Hash Tables, Sets and Dictionaries

Hashing and Collisions

0	1	2	…	m-1
null	null	SoftUni	…	Java

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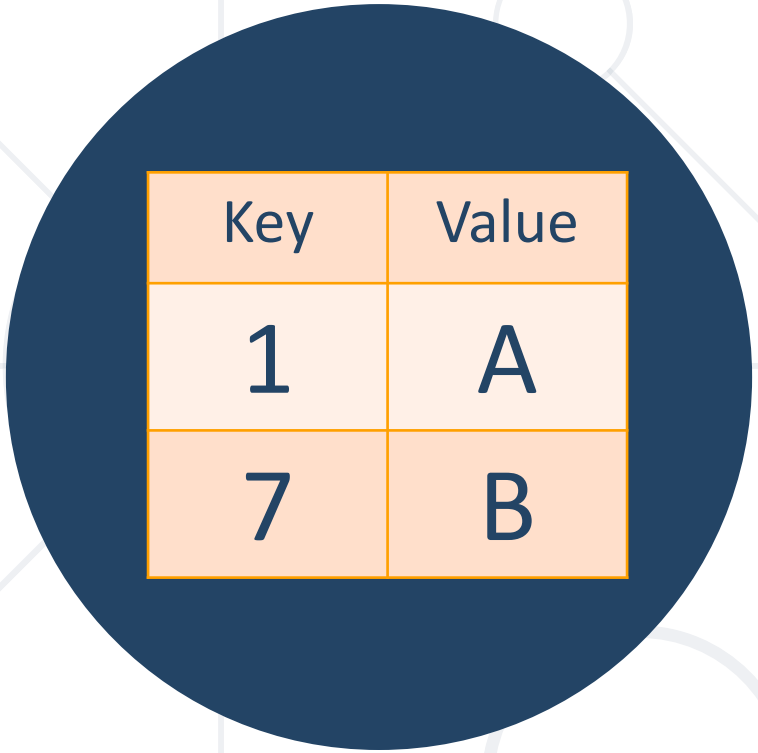


Hash Tables

Hashing and Collision Resolution

Hash Function

- Given a key of any type, convert it to an integer


```
class Person {
  String firstName;
  String lastName;
  int age;
}
```

Hash Function (2)

```
class Person {
  String firstName;
  String lastName;
  int age;

  @Override
  public int hashCode() {
      Hash Function
  }
}
```

Hash Function (3)

```
class Person {
  String firstName;
  String lastName;
  int age;

  @Override
  public int hashCode() {
    return firstName.getHashCode()
           + lastName.getHashCode()
           + Integer.hashCode(age);
  }
}
```

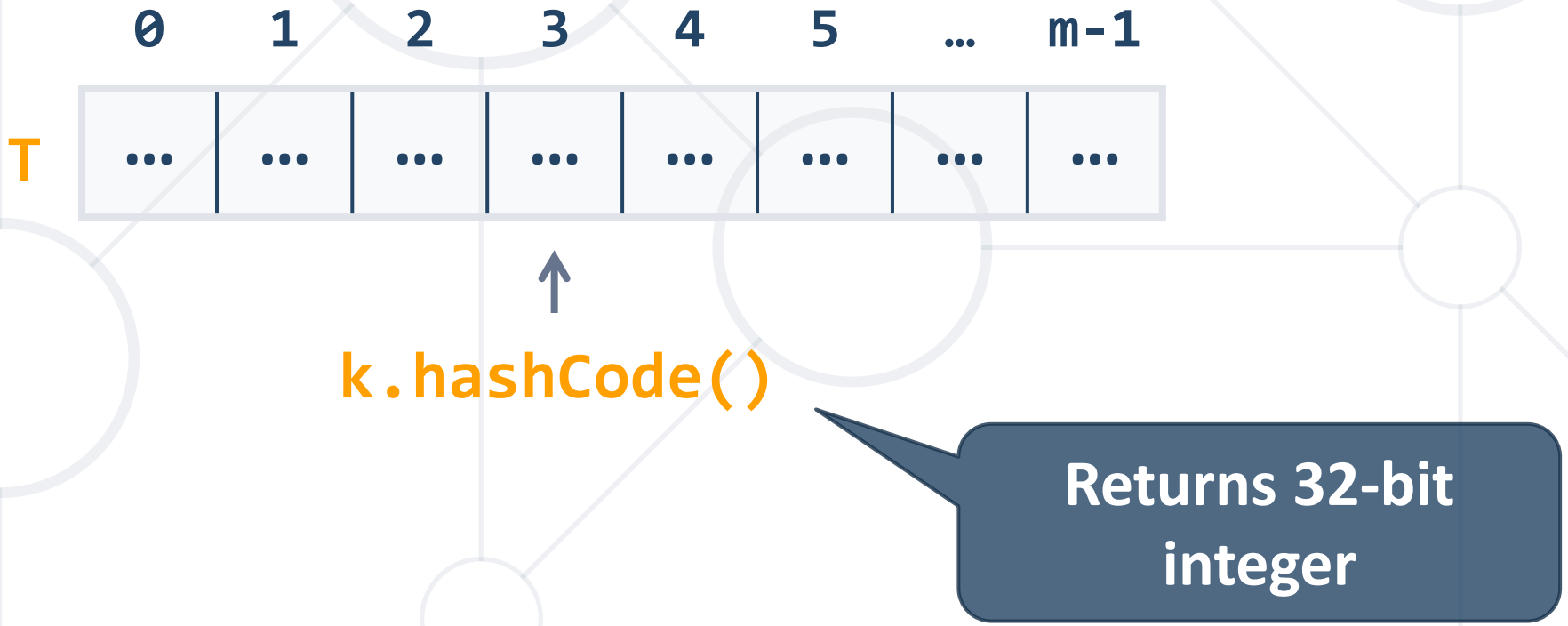
Hash Table

- A hash table is an array that holds a set of **{key, value} pairs**
- The process of mapping a key to a position in a table is called **hashing**

Hash table of size **m**

Hash Functions and Hashing

- A hash table has **m** slots, indexed from **0** to **m-1**
- A hash function converts **keys** into array indices



Hashing Functions

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- Perfect hashing function (PHF)
  - **h(k)**: one-to-one mapping of each key **k** to an integer in the range **[0, m-1]**
  - The PHF maps each key to a **distinct** integer within some manageable range
- Finding a perfect hashing function is impossible in most cases

Hashing Functions (2)

- Good hashing function
  - **Consistent** - equal keys must produce the same hash value
  - **Efficient** - efficient to compute the hash
  - **Uniform** - should uniformly distribute the keys

Hash Functions - Quiz

TIME'S

- Which of the following is **not** property of a **hashCode()** for strings
  - Can return a negative integer
  - Can take time proportional to the length of the string to compute
  - A string and its reverse will have the same hash code
  - Two strings with different hash code values are different strings

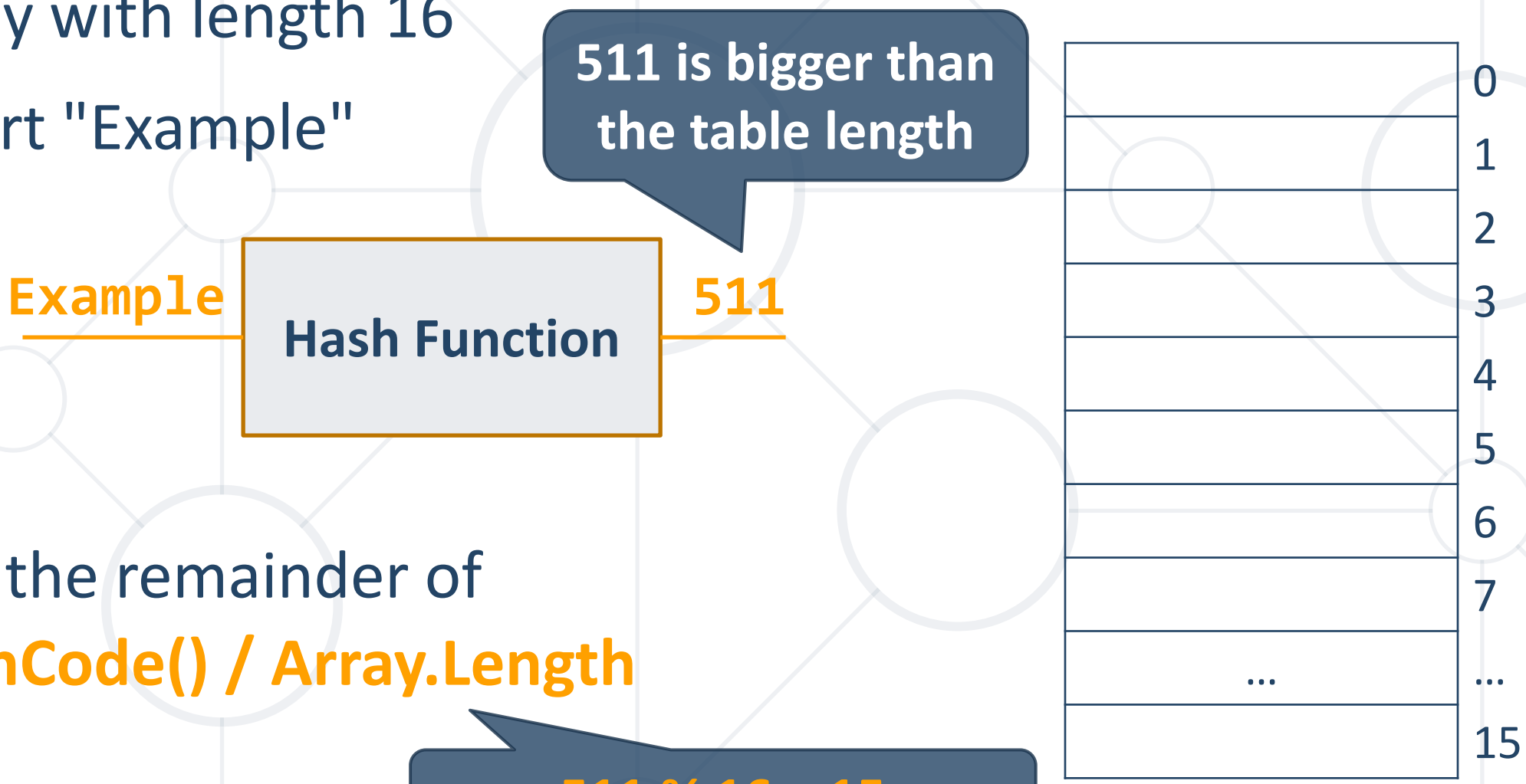
Hash Functions - Answer

- Which of the following is **not** property of a **hashCode()** for strings
  - Can return a negative integer
  - Can take time proportional to the length of the string to compute
  - A string and its reverse will have the same hash code
  - Two strings with different hash code values are different strings

```
"ab".hashCode() != "ba".hashCode()
```

Modular Hashing

- Array with length 16
- Insert "Example"



- Use the remainder of **hashCode() / Array.Length**

511 % 16 = 15

Adding to Hash Table

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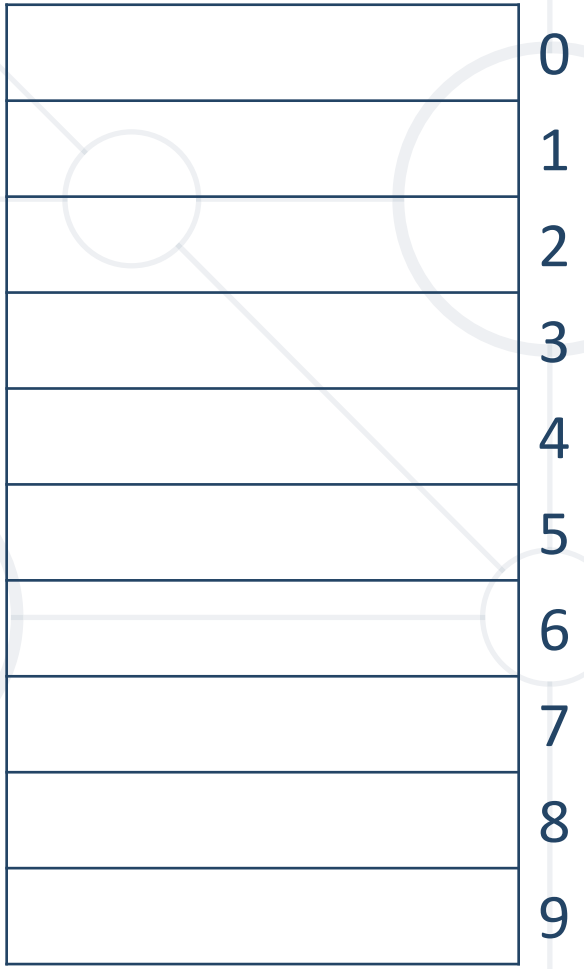
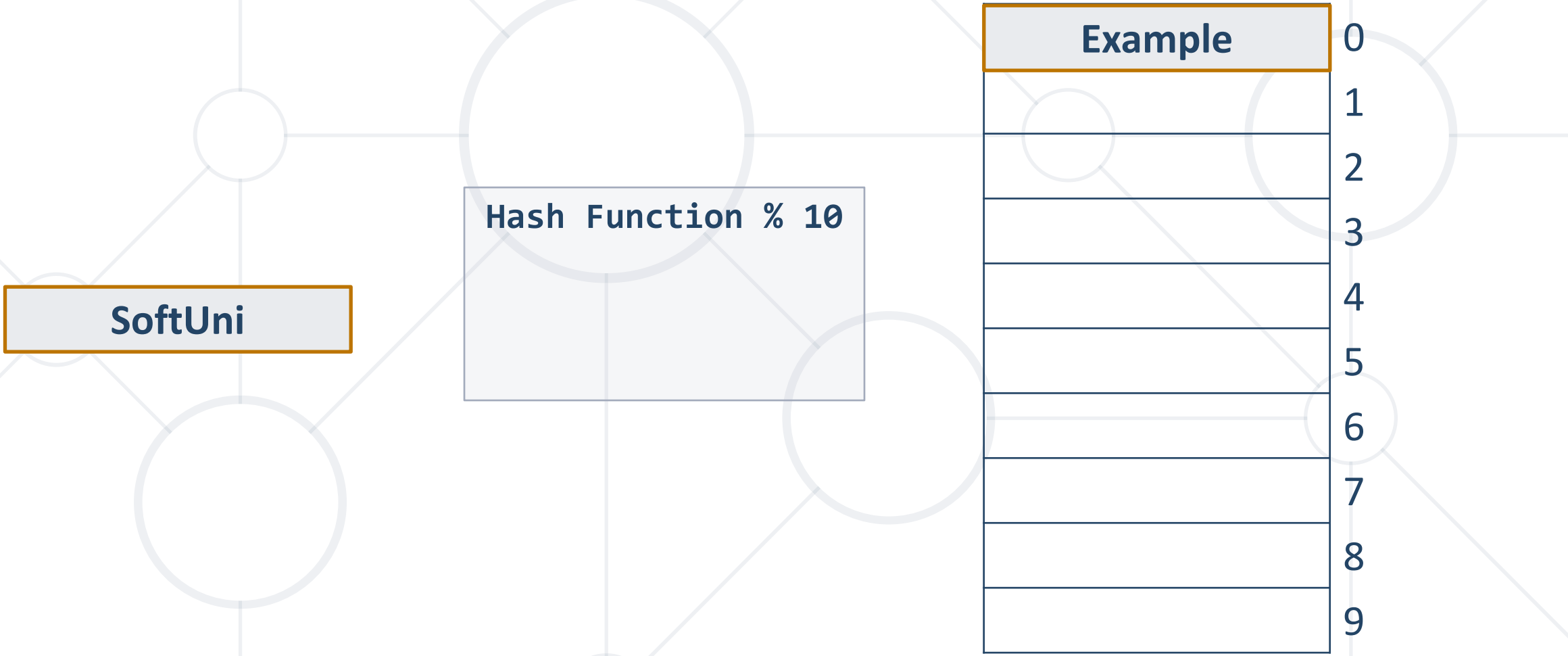
Example

