} \in C \quad (\theta t + (1-\theta)s > 0 \text{ is obvious})$$

Recall  $\frac{\theta x + (1-\theta)y}{\theta t + (1-\theta)s} \in C$  where  $\mu = \frac{\theta t}{\theta t + (1-\theta)s}$

### Linear-fractional functions

Def $^n$ : Linear-fractional function := composition of perspective  $f^n$ 's with an affine  $f^n$ .

: let  $g: \mathbb{R}^n \rightarrow \mathbb{R}^{n+1}$  be affine, i.e.  $g(x) = \begin{bmatrix} A \\ C^T \\ d \end{bmatrix} x + \begin{bmatrix} b \\ 0 \end{bmatrix}$  where  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ ,  
 $C \in \mathbb{R}^{n \times m}$  &  $d \in \mathbb{R}^n$

$f: \mathbb{R}^n \rightarrow \mathbb{R}^m$  given by  $f = P \circ g$ , i.e.  $f(x) = \frac{(Ax+b)}{(C^T x + d)}$ ,  $\text{dom}(f) = \{x \mid C^T x + d > 0\}$

Remark: One can write the linear-fractional  $f^n$  as  $\underline{x} = \begin{bmatrix} A & b \\ C^T & d \end{bmatrix} \in \mathbb{R}^{(m+1) \times (n+1)}$

so that  $(x, 1)$  yields  $(Ax+b, C^T x + d)$ . Now the last component is normalized to one.

: One can think of  $\mathbb{R}^n$  associated with rays in  $\mathbb{R}^{n+1}$  as  $P(z) = \{t(z, 1) \mid t > 0\}$   
 for  $z \in \mathbb{R}^n$  &  $t \in \mathbb{R}$ . Conversely any ray in  $\mathbb{R}^{n+1}$  through the origin  
 can be written as  $P(v) = \{t(v, 1) \mid t > 0\}$ .  $P$  &  $\mathbb{R}^n$  are one-one &  
 onto.  $\square$  (more details).

Claim: The linear-fractional  $f'$  can be expressed as  $f(x) = P^{-1}(\alpha P(x))$   
with  $\alpha \neq 0$  as defined earlier.

Proof:  $P(x)$  for  $x \in \text{dom } f$  i.e. with  $(Tx+d) > 0$  will yield a ray  $\{t(x, 1) | t > 0\}$  which is then acted upon by  $P$  to yield  $t(Ax+b), t(cx+d)$  which upon inversion,  $P^{-1}$ , yields the correctly transformed  $x$  to its trivial really.

Claim: The linear-fractional  $f'$  preserves convexity; so does its inverse.

Proof: Affine map is convex, projection map is convex.

e.g. Conditional probabilities: Let  $u \in \{1, \dots, n\}$ ,  $v \in \{1, \dots, m\}$  be random variables  
 $\& p_{ij} = \Pr(u=i, v=j)$ .  $f_{ij} = \Pr(u=i | v=j) = \frac{p_{ij}}{\sum_k p_{kj}}$   
 is a convex set, if the joint  
 is a convex set.  $\sum_k p_{kj}$

## § 2.4 Generalized Inequalities

### § 2.4.1 Proper cones & generalized inequalities

Def': Proper cone := A cone  $K \subseteq \mathbb{R}^n$  is proper if 

- $K$  is convex
- $K$  is closed
- $K$  is solid (non-empty interior)
- $K$  is pointed (contains no line);  $x \in K, -x \in K \Rightarrow x = 0$

Motivat': Use this to define generalized inequalities.

Def': Generalized (non-strict) inequality :=  $x \leq_K y \Leftrightarrow y - x \in K$

Def': Generalized (strict) inequality :=  $x <_K y \Leftrightarrow y - x \in \text{int } K$

Recall:  $\mathbb{R}_+$  is convex, (closed? solid?) & contains no line (pointed)  $\Rightarrow \mathbb{R}_+$  is a proper cone.

E.g.(1):  $K = \mathbb{R}_+$  yields  $\leq_K$  as  $\leq$  the usual ordering on reals.

E.g.(2):  $S_+^n$  is a proper cone  $\subseteq S^n$ :  $x <_K y \Leftrightarrow y - x \in S_+^n$ .  $K$  is often dropped.

