

Homework-3

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1 Introduction

I analyzed several paper together with our lectures during my homework. I studied to understand those, followed step by step entire of the process. Now, we have to 5 problems and will try to solve step by step these.

2 Analysis of problems

2.1 Problem 1

(1) IP(1248842112488421): In this process only place of values change, not themselves of values. We can figure out it with a little codes snippet.

Listing 1: for to permute IP

```
def permutation(ptext, table, nbit):  
    permuted = ""  
    for i in range(0, nbit):  
        permuted = permuted + ptext[table[i]-1]  
    return permuted
```

Result: 2211448844882211

(2) E(84211248): By changing the parameters of the code snippet above, we can find the E permutation.

Result: 4081028A4251

(3) P(12488421): Similarly;

Result: 41010C4A

(4) Si(110011) for $i = 1; 2; : : : ; 8$: We can use below codes:

Listing 2: for to S-Box

```
for i in range(0,8):
    val = S_BOX[i][3][9]
    sboxTable = sboxTable + dec2bin(val)
```

Result: 1011 0110 1111 0100 1111 1110 0101 1100

2.2 Problem 2

2.2.1

Suppose that $K_1 = K_2 = \dots = K_{16}$. Show that all bits in C_1 are equal and all bits in D_1 are equal

Key scheduling is exactly independent a process from making ciphertext. The first half of the 56-bit key is left, also its second half is right are called as. At the same time, entire of the this process is logic and can be inversion. Now, we can establish the logical connection. K_i values is generated after from come side by side C_i and D_i values, and this process consists two of two steps. (a) Left-shift and (b) PC2. C_i and D_i values are make left-shift for each round. After new values are permuted with PC2. According to our suppose, if entire of the our keys are same, this mean entire of C_i are same and also D_i are same. I encrypted the 'CryptoEn' to explain with an example. The result is as follow.

```

Plaintext giriniz: CryptoEn
C 1 : FFFFFFFF D 1 : FFFFFFFF K 1 : FFFFFFFFFFFFFFFF
C 2 : FFFFFFFF D 2 : FFFFFFFF K 2 : FFFFFFFFFFFFFFFF
C 3 : FFFFFFFF D 3 : FFFFFFFF K 3 : FFFFFFFFFFFFFFFF
C 4 : FFFFFFFF D 4 : FFFFFFFF K 4 : FFFFFFFFFFFFFFFF
C 5 : FFFFFFFF D 5 : FFFFFFFF K 5 : FFFFFFFFFFFFFFFF
C 6 : FFFFFFFF D 6 : FFFFFFFF K 6 : FFFFFFFFFFFFFFFF
C 7 : FFFFFFFF D 7 : FFFFFFFF K 7 : FFFFFFFFFFFFFFFF
C 8 : FFFFFFFF D 8 : FFFFFFFF K 8 : FFFFFFFFFFFFFFFF
C 9 : FFFFFFFF D 9 : FFFFFFFF K 9 : FFFFFFFFFFFFFFFF
C 10 : FFFFFFFF D 10 : FFFFFFFF K 10 : FFFFFFFFFFFFFFFF
C 11 : FFFFFFFF D 11 : FFFFFFFF K 11 : FFFFFFFFFFFFFFFF
C 12 : FFFFFFFF D 12 : FFFFFFFF K 12 : FFFFFFFFFFFFFFFF
C 13 : FFFFFFFF D 13 : FFFFFFFF K 13 : FFFFFFFFFFFFFFFF
C 14 : FFFFFFFF D 14 : FFFFFFFF K 14 : FFFFFFFFFFFFFFFF
C 15 : FFFFFFFF D 15 : FFFFFFFF K 15 : FFFFFFFFFFFFFFFF
C 16 : FFFFFFFF D 16 : FFFFFFFF K 16 : FFFFFFFFFFFFFFFF
Round 1  00BEA4A3  2EA0A9B5  FFFFFFFFFFFFFFFF
Round 2  2EA0A9B5  7520B599  FFFFFFFFFFFFFFFF
Round 3  7520B599  0A02D290  FFFFFFFFFFFFFFFF
Round 4  0A02D290  344FD080  FFFFFFFFFFFFFFFF
Round 5  344FD080  50F9FD50  FFFFFFFFFFFFFFFF
Round 6  50F9FD50  F254D34A  FFFFFFFFFFFFFFFF
Round 7  F254D34A  DB642A76  FFFFFFFFFFFFFFFF
Round 8  DB642A76  189C3893  FFFFFFFFFFFFFFFF
Round 9  189C3893  4F603EE3  FFFFFFFFFFFFFFFF
Round 10  4F603EE3  1A52C542  FFFFFFFFFFFFFFFF
Round 11  1A52C542  2F758704  FFFFFFFFFFFFFFFF
Round 12  2F758704  2AEAA316  FFFFFFFFFFFFFFFF
Round 13  2AEAA316  AE908DF9  FFFFFFFFFFFFFFFF
Round 14  AE908DF9  F310184C  FFFFFFFFFFFFFFFF
Round 15  F310184C  6A1126A6  FFFFFFFFFFFFFFFF
Round 16  9AC8C838  6A1126A6  FFFFFFFFFFFFFFFF
Plaintext : 43727970746F456E
Key       : FEFEFEFEFEFEFEF
Ciphertext: 20CA0AD5618B9456

