Seminar 10 week 10 (2 December 2019 – 6 December 2019)

Toy language syntax: Type ::= int bool string | Ref Type void **Stmt ::= Stmt;Stmt** | **Id= Exp** | Type Id | print(Exp) If Exp Then Stmt1 else Stm2 | openRFile(Exp) | readFile(Exp, id) | closeRFile(Exp) | new(Id, Exp) | wH(Id, Exp) | while Exp Stmt | fork(Stmt) **Value ::= Number** | True |False String | (value, Type) //ref value **Exp** ::= Value | id | rH(Exp) | Exp1 + Exp2 || Exp1 - Exp2 | Exp1 * Exp2 | Exp1 / Exp2 | Exp1 and Exp2 | Exp1 or Exp2 | Exp1 < Exp2 | $| Exp1 \le Exp2$ | Exp1 == Exp2| Exp1 != Exp2 | Exp1 > Exp2

| Exp1 >= Exp2

<u>Types</u>
2: int id: type (given by the programmer)
Type rules
1: int 2:int
1 + 2 : int
is reading as: 1+2 has type int IF 1 has type int and 2 has type int or IF 1 has type int and 2 has type int THEN 1+2 has type int
1: int v:??
1 + v : int
G- type environment, defined as a list of pairs (id:type)
G -1: int G -v:int
where G=[v:int] G -1 + v : int
Type rules for Values
G - Number : int
G -True:bool
G -False:bool
G -String:string
G - val: int
G -(val,type) : Ref type

Type rules for Expressions

(id:t) is in G

G|- id:t

G|- e1: int G|-e2:int

G|-e1 + e2:int

the same rule for -,*,/

G|- e1: bool G|-e2:bool

G|-e1 and e2:bool the same rule for or

G|-e1: int G|-e2:int

G|-e1 < e2:bool

the same rule for <=,==,!=,>, >=

G|- e1 :Ref t1

G|- rH(e1): t1

Type rules for Statements

G|- s1:void,G1 G1|- s2:void,G2

G|- s1;s2: void,G2

G|- id:t G|- exp:t

G|-id=exp: void,G

G|- type id : void, G+[(id:type)]

G - exp:t
G - print(exp): void,G
G - e : bool G - s1:void,G1
G - s2:void,G2
G - if e then s1 else s2 : void,G
G - nop:void, G
G - exp:string G -id:int
G - readFile(exp,id):void,G
G - exp:string
G - openRFile(exp):void, G
G - exp:string
G - closeRFile(exp):void, G
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G - exp:t G - id:Ref t
G - new(id,exp):void, G
G - exp:t G - id:Ref t
G - wH(id,exp):void, G
G - exp:bool G - stmt:void,G1
G - while exp stmt:void, G
G - stmt:void,G1
G - fork(stmt):void, G