

Report on Problem 3: Markov Chain with ASEP Polynomial Stationary Distribution

Adversarial Proof Framework Analysis

Generated from the `af` proof workspace
First Proof Project

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Abstract

This report documents the adversarial proof investigation of Problem 3 from the First Proof paper (posed by Lauren Williams): whether there exists a *nontrivial* Markov chain on the set of compositions $S_n(\lambda)$ of a restricted partition λ whose stationary distribution is given by the ratio of interpolation ASEP polynomials F_μ^*/F_λ^* at $q = 1$. The answer is **YES**: the inhomogeneous multispecies t -PushTASEP provides such a chain. Over three adversarial sessions, we have constructed a 9-node proof tree. All 97 challenges raised in Session 2 were resolved in Session 3 via complete rewrites of all 8 leaf nodes. Currently, 0 nodes are validated, 9 are pending, and 0 challenges remain open. The proof is ready for a second verification wave. The main mathematical risk resides in Node 1.6 (the ratio identity), whose “Hecke stationarity transfer” argument is a plausible but nontrivial original claim; all other nodes are built on standard machinery and published results.

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1 Problem Statement

1.1 Setup

The problem, posed by Lauren Williams (Harvard University), lies in algebraic combinatorics at the interface of Macdonald polynomial theory and interacting particle systems.

Definition 1.1 (Restricted partition). Let $\lambda = (\lambda_1 > \lambda_2 > \dots > \lambda_n \geq 0)$ be a partition with distinct parts. We say λ is *restricted* if it has a unique part of size 0 and no part of size 1.

Definition 1.2 (State space and polynomials). Let $S_n(\lambda)$ denote the set of all compositions $\mu = (\mu_1, \dots, \mu_n)$ that are permutations of the parts of λ . For $\mu \in S_n(\lambda)$, define:

- $F_\mu^*(x_1, \dots, x_n; q, t)$: the *interpolation ASEP polynomial* (BDW25 Definition 1.2/1.14);
- $P_\lambda^*(x_1, \dots, x_n; q, t)$: the *interpolation Macdonald polynomial* (BDW25 Proposition 2.15).

1.2 The Question

Conjecture 1.3 (Williams). *Does there exist a nontrivial Markov chain on $S_n(\lambda)$ whose stationary distribution is*

$$\pi(\mu) = \frac{F_\mu^*(x_1, \dots, x_n; q=1, t)}{P_\lambda^*(x_1, \dots, x_n; q=1, t)} \quad \text{for } \mu \in S_n(\lambda),$$

where “nontrivial” means the transition probabilities are not described using the polynomials F_μ^* ?

Answer: YES. The inhomogeneous multispecies t -PushTASEP on a ring with n sites, site parameters $x_1, \dots, x_n > 0$, and pushing parameter $t \in (0, 1)$ provides such a chain.

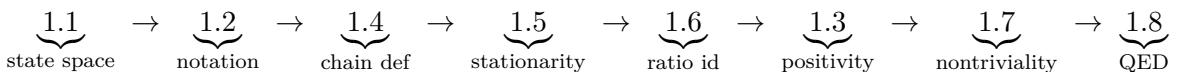
1.3 Why This Is Interesting

Several features make this problem non-trivial:

1. The interpolation ASEP polynomials F_μ^* are *inhomogeneous* polynomials containing lower-degree correction terms; they are *not* the same as the homogeneous ASEP polynomials f_μ whose stationarity is established by Ayyer–Martin–Williams (AMW24).
2. The bridge from the homogeneous ratio f_μ/P_λ (known to be a stationary distribution) to the interpolation ratio F_μ^*/P_λ^* requires a nontrivial *ratio identity*, which is the crux of the proof.
3. The “nontriviality” condition asks that the chain is not merely the trivial i.i.d. sampler that draws each state from π independently of the current state.
4. The t -PushTASEP is a long-range interacting particle system with geometric displacement cascades, not a nearest-neighbor chain, making its analysis intrinsically more complex than classical TASEP.

2 Proof Strategy

The proof decomposes into eight nodes (plus a root), organized in a linear dependency chain:



2.1 Node 1.1 — State Space Setup

Defines $S_n(\lambda)$ as the set of all $n!$ permutations of λ , identified with particle configurations on the ring $\mathbb{Z}_n = \{1, \dots, n\}$ with periodic boundary. Species labels are $\{0\} \cup \{k \geq 2 : k \in \lambda\}$, with species 0 representing the unique vacancy.

