

# Topological Data Analysis with $\epsilon$ -net Induced Lazy Witness Complex

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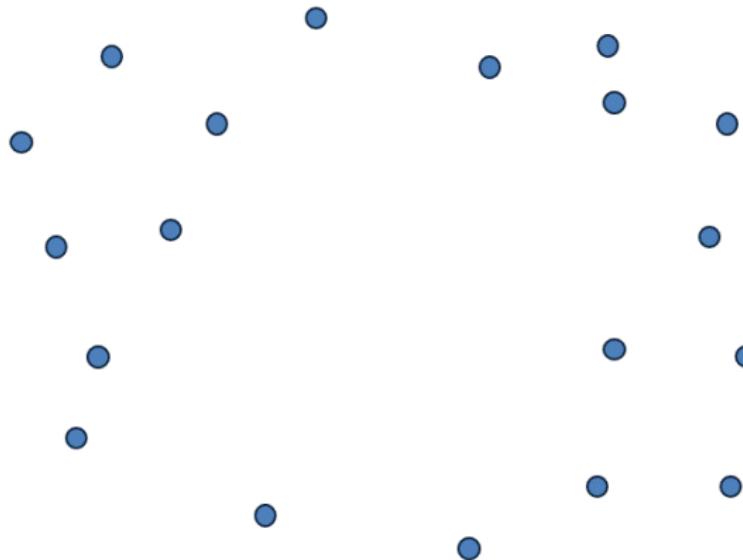
November 25, 2020

# What is this Talk About?

Approximating **persistent topological features**  
from point-clouds  
via (better) **sampling**

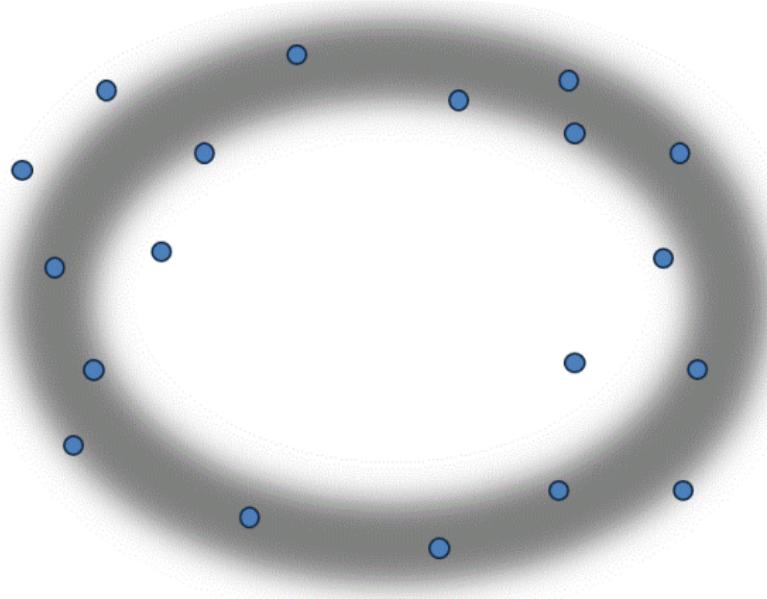
# Topological features

A point-cloud



# Topological features

The underlying space: A Ring in  $\mathbb{R}^2$

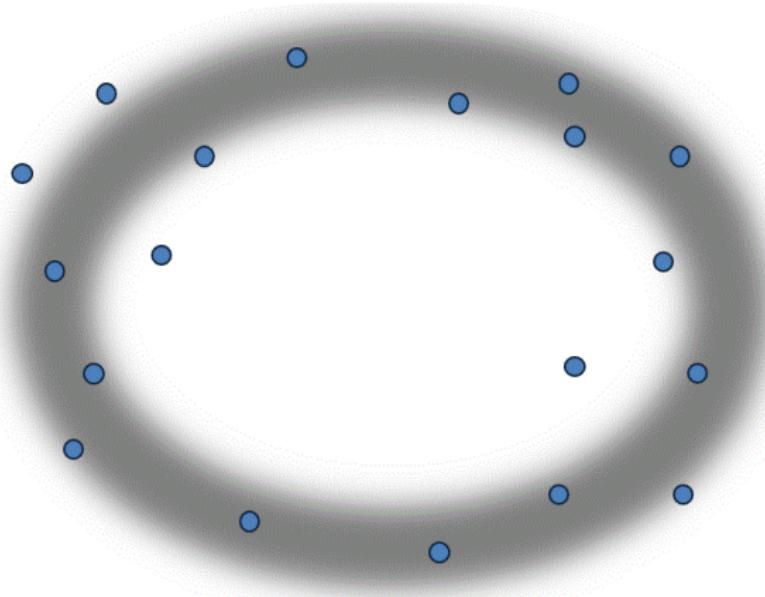


# Topological features

1 Connected component (Top. feat. at dim. 0)

1 cycle, Inner-cycle ~ Outer-cycle (Top. feat. at dim. 1)

