

## Objective:

**At the end of this lecture, students will understand :**

- Time Complexity
- Space Complexity
- Importance of them in Competitive Programming

## Time and Space Complexity:

Sometimes, there are more than one way to solve a problem. We need to learn how to compare the performance of different algorithms and choose the best one to solve a particular problem. While analysing an algorithm, we mostly consider time complexity and space complexity.

## Time and Space Complexity: (Contd.)

- Time complexity of an algorithm quantifies the amount of time taken by an algorithm to run as a function of the length of the input.
- Similarly, Space complexity of an algorithm quantifies the amount of space or memory taken by an algorithm to run as a function of the length of the input.
- Time and space complexity depends on lots of things like hardware, operating system, processors, etc.

## Order of growth:

- Order of growth of an algorithm is a way of saying /predicting how execution time of a program and the space/memory occupied by it changes with the input size.
- The most famous way is the Big-Oh notation. It gives the worst case possibility for an algorithm.

## Order of growth: (Contd.)

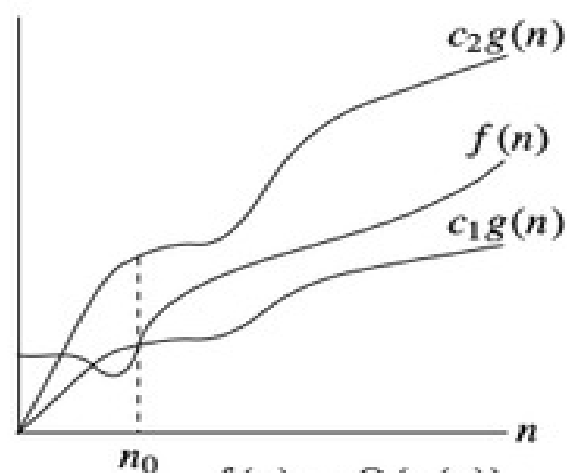
- E.g. a program takes  $O(n)$  time means it takes at most  $n$  operations to compute the answer. Now if your input for  $n$  is  $10^6$ , it takes  $10^6$  operations, while a  $O(n^2)$  algorithm takes  $10^{12}$  operations for the same input.

## Order of growth: (Contd.)

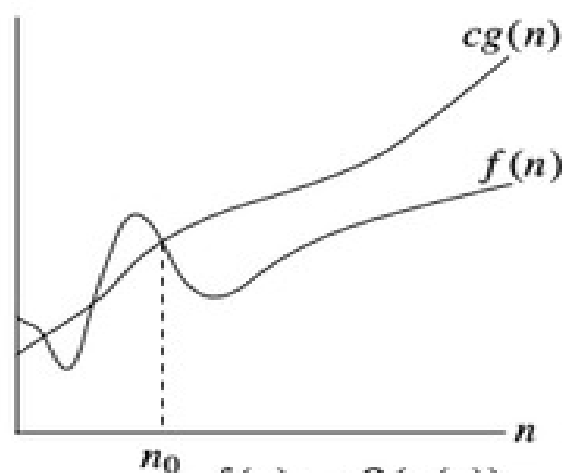
### *O*-notation:

To denote asymptotic upper bound, we use *O*-notation. For a given function  $g(n)$ , we denote by  $O(g(n))$  (pronounced “big-oh of  $g$  of  $n$ ”) the set of functions:

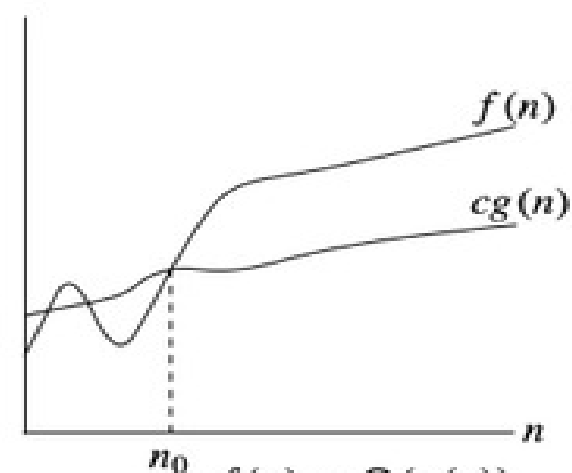
$O(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq c * g(n) \text{ for all } n \geq n_0 \}$



