

ViolationLS

Constraint Based Local Search in CP-SAT

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Constraint-Based Local Search

Constraint Based: inspired by Constraint Programming

Many problems have similar combinatorial structures, don't make users re-implement them!



Constraint-Based Local Search

Local Search: makes local changes to an assignment

Constraints define "how violated" they are in an assignment, solvers first minimize total violation, then the objective.

Constraints typically also define a neighborhood.*



Violation Functions

$$V \left(\sum cx \leq K \right) = \min \left(\sum cx - K, 0 \right)$$



Feasibility Jump

Feasibility Jump = Single-var moves + Guided Local Search



Feasibility Jump

Feasibility Jump changes **one variable at a time**

The variable "jumps" to the value that minimizes the **weighted sum** of constraint violations for that variable.

If no moves reduce weighted violation, **increase the weight of violated constraints**

If feasible, **emit the new solution** and replace objective constraint:
objectiveVar < currentObjectiveValue



Feasibility Jump

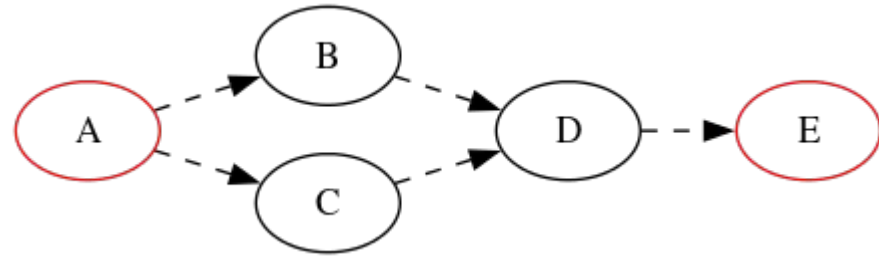
Computing the "Jump Value" is cheap if we assume violation is convex
when changing a single variable:

$$O(\ln(x_{max} - x_{min}))$$



Network Flow Example

Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E



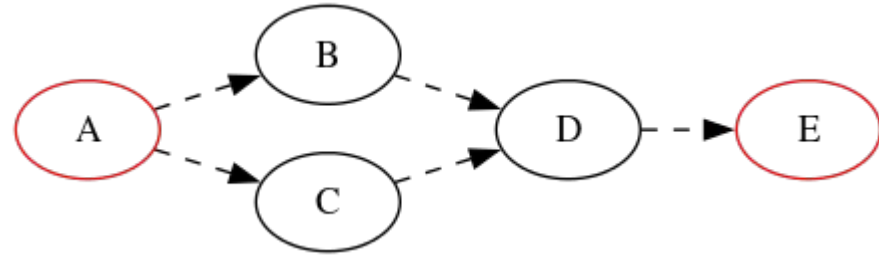
	A	B	C	D	E
Weight	1	1	1	1	1
Violation	1	0	0	0	1

	AB	AC	BD	CD	DE
Score	0	0	-2	-2	0
Value	0	0	0	0	0



Network Flow Example

Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E



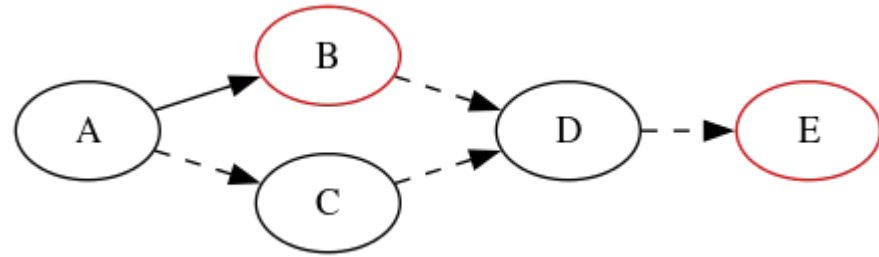
	A	B	C	D	E
Weight	2	1	1	1	2
Violation	1	0	0	0	1

	AB	AC	BD	CD	DE
Score	1	1	-2	-2	1
Value	0	0	0	0	0



Network Flow Example

Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E



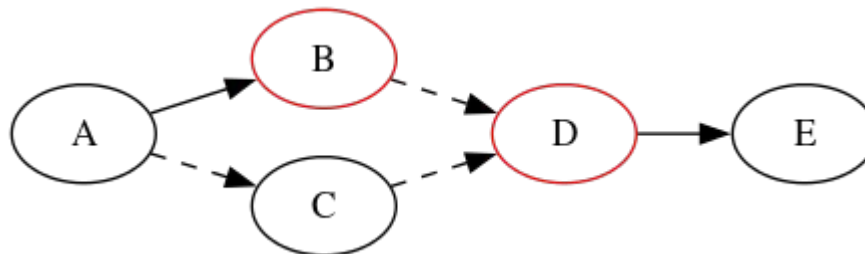
	A	B	C	D	E
Weight	2	1	1	1	2
Violation	0	1	0	0	1