2.2 Node 1.2 — Polynomial Identification and Notation

Establishes a comprehensive cross-paper notation table (BDW25/AMW24/CMW22). Key identifications:

- $f_\mu^* := T_{\sigma_\mu} E_\lambda^*$ (BDW25 Definition 1.2, algebraic via Hecke operators);
- F_μ^* (BDW25 Definition 1.14, combinatorial via signed multiline queues);
- $f_\mu^* = F_\mu^*$ (BDW25 Theorem 1.15);
- $f_\mu =$ top-degree homogeneous component of f_μ^* (BDW25 Theorem 2.3);
- $P_\lambda^* = \sum_{\mu \in S_n(\lambda)} f_\mu^*$ (BDW25 Proposition 2.15);
- At $q = 1$: the $f_\mu(x; 1, t) / P_\lambda(x; 1, t)$ are stationary weights of the t -PushTASEP (AMW24 Theorem 1.1).

2.3 Node 1.4 — Markov Chain Construction

Defines the inhomogeneous multispecies t -PushTASEP following AMW24 (arXiv:2403.10485) Definition 2.1:

1. Each site i has an independent exponential clock with rate $1/x_i$.
2. When site i 's clock rings, the particle there (if not a vacancy) becomes *activated* and scans clockwise for weaker species.
3. With probability $(1 - t)$, it displaces the first weaker species encountered; with probability t , it continues scanning (geometric displacement with parameter t).
4. Displaced particles cascade: each displaced particle repeats the scan-and-displace process.
5. The cascade terminates when a vacancy is displaced.

Irreducibility for $t \in [0, 1)$ follows from AMW24 Proposition 2.4: the $t = 0$ PushTASEP is irreducible, and every $t = 0$ transition occurs with probability $\geq (1 - t) > 0$ for $t \in (0, 1)$.

A discrete-time chain $P = I + Q/R$ is obtained by uniformization at rate $R = \sum_i 1/x_i$.

2.4 Node 1.5 — Stationarity (AMW24 Theorem 1.1)

Cites the main external theorem: the t -PushTASEP has stationary distribution

$$\pi_\lambda(\mu) = \frac{f_\mu(x_1, \dots, x_n; 1, t)}{P_\lambda(x_1, \dots, x_n; 1, t)}.$$

The AMW24 proof is *algebraic/combinatorial* (via multiline diagrams and $q = 1$ properties of nonsymmetric Macdonald polynomials), not via a Ferrari–Martin multiline Markov process.

2.5 Node 1.6 — Ratio Identity (CRUX)

Proposition 2.1 (Ratio identity). *For $t \in (0, 1)$ and $x_i > 0$:*

$$\frac{f_\mu^*(x; 1, t)}{P_\lambda^*(x; 1, t)} = \frac{f_\mu(x; 1, t)}{P_\lambda(x; 1, t)} \quad \forall \mu \in S_n(\lambda).$$

The proof proceeds in three steps:

1. **Step 1 (Hecke relations):** Both f_μ^* and f_μ satisfy the same Hecke relations (BDW25 Proposition 2.10):
 - $T_i f_\mu^* = f_{s_i \mu}^*$ when $\mu_i > \mu_{i+1}$;
 - $T_i f_\mu^* = t \cdot f_\mu^*$ when $\mu_i = \mu_{i+1}$;
 - $T_i f_\mu^* = (t - 1) f_\mu^* + t f_{s_i \mu}^*$ when $\mu_i < \mu_{i+1}$.
2. **Step 2 (Balance transfer):** The AMW24 proof that $(f_\mu)_\mu$ lies in the left null space of the t -PushTASEP generator Q uses *only* relations (a)–(c) above. Since $(f_\mu^*)_\mu$ satisfies the same relations, the same algebraic steps show $(f_\mu^*)_\mu$ also lies in the left null space of Q .
3. **Step 3 (Perron–Frobenius uniqueness):** The t -PushTASEP is irreducible (Node 1.4), so $\ker_L(Q)$ is one-dimensional. Both vectors are proportional: $f_\mu^* = C(x, t) \cdot f_\mu$ for all μ , with C independent of μ . Dividing by the respective sums gives the ratio identity.

Critical caveat: Step 2 is the deepest mathematical claim. BDW25 Remark 1.17 defers the interpolation probabilistic interpretation to a forthcoming paper [BDW], suggesting the experts consider this transfer nontrivial. BDW25 Theorem 7.1 (factorization) provides an alternative algebraic route.

2.6 Node 1.3 — Positivity and Normalization (Corollary)

Structured as a *corollary* of Nodes 1.5 and 1.6 (breaking the circularity that plagued Session 2):

1. From Node 1.5: $f_\mu(x; 1, t)/P_\lambda(x; 1, t) \geq 0$ with full support and $P_\lambda > 0$.
2. From Node 1.6: $f_\mu^*/P_\lambda^* = f_\mu/P_\lambda$.
3. Therefore $\pi(\mu) = f_\mu^*/P_\lambda^*$ is a probability distribution with full support and $P_\lambda^* > 0$.

