

# Red-Black Trees and AA Trees

Node Color, Insertions and Rotations

SoftUni Team  
Technical Trainers



**SoftUni**



Software University

<http://softuni.bg>

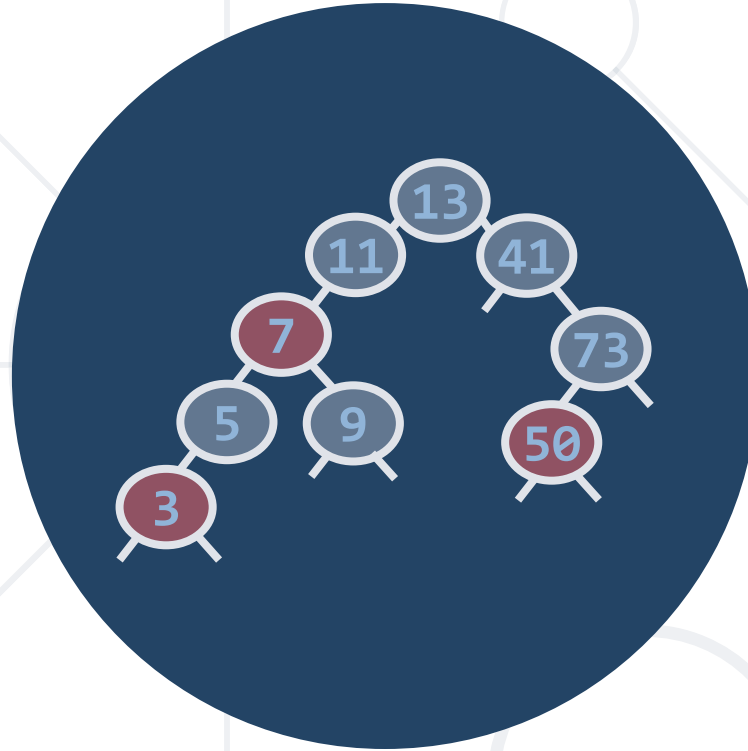
## 1. Red-Black Tree

- Simple Representation of a 2-3 Tree
- Rebalancing Trees
- Rotations
- Insertion Algorithm

## 2. AA Tree

- Insertion Algorithm





# Red-Black Tree

## Simple Representation of a 2-3 Tree

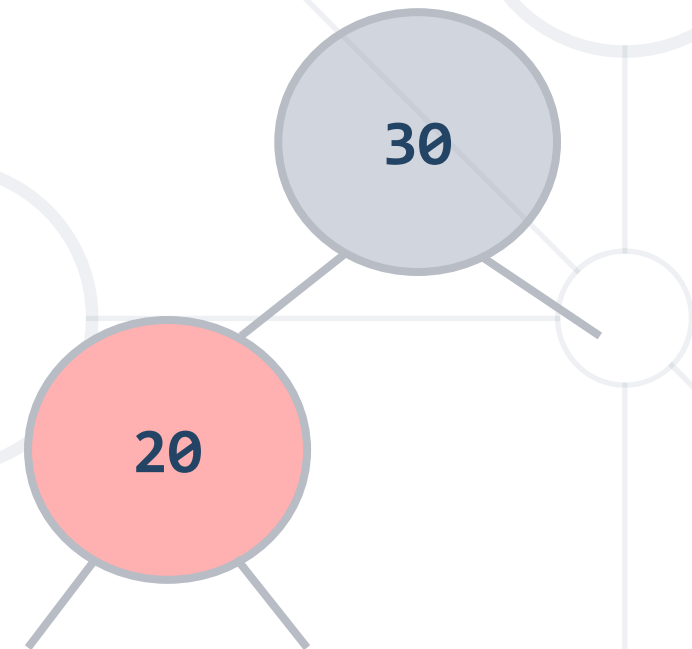
# Why Yet Another Balanced BST?

- We want operations to happen at:
  - $O(\log(n))$  **not**  $O(h)$  where **h** in worst case is **n**
- **AVL** vs **Red-Black** trees:
  - The AVL trees are more balanced that causes more rotations during **insertion** and **deletion**
  - if your application involves many frequent insertions and deletions, then **Red Black** trees should be preferred



# Representing 3-Nodes from 2-3 Tree

- We will represent 3-nodes with a **left-leaning** red nodes
- Nodes with values between the 2 nodes will be to the **right** of the **red** node



# Red-Black Tree Properties

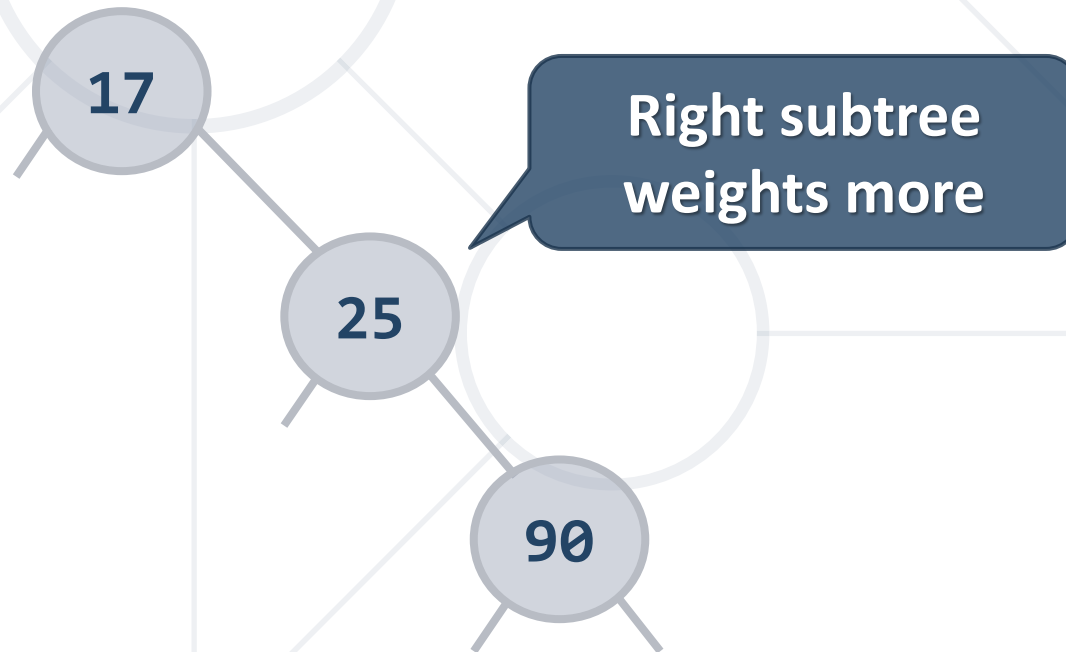
- All **leaves** are black
- The **root** is black
- No node has **two red links** connected to it
- Every path from a given **node** to its **descendant leaf** nodes contains the **same** number of **black** nodes
- Red links **lean** left



# Rebalancing Trees

## Rotations

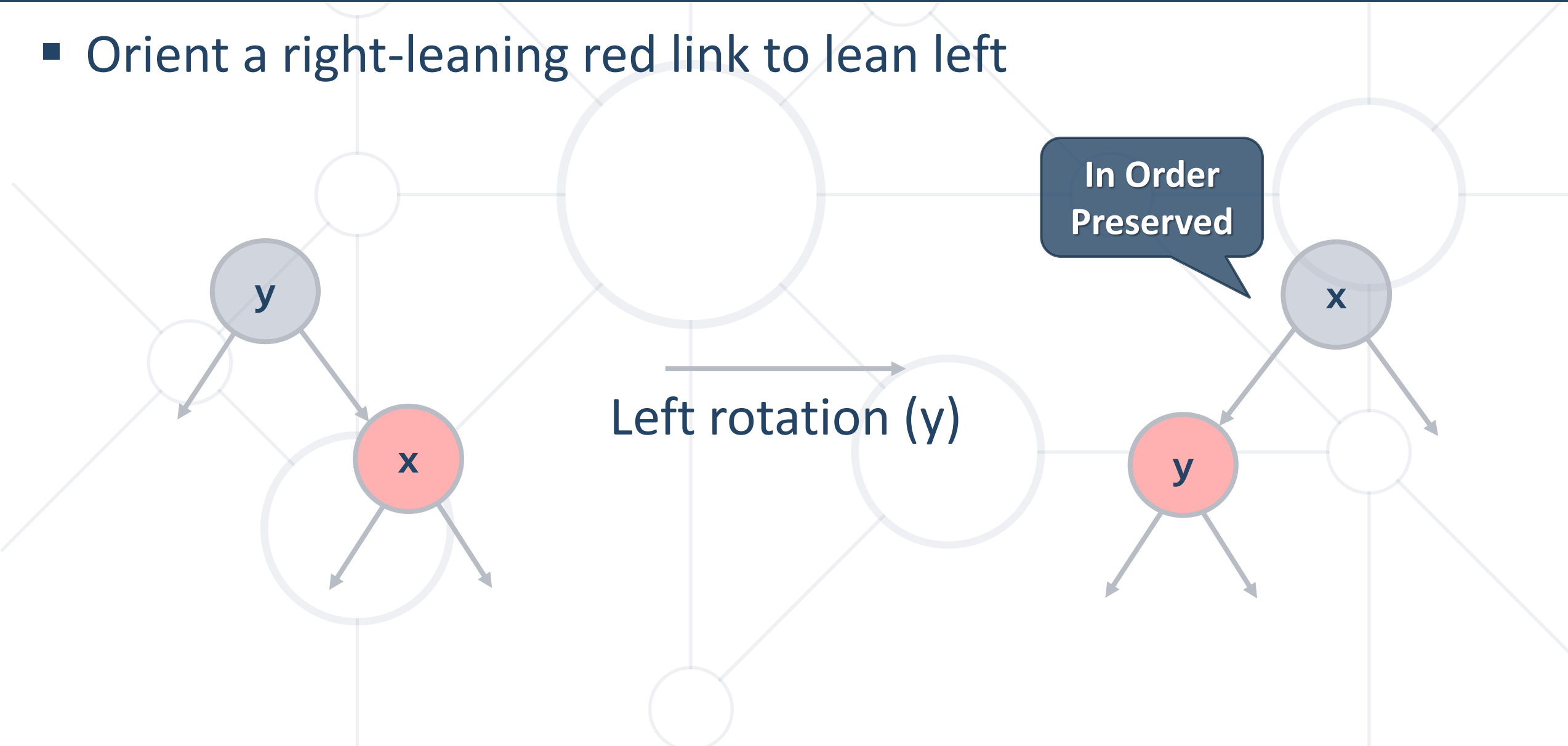
- Rotations are used to correct the balance of a tree
- Balance can be measured in height, depth, size etc. of subtrees





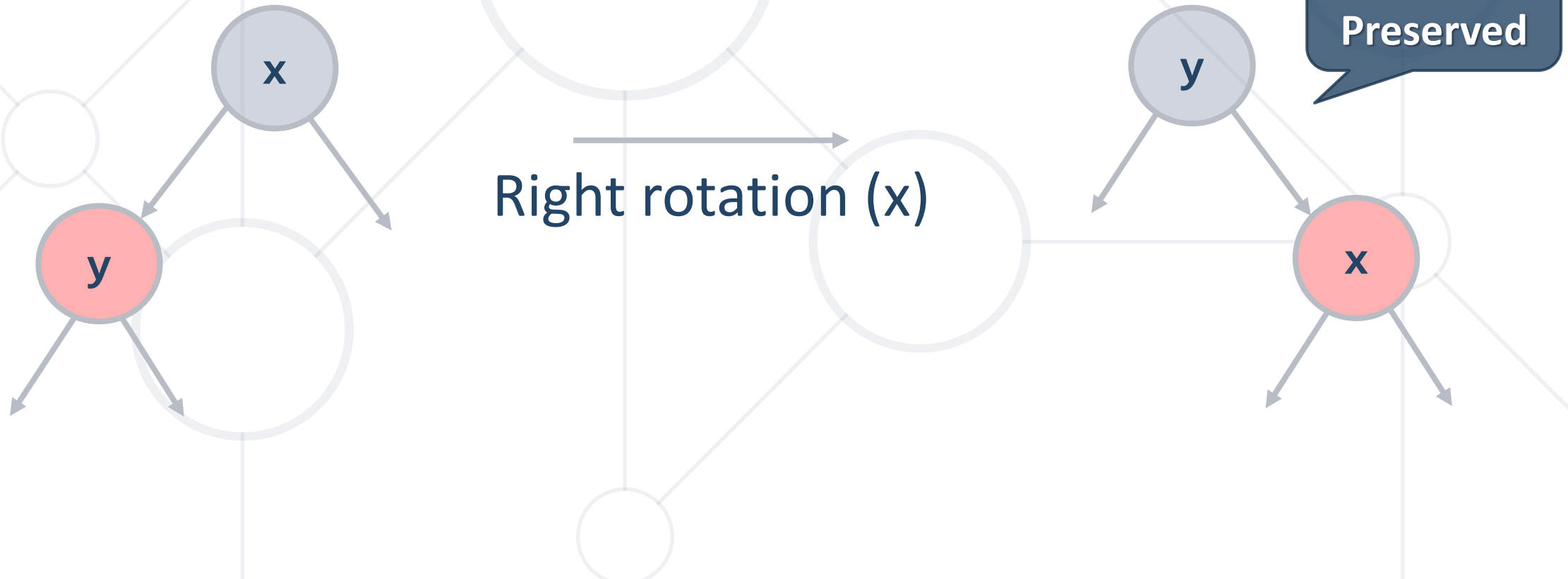
# Left Rotation

- Orient a right-leaning red link to lean left



# Right Rotation

- Orient a left-leaning red link to lean right (temporarily)