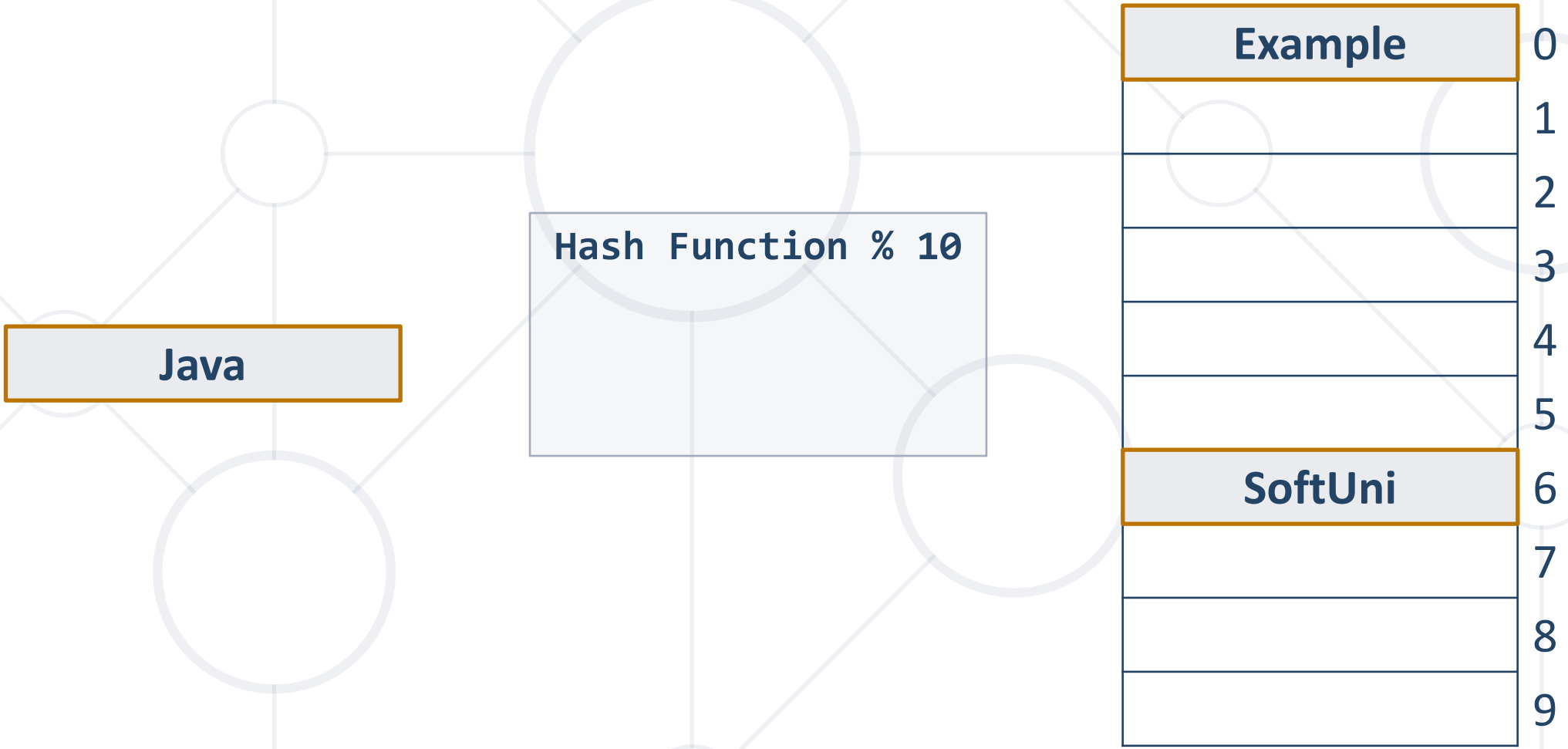
```
Hash Function % 10
```



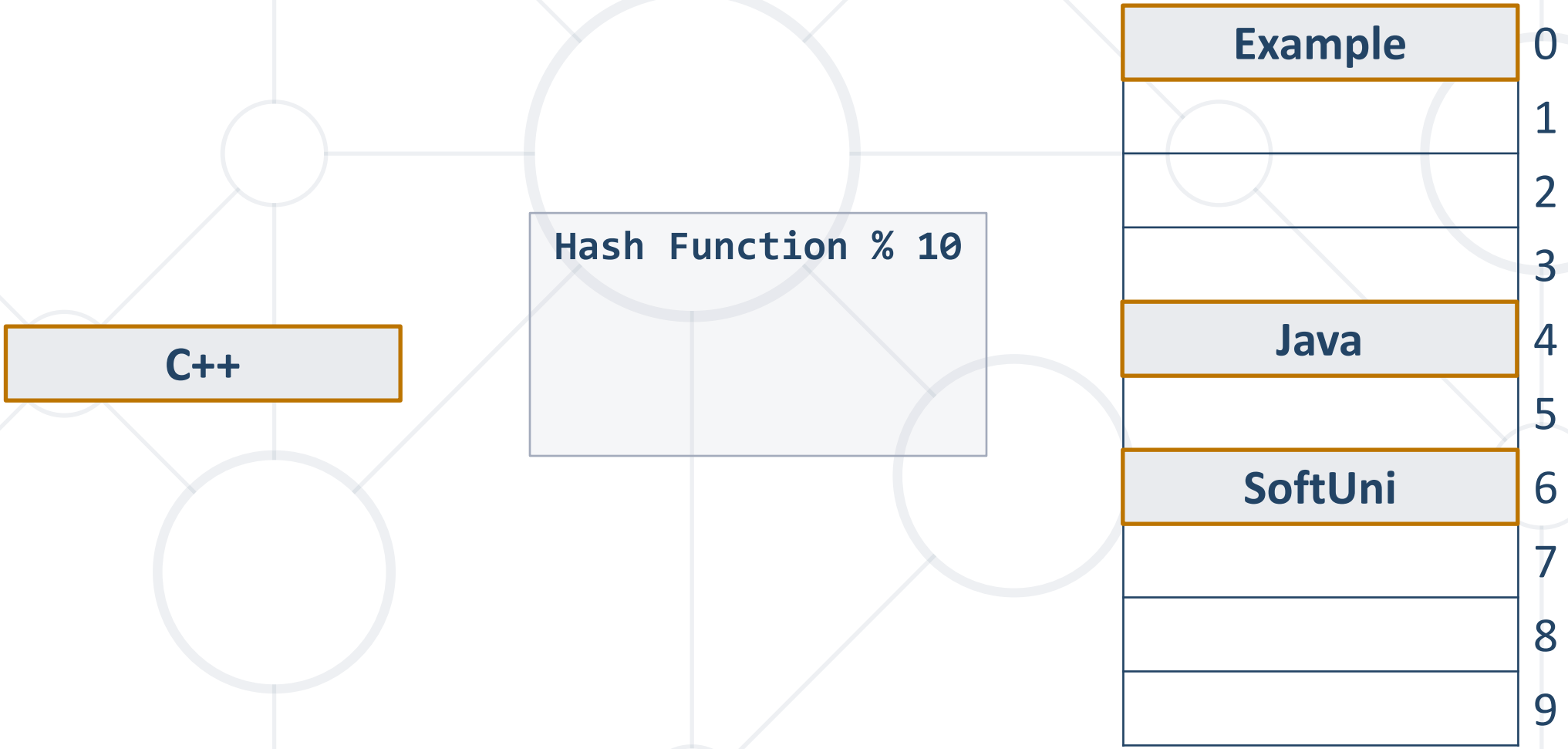
Adding to Hash Table (2)



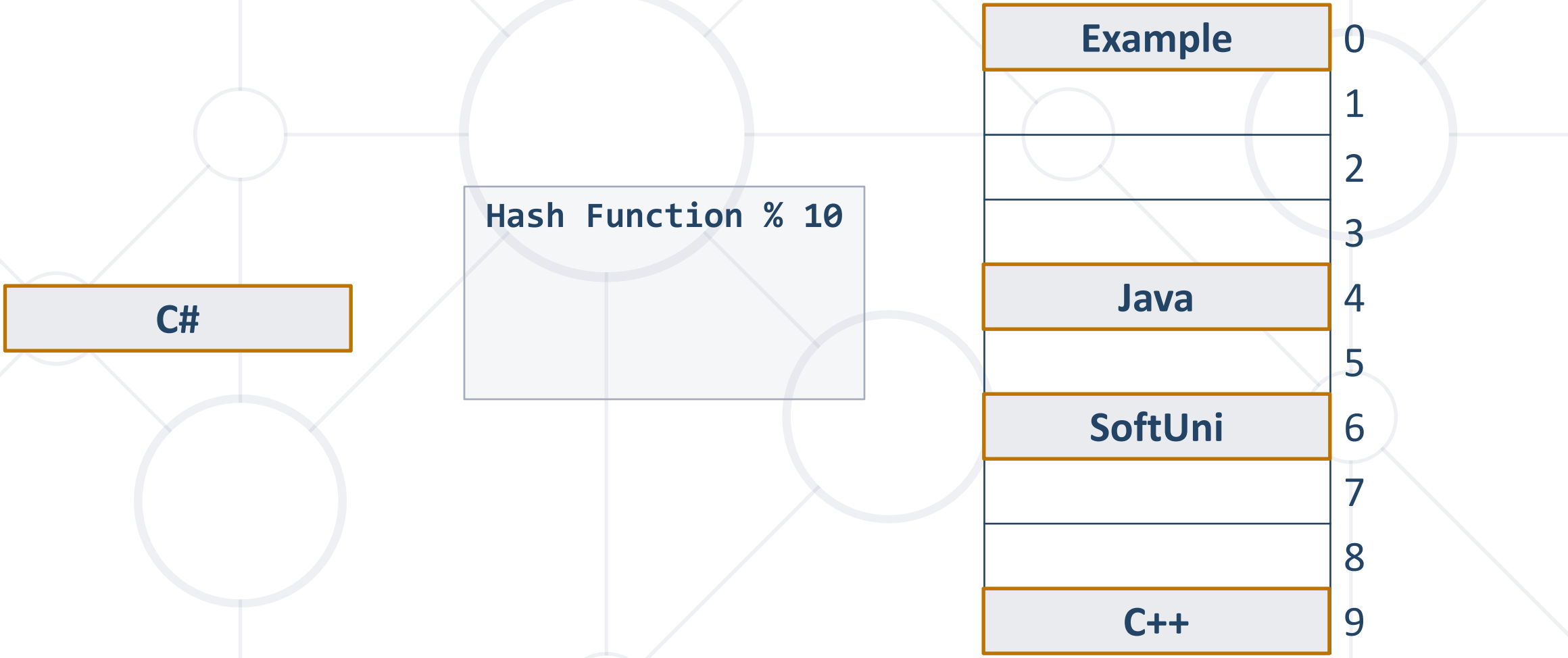
Adding to Hash Table (3)



Adding to Hash Table (4)



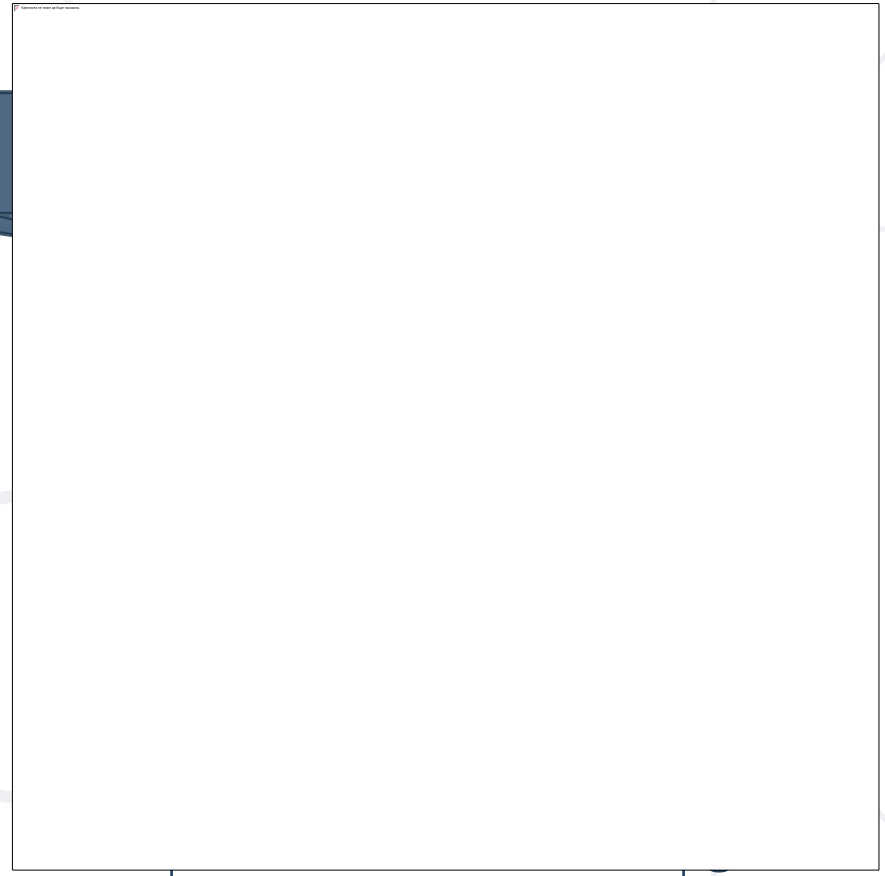
Adding to Hash Table (5)



Adding to Hash Table (6)

Collision

```
Hash Function % 10
```



C++

9

Collisions in a Hash Table

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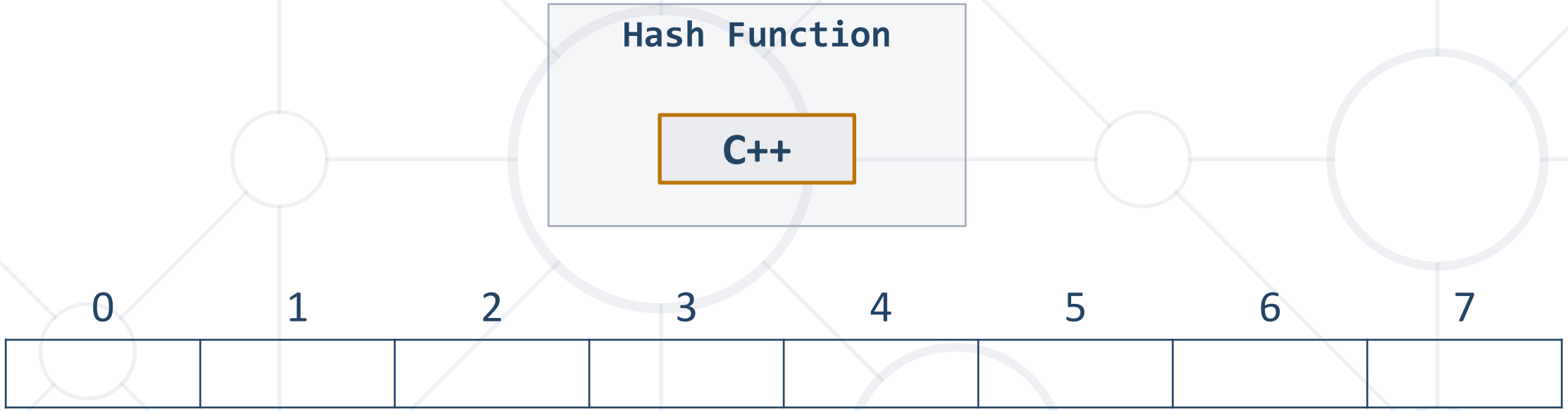
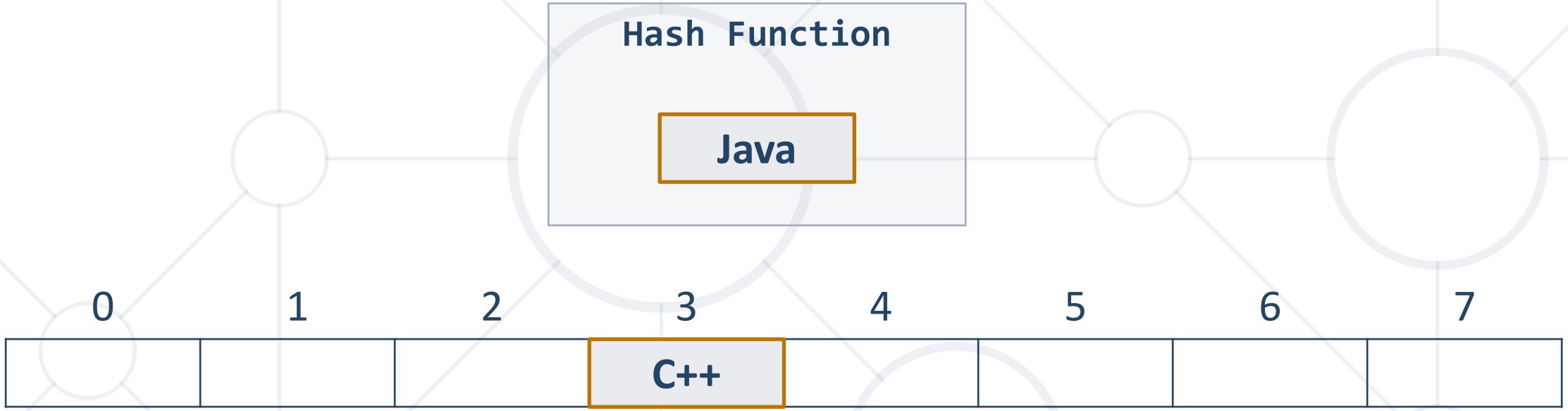
- A **collision** comes when **different keys** have the **same hash value**
  - **$h(k_1) = h(k_2)$ for $k_1 \neq k_2$**
- When the number of collisions is sufficiently small, the hash tables work quite well (fast)
- Several **collisions resolution strategies** exist
  - **Chaining** collided keys (+ values) in a list
  - Using **other slots** in the table (open addressing)
  - Cuckoo hashing
  - Many other