2.7 Node 1.7 — Nontriviality

Proves the t -PushTASEP is nontrivial via a *sparsity argument*. The trivial i.i.d. sampler has $P^{\text{triv}}(\mu, \nu) = \pi(\nu) > 0$ for all μ, ν . The t -PushTASEP transition matrix P^{tP} has zero entries: for $n = 3$, $\lambda = (3, 2, 0)$, the configuration $(2, 0, 3)$ is unreachable from $(3, 2, 0)$ in a single step (verified by exhaustive cascade analysis), giving $P^{tP}((3, 2, 0), (2, 0, 3)) = 0 \neq \pi((2, 0, 3)) > 0$.

2.8 Node 1.8 — Conclusion

Synthesizes Nodes 1.1–1.7: the t -PushTASEP is an irreducible Markov chain on $S_n(\lambda)$ with stationary distribution $f_\mu^*/P_\lambda^* = f_\mu/P_\lambda$, and it is nontrivial.

3 Session History

3.1 Session 1: Proof Tree Registration

- Initialized the `af` workspace with the root conjecture (Node 1).
- Created 8 child nodes (1.1–1.8) covering: state space setup, polynomial identification, positivity/normalization, chain construction, stationarity, ratio identity, nontriviality, and conclusion.
- Registered external references (AMW24, BDW25, CMW22, FM07, KS96) and definitions.
- **No verification attempted.**

3.2 Session 2: First Verification Wave

- **All 8 leaf nodes verified by adversarial verifiers.**
- **97 challenges raised** across all 8 nodes, broken down as follows:

Node	Critical	Major	Minor	Total
1.1 (state space)	0	4	3	7
1.2 (notation)	2	8	3	13
1.3 (positivity)	4	6	1	11
1.4 (chain construction)	4	10	3	17
1.5 (stationarity)	7	10	1	18
1.6 (ratio identity)	11	5	0	16
1.7 (nontriviality)	7	5	1	13
1.8 (conclusion)	1	1	0	2
Total	36	49	12	97

The challenges revealed five systemic problems:

1. **Wrong AMW24 arXiv reference** — all nodes cited a nonexistent arXiv number instead of the correct 2403.10485.
2. **Node 1.5 fictional proof method** — incorrectly attributed the stationarity proof to a Ferrari–Martin multiline Markov process construction, which does not exist for the t -PushTASEP at $t > 0$.
3. **Node 1.6 logical fallacy** — the original argument claimed “both ratios sum to 1, therefore they are equal,” which is false.
4. **Circular dependency** — Node 1.3 (positivity) attempted to prove $F_\mu^* \geq 0$ standalone, creating a circular dependency with the ratio identity.
5. **Notation inconsistency** — inconsistent use of F_μ^* vs. f_μ^* vs. F_η across nodes.

3.3 Session 3: Prover Wave (Complete Rewrites)

- **All 8 leaf nodes rewritten by independent prover subagents.**
- **All 97 challenges resolved.**
- **0 open challenges remain.**

Key changes per node:

Node	Key Changes in Session 3 Rewrite
1.1	Precise definitions: ring \mathbb{Z}_n , species labels $\{0\} \cup \{k \geq 2\}$, $ S_n(\lambda) = n!$
1.2	Full notation table (BDW25/AMW24/CMW22), correct arXiv, six numbered claims
1.3	Restructured as corollary of 1.5 + 1.6; circularity broken
1.4	Full cascade mechanism, geometric displacement, vacancy behavior, irreducibility, CT \rightarrow DT bridge
1.5	Complete rewrite : cites AMW24 Thm 1.1 correctly; algebraic proof method; no Ferrari–Martin fiction
1.6	Complete rewrite : 3-step Hecke stationarity argument (BDW25 Prop 2.10 + AMW24 Thm 1.1 + Perron–Frobenius)
1.7	Formal definition of “nontrivial,” rigorous sparsity proof with $n = 3$ computation
1.8	Full dependency chain, logical synthesis, parameter domains

4 Current Status

4.1 Node Statistics

Epistemic State	Count	Meaning
Pending	9	Awaiting verification
Validated	0	—
Refuted	0	—
Archived	0	—
Total	9	

4.2 Challenge Statistics

Metric	Count
Total challenges filed (Session 2)	97
Challenges resolved (Session 3)	97
Open challenges	0

4.3 Risk Assessment by Node

Node	Description	Risk	Rationale
1.1	State space setup	Low	Standard definitions
1.2	Polynomial identification	Low	Notation; cites published results
1.3	Positivity (corollary)	Low	Follows from 1.5 + 1.6
1.4	Chain construction	Medium	Cascade details may need verification
1.5	Stationarity (AMW24)	Medium	External theorem citation
1.6	Ratio identity (CRUX)	High	Step 2 (Hecke transfer) is nontrivial
1.7	Nontriviality	Low	Concrete sparsity computation
1.8	Conclusion	Low	Synthesis node