	AB	AC	BD	CD	DE
Score	0	-3	0	0	1
Value	1	0	0	0	0



Network Flow Example

Arc vars have domain $\{0, 1\}$;
1 inflow at A; 1 outflow at E

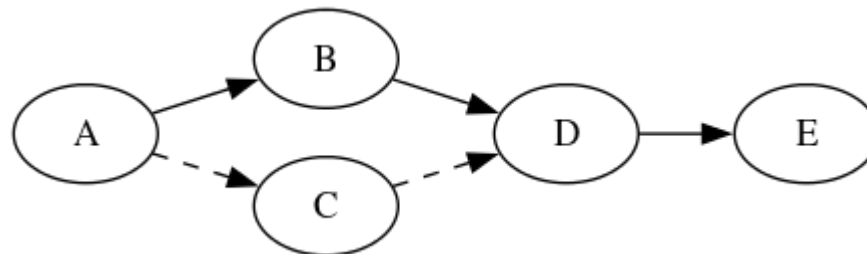


	A	B	C	D	E
Weight	2	1	1	1	2
Violation	0	1	0	1	0

	AB	AC	BD	CD	DE
Score	0	-3	2	0	0
Value	1	0	0	0	1



Network Flow Example



Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E

	A	B	C	D	E
Weight	2	1	1	1	2
Violation	0	0	0	0	0

	AB	AC	BD	CD	DE
Score	0	-1	0	-1	0
Value	1	0	1	0	1



Multi-Variable Moves

Usually local search changes more than 1 variable at a time!

We'd like a way to find chains of moves automatically,
while using the same operations as Feasibility Jump

But I'm lazy: I don't want to write specialized neighbourhoods



Local Search as Planning

Idea: Treat the search for a move as a domain-independent planning problem.

Note: this is just an inspiration, we don't really solve a planning problem!



Novelty

"Novelty" based search: only expands "novel" states

More structured than random exploration

Polynomial work per search



Novelty

What makes a state novel?

Some atomic fact is true in the state that has never been true in any prior one



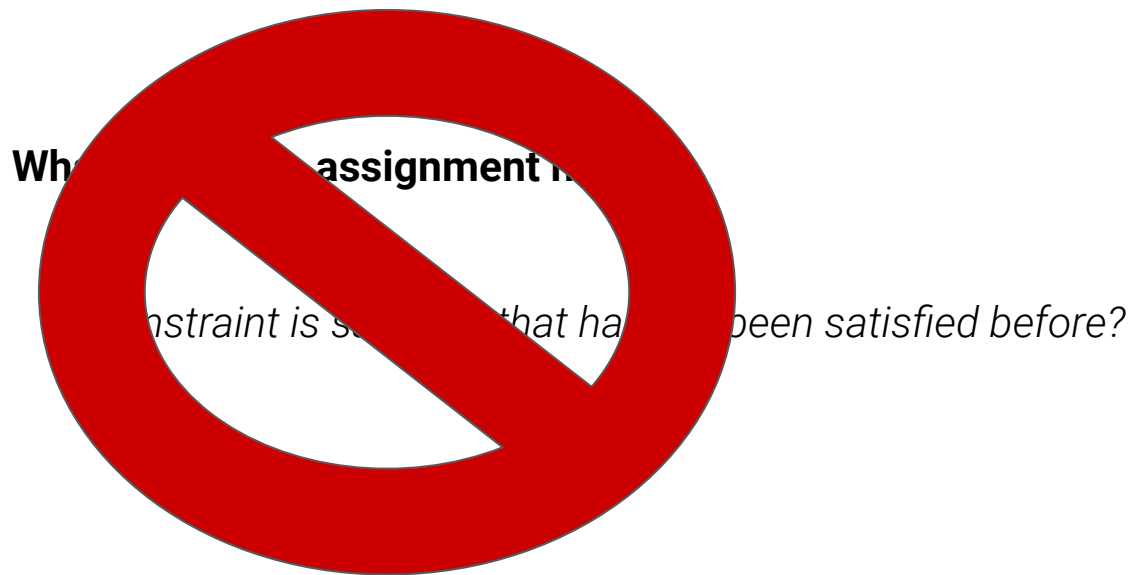
Novelty

What makes an assignment novel?