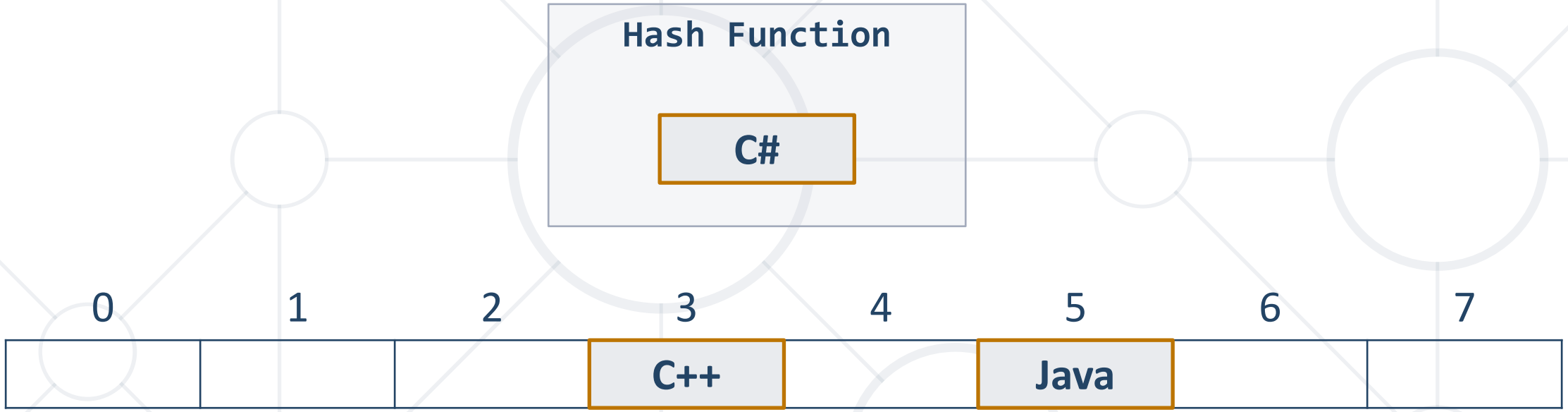
Collision Resolution: Chaining



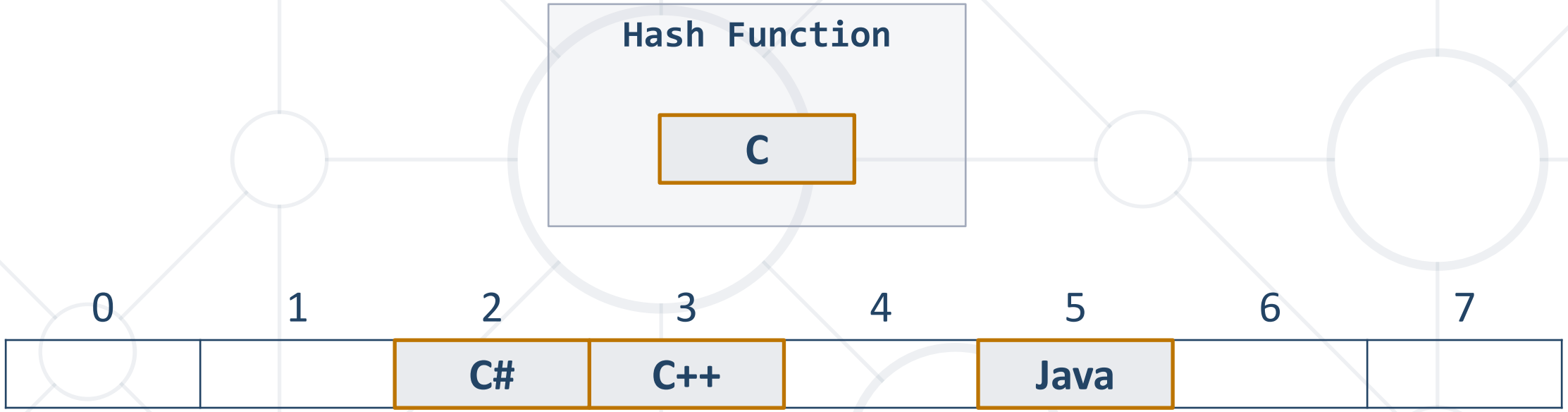
Collision Resolution: Chaining



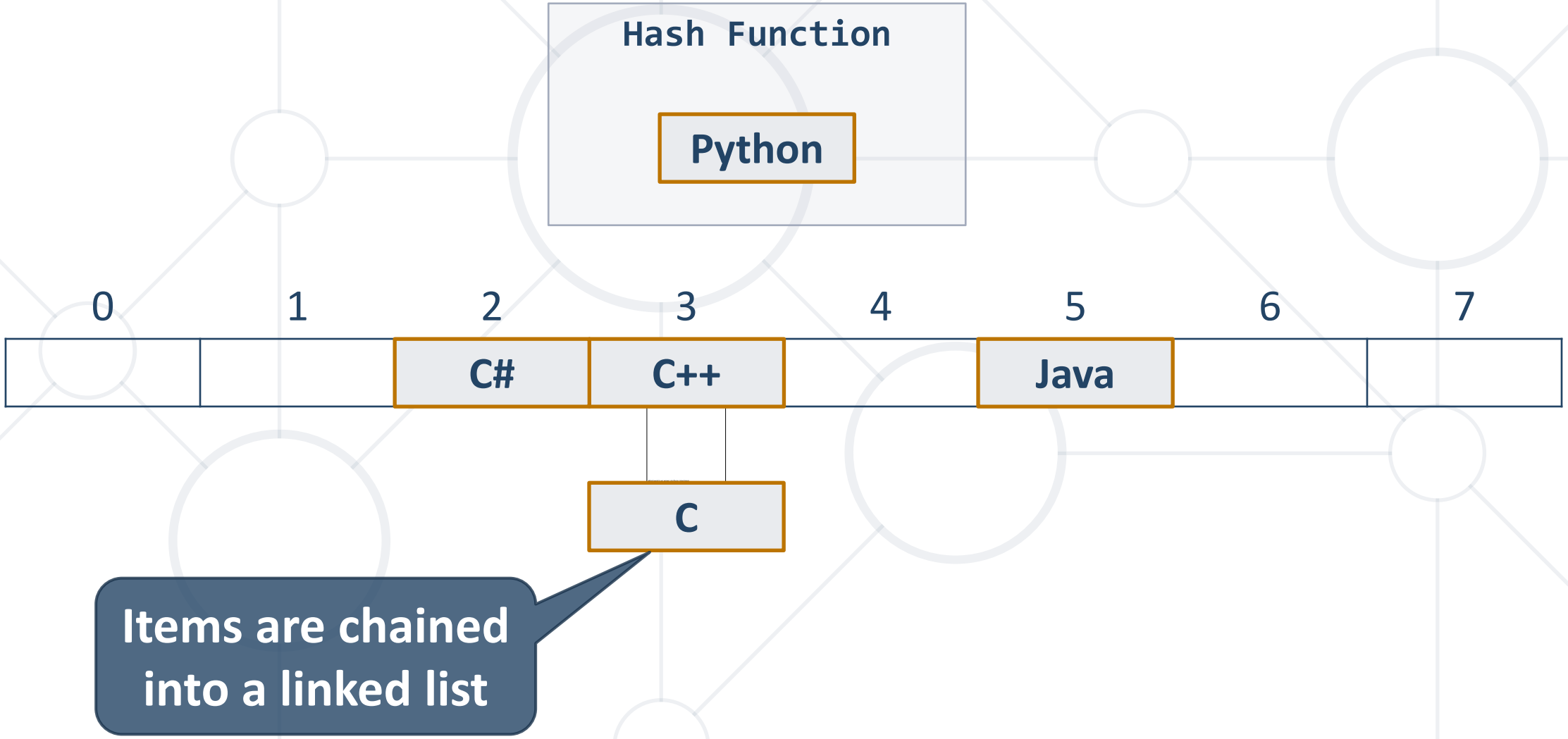
Collision Resolution: Chaining



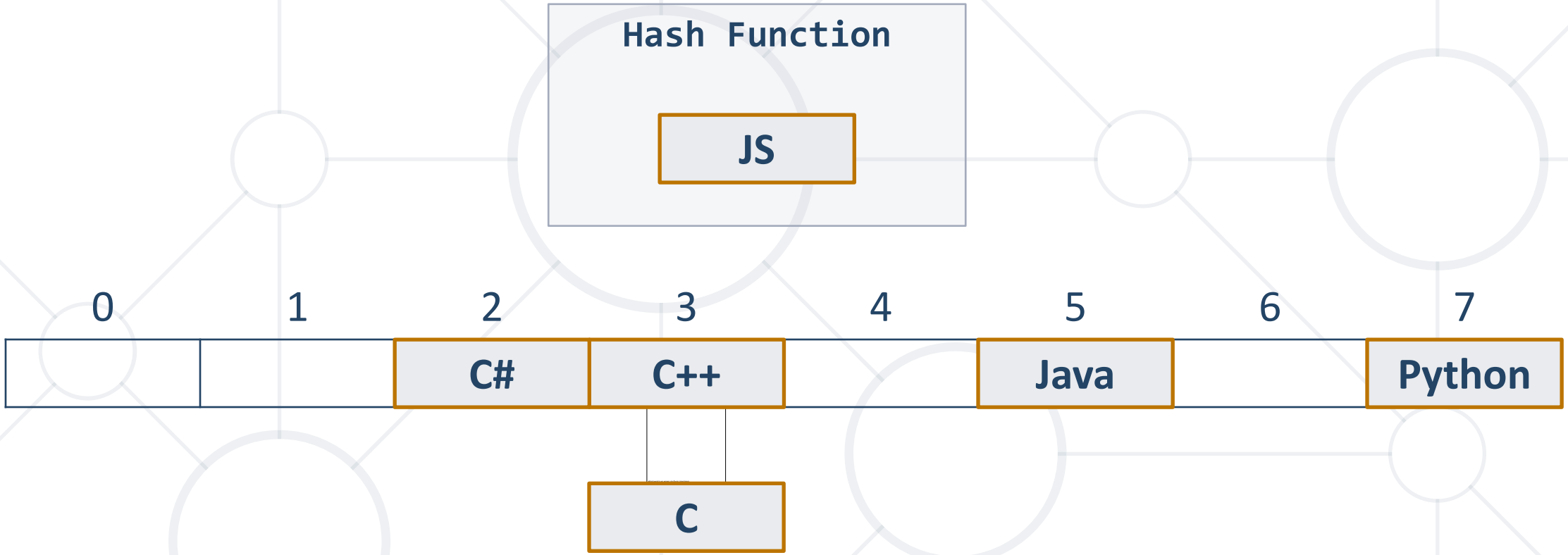
Collision Resolution: Chaining



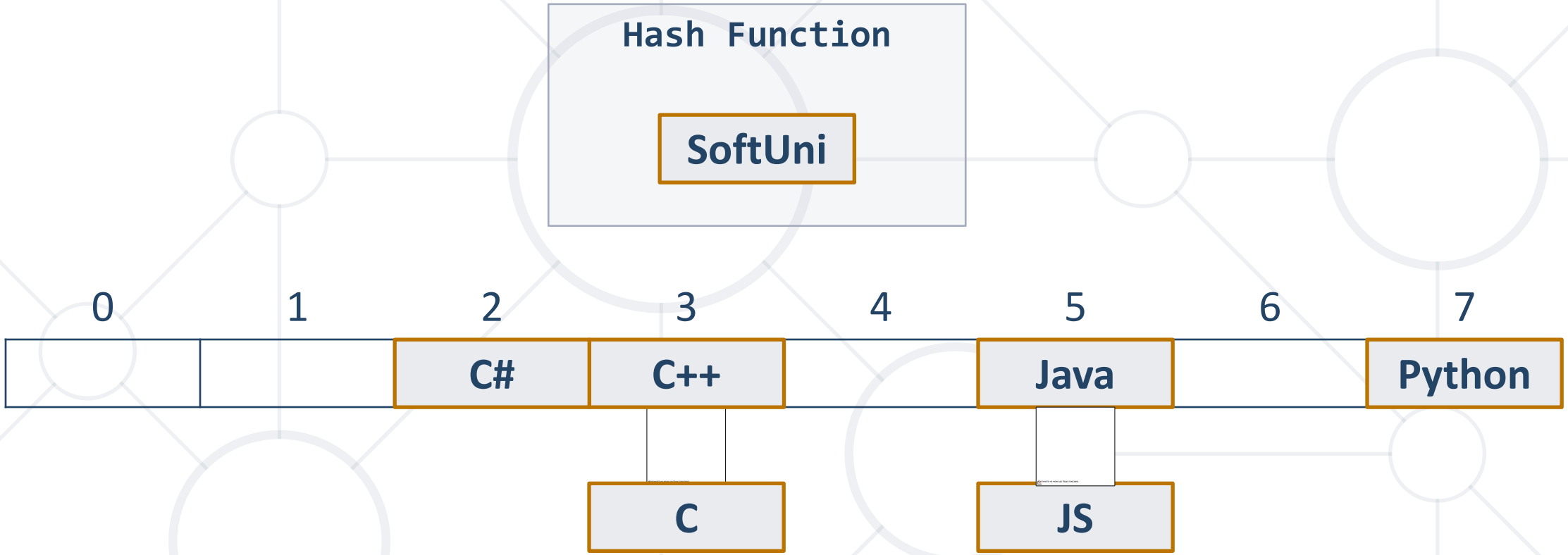
Collision Resolution: Chaining



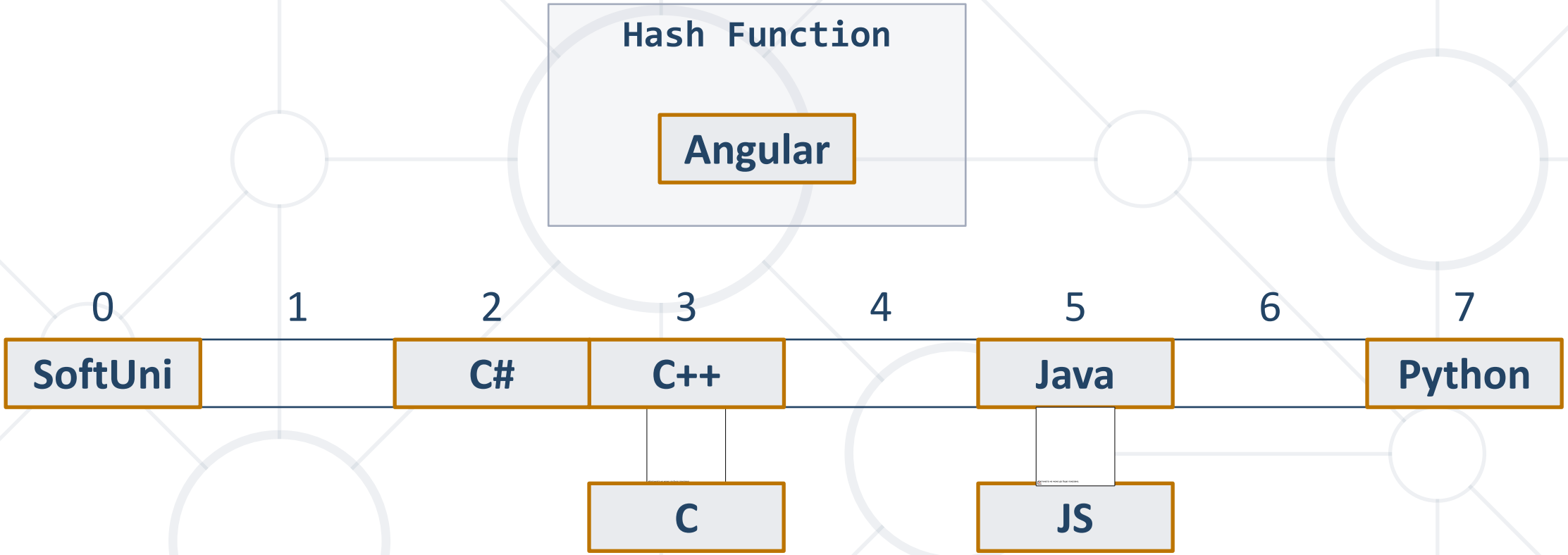
Collision Resolution: Chaining



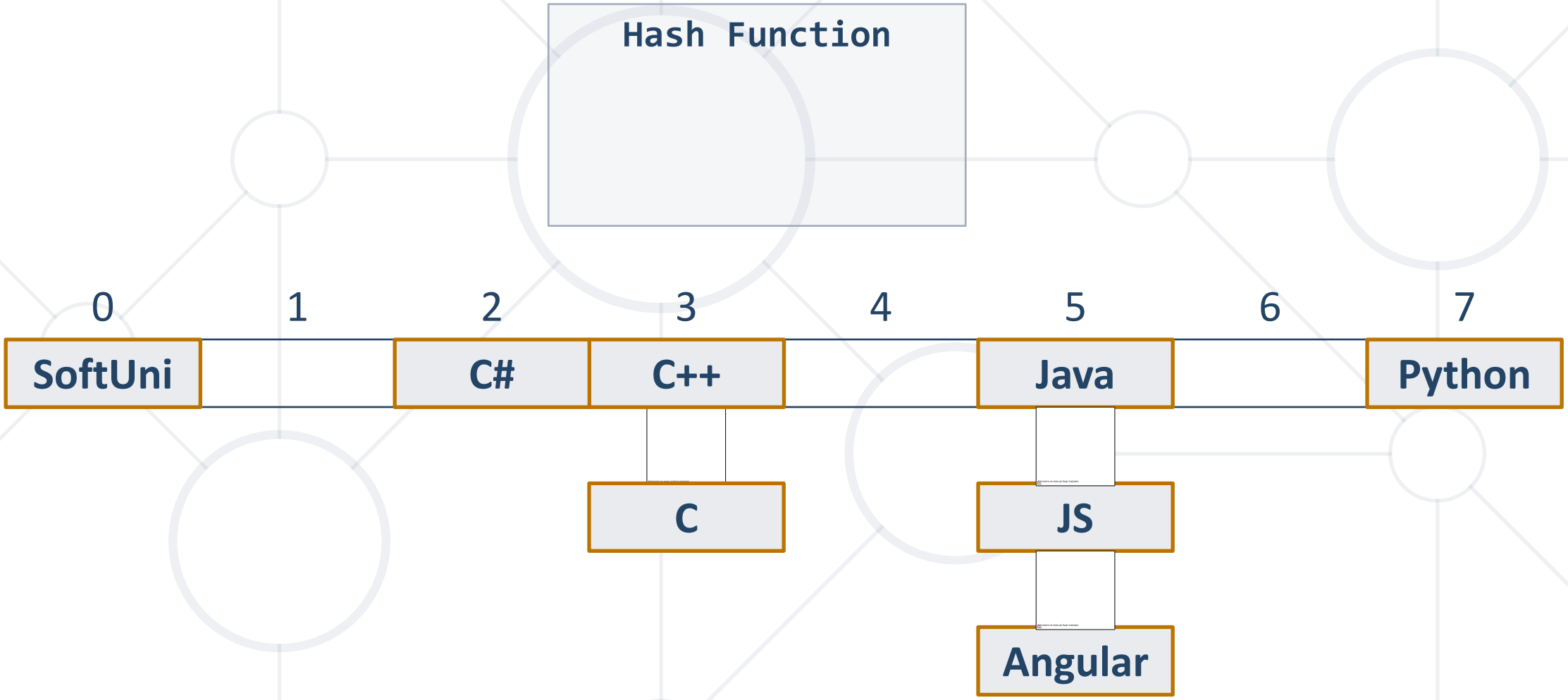
Collision Resolution: Chaining



Collision Resolution: Chaining



Collision Resolution: Chaining



Collision Resolution: Open Addressing

- **Open addressing** as collision resolution strategy means to take another slot in the hash-table in case of collision, e.g.
  - **Linear probing**: take the next empty slot just after the collision
    - **h(key, i) = h(key) + i**
    - where **i** is the attempt number: 0, 1, 2, …
    - **h(key) + 1, h(key) + 2, h(key) + 3**, etc.