0 void (Top. feat. at dim. 2)

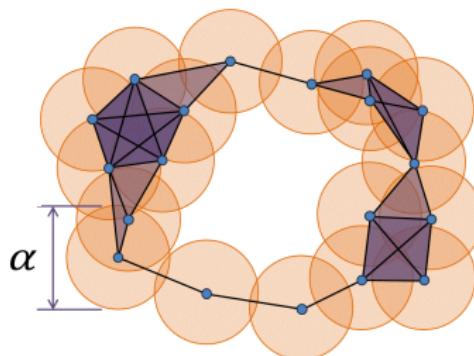


## Formal Representation: Simplicial complex

- Simplicial Complex: A set of simplices (0-simplex: A point, 1-simplex: Edge, 2-simplex: Filled triangle)

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- Simplicial Complex: A set of simplices (0-simplex: A point, 1-simplex: Edge, 2-simplex: Filled triangle)
- Choose a threshold  $\alpha$ .
  - Draw diameter  $\alpha$ -balls around each **point**.
  - Connect two points with an **edge** if their corresponding balls intersect pairwise.
  - Connect three points with a **filled triangle** if their corresponding balls intersect pairwise. And so on.



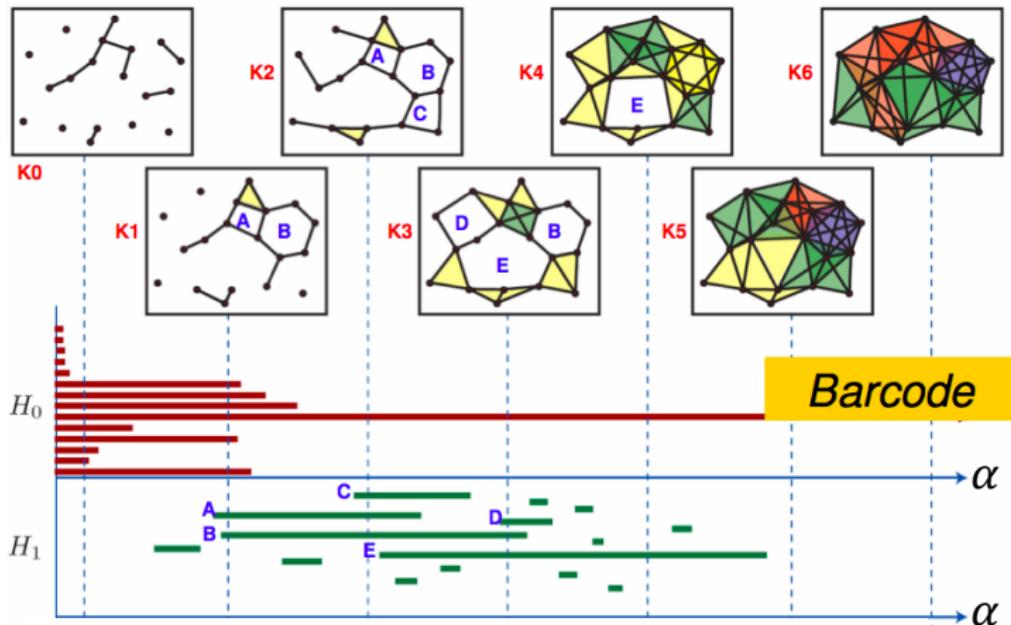
Vietoris-Rips complex at  $\alpha$  ( $R_\alpha$ ).

# Persistent Topological Features

- Issue: Choice of the right value for  $\alpha \rightarrow$  Persistence.

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- Issue: Choice of the right value for  $\alpha \rightarrow$  Persistence.
- Construct simplicial complex at different scales i.e.  $\alpha$ 's  $\rightarrow$  **Filtration**.
- Track appearance (birth) and merge (death) of topological features across scales of  $\alpha \rightarrow$  **Persistent homology**.



# Approximate Simplicial Representations

- **Čech complex** captures the actual topology of the underlying space of the point-cloud, but not feasible to compute → *at most  $(1 + n)^k$  simplices of dimension up to  $k$ .*
- **Vietoris-Rips Complex** is a 2-approximation of the Čech complex → *at most  $(1 + n)^k$  simplices of dimension up to  $k$ .*

Computational bottleneck: Enumerating large number of simplices.

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Computational bottleneck: Enumerating large number of simplices.

## The Central Computational Question of TDA

Can we have approximate simplicial representations which are **computable in reasonable time**, yet **good approximations to Vietoris-Rips or Čech complex?**

# A Computationally Faster Approximation: Lazy Witness Complex and The Question to Solve

## Lazy witness Complex $LW_\alpha(P, L, \nu)$

**Lazy witness Complex**  $LW_\alpha(P, L, \nu)$  of a **point-cloud**  $P$  is a simplicial complex over a **landmark set**  $L$  that consists of  $k$ -simplices  $[v_0 v_1 \cdots v_k]$  whose any two points  $v_i, v_j$  are in  $\alpha + d_\nu$  proximity of some witness point  $w$  ( $d_\nu$  is the distance from  $w$  to its  $\nu$ -th nearest neighbour in  $L$ .)