$f(n) = \Theta(g(n))$   
(a)



$f(n) = O(g(n))$   
(b)



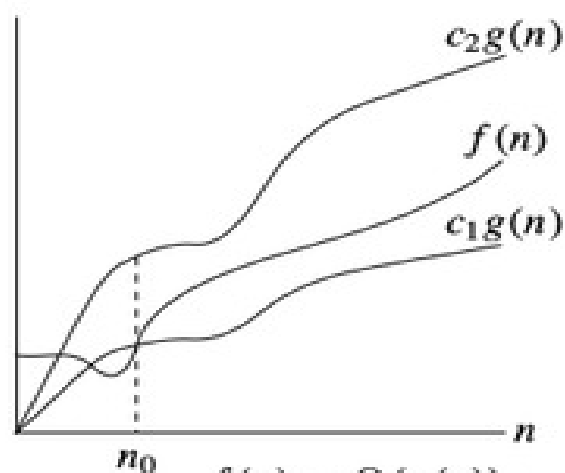
$f(n) = \Omega(g(n))$   
(c)

## Order of growth: (Contd.)

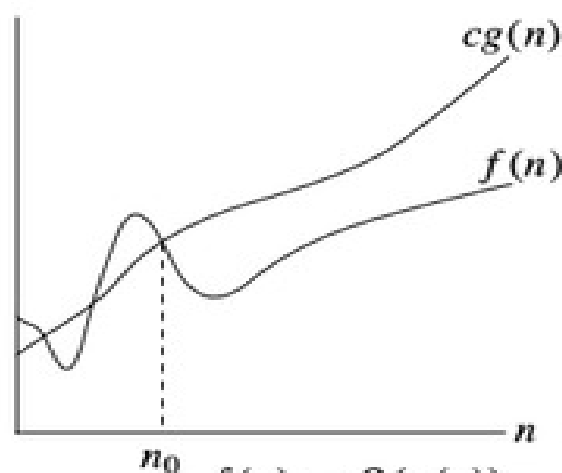
### $\Omega$ -notation:

To denote asymptotic lower bound, we use  $\Omega$ -notation. For a given function  $g(n)$ , we denote by  $\Omega(g(n))$  (pronounced “big-omega of g of n”) the set of functions:

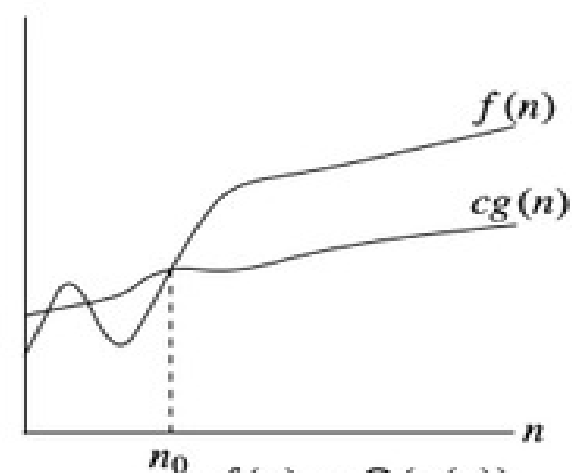
$\Omega(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq c * g(n) \leq f(n) \text{ for all } n \geq n_0 \}$



(a)



(b)