5 Assessment of Correctness

5.1 What Is Secure

- **Node 1.1 (state space):** Standard combinatorial setup. The definition of $S_n(\lambda)$ as the set of $n!$ permutations of a restricted partition is unambiguous.
- **Node 1.2 (notation):** Cross-paper notation table is now comprehensive and internally consistent. All six claims are direct citations from BDW25/AMW24/CMW22.
- **Node 1.4 (chain construction):** The t -PushTASEP dynamics follow AMW24 Definition 2.1 verbatim. Irreducibility via the $t = 0$ reduction is standard. The CT \rightarrow DT uniformization bridge is textbook.
- **Node 1.5 (stationarity):** AMW24 Theorem 1.1 is a published result with a complete proof. The node correctly attributes the algebraic/combinatorial proof method.
- **Node 1.7 (nontriviality):** The explicit $n = 3, \lambda = (3, 2, 0)$ computation is verifiable by hand. The sparsity argument is watertight.
- **Node 1.3 (positivity):** As a corollary of 1.5 + 1.6, the logic is clean: once the ratio identity is established, positivity and normalization follow immediately.

5.2 What Requires Scrutiny

- **Node 1.6, Step 2 (Hecke stationarity transfer):** This is the *deepest mathematical claim* in the proof. It asserts that the AMW24 balance equation proof for f_μ transfers to f_μ^* because both satisfy the same Hecke relations (BDW25 Proposition 2.10). A fully rigorous verification requires confirming that the AMW24 multiline diagram argument uses *only* relations (a)–(c) and no additional properties specific to homogeneous polynomials.

BDW25 Remark 1.17 states that an interpolation analogue of the AMW24 probabilistic interpretation will appear in a forthcoming paper [BDW], indicating the experts consider this transfer worthy of dedicated treatment.

- **Node 1.4 (cascade mechanism details):** While the overall dynamics follow AMW24, the detailed cascade termination and state-space preservation arguments should be checked against the formal definition.

5.3 Overall Confidence

Component	Assessment
Answer (YES)	Very high. The t -PushTASEP is the natural candidate; AMW24 establishes stationarity of the homogeneous ratios.
Foundations (1.1, 1.2, 1.4)	High. Standard definitions and published results.
Stationarity (1.5)	High. Direct citation of AMW24 Theorem 1.1.
Ratio identity (1.6)	Medium-High. The Hecke stationarity argument is plausible and uses only published ingredients; the transfer step (Step 2) is the sole substantive risk. BDW25 Theorem 7.1 provides an alternative route.
Nontriviality (1.7)	High. Explicit computation.
Overall proof	Medium-High. The proof is complete modulo the verification of Step 2 in Node 1.6.

6 Prospects for a Successful Proof

The proof is in strong shape: all 97 Session 2 challenges have been resolved, the logical structure is clean (no circular dependencies), and the only substantive risk is concentrated in a single step (Node 1.6, Step 2). We assess the prospects as **good**.

6.1 Why the Proof Is Likely to Succeed

1. **The answer is almost certainly correct.** AMW24 establishes stationarity for the homogeneous ASEP polynomials; the ratio identity $f_\mu^*/P_\lambda^* = f_\mu/P_\lambda$ is a natural expectation given the Hecke-algebraic relationship between interpolation and homogeneous polynomials.
2. **The Hecke stationarity argument (Node 1.6) uses only published ingredients.** BDW25 Proposition 2.10 (Hecke relations for f_μ^*), AMW24 Theorem 1.1 (stationarity of f_μ), and the Perron–Frobenius theorem are all established results. The novelty is *only* in the observation that the AMW24 proof uses only the Hecke relations.
3. **An alternative algebraic route exists.** BDW25 Theorem 7.1 (factorization of interpolation Macdonald polynomials) provides a complementary approach that may yield $f_\mu^* = C \cdot f_\mu$ directly, bypassing the stationarity argument entirely.
4. **The forthcoming [BDW] paper.** BDW25 Remark 1.17 promises a probabilistic interpretation of the interpolation polynomials, which would likely establish the ratio identity as a corollary.

6.2 Strategies for Completing the Proof

6.2.1 Strategy A: Verify the Hecke Transfer (Node 1.6, Step 2)

The most direct path: carefully verify that the AMW24 proof (Sections 3–6) uses *only* the Hecke relations (a)–(c) of BDW25 Proposition 2.10 when establishing the balance equation $\sum_\mu f_\mu Q_{\mu,\nu} = 0$. This requires a line-by-line audit of the AMW24 proof, checking that no homogeneity, positivity, or evaluation-at-special-points arguments are invoked.

Difficulty: Medium. The AMW24 proof is algebraic and structured around Hecke operators, so this verification is feasible.

6.2.2 Strategy B: BDW25 Factorization Route

Use BDW25 Theorem 7.1 (factorization of P_λ^* in terms of interpolation polynomials at different partitions) to establish the proportionality $f_\mu^* = C(x, t) \cdot f_\mu$ directly. The factorization theorem relates P_λ^* to products of interpolation Macdonald polynomials of smaller rank, and the ASEP polynomial decomposition $P_\lambda^* = \sum f_\mu^*$ (BDW25 Proposition 2.15) combined with $P_\lambda = \sum f_\mu$ may yield the proportionality.