- Lazy witness complex is Vietoris-Rips complex on landmarks  $L$  for  $\nu = 0$ .
- The size of the lazy witness complex is at most  $(1 + |L|)^k$  where  $|L| \ll n$ .

# A Computationally Faster Approximation: Lazy Witness Complex and The Question to Solve

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## New Questions

How to select the landmarks?

How good are the landmarks selected by an algorithm?

Can we obtain any approximation guarantee for the lazy witness complex?

# Our Contributions

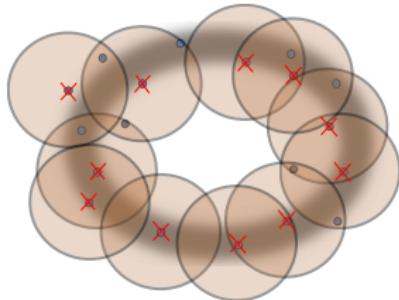
## The Central Concept

We respond to all these questions by reincarnating the idea of  $\epsilon$ -net in TDA.

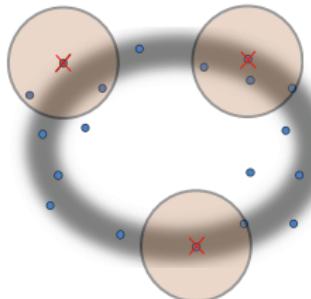
- Q. Can we obtain any approximation guarantee for the lazy witness complex?  
-> **Lazy witness complex induced by an  $\epsilon$ -net is a 3-approximation to the Vietoris-Rips complex.**
- Q. How good are the landmarks selected by an algorithm?  
->  **$\epsilon$ -net is an  $\epsilon$ -approximation of the point cloud.**
- Q. How to select the landmarks?  
-> **We propose three algorithms to construct  $\epsilon$ -net.**

Additionally, we validate these theoretical claims experimentally.

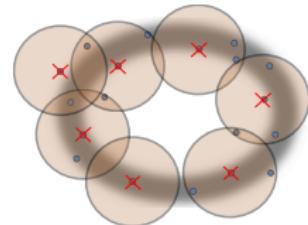
# The Central Concept: $\epsilon$ -Net



$\epsilon$ -sample (Each blue point is within  $\epsilon$  of some red point)



$\epsilon$ -sparse (Each pair of red points are  $\epsilon$ -far from each other)

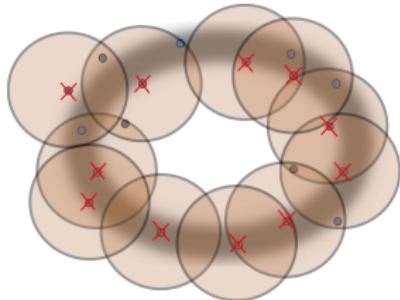


$\epsilon$ -net

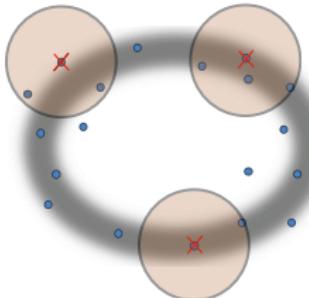
## Definition ( $\epsilon$ -sample)

A set  $L \subseteq P$  is an  $\epsilon$ -sample of  $P$  if the collection  $\{B_\epsilon(x) : x \in L\}$  of  $\epsilon$ -balls of radius  $\epsilon$ -covers  $P$ , i.e.  $P = \bigcup_{x \in L} B_\epsilon(x)$ .

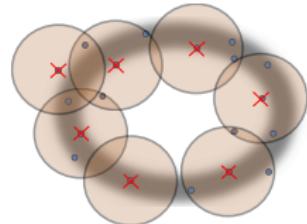
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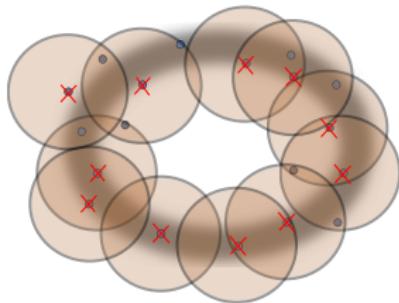


$\epsilon$ -net

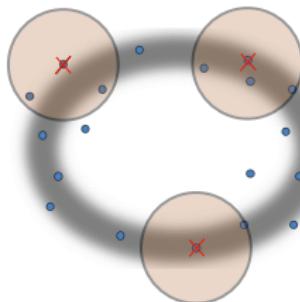
## Definition ( $\epsilon$ -sparse)

A set  $L \subset P$  is  $\epsilon$ -sparse if for all  $x, y \in L$ ,  $d(x, y) > \epsilon$ .

