

Report for assignment 1

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Multiplication of long integers using FFT/IFFT technique

1. Random generation of integers

The digits of the integers are generated randomly by taking the number of digits as a parameter.

2. Obtaining polynomial from integers

Polynomials are obtained by taking the digits of the integers to be the coefficients of the polynomials

E.g. : $5962 = A(x) = 2 + 6x + 9(x^2) + 5(x^3)$

3. Fast Fourier Transform

Consider the above polynomial. Suppose it is to be evaluated at four points. The polynomial is broken into 2 parts, even and odd, and evaluated recursively.

$$A_{\text{even}}(x) = 2 + 9x$$

$$A_{\text{odd}}(x) = 6 + 5x$$

$$A(x) = A_{\text{even}}(x^2) + xA_{\text{odd}}(x^2)$$

If the four point are taken to be the 4th roots of unity(w), the time for evaluating the polynomial reduces to $O(m \log m)$ using the FFT algorithm.

Algorithm 1 FFT

```
1: FFT( $n, A, F$ )
2: if  $n = 1$  then
3:    $F[0] \leftarrow a_0$ 
4:   return
5: end if
6:  $A_{\text{even}} = [a_0 a_2 \dots a_{n-2}]$ 
7:  $A_{\text{odd}} = [a_1 a_3 \dots a_{n-1}]$ 
8:  $FFT(n/2, A_{\text{even}}, EF)$ 
9:  $FFT(n/2, A_{\text{odd}}, OF)$ 
```

```

10: for  $j \leftarrow 0; j \leq n/2; ++j$  do
11:    $F[j] = EF[j] + w^j * OF[j]$ 
12:    $F[j + n/2] = EF[j] - w^j * OF[j]$ 
13: end for

```

4. Inverse Fast Fourier Transform

The IFFT is used to obtain the coefficients of a polynomial from its value at n points, i.e., the n th roots of unity.

Algorithm 2 IFFT

```

1: IFFT( $n, A, F$ )
2: if  $n = 1$  then
3:    $F[0] \leftarrow a_0$ 
4:   return
5: end if
6:  $A_{even} = [a_0 a_2 \dots a_{n-2}]$ 
7:  $A_{odd} = [a_1 a_3 \dots a_{n-1}]$ 
8: IFFT( $n/2, A_{even}, EF$ )
9: IFFT( $n/2, A_{odd}, OF$ )
10: for  $j \leftarrow 0; j \leq n/2; ++j$  do
11:    $F[j] = EF[j] + w^{-j} * OF[j]$ 
12:    $F[j + n/2] = EF[j] - w^{-j} * OF[j]$ 
13: end for

```

The final result is divided by n .

5. Multiplication of two polynomials

Multiplication of 2 polynomials is achieved by multiplying their fast fourier transforms at sufficient number of points and then applying inverse FFT on the result.

Algorithm 3 Mutiplication

```

1: FFT( $m, A, F_A$ )
2: FFT( $m, B, F_B$ )
3: for  $i = 1, m$  do
4:    $F_C[i] = F_A[i] * F_B[i]$ 
5: end for
6: IFFT( $m, F_C, C$ )
7:  $C \leftarrow \frac{1}{m} * C$ 

```

6. Final Result

The final integer is obtained from the coefficients of the above polynomial.
Example :

$$12 * 26$$

$$A(x) = 2 + x$$

$$B(x) = 6 + 2x$$

$$C(x) = A(x) * B(x) = 12 + 10x + 2x^2$$

$$C \leftarrow 312$$

END