E.g.(3): Claim: "cone of polynomials non-negative on  $[0, 1]$ " is proper:  $K = \{c \in \mathbb{R}^n | c_1 + c_2 t + \dots + c_n t^{n-1} \geq 0 \forall t \in [0, 1]\}$

"Proof": cone: let  $c \in K$ , then  $\theta c \in K$  obviously ( $\theta > 0$ ).

convex: let  $c, d \in K$ . To show  $\theta c + (1-\theta)d$  is in  $K$ , we must have

$$\theta c_1 + (1-\theta)d_1 + (\theta c_2 + (1-\theta)d_2)t + \dots + (\theta c_n + (1-\theta)d_n)t^{n-1} \geq 0$$

$$\text{since } (\theta c_1 + (1-\theta)d_1) + ((1-\theta)(d_2 + d_3 + \dots)) \geq 0$$

solid: (not sure) claim: interior is  $\{x | c_1 + c_2 t + \dots + c_n t^{n-1} > 0\}$

Doubt

— Doesn't the notion of "interior" already require some sort of ordering?

pointed: trivial (show no line)

Now for  $c, d \in \mathbb{R}^n$ ,  $c \leq_K d$  iff  $c_1 - d_1 + (c_2 - d_2)t + \dots + (c_n - d_n)t^{n-1} \geq 0 \Leftrightarrow c_1 + c_2 t + \dots + c_n t^{n-1} \geq d_1 + d_2 t + \dots + d_n t^{n-1}$

Properties: (1)  $\leq_K$  is preserved under addition: if  $x \leq_K y$  &  $u \leq_K v \Rightarrow x+u \leq_K y+v$  (§#5)

(2)  $\leq_K$  is transitive: if  $x \leq_K y$  &  $y \leq_K z \Rightarrow x \leq_K z$

(3)  $\leq_K$  is preserved under non-negative scaling: if  $x \leq_K y$  &  $\alpha \geq 0 \Rightarrow \alpha x \leq_K \alpha y$

(4)  $\leq_K$  is reflexive:  $x \leq_K x$

(5)  $\leq_K$  is anti-symmetric: if  $x \leq_K y$  &  $y \leq_K x$  then  $x = y$  (§#5)

(6)  $\leq_K$  is preserved under limits: if  $x_i \leq_K y_i$  for  $i=1, 2, \dots$   $x_i \rightarrow x$  &  $y_i \rightarrow y$  as  $i \rightarrow \infty$ , then  $x \leq_K y$

Proofs: most are simple, others in rough.

(7) if  $x \leq_K y \Rightarrow x \leq_K y$  (8) if  $x \leq_K y$  &  $u \leq_K v \Rightarrow x+u \leq_K y+v$   
 (9) if  $x \leq_K y$  &  $\alpha > 0 \Rightarrow \alpha x \leq_K \alpha y$  (10)  $x \leq_K x$  (11) if  $x \leq_K y \Rightarrow u \leq_K v$  &  $v$  small enough,  $x+u \leq_K y+v$

Doubt  $\leq_K$  instead of  $\leq_K$  (for 9)

not sure what "small" would mean (for 11)

## § 2.4.2 Minimum & minimal elements

Def<sup>n</sup>: Linear ordering := any two elements of a set are comparable.

Remark: In reals, either  $x \leq y$  or  $y \leq x$ . This doesn't hold for generalised inequalities.

Implication: Notions of min & max get complicated.

Def<sup>n</sup>: Minimum Element :=  $x \in S$  is a minimum element if  $\forall y \in S, x \leq y$ .

Imp. Def<sup>n</sup>: Maximum Element := (implicit)

Def<sup>n</sup>: Minimal Element :=  $x \in S$  is a minimal element if  $\forall y \in S, y \leq x$  only if  $y = x$

Claim: A set can have many minimal elements but only one minimum element.

Proof:  $\text{Q.E.D.}$  e.g. unique minimum

$\begin{array}{c} k+1 \\ \vdots \\ k+1 \end{array}$  multiple

All<sup>n</sup> Def<sup>n</sup>: Minimum Element :=  $x \in S$  is minimum iff  $S \subseteq x + K$

All<sup>n</sup> Def<sup>n</sup>: Minimal Element :=  $x \in S$  is minimal iff  $(x - K) \cap S = \{x\}$

Remark: For reals, these coincide.