A constraint is satisfied that has not been satisfied before?



Novelty



Novelty

What makes an assignment novel?

*A constraint is **violated** that has not been **violated** before*

Or (non-discounted) weighted violation is less than any prior assignment



Novelty Jump

How to make "jump if weighted violation decreases" find novel states:

- Discount the weight of satisfied constraints (multiply by epsilon)
- Reset the weight when a constraint becomes violated
- Keep a stack of var changes & backtrack when we have no good moves*
- Clear the stack when sum of (non discounted) scores are positive



Novelty Jump: Why backtrack?

Discounted score can be positive, but real score negative (*which is the point!*)

If we commit such a move we may not find our way back.

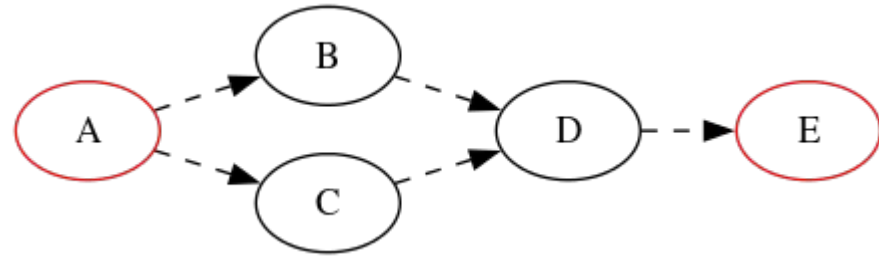
This would be an interesting metaheuristic, but not a compound move.

This interferes with GLS rather than complements it



Novelty Example

Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E



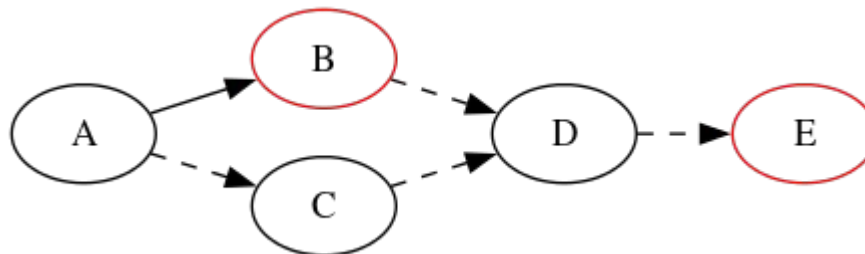
	A	B	C	D	E
Weight	1	ϵ	ϵ	ϵ	1
Violation	1	0	0	0	1

	AB	AC	BD	CD	DE
Score	$1-\epsilon$	$1-\epsilon$	-2ϵ	-2ϵ	$1-\epsilon$
Value	0	0	0	0	0



Novelty Example

Arc vars have domain $\{0, 1\}$;
1 inflow at A; 1 outflow at E



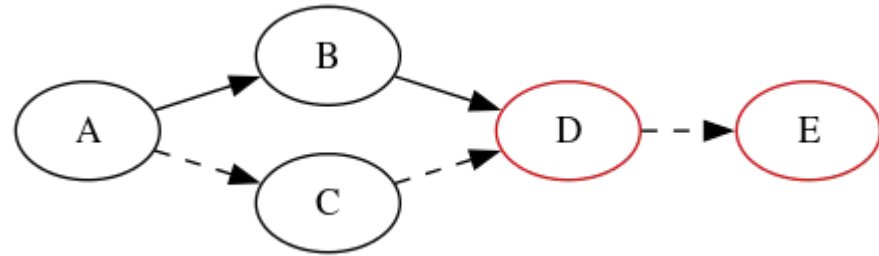
	A	B	C	D	E
Weight	1	1	ϵ	1	1
Violation	0	1	0	0	1

	AB	AC	BD	CD	DE
Score	0	-1-ϵ	$1-\epsilon$	-2ϵ	$1-\epsilon$
Value	1	0	0	0	0



Novelty Example

Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E



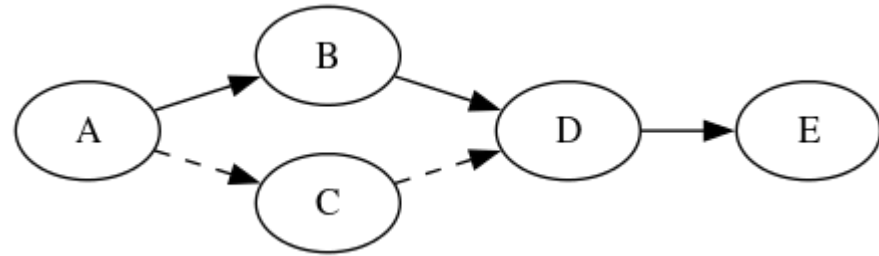
	A	B	C	D	E
Weight	1	1	ϵ	1	1
Violation	0	0	0	1	1