```

Figure 1: if K_i values are the same

2.2.2

Show that there are exactly 4 DES keys for which all round keys are the same. They are called weak DES keys.

If we carefully analyze, we can see that dropped 48-bit of PC2 permutation consists its first half from 1-28th bits and also second half from 29-48th bits. Therefore, we can easily say that first 24 bits of the K_i values formed from PC2 are left and second 24 bits of the K_i values formed from PC2 are right. The results is follow:

```

Plaintext giriniz: CryptoEn
C 1 : 0000000 D 1 : FFFFFFFF K 1 : 000000FFFFFF
C 2 : 0000000 D 2 : FFFFFFFF K 2 : 000000FFFFFF
C 3 : 0000000 D 3 : FFFFFFFF K 3 : 000000FFFFFF
C 4 : 0000000 D 4 : FFFFFFFF K 4 : 000000FFFFFF
C 5 : 0000000 D 5 : FFFFFFFF K 5 : 000000FFFFFF
C 6 : 0000000 D 6 : FFFFFFFF K 6 : 000000FFFFFF
C 7 : 0000000 D 7 : FFFFFFFF K 7 : 000000FFFFFF
C 8 : 0000000 D 8 : FFFFFFFF K 8 : 000000FFFFFF
C 9 : 0000000 D 9 : FFFFFFFF K 9 : 000000FFFFFF
C 10 : 0000000 D 10 : FFFFFFFF K 10 : 000000FFFFFF
C 11 : 0000000 D 11 : FFFFFFFF K 11 : 000000FFFFFF
C 12 : 0000000 D 12 : FFFFFFFF K 12 : 000000FFFFFF
C 13 : 0000000 D 13 : FFFFFFFF K 13 : 000000FFFFFF
C 14 : 0000000 D 14 : FFFFFFFF K 14 : 000000FFFFFF
C 15 : 0000000 D 15 : FFFFFFFF K 15 : 000000FFFFFF
C 16 : 0000000 D 16 : FFFFFFFF K 16 : 000000FFFFFF
Round 1 008EA4A3 AE6939E1 000000FFFFFF
Round 2 AE6939E1 21DBEE9D 000000FFFFFF
Round 3 21DBEE9D 3194F799 000000FFFFFF
Round 4 3194F799 0411E67C 000000FFFFFF
Round 5 0411E67C C8746140 000000FFFFFF
Round 6 C8746140 398C868F 000000FFFFFF
Round 7 398C868F 38654B59 000000FFFFFF
Round 8 38654B59 0EBC0D0A 000000FFFFFF
Round 9 0EBC0D0A 52AE1A19 000000FFFFFF
Round 10 52AE1A19 58C91737 000000FFFFFF
Round 11 58C91737 1BA75F33 000000FFFFFF
Round 12 1BA75F33 682193E3 000000FFFFFF
Round 13 682193E3 2E891679 000000FFFFFF
Round 14 2E891679 29EBFFC6 000000FFFFFF
Round 15 29EBFFC6 34998097 000000FFFFFF
Round 16 34998097 000000FFFFFF
Plaintext : 43727970746F456E
Key : 1F1F1F1F0E0E0E0E
Ciphertext: 73178335A285512F