E.g. (1) consider:  $\mathbb{R}_+^n$  (component wise inequality inducing cone).  $x \leq y$  would mean  $y$  is on top & right of  $x$ . If  $x$  is a minimum element then all points of the set lie on top & right of  $x$ . If  $x$  is a minimal element, then no point of the set lies below & on the left of the set.

(2) Set of symmetric Matrices:  $\forall A \in S_{++}^n$ , consider an ellipsoid (centred at origin)

$$E_A = \{x \mid x^T A^{-1} x \leq 1\}$$

Define:  $A \leq B$  iff  $E_A \subseteq E_B$  (assumed for now; not worrying about cones)

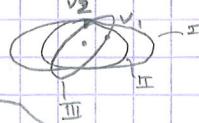
$$: S = \{P \in S_{++}^n \mid v_i^T P^{-1} v_i \leq 1, i=1 \dots n\} \text{ for } v_1 \dots v_n \in \mathbb{R}^n \text{ fixed.}$$

Remark: This corresponds to the set of ellipsoids that contain the points  $v_1 \dots v_n$

Claim: The set  $S$  has no minimum element

Justify:  $\because \exists$  intersecting ellipsoids  $\Rightarrow \exists$  ellipsoids that are not comparable.  
A minimum element must be comparable (smaller) than all elements of a set.

Illustration: Minimal element.



I is not minimal:  $\exists$  a "smaller" ellipsoid containing  $v_1, v_2$ .  
II for eg. similarly for III.  
III is a minimal element.  $\nexists$  (usually) an ellipsoid contained in III (< than III) that enclosed  $v_1 \& v_2$ .

## § 2.5 Separating and Supporting Hyperplanes

### § 2.5.1 Separating hyperplane theorem

Thm: For convex sets  $C \& D$ , s.t.  $C \cap D = \emptyset$ ,  $\exists a \neq 0 \& b$  s.t.  $a^T x \leq b \forall x \in C \& a^T x \geq b \forall x \in D$ .

: For (same)  $C \& D$ , (the affine-fn)  $a^T x - b$  is nonpositive on  $C$  & nonnegative on  $D$ .

Def<sup>n</sup> (Impl): Separating Hyperplane := The hyperplane  $\{x \mid a^T x = b\}$

Proof: (special case)

Recall: Euclidean Distance :=  $\text{dist}(C, D) := \inf \{ \|u - v\|_2 \mid u \in C, v \in D\}$

Assume: Euclidean distance is positive

(b)  $\exists$  points  $c \in C$   $d \in D$  s.t. the minimum is achieved,  $\|c - d\|_2 = \text{dist}(C, D)$ .

Def<sup>n</sup>:  $a := d - c$ ,  $b := \|d\|_2^2 - \|c\|_2^2$

objectiv: Show that the affine fn  $f(x) = a^T x - b = (d - c)^T [x - (\frac{c+d}{2})] (d+c)$  #6  
is non-positive on  $C$  & non-negative on  $D$ .

remark: This is a hyperplane perpendicular to the line segment between  $c \& d$ , passing through its midpoint.

"proof start":  $f(u) = (d - c)^T (u - d + (\frac{c+d}{2})(d - c)) = (d - c)^T (u - d) + \frac{1}{2} \|d - c\|_2^2$

Dis non-neg

Suppose  $u \in D$  was s.t.  $f(u) < 0$  (on the contrary).

$$\Rightarrow (d - c)^T (u - d) < 0.$$

$$\text{NB: } \frac{\partial}{\partial t} \|d + t(u - d) - c\|_2^2 \Big|_{t=0} = 2(d - c)^T (u - d) < 0$$

#6

$\Rightarrow$  for small  $t > 0$  we must have  $\|d + t(u-d) - c\|_2 < \|d - c\|_2$   
viz.  $d + t(u-d)$  is closer to  $c$  than  $d$  is.

$$tu + (1-t)d$$

Now this point is a convex combination of points  $u \& d \Rightarrow$  must be in  $D$ . But that's a contradiction since  $d$  was assumed closest to  $c$  in set  $D$ .  $\square$