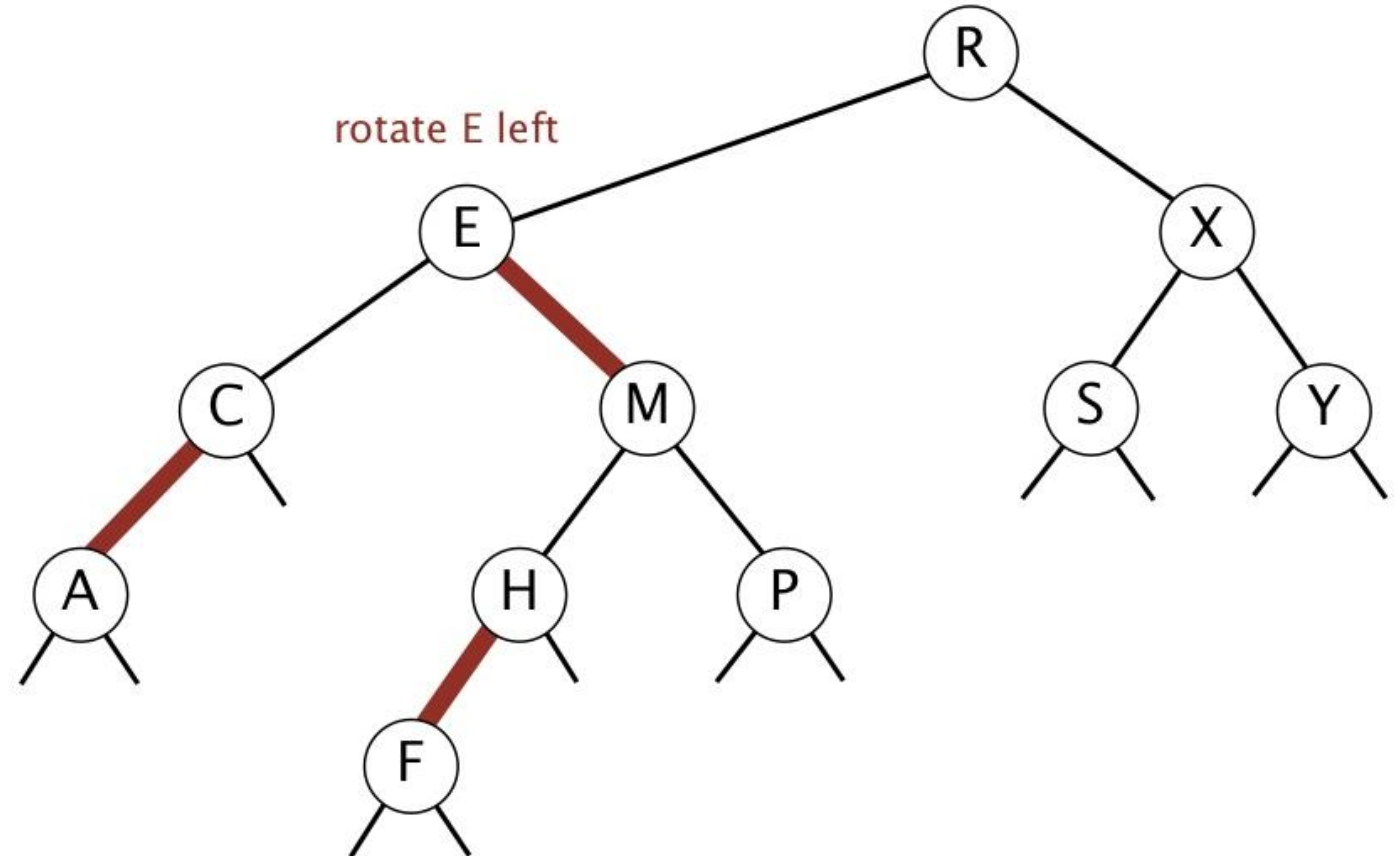
# Rotations - Quiz

A. REXCMSYAHPPF

B. RMXEHSYCFPA

C. RMXEPSYCHAF

D. RCXAESYMHPPF



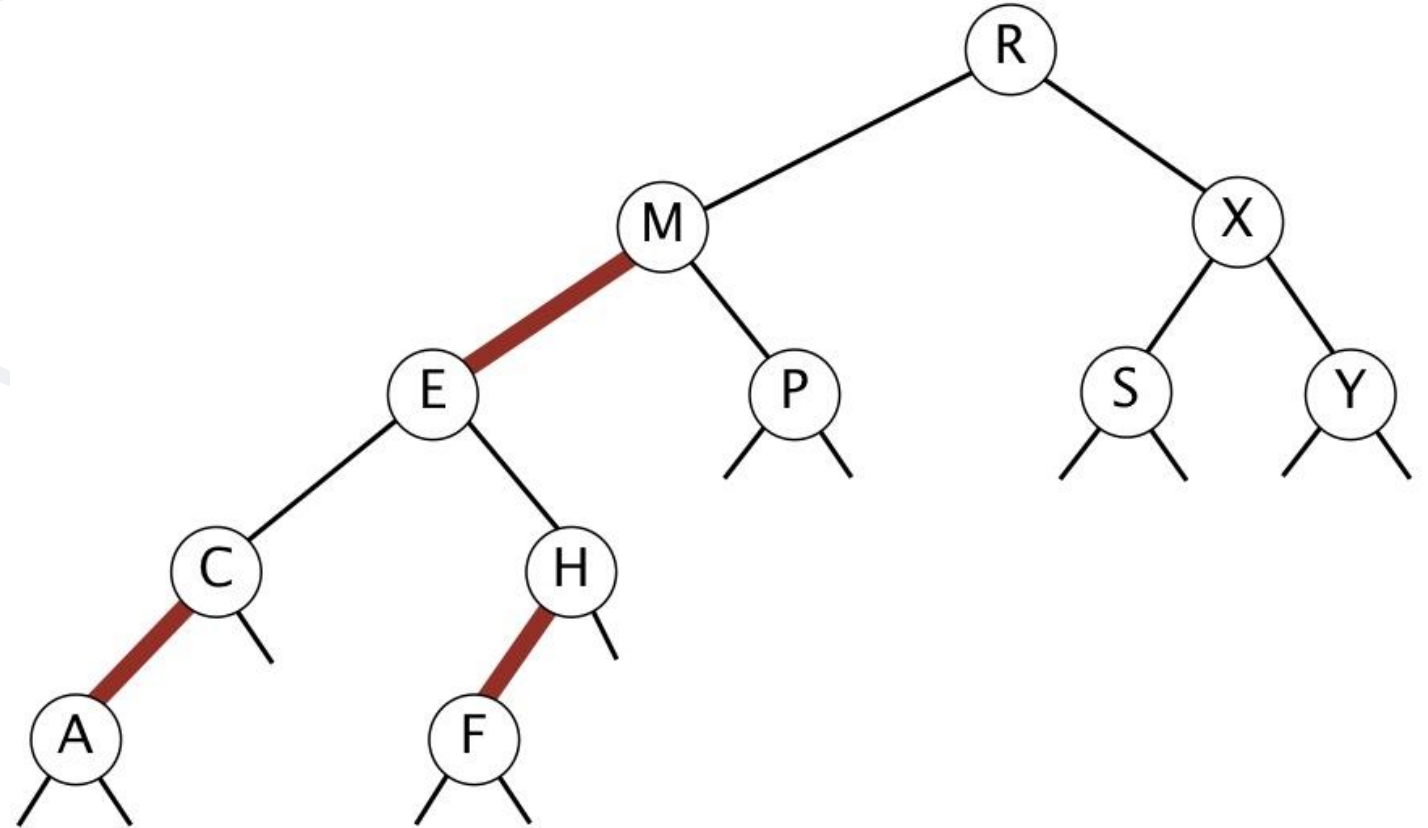
# Rotations - Answer

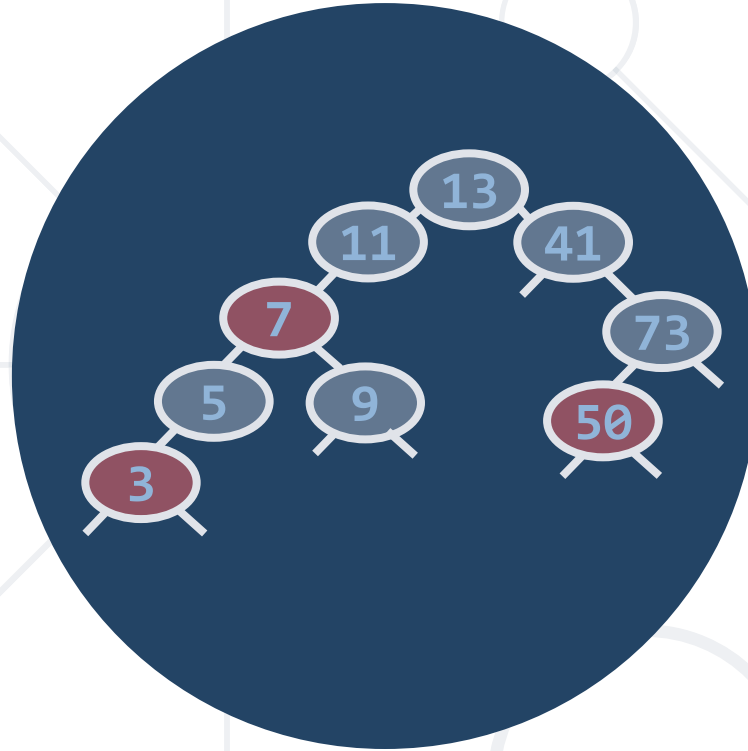
A. REXCMSYAHPPF

B. RMXEHSYCFFA

C. **RMXEPSYCHAF**

D. RCXAESYMHPPF





# Red-Black Tree

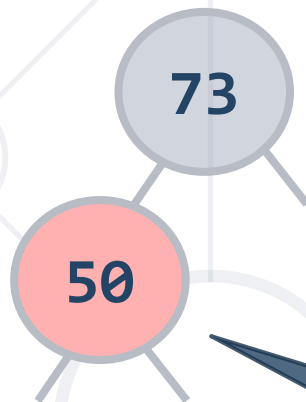
## Insertion Algorithm

# Insertion Algorithm

- **Locate** the node position
- Create new **red** node
- **Add** the new node to the tree
- **Balance** the tree if needed

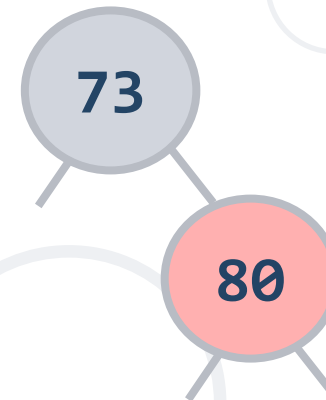
# Insertion

- Insert into a single 2-node:
- Smaller element

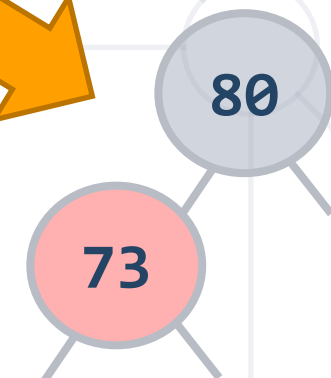


The red node is leaning left

- Larger element

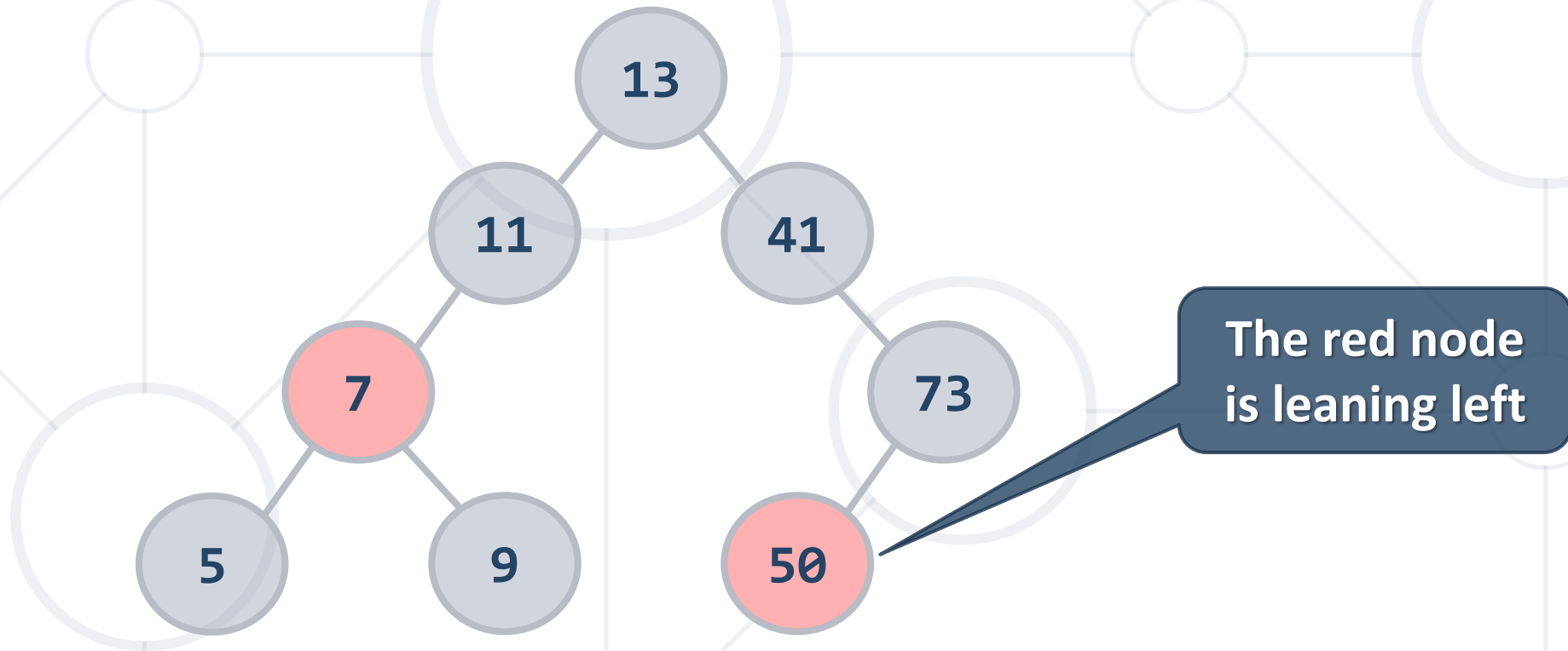


The red node is leaning right, we need left rotation



# Insertion (2)

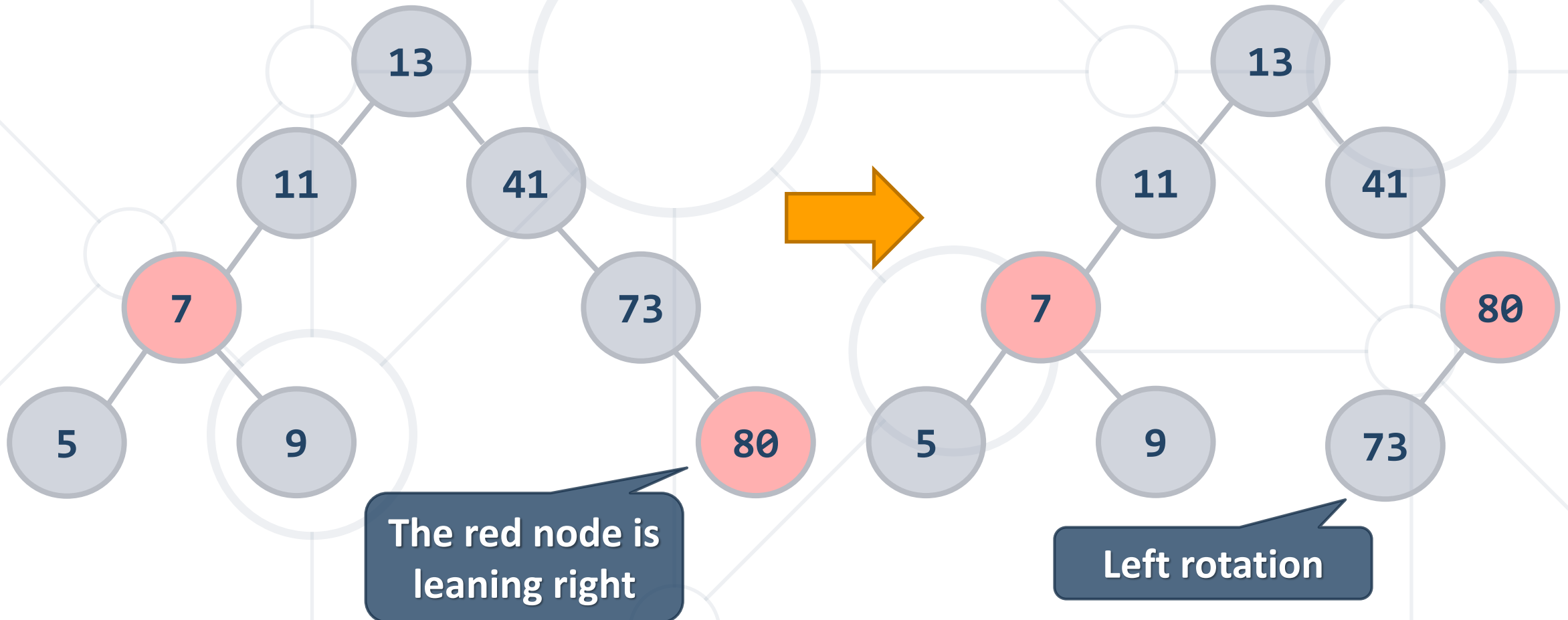
- Insert **smaller** item into a 2-node at the bottom:





## Insertion (3)

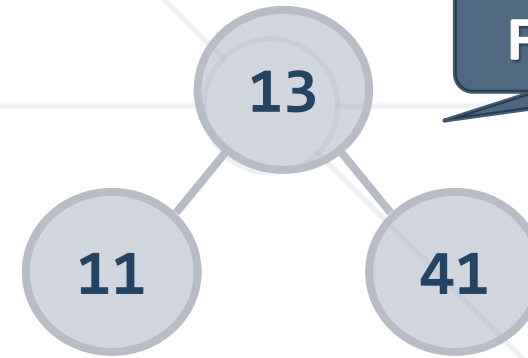
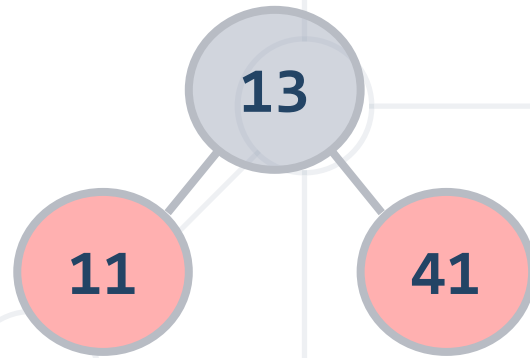
- Insert **larger** item into a 2-node at the bottom:



- 3 cases:
  - The element is **smaller** than both keys
  - The element is **larger** than both keys
  - The element is **between** the 2 keys

# Insertion Into 3-Node (2)

- **Larger** than both keys:

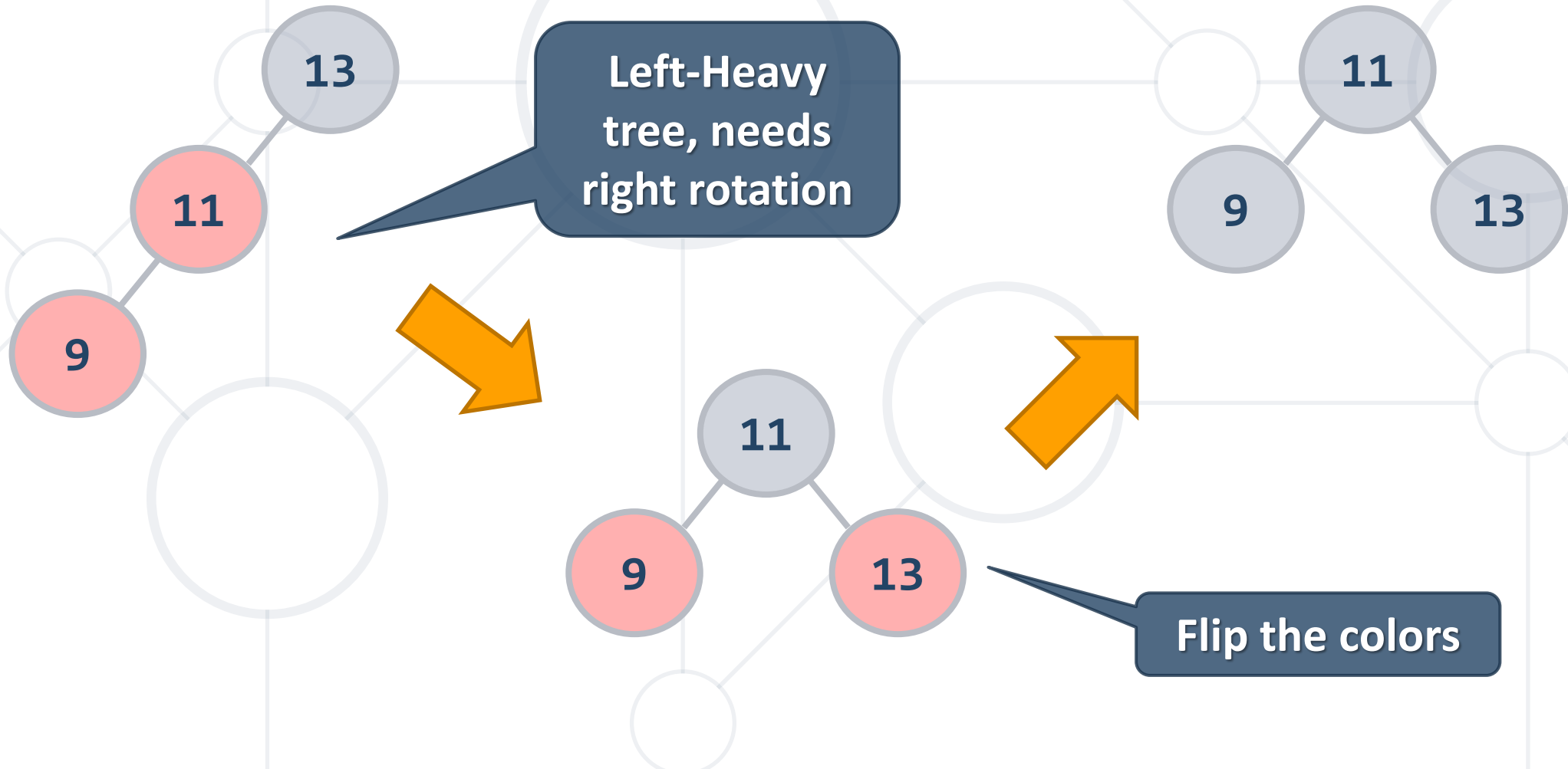


Flip the colors

- Flipping the colors **increases** the **tree height**, which maintains the 1-1 correspondence to 2-3 trees

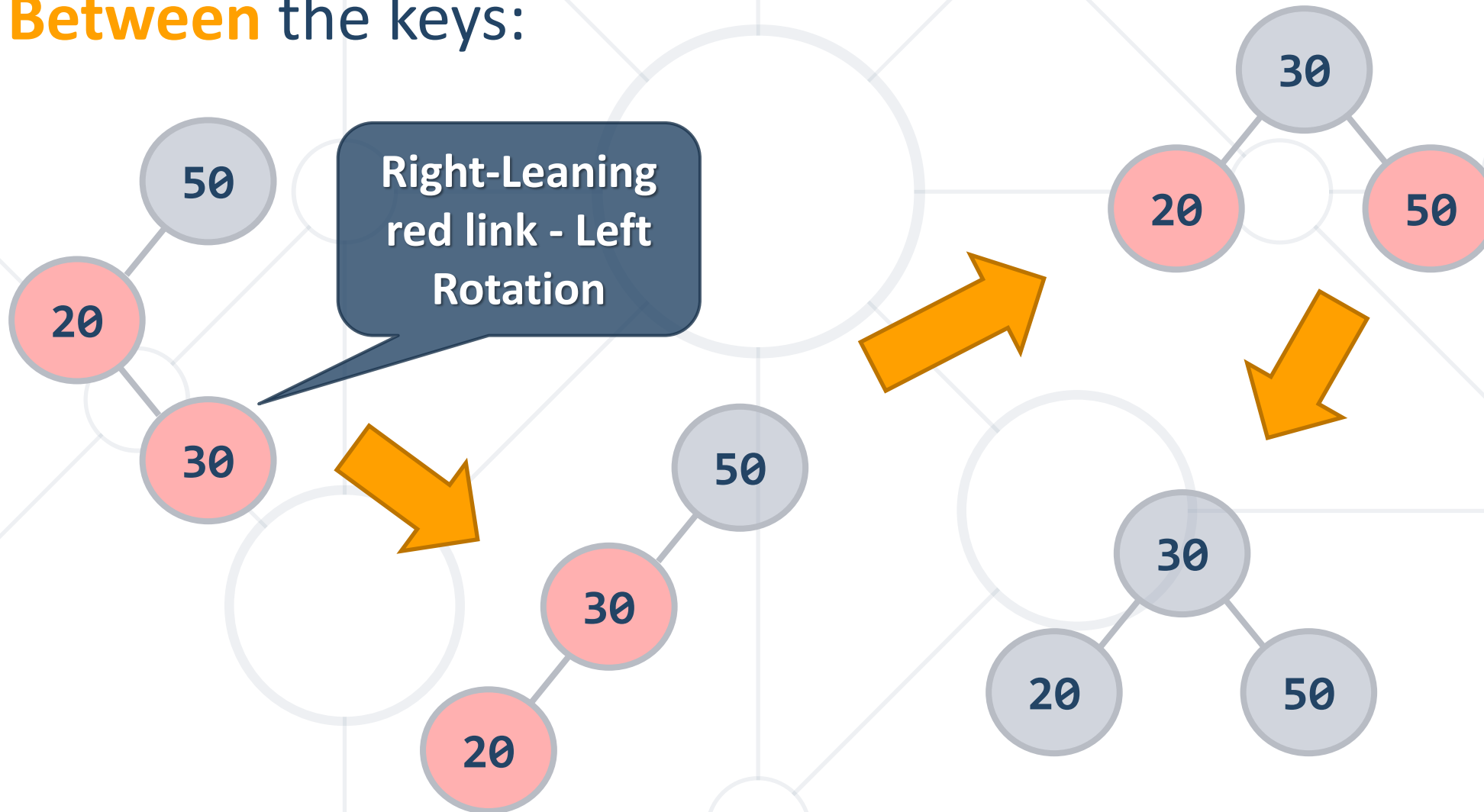
# Insertion Into 3-Node (3)

- **Smaller** than both keys:



# Insertion Into 3-Node (4)

- Between the keys:



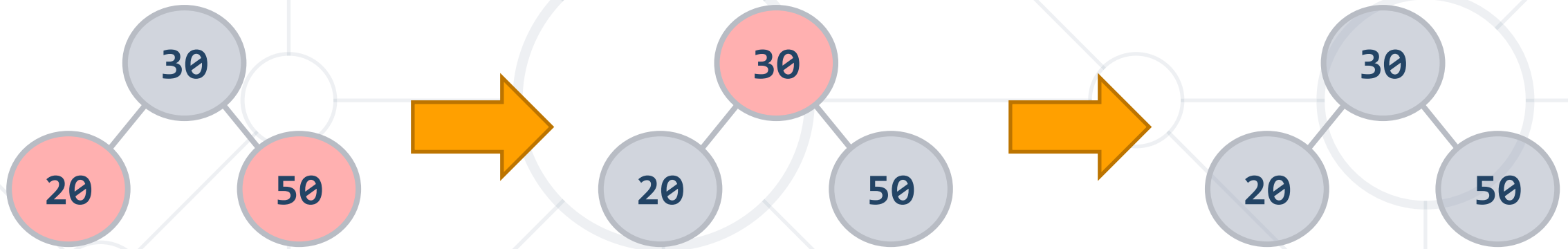
- Flipping the colors should also change the **parent color** to **red**

```
void flipColors(Node<T> node) {  
    node.color = RED;  
    node.left.color = BLACK;  
    node.right.color = BLACK;  
}
```

- Preserves perfect **black balance** in the tree!

# Keeping Black Root

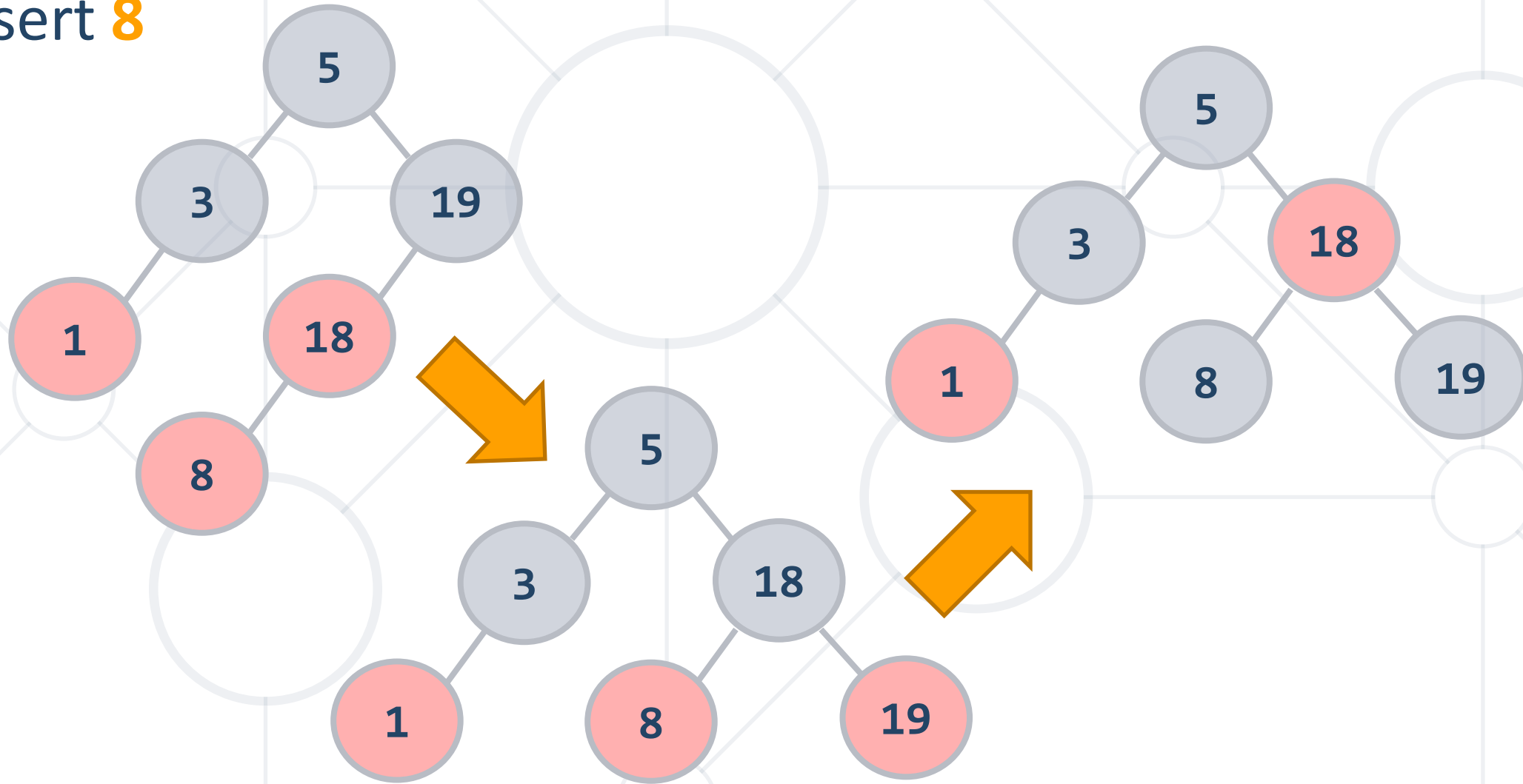
- Insert on a single node (root):