Collision Resolution: Open Addressing (2)

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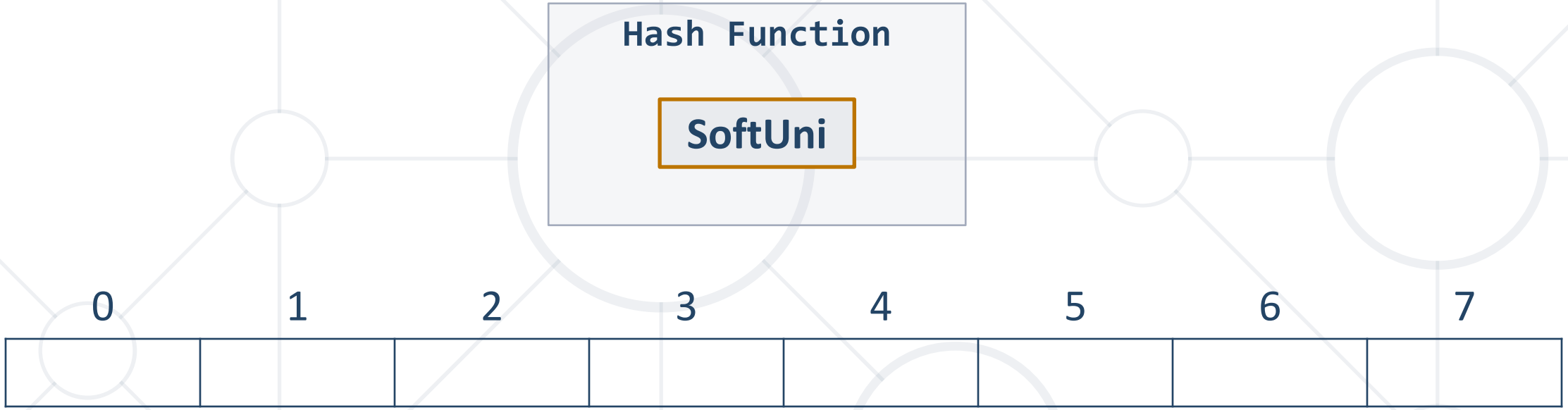
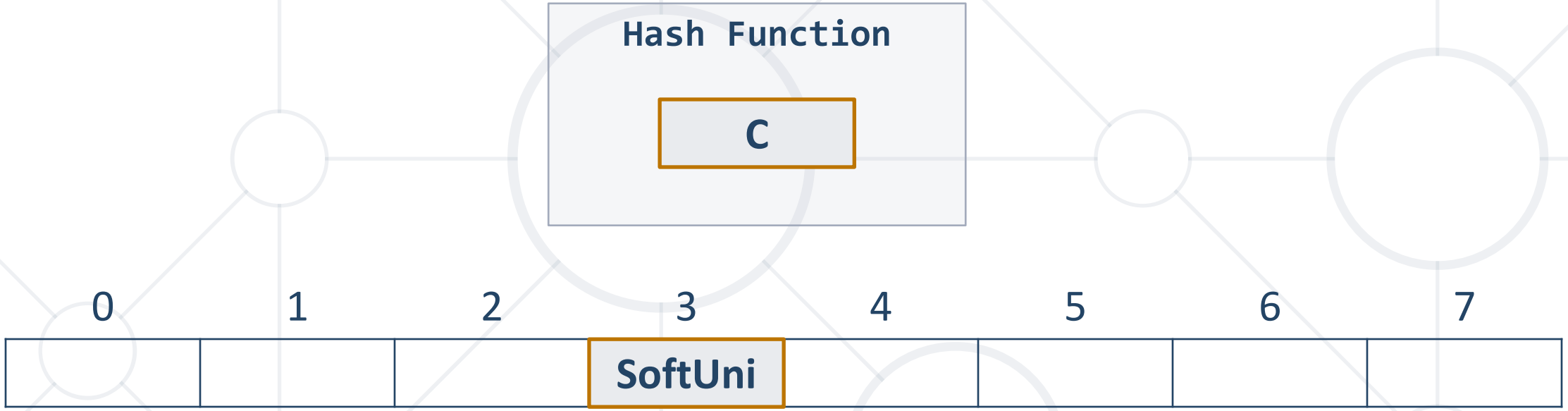
- **Quadratic probing**: the i^{th} next slot is calculated by a quadratic polynomial (c_1 and c_2 are some constants)
  - $h(key, i) = h(key) + c_1*i + c_2*i^2$
  - $h(key) + 1^2$, $h(key) + 2^2$, $h(key) + 3^2$, etc.
- **Re-hashing**: use separate (second) hash-function for collisions
  - $h(key, i) = h_1(key) + i*h_2(key)$

