



Objective:

At the end of this lecture, students will understand:

- Time Complexity
- Space Complexity
- Importance of them in Competitive Programming





Time and Space Complexity:

Sometimes, there are more than one way to solve a problem. We need to learn how to compare the performance different algorithms and choose the best one to solve a particular problem. While analysing an algorithm, we mostly consider time complexity and space complexity.





Time and Space Complexity: (Contd.)

- Time complexity of an algorithm quantifies the amount of time taken by an algorithm to run as a function of the length of the input.
- Similarly, Space complexity of an algorithm quantifies the amount of space or memory taken by an algorithm to run as a function of the length of the input.
- Time and space complexity depends on lots of things like hardware, operating system, processors, etc.





Order of growth:

- Order of growth of an algorithm is a way of saying /predicting how execution time of a program and the space/memory occupied by it changes with the input size.
- The most famous way is the Big-Oh notation. It gives the worst case possibility for an algorithm.





• E.g. a program takes O(n) time means it takes at most n operations to compute the answer. Now if your input for n is 10⁶, it takes 10⁶ operations, while a O(n²) algorithm takes 10¹² operations for the same input.

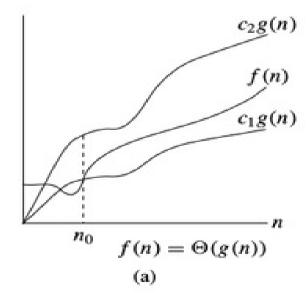


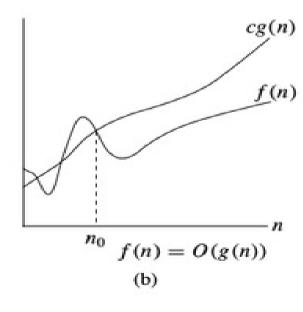


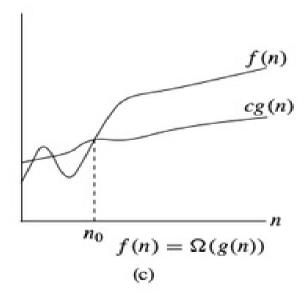
O-notation:

To denote asymptotic upper bound, we use O-notation. For a given function g(n), we denote by O(g(n)) (pronounced "big-oh of g of n") the set of functions:

 $O(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq c * g(n) \text{ for all } n \geq n_0 \}$







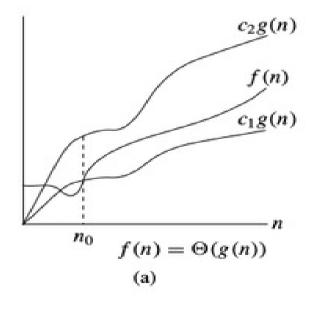


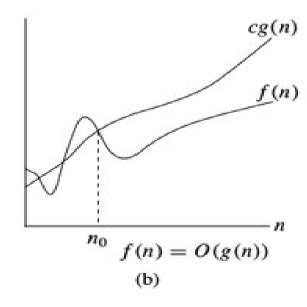


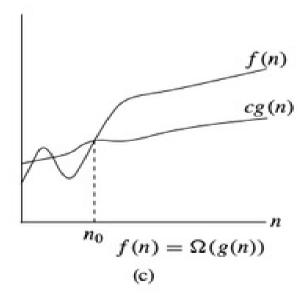
Ω -notation:

To denote asymptotic lower bound, we use Ω -notation. For a given function g(n), we denote by $\Omega(g(n))$ (pronounced "big-omega of g of n") the set of functions:

 $\Omega(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq c * g(n) \leq f(n) \text{ for all } n \geq n_0 \}$







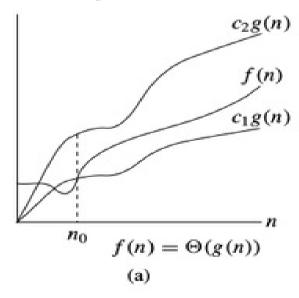


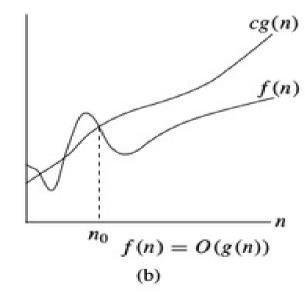


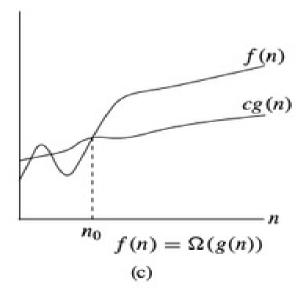
Θ -notation:

To denote asymptotic tight bound, we use Θ -notation. For a given function g(n), we denote by $\Theta(g(n))$ (pronounced "big-theta of g of n") the set of functions:

 $\Theta(g(n)) = \{ f(n) : \text{there exist positive constants } c_1, \ c_2 \text{ and } n_0 \text{ such that } 0 \le c_1 * g(n) \le f(n) \le c_2 * g(n) \text{ for all } n > n_0 \}$





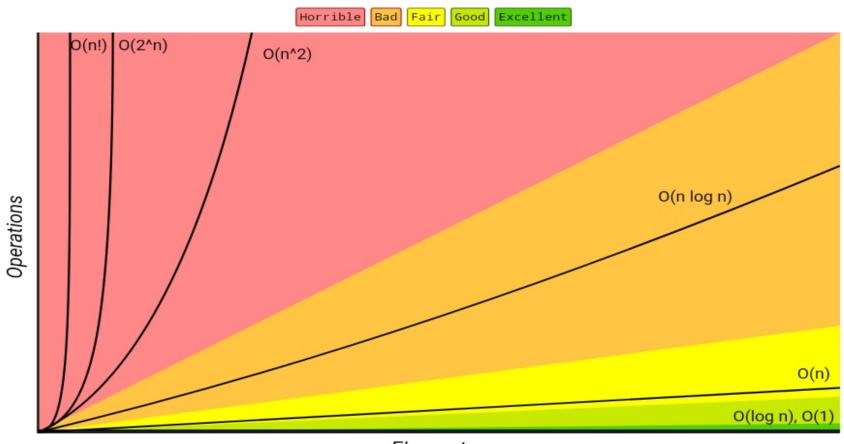






N vs. Growth of Complexity:

Big-O Complexity Chart



Elements





Complexity O(1):

Time complexity of a function (or set of statements) is considered as O(1) if it doesn't contain loop, recursion and call to any other non-constant time function.

- For example swap() function has O(1) time complexity.
- A loop or recursion that runs a constant number of times is also considered as O(1). E.g. the following loop is O(1).

```
for (int i = 1; i <= c; i++) { // Here c is a constant
  // some O(1) expressions
}</pre>
```





Complexity O(n):

Time Complexity of a loop is considered as O(n) if the loop variables is incremented / decremented by a constant amount. For example following functions have O(n) time complexity.

```
for (int i = 1; i <= n; i += c) {  // Here c is a +ve integer constant
  // some O(1) expressions
}
for (int i = n; i > 0; i -= c) {
  // some O(1) expressions
}
```





Complexity $O(n^c)$:

Time complexity of nested loops is equal to the number of times the innermost statement is executed. For example the following sample loops have $O(n^2)$ time complexity

```
for (int i = 1; i <=n; i += c) {
    for (int i = n; i > 0; i -= c) {
      for (int j = 1; j <=n; j += c) {
            // some O(1) expressions
      }
    }
}</pre>
```

For example Selection sort and Insertion Sort have O(n²) time complexity.





Complexity O(logn):

Time Complexity of a loop is considered as O(Log_cn) if the loop variables is divided / multiplied by a constant amount.

```
for (int i = 1; i <=n; i *= c) {
    // some O(1) expressions
}
for (int i = n; i > 0; i /= c) {
    // some O(1) expressions
}
```

For example Binary Search has O(Log₂n) time complexity.





Complexity O(LogLogn):

Time Complexity of a loop is considered as O(LogLog_cn) if the loop variables is reduced / increased exponentially by a const amount.

```
for (int i = 2; i <=n; i = pow(i, c)) {  // c is a const greater than 1
  // some O(1) expressions
}
// Here fun is sqrt or cuberoot or any other constant root
for (int i = n; i > 1; i = fun(i)) {
  // some O(1) expressions
}
```





Time complexities of consecutive loops:

Time Complexity When there are consecutive loops, we calculate time complexity as sum of time complexities of individual loops.

```
for (int i = 1; i <=m; i += c) {
    // some O(1) expressions
}
for (int i = 1; i <=n; i += c) {
    // some O(1) expressions
}</pre>
```

```
Time complexity of the code is O(m) + O(n) which is O(m + n)
```

If m == n, the time complexity becomes O(2n) which is O(n).

Quiz!



```
for (int i = 0; i < n; i++) {
     for (int j = 0; j < m; j++) {
        cout<<"Hello"<<endl;
     }
     }
A. O(n)
B. O(m)
C. O(n+m)</pre>
```

D. O(nm)



```
for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++) {
        cout<<"Hello"<<endl;
    }
    }

A. O(n)
B. O(m)
C. O(n+m)
```



```
for (int j = 0; j < n; j++) {
    for (int i = 0; i < n; i++)
        cout << "First loop" << endl;
    }
for (int i = 0; i < n; i++) {
        cout << "Second loop" << endl;
    }
for (int i = 0; i < n; i++) {
    cout << "Third loop" << endl;
    }
}</pre>
```