- Each time the root switches colors, the height of the tree is increased

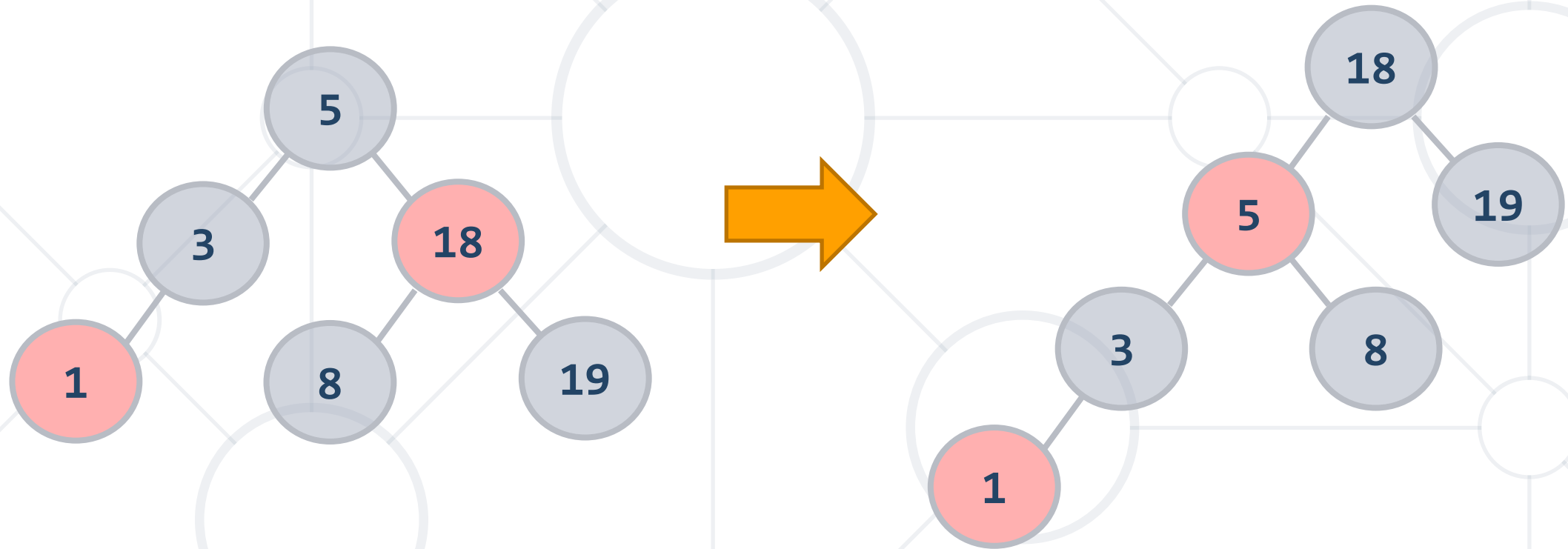
# Insert Into 3-Node at the Bottom

- Insert 8

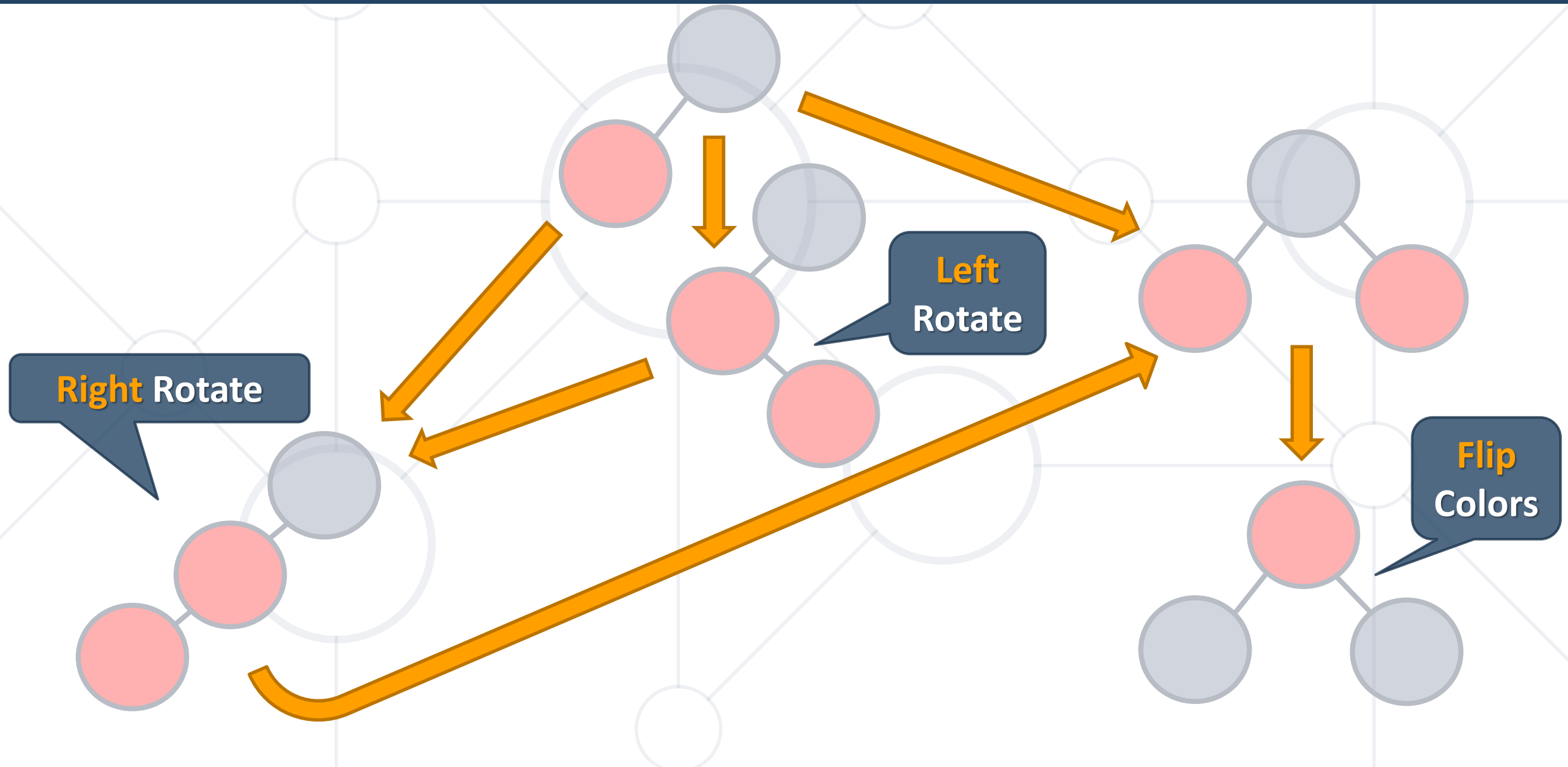


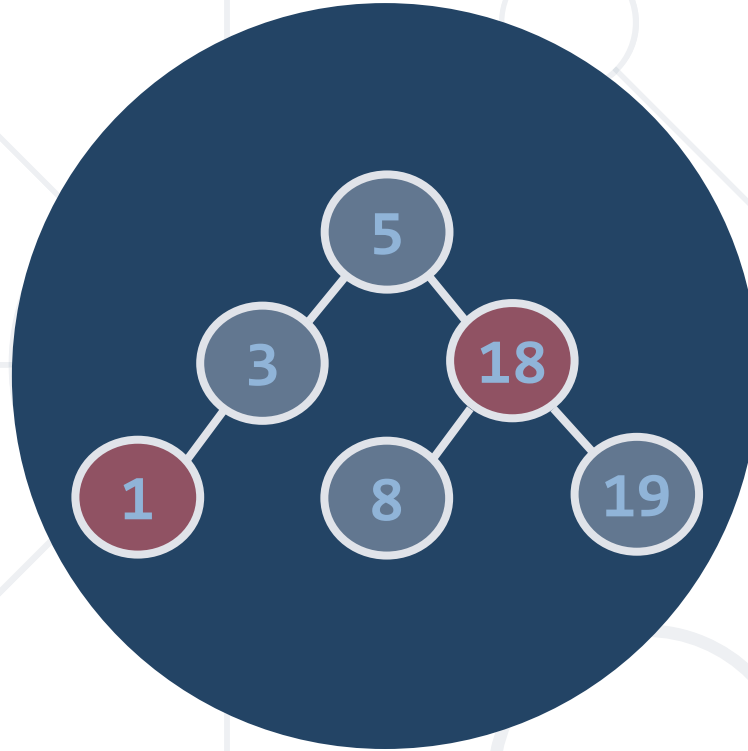


# Insert Into 3-Node at the Bottom (2)



# Overall Insertion Process





# Red-Black Tree

## Insertion Implementation

# Changes to the BST Class

```
class RedBlackTree<T> {  
    private static final boolean RED = true;  
    private static final boolean BLACK = false;  
  
    private static class Node<T> {  
        public boolean color;  
        public Node(T value, bool color) {  
            // TODO: Add setup logic here  
        }  
    }  
}
```

# Changes to the BST Class (2)

```
class RedBlackTree<T> {  
    private boolean isRed(Node<T> node)  
    private Node<T> rotateLeft(Node<T> node)  
    private Node<T> rotateRight(Node<T> node)  
    private void flipColors(Node<T> node)  
}
```

# Rotate Right

```
private Node<T> rotateRight(Node<T> node) {  
    Node<T> temp = node.left;  
    node.left = temp.right;  
    temp.right = node;  
    temp.color = node.color;  
    node.color = RED;  
    node.count = 1 + count(node.left) + count(node.right);  
  
    return temp;  
}
```

# Rotate Left

```
private Node<T> rotateLeft(Node<T> node) {  
    Node<T> temp = node.right;  
    node.right = temp.left;  
    temp.left = node;  
  
    // Same operations as rotateRight()  
  
    return temp;  
}
```

```
private boolean isRed(Node<T> node) {  
    if (node == null) return false;  
    return node.color;  
}
```

```
public void insert(T element) {  
    this.root = this.insert(element, this.root);  
    this.root.color = BLACK;  
}
```



# Insert(2)

```
private Node<T> insert(T element, Node<T> node) {
    if (node == null) node = new Node<>(element, RED);

    // Recursive calls to go left or right

    if (this.isRed(node.right) && !this.isRed(node.left))
        node = this.rotateLeft(node);
    if (this.isRed(node.left) && this.isRed(node.left.left))
        node = this.rotateRight(node);
    if (this.isRed(node.left) && this.isRed(node.right))
        this.flipColors(node);


    // Increase count
}
```

# Red-Black Tree - Quiz

TIME'S

- Suppose that you insert  $n$  keys in ascending order into a red-black BST. What is the height of the resulting tree?
  - Constant
  - Logarithmic
  - Linear
  - Linearithmic

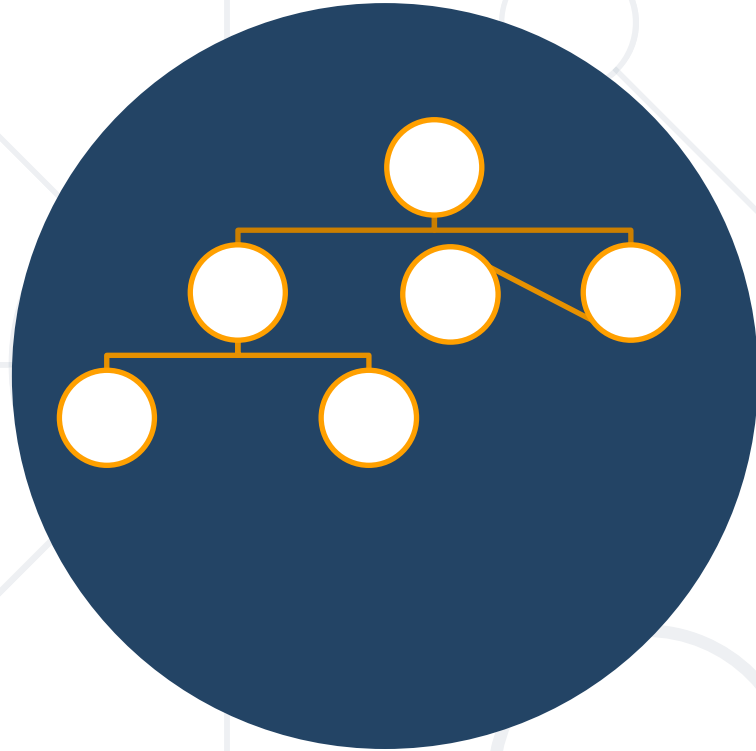
# Red-Black Tree - Answer

- Suppose that you insert  $n$  keys in ascending order into a red-black BST. What is the height of the resulting tree?
  - Constant
  - **Logarithmic** 
  - Linear
  - Linearithmic

The height of any red-black BST on  $n$  keys (regardless of the order of insertion) is guaranteed to be between  $\log_2 n$  and  $2\log_2 n$

# Red-Black Tree - Summary

Structure	Worst case			Average case	
	Search	Insert	Delete	Search Hit	Insert
BST	$N$	$N$	$N$	$1.39 \lg N$	$1.39 \lg N$
2-3 Tree	$c \lg N$	$c \lg N$	$c \lg N$	$c \lg N$	$c \lg N$
Red-Black	$2 \lg N$	$2 \lg N$	$2 \lg N$	$\lg N$	$\lg N$



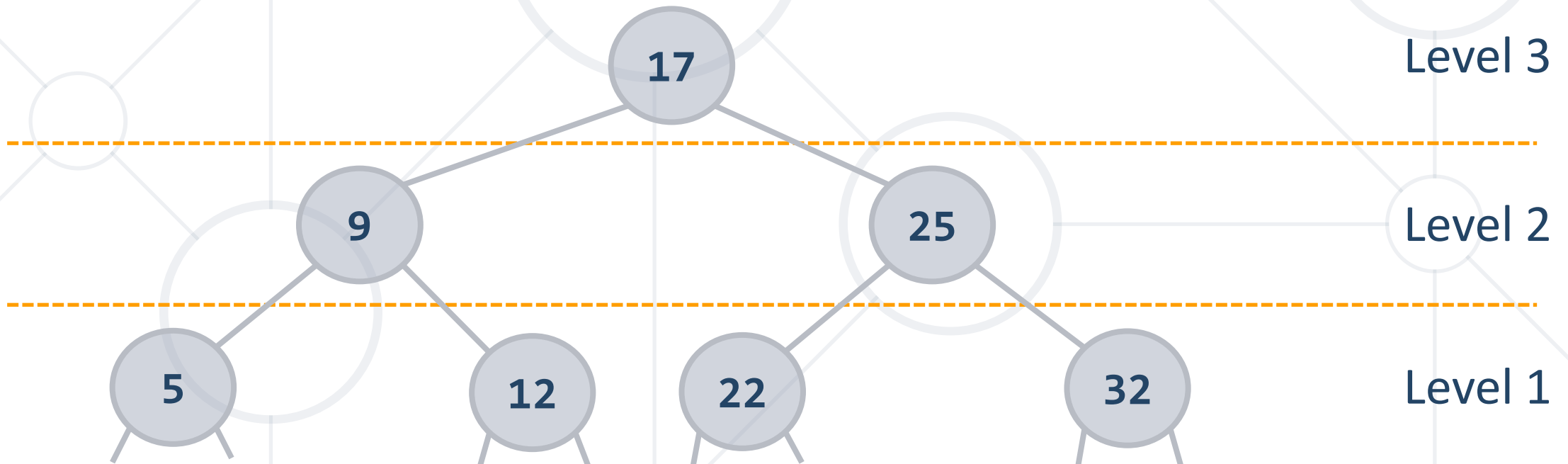
# AA Tree Definition

# Why AA Trees

- **Red-Black** vs **AA** trees:
  - The implementation and number of rotation cases in Red-Black Trees is **complex**. AA trees **simplifies** the algorithm
  - It eliminates **half** of the restructuring process by eliminating half of the **rotation** cases, which is easier to code
  - It **simplifies** the deletion process by removing multiple cases



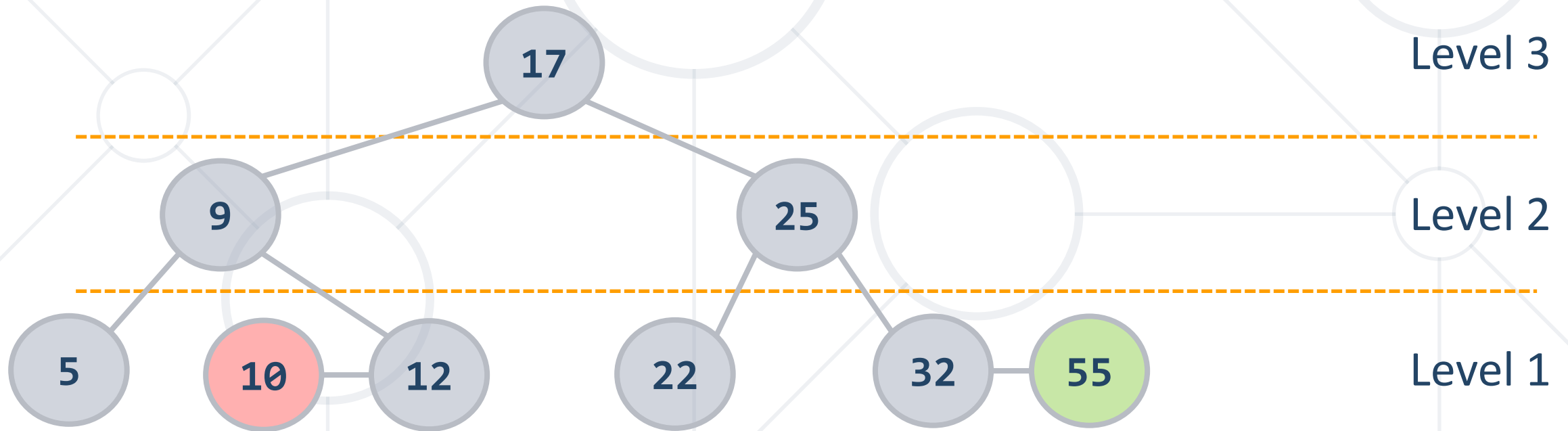
- Utilizes the concept of **levels**
- **Level** - the **number of left links** on the path to a **null** node



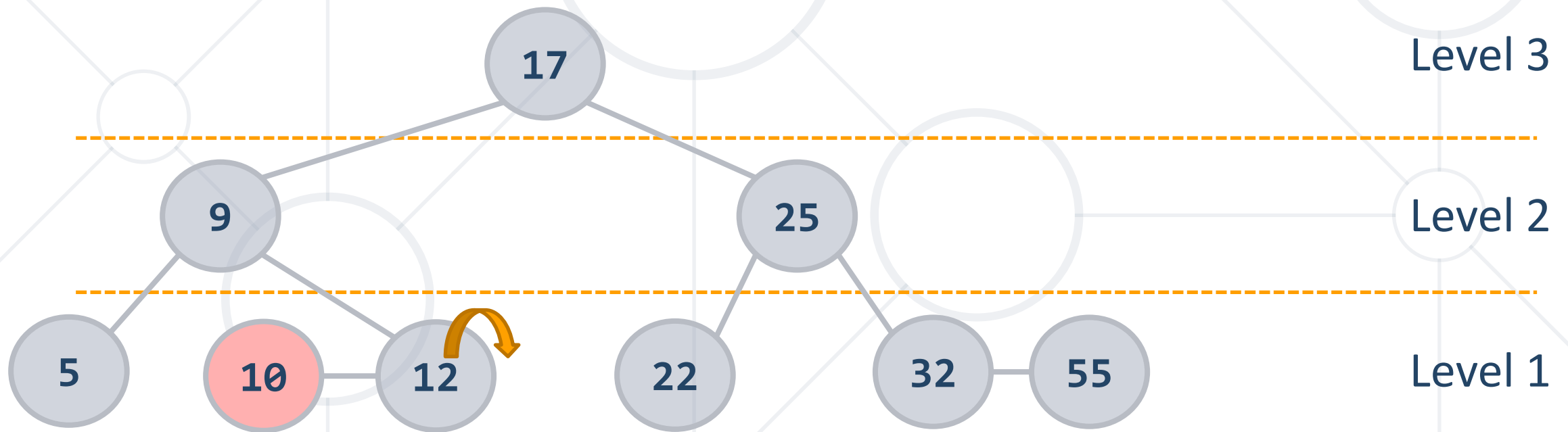
- AA tree invariants
  - The **level** of every **leaf node** is **one**
  - Every **left child** has **level one less** than its **parent**
  - Every **right child** has **level equal** to or **one less** than its **parent**
  - **Right grandchildren** have **levels less** than their **grandparents**
  - Every node of level greater than one **has two children**



- **Right** horizontal links **are possible**
- **Left** horizontal links **are not allowed**