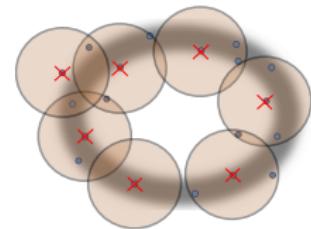
# The Central Concept: $\epsilon$ -Net



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$\epsilon$ -net

## $\epsilon$ -net

A subset ( $L$ ) of points which is  $\epsilon$ -sparse and  $\epsilon$ -sample of the point-cloud ( $P$ ).

# Approximation Guarantee: $\epsilon$ -Net Induced Lazy Witness Complex

## Approximating the Vietoris-Rips Complex

If the landmark set  $L$  is an  $\epsilon$ -net of the point cloud  $P$ , the lazy witness complex at  $\alpha$  and  $\nu = 1$  is 3-approximation of the Vietoris-Rips complex of  $L$  for  $\alpha \geq 2\epsilon$ . Mathematically,

$$R_{\alpha/3}(L) \subseteq LW_\alpha(P, L, 1) \subseteq R_{3\alpha}(L) \quad \forall \alpha \geq 2\epsilon.$$

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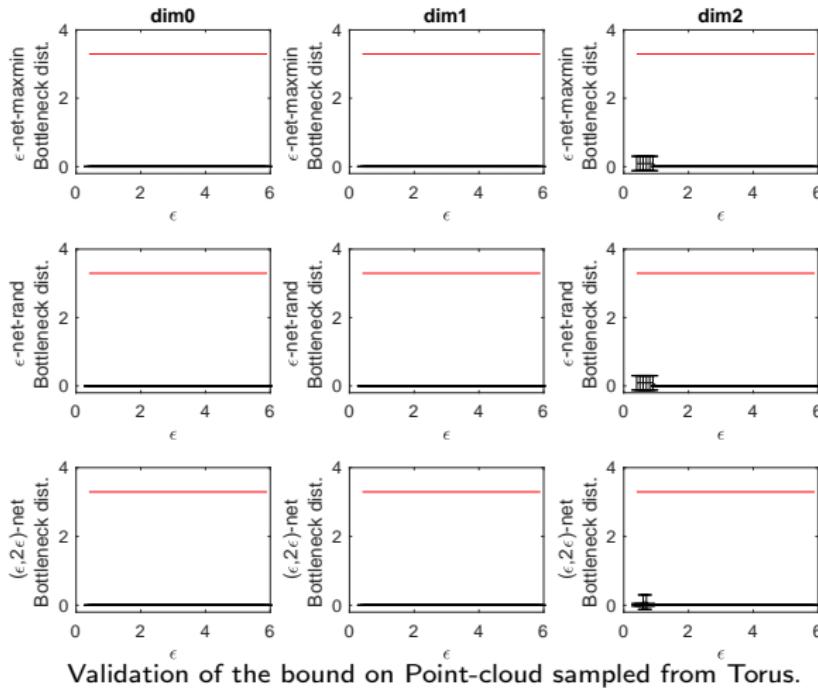
$$R_{\alpha/3}(L) \subseteq LW_\alpha(P, L, 1) \subseteq R_{3\alpha}(L) \quad \forall \alpha \geq 2\epsilon.$$

## Approximating the Persistent topological feature

If we compare the bars (in the barcode) appearing after  $2\epsilon$ , barcodes (log-scale) of the lazy witness filtration and the Vietoris-Rips filtration are  $3 \log 3$ -approximations of each other. (By weak-stability theorem<sup>a</sup>)

<sup>a</sup>Chazal et al., “Proximity of persistence modules and their diagrams”.

# Experimental Validation of Approximation Guarantee



## Point-cloud Approximation Guarantee

The Hausdorff distance between the point cloud and its  $\epsilon$ -net is at most  $\epsilon$ .

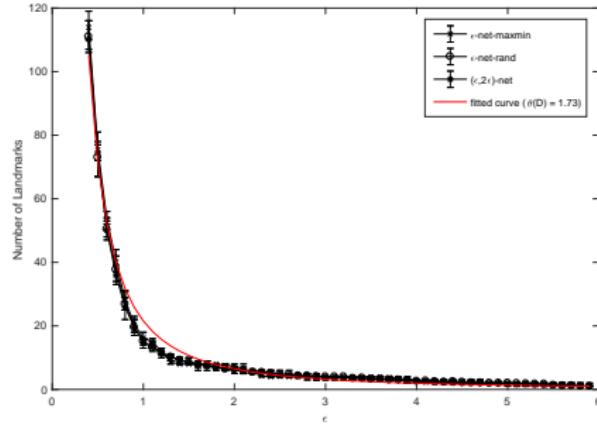
# Quality of Landmarks: Properties of $\epsilon$ -Net