```

(a) 1F1F1F1F 0E0E0E0E

```

Plaintext giriniz: CryptoEn
C 1 : FFFFFFFF D 1 : 00000000 K 1 : FFFFFFFF000000
C 2 : FFFFFFFF D 2 : 00000000 K 2 : FFFFFFFF000000
C 3 : FFFFFFFF D 3 : 00000000 K 3 : FFFFFFFF000000
C 4 : FFFFFFFF D 4 : 00000000 K 4 : FFFFFFFF000000
C 5 : FFFFFFFF D 5 : 00000000 K 5 : FFFFFFFF000000
C 6 : FFFFFFFF D 6 : 00000000 K 6 : FFFFFFFF000000
C 7 : FFFFFFFF D 7 : 00000000 K 7 : FFFFFFFF000000
C 8 : FFFFFFFF D 8 : 00000000 K 8 : FFFFFFFF000000
C 9 : FFFFFFFF D 9 : 00000000 K 9 : FFFFFFFF000000
C 10 : FFFFFFFF D 10 : 00000000 K 10 : FFFFFFFF000000
C 11 : FFFFFFFF D 11 : 00000000 K 11 : FFFFFFFF000000
C 12 : FFFFFFFF D 12 : 00000000 K 12 : FFFFFFFF000000
C 13 : FFFFFFFF D 13 : 00000000 K 13 : FFFFFFFF000000
C 14 : FFFFFFFF D 14 : 00000000 K 14 : FFFFFFFF000000
C 15 : FFFFFFFF D 15 : 00000000 K 15 : FFFFFFFF000000
C 16 : FFFFFFFF D 16 : 00000000 K 16 : FFFFFFFF000000
Round 1 008EA4A3 1DB6819C FFFFFFFF000000
Round 2 1DB6819C 1FC48C8D FFFFFFFF000000
Round 3 1FC48C8D D8EAC44D FFFFFFFF000000
Round 4 D8EAC44D A613ACFB FFFFFFFF000000
Round 5 A613ACFB 18BAC219 FFFFFFFF000000
Round 6 18BAC219 9775BDCB FFFFFFFF000000
Round 7 9775BDCB 31204E4E FFFFFFFF000000
Round 8 31204E4E DBA9CEA7 FFFFFFFF000000
Round 9 DBA9CEA7 F6C0FEF9 FFFFFFFF000000
Round 10 F6C0FEF9 83BD9A17 FFFFFFFF000000
Round 11 83BD9A17 DD8DA216 FFFFFFFF000000
Round 12 DD8DA216 1C101C56 FFFFFFFF000000
Round 13 1C101C56 348F42E0 FFFFFFFF000000
Round 14 348F42E0 95C852EE FFFFFFFF000000
Round 15 95C852EE 7DAE4B27 FFFFFFFF000000
Round 16 7DAE4B27 00000000 FFFFFFFF000000
Plaintext : 43727970746F456E
Key : E0E0E0E0F1F1F1F1
Ciphertext: CB6AB6ACC0A38C31