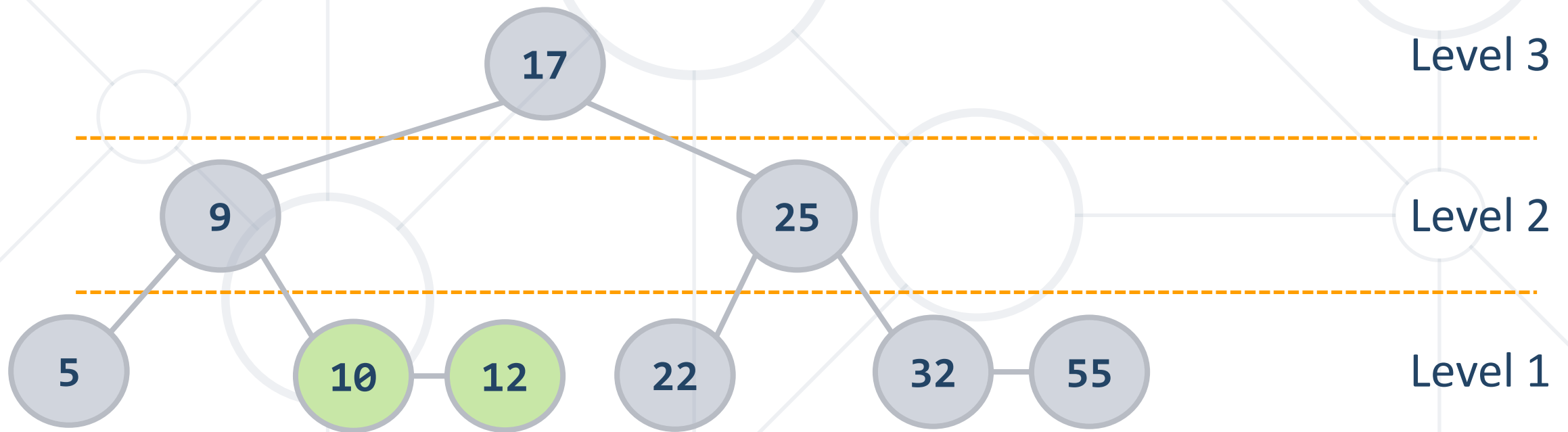


- Skew operation is a single **right** rotation
- Skew when an **insertion** or **deletion** creates a horizontal **left link**

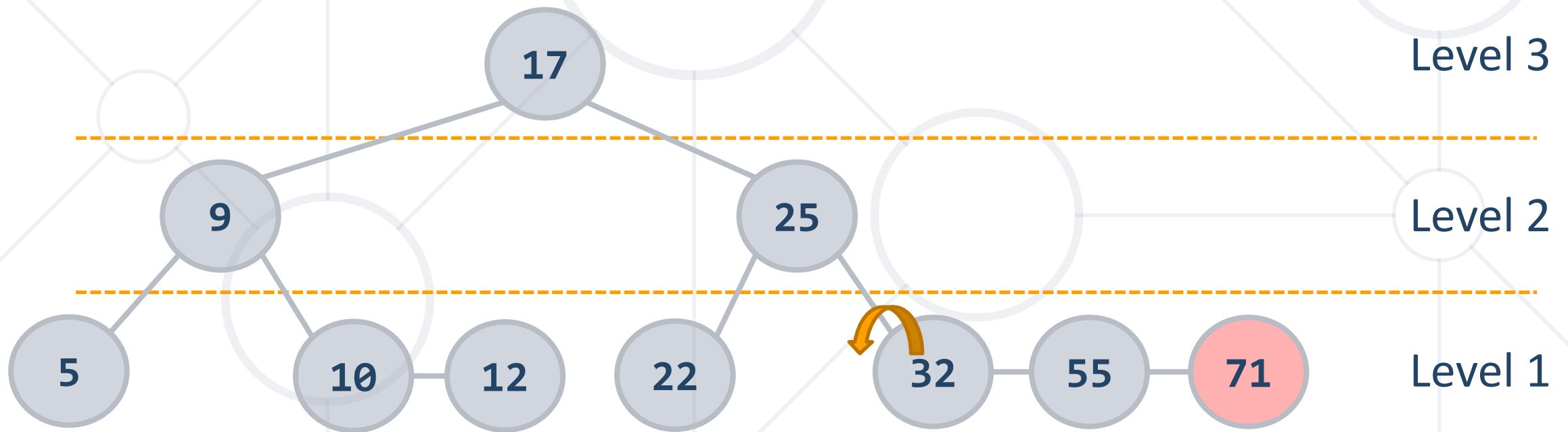


# Skew (2)

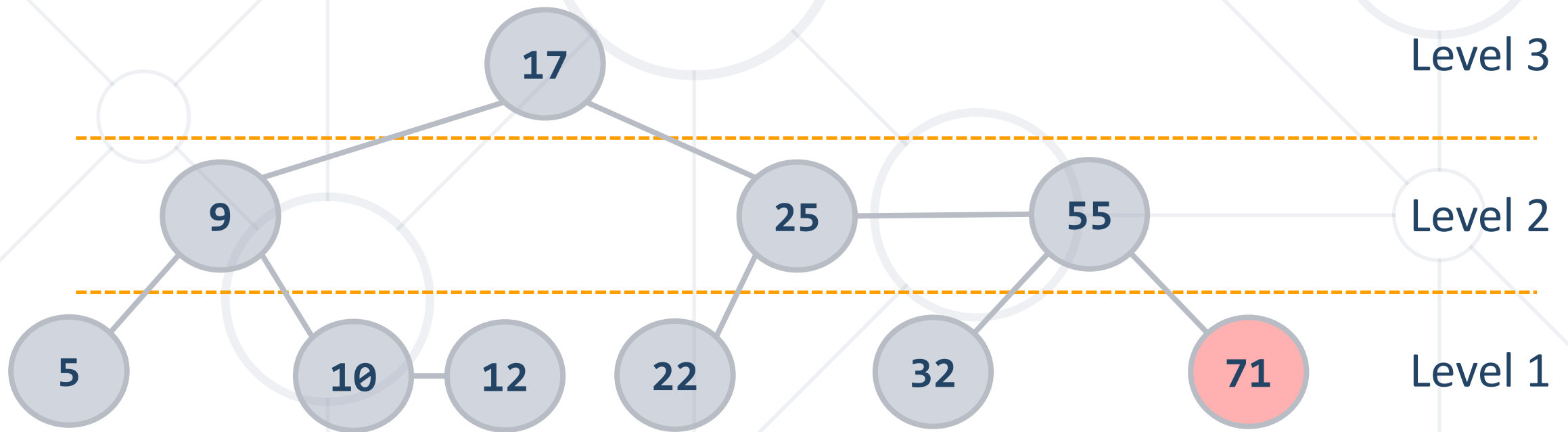
- Skew operation is a **single right rotation**
- Skew when an insertion or deletion **creates a horizontal left link**

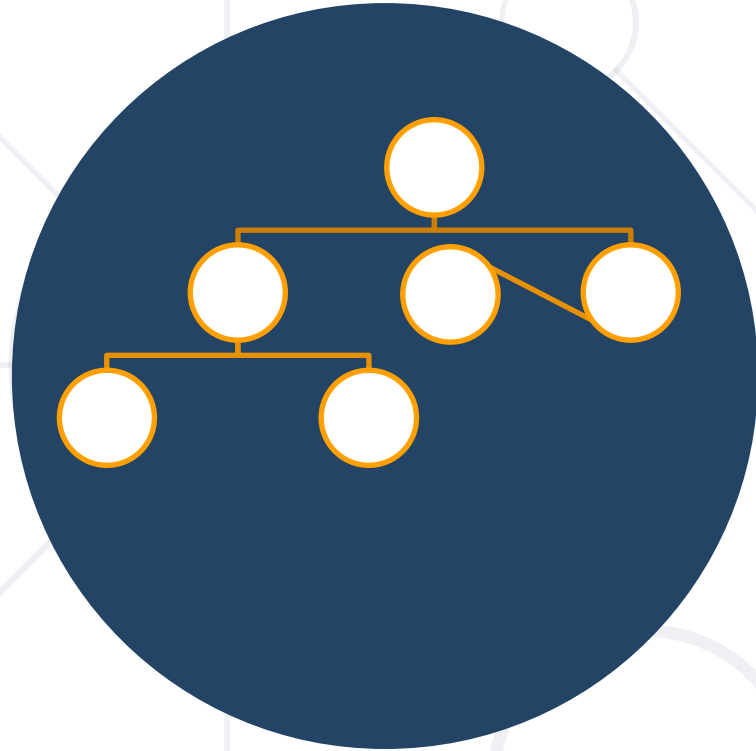


- Split operation is a **single left rotation**
- Split when an insertion or deletion **two consecutive right horizontal links**



- Split operation is a **single left rotation**
- Split when an insertion or deletion **two consecutive right horizontal links**