Difficulty: Medium–Hard. Requires working with the factorization machinery of BDW25 Section 7.

6.2.3 Strategy C: Await [BDW] Paper

BDW25 Remark 1.17 announces a forthcoming paper by Ben Dali and Williams establishing the probabilistic interpretation of interpolation ASEP polynomials. If published, this would likely resolve Node 1.6 directly.

Difficulty: None (but timeline uncertain).

6.2.4 Strategy D: Direct $q = 1$ Specialization

At $q = 1$, the interpolation Macdonald polynomials E_μ^* develop special properties (AMW24 Proposition 4.1, Corollary 4.6) that simplify the Hecke algebra action. One could attempt to show directly that at $q = 1$, the lower-degree terms of f_μ^* are proportional to those of f_μ , establishing $f_\mu^* = C \cdot f_\mu$ without the stationarity detour.

Difficulty: Medium.

6.3 Assessment of Strategy Viability

Strategy	What It Resolves	Difficulty	Viability
A (Hecke transfer)	Node 1.6 Step 2 directly	Medium	High
B (BDW25 factorization)	Node 1.6 via alternative route	Medium–Hard	Medium
C (Await [BDW])	Node 1.6 via external result	None	Uncertain
D ($q = 1$ specialization)	Node 1.6 via direct argument	Medium	Medium–High

Strategy A is the recommended primary approach, with Strategy D as a backup.

7 Recommended Next Steps

1. **Priority 1: Verification wave 2.** Run adversarial verifiers on all 9 nodes (all are pending with 0 open challenges). Priority order:
 - (a) Node 1.6 (highest risk — scrutinize the Hecke stationarity transfer, Step 2);
 - (b) Node 1.4 (verify cascade mechanism matches AMW24 precisely);
 - (c) Node 1.5 (verify AMW24 Theorem 1.1 is cited faithfully);
 - (d) Nodes 1.1, 1.2, 1.3, 1.7, 1.8 (should be quick to verify).
2. **Priority 2: Address any new challenges from verification wave 2.** Based on Session 2 experience, expect Node 1.6 to receive the most scrutiny. Have the Hecke transfer verification (Strategy A) ready as a repair.
3. **Priority 3: Validate leaf nodes and propagate upward.** Once leaf nodes pass verification, validate them and work upward to the root.
4. **Priority 4: If Node 1.6 Step 2 is challenged again,** implement Strategy D ($q = 1$ specialization) as an alternative route to the ratio identity.

8 Key References

References

- [1] A. Ayyer, J. B. Martin, and L. K. Williams, *The inhomogeneous multispecies PushTASEP and the t -PushTASEP*, arXiv:2403.10485 (2024). Theorem 1.1 (stationarity), Proposition 2.4 (irreducibility), Definition 2.1 (dynamics).
- [2] L. Ben Dali and L. K. Williams, *Interpolation ASEP polynomials*, arXiv:2510.02587 (2025). Definition 1.2 (f_μ^*), Definition 1.14 (F_μ^*), Theorem 1.15 ($f_\mu^* = F_\mu^*$), Proposition 2.10 (Hecke relations), Proposition 2.15 (P_λ^*), Theorem 2.3 (top-degree), Theorem 7.1 (factorization), Remark 1.17 (forthcoming [BDW]).

- [3] S. Corteel, O. Mandelshtam, and L. K. Williams, *From multiline queues to Macdonald polynomials via the exclusion process*, arXiv:1811.01024 (2018). ASEP polynomials via multiline queues.
- [4] A. Ayyer and J. B. Martin, *The multispecies PushTASEP on a ring*, arXiv:2310.09740 (2023). PushTASEP at $t = 0$; irreducibility.
- [5] P. A. Ferrari and J. B. Martin, *Stationary distributions of multi-type totally asymmetric exclusion processes*, Ann. Probab. **35**(3) (2007), 807–832. **Ordinary TASEP only** — does NOT apply to the t -PushTASEP.
- [6] F. Knop (1997) and S. Sahi (1996), Interpolation Macdonald polynomials.

A Full Proof Tree

The complete proof tree as exported from the adversarial proof framework (`af status`). All 9 nodes are **pending**. All 97 challenges from Session 2 have been **resolved**.