## Point-cloud Approximation Guarantee

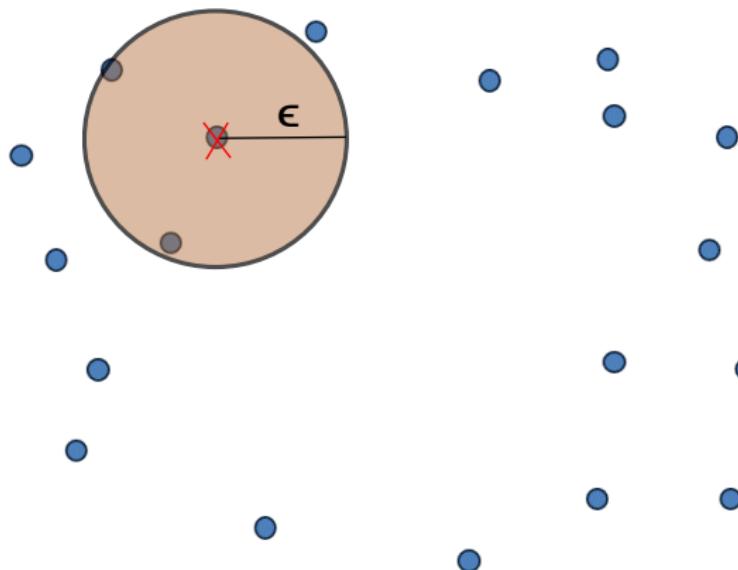
The Hausdorff distance between the point cloud and its  $\epsilon$ -net is at most  $\epsilon$ .

## Size of an $\epsilon$ -Net

The number of points in an  $\epsilon$ -net is at most  $(\frac{\Delta}{\epsilon})^{\theta(D)}$  for  $P \in \mathbb{R}^D$  of diameter  $\Delta$ .

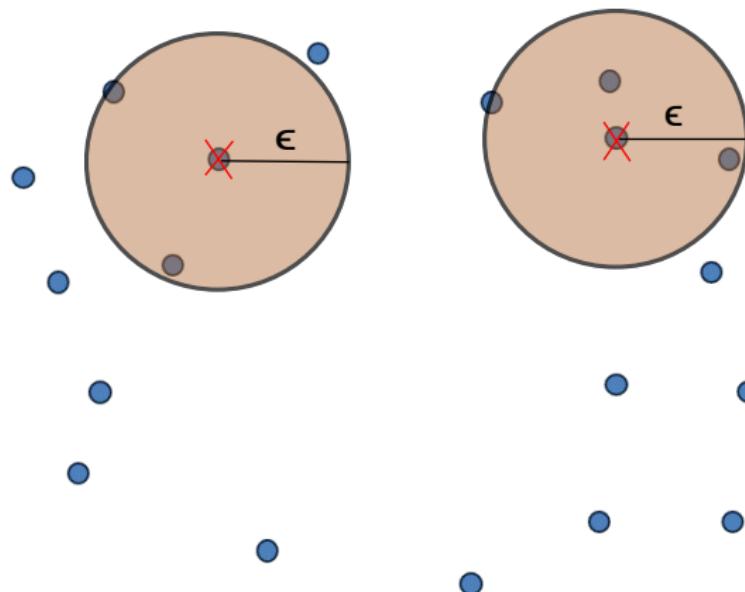


## Algorithm: $\epsilon$ -net-rand



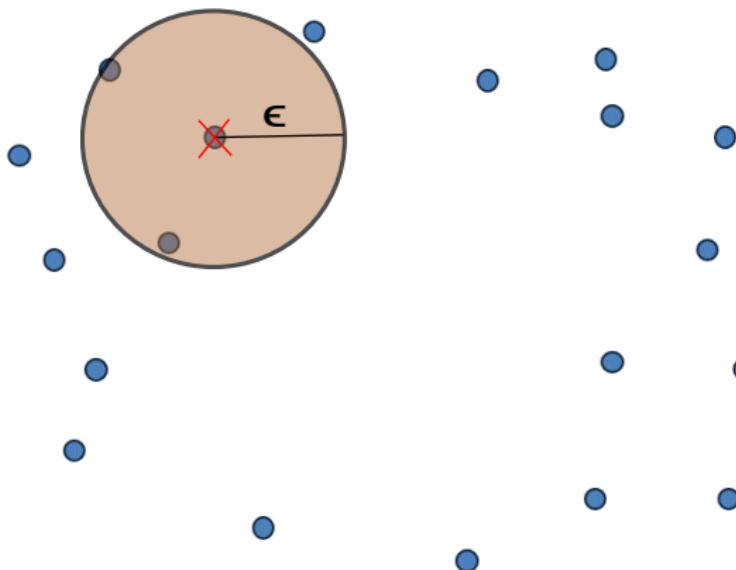
First landmark: Select uniformly at random from the point-cloud. Mark points in its  $\epsilon$ -ball.