```

(b) E0E0E0E0 F1F1F1F1

```

Plaintext giriniz: CryptoEn
C 1 : FFFFFFFF D 1 : FFFFFFFF K 1 : FFFFFFFF000000
C 2 : FFFFFFFF D 2 : FFFFFFFF K 2 : FFFFFFFF000000
C 3 : FFFFFFFF D 3 : FFFFFFFF K 3 : FFFFFFFF000000
C 4 : FFFFFFFF D 4 : FFFFFFFF K 4 : FFFFFFFF000000
C 5 : FFFFFFFF D 5 : FFFFFFFF K 5 : FFFFFFFF000000
C 6 : FFFFFFFF D 6 : FFFFFFFF K 6 : FFFFFFFF000000
C 7 : FFFFFFFF D 7 : FFFFFFFF K 7 : FFFFFFFF000000
C 8 : FFFFFFFF D 8 : FFFFFFFF K 8 : FFFFFFFF000000
C 9 : FFFFFFFF D 9 : FFFFFFFF K 9 : FFFFFFFF000000
C 10 : FFFFFFFF D 10 : FFFFFFFF K 10 : FFFFFFFF000000
C 11 : FFFFFFFF D 11 : FFFFFFFF K 11 : FFFFFFFF000000
C 12 : FFFFFFFF D 12 : FFFFFFFF K 12 : FFFFFFFF000000
C 13 : FFFFFFFF D 13 : FFFFFFFF K 13 : FFFFFFFF000000
C 14 : FFFFFFFF D 14 : FFFFFFFF K 14 : FFFFFFFF000000
C 15 : FFFFFFFF D 15 : FFFFFFFF K 15 : FFFFFFFF000000
C 16 : FFFFFFFF D 16 : FFFFFFFF K 16 : FFFFFFFF000000
Round 1 008EA4A3 2EA0A9B5 FFFFFFFF000000
Round 2 2EA0A9B5 7520B599 FFFFFFFF000000
Round 3 7520B599 0A02D290 FFFFFFFF000000
Round 4 0A02D290 34FD0880 FFFFFFFF000000
Round 5 34FD0880 50F9FD50 FFFFFFFF000000
Round 6 50F9FD50 F254D34A FFFFFFFF000000
Round 7 F254D34A DB642A76 FFFFFFFF000000
Round 8 DB642A76 189C3893 FFFFFFFF000000
Round 9 189C3893 4F603EE3 FFFFFFFF000000
Round 10 4F603EE3 1A52C542 FFFFFFFF000000
Round 11 1A52C542 2F758704 FFFFFFFF000000
Round 12 2F758704 2AEA3316 FFFFFFFF000000
Round 13 2AEA3316 AE908DF9 FFFFFFFF000000
Round 14 AE908DF9 F310184C FFFFFFFF000000
Round 15 F310184C 6A1126A6 FFFFFFFF000000
Round 16 6A1126A6 00000000 FFFFFFFF000000
Plaintext : 43727970746F456E
Key : FEFEFEFEFEFEFEFE
Ciphertext: 20CA0A05618B9456

```

(c) FEFEFEFEFEFEFEFE

```

Plaintext giriniz: CryptoEn
C 1 : 00000000 D 1 : 00000000 K 1 : 00000000000000
C 2 : 00000000 D 2 : 00000000 K 2 : 00000000000000
C 3 : 00000000 D 3 : 00000000 K 3 : 00000000000000
C 4 : 00000000 D 4 : 00000000 K 4 : 00000000000000
C 5 : 00000000 D 5 : 00000000 K 5 : 00000000000000
C 6 : 00000000 D 6 : 00000000 K 6 : 00000000000000
C 7 : 00000000 D 7 : 00000000 K 7 : 00000000000000
C 8 : 00000000 D 8 : 00000000 K 8 : 00000000000000
C 9 : 00000000 D 9 : 00000000 K 9 : 00000000000000
C 10 : 00000000 D 10 : 00000000 K 10 : 00000000000000
C 11 : 00000000 D 11 : 00000000 K 11 : 00000000000000
C 12 : 00000000 D 12 : 00000000 K 12 : 00000000000000
C 13 : 00000000 D 13 : 00000000 K 13 : 00000000000000
C 14 : 00000000 D 14 : 00000000 K 14 : 00000000000000
C 15 : 00000000 D 15 : 00000000 K 15 : 00000000000000
C 16 : 00000000 D 16 : 00000000 K 16 : 00000000000000
Round 1 008EA4A3 9D7F11C8 00000000000000
Round 2 9D7F11C8 4C8F5206 00000000000000
Round 3 4C8F5206 6489AB94 00000000000000
Round 4 6489AB94 96770C19 00000000000000
Round 5 96770C19 D9207B6C 00000000000000
Round 6 D9207B6C 261E981A 00000000000000
Round 7 261E981A 95C39B62 00000000000000
Round 8 95C39B62 D9A3C6AA 00000000000000
Round 9 D9A3C6AA 47D52BD4 00000000000000
Round 10 47D52BD4 02DFCDF6 00000000000000
Round 11 02DFCDF6 5A379C20 00000000000000
Round 12 5A379C20 2B02E202 00000000000000
Round 13 2B02E202 030B1FA6 00000000000000
Round 14 030B1FA6 F6485C1A 00000000000000
Round 15 F6485C1A AF07F530 00000000000000
Round 16 2197D09B AF07F530 00000000000000
Plaintext : 43727970746F456E
Key : 0101010101010101
Ciphertext: F0B18C851FCA0C9D