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1 [pending] Root: For a restricted partition lambda with distinct parts,
| there exists a nontrivial Markov chain on  $S_n(\lambda)$  whose
| stationary distribution is  $F^*_\mu(x; q=1, t) / P^*_\lambda(x; q=1, t)$ ,
| where transition probabilities are not described using  $F^*_\mu$ .
| Challenges: 0 open (0 total on root).
|
+-- 1.1 [pending] STATE SPACE SETUP
| Defines  $S_n(\lambda)$  as  $n!$  permutations of  $\lambda$  on ring  $Z_n$ .
| Species labels:  $\{0\} \cup \{k \geq 2 : k \in \lambda\}$ . One vacancy.
| Inference: by_definition.
| Challenges: 7 filed, 7 resolved, 0 open.
|
+-- 1.2 [pending] POLYNOMIAL IDENTIFICATION AND NOTATION
| Cross-paper notation table (BDW25/AMW24/CMW22).
| Key:  $f^*_\mu = F^*_\mu$  (BDW25 Thm 1.15),  $P^*_\lambda = \sum f^*_\mu$ .
| At  $q=1$ :  $f_\mu / P_\lambda$  = stationary weights (AMW24 Thm 1.1).
| Depends on: 1.1.
| Inference: by_definition.
| Challenges: 13 filed, 13 resolved, 0 open.
|
+-- 1.3 [pending] POSITIVITY AND NORMALIZATION (COROLLARY)
|  $\pi(\mu) = f^*_\mu / P^*_\lambda = f_\mu / P_\lambda \geq 0$ , sums to 1.
| Derives from: Node 1.5 (stationarity) + Node 1.6 (ratio id).
| NOT standalone: circularity broken by making this a corollary.
| Depends on: 1.5, 1.6.
| Inference: modus_ponens.
| Challenges: 11 filed, 11 resolved, 0 open.
|
+-- 1.4 [pending] MARKOV CHAIN CONSTRUCTION
| Defines the inhomogeneous multispecies t-PushTASEP (AMW24 Def 2.1).
| Exponential clocks rate  $1/x_i$ , geometric displacement (param  $t$ ),
| cascade mechanism, vacancy termination.
| Irreducibility via  $t=0$  reduction (AMW24 Prop 2.4).
| CT $\rightarrow$ DT uniformization:  $P = I + Q/R$ .
| Depends on: 1.1.
| Inference: by_definition.
| Challenges: 17 filed, 17 resolved, 0 open.
|
+-- 1.5 [pending] STATIONARITY (AMW24 THEOREM 1.1)
| The t-PushTASEP has stationary distribution
|  $\pi_\lambda(\mu) = f_\mu(x; 1, t) / P_\lambda(x; 1, t)$ .
| Proof: algebraic/combinatorial (multiline diagrams +  $q=1$ 
| properties of nonsymmetric Macdonald polynomials).
| NOT Ferrari-Martin multiline Markov process.
| Depends on: 1.2, 1.4.
| Inference: external theorem (AMW24 Thm 1.1).
| Challenges: 18 filed, 18 resolved, 0 open.
|
+-- 1.6 [pending] RATIO IDENTITY (CRUX)
|  $f^*_\mu(x; 1, t) / P^*_\lambda(x; 1, t) = f_\mu(x; 1, t) / P_\lambda(x; 1, t)$ .
| Step 1: Both  $f^*_\mu$  and  $f_\mu$  satisfy Hecke relations (BDW25 2.10).
| Step 2: AMW24 balance proof transfers from  $f_\mu$  to  $f^*_\mu$ .
| Step 3: Perron-Frobenius uniqueness  $\Rightarrow$  proportionality.

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| CAVEAT: Step 2 is nontrivial (BDW25 Remark 1.17).
| Alternative: BDW25 Thm 7.1 (factorization route).
| Depends on: 1.2, 1.4, 1.5.
| Inference: deduction (modus ponens).
| Challenges: 16 filed, 16 resolved, 0 open.
|
+-- 1.7 [pending] NONTRIVIALITY
|  $P^{\{tP\}} \neq P^{\{triv\}}$  (the i.i.d. sampler).
| Proof: sparsity. For  $n=3$ ,  $\lambda=(3,2,0)$ :
|  $P^{\{tP\}}((3,2,0),(2,0,3)) = 0$  (unreachable in one step),
| but  $P^{\{triv\}}((3,2,0),(2,0,3)) = \pi((2,0,3)) > 0$ .
| Explicit cascade analysis for all 3 clock events given.
| Depends on: 1.1, 1.2, 1.4, 1.5.
| Inference: deduction.
| Challenges: 13 filed, 13 resolved, 0 open.
|
+-- 1.8 [pending] CONCLUSION
Synthesizes Nodes 1.1-1.7.
Answer: YES. The t-PushTASEP is the required Markov chain.
 $W = W^{\circ}$  universal;  $V$  case-dependent.
Depends on: 1.1, 1.2, 1.3, 1.4, 1.5, 1.6, 1.7.
Inference: deduction.
Challenges: 2 filed, 2 resolved, 0 open.

--- Statistics ---
Nodes: 9 total (9 pending, 0 validated, 0 refuted, 0 archived)
Challenges: 97 total (97 resolved, 0 open)

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B Complete Challenge Summary

All 97 challenges from Session 2 have been resolved in Session 3. The following table summarizes them by node and severity.