## Algorithm: $\epsilon$ -net-rand



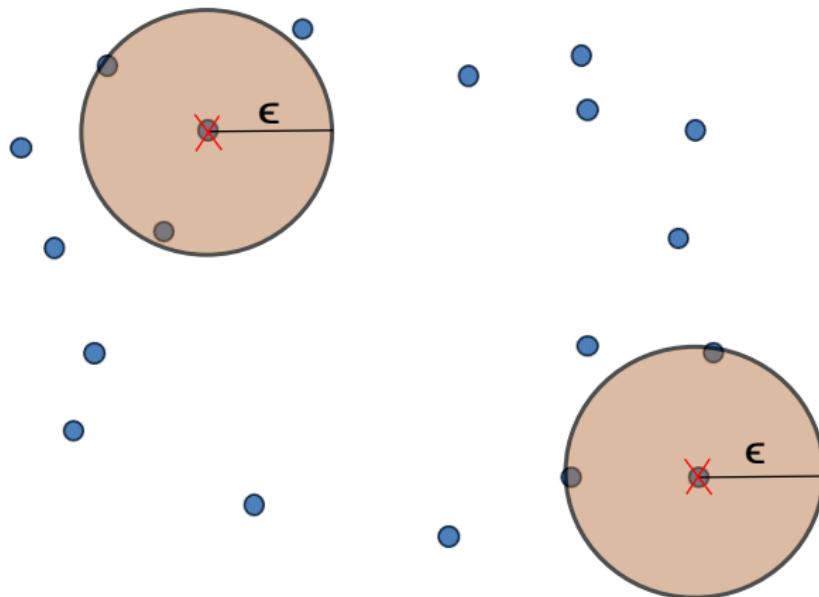
Next landmark: Select u.a.r from the set of unmarked points. Mark points in its  $\epsilon$ -ball. And so on.

## Algorithm: $\epsilon$ -net-maxmin



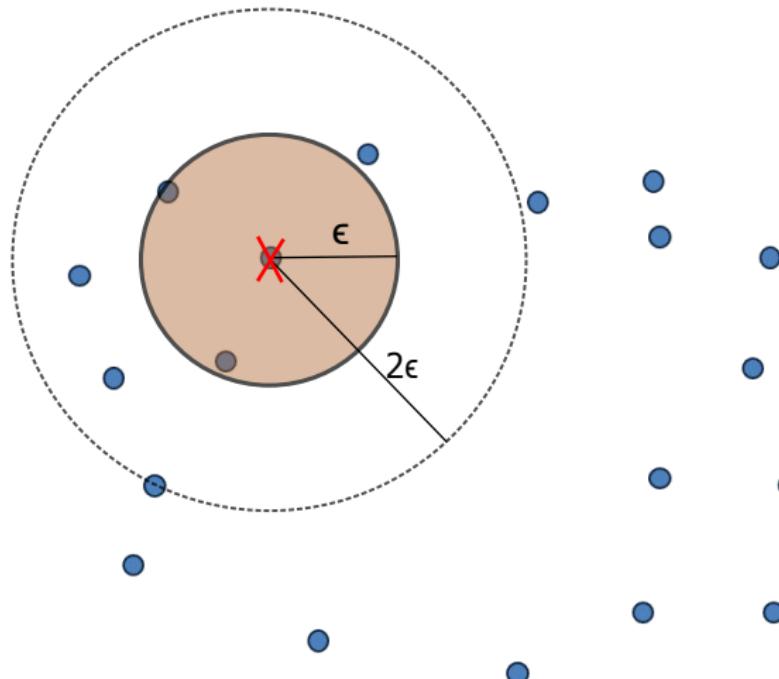
First landmark: Select u.a.r. from the point-cloud. Mark points in its  $\epsilon$ -ball.

## Algorithm: $\epsilon$ -net-maxmin



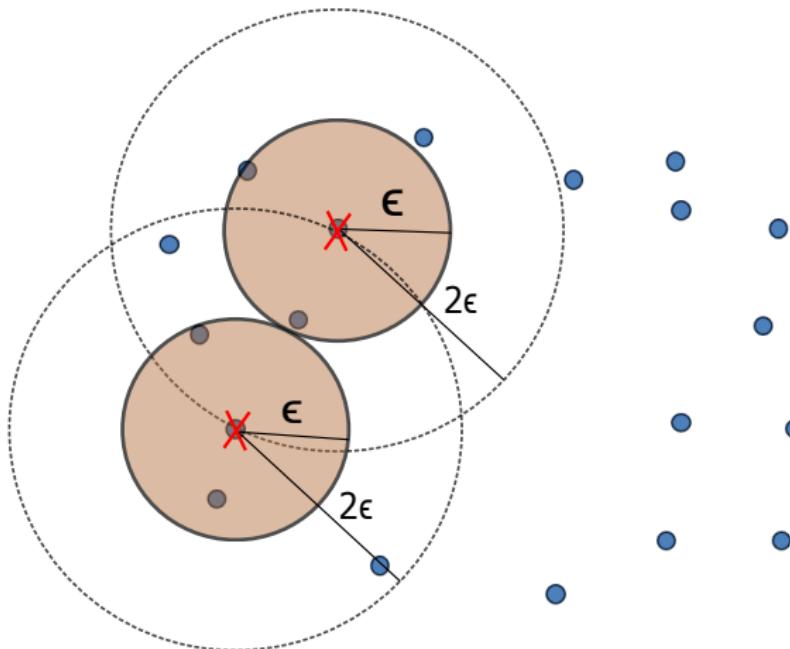
Next landmark: Select the point that is farthest from the current set of landmarks. And so on.