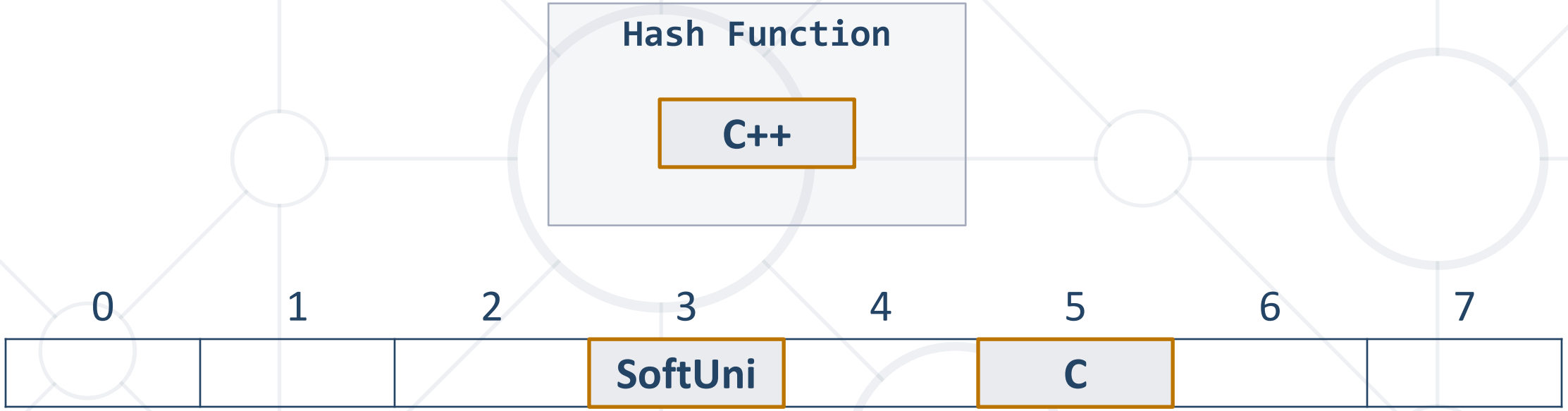
Collision Resolution: Linear Probing



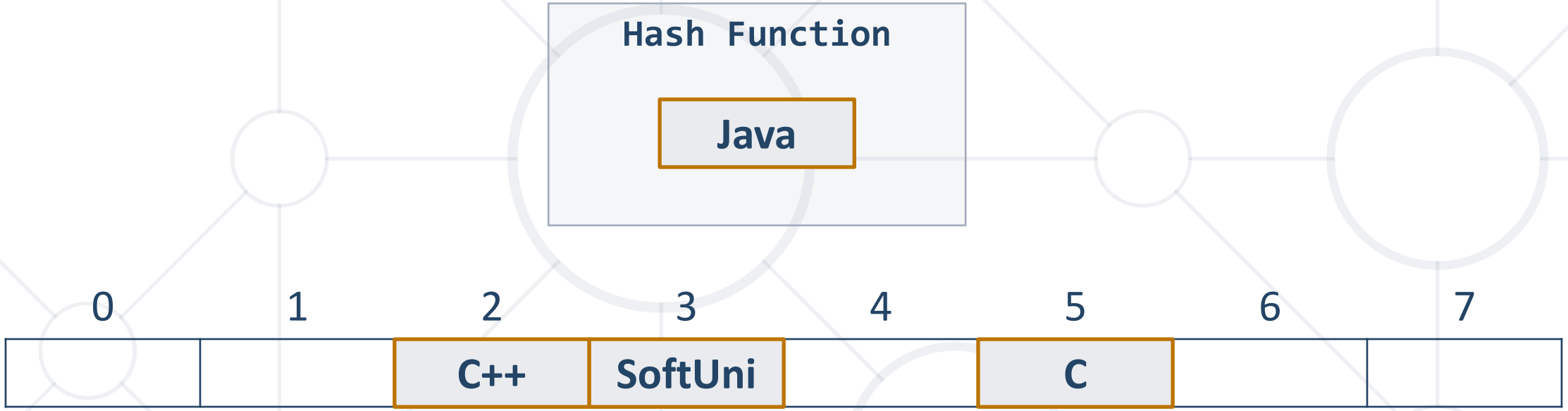
Collision Resolution: Linear Probing



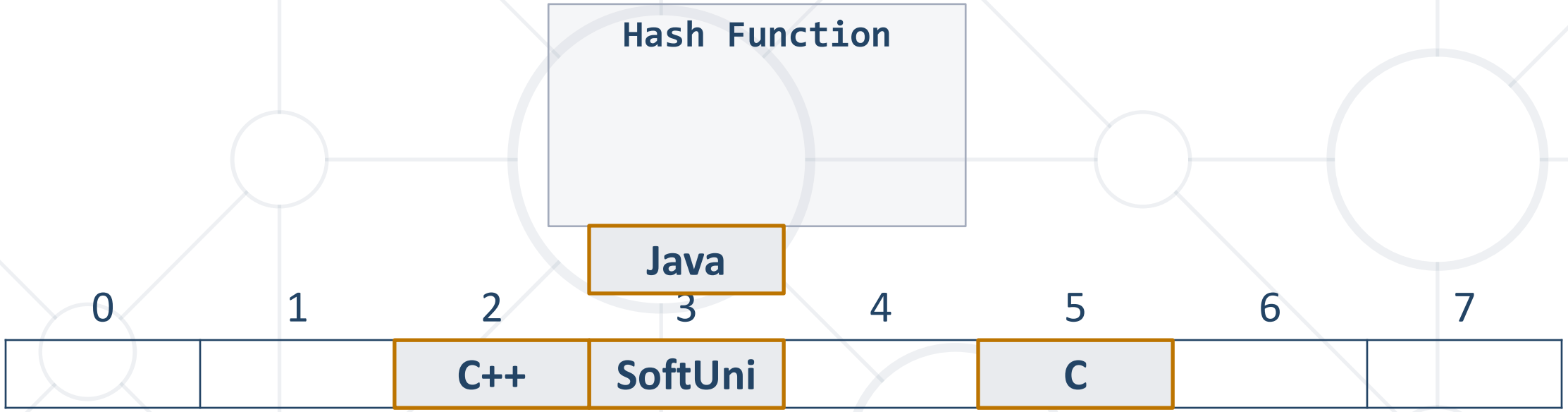
Collision Resolution: Linear Probing



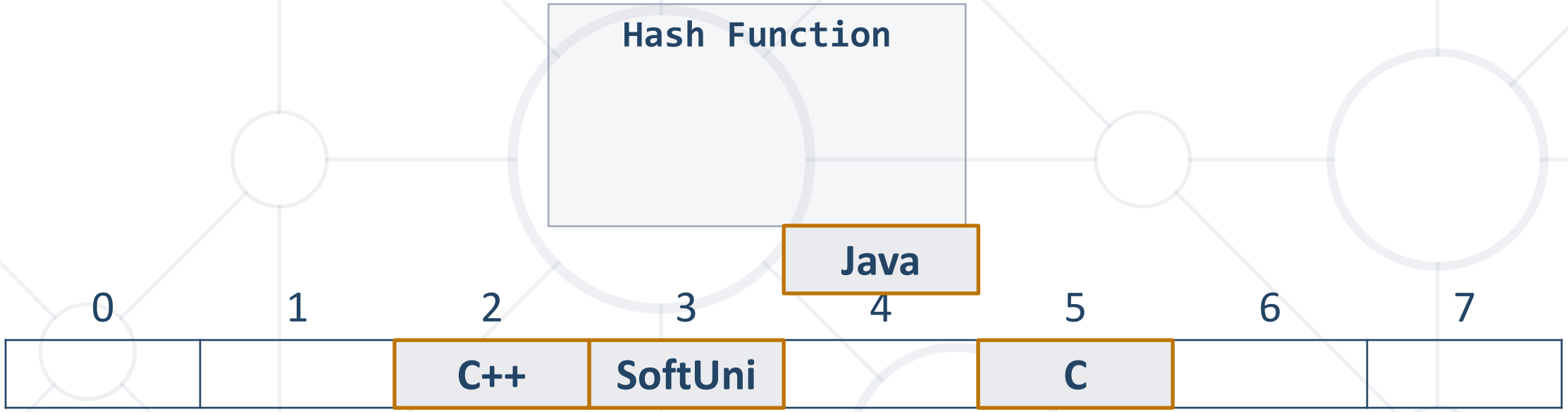
Collision Resolution: Linear Probing



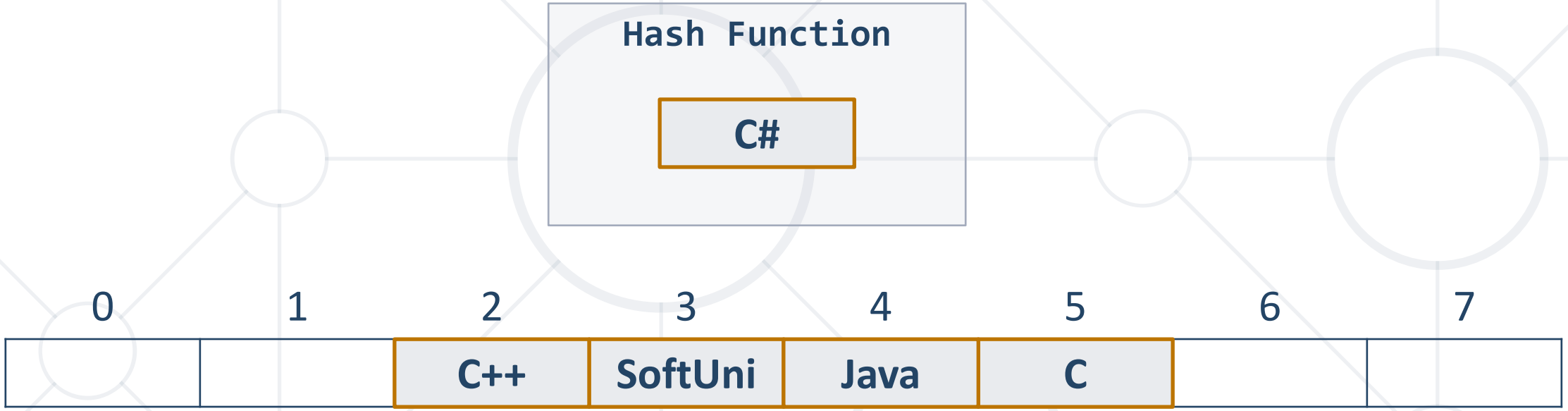
Collision Resolution: Linear Probing



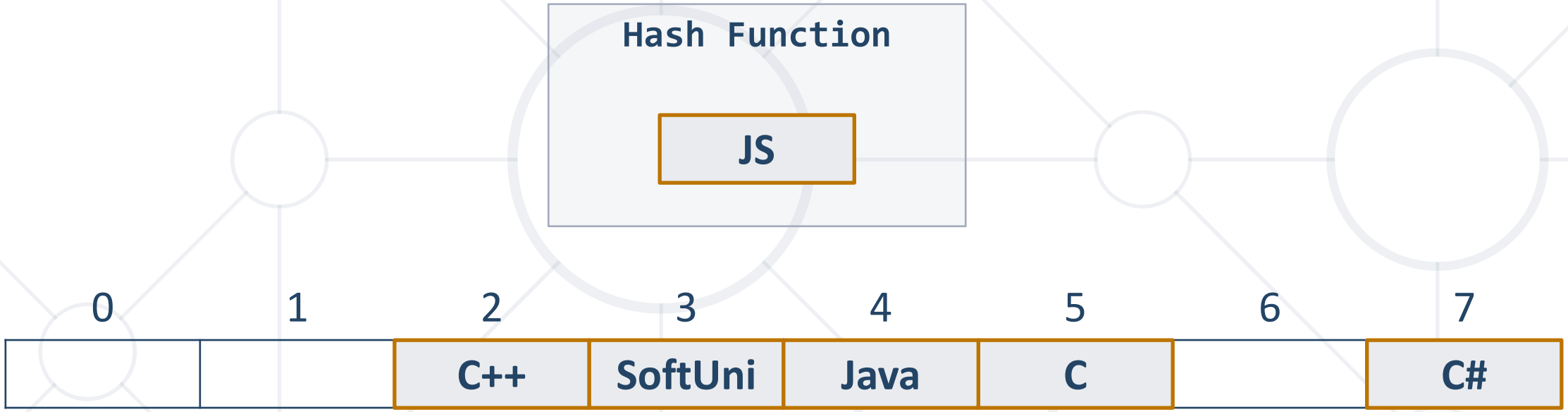
Collision Resolution: Linear Probing



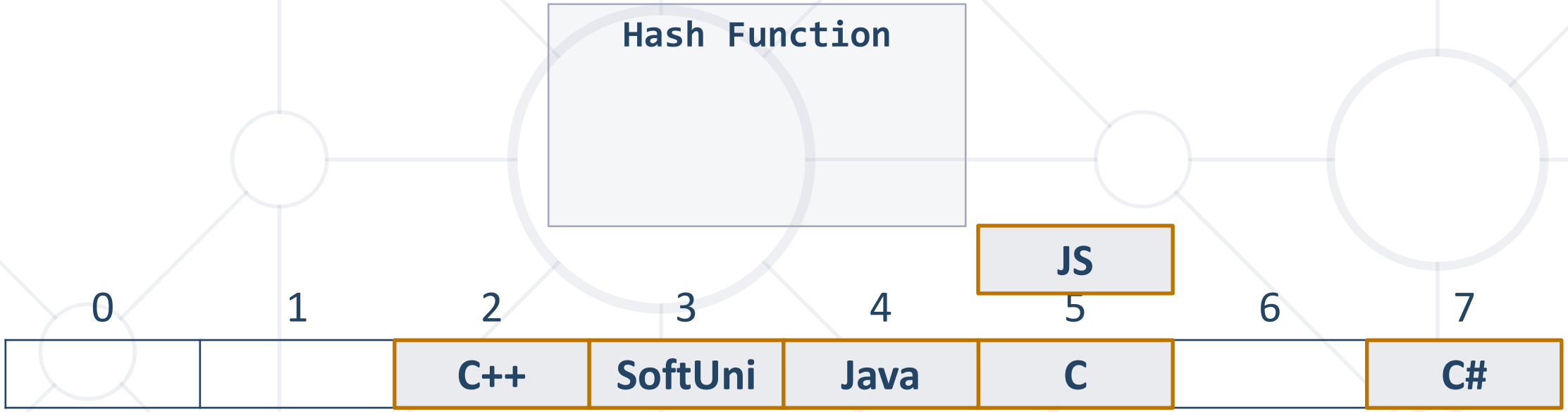
Collision Resolution: Linear Probing



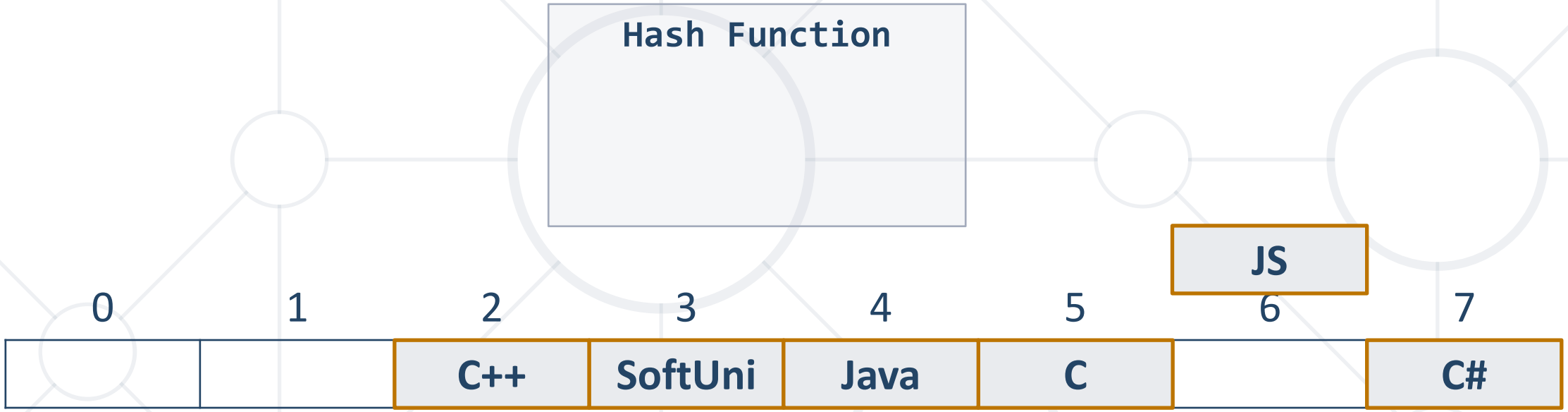
Collision Resolution: Linear Probing



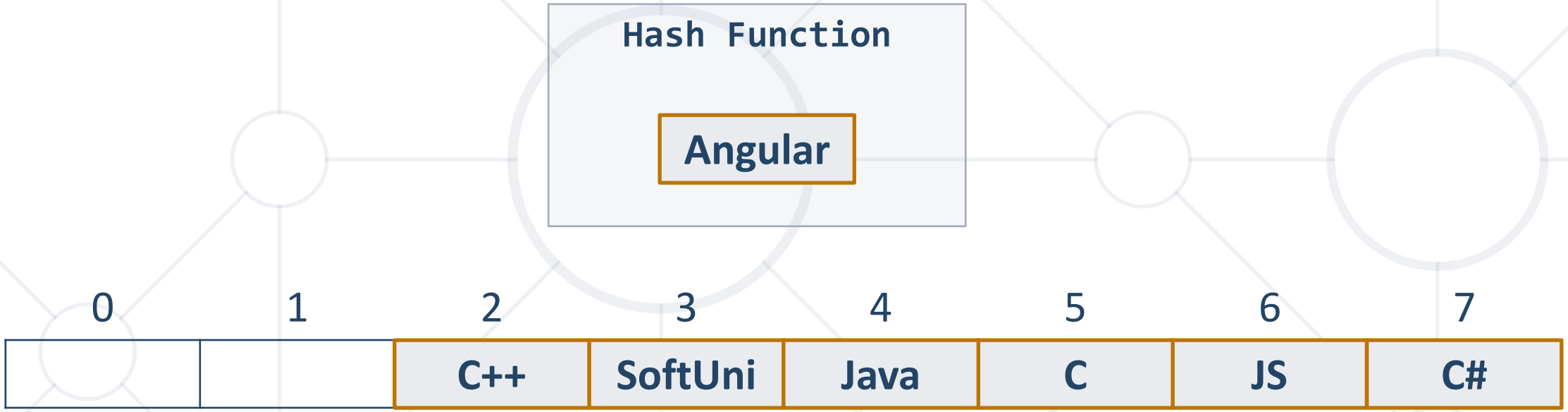
Collision Resolution: Linear Probing



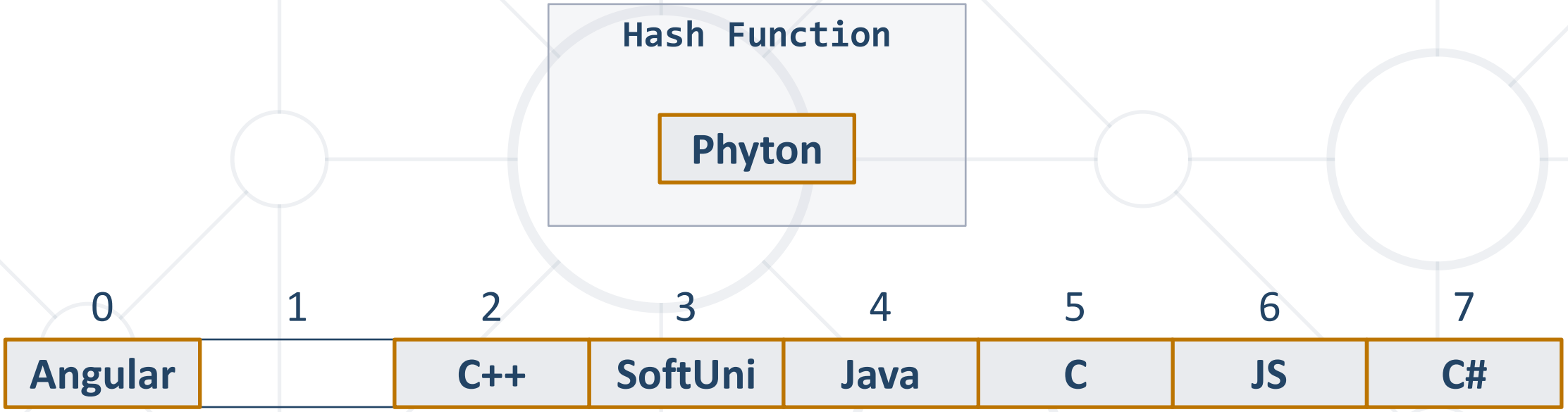
Collision Resolution: Linear Probing



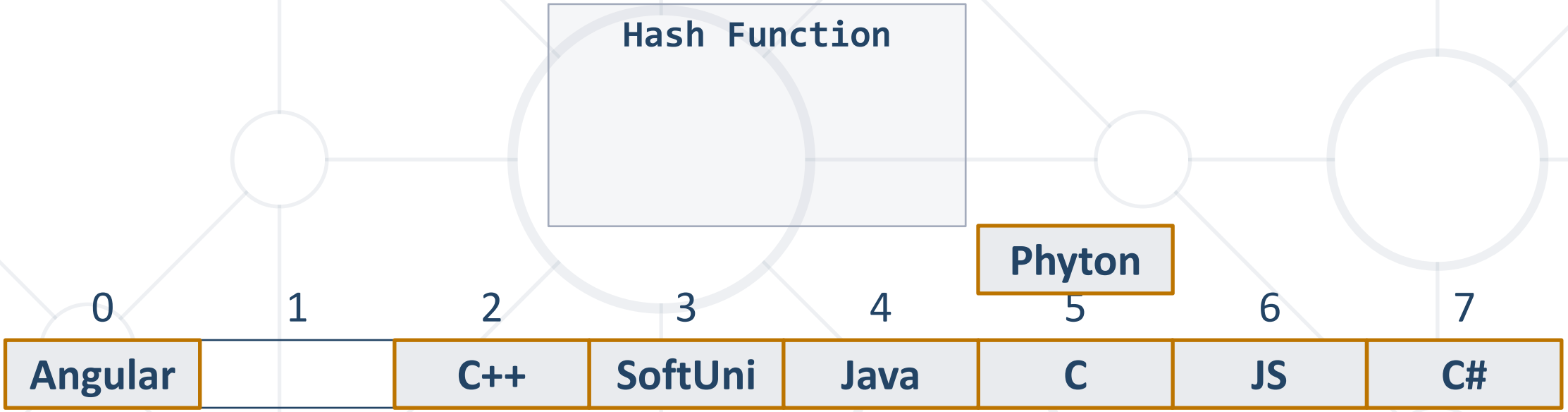
Collision Resolution: Linear Probing



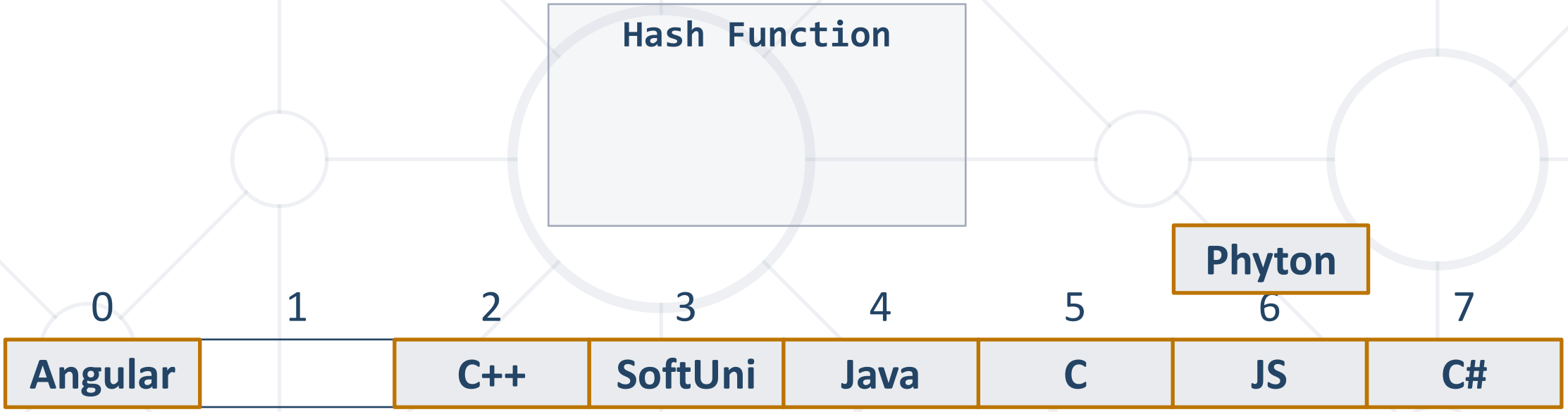
Collision Resolution: Linear Probing



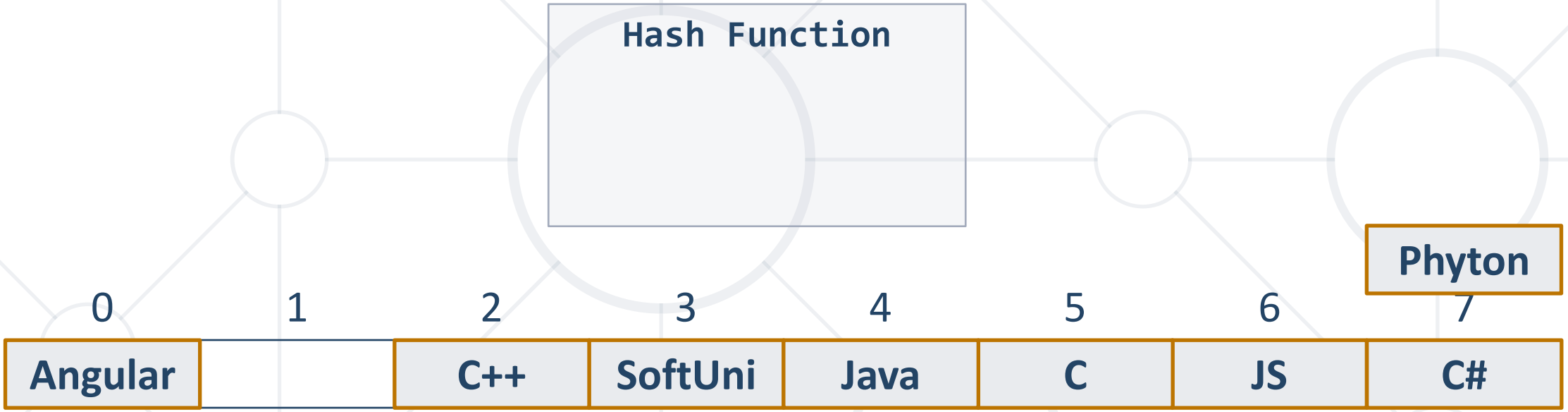
Collision Resolution: Linear Probing



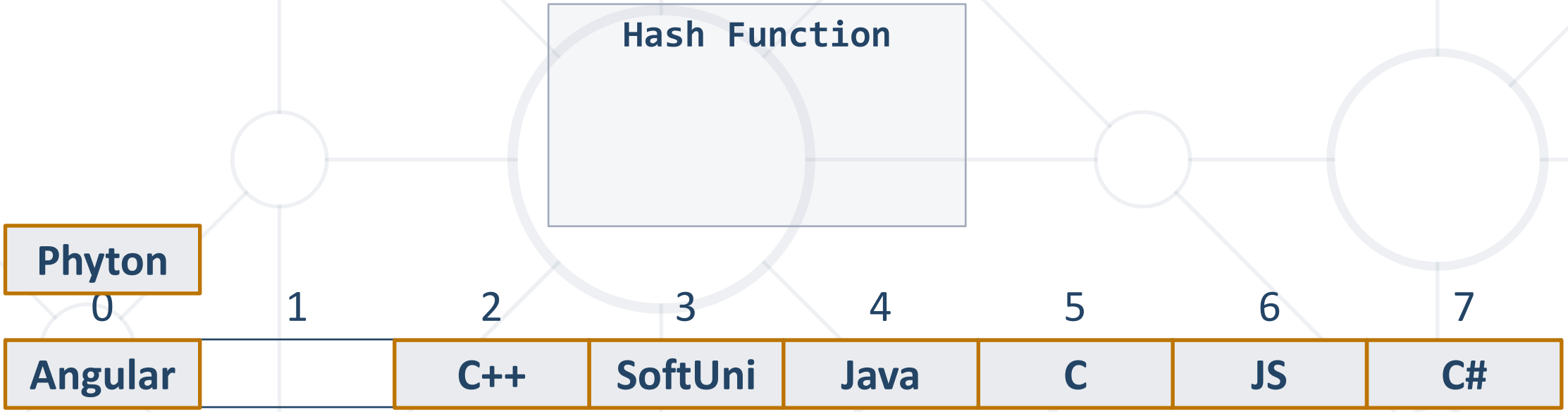
Collision Resolution: Linear Probing



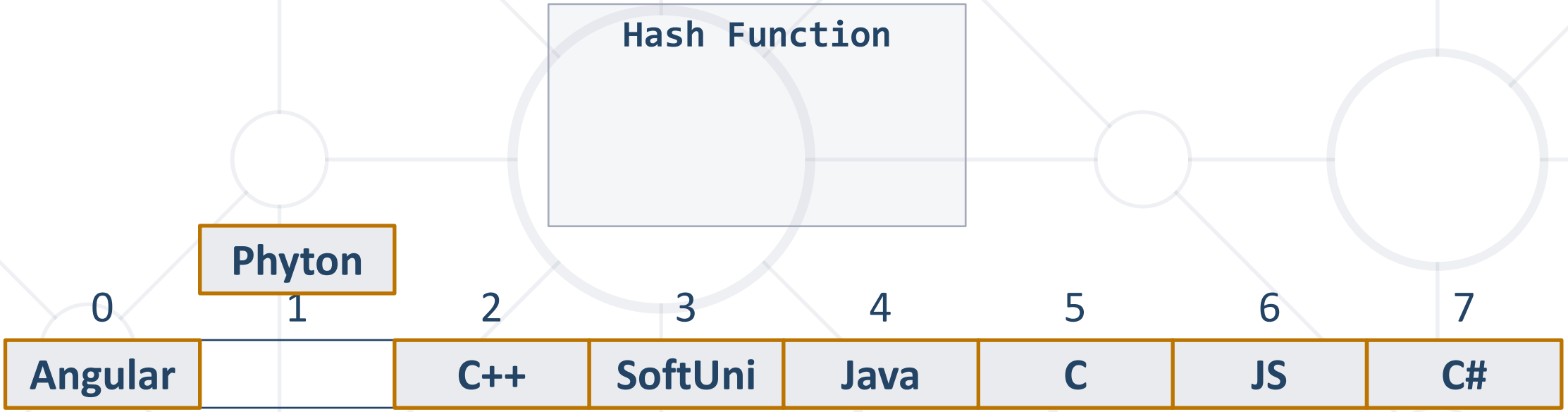
Collision Resolution: Linear Probing



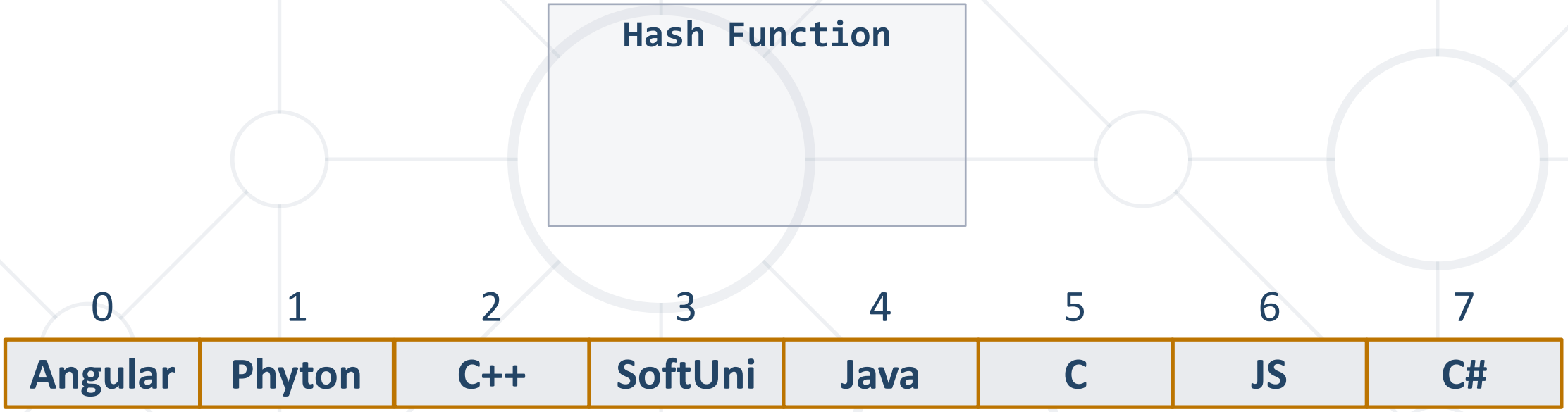
Collision Resolution: Linear Probing



Collision Resolution: Linear Probing



Collision Resolution: Linear Probing



Linear Probing - Quiz

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TIME:

- What is the average running time of delete in linear-probing hash table? Your hash function satisfies the uniform hashing assumption and that the hash table is at most 50% full.
  - O(1)
  - O(log N)
  - O(N)
  - O(N log N)

Linear Probing - Answer