- A. O(n)
- B. O(n^2)
- C. Càn't be determined
- D. None of these



```
for (int j = 0; j < n; j++) {
    for (int i = 0; i < n; i++)
        cout << "First loop" << endl;
    }
for (int i = 0; i < n; i++) {
        cout << "Second loop" << endl;
    }
for (int i = 0; i < n; i++) {
        cout << "Third loop" << endl;
    }

O(n)
O(n^2)
Can't be determined</pre>
```

D. None of these



```
while (var < n){
    cout << pie << endl;
    for (int j = 1; j < n; j += 2){
        sum += 1;
    }
    var *= 3;
}

A. O(n)
B. O(logn)
C. O(nlog<sub>3</sub>n)
D. O(n<sup>2</sup>)
```



```
while (var < n){</pre>
       cout << pie << endl;</pre>
       for (int j = 1; j < n; j += 2) \begin{cases} ln \text{ the outer loop we have} \\ var *= 3, \text{ therefore outer} \end{cases}
          sum += 1;
       var *= 3;
 A. O(n)
 B. O(logn)
 C. O(nlog<sub>3</sub>n)
```

loop will have log₃(n) complexity but the inner loop proceeds 1 to n, therefore ultimately the whole code will have complexity nlog₃n.



```
for (int i = 0; i < n; i += 3) {
      for (int j = 0; j < m; j += 2) {
          cout << "Hello" << endl;
      }
    }
A. O(n)
B. O(m)
C. O(n+m)</pre>
```

D. O(nm)



```
for (int i = 0; i < n; i += 3) {
    for (int j = 0; j < m; j += 2) {
        cout << "Hello" << endl;
    }
    }
A. O(n)
B. O(m)
C. O(n+m)</pre>
```



```
for (int i = n/2; i <= n; i++) {
    for (int j = 1; j < m; j *= 2) {
        cout<<"Hello"<<endl;
```

- A. O(nlog₂m) B. O(nm)
- C. $O(\log_2 n^* \log_2 m)$
- D. None of these



```
for (int i = n/2; i <= n; i++) {
    for (int j = 1; j < m; j *= 2) {
        cout<<"Hello"<<endl;
```

- ⟨⟨⟨⟨ O(nlog₂m⟩⟩ B. O(nm)
- C. $O(\log_2 n*\log_2 m)$
- D. None of these



```
int i = 1, j = 1;
for (; i < n; i++) {
while(j < i) {
        cout<<"Hello"<<endl;
```

- A. O(log₂n) B. O(n)
- C. O(nlog₂n)
- D. None of the above



```
int i = 1, j = 1;
for (; i < n; i++) {
    while(j < i) {
        cout<<"Hello"<<endl;
```

A. $O(log_2n)$ B. O(n)

C. O(nĺog₂n)

D. None of the above



```
for (int i = 1; i < n; i += i) {
      for (int j = 0; j < n; j++) {
          cout<<"Hello"<<endl;
      }
    }
A. O(logn)
B. O(nlogn)
C. O(n)</pre>
```



```
for (int i = 1; i < n; i += i) {
     for (int j = 0; j < n; j++) {
        cout<<"Hello"<<endl;
     }
     }

A. O(logn)
B. O(nlogn)
C. O(n)</pre>
```



```
for (int i = 1; i < n; i *= 3) {
     for (int j = 1; j < n; j += 2) {
        cout << "Hello" << endl;
     }
    }

A. O(log<sub>3</sub>n)
B. O(nlog<sub>3</sub>n)
C. O(n)
D. O(n<sup>2</sup>)
```



```
for (int i = 1; i < n; i *= 3) {
     for (int j = 1; j < n; j += 2) {
        cout << "Hello" << endl;
     }
   }

A. O(log<sub>3</sub>n)
5. O(nlog<sub>3</sub>n)
C. O(n)
```



```
int i = 1, j = 1;
for (; i < n; i++) {
while(j < i) {
 cout < "Hello" < endl;
```

- A. O(logn) B. O(n)
- C. O(nlogn)
- D. None of the above



```
int i = 1, j = 1;
for (; i < n; i++) {
while(j < i) {
 cout < "Hello" < endl;
```

- A O(logn) ら、O(n)
- C. O(nĺogn)
- D. None of the above



```
for (int i = 1; i < n; i += 2) {
    for (int j = 1; j < n; j *= 3) {
        cout << "Hello" << endl;
    }
    }

A. O(log<sub>3</sub>n)
B. O(nlog<sub>3</sub>n)
C. O(n)
D. O(n^2)
```



```
for (int i = 1; i < n; i += 2) {
    for (int j = 1; j < n; j *= 3) {
        cout << "Hello" << endl;
    }
    }

A. O(log<sub>3</sub>n)
B. O(nlog<sub>3</sub>n)
C. O(n)
```