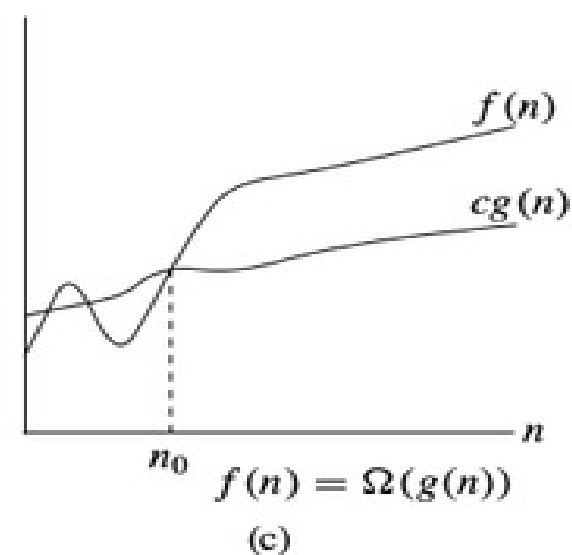
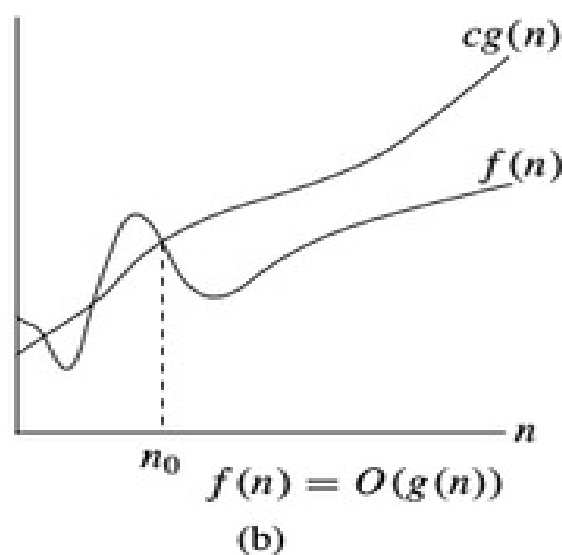
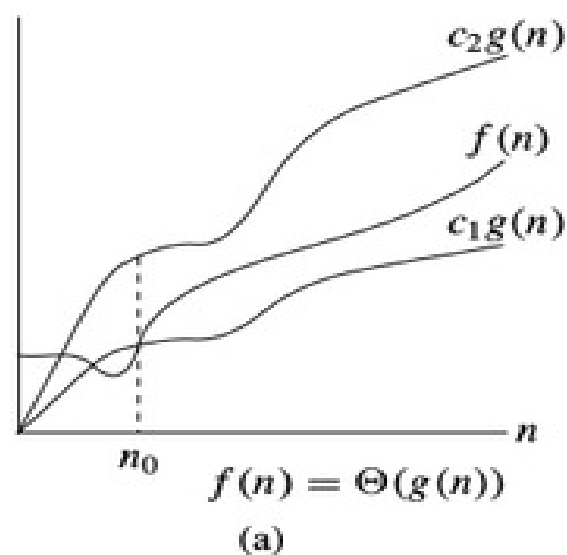
(c)

## Order of growth: (Contd.)

$\Theta$ -notation:

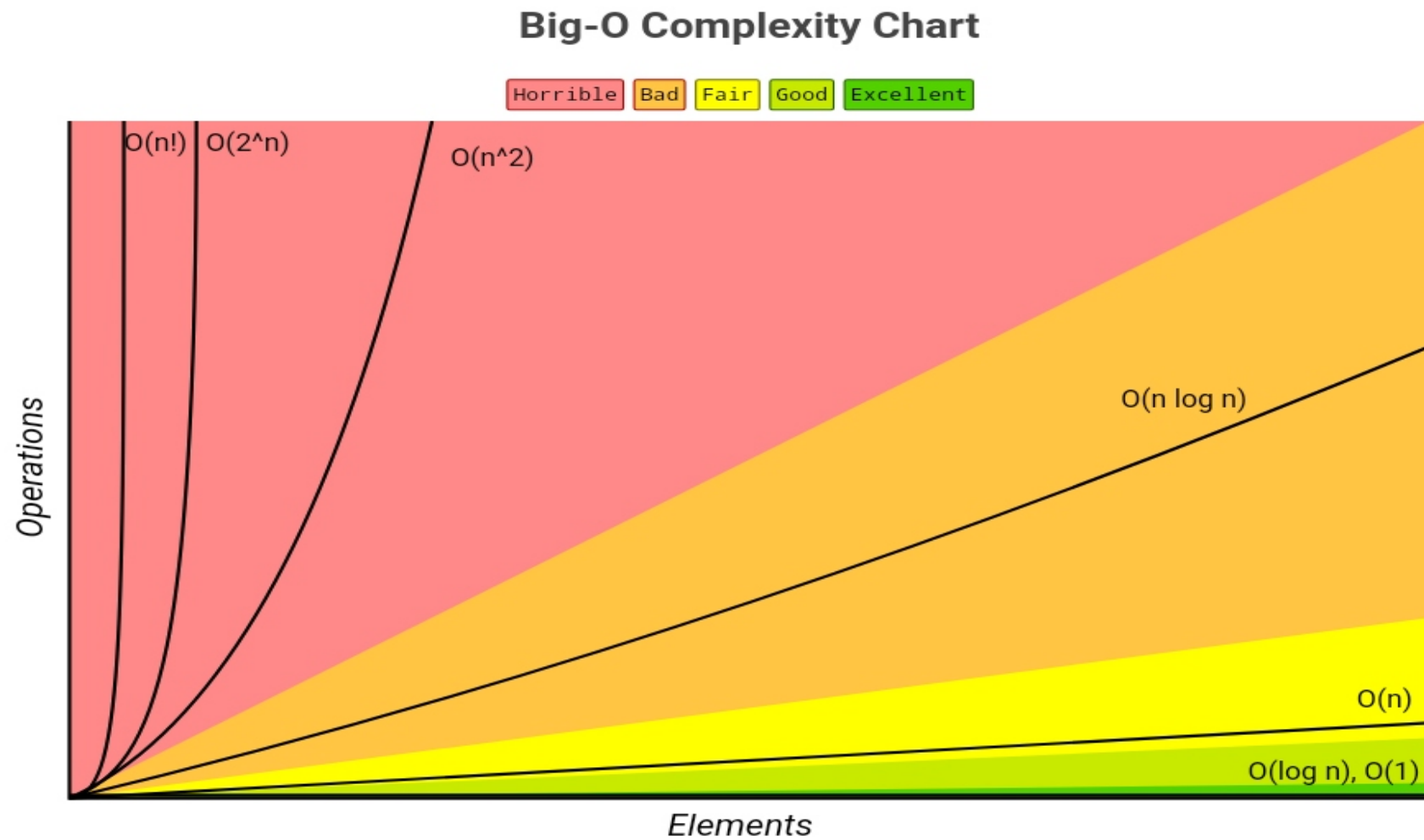
To denote asymptotic tight bound, we use  $\Theta$ -notation. For a given function  $g(n)$ , we denote by  $\Theta(g(n))$  (pronounced "big-theta of g of n") the set of functions:

$\Theta(g(n)) = \{ f(n) : \text{there exist positive constants } c_1, c_2 \text{ and } n_0 \text{ such that } 0 \leq c_1 * g(n) \leq f(n) \leq c_2 * g(n) \text{ for all } n > n_0 \}$





# N vs. Growth of Complexity:



## Complexity $O(1)$ :

Time complexity of a function (or set of statements) is considered as  $O(1)$  if it doesn't contain loop, recursion and call to any other non-constant time function.

- For example swap() function has  $O(1)$  time complexity.
- A loop or recursion that runs a constant number of times is also considered as  $O(1)$ . E.g. the following loop is  $O(1)$ .

```
for (int i = 1; i <= c; i++) {           // Here c is a constant  
    // some O(1) expressions  
}
```

## Complexity $O(n)$ :

Time Complexity of a loop is considered as  $O(n)$  if the loop variable is incremented / decremented by a constant amount. For example following functions have  $O(n)$  time complexity.

```
for (int i = 1; i <= n; i += c) { // Here c is a +ve integer constant  
    // some O(1) expressions  
}  
  
for (int i = n; i > 0; i -= c) {  
    // some O(1) expressions  
}
```

## Complexity $O(n^c)$ :

Time complexity of nested loops is equal to the number of times the innermost statement is executed. For example the following sample loops have  $O(n^2)$  time complexity

```
for (int i = 1; i <= n; i += c) {  
    for (int j = 1; j <= n; j += c) {  
        // some  $O(1)$  expressions  
    }  
}
```

```
for (int i = n; i > 0; i -= c) {  
    for (int j = i+1; j <= n; j += c) {  
        // some  $O(1)$  expressions  
    }  
}
```