- What is the average running time of delete in linear-probing hash table? Your hash function satisfies the uniform hashing assumption and that the hash table is at most 50% full.
  - O(1)
  - O(log N)
  - O(N)
  - O(N log N)

Hash Table Performance

- The hash-table performance depends on the probability of collisions
  - **Less collisions** → **faster** add / find / delete operations
  - **Collisions resolution** algorithm
  - **Fill factor** (used buckets / all buckets)

Hash Tables Efficiency

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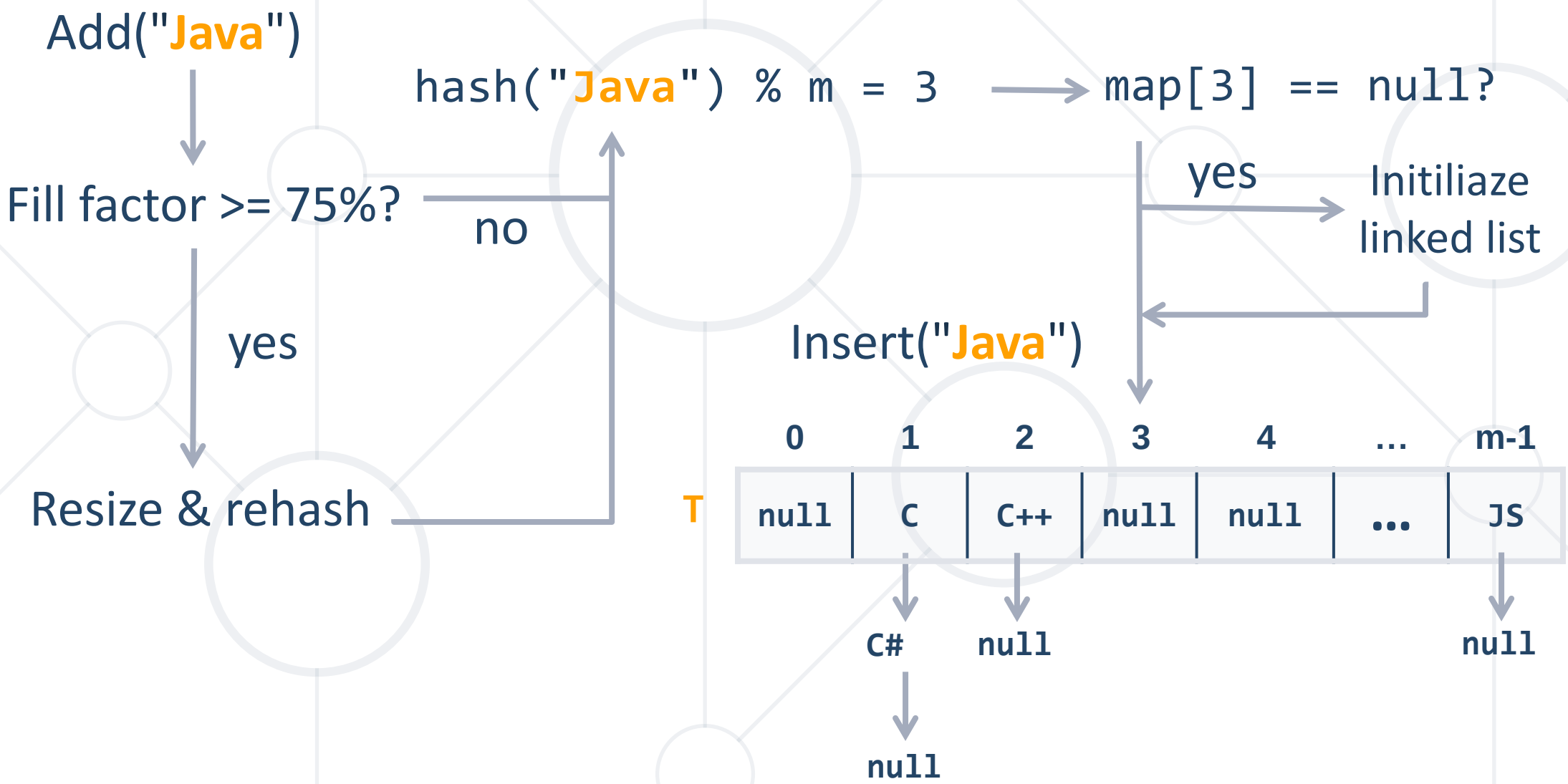
- **Add** / **Find** / **Delete** take just few primitive operations
  - Speed does not depend on the size of the hash-table
  - Amortized complexity **O(1)** – constant time
- Example:
  - Finding an element in a **hash-table** holding **1 000 000 elements** takes average just **1-2 steps**
  - Finding an element in an **array** holding **1 000 000 elements** takes average **500 000 steps**

How Big the Hash-Table Should Be?

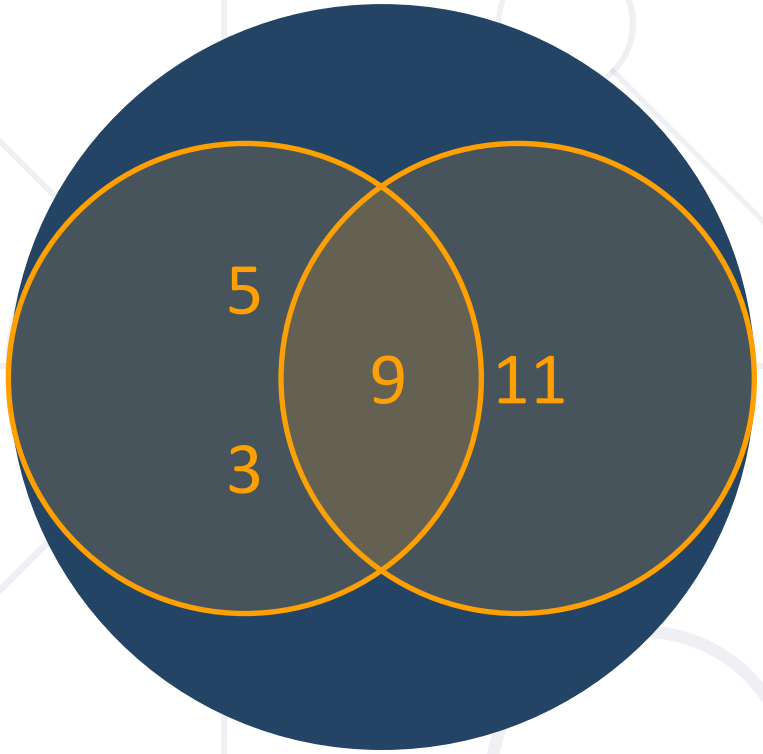
- The **load factor** (fill factor) **= used cells / all cells**
  - How much the hash table is filled, e.g. 65%
- Smaller fill factor leads to less collisions (faster average seek time)
- Recommended fill factors:
  - When **chaining** is used as collision resolution → less than **75%**
  - When **open addressing** is used → less than **50%**

Adding Item to Hash Table With Chaining



Lab Exercise

Implement a Hash-Table with Chaining



Sets and Bags

Set Operations

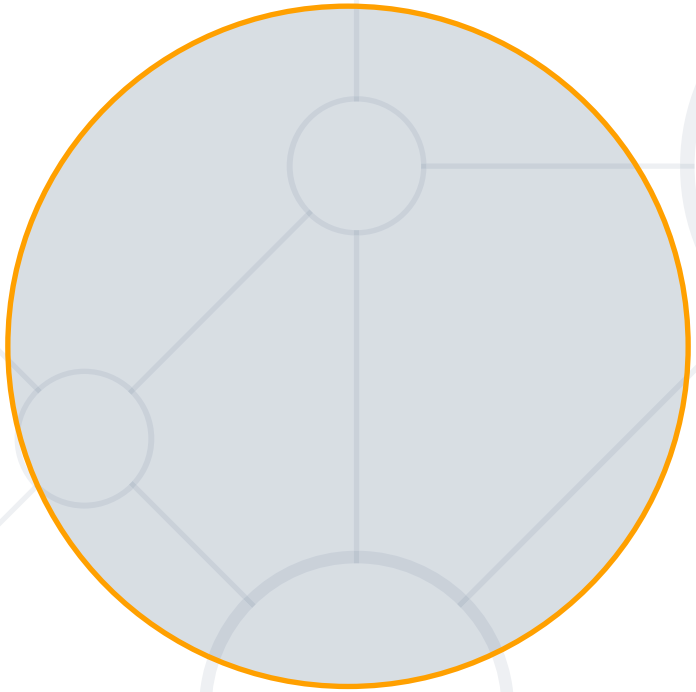
Set and Bag ADTs

- The abstract data type (ADT) "**set**" keeps a set of elements with no duplicates
- Sets with duplicates are also known as ADT "**bag**"
- Set specific operations:
  - **unionWith(set)**
  - **intersectWith(set)**
  - **exceptWith(set)**
  - **symmetricExceptWith(set)**