# AA Tree

## Insertion Algorithm

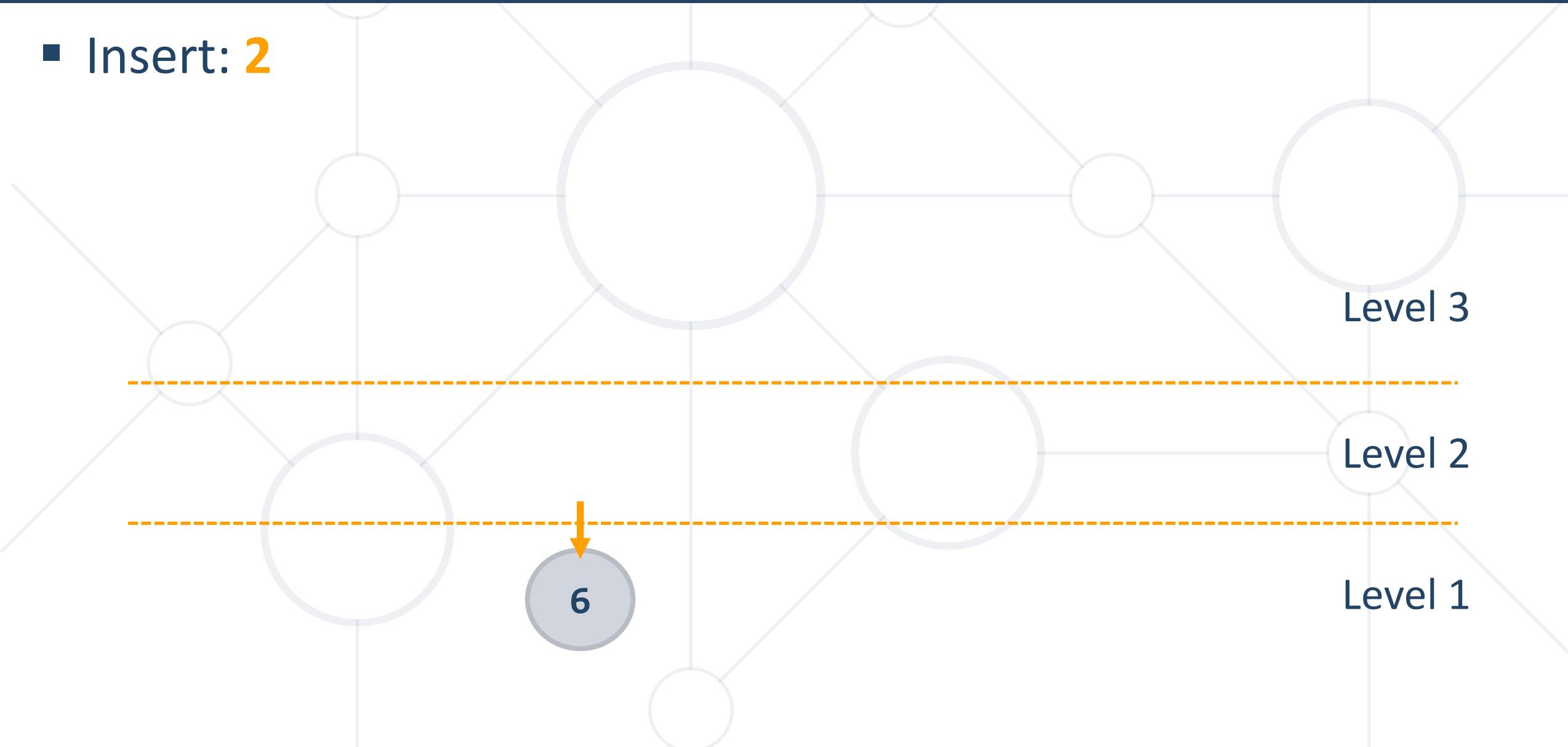
# AA Tree Insertion #1

- Insert: **6**
- **New nodes** are always inserted at **Level 1**



# AA Tree Insertion #1

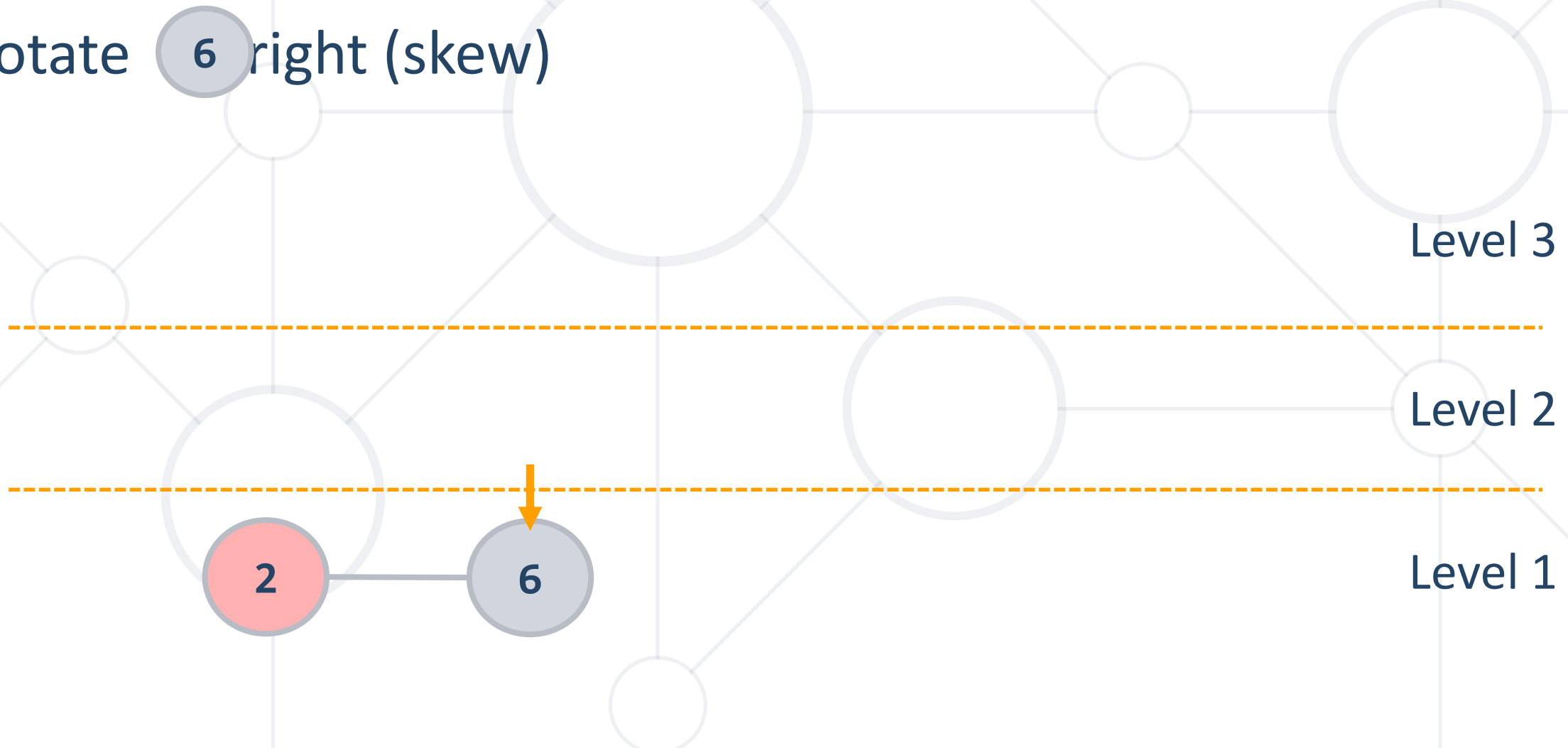
- Insert: **2**





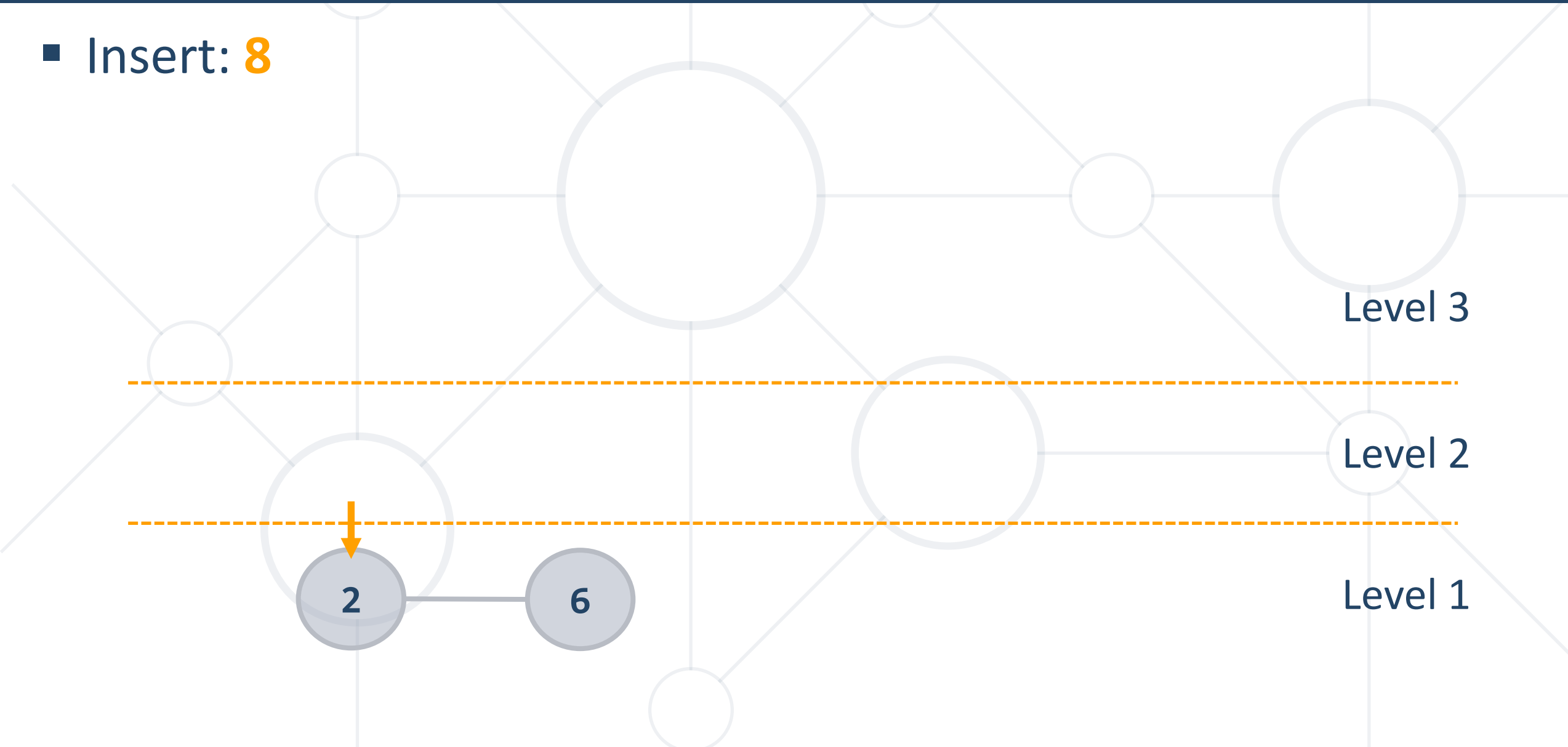
# AA Tree Insertion #1

- **Left horizontal** link is **not allowed**
- Rotate **6** right (skew)



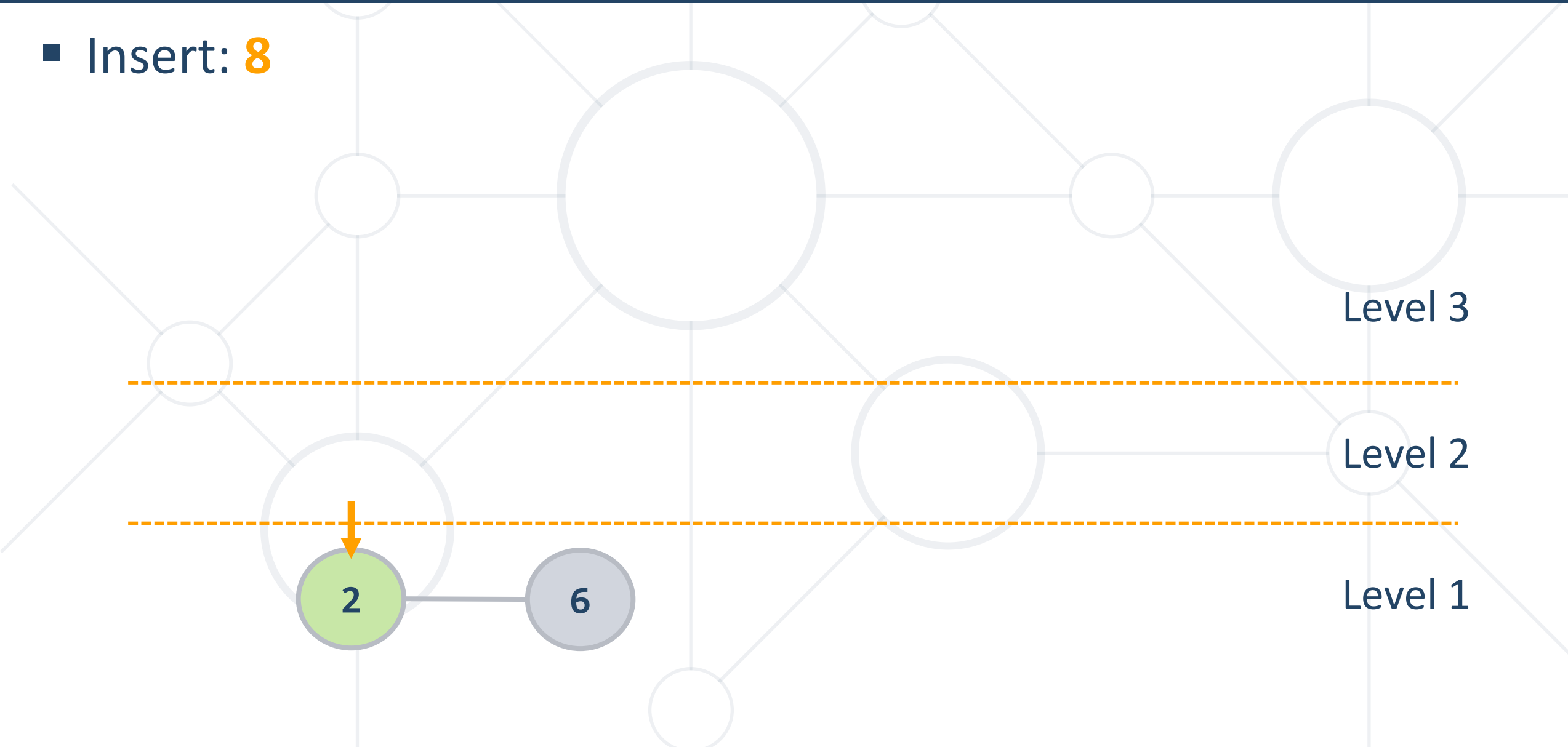
# AA Tree Insertion #1

- Insert: 8



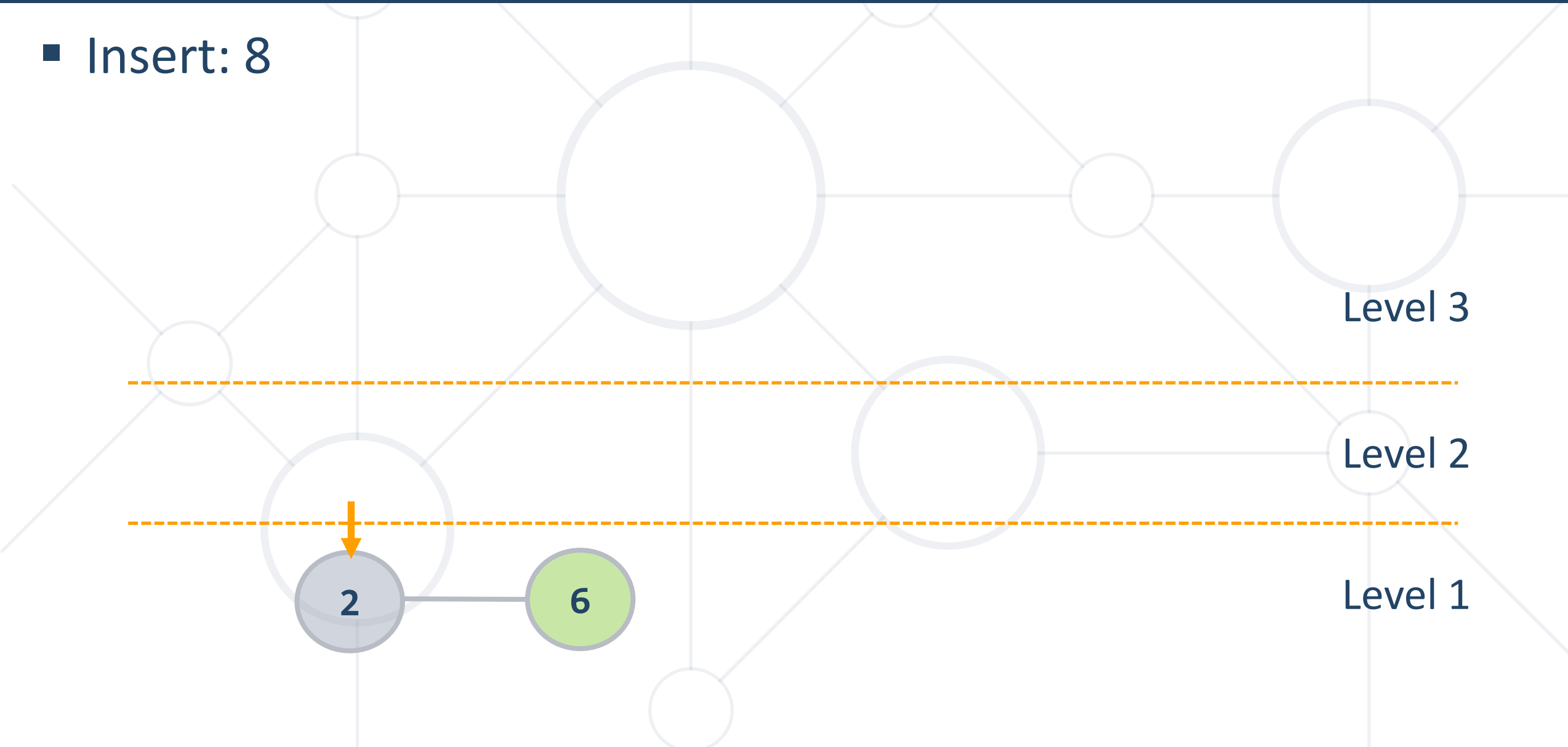
# AA Tree Insertion #1

- Insert: 8



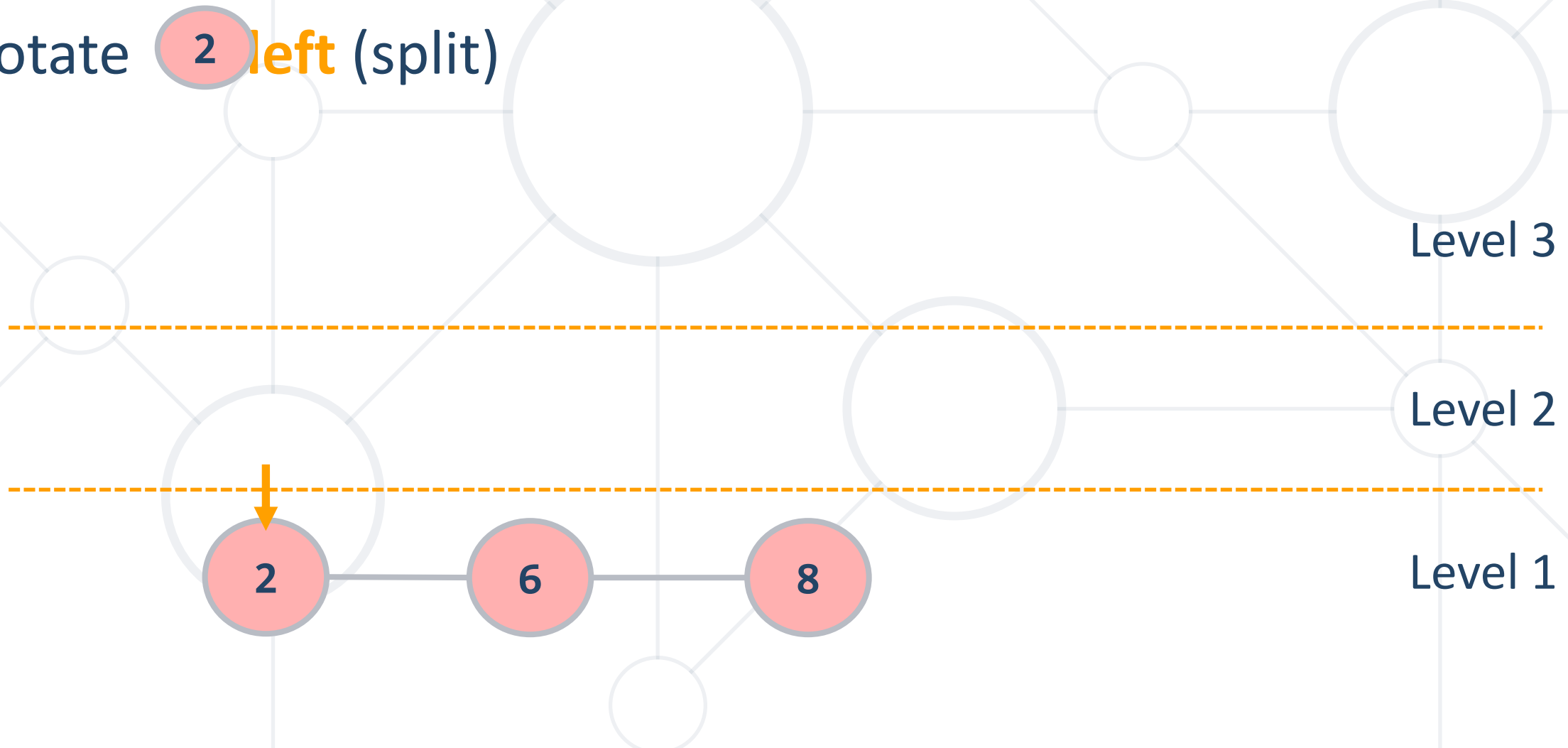
# AA Tree Insertion #1

- Insert: 8



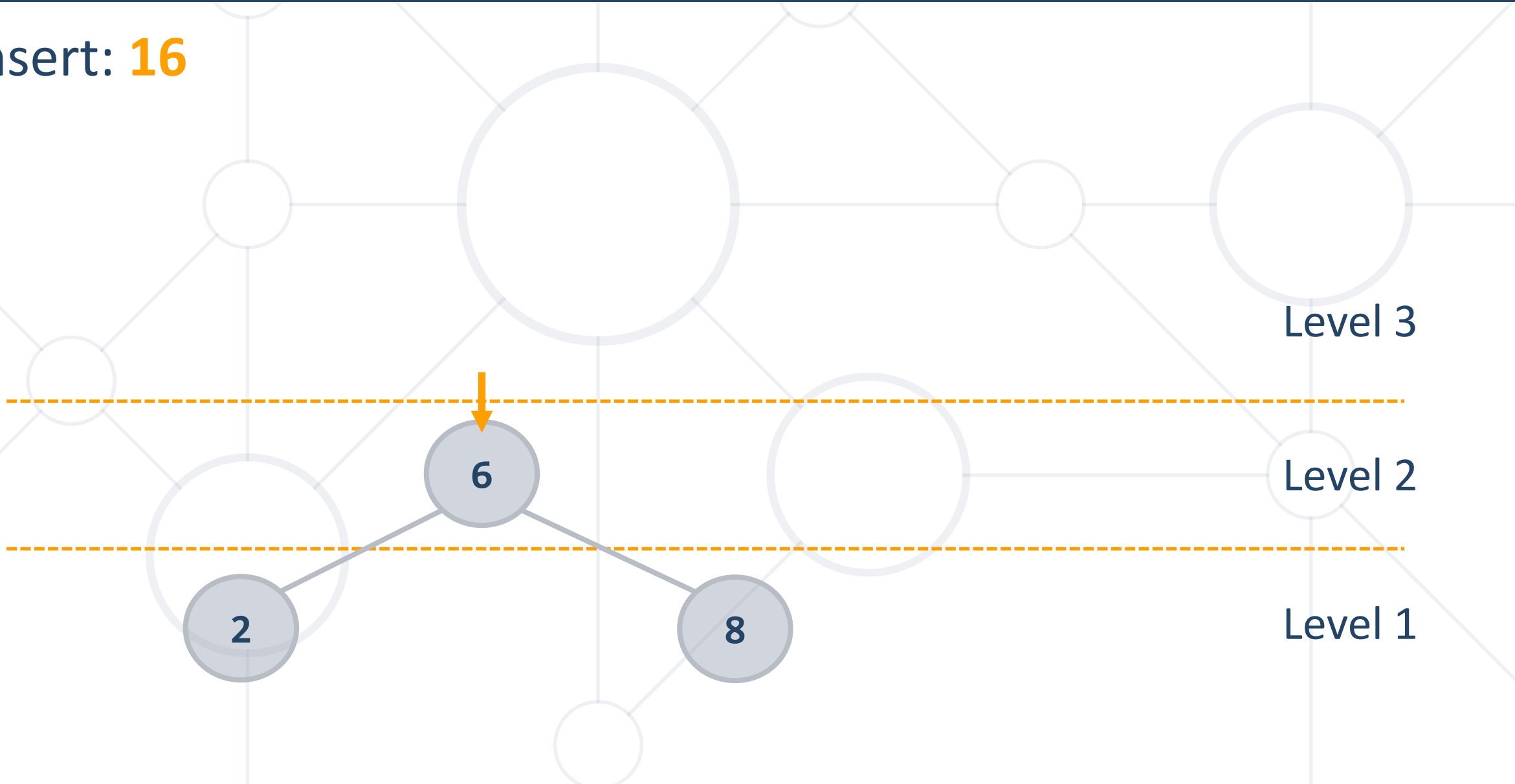
# AA Tree Insertion #1

- Two consecutive right horizontal links
- Rotate **2** left (split)