## Algorithm: $(\epsilon, 2\epsilon)$ -net



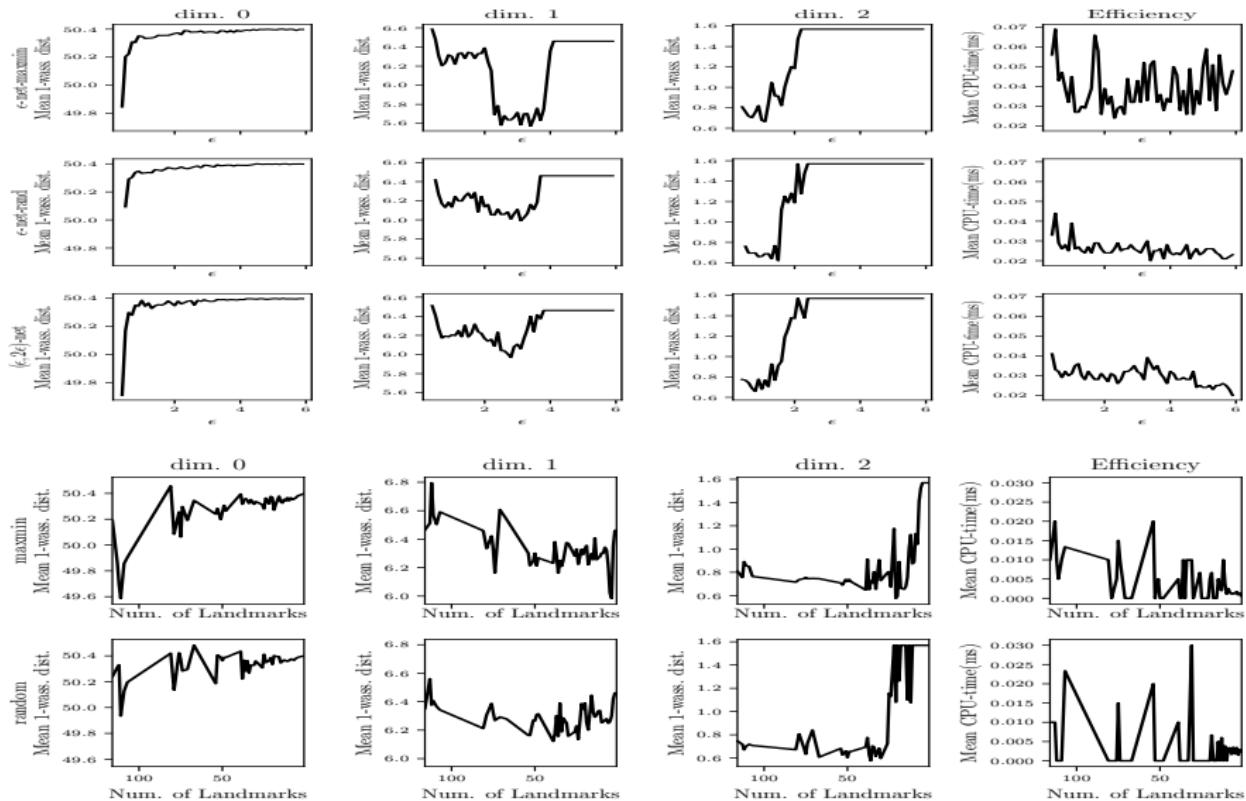
First landmark: Select u.a.r from the point-cloud. Mark points in its  $\epsilon$ -ball.

## Algorithm: $(\epsilon, 2\epsilon)$ -net



Second landmark: Select u.a.r from the unmarked points in the  $(\epsilon, 2\epsilon)$  envelope of the current set of landmarks. And so on.