```

(d) 01010101 01010101

Figure 2: weak Keys

2.2.3

Determine these 4 weak DES keys.

The 4 weak keys were published by NIST in January 2012¹. These are:

- 01010101 01010101
- FEF EFEFE FEF EFEFE
- E0E0E0E0 F1F1F1F1
- 1F1F1F1F 0E0E0E0E

2.3 Problem 3

There are other modes of block cipher besides the ones we have learned. One of these modes is named Plaintext Block Chaining (PBC) Mode. On the encryption side, the following is executed to obtain the n th ciphertext: $C_n := E_k(M_n) \oplus M_{n-1}$. Suppose that we need to encrypt M_1, \dots, M_5 using the PBC mode. Show the explicit formulas to obtain C_1, \dots, C_5 . What do you need to use for M_0 ? Also, show the steps on the decryption side to obtain M_1, \dots, M_5 .

M_0 is Initialization Vector (IV). In this case, we can constitute encryption and decryption formulas:

- for encryption:
 - $C_1 = E_k(M_1) \oplus IV$
 - $C_2 = E_k(M_2) \oplus M_1$
 - $C_3 = E_k(M_3) \oplus M_2$
 - $C_4 = E_k(M_4) \oplus M_3$
 - $C_5 = E_k(M_5) \oplus M_4$

Interestingly, we continue toward from start to end in order to decryption, because decryption process continues toward from end to start in some algorithm.

¹Recommendation for the Triple Data Encryption Algorithm (TDEA) Block Cipher, NIST, January 2012, p.11

- for decryption:
 - $M_1 = D_k(C_1) \oplus IV$
 - $M_2 = D_k(C_2) \oplus C_1$
 - $M_3 = D_k(C_3) \oplus C_2$
 - $M_4 = D_k(C_4) \oplus C_3$
 - $M_5 = D_k(C_5) \oplus C_4$

2.4 Problem 4

Consider the AES/Rijndael algorithm and its Galois field $GF(2^8)$

2.4.1 Compute the sum (a7) + (5c) in $GF(2^8)$

At first, we should convert these hex values to binary numbers.

$$(a7) = 1\ 0\ 1\ 0\ 0\ 1\ 1\ 1$$

$$(5c) = 0\ 1\ 0\ 1\ 1\ 1\ 0\ 0$$

After, we should write to binary numbers as polynomial and should sum to this polynomials (indeed XOR);

$$\begin{aligned} & (x^7 + x^5 + x^2 + x^1 + 1) \oplus (x^6 + x^4 + x^3 + x^2) \\ &= x^7 + x^6 + x^5 + x^4 + x^3 + x^1 + 1 \end{aligned}$$

Now, we should convert these polynomials to binary numbers.

$$= \underbrace{1111}_f \underbrace{1011}_b$$

$$\underline{\text{Result}} = (a7) + (5c) = (fb)$$

2.4.2 Compute the product (a7) × (5c) in GF(2⁸)

At first, we should convert these hex values to binary numbers.

$$(a7) = 1\ 0\ 1\ 0\ 0\ 1\ 1\ 1$$

$$(5c) = 0\ 1\ 0\ 1\ 1\ 1\ 0\ 0$$

Besides, I want to solve with two method.