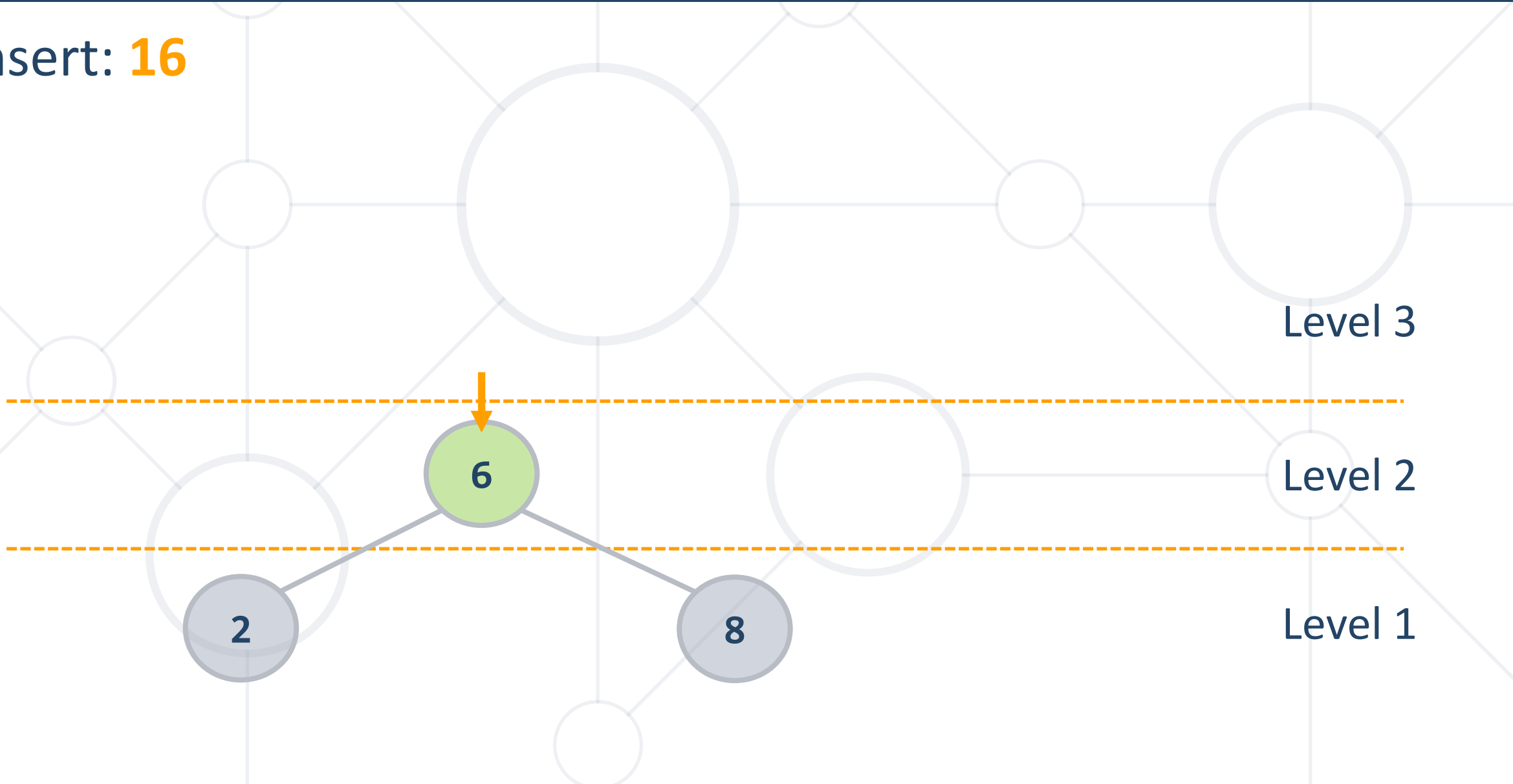
# AA Tree Insertion #1

- Insert: **16**



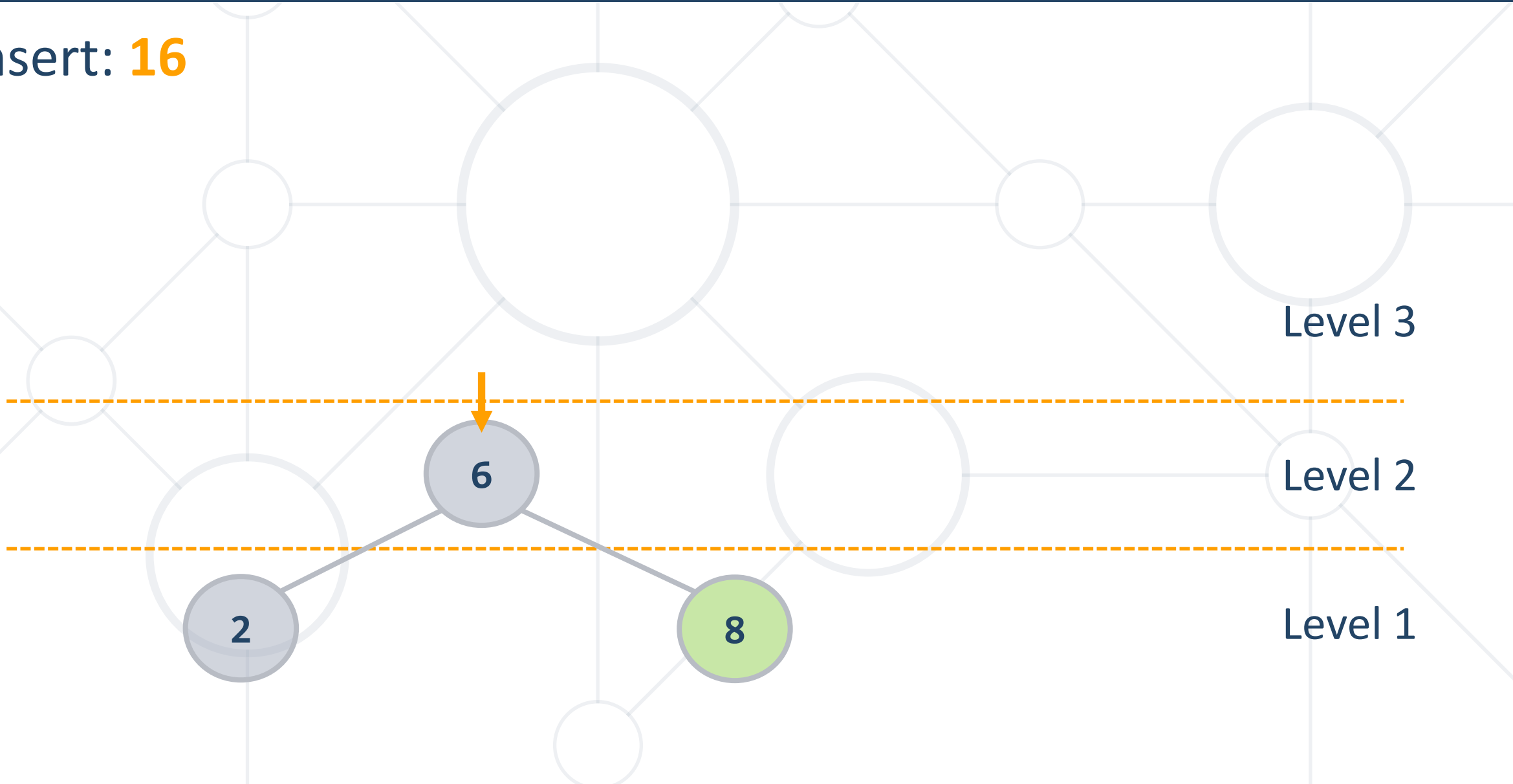
# AA Tree Insertion #1

- Insert: **16**



# AA Tree Insertion #1

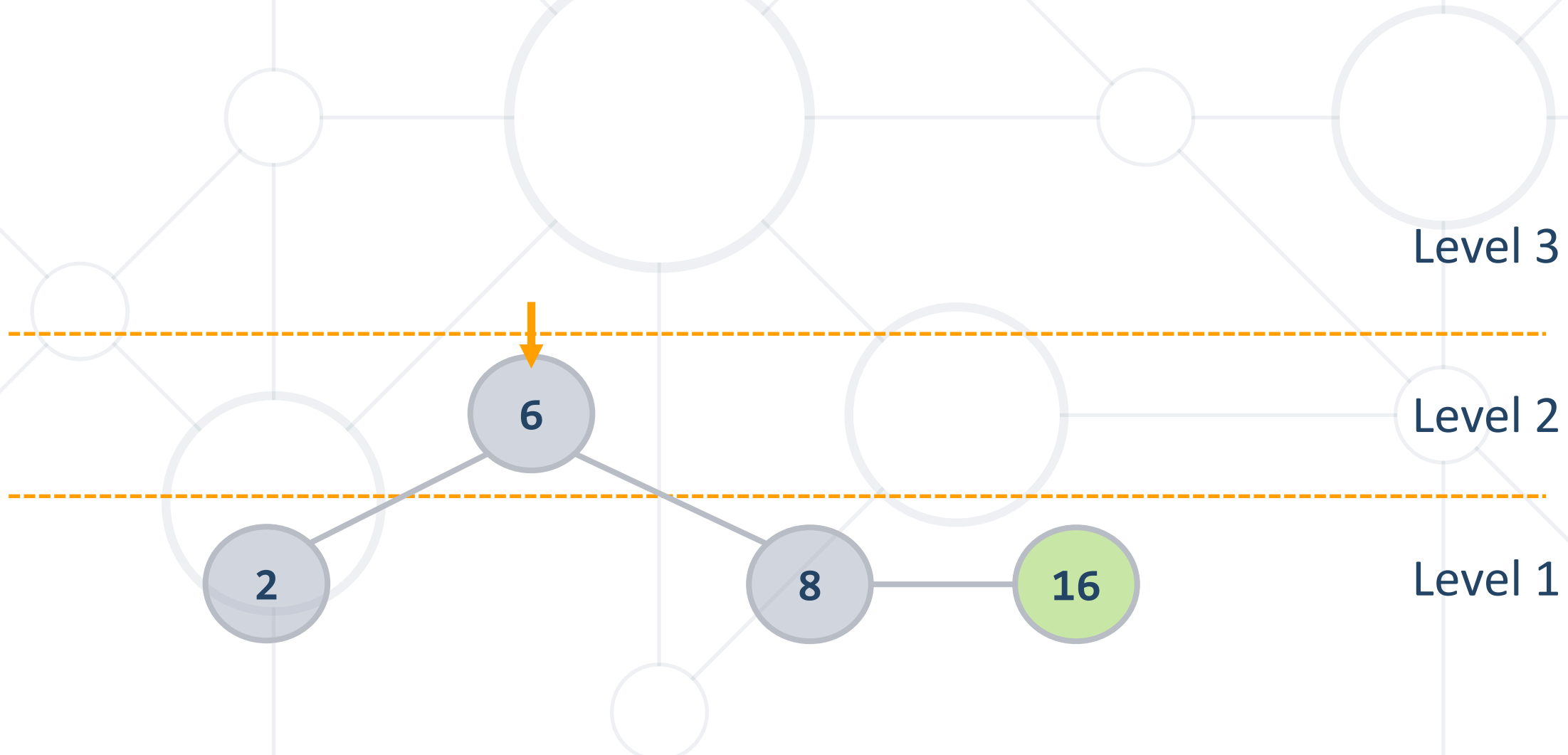
- Insert: **16**





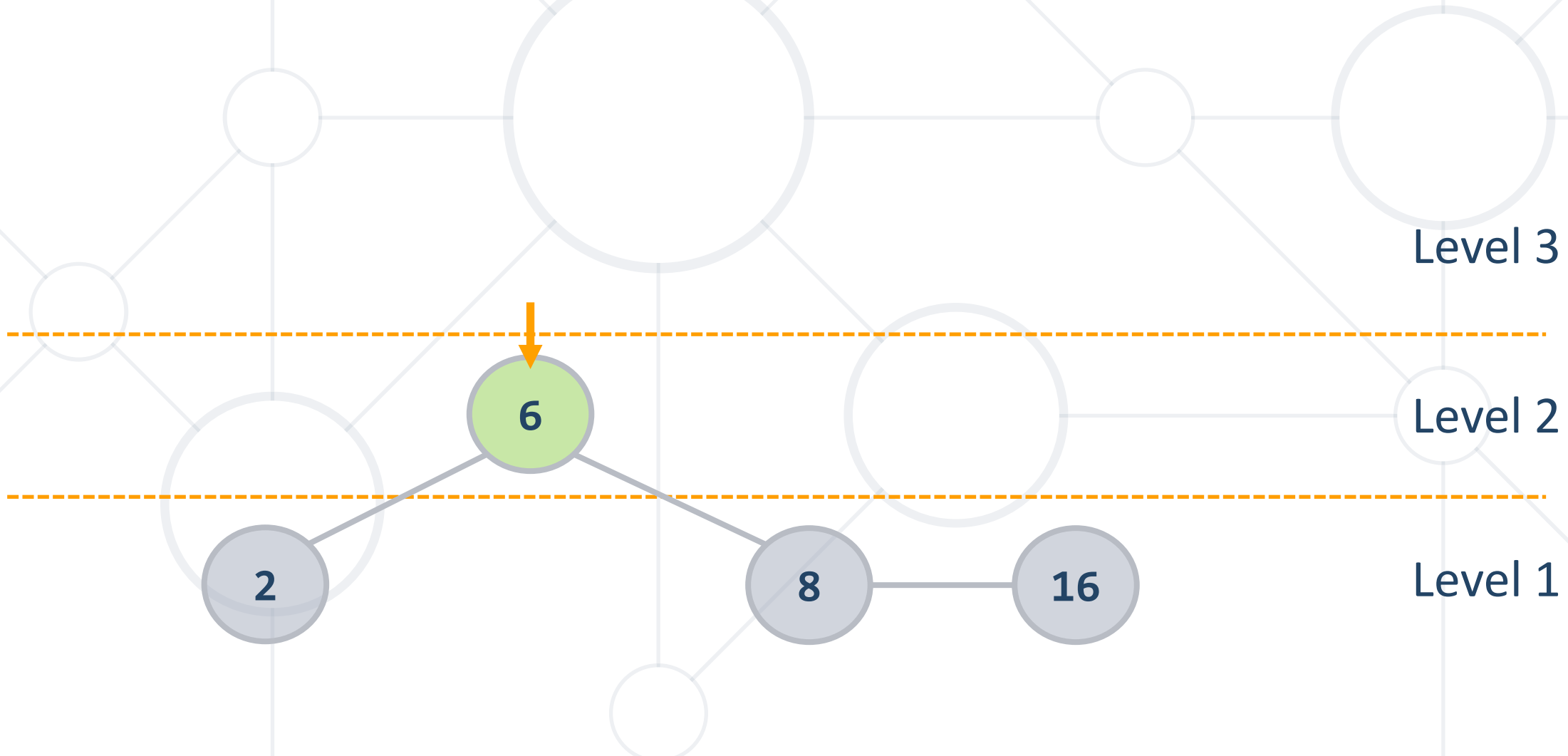
# AA Tree Insertion #1

- Insert: **16**



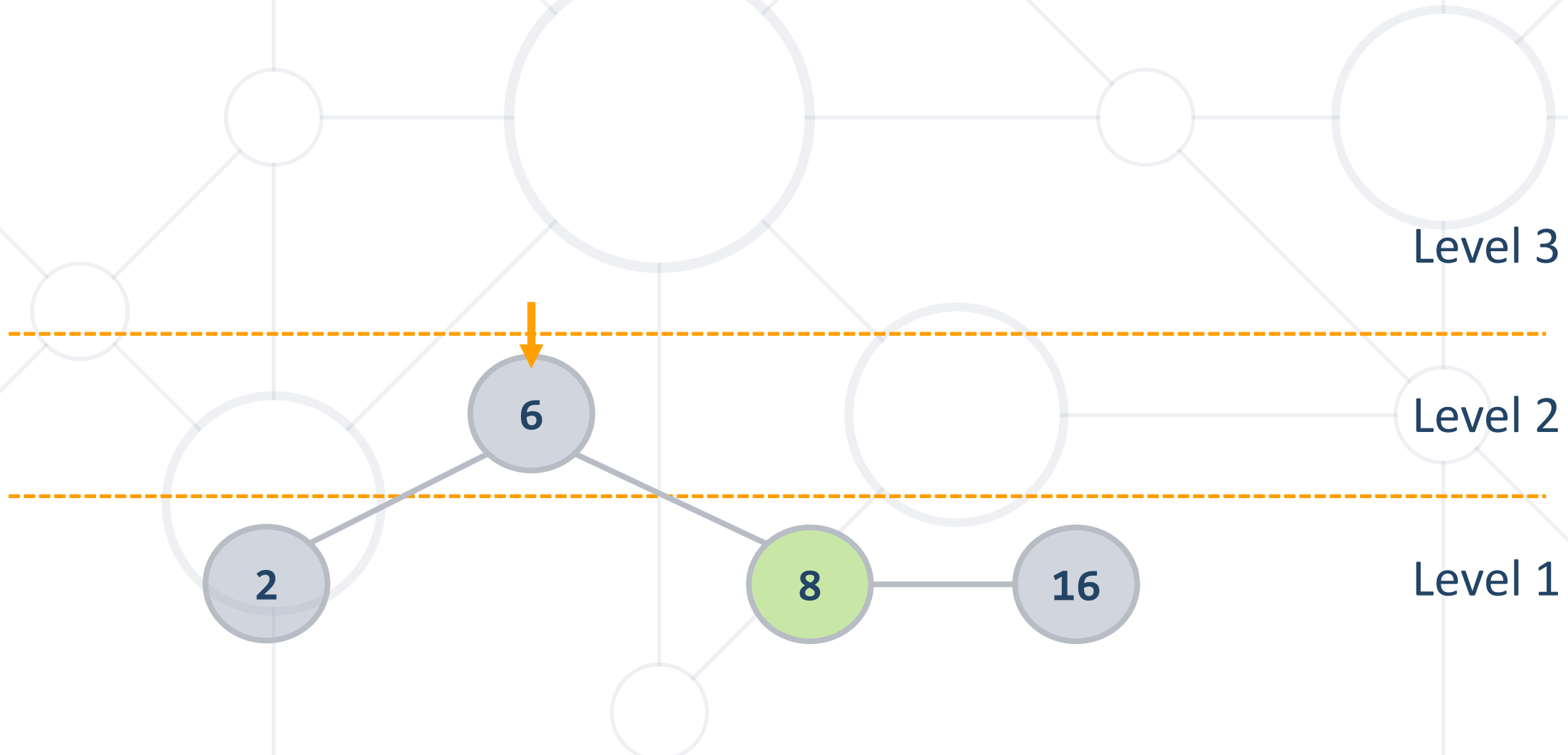
# AA Tree Insertion #1

- Insert: **10**



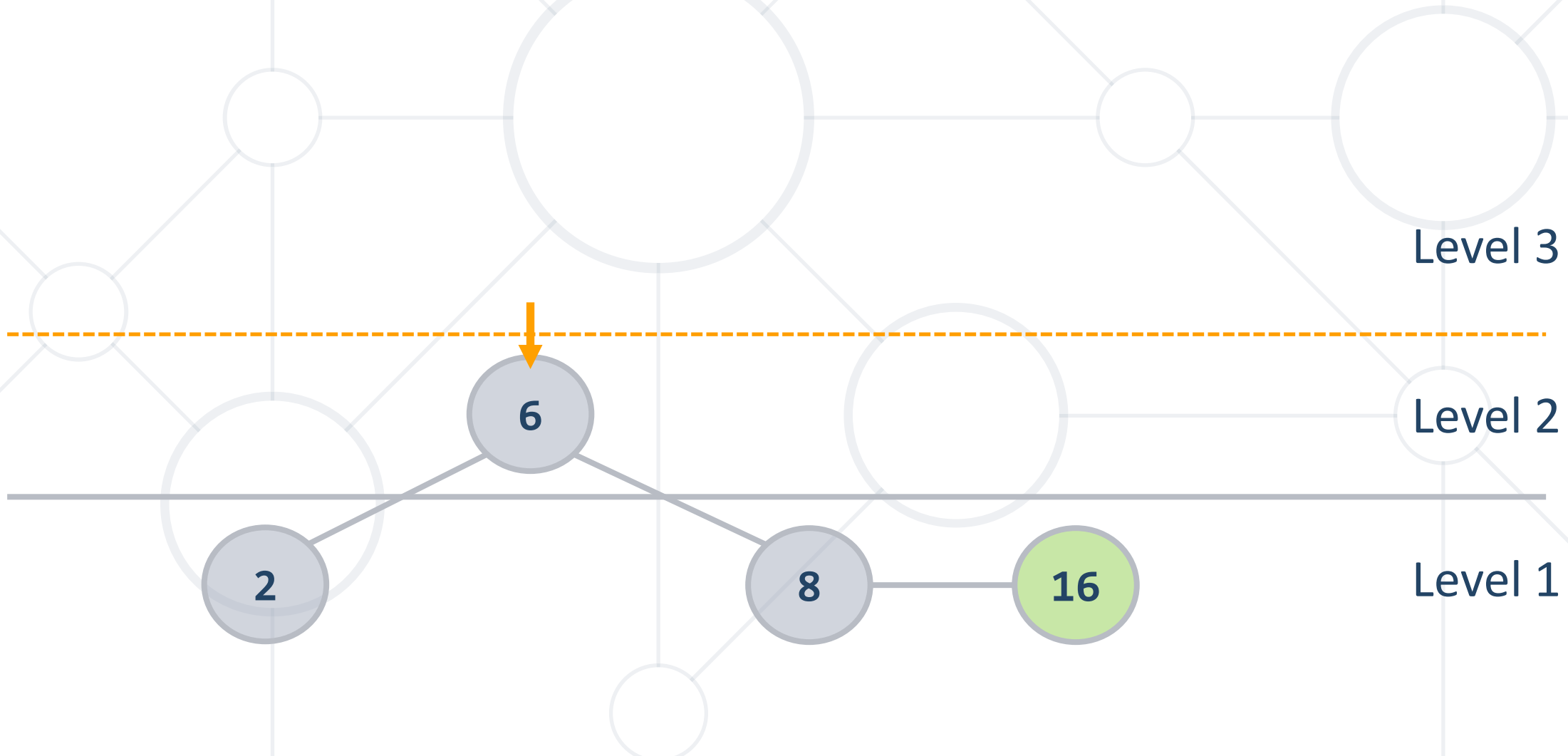
# AA Tree Insertion #1

- Insert: 10



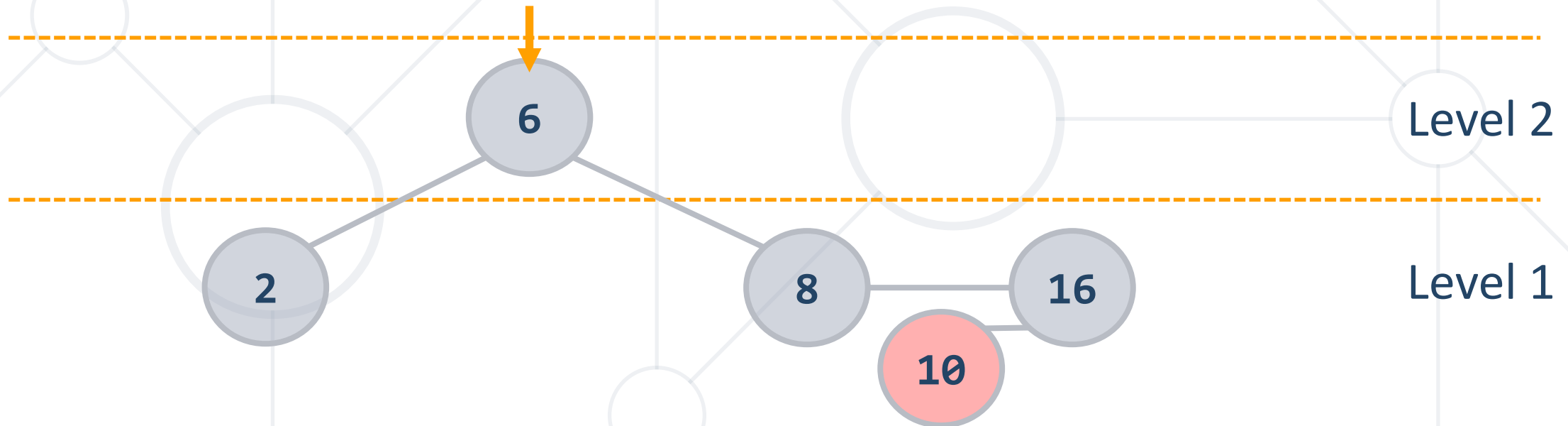
# AA Tree Insertion #1

- Insert: **10**



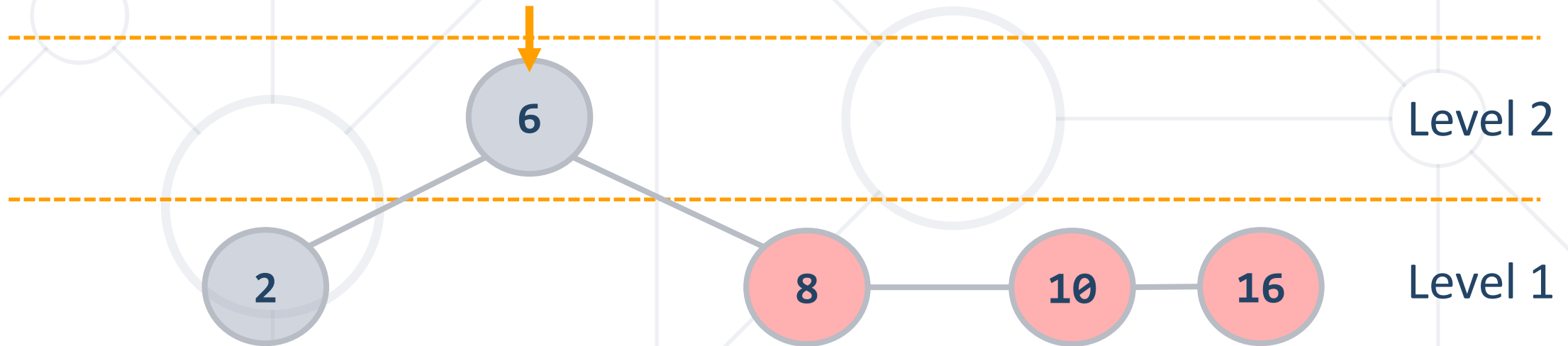
# AA Tree Insertion #1

- Horizontal left link not allowed
- Rotate **16** right (skew)



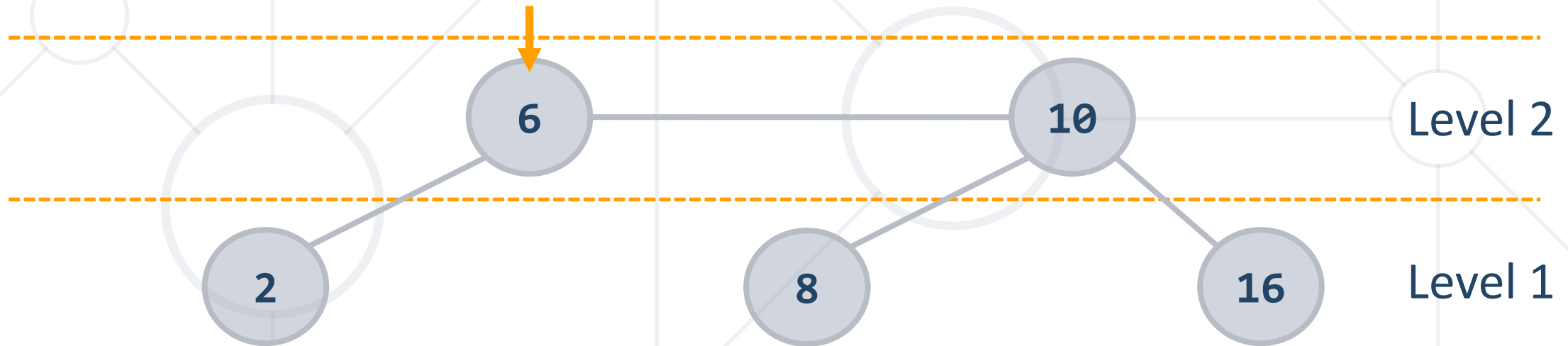
# AA Tree Insertion #1

- Two consecutive right horizontal links