# Experimental Evaluation: Effectiveness-Efficiency



# Take Away

- If the landmarks is an  $\epsilon$ -net, we know about the quality of the-
  - landmarks
  - lazy witness approximation
  - approximated persistent topological features.
- Use  $\epsilon$ -net as landmarks.
- *You have a point-cloud dataset? -> Apply Topological Data Analysis!*

# Thank You!

# Supplementary Slides

## Experiments: Datasets

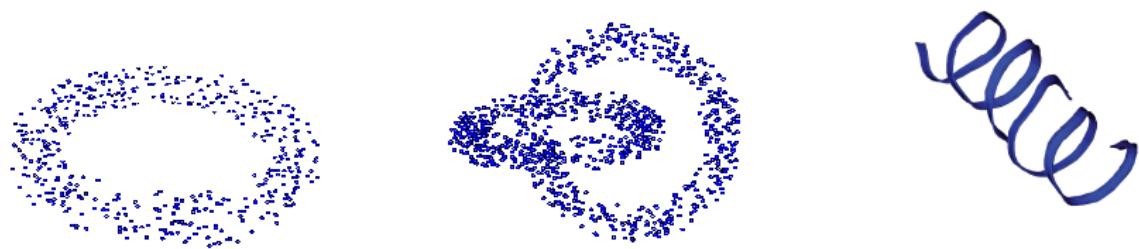


Figure 1: (left) Torus, (middle) Tangled-torus, and (right) 1grm Dataset

# Relation to Maxmin and Random Landmark Selection Algorithms

- Given the number of landmarks  $K > 1$ , the set of landmarks selected by the algorithm random/maxmin is  $\delta$ -sparse where  $\delta$  is the minimum of the pairwise distances among the landmarks.
- The choice of  $K$  may not necessarily make the landmarks a  $\delta$ -sample of the point cloud.

# Algorithm complexity

- $\epsilon$ -net-rand:  $O\left(\frac{1}{\epsilon^D}\right)$
- $\epsilon$ -net-maxmin:  $O\left(\frac{n}{\epsilon^D}\right)$
- $(\epsilon, 2\epsilon)$ -net:  $O\left(\frac{1}{\epsilon^D}\right) \log\left(\frac{1}{\epsilon}\right)$

# Experimental Evaluation: Stability

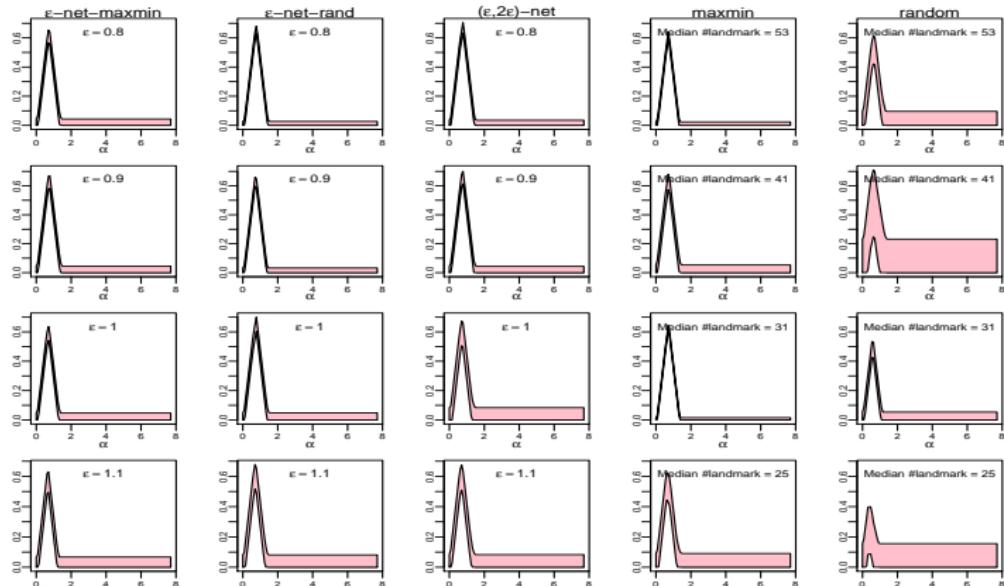


Figure 2: 95% confidence band of the rank one persistence landscape at dimension 1 of the lazy witness filtration induced by the landmark selection algorithms on Tangled-torus dataset.

# Whats Next?

The topological approximation guarantee is

- with respect to the Vietoris-Rips complex on  **$\epsilon$ -net landmarks** chosen from a **point-cloud input**.

Next up -

- Better guarantee w.r.t Vietoris-Rips complex on point-cloud:-
  - Improved the approximation guarantee from 3-approximation of  $R_\alpha(L)$  to  $\frac{3\log(c)}{2}$ -approximation of  $R_\alpha(P)$  for  $c \geq 2$ .
- Graph data <sup>1</sup>:-
  - Defined  $\epsilon$ -net for graphs.
  - Devised algorithms for computing  $\epsilon$ -net of graphs.
  - Potential applications: Graph clustering, Graph visualization, Graph classification.
- Comparison with Sparse-Rips and Graph Induced filtration (A weakness!).

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<sup>1</sup>To appear at ECML-PKDD'19 workshop on Applied Topological Data Analysis