Assume:

E.g. Separation of an affine  $\&$  a convex set: (1)  $C$  is convex, say, &  $D$  is affine,  $= \{Fu+g \mid u \in \mathbb{R}^m\}$   
where  $F \in \mathbb{R}^{m \times m}$ . (2)  $C \& D$  are disjoint

Concl': By the hyperplane theorem  $\exists$  an  $a \neq 0 \& b$  s.t.  $a^T x \leq b \Leftrightarrow x \in C$   
 $a^T x > b \Leftrightarrow x \in D$

Explicit: For  $x \in D$ ,  $x = Fu + g \Rightarrow a^T Fu > b - a^T g \Leftrightarrow u \in \mathbb{R}^m$ .  
(Plane constn) A linear  $f'$  is bounded below only when it's zero (in  $\mathbb{R}^m$ ).  $\square \#7$

$$\Rightarrow a^T F = 0 \Rightarrow a^T g > b \Rightarrow b \leq a^T g$$

We obtain a from  $a^T F = 0$  & use  $b = a^T g$ . This specifies one possible plane,  $a^T x = b$ .

### Def: Strict separation of $C \& D$

$\Leftrightarrow$  (For the aforesaid hyperplane)  $a^T x < b \Leftrightarrow x \in C \& a^T x > b \Leftrightarrow x \in D$

E.g. Strict separation of a point & a closed convex set. [Claim implies]: "Proof" that a closed set

Consider: (1)  $C$  is a closed convex set  
(2) a point  $x_0 \notin C$ .

is the intersection of all halfspaces that contain it.

Claim:  $\exists$  a hyperplane, strictly separating  $x_0$  from  $C$ .

Assume:  $C$  is convex closed

Proof: consider a set  $B := \{x + u \mid \|u\|_2 \leq \epsilon\}$  for some  $\epsilon > 0$ .

$S$  is the set intersectn

of halfspaces.  
by thm:  $\exists a \neq 0 \& b$  s.t.  $a^T x \leq b \quad x \in C$

$$a^T x > b \quad x \in B$$

$\square \#8$

N.B: If  $x \in C$  obviously

$x \in S$ .

This means  $a^T(x_0 + u) \geq b + \epsilon \|u\|_2$   $\forall \|u\|_2 \leq \epsilon$ .

We show:  $x \in S \Rightarrow x \in C$ .  
assume:  $x \in S$ , but

$x \notin C$ . By strict separation,  $\exists$  a

hyperplane  $b/w x \& C$ .

$\Rightarrow \exists$  a halfspace containing  $C$  but not  $x$ .  $\Rightarrow x \notin S$ , contradiction!

Now  $a^T x \leq b$  for  $x \in C \Rightarrow a^T x < b + \frac{\epsilon \|u\|_2}{2} = b'$

So we have strict separation with  $a^T x = b' (= b + \frac{\epsilon \|u\|_2}{2})$

### Converse separating hyperplane thms

False statement: If  $\exists$  a hyperplane (separating)  $\Rightarrow C \cap D = \emptyset$ .

Counter e.g.:  $C = D = \{0\} \subset \mathbb{R}^3$ . The  $x=0$  separates  $C \& D$  but  $C \cap D = \{0\} \neq \emptyset$ .

(Note, separates  $\neq$  strictly separated)

Converse thm: If (1)  $C \& D$  are convex with  $C$  open  $\Rightarrow C \cap D = \emptyset$ , [e.g. Theorem of alternatives]

(2)  $\exists$  a separating hyperplane  $b/w C \& D$

[for strict linear inequalities]

$\exists$  an affine  $f' : f' \leq 0$  in  $C$  &  $\geq 0$  in  $D$

target:

"Proof":

combined thm: Consider convex sets  $C \& D$ , with at least one open, are disjoint  $\Leftrightarrow \exists$  a separating hyperplane.

Find necessary & sufficient conditions for solvability of

$$Ax \leq b$$

$\Rightarrow$  N.B:  $f$  must be  $< 0$  in  $C$ .  $\because$  If  $f=0$  at some point, then  $f$  would become positive somewhere near that point (along some dir, since it's like a linear  $f'$ , the sign must change if  $f=0$  somewhere, no direction is ruled out really because of openness of  $C$  (roughly speaking)). which is a contradiction.