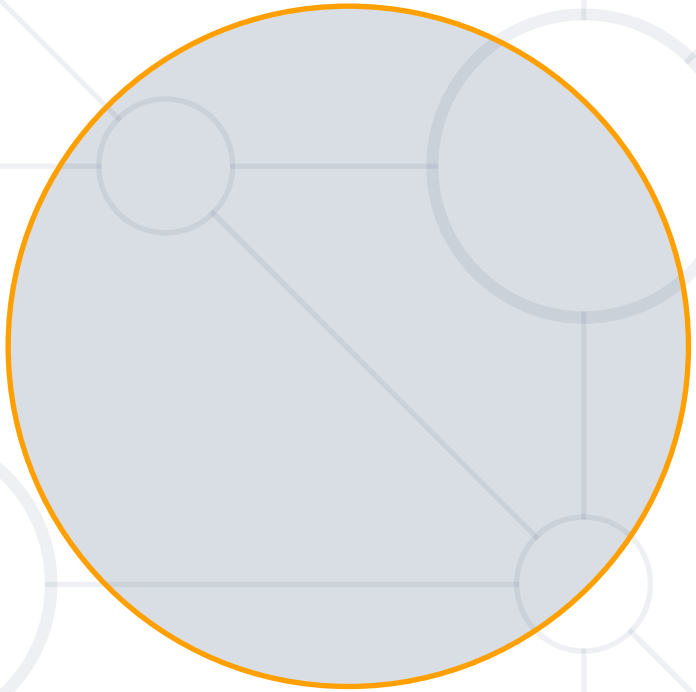
Known as relative complement in math

Known as symmetric difference

Union

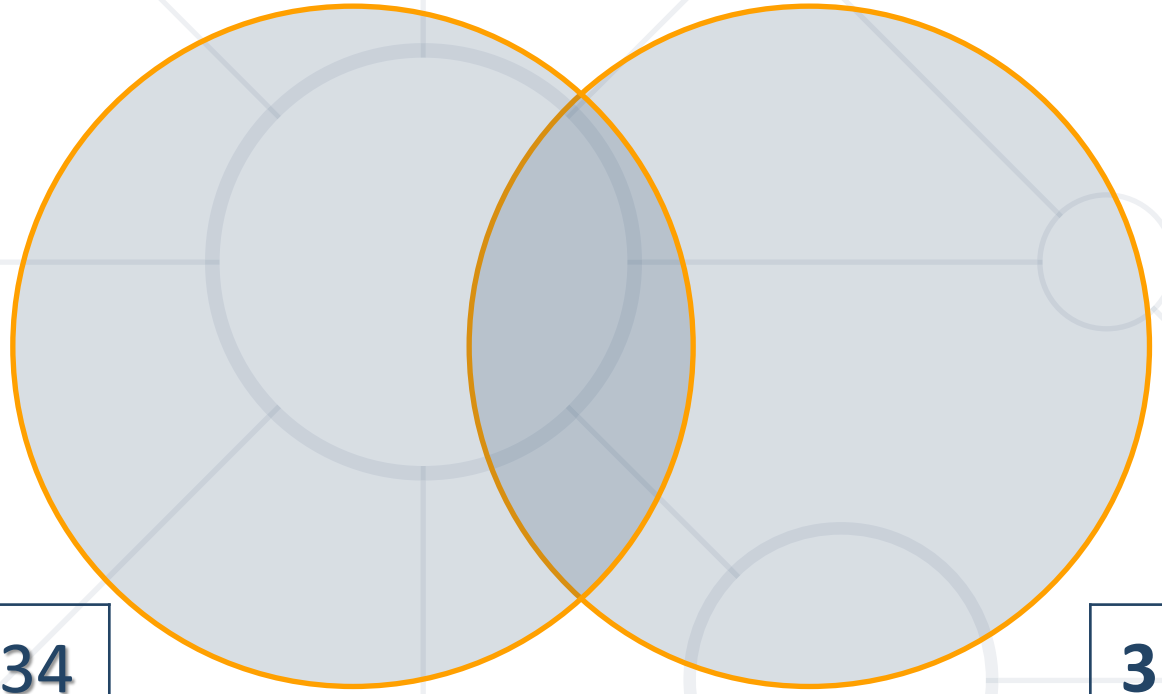


1	5	2	11	8	34



3	2	9	34	8	0

Union

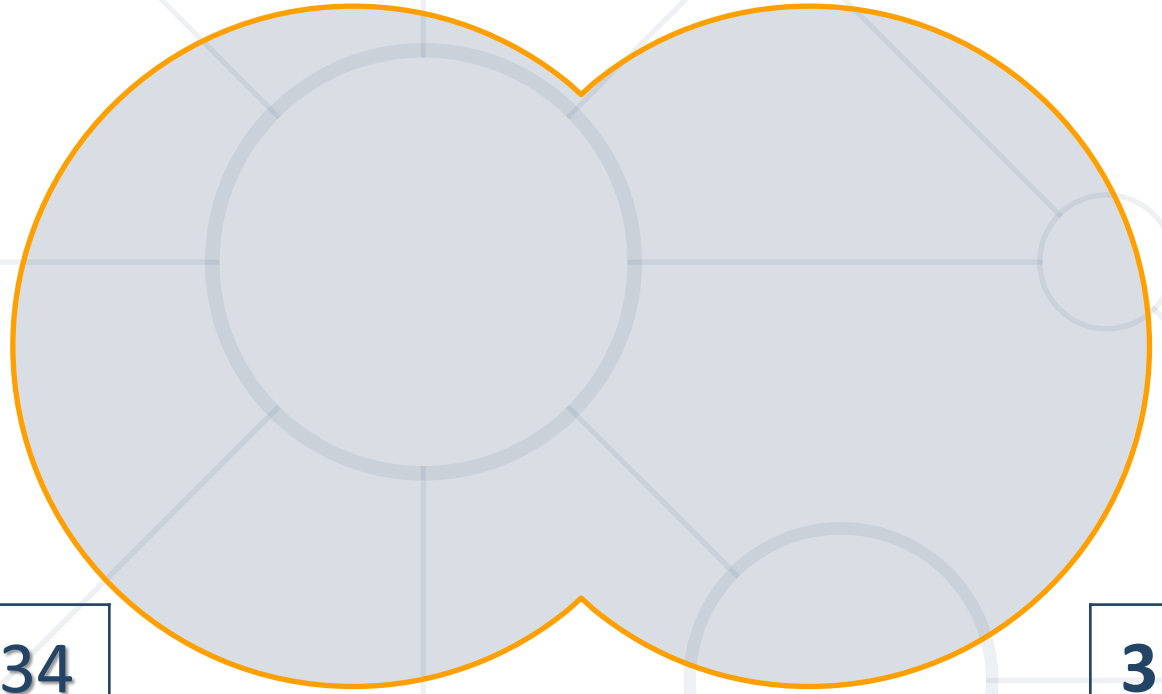


1	5	2	11	8	34

3	2	9	34	8	0

1	5	2	11	8	34	3	9	0

Union

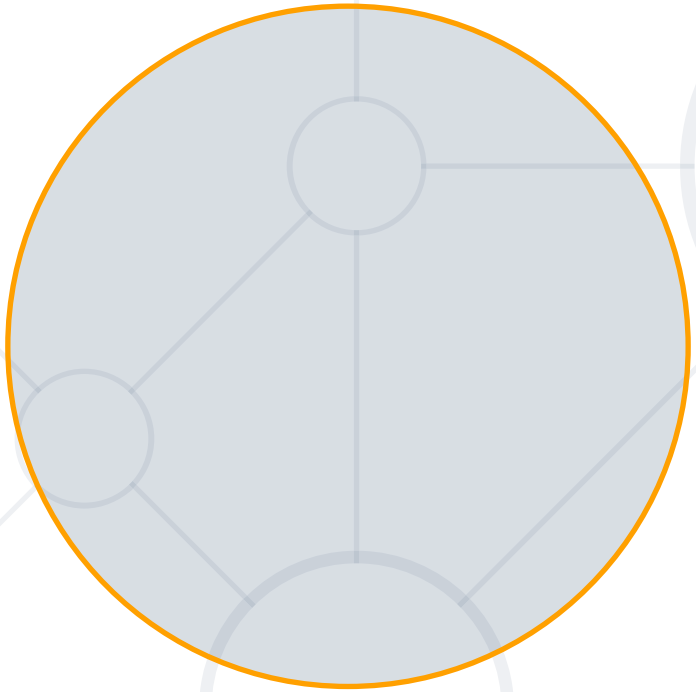


1	5	2	11	8	34

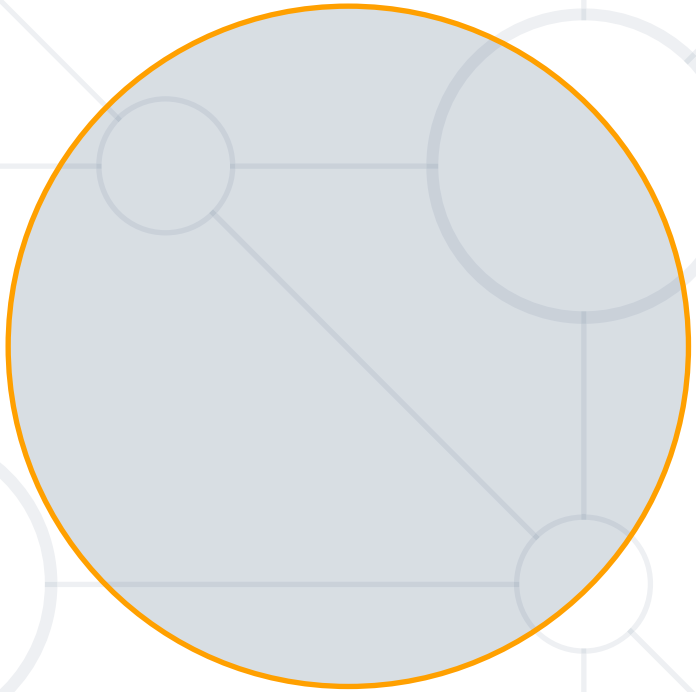
3	2	9	34	8	0

1	5	2	11	8	34	3	9	0

Intersects

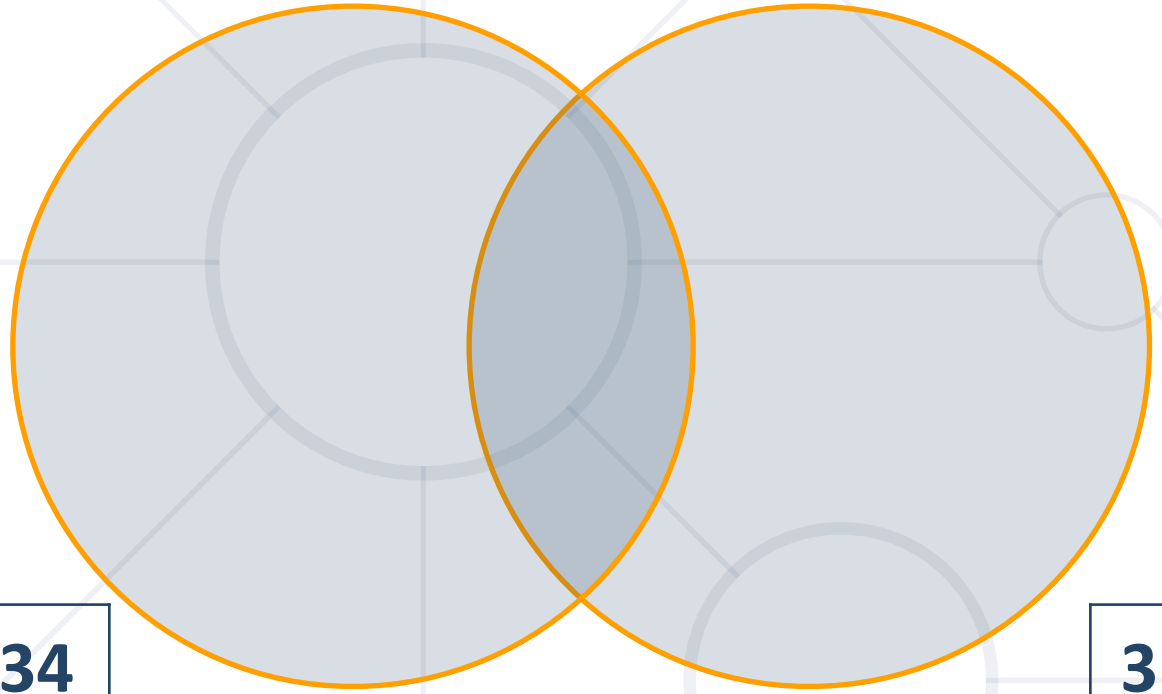


1	5	2	11	8	34



3	2	9	34	8	0

Intersects

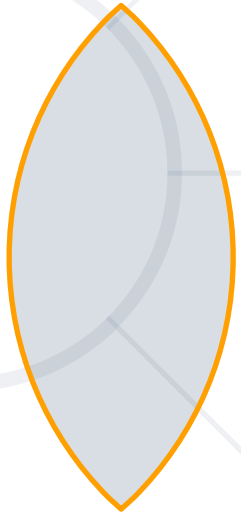


1	5	2	11	8	34

3	2	9	34	8	0

2	8	34

Intersects

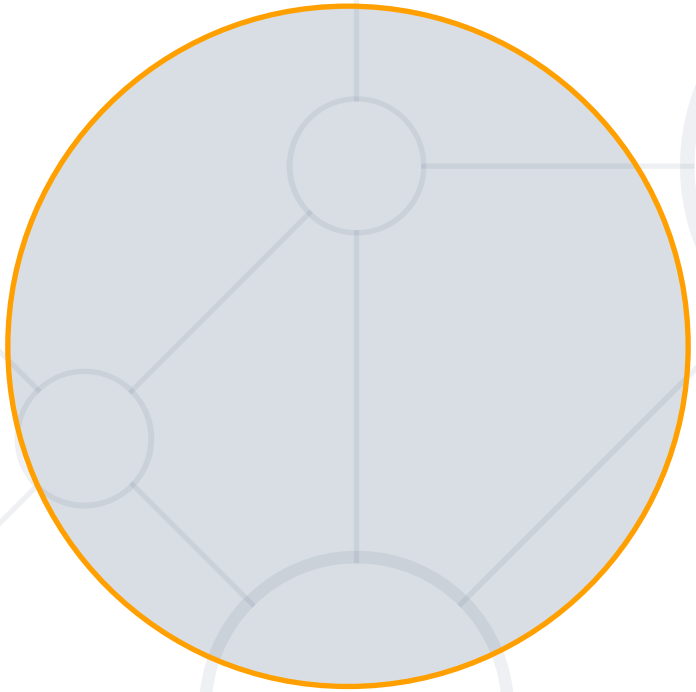


1	5	2	11	8	34

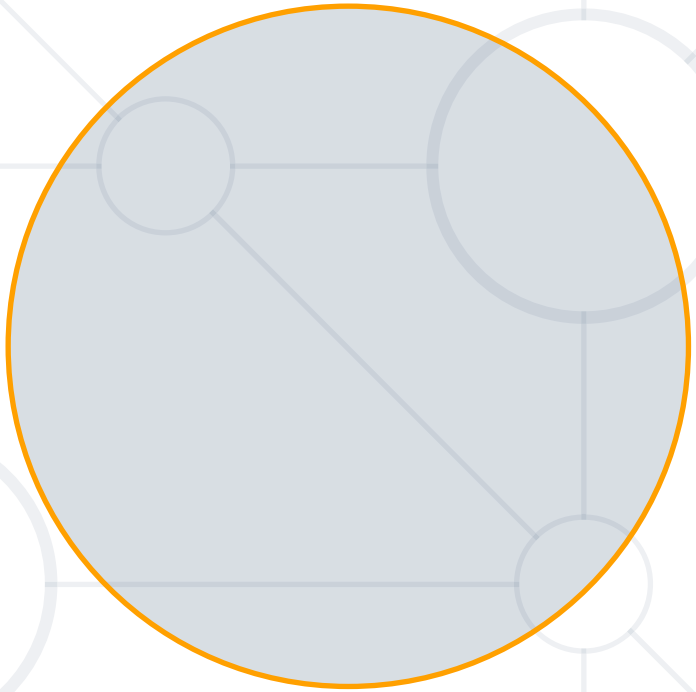
2	8	34

3	2	9	34	8	0

Except

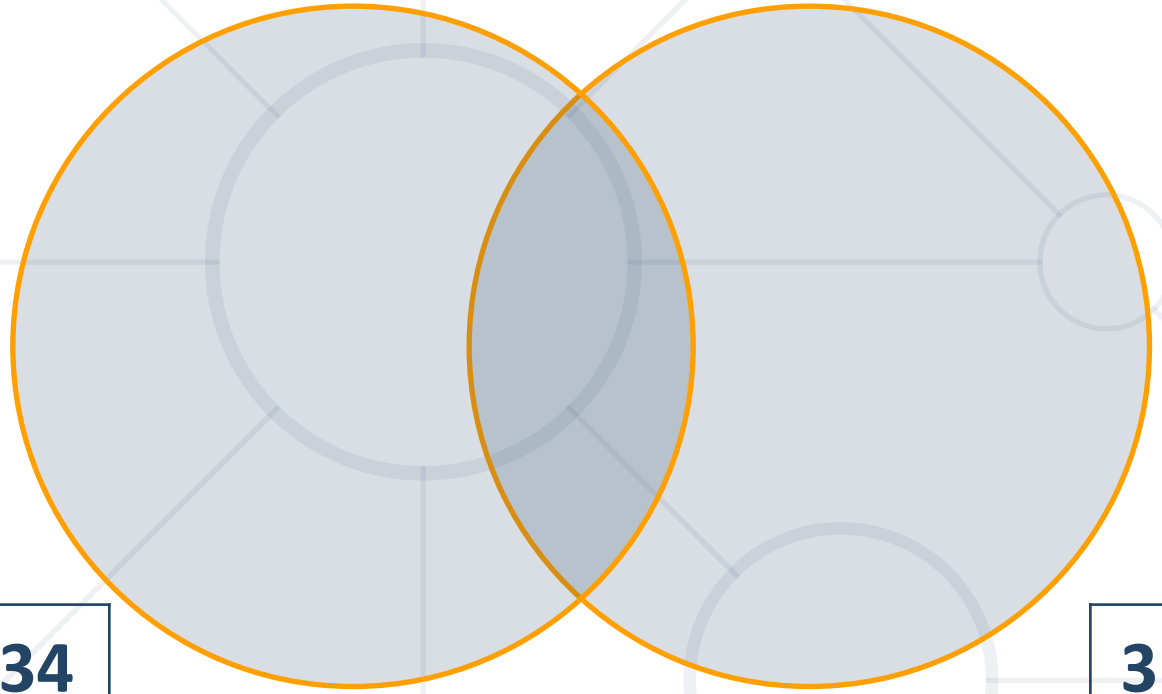


1	5	2	11	8	34



3	2	9	34	8	0

Except

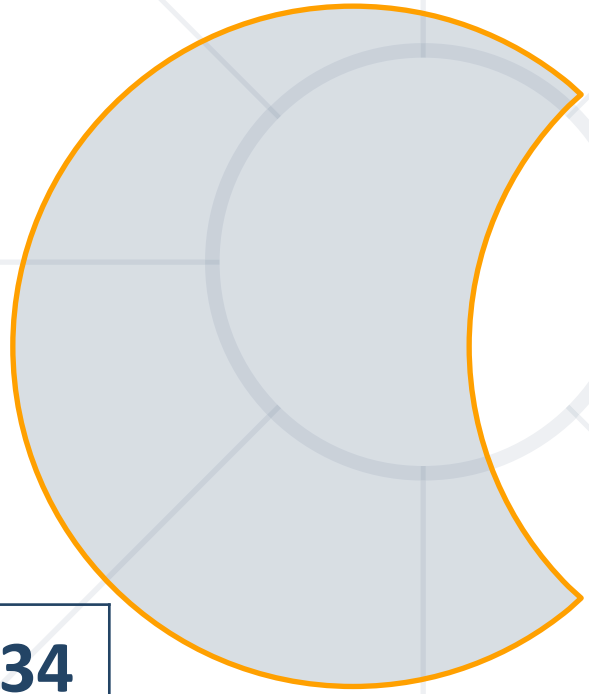


1	5	2	11	8	34

3	2	9	34	8	0

1	5	11

Except

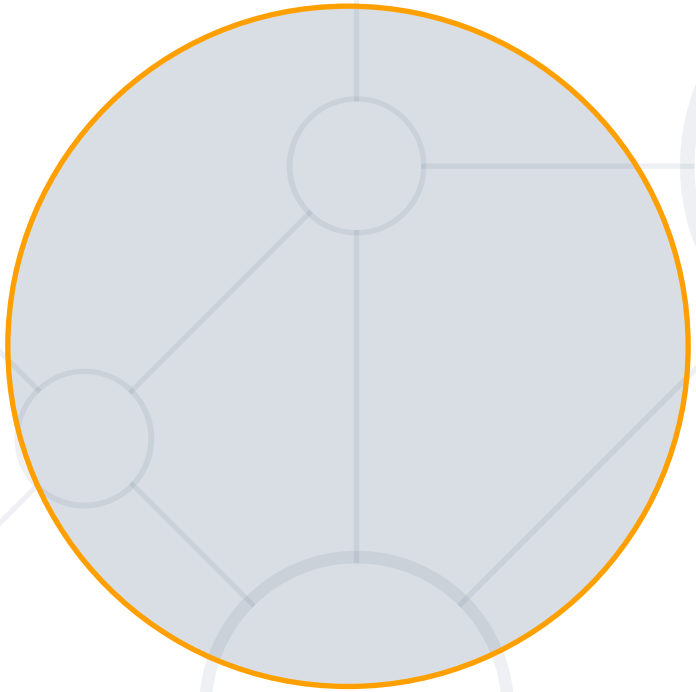


1	5	2	11	8	34

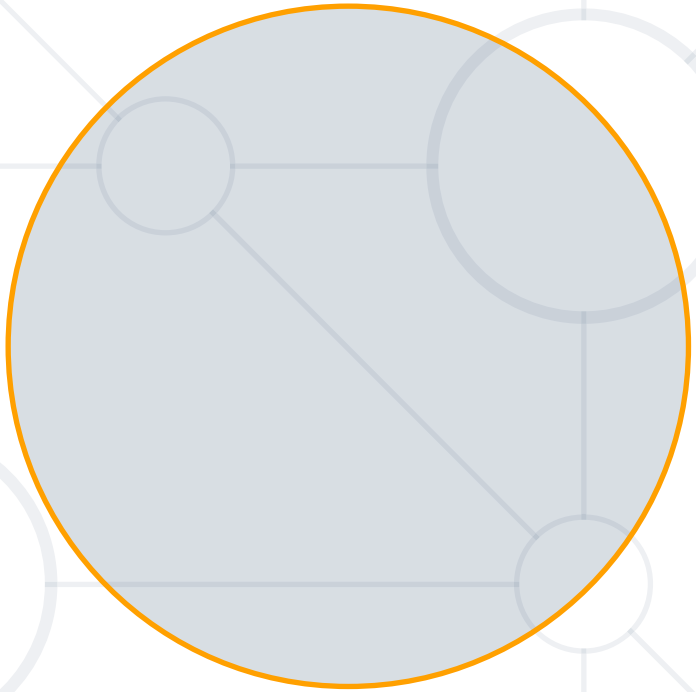
1	5	11

3	2	9	34	8	0

Symmetric Except

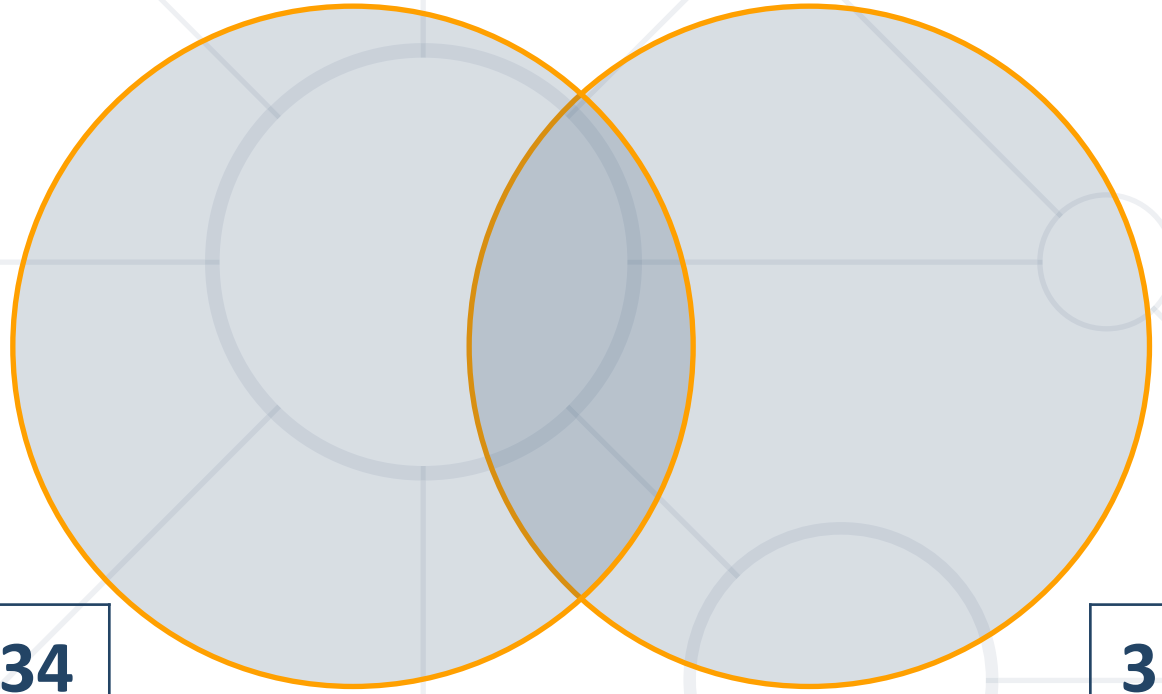


1	5	2	11	8	34



3	2	9	34	8	0

Symmetric Except

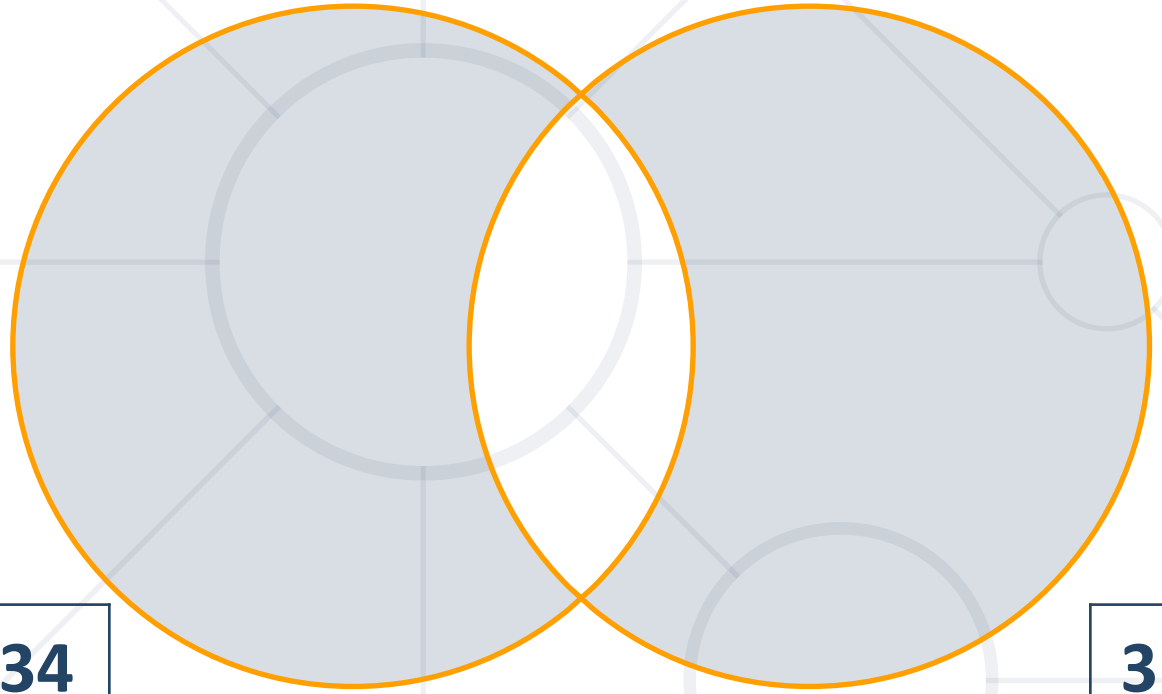


1	5	2	11	8	34

1	5	11	3	9	0

3	2	9	34	8	0

Symmetric Except



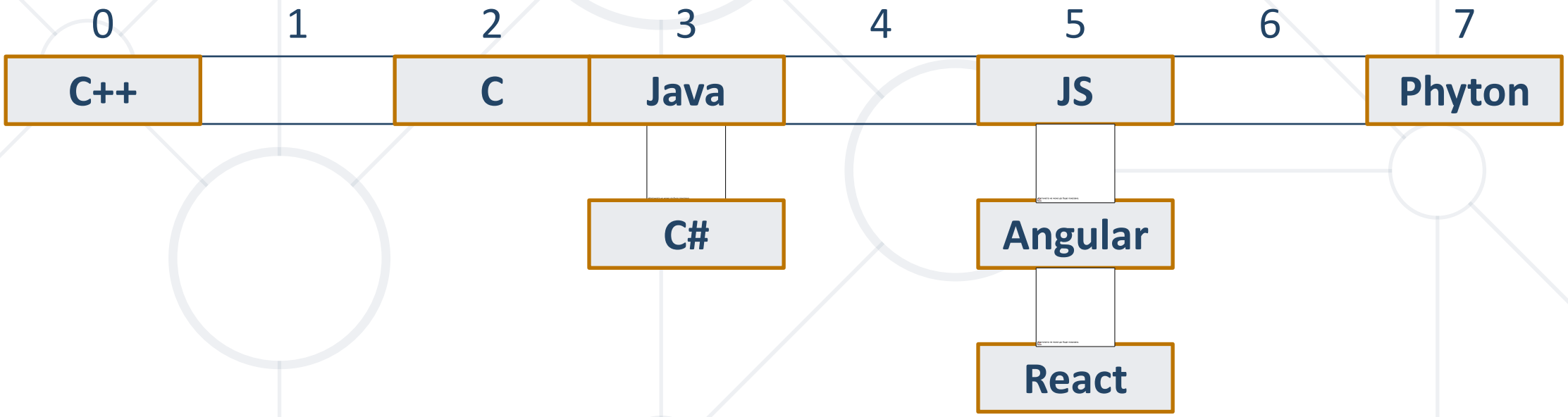
1	5	2	11	8	34

3	2	9	34	8	0

1	5	11	3	9	0

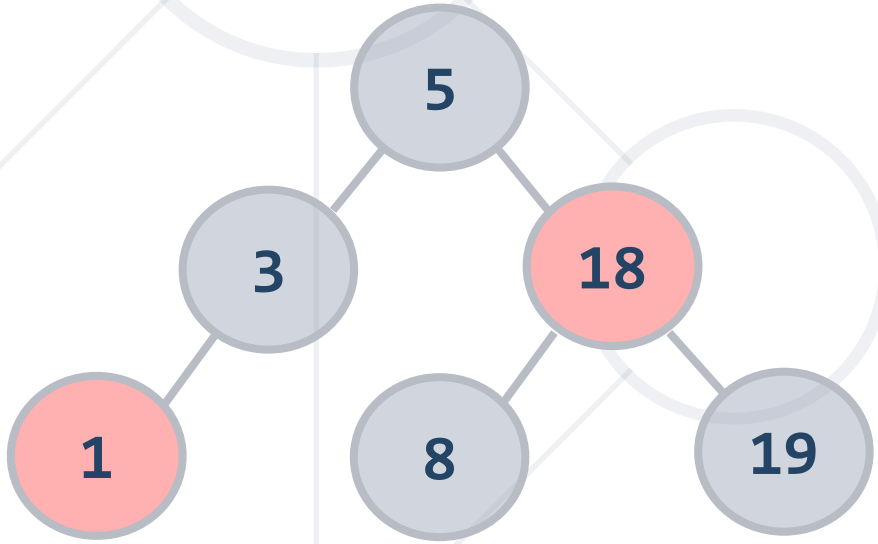
HashSet<T>

- **HashSet<T>** implements ADT **set** by hash table
  - Elements are in no particular order
- All major operations are fast: **Add** / **Delete** / **Contains**



TreeSet<T>

- **TreeSet<T>** implements ADT **set** by balanced search tree (red-black tree)
  - Elements are sorted in increasing order



Sets - Quiz

TIME:

- For given sets - {1, 2, 3, 4, 5} and {3, 4, 5, 6, 7}, what is the operation that will give us the following result: {1, 2, 6, 7}
  - Union
  - Intersects
  - Except
  - SymmetricExcept

Sets - Answer

- For given sets - {1, 2, 3, 4, 5} and {3, 4, 5, 6, 7}, what is the operation that will give us the following result: {1, 2, 6, 7}
  - Union
  - Intersects
  - Except
  - SymmetricExcept

Comparing Keys

Using Custom Key Classes

Comparasion Methods

- **Map<Key, Value>** relies on
  - **Object.equals()** – for comparing the keys
  - **Object.hashCode()** – for calculating the hash codes of the keys
- **TreeMap<Key, Value>** relies on **Comparable<Key>** for ordering the keys

Implementing equals() and hashCode()

Software University

```
public class Point {
  public int x;
  public int y;
  public boolean equals(Object obj) {
    if (this == obj) return true;
    if (o == null || getClass() != o.getClass()) return false;
    Point p = (Point)obj;
    return (x == p.x) && (y == p.y);
  }
  public int hashCode() {
      return (x << 16 | y >> 16) ^ y;
  }
}
```

Implementing Comparable<T>

Software University

```
public class Point implements Comparable<Point> {
  public int x;
  public int y;