**For example Selection sort and Insertion Sort have  $O(n^2)$  time complexity.**

## Complexity $O(\log n)$ :

Time Complexity of a loop is considered as  $O(\log_c n)$  if the loop variables is divided / multiplied by a constant amount.

```
for (int i = 1; i <= n; i *= c) {  
    // some  $O(1)$  expressions  
}
```

```
for (int i = n; i > 0; i /= c) {  
    // some  $O(1)$  expressions  
}
```

**For example Binary Search has  $O(\log_2 n)$  time complexity.**

## Complexity $O(\text{LogLog}n)$ :

Time Complexity of a loop is considered as  $O(\text{LogLog}_c n)$  if the loop variables is reduced / increased exponentially by a const amount.

```
for (int i = 2; i <= n; i = pow(i, c)) { // c is a const greater than 1  
    // some O(1) expressions  
    }  
  
// Here fun is sqrt or cuberoot or any other constant root  
for (int i = n; i > 1; i = fun(i)) {  
    // some O(1) expressions  
    }
```

## Time complexities of consecutive loops:

Time Complexity When there are consecutive loops, we calculate time complexity as sum of time complexities of individual loops.

```
for (int i = 1; i <= m; i += c) {  
    // some O(1) expressions  
}
```

```
for (int i = 1; i <= n; i += c) {  
    // some O(1) expressions  
}
```

**Time complexity of the code is  $O(m) + O(n)$  which is  $O(m + n)$**

**If  $m == n$ , the time complexity becomes  $O(2n)$  which is  $O(n)$ .**

**Quiz!**




1. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 0; i < n; i += 1) {  
    for (int j = 0; j < m; j += 1) {  
        print "Hello"  
    }  
}
```

- A.  $O(n)$
- B.  $O(m)$
- C.  $O(n+m)$
- D.  $O(nm)$

1. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 0; i < n; i += 1) {  
    for (int j = 0; j < m; j += 1) {  
        print "Hello"  
    }  
}
```

- A.  $O(n)$
- B.  $O(m)$
- C.  $O(n+m)$
-  D.  $O(nm)$


2. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int j = 0; j < n; j += 1) {  
    for (int i = 0; i < n; i += 1)  
        print "First loop"  
}  
for (int i = 0; i < n; i +=1) {  
    print "Second loop"  
}  
for (int i = 0; i < n; i += 1) {  
    print "Third loop"  
}
```

- A.  $O(n)$
- B.  $O(n^2)$
- C. Can't be determined
- D. None of these

2. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int j = 0; j < n; j += 1) {  
    for (int i = 0; i < n; i += 1)  
        print "First loop"  
}  
for (int i = 0; i < n; i +=1) {  
    print "Second loop"  
}  
for (int i = 0; i < n; i += 1) {  
    print "Third loop"  
}
```

- A.  $O(n)$
-  B.  $O(n^2)$
- C. Can't be determined
- D. None of these

3. What is the complexity of the following snippet ? ( n can be any integer )


```
while (var < n){  
    cout << pie << endl;  
    for (int j = 1; j < n; j+=2){  
        sum += 1  
    }  
    var *= 3  
}
```

- A.  $O(n)$
- B.  $O(\log n)$
- C.  $O(n \log_3 n)$
- D.  $O(n^2)$

3. What is the complexity of the following snippet ? ( n can be any integer )

```
while (var < n){  
    cout << pie << endl;  
    for (int j = 1; j < n; j+=2){  
        sum += 1  
    }  
    var *= 3  
}
```

In the outer loop we have  $\text{var} *= 3$ , therefore outer loop will have  $\log_3(n)$  complexity but the inner loop proceeds 1 to n, therefore ultimately the whole code will have complexity  $n \log_3 n$ .