- Rotate **8** left (split)



# AA Tree Insertion #1

- Two consecutive right horizontal links
- Rotate **8** left (split)




TIME'S

- Consider **web application** in which
  - **searches** are **far more frequent** than **insertions/deletions**
- Which of the following do you prefer:
  - AVL
  - Linked List
  - Red-Black
  - B+



TIME'S UP!

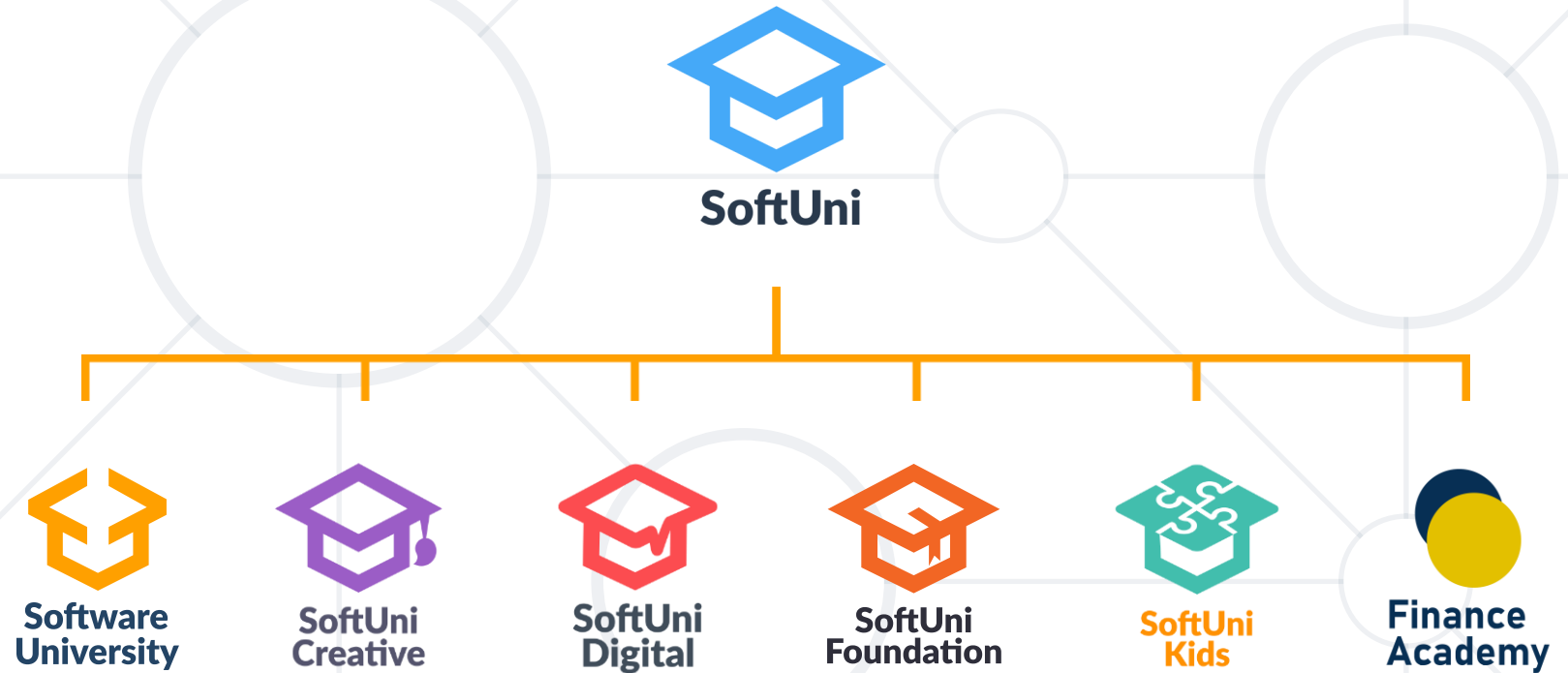
- Consider web application in which
  - searches are far more frequent than insertions/deletions
- Which of the following do you prefer:
  - **AVL** 
  - Linked List
  - Red-Black
  - B+

AVL trees are more rigidly balanced, so they have faster search

- Red-Black Trees are widely used
  - Insertion is easy
  - Balance by color
  - Color is a single byte
- AA Trees
  - Insertion algorithm



# Questions?



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