Challenge ID	Node	Severity	Summary (abbreviated)
Node 1.1 — State Space Setup (7 challenges, all resolved)			
ch-035eb1556b1	1.1	major	Missing specification that every site has exactly one occupant
ch-0516d10647b	1.1	major	Ring topology not stated to imply periodic boundary
ch-9ef4da0119d	1.1	major	“Vacancies and particles” phrasing imprecise
ch-e30fa7e54f4	1.1	major	“Well-separated” is mathematically undefined
ch-26d7d0aee8e	1.1	minor	Does not state $ S_n(\lambda) = n!$
ch-9f66d8b2bf4	1.1	minor	Inference type should be by_definition
ch-f87bbe81877	1.1	minor	“ n sites of a ring” phrasing ambiguous
Node 1.2 — Polynomial Identification (13 challenges, all resolved)			
ch-3281658e3b5	1.2	critical	Wrong arXiv for AMW24
ch-6bf8ddf12b6	1.2	critical	Wrong arXiv for AMW24 (duplicate)
ch-331c0d2575b	1.2	major	P_λ as partition function not established
ch-9b47dae9fa2	1.2	major	Interpolation vs. homogeneous conflation
ch-a7d6a8c27f2	1.2	major	P_λ partition function claim
ch-bbd61e7e621	1.2	major	Notation mismatch F_μ^* vs f_μ^*
ch-bd836338c75	1.2	major	Notation confusion AMW24 vs BDW25
ch-d37cd6df09e	1.2	major	Notation mismatch in identities
ch-daea0680bca	1.2	major	AMW24 uses F_η (no asterisk)
ch-df660cc9766	1.2	major	Conflation interpolation/homogeneous
ch-34f72af7e1d	1.2	minor	Missing domain of validity
ch-9dc89a4a93d	1.2	minor	Imprecise E_λ^* description
ch-f566b5b1fc6	1.2	minor	Missing dependency on Node 1.1
Node 1.3 — Positivity and Normalization (11 challenges, all resolved)			
ch-25c6c5abfb4	1.3	critical	Circular positivity argument
ch-35912e69e9b	1.3	critical	Undeclared circular dependency
ch-51f3984ac6d	1.3	critical	Circular positivity: $F_\mu^* \geq 0$ standalone
ch-8f12f73248b	1.3	critical	Missing dependency on Node 1.6
ch-1e86d0c492e	1.3	major	Hecke symmetrization identity source

Challenge ID	Node	Severity	Summary (abbreviated)
ch-27ac074eac1	1.3	major	Division by P_λ^* requires $P_\lambda^* > 0$
ch-32e4972f5d3	1.3	major	Conclusion presented as assumption
ch-3abeea04155	1.3	major	Hecke symmetrization identity at $q = 1$
ch-c45e58fbcb9a	1.3	major	No proof strategy for nonnegativity
ch-d3617c80310	1.3	major	$P_\lambda^* > 0$ not established
ch-c7c400820a7	1.3	minor	Edge cases not addressed
Node 1.4 — Chain Construction (17 challenges, all resolved)			
ch-14311e6596c	1.4	critical	Ambiguous “next weaker particle”
ch-1ec04fc2163	1.4	critical	Fundamentally incomplete dynamics
ch-41d4aa1ee0c	1.4	critical	Cascade mechanism omitted
ch-f020555a39e	1.4	critical	Ambiguous “next weaker particle”
ch-171bbc0fce3	1.4	major	Vacancy behavior unspecified
ch-39bf563cf04	1.4	major	State space preservation not verified
ch-4653a7ed47b	1.4	major	Irreducibility not addressed
ch-4cf17322d27	1.4	major	Missing parameter domains
ch-50b37b860d3	1.4	major	Irreducibility not addressed
ch-74a7074b0a7	1.4	major	CT vs DT not addressed
ch-a1e37bfece4	1.4	major	CT vs DT bridge missing
ch-cb75c248d3c	1.4	major	State space preservation
ch-da6b8df72d3	1.4	major	Vacancy behavior unspecified
ch-e172d2fb63a	1.4	major	No theorem citation
ch-31ae7af3f1e	1.4	minor	Misleading phrase
ch-6ead9e2ddfb	1.4	minor	Missing dependency on 1.1
ch-7274c0c9ec2	1.4	minor	Wrong inference type
Node 1.5 — Stationarity (18 challenges, all resolved)			
ch-1c4cf108384	1.5	critical	Misattribution to Ferrari–Martin
ch-19bf58e0220	1.5	critical	Marginal-to-polynomial gap
ch-6670d121be2	1.5	critical	No actual proof structure
ch-76d4c8c4a62	1.5	critical	Misattribution of proof method
ch-934661602e9	1.5	critical	False product-form claim

