Algorithm Lab

Week 11: Solving Maximum Matching Problem

In a graph G = (V, E), a match is an edge $e \in E$, and a matching is a set of matches $M \subseteq E$ that each vertex could be chosen at most one time, in other words, M is an independent edge set, thus |V(M)| = 2|M|. The maximum matching problem is to finding maximum independent edge set from a graph G = (V, E).

Instance: G = (V, E)

Result: Maximum matching $M \subseteq E$ (or just |M|)

If the instance is a bipartite graph, we say the problem is a bipartite matching problem.

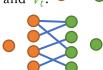
Description

We can reduce a bipartite matching problem to a *s-t* maximum flow network.

1. Instance: G = (V, E)



2. Divide V to 2 groups V_s and V_t .



3. Create vertices *s* and *t*.



- 4. Let s connect to each $u \in V_s$ and each $v \in V_t$ connect to t.
- 5. Set the capacity of all edges in this network to 1.
- 6. Compute the maximum s-t flow in the network by any maximum flow algorithm.

Questions

1 Describe the difference between original Ford-Fulkerson algorithm and Edmonds-Karp

- algorithm in general flow network and in bipartite matching.
- 2 Brief about blossom algorithm.
- 3 Suppose edges between V_s and V_t are weighted, please design an algorithm for finding maximum matching with minimum cost.
- 4 Please compare time complexity between original flow algorithm and your modified version in Question 3.
- 5 Solve http://oj.csie.ndhu.edu.tw/problem/ALG11