This means  $C \& D$  must be disjoint " $f$  is negative on  $C$  & non-positive on  $D$ ".

NB: The inequality is unfeasible  $\Leftrightarrow$  the (convex) sets  $C = \{b - Ax \mid x \in \mathbb{R}^n\}$  &  $D = \mathbb{R}_{++}^m = \{y \in \mathbb{R}^m \mid y \geq 0\}$  do not intersect.

NB2: D is open, C is affine

Concl': From the thm, C & D are disjoint  $\Leftrightarrow \exists$  a separating hyperplane, i.e.

$$\exists a \in \mathbb{R}^m \text{ & } \mu \in \mathbb{R} \text{ s.t. } x^T y \leq \mu \text{ on } C$$

Simplif': The first yields  $x^T(b - Ax) \leq \mu + x \Rightarrow Ax^T = 0$  (else you can't have  $x^T y \geq \mu$  on D)

&  $x^T b \leq \mu$  bound a linear fn from below

The second yields  $x^T y \geq \mu + y^T 0 \Rightarrow \mu \leq 0 \text{ & } x^T 0, \lambda \geq 0, \lambda \neq 0$ .  $\square \#9$   
(rough)

Putting them together we have

Final Result: The inequality is unfeasible  $\Leftrightarrow \exists a \in \mathbb{R}^m \text{ s.t. } \lambda \geq 0, \lambda \neq 0, A^T \lambda = 0, \lambda^T b \leq 0$

Remark: This is also a system of linear eq's in  $\lambda$ .

: For a given A & b, exactly one of these sets of equations is solvable.

Nomencl: These form a pair of alternative eq's.

## § 2.5.2 Supporting hyperplanes

Def': Boundary := for  $C \subseteq \mathbb{R}^n$ ,  $\partial C := \text{cl } C \setminus \text{int } C$ .

(Imp) Def': Supporting hyperplane := for  $x_0 \in \partial C$ , for  $a \neq 0$ , if  $a^T x \leq a^T x_0 \forall x \in C$  then the hyperplane  $\{x \mid a^T x = a^T x_0\}$  is a supporting hyperplane.

Remark: Geometrically, it's a tangent to the set.

Thm: Supporting Hyperplane Theorem: For a non-empty convex set C & any point  $x_0 \in \partial C$ ,  $\exists$  a supporting plane at  $x_0$

Proof: (1) If  $\text{int}(C)$  is non-empty, then result follows from the hyperplane thm applied to  $\text{int}(C) \setminus \{x_0\}$ .

(2) If  $\text{int}(C)$  is empty  $\Rightarrow C$  has an affine dimension less than n. Any hyperplane containing this affine set, contains C &  $x_0$ . This then is a (trivial) supporting hyperplane.

Inversethm: If a set, C, is closed,  $\text{int}(C) \neq \emptyset$  & it has a supporting hyperplane  $\partial C$  then it's convex.

## § 2.6 Dual cones and generalized inequalities.

19 Oct 2016

### § 2.6.1 Dual cones



Def': Dual Cone :=  $K^* = \{y \mid x^T y \geq 0 \forall x \in K\}$  where K is a cone.

Claim:  $K^*$  is a cone & is always convex (even when K is not).

NB: Geometrically  $K^* = \{y \mid -y \text{ is the normal of a hyperplane that supports } K \text{ at the origin}\}$

e.g. (1) Subspace: dual cone of a subspace  $V \subseteq \mathbb{R}^n$  is its orthnormal complement  $V^\perp = \{y \mid v^T y = 0 \forall v \in V\}$

(2) Non-negative orthant: (should help both geometrically & algebraically) the cone  $\mathbb{R}_+^n$  is its own dual:  $x^T y \geq 0 \Leftrightarrow y \geq 0$ .