	AB	AC	BD	CD	DE
Score	0	-1-ϵ	0	1-ϵ	2
Value	1	0	1	0	0



Novelty Example

Arc vars have domain $\{0,1\}$;
1 inflow at A; 1 outflow at E



	A	B	C	D	E
Weight	1	1	ε	1	1
Violation	0	0	0	0	0

	AB	AC	BD	CD	DE
Score	0	$-1-\varepsilon$	-2	$-1-\varepsilon$	0
Value	1	0	1	0	1



Is Novelty Jump always better?

No! But you don't have to pick!

ViolationLS:

- If there is a new better solution from some other worker:
 - Set incumbent to the new solution
 - Randomly pick Novelty Jump or Feasibility Jump
- Do some iterations
- Repeat



Results

Solver	ViolationLS	Feasibility Jump	Yuck	fzn-oscar-cbls
ViolationLS	-	650.6	652.28	734.64
Feasibility Jump	182.00	-	523.84	617.24
Yuck	209.52	268.96	-	435.72
fzn-oscar-cbls	101.96	135.36	147.6	-

#instances (of 1216) where ViolationLS beats Yuck

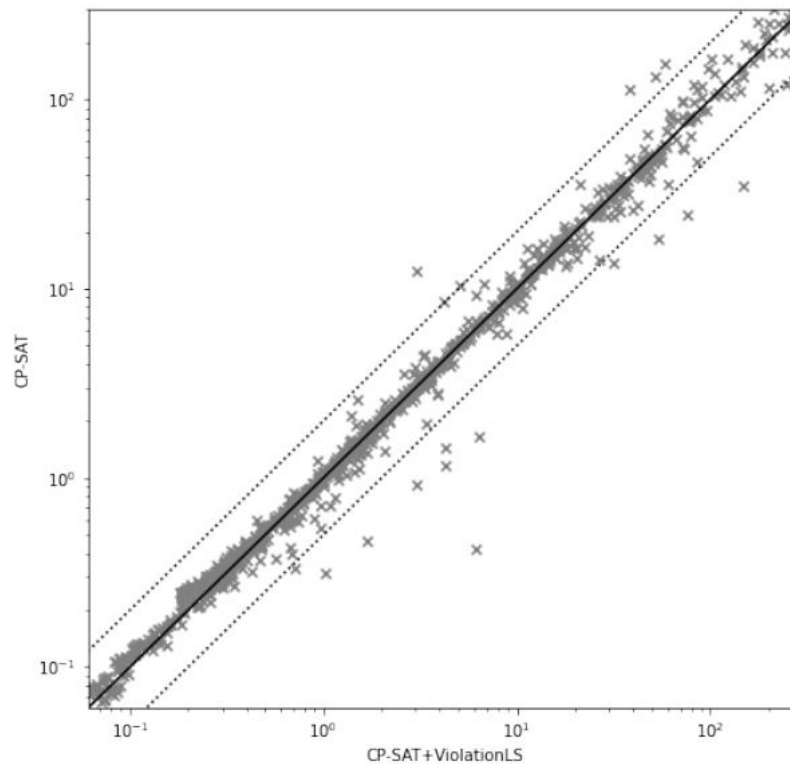
#instances where row solver beats column (1216 mzn challenge instances)

Using 8 cores, except fzn-oscar-cbls which is single-threaded

Averaged over 5 seeds per solver



Results



Solve time of CP-SAT with and without ViolationLS
60.5% faster; 39.5% slower



Results

CP-SAT+	ViolationLS	Feasibility Jump	Baseline
ViolationLS	-	609.20	659.04
Feasibility Jump	588.36	-	651.80
Baseline	538.32	545.48	-

#instances where column solver beats row (1216 mzn challenge instances)
Using 8 cores



Conclusions

- CBLS makes a powerful primal heuristic for a state-of-the-art CP solver
- Efficiently implemented, generic LS moves can beat constraint-specific ones
- Novelty is a good exploration criterion that integrates well with GLS