First method

$$\begin{aligned} & (x^7 + x^5 + x^2 + x^1 + 1) \times (x^6 + x^4 + x^3 + x^2) \\ = & x^{13} + x^{11} + x^{10} + x^9 + x^{11} + x^9 + x^8 + x^7 + x^8 + x^6 + x^5 + x^4 + x^7 \\ & + x^5 + x^4 + x^3 + x^6 + x^4 + x^3 + x^2 \\ = & x^{13} + x^{10} + x^4 + x^2 \end{aligned}$$

And now, I will divide to $p(x) = x^8 + x^4 + x^3 + x + 1$

$$\begin{array}{r|l} x^{13} + x^{10} & + x^4 + x^2 \\ - x^{13} & - x^9 - x^8 - x^6 - x^5 \\ \hline x^{10} - x^9 - x^8 - x^6 - x^5 & + x^4 + x^2 \\ - x^{10} & - x^6 - x^5 - x^3 - x^2 \\ \hline - x^9 - x^8 - 2x^6 - 2x^5 & + x^4 - x^3 \\ x^9 & + x^5 + x^4 + x^2 + x \\ \hline - x^8 - 2x^6 - x^5 + 2x^4 - x^3 + x^2 & + x \\ x^8 & + x^4 + x^3 + x + 1 \\ \hline - 2x^6 - x^5 + 3x^4 & + x^2 + 2x + 1 \\ \hline = x^5 + x^4 + x^2 + 1 \end{array}$$

Now, we will convert it to binary numbers;

$$= \underbrace{0011}_3 \underbrace{0101}_5$$

Result: (a7) x (5c) = (35)

Second method

At first, we are convert to $x^{13} + x^{10} + x^4 + x^2$ polynomial to binary numbers and will write its under;

1 0 0 1 0 0 0 0 0 1 0 1 0 0

1 0 0 0 1 1 0 1 1

0 0 0 1 1 1 0 1 1 1 0 1 0 0

0 0 0 1 0 0 0 1 1 0 1 1

0 0 0 0 1 1 0 0 0 1 1 0

0 0 0 0 1 0 0 0 1 1 0 1 1

0 0 0 0 0 1 0 0 1 0 1 1 1 0

0 0 0 0 0 1 0 0 0 1 1 0 1 1

0 0 0 0 0 0 0 0 1 1 0 1 0 1

We take to last 8 numbers and convert it hex number;

$\underbrace{0011}_3 \underbrace{0101}_5$

Result: (a7) x (5c) = (35)

2.4.3 Compute S(a7)

I will explain in details solving of this question in next question, for this reason I prefer to solve by looking to the table.

Result: S(a7) = (5c)

		y															
		0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
x	0	63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	76
	1	ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
	2	b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
	3	04	c7	23	c3	18	96	05	9a	07	12	80	e2	eb	27	b2	75
	4	09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
	5	53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
	6	d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
	7	51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
	8	cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
	9	60	81	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
	a	e0	32	3a	0a	49	06	24	5c	c2	d3	ac	62	91	95	e4	79
	b	e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
	c	ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
	d	70	3e	b5	66	48	03	f6	0e	61	35	57	b9	86	c1	1d	9e
	e	e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
	f	8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

Figure 3: The SubByte S-Box

2.4.4 Compute $S^{-1}(5c)$

We will try to find on the $c = A.b + d$ formula. In this formula, if $a(x) \neq 0$ we will compute multiplicative inverse. After, the value found will equal to $b(x)$.

2.5 Problem 5

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