(3) Positive semidefinite cone: Consider  $S^n$  & use  $u(x,y)$  inner product.

claim:  $S^n$  is self-dual, i.e. for  $X, Y \in S^n$ ,  $u(X, Y) \geq 0 \Leftrightarrow X \geq 0 \Leftrightarrow Y \geq 0$

Proof: Suppose  $Y \notin S^n$ .  $\Rightarrow \exists q \in \mathbb{R}^n$  s.t.  $q^T Y q = u(q, Y) < 0$ .  
 $Y \notin S^n \Leftrightarrow \square \#10$

Suppose  $X, Y \in S^n$ . We can write  $X = \sum \lambda_i q_i q_i^T$  with  $\lambda_i \geq 0$

$$\Rightarrow u(Y, X) = \sum \lambda_i (Y^T q_i q_i^T) = \sum \lambda_i q_i^T Y q_i \geq 0$$

$\therefore Y \geq 0 \Leftrightarrow \lambda_i \geq 0$

(4) Dual of a norm cone:

Assume:  $\|\cdot\|$  is a norm on  $\mathbb{R}^n$

Recall: A cone  $K := \{(x, t) \in \mathbb{R}^{n+1} \mid \|x\| \leq t\}$

Claim: Dual cone  $K^* = \{(u, v) \in \mathbb{R}^{n+1} \mid \|u\|_\infty \leq v\}$  where  $\|u\|_\infty = \sup \{u^T x \mid \|x\| \leq 1\}$

Proof Idea: Show that

$x^T u + t v \geq 0$  whenever  $\|x\| \leq t \Leftrightarrow \|x\|_\infty \leq v$

(from the Def' of a cone, we have  $(x^T u + t v) \geq 0$  for all points in the domain)

Proof: Start with LHS  $\geq$  RHS; Let  $\|u\|_\infty \leq v$ , &  $\|x\| \leq t$  for a  $t > 0$  (for  $t=0$ ,  $x=0$ )  
 $\Rightarrow u^T x + v^T t \geq 0$

Recall: Def<sup>\*</sup> of dual norm: =  $\sup\{u^T x \mid \|x\| \leq 1\} = \|u\|_*$   
 Using the def<sup>\*</sup> of dual norm & that  $\left\|\frac{-x}{t}\right\| \leq 1$ , we have  
 $u^T \left(\frac{-x}{t}\right) \leq \|u\|_* \leq v \Rightarrow u^T x + vt \geq 0$

-11-

Next LHS  $\Rightarrow$  RHS; Let  $\|u\|_* > v$  (i.e. RHS doesn't hold).  
 Now from def<sup>\*</sup> of the dual norm  $\exists$  an  $x$  with  $\|x\| \leq 1$   
 s.t.  $x^T u > v$ . For  $t=1$ , we have  $u^T(-x) + v < 0$   
 which violates the LHS (for  $t=1$ ). !RHS  $\Rightarrow$  !LHS  
 $LHS \Rightarrow RHS$ .

Claim: Properties of the dual cone (a)  $K^*$  is closed & convex (b)  $K_1 \subseteq K_2 \Rightarrow K_2^* \subseteq K_1^*$   
 (c) If  $K$  has a non-empty interior, then  $K^*$  is pointed  
 (d) If closure of  $K$  is pointed then  $K^*$  has non-empty interior  
 (e)  $K^{**}$  is the closure of the convex hull of  $K$ .  
 (Hence if  $K$  is convex & closed,  $K^{**} = K$ )

~ TODO: Ex 2.31

## § 2.6.2 Dual Generalised Inequalities

Assume:  $K$  is a proper convex cone.

NB:  $g_t$  induces a generalized inequality  $\leq_K$

NB: The dual cone  $K^*$  is also proper & convex.

: It induces a generalized inequality  $\leq_{K^*}$

Def<sup>\*</sup>: Generalized inequality  $\leq_{K^*}$  is the "dual of"  
 the generalized inequality  $\leq_K$

Properties: (a)  $x \leq_K y \Leftrightarrow x^T z \leq y^T z \forall z \in K$   
 (b)  $x \leq_{K^*} y \Leftrightarrow x^T z \leq y^T z + \lambda \sum_{z \in K^*} 0, \forall \lambda \geq 0$

NB: Since  $K^{**} = K$ , dual associated with  
 $\leq_{K^*}$  is  $\leq_K$ ;  $\Rightarrow$  properties hold with  
 $K^*$  &  $K$  swapped.

E.g. 2.26 Theorem of alternatives for linear strict generalized inequalities

Assume:  $K \subseteq \mathbb{R}^m$  is a proper cone.