- A.  $O(n)$
- B.  $O(\log n)$
-  C.  $O(n \log_3 n)$
- D.  $O(n^2)$


4. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 0; i < n; i += 3) {  
    for (int j = 0; j < m; j += 2) {  
        print "Hello"  
    }  
}
```

- A.  $O(n)$
- B.  $O(m)$
- C.  $O(n+m)$
- D.  $O(nm)$

4. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 0; i < n; i += 3) {  
    for (int j = 0; j < m; j += 2) {  
        print "Hello"  
    }  
}
```

- A.  $O(n)$
- B.  $O(m)$
- C.  $O(n+m)$
-  D.  $O(nm)$




5. What is the complexity of the following snippet ( n can be any integer )?

```
for (int i = n/2; i <= n; i += 1) {  
    for (int j = 1; j < m; j *= 2) {  
        print "Hello"  
    }  
}
```

- A.  $O(n \log_2 m)$
- B.  $O(nm)$
- C.  $O(\log_2 n * \log_2 m)$
- D. None of these

5. What is the complexity of the following snippet ( n can be any integer )?

```
for (int i = n/2; i <= n; i += 1) {  
    for (int j = 1; j < m; j *= 2) {  
        print "Hello"  
    }  
}
```

-  A.  $O(n \log_2 m)$
- B.  $O(nm)$
- C.  $O(\log_2 n * \log_2 m)$
- D. None of these


6. What is the complexity of the following snippet ? ( n can be any integer )

```
int i = 1, j = 1
for (; i < n; i += 1) {
    while(j < i) {
        print "Hello"
        j *= 2
    }
    j = 1
}
```

- A.  $O(\log_2 n)$
- B.  $O(n)$
- C.  $O(n \log_2 n)$
- D. None of the above

6. What is the complexity of the following snippet ? ( n can be any integer )

```
int i = 1, j = 1
for (; i < n; i += 1) {
    while(j < i) {
        print "Hello"
        j *= 2
    }
    j = 1
}
```

- A.  $O(\log_2 n)$
- B.  $O(n)$
-  C.  $O(n \log_2 n)$
- D. None of the above


7. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 1; i < n; i += i) {  
    for (int j = 0; j < n; j += 1) {  
        print "Hello"  
    }  
}
```

- A.  $O(\log n)$
- B.  $O(n \log n)$
- C.  $O(n)$
- D.  $O(n^2)$

7. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 1; i < n; i += i) {  
    for (int j = 0; j < n; j += 1) {  
        print "Hello"  
    }  
}
```

- A.  $O(\log n)$
-  B.  $O(n \log n)$
- C.  $O(n)$
- D.  $O(n^2)$


8. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 1; i < n; i *= 3) {  
    for (int j = 1; j < n; j += 2) {  
        print "Hello"  
    }  
}
```

- A.  $O(\log_3 n)$
- B.  $O(n \log_3 n)$
- C.  $O(n)$
- D.  $O(n^2)$

8. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 1; i < n; i *= 3) {  
    for (int j = 1; j < n; j += 2) {  
        print "Hello"  
    }  
}
```

- A.  $O(\log_3 n)$
-  B.  $O(n \log_3 n)$
- C.  $O(n)$
- D.  $O(n^2)$




9. What is the complexity of the following snippet ? ( n can be any integer )

```
int i = 1 , j = 1
for (; i < n; i +=1) {
    while(j < i) {
        print "Hello"
        j *= 2
    }
}
```

- A.  $O(\log n)$
- B.  $O(n)$
- C.  $O(n \log n)$
- D. None of the above

9. What is the complexity of the following snippet ? ( n can be any integer )

```
int i = 1 , j = 1
for (; i < n; i +=1) {
    while(j < i) {
        print "Hello"
        j *= 2
    }
}
```

- A.  $O(\log n)$
-  B.  $O(n)$
- C.  $O(n \log n)$
- D. None of the above


10. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 1; i < n; i += 2) {  
    for (int j = 1; j < n; j *= 3) {  
        print "Hello"  
    }  
}
```

- A.  $O(\log_3 n)$
- B.  $O(n \log_3 n)$
- C.  $O(n)$
- D.  $O(n^2)$

10. What is the complexity of the following snippet ? ( n can be any integer )

```
for (int i = 1; i < n; i += 2) {  
    for (int j = 1; j < n; j *= 3) {  
        print "Hello"  
    }  
}
```

- A.  $O(\log_3 n)$
-  B.  $O(n \log_3 n)$
- C.  $O(n)$
- D.  $O(n^2)$