Challenge ID	Node	Severity	Summary (abbreviated)
ch-9658da4177f	1.5	critical	False product-form claim
ch-e017b73e220	1.5	critical	No actual proof structure
ch-053def23083	1.5	major	Uniqueness not established
ch-26806410fb8	1.5	major	Notation inconsistency
ch-3e9c0d7d87e	1.5	major	Uniqueness of stationary distrib.
ch-4139786e2b5	1.5	major	$q = 1$ specialization not justified
ch-4aad6c675a9	1.5	major	Wrong inference type
ch-6dbd5eeae65	1.5	major	No citation to specific theorem
ch-974ea3da1b7	1.5	major	Missing parameter domain
ch-c7fac73aacf	1.5	major	Missing dependency on 1.4
ch-d3321b98205	1.5	major	No specific theorem citation
ch-fa80af229fc	1.5	major	Missing dependency on 1.4
ch-ae87a001369	1.5	minor	Ambiguous phrase
Node 1.6 — Ratio Identity (16 challenges, all resolved)			
ch-0f3e80c57dd	1.6	critical	Wrong characterization of lower-order terms
ch-1b65900e006	1.6	critical	Missing stationarity proof
ch-3cd11cc979f	1.6	critical	Fatal logical fallacy (“both sum to 1”)
ch-3dbba9297a9	1.6	critical	Open research problem
ch-430bb8a8557	1.6	critical	Vague Hecke symmetrization
ch-4b9c28e7647	1.6	critical	Vague Hecke symmetrization
ch-710ff6d7316d	1.6	critical	BDW25 §7 factorization relevance
ch-8f48027c9c1	1.6	critical	Wrong inference type
ch-9505e439a36	1.6	critical	Fatal logical fallacy
ch-d137782ed8d	1.6	critical	Ratio identity is open problem
ch-e00bb70559e	1.6	critical	Unproven prerequisite
ch-54dae5b9c11	1.6	major	Conflation of notation
ch-7fd14d62c5c	1.6	major	Alternative strategies identified but not pursued
ch-e4f2f63360a	1.6	major	No concrete verification for $n = 2$
ch-fef9b62e8eb	1.6	major	Undeclared dependency on 1.4, 1.5
ch-2d5406cea99	1.6	major	Missing parameter domain

Challenge ID	Node	Severity	Summary (abbreviated)
Node 1.7 — Nontriviality (13 challenges, all resolved)			
ch-0781b5a14b8	1.7	critical	Argument is not a proof
ch-0e46adbdfbd	1.7	critical	Transition probs may be described using F_μ^*
ch-0fe08ac59a5	1.7	critical	“Described using” is undefined
ch-59a456f3dad	1.7	critical	Not a proof: “local rates” non-sequitur
ch-6f040c21803	1.7	critical	False claim of locality
ch-dda1f9555ce	1.7	critical	What does “described using” mean?
ch-f044bfe282d	1.7	critical	False claim of locality
ch-38532088086	1.7	major	Undeclared dependency on 1.4
ch-446df132c99	1.7	major	F_μ^* are not symmetric-function stationary weights
ch-53a810054e3	1.7	major	No engagement with Hecke connection
ch-a4ba65ae810	1.7	major	F_μ^* are not symmetric-function stationary weights
ch-ae20e2d87b0	1.7	major	Wrong inference type
ch-4068a87ead2	1.7	minor	Missing parameter domain
Node 1.8 — Conclusion (2 challenges, all resolved)			
ch-0d50b43ab6a	1.8	critical	Conclusion depends on all siblings
ch-8ffcc34da1f	1.8	major	Missing parameter domains

C Lessons Learned

1. **The Hecke stationarity argument (Node 1.6) is a plausible original mathematical argument** — not just a citation of existing work. It uses BDW25 Proposition 2.10 + AMW24 Theorem 1.1 + Perron–Frobenius in a novel combination. This makes it the most vulnerable node to further challenges.
2. **Restructuring Node 1.3 as a corollary of 1.5 + 1.6 elegantly breaks the circularity** without needing to prove the open problem of $F_\mu^* \geq 0$ directly.
3. **The nontriviality argument via sparsity (Node 1.7) is clean and concrete** — exhibiting a zero in the transition matrix is much stronger than informal “local vs global” heuristics.
4. **Notation discipline pays off** — the cross-paper notation table in Node 1.2 makes all subsequent nodes readable and prevents notation-related challenges.
5. **Provers should always resolve ALL challenges**, not just critical ones. Leaving minor challenges open accumulates technical debt.
6. **Session 2 revealed that LLM provers have a tendency to fabricate proof methods** (e.g., attributing the stationarity proof to Ferrari–Martin). Adversarial verification is essential for catching such errors.
7. **The most common challenge type was “notation inconsistency”** (across nodes 1.2, 1.3, 1.5, 1.6). Establishing a canonical notation table early (Node 1.2) is critical.