  public int compareTo(Point other) {
    if (x != other.x) {
      return this.X.CompareTo(other.x);
    }
    else {
      return this.y.CompareTo(other.y);
    }
  }
}
```

Maps (Dictionaries)

Definition and Operations

The Dictionary (Map) ADT

- The abstract data type (ADT) "**dictionary**" maps key to values
  - Also known as "**map**" or "**associative array**"
  - Holds a set of **{key, value} pairs**
- Many implementations
  - Hash table, balanced tree, list, array, ...

key	value
John Smith	+1-555-8976
Sam Doe	+1-555-5030

ADT Map – Example

- Sample dictionary:

Key	Value
Java	Modern general-purpose object-oriented programming language
PHP	Popular server-side scripting language for Web development
compiler	Software that transforms a computer program to executable machine code
…	…

Map <Key, Value>

- Major operations:
  - **add(key, value)** – adds an element by key + value
  - **remove(key)** – removes a value by key
  - **get(key)** – returns the value by key
  - **keys** – returns a collection of all keys (in order of entry)
  - **values** – returns a collection of all values (in order of entry)

Map<Key, Value> (2)

- Major operations:
  - **containsKey(key)** – checks if given key exists in the dictionary
  - **containsValue(value)** – checks whether the dictionary contains given value
    - Warning: slow operation – **O(n)**

TreeMap<Key, Value>

- **TreeMap<Key, Value>** implements the ADT "dictionary" as self-balancing search tree
  - Elements are arranged in the tree ordered by key
  - Traversing the tree returns the elements in increasing order
  - **add** / **find** / **delete** perform **log N** operations
- Use **TreeMap<Key, Value>** when you need the elements sorted by key
  - Otherwise use **Map<Key, Value>** – it has better performance

Maps - Quiz

TIME'S

- Which built-in implementation of **Map<Key, Value>** sorts the items by value?
  - HashMap<Key, Value>
  - TreeMap<Key, Value>
  - None

Maps - Answer

- Which built-in implementation of **Map<Key, Value>** sorts the items by value?
  - HashMap<Key, Value>
  - TreeMap<Key, Value>
  - None

Hash Tables - Quiz

TIME'S

- Which is the main reason to use a hash table instead of a red-black BST?
  - Supports more operations efficiently
  - Better worst-case performance guarantee
  - Better performance in practice on typical inputs

Hash Tables - Answer

Software University

- Which is the main reason to use a hash table instead of a red-black BST?
  - Supports more operations efficiently
  - Better worst-case performance guarantee
  - Better performance in practice on typical inputs

Summary

- **Hash-tables** map keys to values
  - Rely on hash-functions to distribute the keys in the table
  - Collisions needs resolution algorithm (e.g. chaining)
  - Very fast add / find / delete – `O(1)`
- **Sets** hold a group of elements
- **Maps** map key to value

Questions?

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